SFJ: An implementation of Semantic Featherweight Java

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Abstract. The abstract should briefly summarize the contents of the paper in 15-250 words.

Keywords: nominal and structural subtyping \cdot Featherweight Java \cdot object-oriented languages \cdot semantic subtyping \cdot type theory.

1 Introduction

A typing system for a programming language is a set of deduction rules that allow

$$\tau ::= \alpha \mid \mu \tag{1}$$

2 Semantic Subtyping for Featherweight Java

3 SFJ: Design and Implementation

Since we want our types to represent sets of values, we restrict our types to finite trees whose leaves are constants with no cycles. For example a recursive type $\alpha = [a : \alpha]$ would be an infinite tree **new** $C(\mathbf{new}\ C(\cdots))$. Similarly the types $\alpha = [b : \beta]$, $\beta = [a = \alpha]$ would also be impossible to instantiate.

Therefore, to not allow this, when processing the AST of a program, we do not allow the type of the class we are defining to be a field and we mark any classes with only basic types in it's fields as resolved and otherwise if it contains a class type, we mark it as unresolved. After processing the whole AST, we perform the following algorithm to decide whether the type definitions in the program are valid.

boolean resolutionOccured = false
do
 for class that is unresolved:
 boolean resolved = true

for field that contains a class type:
 if class type is not resolved

resolved = false

Given that we now know that all the types in the program are valid, we can create the subtyping relation which will be a map of types to a set of its subtypes. For every program, we assume the initial subtyping relation for all of our basic types. We would like to highlight that Int is not a subtype of Float as it cannot represent the whole domain of Int accurately therefore the domain is not fully contained in Float. Similarly so for Double and Long.

```
Boolean = \{Boolean\}
Double = \{Double, Float, Int, Short, Byte\}
Float = \{Float, Short, Byte\}
Long = \{Long, Int, Short, Byte\}
Int = \{Int, Short, Byte\}
Short = \{Short, Byte\}
Byte = \{Byte\}
```

Knowing that all our types are finite trees with leaves as constants and given this initial relation, we can now confidently create a subtyping relation for all class types using the following algorithm.

```
function generateRelation(classes):
   List<class> untyped = []

for class in classes:
    if addClass(class) is false:
        untyped.add(class)

if untyped is not []:
        generateRelation(untyped)

function addClass(class):
   for existing class type in relation:
        if tryAddSubtype(class, existingClass) is false:
            return false

        tryAddSubtype(existingClass, class)
```

```
add class to its own subtype relation
function tryAddSubtype(class, other):
    boolean flag = true
    for field in class:
        if field contains type not in relation:
            return false
        if other does not have field:
            flag = false
        else:
            if other.field.types not fully contains field.types:
                flag = false
    for method in class:
        if method contains type not in relation:
            return false
        if other does not have method:
            flag = false
        else:
            if other.method.types not fully contains method.types:
                flag = false
    if flag == true:
        add class to other subtype relation
```

The following bibliography provides a sample reference list with entries for journal articles [1].

Related Work and Conclusion

References

1. Dardha, O., Gorla, D., Varacca, D.: Semantic Subtyping for Objects and Classes. Computer Journal 60(5), 636-656 (2017). https://doi.org/10.1093/comjnl/bxw080