

$$\textcircled{1} \quad T(n) = 3T(n/2) + n^2.$$

$$\boxed{a=3} \quad \boxed{b=2} \quad \boxed{k=2} \quad \boxed{p=0}$$

$$b^k = 2^2 = 4$$

$$a < b^k \rightarrow \text{Case (3).}$$

$$\boxed{p=0} \rightarrow T(n) = O(n^k \log^p n)$$

$$\rightarrow O(n^2 \log^0 n)$$

$$\rightarrow \boxed{O(n^2)}$$

$$\textcircled{2} \quad T(n) = 7T(n/2) + n^2.$$

$$\boxed{a=7} \quad \boxed{b=2} \quad \boxed{k=2} \quad \boxed{p=0}$$

$$b^k = 2^2 = 4$$

$$a > b^k \rightarrow 7 > 4.$$

$a > b^k \rightarrow \text{case (1)}$

$$T(n) = O(n^{\log_b a})$$

$$T(n) = O(n^{\log_2 7})$$

$$= O(n^{2.801})$$

$$\boxed{T(n) = O(n^3)}$$

③ $T(n) = 4T(n/2) + n^2$

$$\boxed{a=4} \quad \boxed{b=2} \quad \boxed{k=2} \quad \boxed{p=0}$$

$$b^k = 2^2 = 4.$$

$a = b^k \rightarrow \text{case : (2)}$

$$p=0 \rightarrow T(n) = O(n^{\log_b a} \cdot \log^{p+1} n)$$

$$T(n) = O(n^{\log_2 4} \cdot \log^1 n)$$

$$T(n) = O(n^2 \cdot \log^1 n) \rightarrow \boxed{O(n^2 \log n)}$$

$$(4) \quad T(n) = 3T(n/4) + n \log n.$$

$$\boxed{a=3} \quad \boxed{b=4} \quad \boxed{k=1} \quad \boxed{p=1}$$

$$a < b^k \rightarrow \text{case: (3)}$$

$$\boxed{p=1} \rightarrow T(n) = O(n^k \log^p n)$$

$$T(n) = O(n^1 \log^1 n)$$

$$\boxed{T(n) = O(n \log n)}$$

$$(5) \quad T(n) = 4T(n/2) + \log n.$$

$$\boxed{a=4} \quad \boxed{b=2} \quad \boxed{k=1} \quad \boxed{p=1}$$

$$a > b^k \rightarrow \text{case (1)}$$

$$T(n) = O(n \log_b a)$$

$$T(n) = O(n \log_2 4) = \boxed{O(n^2)}$$

⑥ $T(n) = T(n-1) + n$

Hence $a=1$ $b=1$ $k=1$ $p=0$

$b > 1 \rightarrow$ For master's theorem.

\rightarrow So it doesn't applicable for master's theorem.

⑦ $T(n) = 4T(n/2) + n^2 \log n$.

$a=4$ $b=2$ $k=2$ $p=1$

$$b^k = 2^2 = 4$$

$a = b^k \rightarrow$ case (2).

$$\begin{aligned} T(n) &= O(n^{\log_b a} \cdot \log^{p+1} n) \\ &= O(n^{\log_2 4} \cdot \log^2 n) \end{aligned}$$

$$T(n) = O(n^2 \cdot \log^2 n)$$

$$(8) \quad 5T(n/2) + n^2 \log n$$

$$\boxed{a=5} \quad \boxed{b=2} \quad \boxed{k=2} \quad \boxed{p=1}$$

$$2^2 \leq 5 \rightarrow \text{case (1)} \rightarrow a > b^k$$

$$\boxed{p=1} \rightarrow T(n) = O(n^k \log_b a)$$

$$= O(n^2 \log_2 5)$$

$$\boxed{T(n) = O(n \log_2 5)}$$

$$(9) \quad T(n) = 3T(n/3) + n/\log n$$

$$\boxed{a=3} \quad \boxed{b=3}$$

\rightarrow Here $f(n) = n/\log n$ is not polynomially smaller than $n \log_3 3^{-\epsilon}$ for any $\epsilon > 0$.

\rightarrow So m-T not applicable.

$$(i) \quad T(n) = 2T(n/4) + C$$

$$\boxed{a=2} \quad \boxed{b=4} \quad \boxed{k=0}$$

$$b^k = 4^0 \geq 1 \quad 1 < 2 \quad b^k < a \text{ (or) } a > b^k$$

$$a^0 > b^k \rightarrow \text{case (i)}$$

$$T(n) = \Theta(n \log_b a)$$

$$= \Theta(n \log_4 2)$$

$$= \Theta(n \log 4^2) \rightarrow \Theta(n^{0.5})$$

$$\boxed{T(n) = \Theta(n^{1/2})}$$

$$(ii) \quad T(n) = T(n/4) + \log n.$$

$$\boxed{a=1} \quad \boxed{b=4} \quad \boxed{k=0} \quad \boxed{p=1}$$

$$b^k = 4^0 = 1 \quad \boxed{a=1} \quad \boxed{b^k=1}$$

$a \leq b^k \rightarrow \text{case: (2)}$

$$\boxed{p=1} \rightarrow T(n) = O(n^k \log^p n)$$

$$T(n) = O(n^{\log_b a} \log^{p+1} n)$$

$$= O(n^{\log_4 1} \log^2 n)$$

$$= O(n^0 \log^2 n)$$

$$\boxed{T(n) = O(\log^2 n)}$$

$$(2) \quad T(n) = T(n/2) + T(n/4) + n^2$$

\rightarrow Master theorem doesn't apply.

\rightarrow Recursion Tree $T(n) = O(n^2)$.

$$(3) \quad 2T(n/4) + \log n.$$

$$\boxed{a=2}$$

$$\boxed{b=4}$$

$$\boxed{k=0}$$

$$\boxed{p=1}$$

$$b^k = (4)^0 = 1$$

$$b^k \leq a = a > b^k \rightarrow \text{case ①.}$$

$$T(n) = O(n \log_b^a)$$

$$T(n) = O(n \log 4^2)$$

$$T(n) = O(n^{\frac{1}{2}})$$

$$(14) \quad T(n) = 3T(n/3) + n \log n.$$

$$a=3$$

$$b=3$$

$$k=1$$

$$p=1$$

$$b^k = 3^1 = 3$$

$$a = k^k$$

$$\rightarrow \text{case ②.}$$

$$[p=1] \rightarrow T(n) = O(n^{\log_b a} \cdot \log^{p+1} n)$$

$$T(n) = O(n \log 3^3 \cdot \log^2 n).$$

$$T(n) = O(n \cdot \log^2 n)$$

$$T(n) = O(n \cdot \log^2 n)$$

$$(15) \quad T(n) = 8T(n - \sqrt{n})/4 + n^2$$

$$a=8 \quad b=4 \quad k=0 \quad p=2$$

Master theorem doesn't apply.

$$(16) \quad T(n) = 2T(n/4) + \sqrt{n}$$

$$a=2 \quad b=4 \quad k=0 \quad p=\frac{1}{2}$$

$$a=2 \rightarrow b^k = 4^0 = 1$$

$$a > b^k \rightarrow O(n \log_b a) \rightarrow \text{case 1}$$

$$T(n) = O(n \log_4 2)$$

$$= O(n^{0.5})$$

$$T(n) = O(n)$$

$$(17) \quad T(n) = 2T(n/4) + n^{0.51}$$

$$\boxed{a=2} \quad \boxed{b=4} \quad \boxed{k=0.51} \quad \boxed{p=0}$$

$$a \neq 2$$

$$b^k = 4^{0.51} \rightarrow 2.02791$$

$$b^k > a$$

$$a < b^k \Rightarrow \text{case (3)}$$

$$\boxed{p=0} \rightarrow T(n) = O(n^k \log^p n)$$

$$T(n) = O(n^{0.51} \log^0 n)$$

$$\boxed{T(n) = O(n^{0.51})}$$

$$(18) \quad T(n) = 16T(n/4) + n!$$

$$\boxed{a=16} \quad \boxed{b=4} \quad \boxed{f(n) = n!}$$

$$\log_b a = \log_4 16 = 2.$$

$$n^{\log_b a} = n^2$$

Third case of the theorem:

$$f(n) = \Omega(n^{\log_a b + \epsilon}) \rightarrow c > 0$$

$$T(n) = \Theta(f(n))$$

$$n! = \Omega(n^{2+\epsilon})$$

$$16 f(n/4) \leq c f(n)$$

$$16 (n/4)! \leq n!$$

$\boxed{c=0.5} \rightarrow$ Then the condition satisfies

So case (3) $\Rightarrow T(n) = \Theta(n!)$

(19) $T(n) = 3T(n/2) + n$

$$\boxed{a=3}$$

$$\boxed{b=2}$$

$$\boxed{k=0}$$

$$\boxed{p=1}$$

$$b^k = 2^0 = 1$$

$$a > b^k \rightarrow \text{case (1)}$$

$$T(n) = \Theta(n \log_b a)$$

$$T(n) = \Theta(n \log_2 3)$$

$$T(n) = \Theta(n^{\log_2 3})$$

$$(20) \quad T(n) = 4T(n/2) + cn.$$

$$a=4$$

$$b=2$$

$$p=1$$

$$k=0$$

$$b^k = 2^0 = 1$$

$$a > b^k \rightarrow \text{case (1)}$$

$$T(n) = \Theta(n \log_b a)$$

$$= \Theta(n^{\log_2 4})$$

$$T(n) = \Theta(n^2)$$

$$(21) \quad T(n) = 3T(n/3) + n/2$$

$$\boxed{a=3} \quad \boxed{b=3} \quad \boxed{k=0} \quad \boxed{p=0}$$

$$b^k = 3^0 = 3$$

$$\boxed{a \leq b^k} \Rightarrow \text{case 2}$$

$$\boxed{p=0} \rightarrow T(n) = O(n^{\log_b a} \cdot \log^{p+1} n)$$

$$T(n) = O(n^{\log_3 3} \cdot \log^1 n)$$

$$= O(n^1 \log n)$$

$$\boxed{T(n) = O(n \log n)}$$

$$(22) \quad T(n) = 4T(n/2) + n/\log n$$

$$\boxed{a=4} \quad \boxed{b=2} \quad \boxed{k=1} \quad \boxed{p=0}$$

$$b^k = 2^1 = 2$$

$a > b^4 \rightarrow \text{case (1)}$

$$\rightarrow O(n \log_b a)$$

$$\rightarrow O(n \log_2 4)$$

$$\rightarrow O(n^2)$$

$$T(n) = O(n^2)$$

$$(23) \quad T(n) = 7T(n/3) + n^2$$

$$a=7$$

$$b=3$$

$$k=2$$

$$p=0$$

$$a=7, \quad b^k = 3^2 = 9.$$

$$a < b^k \rightarrow \text{case : (3)}$$

$$[p=0] \rightarrow O(n^{\cancel{k}} \log^p n)$$

$$T(n) = O(n^k \log^p n)$$

$$T(n) = O(n^2 \cdot \log^0 2) \rightarrow O(n^2)$$

$$(24) \quad T(n) = 8T(n/3) + 2^n.$$

$$\boxed{a=8} \quad \boxed{b=3} \quad f(n) = 2^n.$$

$$8 f(n/3) \triangleq c(2^n)$$

$$8(2^{n/3}) \triangleq c(2^n)$$

$$\boxed{c=0.5}$$

$$8(2^{n/3}) \triangleq 2^n$$

Case - (3)

$$\boxed{T(n) = O(2^n)}$$

$$(25) \quad T(n) = 16T(n/4) + n$$

$$\boxed{a=16} \quad \boxed{b=4} \quad \boxed{k=1} \quad \boxed{p=0}$$

$$b^k = 4^1 = 4$$

$$\rightarrow a > b^k = O(n \log_b a) \rightarrow \text{case ①}$$

$$\rightarrow O(n \log_4 16)$$

$$\rightarrow \boxed{T(n) = O(n^2)}$$