# ACM ICPC REGIONAL 2010

## 1. Generales

# 1.1. LIS en O(nlgn).

```
vector<int> LIS(vector<int> X) {
   int n = X.size(), L = 0, M[n+1], P[n];
   int lo, hi, mi;

L = 0;
   M[0] = 0;

for(int i=0, j; i<n; i++) {
   lo = 0; hi = L;

   while(lo!=hi) {
      mi = (lo+hi+1)/2;

      if(X[M[mi]]<X[i]) lo = mi;
      else hi = mi-1;
   }

   j = lo;</pre>
```

# 1.2. Problema de Josephus.

```
int survivor(int n, int m) {
   for (int s=0,i=1;i<=n;++i) s = (s+m)%i;</pre>
```

# 1.3. Lectura rápida de enteros.

```
void readInt(int &n) {
   int sign = 1;
   char c;
  bool found = false;
```

```
P[i] = M[j];

if(j==L || X[i]<X[M[j+1]]) {
    M[j+1] = i;
    L = max(L, j+1);
  }
}
int a[L];

for(int i=L-1, j=M[L]; i>=0; i--) {
    a[i] = X[j];
    j = P[j];
}

return vector<int>(a, a+L);
```

```
return (s+1);
}

n = 0;
while(true){
```

c = getc(stdin);

```
switch(c) {
   case '-' :
      sign = -1;
      found = true;
      break;
   case '_':
      if(found) goto jump;
      break;
   case '\n':
      if(found) goto jump;
      break;
```

#### 1.4. Contar inversiones.

```
#define MAX_SIZE 100000
int A[MAX_SIZE],C[MAX_SIZE],pos1,pos2,sz;

long long countInversions(int a, int b) {
   if (a==b) return 0;

   int c = ((a+b)>>1);
   long long aux = countInversions(a,c)+countInversions(c+1,b);
   pos1 = a; pos2 = c+1; sz = 0;

while(pos1<=c && pos2<=b) {
   if (A[pos1]<A[pos2]) C[sz] = A[pos1++];
   else{</pre>
```

# 1.5. Números dada la suma de pares.

```
bool solve(int N, int sums[], int ans[]) {
  int M = N*(N-1)/2;
  multiset<int> S;
  multiset<int> :: iterator it;

  sort(sums,sums+M);

  for(int i = 2;i<M;++i) {
    if((sums[0]+sums[1]-sums[i])%2!=0) continue;

  ans[0] = (sums[0]+sums[1]-sums[i])/2;</pre>
```

```
default:
         if(c>='0' && c<='9'){
            n = n * 10 + c - '0';
            found = true;
          }else goto jump;
         break;
jump:
   n *= sign;
      C[sz] = A[pos2++];
      aux += c-pos1+1;
   ++sz;
if (pos1>c) memcpy (C+sz, A+pos2, (b-pos2+1) *sizeof(int));
else memcpy(C+sz,A+pos1,(c-pos1+1)*sizeof(int));
sz = b-a+1;
memcpy(A+a,C,sz*sizeof(int));
return aux;
   S = multiset<int>(sums,sums+M);
   bool valid = true;
   for(int j = 1; j<N && valid; ++j) {</pre>
      ans[j] = (*S.begin())-ans[0];
      for (int k = 0; k< j && valid; ++k) {</pre>
         it = S.find(ans[k]+ans[j]);
         if(it==S.end()) valid = false;
```

```
else S.erase(it);
}

if(valid) return true;
```

return false;
}

#### 2. Grafos

#### 2.1. Ciclo de Euler.

```
// Las listas de adyacencia se deben ordenar de forma ascendente para
// obtener el ciclo lexicografico minimo deacuerdo a la numeracion
// de las aristas
#define MAX_V 44
#define MAX_E 1995
int N, deg[MAX_V], eu[MAX_E], ev[MAX_E];
list<int> G[MAX_V],L;
bool visited[MAX_V];
stack<int> S;
queue<int> Q;
bool connected() {
   int cont = 0;
   Q.push(0);
   memset (visited, false, sizeof (visited));
   visited[0] = true;
   while(!Q.empty()){
      int v = Q.front(); Q.pop();
      ++cont;
      for(list<int>::iterator it = G[v].begin();it!=G[v].end();++it){
            int e = *it;
            int w = eu[e] == v? ev[e] : eu[e];
         if(!visited[w]){
            visited[w] = true;
            Q.push(w);
   return cont == N:
```

```
bool eulerian(){
  if(!connected()) return false;
   for (int v = 0; v < N; ++v)
     if(deg[v]&1)
         return false;
   return true;
void take_edge(int v, int w) {
  --deg[v]; --deg[w];
  int e = G[v].front();
  G[v].pop_front();
   for(list<int>::iterator it = G[w].begin();it!=G[w].end();++it){
      if(*it==e){
         G[w].erase(it);
         break;
void euler(int v) {
  while(true) {
      if(G[v].empty()) break;
      int e = G[v].front();
      int w = eu[e] == v? ev[e] : eu[e];
      S.push(e);
      take_edge(v,w);
      v = w;
```

```
bool find_cycle(int s) {
   if(!eulerian()) return false;

int v = s,e;
   L.clear();

do{
     euler(v);
     e = S.top(); S.pop();
     L.push_back(e);

   v = eu[e] == v? ev[e] : eu[e];
} while(!S.empty());

return true;
```

#### 2.2. Union-Find.

```
#define MAX_SIZE 26
int parent[MAX_SIZE], rank[MAX_SIZE];

void Make_Set(const int x) {
   parent[x] = x; rank[x] = 0;
}

int Find(const int x) {
   if(parent[x]!=x) parent[x] = Find(parent[x]);
   return parent[x];
}
```

#### 2.3. Punto de articulación.

```
#define SZ 100
bool M[SZ][SZ];
int N,colour[SZ],dfsNum[SZ],num,pos[SZ],leastAncestor[SZ],parent[SZ];
int dfs(int u) {
   int ans = 0,cont = 0,v;
   stack<int> S;
   S.push(u);
```

```
void print_cycle(int s){
  if(!find_cycle(s)) printf("-1\n");
      bool first = true;
      reverse(L.begin(), L.end());
      for(list<int>::iterator e = L.begin();e!=L.end();++e){
            if(!first) printf("_");
            first = false;
         printf("%d",1+(*e));
      printf("\n");
void Union(const int x, const int y) {
  int PX = Find(x), PY = Find(y);
  if(rank[PX]>rank[PY]) parent[PY] = PX;
      parent[PX] = PY;
      if(rank[PX] == rank[PY]) ++ rank[PY];
   while(!S.empty()){
      v = S.top();
      if(colour[v] == 0) {
         colour[v] = 1;
         dfsNum[v] = num++;
         leastAncestor[v] = num;
      for(;pos[v]<N;++pos[v]){</pre>
         if(M[v][pos[v]] && pos[v]!=parent[v]) {
```

## 2.4. Detección de puentes.

```
#define SZ 100
bool M[SZ][SZ];
int N,colour[SZ],dfsNum[SZ],num,pos[SZ],leastAncestor[SZ],parent[SZ];

void dfs(int u) {
   int v;
   stack<int> S;
   S.push(u);

   while(!S.empty()) {
      v = S.top();
      if(colour[v] == 0) {
        colour[v] = 1;
        dfsNum[v] = num++;
        leastAncestor[v] = num;
   }
```

```
for(int j = 0; j<N; j++)
         if(M[i][j] && parent[j]==i && leastAncestor[j]>=dfsNum[i]) {
            printf("%d\n",i);
            ++ans;
            break;
   return ans;
void Articulation_points() {
   memset(colour, 0, sizeof(colour));
   memset(pos, 0, sizeof(pos));
   memset (parent, -1, sizeof (parent));
   num = 0;
   int total = 0;
   for(int i = 0;i<N;++i) if(colour[i]==0) total += dfs(i);</pre>
   printf("#_Articulation_Points_:_%d\n",total);
      for(;pos[v]<N;++pos[v]){</pre>
         if(M[v][pos[v]] && pos[v]!=parent[v]){
            if(colour[pos[v]]==0){
               parent[pos[v]] = v;
               S.push(pos[v]);
               break;
            }else leastAncestor[v] <?= dfsNum[pos[v]];</pre>
      if (pos[v]==N) {
         colour[v] = 2;
         S.pop();
         if(v!=u) leastAncestor[parent[v]] <?= leastAncestor[v];</pre>
```

```
void Bridge_detection() {
   memset(colour,0,sizeof(colour));
   memset(pos,0,sizeof(pos));
   memset(parent,-1,sizeof(parent));
   num = 0;
   int ans = 0;
   for(int i = 0;i<N;i++) if(colour[i]==0) dfs(i);</pre>
```

## 2.5. Componentes biconexas (Tarjan).

```
#define MAXN 100000
int V;
vector<int> adj[MAXN];
int dfn[MAXN],low[MAXN];
vector< vector<int> > C;
stack< pair<int, int> > stk;
void cache_bc(int x, int y) {
   vector<int> com;
   int tx, ty;
      tx = stk.top().first, ty = stk.top().second;
      stk.pop();
      com.push_back(tx), com.push_back(ty);
   }while(tx!=x || ty!=y);
   C.push_back(com);
void DFS(int cur, int prev, int number) {
   dfn[cur] = low[cur] = number;
   for(int i = adj[cur].size()-1;i>=0;--i){
      int next = adj[cur][i];
      if(next==prev) continue;
```

```
for (int i = 0; i < N; i++)</pre>
      for(int j = 0; j<N; j++)
         if(parent[j]==i && leastAncestor[j]>dfsNum[i]){
            printf("%d_-_%d\n",i,j);
            ++ans;
  printf("%d_bridges\n",ans);
      if (dfn[next] ==-1) {
         stk.push(make_pair(cur,next));
         DFS (next, cur, number+1);
         low[cur] = min(low[cur], low[next]);
         if(low[next]>=dfn[cur]) cache_bc(cur,next);
      }else low[cur] = min(low[cur],dfn[next]);
void biconnected_components() {
   memset (dfn, -1, sizeof (dfn));
  C.clear();
  DFS(0,0,0);
   int comp = C.size();
  printf("%d\n",comp);
   for (int i = 0; i < comp; ++i) {</pre>
      sort(C[i].begin(),C[i].end());
      C[i].erase(unique(C[i].begin(),C[i].end()),C[i].end());
      int m = C[i].size();
      for(int j = 0; j<m; ++j) printf("%d_", 1+C[i][j]);</pre>
      printf("\n");
```

#### 2.6. Componentes fuertemente conexas (Tarjan).

```
#define MAX_V 100000
vector<int> L[MAX_V],C[MAX_V];
int V, dfsPos, dfsNum[MAX_V], lowlink[MAX_V], num_scc;
bool in_stack[MAX_V];
stack<int> S;
void tarjan(int v) {
   dfsNum[v] = lowlink[v] = dfsPos++;
   S.push(v); in_stack[v] = true;
   for(int i = L[v].size()-1;i>=0;--i){
      int w = L[v][i];
      if(dfsNum[w] == -1) {
         tarjan(w);
         lowlink[v] = min(lowlink[v],lowlink[w]);
      }else if(in_stack[w]) lowlink[v] = min(lowlink[v], lowlink[w]);
   if(dfsNum[v] == lowlink[v]) {
      vector<int> com;
      int aux;
```

# 2.7. Ciclo de peso promedio mínimo (Karp).

```
#define MAX_V 676

vector< pair<int, int> > L[MAX_V+1];
int dist[MAX_V+1][MAX_V+2];

void karp(int n) {
    for(int i = 0;i<n;++i)
        if(!L[i].empty())
            L[n].push_back(make_pair(i,0));
        ++n;

    for(int i = 0;i<n;++i)
        fill(dist[i],dist[i]+(n+1),INT_MAX);

    dist[n-1][0] = 0;</pre>
```

```
do√
         aux = S.top(); S.pop();
         com.push_back(aux);
         in_stack[aux] = false;
      }while (aux!=v);
      C[num\_scc] = com;
      ++num_scc;
void build_scc() {
  memset (dfsNum, -1, sizeof (dfsNum));
   memset(in_stack, false, sizeof(in_stack));
   dfsPos = num_scc = 0;
   for(int i = 0;i<V;++i)</pre>
      if(dfsNum[i]==-1)
         tarjan(i);
   for (int k = 1; k \le n; ++k) for (int u = 0; u \le n; ++u) {
      if (dist[u][k-1]==INT_MAX) continue;
      for(int i = L[u].size()-1;i>=0;--i)
         dist[L[u][i].first][k] = min(dist[L[u][i].first][k],
                                   dist[u][k-1]+L[u][i].second);
  bool flag = true;
   for(int i = 0;i<n && flag;++i)</pre>
      if (dist[i][n]!=INT_MAX)
         flag = false;
   if(flag){
      //El grafo es aciclico
```

return;

```
double ans = le15;
for(int u = 0;u+1<n;++u) {
   if(dist[u][n]==INT_MAX) continue;
   double W = -le15;</pre>
```

#### 2.8. Minimum cost arborescence.

```
#define MAX V 1000
typedef int edge_cost;
edge_cost INF = INT_MAX;
int V, root, prev[MAX_V];
bool adj[MAX_V][MAX_V];
edge_cost G[MAX_V][MAX_V], MCA;
bool visited[MAX_V], cycle[MAX_V];
void add_edge(int u, int v, edge_cost c){
   if(adj[u][v]) G[u][v] = min(G[u][v],c);
   else G[u][v] = c;
   adj[u][v] = true;
void dfs(int v) {
   visited[v] = true;
   for (int i = 0; i < V; ++i)</pre>
      if(!visited[i] && adj[v][i])
         dfs(i);
bool check() {
   memset(visited, false, sizeof(visited));
   dfs(root);
   for (int i = 0; i < V; ++i)</pre>
      if(!visited[i])
         return false;
   return true;
```

```
for(int k = 0; k < n; ++ k)
    if (dist[u][k]!=INT_MAX)
        W = max(W, (double) (dist[u][n]-dist[u][k]) / (n-k));
    ans = min(ans, W);
}</pre>
```

```
int exist_cycle(){
  prev[root] = root;
   for (int i = 0; i < V; ++i) {</pre>
      if(!cycle[i] && i!=root){
         prev[i] = i; G[i][i] = INF;
         for (int j = 0; j<V; ++j)</pre>
             if(!cycle[j] && adj[j][i] && G[j][i] <G[prev[i]][i])</pre>
                prev[i] = j;
   for (int i = 0, j; i < V; ++i) {</pre>
      if(cycle[i]) continue;
      memset (visited, false, sizeof (visited));
      j = i;
      while(!visited[j]){
         visited[j] = true;
         j = prev[j];
      if(j==root) continue;
      return j;
   return -1;
void update(int v) {
  MCA += G[prev[v]][v];
```

```
for(int i = prev[v];i!=v;i = prev[i]){
  MCA += G[prev[i]][i];
   cycle[i] = true;
for (int i = 0; i < V; ++i)</pre>
   if(!cycle[i] && adj[i][v])
      G[i][v] -= G[prev[v]][v];
for(int j = prev[v]; j!=v; j = prev[j]) {
   for(int i = 0; i<V; ++i) {</pre>
      if(cycle[i]) continue;
      if(adj[i][j]){
         if(adj[i][v]) G[i][v] = min(G[i][v],G[i][j]-G[prev[j]][j]);
         else G[i][v] = G[i][j]-G[prev[j]][j];
         adj[i][v] = true;
      if(adj[j][i]){
         if(adj[v][i]) G[v][i] = min(G[v][i],G[j][i]);
         else G[v][i] = G[j][i];
         adj[v][i] = true;
```

## 2.9. Ordenamiento Topológico.

```
#define MAX_v 100000
#define MAX_E 100000

int V,E,indeg[MAX_V],topo_pos[MAX_V];
int last[MAX_V],next[MAX_E],to[MAX_E];
int Q[MAX_V],head,tail;

void init() {
    memset(indeg,0,sizeof(indeg));
    memset(last,-1,sizeof(last));
}

void add_edge(int u, int v) {
    to[E] = v, next[E] = last[u], last[u] = E; ++E;
    ++indeg[v];
}

void topological_sort() {
    head = tail = 0;
```

```
bool min_cost_arborescence(int _root) {
   root = _root;
   if(!check()) return false;
   memset (cycle, false, sizeof (cycle));
  MCA = 0;
   int v;
   while((v = exist_cycle())!=-1)
      update(v);
   for (int i = 0; i < V; ++i)</pre>
      if(i!=root && !cycle[i])
         MCA += G[prev[i]][i];
   return true;
   for (int i = 0; i < V; ++i) {</pre>
      if(indeg[i]==0){
         topo_pos[i] = tail;
         Q[tail++] = i;
   while (head!=tail) {
      int u = Q[head++];
      for(int e = last[u], v; e! =-1; e = next[e]) {
         v = to[e];
         --indeg[v];
         if(indeg[v]==0){
            topo_pos[v] = tail;
            Q[tail++] = v;
```

```
}
```

## 2.10. Diámetro de un árbol.

```
#define MAX_SIZE 100
bool visited[MAX_SIZE];
int prev[MAX_SIZE];
int most_distant(int s) {
    queue<int> Q;
    Q.push(s);

    memset(visited, false, sizeof(visited));
    visited[s] = true;
    prev[s] = -1;
    int ans = s;

while(!Q.empty()) {
    int aux = Q.front();
```

# 2.11. Stable marriage.

```
Q.pop();
ans = aux;

for(int i=L[aux].size()-1;i>=0;--i){
    int v = L[aux][i];
    if(visited[v]) continue;
    visited[v] = true;
    Q.push(v);
    prev[v] = aux;
}

return ans;
}
```

# 2.12. Bipartite matching (Hopcroft Karp).

```
#define MAX_V1 50000
#define MAX_V2 50000
#define MAX_E 150000
int V1, V2, 1 [MAX_V1], r [MAX_V2];
int E, to[MAX_E], next[MAX_E], last[MAX_V1];
bool visited[MAX_V1];
void init(){
      memset(last,-1, sizeof(int) *V1);
      E = 0;
void add_edge(int u, int v) {
      to[E] = v, next[E] = last[u]; last[u] = E; ++E;
bool pairup(int u) {
   if (visited[u]) return false;
   visited[u] = true;
   for(int e = last[u];e!=-1;e = next[e]){
      int v = to[e];
      if(r[v]==-1 || pairup(r[v])){
         l[u] = v;
         r[v] = u;
```

## 2.13. Algoritmo húngaro.

```
#define MAX_V 500

int V,cost[MAX_V][MAX_V];
int lx[MAX_V],ly[MAX_V];
int max_match,xy[MAX_V],yx[MAX_V],prev[MAX_V];
bool S[MAX_V],T[MAX_V];
int slack[MAX_V],slackx[MAX_V];
int q[MAX_V],head,tail;

void init_labels(){
    memset(lx,0,sizeof(lx));
    memset(ly,0,sizeof(ly));
```

```
return true;
   return false;
int hopcroft_karp(){
  bool change = true;
   memset(1,-1,sizeof(int)*V1);
   memset(r,-1,sizeof(int)*V2);
   while (change) {
      change = false;
      memset (visited, false, sizeof (bool) *V1);
      for(int i = 0; i<V1; ++i)</pre>
         if(l[i]==-1) change |= pairup(i);
   int ret = 0;
   for (int i = 0; i < V1; ++i)</pre>
      if(l[i]!=-1) ++ret;
   return ret;
   for (int x = 0; x < V; ++x)
      for (int y = 0; y < V; ++y)
         lx[x] = max(lx[x], cost[x][y]);
void update_labels() {
   int x,y,delta = INT_MAX;
   for(y = 0;y<V;++y) if(!T[y]) delta = min(delta,slack[y]);</pre>
   for (x = 0; x < V; ++x) if (S[x]) 1x[x] -= delta;
   for(y = 0;y<V;++y) if(T[y]) ly[y] += delta;</pre>
```

```
for(y = 0;y<V;++y) if(!T[y]) slack[y] -= delta;</pre>
void add_to_tree(int x, int prevx) {
   S[x] = true;
   prev[x] = prevx;
   for (int y = 0; y < V; ++y) {
      if(lx[x]+ly[y]-cost[x][y]<slack[y]){
          slack[y] = lx[x]+ly[y]-cost[x][y];
          slackx[y] = x;
void augment(){
   int x,y,root;
   head = tail = 0;
   memset(S, false, sizeof(S));
   memset(T, false, sizeof(T));
   memset(prev,-1, sizeof(prev));
   for (x = 0; x < V; ++x) {
      if (xy[x]==-1) {
         q[tail++] = root = x;
         prev[root] = -2;
         S[root] = true;
         break;
   for (y = 0; y < V; ++y) {
      slack[y] = lx[root]+ly[y]-cost[root][y];
      slackx[y] = root;
   while(true) {
      while(head<tail){</pre>
         x = q[head++];
          for (y = 0; y < V; ++y)  {
             if(cost[x][y] == lx[x] + ly[y] && !T[y]) {
                if (yx[y] ==-1) break;
                T[y] = true;
                q[tail++] = yx[y];
```

```
add_to_tree(yx[y],x);
         if(y<V) break;</pre>
      if(y<V) break;</pre>
      update labels();
      head = tail = 0;
      for(y = 0;y<V;++y){
         if(!T[y] && slack[y]==0){
            if (yx[y] ==-1) {
                x = slackx[y];
                break;
            T[y] = true;
            if(!S[yx[y]]){
                q[tail++] = yx[y];
                add_to_tree(yx[y],slackx[y]);
      if(y<V) break;</pre>
   ++max_match;
   for (int cx = x, cy = y, ty; cx!=-2; cx = prev[cx], cy = ty) {
      ty = xy[cx];
      yx[cy] = cx;
      xy[cx] = cy;
int hungarian(){
  int ret = 0;
  max_match = 0;
  memset(xy,-1,sizeof(xy));
  memset (yx, -1, sizeof(yx));
```

```
init_labels();
for(int i = 0;i<V;++i) augment();
for(int x = 0;x<V;++x) ret += cost[x][xy[x]];</pre>
```

# 2.14. General matching (Gabow).

```
#define MAXV 200
#define MAXE 19900
int prev_edge[MAXE], v[MAXE], w[MAXE], last_edge[MAXV];
int type[MAXV], label[MAXV], first[MAXV], mate[MAXV], nedges;
bool g_flag[MAXV], g_souter[MAXV];
void g_init(){
   nedges = 0;
   memset(last_edge,-1, sizeof(last_edge));
void q_edge(int a, int b) {
   prev_edge[nedges] = last_edge[a];
   v[nedges] = a;
   w[nedges] = b;
   last_edge[a] = nedges++;
   prev_edge[nedges] = last_edge[b];
  v[nedges] = b;
   w[nedges] = a;
   last_edge[b] = nedges++;
void g_label(int v, int join, int edge, queue<int> &outer) {
   if(v==join) return;
   if(label[v] ==-1) outer.push(v);
   label[v] = edge;
   type[v] = 1;
   first[v] = join;
   g_label(first[label[mate[v]]], join, edge, outer);
void q_augment(int _v, int _w){
   int t = mate[_v];
  mate[\_v] = \_w;
```

```
return ret;
  if (mate[t]!=_v) return;
  if(label[ v]==-1) return;
  if(type[_v]==0){
      mate[t] = label[_v];
      g_augment(label[_v], t);
   }else if(type[_v]==1) {
      g_augment(v[label[_v]], w[label[_v]]);
      g_augment(w[label[_v]], v[label[_v]]);
int gabow(int n) {
  memset (mate, -1, sizeof (mate));
  memset(first,-1,sizeof(first));
  int u = 0, ret = 0;
   for (int z = 0; z < n; ++z) {
      if (mate[z]!=-1) continue;
      memset(label,-1, sizeof(label));
      memset(type,-1,sizeof(type));
      memset(g_souter, 0, sizeof(g_souter));
      label[z] = -1; type[z] = 0;
      queue<int> outer;
      outer.push(z);
      bool done = false;
      while(!outer.empty()){
         int x = outer.front(); outer.pop();
         if(q_souter[x]) continue;
         q_souter[x] = true;
         for(int i = last_edge[x];i!=-1;i = prev_edge[i]) {
```

```
if (mate[w[i]] == -1 && w[i]! = z) {
  mate[w[i]] = x;
  g_augment(x, w[i]);
   ++ret;
  done = true;
  break;
if(type[w[i]]==-1){
  int v = mate[w[i]];
  if(type[v] == -1){
     type[v] = 0;
     label[v] = x;
     outer.push(v);
      first[v] = w[i];
  continue;
int r = first[x], s = first[w[i]];
if(r==s) continue;
```

## 2.15. Flujo máximo (Dinic).

```
struct flow_graph{
   int MAX_V,E,s,t,head,tail;
   int *cap,*to,*next,*last,*dist,*q,*now;

flow_graph() {}

flow_graph(int V, int MAX_E) {
    MAX_V = V; E = 0;
    cap = new int[2*MAX_E], to = new int[2*MAX_E], next = new int[2*MAX_E];
    last = new int[MAX_V], q = new int[MAX_V];
    dist = new int[MAX_V], now = new int[MAX_V];
    fill(last,last+MAX_V,-1);
}

void clear() {
    fill(last,last+MAX_V,-1);
    E = 0;
}
```

```
memset(g_flag,0,sizeof(g_flag));
g_flag[r] = g_flag[s] = true;

while(true){
    if(s!=-1) swap(r, s);
    r = first[label[mate[r]]];
    if(g_flag[r]) break;
    g_flag[r] = true;
}

g_label(first[x], r, i, outer);
g_label(first[w[i]], r, i, outer);

for(int c = 0;c<n;++c)
    if(type[c]!=-1 && first[c]!=-1 && type[first[c]]!=-1)
        first[c] = r;
}
if(done) break;
}
return ret;
</pre>
```

```
void add_edge(int u, int v, int uv, int vu = 0){
    to[E] = v, cap[E] = uv, next[E] = last[u]; last[u] = E++;
    to[E] = u, cap[E] = vu, next[E] = last[v]; last[v] = E++;
}

bool bfs() {
    fill(dist, dist+MAX_V, -1);
    head = tail = 0;

    q[tail] = t; ++tail;
    dist[t] = 0;

while(head<tail) {
    int v = q[head]; ++head;

    for(int e = last[v]; e!=-1; e = next[e]) {
        if(cap[e^1]>0 && dist[to[e]]==-1) {
            q[tail] = to[e]; ++tail;
            dist[to[e]] = dist[v]+1;
            ]
```

```
}
}

return dist[s]!=-1;

int dfs(int v, int f) {
    if(v==t) return f;

for(int &e = now[v];e!=-1;e = next[e]) {
      if(cap[e]>0 && dist[to[e]]==dist[v]-1) {
        int ret = dfs(to[e],min(f,cap[e]));

    if(ret>0) {
        cap[e] -= ret;
        cap[e^1] += ret;
        return ret;
    }
}
```

#### 2.16. Flujo máximo - Costo Mínimo (Succesive Shortest Path).

```
#define MAX_V 350
#define MAX_E 2*12500

typedef int cap_type;
typedef long long cost_type;
const cost_type INF = LLONG_MAX;

int V,E,prev[MAX_V],last[MAX_V],to[MAX_E],next[MAX_E];
bool visited[MAX_V];
cap_type flowVal, cap[MAX_E];
cost_type flowCost,cost[MAX_E],dist[MAX_V],pot[MAX_V];

void init(int _V){
    memset(last,-1,sizeof(last));
    V = _V; E = 0;
}

void add_edge(int u, int v, cap_type capacity, cost_type cst){
    to[E] = v, cap[E] = capacity;
    cost[E] = cst, next[E] = last[u];
```

```
return 0;
}

long long max_flow(int source, int sink){
    s = source; t = sink;
    long long f = 0;
    int x;

    while(bfs()) {
        for(int i = 0;i<MAX_V;++i) now[i] = last[i];

        while(true) {
            x = dfs(s,INT_MAX);
            if(x==0) break;
            f += x;
        }
    }
}

return f;
}</pre>
```

```
last[u] = E++;
to[E] = u, cap[E] = 0;
cost[E] = -cst, next[E] = last[v];
last[v] = E++;
}

bool BellmanFord(int s, int t) {
  bool stop = false;
  for(int i = 0;i<V;++i) dist[i] = INF;
  dist[s] = 0;

  for(int i = 1;i<=V && !stop;++i) {
    stop = true;

    for(int j = 0;j<E;++j) {
        int u = to[j^1], v = to[j];

        if(cap[j]>0 && dist[u]!=INF && dist[u]+cost[j]<dist[v]) {
            stop = false;
            dist[v] = dist[u]+cost[j];</pre>
```

```
}
   for(int i = 0;i<V;++i) if (dist[i]!=INF) pot[i] = dist[i];</pre>
   return stop;
void mcmf(int s, int t) {
   flowVal = flowCost = 0;
   memset (pot, 0, sizeof (pot));
   if(!BellmanFord(s,t)){
      printf("Ciclo_negativo_de_capacidad_infinita");
      return;
   while(true) {
      memset (prev, -1, sizeof (prev));
      memset (visited, false, sizeof (visited));
      for(int i = 0;i<V;++i) dist[i] = INF;</pre>
      priority_queue< pair<cost_type, int> > Q;
      Q.push(make_pair(0,s));
      dist[s] = prev[s] = 0;
      while(!Q.empty()){
         int aux = Q.top().second;
         Q.pop();
```

# 2.17. Flujo máximo (Dinic + Lower Bounds).

```
struct flow_graph{
  int V,E,s,t;
  int *flow,*low,*cap,*to,*next,*last,*delta;
  int *dist,*q,*now,head,tail;

flow_graph() {}

flow_graph(int V, int E) {
    (*this).V = V; (*this).E = 0;
    int TE = 2*(E+V+1);
    flow = new int[TE]; low = new int[TE]; cap = new int[TE];
    to = new int[TE]; next = new int[TE];
```

```
if(visited[aux]) continue;
  visited[aux] = true;
   for(int e = last[aux];e!=-1;e = next[e]){
      if(cap[e]<=0) continue;</pre>
      cost_type new_dist = dist[aux]+cost[e]+pot[aux]-pot[to[e]];
      if(new_dist<dist[to[e]]){</pre>
         dist[to[e]] = new_dist;
         prev[to[e]] = e;
         Q.push(make_pair(-new_dist, to[e]));
if (prev[t]==-1) break;
cap_type f = cap[prev[t]];
for(int i = t;i!=s;i = to[prev[i]^1]) f = min(f,cap[prev[i]]);
for(int i = t;i!=s;i = to[prev[i]^1]){
   cap[prev[i]] -= f;
   cap[prev[i]^1] += f;
flowVal += f;
flowCost += f*(dist[t]-pot[s]+pot[t]);
for(int i = 0;i<V;++i) if (prev[i]!=-1) pot[i] += dist[i];</pre>
```

```
last = new int[V+2]; delta = new int[V];
dist = new int[V+2]; q = new int[V+2]; now = new int[V+2];
}

void clear(int V) {
   (*this).V = V; (*this).E = 0;
   fill(last,last+V,-1);
}

void add_edge(int a, int b, int l, int u) {
   to[E] = b; low[E] = 1; cap[E] = u; flow[E] = 0;
   next[E] = last[a]; last[a] = E++;
```

```
to[E] = a; low[E] = 0; cap[E] = 0; flow[E] = 0;
  next[E] = last[b]; last[b] = E++;
bool bfs(){
  fill(dist,dist+V+2,-1);
  head = tail = 0;
  g[tail] = t; ++tail;
  dist[t] = 0;
   while(head<tail){</pre>
      int v = q[head]; ++head;
      for(int e = last[v];e!=-1;e = next[e]){
         if(cap[e^1]>flow[e^1] && dist[to[e]]==-1) {
           q[tail] = to[e]; ++tail;
           dist[to[e]] = dist[v]+1;
  return dist[s]!=-1;
int dfs(int v, int f) {
  if(v==t) return f;
  for(int &e = now[v];e!=-1;e = next[e]){
      if(cap[e]>flow[e] && dist[to[e]] == dist[v]-1) {
         int ret = dfs(to[e],min(f,cap[e]-flow[e]));
         if(ret>0){
            flow[e] += ret;
            flow[e^1] -= ret;
            return ret;
  return 0;
int max_flow(int source, int sink) {
  fill(delta, delta+V, 0);
```

```
for(int e = 0; e < E; e += 2) {</pre>
   delta[to[e^1]] -= low[e];
   delta[to[e]] += low[e];
   cap[e] -= low[e];
last[V] = last[V+1] = -1;
int sum = 0;
for(int i = 0;i<V;++i){</pre>
  if(delta[i]>0){
      add_edge(V,i,0,delta[i]);
      sum += delta[i];
   if(delta[i]<0) add_edge(i,V+1,0,-delta[i]);</pre>
add_edge(sink, source, 0, INT_MAX);
s = V; t = V+1;
int f = 0, df;
fill(flow, flow+E, 0);
while(bfs()){
   for (int i = V+1;i>=0;--i) now[i] = last[i];
   while(true){
      df = dfs(s,INT_MAX);
      if (df==0) break;
      f += df;
if(f!=sum) return -1;
for(int e = 0;e<E;e += 2){</pre>
  cap[e] += low[e];
   flow[e] += low[e];
   flow[e^1] -= low[e];
   cap[e^1] -= low[e];
s = source; t = sink;
```

```
last[s] = next[last[s]];
last[t] = next[last[t]];
E -= 2;
while(bfs()) {
   for(int i = V-1;i>=0;--i) now[i] = last[i];
   while(true) {
      df = dfs(s,INT_MAX);
}
```

# 2.18. Corte mínimo de un grafo (Stoer - Wagner).

```
#define MAX_V 500
int M[MAX_V][MAX_V], w[MAX_V];
bool A[MAX_V], merged[MAX_V];
int minCut(int n) {
   int best = INT_MAX;
   for(int i=1;i<n;++i) merged[i] = false;</pre>
   merged[0] = true;
   for (int phase=1;phase<n;++phase) {</pre>
      A[0] = true;
      for (int i=1; i < n; ++i) {</pre>
          if(merged[i]) continue;
         A[i] = false;
          w[i] = M[0][i];
      int prev = 0,next;
      for(int i=n-1-phase;i>=0;--i){
          // hallar siguiente vrtice que no esta en A
          next = -1;
          for (int j=1; j<n; ++j)</pre>
             if(!A[j] && (next==-1 || w[j]>w[next]))
```

# 2.19. Graph Facts (No dirigidos).

Un grafo es bipartito si y solo si no contiene ciclos de longitud impar. Todos los arboles son bipartitos.

```
if(df==0) break;
    f += df;
}

return f;
};
```

```
next = j;

A[next] = true;

if(i>0) {
    prev = next;

    // actualiza los pesos
    for(int j=1;j<n;++j)
        if(!A[j]) w[j] += M[next][j];
    }

if(best>w[next]) best = w[next];

// mezcla s y t
for(int i=0;i<n;++i) {
    M[i][prev] += M[next][i];
    M[prev][i] += M[next][i];
}

merged[next] = true;
}

return best;</pre>
```

Las aristas que forman un ciclo, se encuentran en una misma componente biconexa.

# 3.1. Knuth-Morris-Pratt.

```
#define MAX_L 70
int f[MAX_L];

void prefixFunction(string P) {
   int n = P.size(), k = 0;
   f[0] = 0;

   for(int i=1;i<n;++i) {
      while(k>0 && P[k]!=P[i]) k = f[k-1];
      if(P[k]==P[i]) ++k;
      f[i] = k;
   }
}
```

# 3.2. Suffix array.

```
#define MAX_LEN 40000
#define ALPH_SIZE 123

char A[MAX_LEN+1];
int N,pos[MAX_LEN],rank[MAX_LEN];
int cont[MAX_LEN],next[MAX_LEN];
bool bh[MAX_LEN+1],b2h[MAX_LEN+1];

void build_suffix_array() {
    N = strlen(A);

    memset(cont,0,sizeof(cont));

    for(int i = 0;i<N;++i) ++cont[A[i]];
    for(int i = 1;i<ALPH_SIZE;++i) cont[i] += cont[i-1];
    for(int i = 0;i<N;++i) pos[--cont[A[i]]] = i;

    for(int i = 0;i<N;++i) {
        bh[i] = (i==0 || A[pos[i]]!=A[pos[i-1]]);
        b2h[i] = false;
    }
}</pre>
```

#### 3. Cadenas

```
int KMP(string P, string T){
   int n = P.size(), L = T.size(), k = 0, ans = 0;
   for (int i=0; i<L; ++i) {</pre>
      while (k>0 \&\& P[k]!=T[i]) k = f[k-1];
      if(P[k]==T[i]) ++k;
      if(k==n){
         ++ans;
         k = f[k-1];
   return ans;
   for (int H = 1; H<N; H <<= 1) {</pre>
      int buckets = 0;
      for(int i = 0, j; i < N; i = j) {</pre>
         j = i+1;
         while(j<N && !bh[j]) ++j;
         next[i] = j;
         ++buckets;
      if(buckets==N) break;
      for(int i = 0; i<N; i = next[i]) {</pre>
         cont[i] = 0;
         for(int j = i; j<next[i];++j)</pre>
             rank[pos[j]] = i;
      ++cont[rank[N-H]];
      b2h[rank[N-H]] = true;
```

```
for(int i = 0; i < N; i = next[i]) {</pre>
          for (int j = i; j<next[i]; ++j) {</pre>
             int s = pos[j]-H;
             if(s>=0){
                int head = rank[s];
                rank[s] = head+cont[head];
                ++cont[head];
                b2h[rank[s]] = true;
          for (int j = i; j < next[i]; ++j) {</pre>
             int s = pos[j]-H;
             if(s>=0 && b2h[rank[s]]){
                for (int k = rank[s]+1;!bh[k] && b2h[k];++k)
                   b2h[k] = false;
      for (int i = 0; i < N; ++i) {</pre>
         pos[rank[i]] = i;
         bh[i] = b2h[i];
   for(int i = 0;i<N;++i) rank[pos[i]] = i;</pre>
int height[MAX_LEN];
// height[i] = lcp(pos[i],pos[i-1])
// Complejidad : O(n)
void getHeight() {
   height[0] = 0;
   for (int i = 0, h = 0; i < N; ++i) {</pre>
      if(rank[i]>0){
3.3. Trie.
const int ALPH_SIZE = 58;
```

```
#define LOG2 LEN 16
int RMQ[MAX_LEN][LOG2_LEN];
// Complejidad : O(nlgn)
void initialize_rmq() {
   for(int i = 0;i<N;++i) RMQ[i][0] = height[i];</pre>
   for (int j = 1; (1 << j) <= N; ++j) {
      for(int i = 0;i+(1<<j)-1<N;++i){</pre>
         if(RMQ[i][j-1]<=RMQ[i+(1<<(j-1))][j-1])</pre>
             RMQ[i][j] = RMQ[i][j-1];
             RMQ[i][j] = RMQ[i+(1<<(j-1))][j-1];
// lcp(pos[x],pos[y])
int lcp(int x, int y) {
   if(x==y) return N-rank[x];
   if(x>y) swap(x,y);
   int log = 0;
   while((1<<log)<=(y-x)) ++log;</pre>
   --log;
   return min(RMQ[x+1][log], RMQ[y-(1<<log)+1][log]);</pre>
```

int words; // numero de palabras que terminan en el nodo

int j = pos[rank[i]-1];

height[rank[i]] = h;

**if**(h>0) --h;

struct Node {

while (i+h<N && j+h<N && A[i+h]==A[j+h]) ++h;

// Queries para el Longest Common Prefix usando una Sparse Table.

```
int prefixes; // numero de palabras que tienen como prefijo el camino al nodo
   vector<Node*> links; // enlaces a los nodos hijos
   Node(){
      words = prefixes = 0;
      links.resize(ALPH_SIZE, NULL);
};
class Trie{
   public :
  Trie(){
      myRoot = new Node();
      myCount = 1;
   bool contains (const string& s) const;
   int nodeCount() const;
   int countWords(const string& s) const;
   int countPrefixes(const string& s) const;
   int countRepeated(Node* t) const;
   void printAllWords(const Node* t, const string& s) const;
   void insert(const string s);
   private :
   Node* myRoot; // raiz del trie
   int myCount; // # nodos del trie
bool Trie::contains(const string& s) const{
   Node* t = myRoot;
   int len = s.size();
   for (int k=0; k<len; ++k) {</pre>
         if(t==NULL) return false;
      t = t \rightarrow links[s[k] - 'A'];
   if(t==NULL) return false;
   return (t->words > 0);
int Trie::nodeCount() const{
```

```
return myCount;
int Trie::countWords(const string& s) const{
   int len = s.size();
   Node \star t = myRoot;
   for (int k=0; k<len; ++k) {</pre>
      if(t->links[s[k]-'A']==NULL) return 0;
      t = t \rightarrow links[s[k] - 'A'];
   return t->words;
int Trie::countPrefixes(const string& s) const{
   int len = s.size();
   Node \star t = myRoot;
   for (int k=0; k<len; ++k) {</pre>
      if(t->links[s[k]-'A']==NULL) return 0;
      t = t \rightarrow links[s[k] - 'A'];
   return t->prefixes;
void Trie::printAllWords(const Node* t = myRoot, const string& s = "") const{
   if(t->words > 0) cout<<s<<endl;</pre>
   for (int k=0; k<ALPH_SIZE; ++k)</pre>
      if(t->links[k]) printAllWords(t->links[k],s+char(k+'A'));
void Trie::insert(const string s){
   int len = s.size();
   Node * t = myRoot;
   for (int k=0; k<len; ++k) {
          if(t->links[s[k]-'A']==NULL) {
             t \rightarrow links[s[k] - 'A'] = new Node();
             ++myCount;
          t = t \rightarrow links[s[k] - 'A'];
          ++(t->prefixes);
```

```
++(t->words);
}
int Trie::countRepeated(Node* t = myRoot) const{
    int aux = 0;
    if((t->words)>1) ++aux;
    for(int k=0;k<ALPH_SIZE;++k)
    if(t->links[k]) aux += countRepeated(t->links[k]);
    return aux;
}
string test[] = {"tree","trie","algo","assoc","all","also"};
```

#### 3.4. Aho-Corasick.

```
struct No {
   int fail:
   vector< pair<int,int> > out; // num e tamanho do padrao
   //bool marc; // p/ decisao
   map<char, int> lista;
   int next; // aponta para o proximo sufixo que tenha out.size > 0
};
No arvore[1000003]; // quantida maxima de nos
//bool encontrado[1005]; // quantidade maxima de padroes, p/ decisao
int qtdNos, qtdPadroes;
// Funcao para inicializar
void inic() {
   arvore[0].fail = -1;
   arvore[0].lista.clear();
   arvore[0].out.clear();
   arvore[0].next = -1;
   qtdNos = 1;
   qtdPadroes = 0;
   //arvore[0].marc = false; // p/ decisao
   //memset(encontrado, false, sizeof(encontrado)); // p/ decisao
// Funcao para adicionar um padrao
void adicionar(char *padrao) {
   int no = 0, len = 0;
```

```
int main(){
  Trie* myTrie;
  myTrie = new Trie();
   for(int i=0;i<6;++i){</pre>
      myTrie->insert(test[i]);
      cout<<myTrie->nodeCount()<<endl;</pre>
      myTrie->printAllWords();
      cout << endl;
   delete myTrie;
   return 0;
   for (int i = 0 ; padrao[i] ; i++, len++) {
      if (arvore[no].lista.find(padrao[i]) == arvore[no].lista.end()) {
         arvore[qtdNos].lista.clear(); arvore[qtdNos].out.clear();
         //arvore[qtdNos].marc = false; // p/ decisao
         arvore[no].lista[padrao[i]] = qtdNos;
         no = qtdNos++;
      } else no = arvore[no].lista[padrao[i]];
   arvore[no].out.push_back(pair<int,int>(qtdPadroes++,len));
// Ativar Aho-corasick, ajustando funcoes de falha
void ativar() {
  int no, v, f, w;
   queue<int> fila;
   for (map<char,int>::iterator it = arvore[0].lista.begin();
      it != arvore[0].lista.end(); it++) {
      arvore[no = it->second].fail = 0;
      arvore[no].next = arvore[0].out.size() ? 0 : -1;
      fila.push(no);
   while (!fila.empty()) {
      no = fila.front(); fila.pop();
      for (map<char, int>::iterator it=arvore[no].lista.begin();
          it!=arvore[no].lista.end(); it++) {
         char c = it->first;
```

```
v = it->second;
fila.push(v);
f = arvore[no].fail;
while (arvore[f].lista.find(c) == arvore[f].lista.end()) {
    if (f == 0) { arvore[0].lista[c] = 0; break; }
    f = arvore[f].fail;
}
w = arvore[f].lista[c];
arvore[v].fail = w;
arvore[v].next = arvore[w].out.size() ? w : arvore[w].next;
}
}
}
// Buscar padroes no aho-corasik
void buscar(char *input) {
  int v, no = 0;
  for (int i = 0 ; input[i] ; i++) {
    while (arvore[no].lista.find(input[i]) == arvore[no].lista.end()) {
```

### 3.5. Rotación lexicográfica mínima.

```
char s[100001];
scanf("%s",s);

int N = strlen(s),ans = 0,p = 1,1 = 0;

while(p<N && ans+1+1<N)(
    if(s[ans+1]==s[(p+1)*N]) ++1;
    else if(s[ans+1]<s[(p+1)*N]){
        p = p+1+1;
    }
}</pre>
```

# 4. Geometría

# 4.1. Punto y Línea.

```
const double eps = 1e-9;
struct point{
   double x,y;

   point(){}
```

```
1 = 0;
}else{
    ans = max(ans+1+1,p);
    p = ans+1;
    1 = 0;
}
printf("%d\n",ans);
```

```
point (double _x, double _y) {
    x = _x; y = _y;
}

point operator + (const point &p) const{
    return point (x+p.x,y+p.y);
}
```

```
point operator - (const point &p) const{
    return point(x-p.x,y-p.y);
}

point operator * (double v) const{
    return point(x*v,y*v);
}

point perp() {
    return point(-y,x);
}

point normal() {
    return point(-y/abs(),x/abs());
}

double dot(const point &p) const{
    return x*p.x+y*p.y;
}

double abs2() const{
    return dot(*this);
```

# 4.2. Área y orientación de un triángulo.

```
double signed_area(const point &p1, const point &p2, const point &p3) {
    return (p1.x*p2.y+p2.x*p3.y+p3.x*p1.y-p1.y*p2.x-p2.y*p3.x-p3.y*p1.x)/2;
}
```

# 4.3. Fórmulas de triángulos.

```
double AreaHeron(double const &a, double const &b, double const &c) {
   double s=(a+b+c)/2;
   return sqrt(s*(s-a)*(s-b)*(s-c));
}

double Circumradius(const double &a, const double &b, const double &c) {
   return a*b*c/4/AreaHeron(a,b,c);
}
```

```
double abs() const{
      return sqrt(abs2());
  bool operator < (const point &p) const{</pre>
      if(fabs(x-p.x)>eps) return x<p.x;</pre>
      return y>p.y;
};
struct line{
  point p1,p2;
  line(){}
  line(point _p1, point _p2) {
     p1 = _p1; p2 = _p2;
     if(p1.x>p2.x) swap(p1,p2);
};
bool ccw(const point &p1, const point &p2, const point &p3) {
  return signed_area(p1,p2,p3)>-eps;
double Circumradius(const point &P1, const point &P2, const point &P3) {
  return (P2-P1).abs()*(P3-P1).abs()*(P3-P2).abs()/4/fabs(signed_area(P1,P2,P3));
double Inradius (const double &a, const double &b, const double &c) {
  return 2*AreaHeron(a,b,c)/(a+b+c);
```

#### 4.4. Orientación de un polígono.

bool in = 0;

```
//verdadero : sentido anti-horario, Complejidad : O(n)
                                                                                                     ind = i;
bool ccw(const vector<point> &poly) {
                                                                                                     x = poly[i].x;
   //primero hallamos el punto inferior ubicado ms a la derecha
                                                                                                     y = poly[i].y;
   int ind = 0,n = poly.size();
   double x = poly[0].x,y = poly[0].y;
                                                                                                  if (ind==0) return ccw(poly[n-1],poly[0],poly[1]);
   for (int i=1;i<n;i++) {</pre>
                                                                                                 return ccw(poly[ind-1],poly[ind],poly[(ind+1)%n]);
      if (poly[i].y>y) continue;
      if (fabs(poly[i].y-y)<eps && poly[i].x<x) continue;</pre>
4.5. Área con signo.
//valor positivo : vrtices orientados en sentido antihorario
                                                                                                     for(int i=1;i<=n;++i)</pre>
//valor negativo : vrtices orientados en sentido horario
double signed_area(const vector<point> &poly) {
                                                                                                           S += poly[i%n].x*(poly[(i+1)%n].y-poly[i-1].y);
      int n = poly.size();
                                                                                                     S /= 2;
      if(n<3) return 0.0;
                                                                                                     return S;
      double S = 0.0;
4.6. Punto dentro de un polígono.
bool PointInsideConvexPolygon(const point &P, vector<point> &poly) {
   int n = poly.size();
                                                                                                  for(int i = 0, j = n-1; i < n; j = i++) {</pre>
   if(!ccw(poly)) reverse(poly.begin(),poly.end());
                                                                                                     double dx = poly[j].x-poly[i].x;
                                                                                                     double dy = poly[j].y-poly[i].y;
   for (int i=1;i<=n;++i)</pre>
      if(!ccw(poly[i-1],poly[i%n],P))
                                                                                                     if((poly[i].y<=P.y+eps && P.y<poly[j].y) ||</pre>
         return false;
                                                                                                        (poly[j].y<=P.y+eps && P.y<poly[i].y))</pre>
                                                                                                        if(P.x-eps<dx*(P.y-poly[i].y)/dy+poly[i].x)</pre>
   return true;
                                                                                                           in \hat{}=1;
bool PointInsidePolygon(const point &P, const vector<point> &poly) {
                                                                                                  return in;
   int n = poly.size();
```

#### 4.7. Distancia desde un punto.

```
//Distancia de un punto a una recta infinita
double PointToLineDist(const point &P, const line &L) {
    return 2*fabs(signed_area(L.pl,L.p2,P))/(L.p2-L.pl).abs();
}

//Distancia de un punto a un segmento de recta
double PointToSegmentDist(const point &P, const line &L) {
    point v=L.p2-L.pl, w=P-L.pl;
```

#### 4.8. Intersección de líneas.

```
//verdadero : s se intersectan, I : punto de interseccin
bool lineIntersection(line &L1, line &L2, P &I) {
   point n = (L2.p2-L2.p1).perp();

   double denom = n.dot(L1.p2-L1.p1);
   if(fabs(denom)<eps) return false; // las rectas son paralelas</pre>
```

## 4.9. Convex Hull (Monotone Chain).

```
vector<point> ConvexHull(vector<point> P) {
    sort(P.begin(), P.end());
    int n = P.size(), k = 0;
    point H[2*n];

    for(int i=0;i<n;++i) {
        while(k>=2 && !ccw(H[k-2],H[k-1],P[i])) --k;
        H[k++] = P[i];
    }
}
```

#### 4.10. Teorema de Pick.

```
El Teorema de Pick nos dice que : A=I+B/2-1, donde,

A = Area de un poligono de coordenadas enteras

I = Nmero de puntos enteros en su interior

B = Nmero de puntos enteros sobre sus bordes
```

```
double aux1=w.dot(v);
   if(aux1<eps) return (P-L.p1).abs();</pre>
   double aux2=v.dot(v);
   if(aux2<=aux1+eps) return (P-L.p2).abs();</pre>
   return PointToLineDist(P,L);
   double t = n.dot(L2.p1-L1.p1)/denom;
   I = L1.p1 + (L1.p2-L1.p1) *t;
   return true;
   for (int i=n-2, t=k; i>=0; --i) {
      while (k>t \&\& !ccw(H[k-2],H[k-1],P[i])) --k;
      H[k++] = P[i];
   return vector<point> (H,H+k);
Haciendo un cambio en la frmula : I=(2A-B+2)/2, tenemos una forma de calcular
el numero de puntos enteros en el interior del poligono
int IntegerPointsOnSegment(const point &P1, const point &P2) {
   point P=P1-P2;
```

P.x=abs(P.x); P.y=abs(P.y);

```
if(P.x==0) return P.y;
if(P.y==0) return P.x;
return (__gcd(P.x,P.y));
```

## 4.11. Par de puntos más cercano.

```
#define MAX_N 100000
#define px second
#define py first
typedef pair<long long, long long> point;

int N;
point P[MAX_N];
set<point> box;

bool compare_x(point a, point b) { return a.px<b.px; }

inline double dist(point a, point b) {
   return sqrt((a.px-b.px)*(a.px-b.px)+(a.py-b.py)*(a.py-b.py));
}

double closest_pair() {
   if(N<=1) return -1;</pre>
```

# 4.12. Unión de rectángulos (Área).

```
#define MAX_N 10000

struct event{
   int ind;
   bool type;

   event(){};
   event(int ind, int type) : ind(ind), type(type) {};
};

struct point{
   int x,y;
};
```

Se asume que los vertices tienen coordenadas enteras. Sumar el valor de esta funcion para todas las aristas para obtener el numero total de punto en el borde del poligono.

```
sort(P,P+N,compare_x);

double ret = dist(P[0],P[1]);
box.insert(P[0]);

set<point> :: iterator it;

for(int i = 1,left = 0;i<N;++i) {
    while(left<i && P[i].px-P[left].px>ret) box.erase(P[left++]);
    for(it = box.lower_bound(make_pair(P[i].py-ret,P[i].px-ret));
        it!=box.end() && P[i].py+ret>=(*it).py;++it)
            ret = min(ret, dist(P[i],*it));
        box.insert(P[i]);
}

return ret;
}
```

```
int N;
point rects[MAX_N][2];
// rects[i][0] : esquina inferior izquierda
// rects[i][1] : esquina superior derecha
event events_v[2*MAX_N], events_h[2*MAX_N];
bool in_set[MAX_N];

bool compare_x(event a, event b) {
    return rects[a.ind][a.type].x<rects[b.ind][b.type].x;
}
bool compare_y(event a, event b) {
    return rects[a.ind][a.type].y<rects[b.ind][b.type].y;
}</pre>
```

```
long long union_area() {
   int e = 0;
   for (int i = 0; i < N; ++i) {</pre>
      events_v[e] = event(i,0);
      events_h[e] = event(i,0);
      events_v[e] = event(i,1);
      events_h[e] = event(i,1);
      ++e;
   sort(events_v,events_v+e,compare_x);
   sort(events_h, events_h+e, compare_y);
   memset(in_set, false, sizeof(in_set));
   in_set[events_v[0].ind] = true;
   long long area = 0;
   int prev_ind = events_v[0].ind, cur_ind;
   int prev_type = events_v[0].type, cur_type;
   for (int i = 1; i < e; ++i) {</pre>
      cur_ind = events_v[i].ind; cur_type = events_v[i].type;
      int cont = 0, dx = rects[cur_ind][cur_type].x-rects[prev_ind][prev_type].x;
```

# 5. Matemática

# 5.1. Algoritmo de Euclides.

```
struct EuclidReturn{
   int u,v,d;

   EuclidReturn(int _u, int _v, int _d) {
      u = _u; v = _v; d = _d;
   }
};

EuclidReturn Extended_Euclid(int a, int b) {
   if(b==0) return EuclidReturn(1,0,a);
   EuclidReturn aux = Extended_Euclid(b,a%b);
   int v = aux.u-(a/b)*aux.v;
   return EuclidReturn(aux.v,v,aux.d);
}
```

```
// ax = b (mod n)
int solveMod(int a,int b,int n) {
    EuclidReturn aux = Extended_Euclid(a,n);
    if(b%aux.d==0) return ((aux.u * (b/aux.d))%n+n)%n;
    return -1;// no hay solucuin
}

// ax = 1 (mod n)
int modular_inverse(int a, int n) {
    EuclidReturn aux = Extended_Euclid(a,n);
    return ((aux.u/aux.d)%n+n)%n;
}
```

## 5.2. Criba para la función phi de Euler.

```
fill(factors, factors+N+1,0);
phi[1] = 1;

for(int i = 2;i<=N;i++){
   if(factors[i]==0){
     factors[i] = i;
     phi[i] = i-1;

   if(i<=sqrt(N)) for(int j = i*i;j<=N;j += i) factors[j] = i;
}else{
   int aux = i,exp = 0;</pre>
```

#### 5.3. Teorema chino del resto.

```
// rem y mod tienen el mismo nmero de elementos
long long chinese_remainder(vector<int> rem, vector<int> mod) {
   long long ans = rem[0],m = mod[0];
   int n = rem.size();

   for(int i=1;i<n;++i) {
     int a = modular_inverse(m,mod[i]);
}</pre>
```

#### 5.4. Número combinatorio.

```
long long comb(int n, int m) {
   if(m>n-m) m = n-m;

long long C = 1;
   //C^{n}_{ii} -> C^{n}_{ii+1}
   for(int i=0;i<m;++i) C = C*(n-i)/(1+i);
   return C;
}

Cuando n y m son grandes y se pide comb(n,m)%MOD, donde MOD es un numero primo, se puede usar el Teorema de Lucas.

#define MOD 3571

int C[MOD][MOD];</pre>
```

```
while (aux%factors[i] == 0) {
         aux /= factors[i];
          ++exp;
      phi[i] = 1;
      for(int j = 0; j < exp; ++j) phi[i] *= factors[i];</pre>
      phi[i] -= phi[i]/factors[i];
      phi[i] *= phi[aux];
      int b = modular_inverse(mod[i],m);
      ans = (ans*b*mod[i]+rem[i]*a*m)%(m*mod[i]);
      m *= mod[i];
   return ans;
void FillLucasTable() {
   memset(C, 0, sizeof(C));
   for(int i=0;i<MOD;++i) C[i][0] = 1;</pre>
   for(int i=1;i<MOD;++i) C[i][i] = 1;</pre>
   for (int i=2;i<MOD;++i)</pre>
      for (int j=1; j<i; ++j)</pre>
         C[i][j] = (C[i-1][j]+C[i-1][j-1])%MOD;
int comb(int n, int k){
   long long ans = 1;
   while (n!=0) {
      int ni = n%MOD,ki = k%MOD;
      n /= MOD; k /= MOD;
```

```
ans = (ans*C[ni][ki])%MOD;
}
```

## 5.5. Test de Miller-Rabin.

```
typedef unsigned long long ULL;
ULL mulmod(ULL a, ULL b, ULL c) {
    ULL x = 0,y = a%c;
```

```
while (b>0) {
    if (b&1) x = (x+y)%c;
    y = (y<<1)%c;
    b >>= 1;
}

return x;
}

ULL pow(ULL a, ULL b, ULL c) {
    ULL x = 1, y = a;

while (b>0) {
    if (b&1) x = mulmod(x,y,c);
    y = mulmod(y,y,c);
    b >>= 1;
```

#### 5.6. Polinomios.

return x;

```
vector<int> add(vector<int> &a, vector<int> &b) {
   int n = a.size(),m = b.size(),sz = max(n,m);
   vector<int> c(sz,0);

for(int i = 0;i<n;++i) c[i] += a[i];
   for(int i = 0;i<m;++i) c[i] += b[i];

while(sz>1 && c[sz-1]==0) {
    c.pop_back();
```

```
return (int) ans;
bool miller_rabin(ULL p, int it) {
  if(p<2) return false;</pre>
  if(p==2) return true;
  if((p&1)==0) return false;
  ULL s = p-1;
  while(s%2==0) s >>= 1;
  while (it--) {
     ULL a = rand()%(p-1)+1, temp = s;
     ULL mod = pow(a,temp,p);
      if (mod==-1 || mod==1) continue;
      while(temp!=p-1 && mod!=p-1){
        mod = mulmod(mod, mod, p);
         temp <<= 1;
      if (mod!=p-1) return false;
   return true;
      --sz;
   return c;
vector<int> multiply(vector<int> &a, vector<int> &b){
  int n = a.size(), m = b.size(), sz = n+m-1;
  vector<int> c(sz,0);
```

```
for(int i = 0;i<n;++i)
    for(int j = 0;j<m;++j)
        c[i+j] += a[i]*b[j];

while(sz>1 && c[sz-1]==0) {
    c.pop_back();
    --sz;
}

return c;
```

#### 5.7. Fast Fourier Transform.

```
#define lowbit(x) (((x) ^(x-1)) & (x)
typedef complex<long double> Complex;
void FFT(vector<Complex> &A, int s){
   int n = A.size(), p = 0;
   while (n>1) {
      ++p;
      n >>= 1;
   n = (1 << p);
   vector<Complex> a = A;
   for (int i = 0; i < n; ++i) {</pre>
      int rev = 0;
      for(int j = 0; j<p; ++j) {
         rev <<= 1;
         rev |= ((i >> j) \& 1);
      A[i] = a[rev];
   Complex w,wn;
   for (int i = 1; i <= p; ++i) {</pre>
      int M = (1 << i), K = (M >> 1);
      wn = Complex(cos(s*2.0*M_PI/(double)M), sin(s*2.0*M_PI/(double)M));
```

```
bool is_root(vector<int> &P, int r) {
   int n = P.size();
  long long y = 0;
   for (int i = 0; i < n; ++i) {</pre>
      if (abs (y-P[i])%r!=0) return false;
      y = (y-P[i])/r;
   return y==0;
      for(int j = 0; j<n; j += M) {
         w = Complex(1.0, 0.0);
         for (int 1 = \dot{7}; 1 < K + \dot{7}; ++1) {
             Complex t = w;
             t *= A[1 + K];
             Complex u = A[1];
            A[1] += t;
             u -= t;
            A[1 + K] = u;
             w \star = wn;
   if(s==-1){
      for (int i = 0; i < n; ++i)</pre>
         A[i] /= (double)n;
vector<Complex> FFT_Multiply(vector<Complex> &P, vector<Complex> &Q){
   int n = P.size()+O.size();
   while(n!=lowbit(n)) n += lowbit(n);
  P.resize(n,0);
  Q.resize(n,0);
  FFT (P, 1);
  FFT(Q,1);
```

```
vector<Complex> R;
for(int i=0;i<n;i++) R.push_back(P[i]*Q[i]);
FFT(R,-1);
return R;</pre>
```

```
}
// Para multiplicacin de enteros grandes
const long long B = 100000;
const int D = 5;
```

#### 6. Estructuras de datos

## 6.1. **BIT.**

```
#define MAX_SIZE 20001
//los indices que se pueden usar van desde 1 hasta MAX_SIZE-1

void update(long long T[], int idx, int val) {
   for(;idx<MAX_SIZE;idx+=(idx & -idx)) T[idx]+=val;
}

long long f(long long T[], int idx) {
   long long sum = T[idx];

   if(idx>0) {
      int z = idx-(idx & -idx);
      --idx;
   }
```

## 6.2. Range Minimum Query.

```
#define MAX_N 100000
#define LOG2_MAXN 16
long long A[MAX_N];
int N,ind[(1<<(LOG2_MAXN+2))];

void initialize(int node, int s, int e) {
   if(s==e) ind[node] = s;
   else{
      initialize(2*node+1, s, (s+e)/2);
      initialize(2*node+2, (s+e)/2+1, e);
      if(A[ind[2*node+1]]<=A[ind[2*node+2]]) ind[node] = ind[2*node+1];
      else ind[node] = ind[2*node+2];
   }
}</pre>
```

```
while(idx!=z) {
    sum -= T[idx];
    idx -= (idx & -idx);
}

return sum;
}

long long F(long long T[], int idx) {
    long long sum = 0;
    for(;idx>0;idx -= (idx & -idx)) sum += T[idx];
    return sum;
}
```

```
int query(int node, int s, int e, int a, int b) {
   if(b<s || a>e) return -1;
   if(a<=s && e<=b) return ind[node];

int ind1 = query(2*node+1,s,(s+e)/2,a,b);
   int ind2 = query(2*node+2,(s+e)/2+1,e,a,b);

if(ind1==-1) return ind2;
   if(ind2==-1) return ind1;
   if(A[ind1]<=A[ind2]) return ind1;
   return ind2;
}</pre>
```

#### 6.3. Lowest Common Ancestor.

```
#define MAX N 100000
#define LOG2_MAXN 16
// NOTA : memset(parent,-1,sizeof(parent));
int N, parent[MAX_N], L[MAX_N];
int P[MAX_N][LOG2_MAXN];
int get_level(int u) {
   if(L[u]!=-1) return L[u];
   else if(parent[u]==-1) return 0;
   return 1+get_level(parent[u]);
void init(){
      memset(L,-1,sizeof(L));
      for(int i = 0;i<N;++i) L[i] = get_level(i);</pre>
      memset (P,-1, sizeof (P));
       for(int i = 0; i < N; ++i) P[i][0] = parent[i];</pre>
      for(int \dot{j} = 1; (1 << \dot{j}) < N; + + \dot{j})
          for (int i = 0; i < N; ++i)</pre>
                if(P[i][j-1]!=-1)
                       P[i][j] = P[P[i][j-1]][j-1];
```

# 6.4. Maximum Sum Segment Query.

```
#define MAX_N 100000
#define LOG2_MAXN 16
const long long INF = 1000000001LL;

int N,a[MAX_N];
long long c[MAX_N+1],int_min[1<<(LOG2_MAXN+2)],int_max[1<<(LOG2_MAXN+2)];
long long int_best[1<<(LOG2_MAXN+2)];

void build_tree(int node, int lo, int hi) {
    if(lo==hi) {
        int_min[node] = c[lo-1];
        int_max[node] = c[lo];
        int_best[node] = c[lo]-c[lo-1];</pre>
```

```
int LCA(int p, int q) {
    if(L[p]<L[q]) swap(p,q);

int log = 1;
    while((1<<log)<=L[p]) ++log;
    --log;

for(int i = log;i>=0;--i)
    if(L[p]-(1<<i)>=L[q])
        p = P[p][i];

if(p==q) return p;

for(int i = log;i>=0;--i) {
    if(P[p][i]!=-1 && P[p][i]!=P[q][i]) {
        p = P[q][i];
        q = P[q][i];
    }
}

return parent[p];
}
```

```
void init(){
   c[0] = 0;
   for(int i = 0;i<N;++i) c[i+1] = c[i]+a[i];</pre>
   build_tree(0,0,N);
long long minPrefix;
int s,e;
long long tree_query(int node, int lo, int hi) {
   if (s<=lo && hi<=e) {
      long long ret = int_best[node];
      if (minPrefix!=INF) ret = max(ret,int_max[node]-minPrefix);
      minPrefix = min(minPrefix,int_min[node]);
      return ret;
   }else{
6.5. Treap.
long long seed = 47;
long long rand() {
   seed = (seed * 279470273) % 4294967291LL;
   return seed;
typedef int treap_type;
class treap{
   public:
   treap_type value;
   long long priority;
   treap *left, *right;
   int sons;
   treap(treap_type value) : left(NULL), right(NULL), value(value), sons(0){
      priority = rand();
   ~treap(){
      if(left) delete left;
```

```
int mi = (lo+hi)>>1;
      if (mi<s) return tree_query(2*node+2, mi+1, hi);</pre>
      else if(mi>=e) return tree_query(2*node+1,lo,mi);
         long long val1 = tree_query(2*node+1,lo,mi);
        long long val2 = tree_query(2*node+2,mi+1,hi);
         return max(val1,val2);
// Los indices van de 1 a N
long long solve_msq(int x, int y) {
  minPrefix = INF;
  s = x; e = y;
  return tree_query(0,0,N);
      if(right) delete right;
};
treap* find(treap* t, treap_type val){
  if(!t) return NULL;
  if(val == t->value) return t;
  if(val < t->value) return find(t->left, val);
  if(val > t->value) return find(t->right, val);
inline void rotate_to_right(treap* &t) {
  treap* n = t->left;
  t->left = n->right;
  n->right = t;
  t = n;
inline void rotate_to_left(treap* &t){
  treap* n = t->right;
  t->right = n->left;
  n->left = t:
```

```
t = n;
void fix_augment(treap* t){
   if(!t) return;
   t->sons = (t->left ? t->left->sons + 1 : 0) +
      (t->right ? t->right->sons + 1 : 0);
void insert(treap* &t, treap type val){
   if(!t) t = new treap(val);
   else insert(val <= t->value ? t->left : t->right, val);
   if(t->left && t->left->priority > t->priority)
      rotate_to_right(t);
   else if(t->right && t->right->priority > t->priority)
      rotate_to_left(t);
   fix_augment(t->left); fix_augment(t->right); fix_augment(t);
inline long long get_priority(treap* t){
   return t ? t->priority : -1;
void erase(treap* &t, treap_type val) {
   if(!t) return;
```

## 7.1. Exponenciación de matrices.

```
#define MAX_SIZE 100
#define MOD 10000
int size;

struct M{
   long long X[MAX_SIZE][MAX_SIZE];
   M() {}
}M0,aux1,aux2;

void mult(M &m, M &m1, M &m2) {
   memset(m.X,0,sizeof(m.X));
```

```
if(t->value != val) erase(val < t->value ? t->left : t->right, val);
     if(!t->left && !t->right){
         delete t;
        t = NULL;
      }else{
         if(get_priority(t->left) < get_priority(t->right))
            rotate_to_left(t);
            rotate_to_right(t);
         erase(t, val);
   fix_augment(t->left); fix_augment(t->right); fix_augment(t);
int getKth(treap* &t, int K) {
   int left = (t->left==NULL? 0 : 1+t->left->sons);
   int right = (t->right==NULL? 0 : 1+t->right->sons);
   if(1+left==K) return t->value;
  else if(left<K) return getKth(t->right,K-1-left);
   return getKth(t->left,K);
```

## 7. Matrices

```
aux1=exp(n/2);
for(int i=0;i<size;i++) fill(aux2.X[i],aux2.X[i]+size,0);
for(int i=0;i<size;i++) aux2.X[i][i]=1;

mult(aux2,aux1,aux1);

if(n%2==1) {
    mult(aux1,aux2,M0);
    return aux1;
}

return aux2;
}

// para exponente n escrito en base 2<=b<=10

M exp(string &n, int b) {
    M P[b+1];</pre>
```

#### 7.2. Determinante.

```
#define MAX_SIZE 500
int size;

struct M{
    double X[MAX_SIZE][MAX_SIZE];
    M() {}
};

const double eps = 1e-7;

double determinant (M M0) {
    double ans = 1, aux;
    bool found;

for(int i = 0,r = 0;i<size;i++) {
    found = false;

    for(int j = r;j<size;j++)
        if(fabs(M0.X[j][i])>eps) {
        found = true;
    }
}
```

```
for(int i=0;i<=b;++i) P[i] = exp(i);</pre>
int L = n.size();
M ret;
memset(ret.X,0,sizeof(ret.X));
for(int i=0;i<size;++i) ret.X[i][i] = 1;</pre>
int aux = 0;
for (int i=0; i<L; ++i) {</pre>
   int x = n[i] - '0';
   M0 = ret;
   ret = exp(b);
   aux1 = ret;
   mult(ret,aux1,P[x]);
return ret;
          if(j>r) ans = -ans;
          else break;
          for(int k = 0; k < size; k++) swap(M0.X[r][k], M0.X[j][k]);</pre>
          break;
   if (found) {
      for (int j = r+1; j < size; j++) {</pre>
          aux = M0.X[j][i]/M0.X[r][i];
          for(int k = i; k < size; k++) M0.X[j][k] -= aux*M0.X[r][k];
      r++;
    }else return 0;
for(int i = 0;i<size;i++) ans *= M0.X[i][i];</pre>
return ans;
```

## 7.3. Elimación gaussiana módulo MOD.

```
#define MAX_R 500
#define MAX_C 500
int R, C, MOD;
struct M{
   int X[MAX_R][MAX_C];
   M(){}
};
//cuidado con overflow
int exp(int a, int n) {
   if(n==0) return 1;
   if(n==1) return a;
   int aux=exp(a,n/2);
   if(n&1) return ((long long)a*(aux*aux)%MOD)%MOD;
   return (aux*aux)%MOD;
void GaussianElimination(M &M0){
   int aux:
   bool found;
   for(int I = 0, r = 0; r < R && i < C; ++i) {</pre>
      found=false;
      for(int j = r; j<R; ++j) {</pre>
         if(M0.X[j][i]>0){
             found=true;
```

#### 8. Mathematical facts

## 8.1. Números de Catalan. están definidos por la recurrencia:

$$C_{n+1} = \sum_{i=0}^{n} C_i C_{n-i}$$

Una fórmula cerrada para los números de Catalán es:

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n+1}$$

8.2. Números de Stirling de primera clase. son el número de permutaciones de n elementos con exactamente k ciclos disjuntos.

8.3. Números de Stirling de segunda clase. son el número de formas de dividir n elementos en k conjuntos.

$${n \brace k} = k {n-1 \brace k} + {n-1 \brace k-1}$$

Además:

$${n \brace k} = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} {k \choose j} j^n$$

8.4. **Números de Bell.** cuentan el número de formas de dividir n elementos en subconjuntos.

$$\mathcal{B}_{n+1} = \sum_{k=0}^{n} \binom{n}{k} \mathcal{B}_k$$

X	5	6	7	8	9	10	11	12
$\mathcal{B}_x$	52	203	877	4.140	21.147	115.975	678.570	4.213.597

8.5. Funciones generatrices. Una lista de funciones generatrices para secuencias útiles:

$(1,1,1,1,1,1,\ldots)$	$\frac{1}{1-z}$
$(1,-1,1,-1,1,-1,\ldots)$	$\frac{1}{1+z}$
$(1,0,1,0,1,0,\ldots)$	$\frac{1}{1-z^2}$
$(1,0,\ldots,0,1,0,1,0,\ldots,0,1,0,\ldots)$	$\frac{1}{1-z^2}$
$(1,2,3,4,5,6,\ldots)$	$\frac{1}{(1-z)^2}$
$(1, \binom{m+1}{m}, \binom{m+2}{m}, \binom{m+3}{m}, \dots)$	$\frac{1}{(1-z)^{m+1}}$
$(1,c,\binom{c+1}{2},\binom{c+2}{3},\ldots)$	$\frac{1}{(1-z)^c}$
$(1,c,c^2,c^3,\ldots)$	$\frac{1}{1-cz}$
$(0,1,\frac{1}{2},\frac{1}{3},\frac{1}{4},\ldots)$	$\ln \frac{1}{1-z}$

Truco de manipulación:

$$\frac{1}{1-z}G(z) = \sum_{n} \sum_{k \le n} g_k z^n$$

8.6. The twelvefold way. ¿Cuántas funciones  $f: N \to X$  hay?

N	X	Any $f$	Injective	Surjective
dist.	dist.	$x^n$	$(x)_n$	$x!\binom{n}{x}$
indist.	dist.	$\binom{x+n-1}{n}$	$\binom{x}{n}$	$\binom{n-1}{n-x}$
dist.	indist.	$\binom{n}{1} + \ldots + \binom{n}{x}$	$[n \le x]$	$\binom{n}{k}$
indist.	indist.	$p_1(n) + \dots p_x(n)$	$[n \le x]$	$p_x(n)$

Where  $\binom{a}{b} = \frac{1}{b!}(a)_b$  and  $p_x(n)$  is the number of ways to partition the integer n using x summands.

8.7. **Teorema de Euler.** si un grafo conexo, plano es dibujado sobre un plano sin intersección de aristas, y siendo v el número de vértices, e el de aristas y f la cantidad de caras (regiones conectadas por aristas, incluyendo la región externa e infinita), entonces

$$v - e + f = 2$$

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