

# Linked List

- Karun Karthik

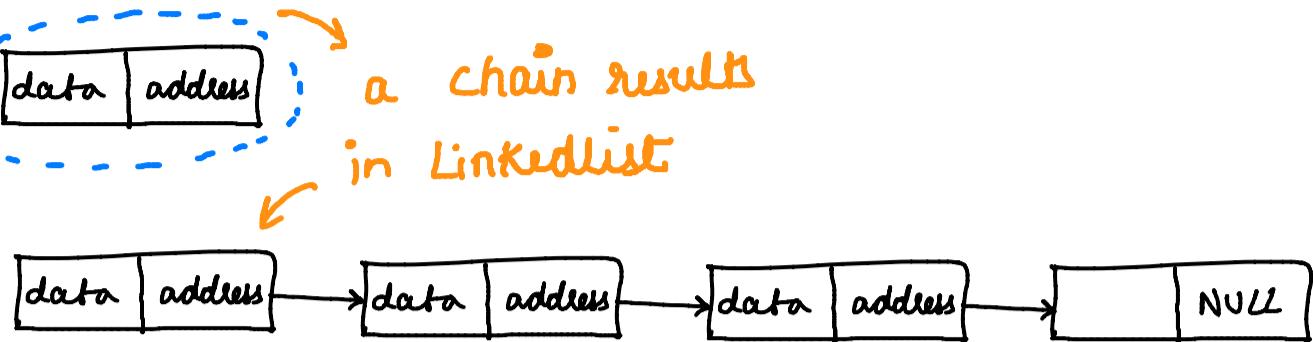
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# Linked List

LinkedList is linear data structure, which consists of a group of nodes in a sequence.

Class Node  
→ data  
→ Node\*



## Advantages

1. Dynamic nature
2. Optimal insertion & deletion
3. Stacks and queues can be easily implemented
4. No memory wastage

## Disadvantages

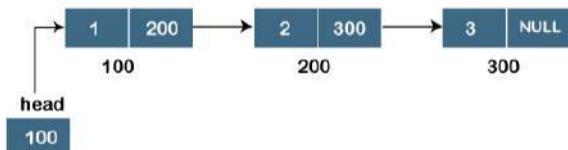
1. More memory usage due to address pointer.
2. Slow traversal compared to arrays.
3. No reverse traversal in singly linked list
4. No random access.

## Real life Applications

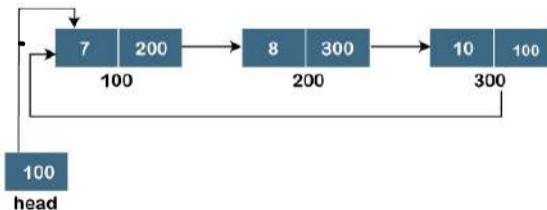
1. Previous & next page in browser
2. Image Viewer
3. Music player

## Type

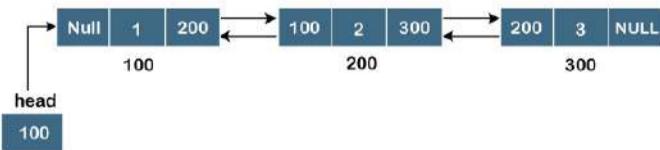
### 1. Singly linkedlist



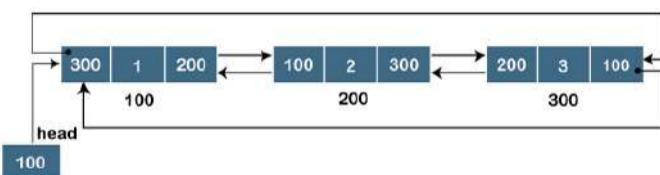
### 3. Circular Linkedlist



### 2. Doubly linkedlist



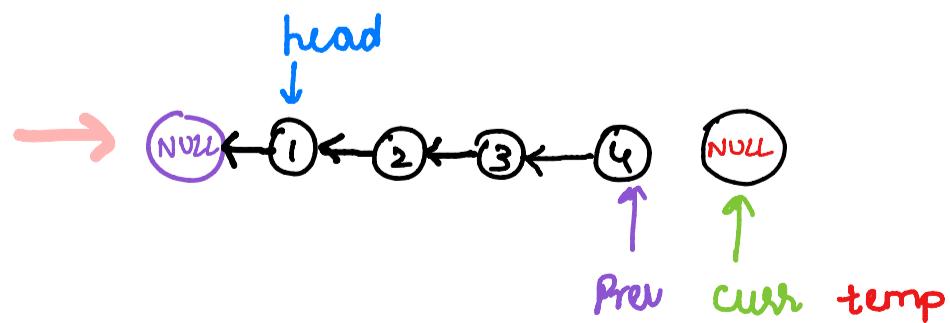
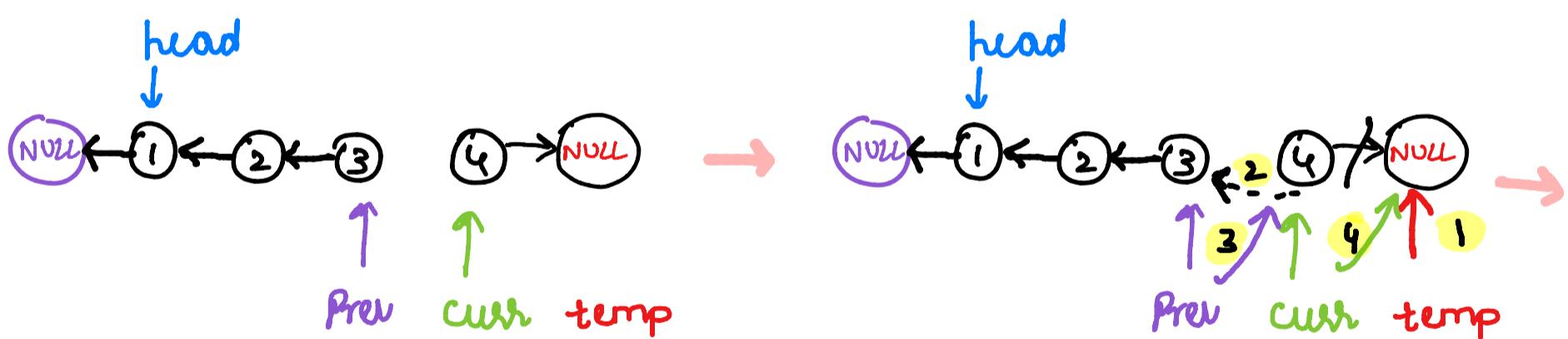
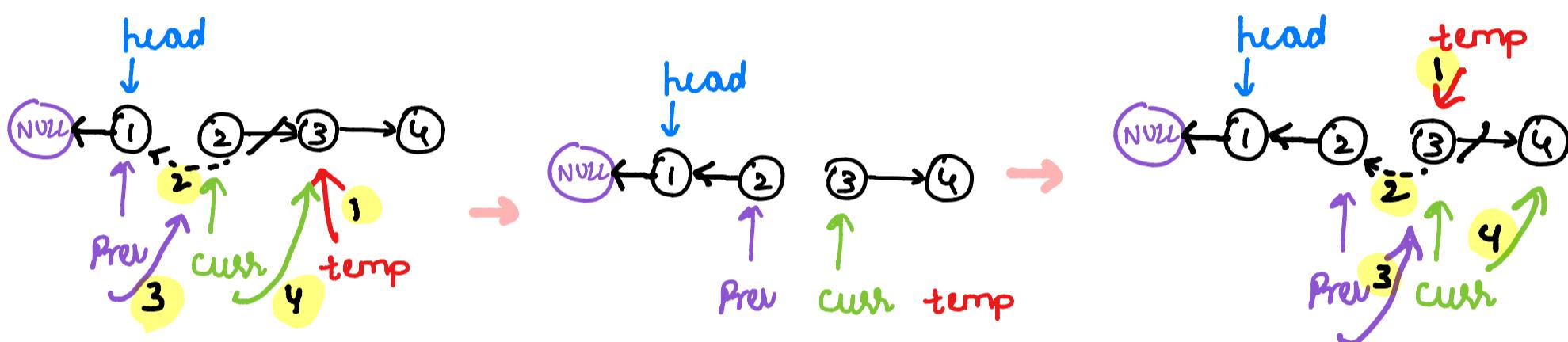
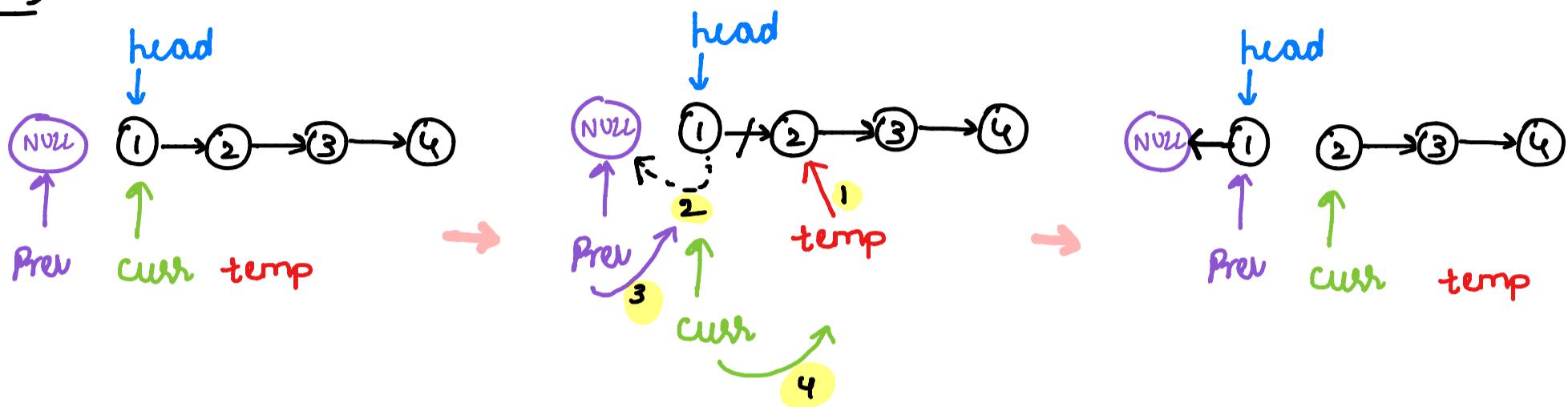
### 4. Doubly circular linkedlist



① Reverse a linkedlist → gives a linkedlist, returns reversed list.

Eg  $1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \Rightarrow 4 \rightarrow 3 \rightarrow 2 \rightarrow 1$

Sol)



→ As curr points to NULL, we reached end of Linkedlist  
→ Return Prev as it is the starting pointer of reversed list.

## Recursive →



curr, prev

(1, NULL):

newNode = 1 → next = 2

1 → next = null

call (2, 1)

curr, prev

(2, 1):

newNode = 2 → next = 3

2 → next = 1

call (3, 2)

curr, prev

(3, 2):

newNode = 3 → next = 4

3 → next = 2

call (4, 3)

curr, prev

(4, 3):

newNode = 4 → next = null

4 → next = 3

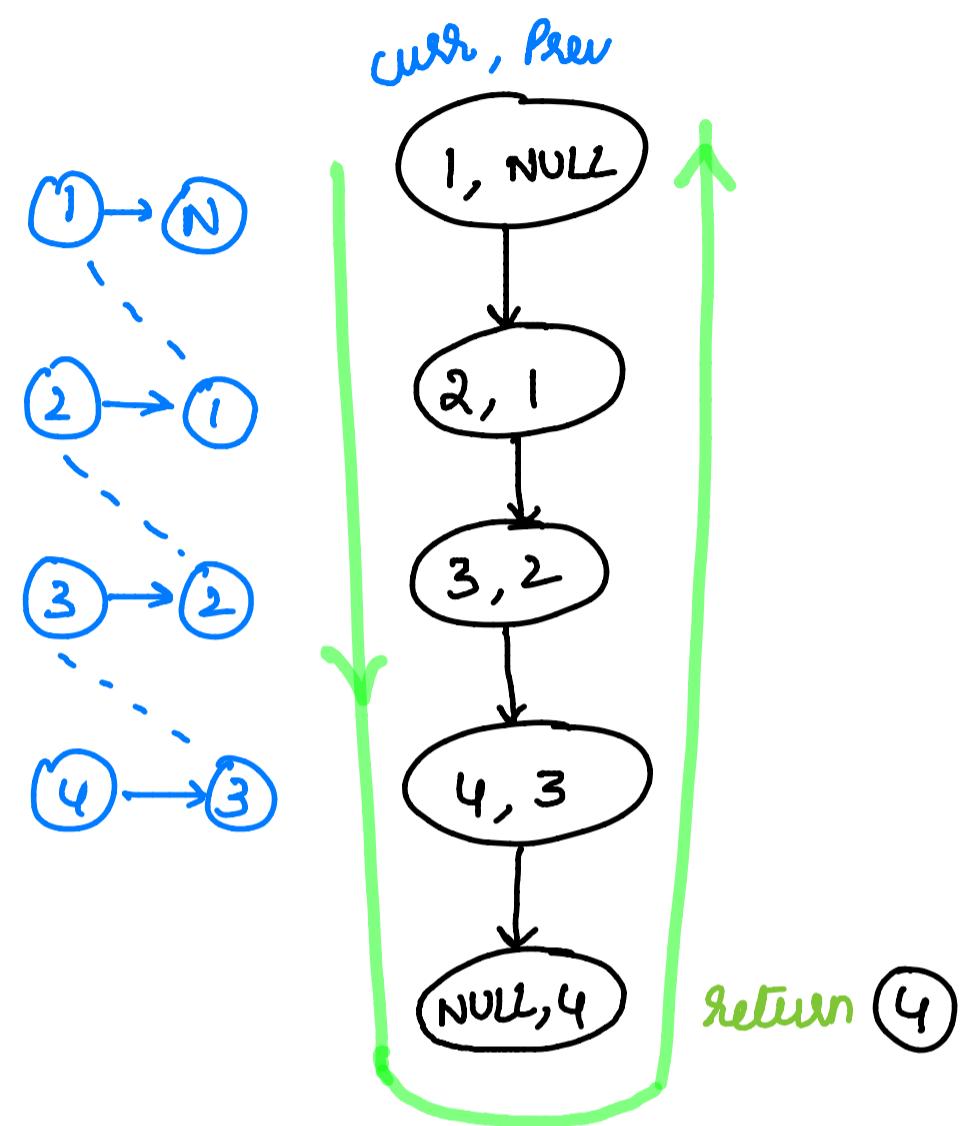
call (NULL, 4)

curr, prev

(NULL, 4):

as curr == NULL

return prev



func (curr, prev):

if curr == NULL

return prev

newNode = curr → next

curr → next = prev

recursively call newNode & curr as  
as curr → prev

Code →

$Tc \rightarrow O(n)$

$Sc \rightarrow O(1)$

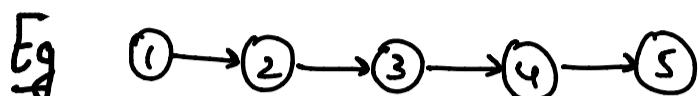


```
1
2 // Iterative ->
3 class Solution {
4 public:
5     ListNode* reverseList(ListNode* head) {
6         ListNode *prev = NULL, *curr=head, *temp;
7         while(curr){
8             temp = curr->next;
9             curr->next = prev;
10            prev = curr;
11            curr = temp;
12        }
13        return prev;
14    }
15 };
16
17
18 // Recursive ->
19 class Solution {
20 public:
21     ListNode* reverseLinker(ListNode* curr, ListNode* prev) {
22         if(curr==NULL)
23             return prev;
24         ListNode* newNode = curr->next;
25         curr->next = prev;
26         return helper(newNode, curr);
27     }
28
29     ListNode* reverseLinker(ListNode* head) {
30         return helper (head, NULL);
31     }
32 };
```

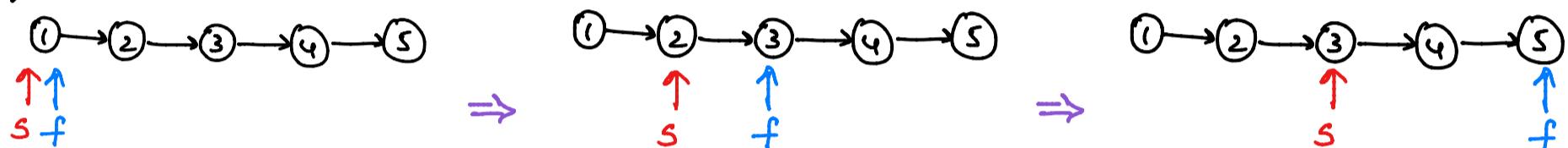
② Middle of Linkedlist → Given the head, return middle node.

Approach 1 → traverse the list & find no of nodes & return mid

Approach 2 → use 2 pointers → slow (moves by 1) } By time  
fast (moves by 2) } fast reaches end, slow points to the mid.



Res ⇒ 



code

as fast reached end  
return slow.



```
● ● ●  
1 class Solution {  
2 public:  
3     ListNode* middleNode(ListNode* head) {  
4         if(head == NULL)  
5             return head;  
6         ListNode* slow = head, *fast = head;  
7  
8         // Traverse the LinkedList  
9         while(fast != NULL && fast -> next != NULL)  
10        {  
11            slow = slow -> next;  
12            fast = fast -> next -> next;  
13        }  
14  
15        return slow;  
16    }  
17}
```

TC → O(n)

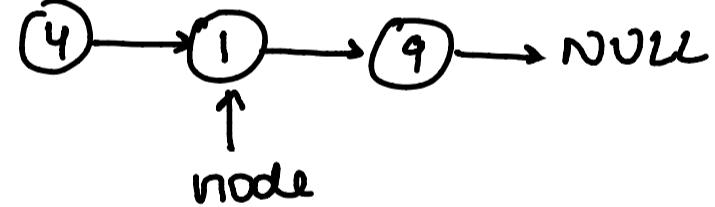
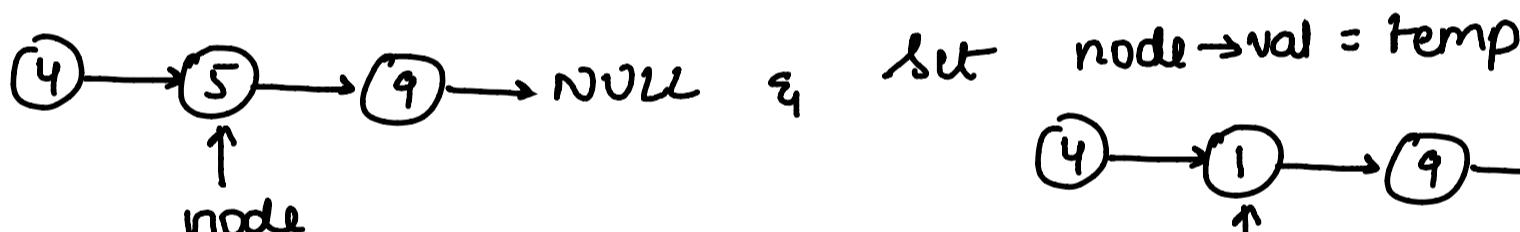
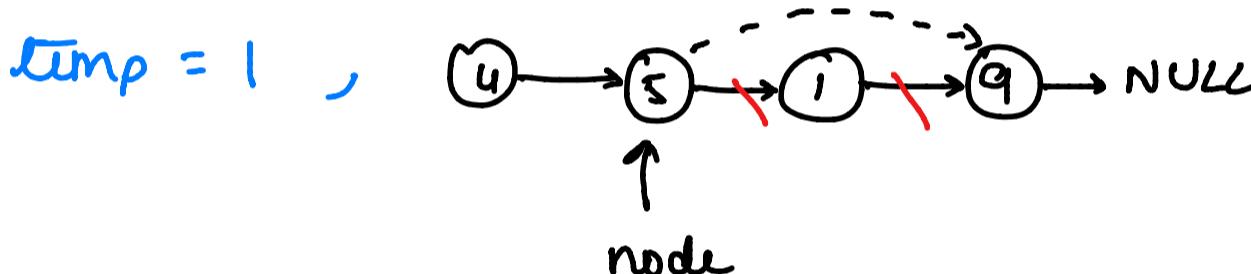
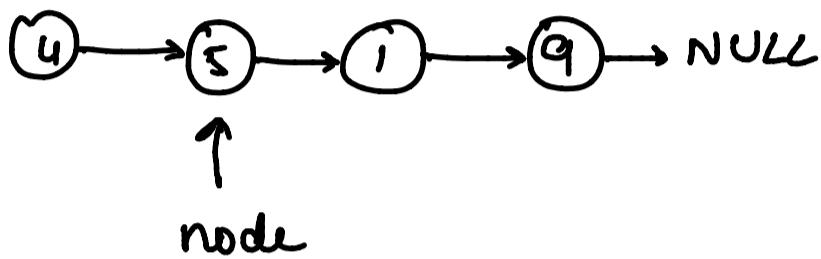
SC → O(1)

### ③ Delete node in a linkedlist →

given a linkedlist's node, delete it.

- copy node's next node's val into a temp variable
- skip the node→next node
- copy the temp variable's value into the node.

Eg

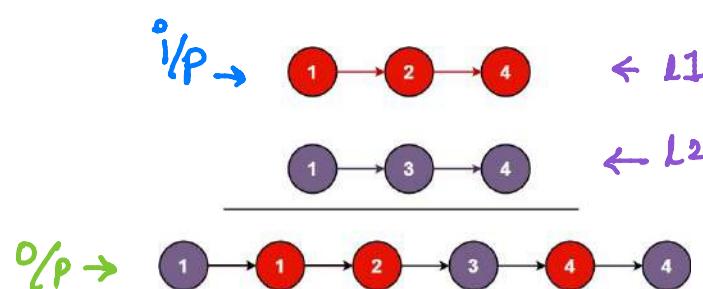


code →

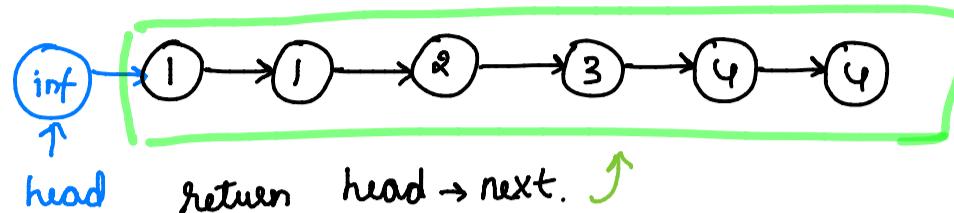
```
1 class Solution {
2 public:
3     void deleteNode(ListNode* node) {
4         int val = node->next->val;
5         node->next = node->next->next;
6         node->val = val;
7     }
8 };
```

Tc → O(1)  
Sc → O(1)

#### ④ Merge two sorted lists



→ Take a dummyNode & chain the next node which contains smaller value of  $L_1$  &  $L_2$



Code →

```

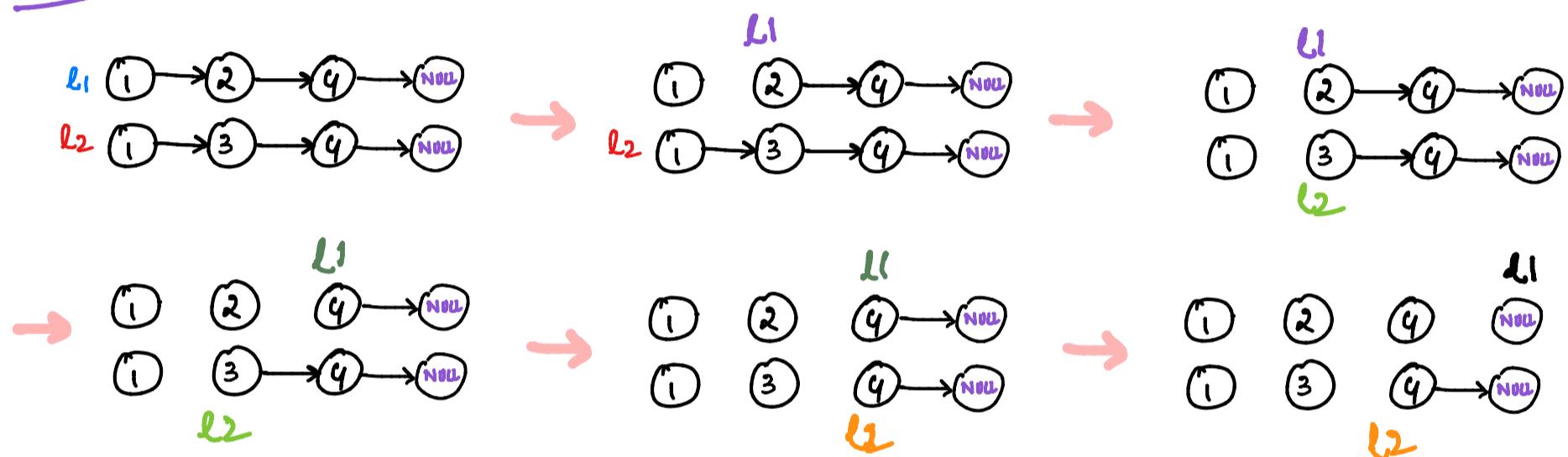
1  class Solution {
2  public:
3      ListNode* mergeTwoLists(ListNode* l1, ListNode* l2) {
4
5          if( l1 == NULL ) return l2;
6          if( l2 == NULL ) return l1;
7
8          ListNode* dummy = new ListNode(-101);
9          ListNode* head = dummy;
10
11         // Traverse the lists
12         while( l1 != NULL && l2 != NULL )
13         {
14             if( l1->val < l2->val )
15             {
16                 ListNode* newnode = new ListNode(l1->val);
17                 dummy->next = newnode;
18                 l1 = l1->next;
19             }
20             else
21             {
22                 ListNode* newnode = new ListNode(l2->val);
23                 dummy->next = newnode;
24                 l2 = l2->next;
25             }
26             dummy = dummy->next;
27         }
28
29         /* If a particular list is NULL, then directly chain
30         the other */
31         if(l1!=NULL) dummy->next = l1;
32         if(l2!=NULL) dummy->next = l2;
33
34         return head->next;
35     }
36 };

```

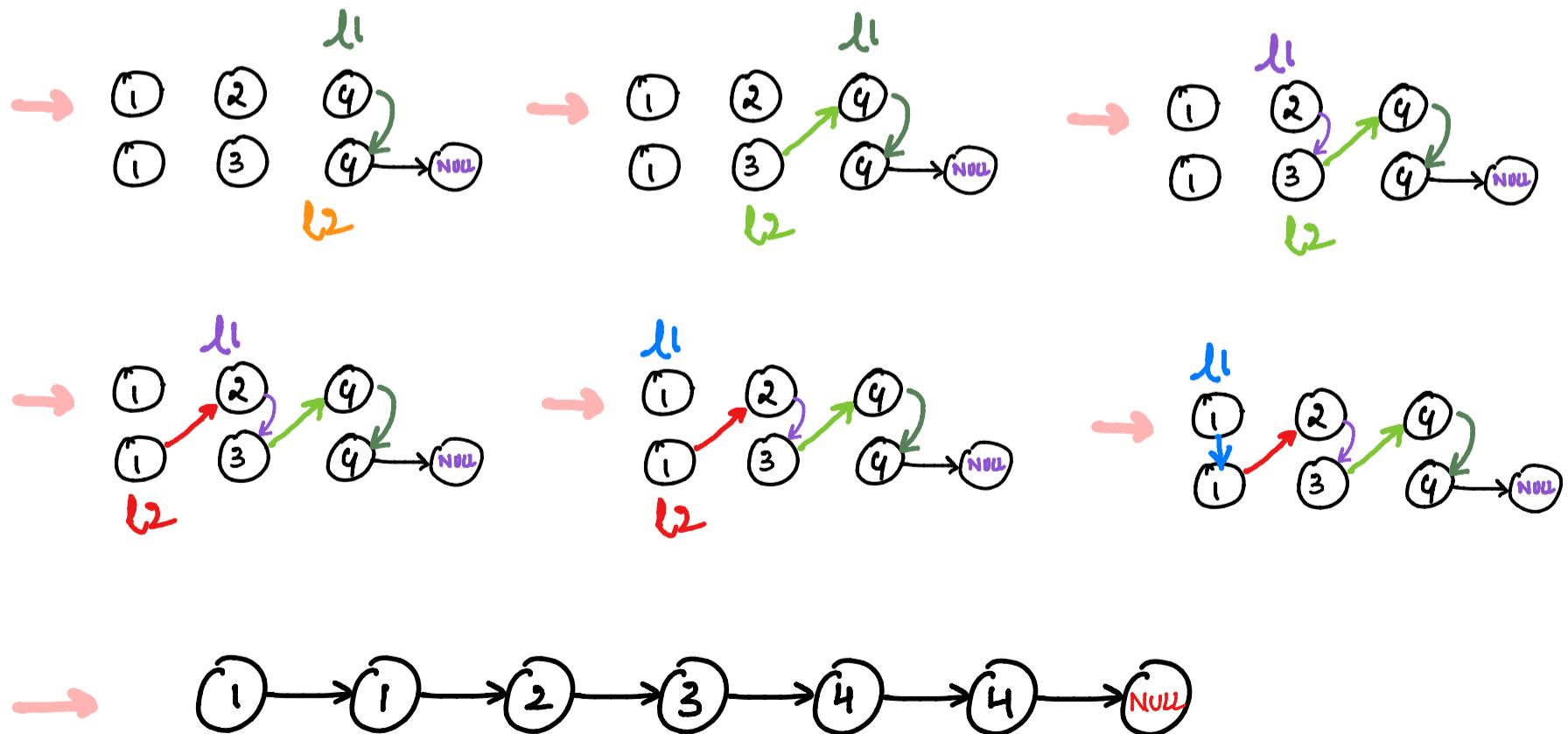
## Recursive Code →

```
class Solution {  
public:  
    ListNode* mergeTwoLists(ListNode* l1, ListNode* l2) {  
  
        if (l1 == NULL) return l2;  
        if (l2 == NULL) return l1;  
  
        // compare the starting values and link the nodes  
        if (l1->val <= l2->val) {  
            l1->next = mergeTwoLists(l1->next, l2);  
            return l1;  
        } else {  
            l2->next = mergeTwoLists(l1, l2->next);  
            return l2;  
        }  
    }  
};
```

## Dry Run

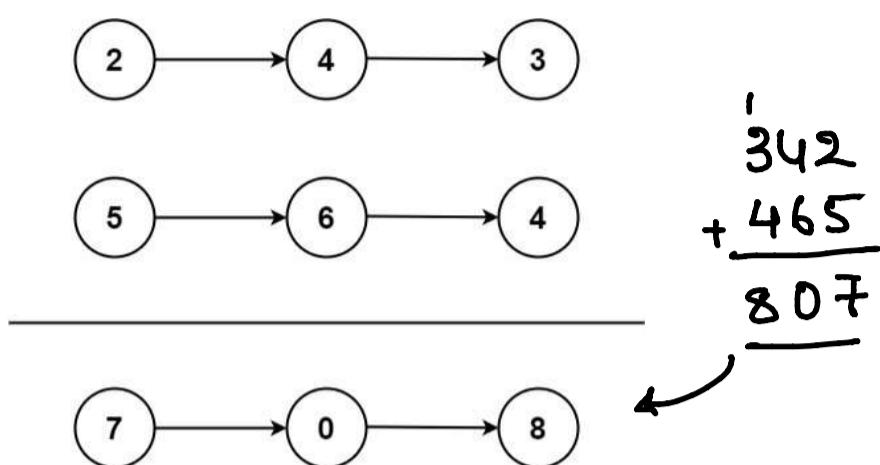


As  $l_1$  is null, return  $l_2$



## ⑤ Add two Numbers

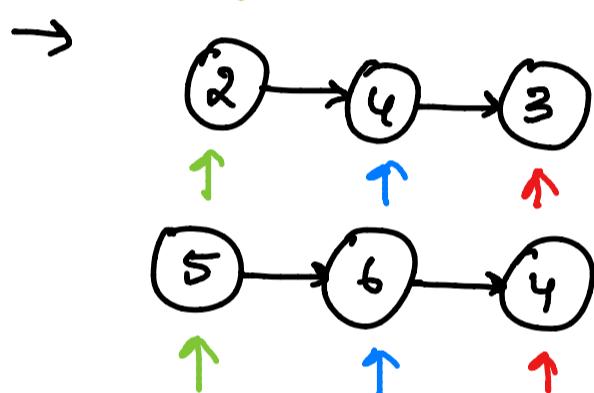
Given 2 lists in reverse order, add them and return the sum.



$Tc \rightarrow O(m+n)$

$Sc \rightarrow O(\max(m,n))$

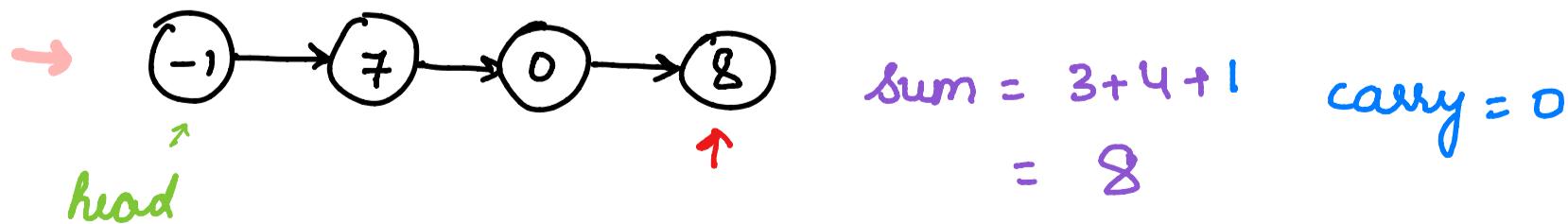
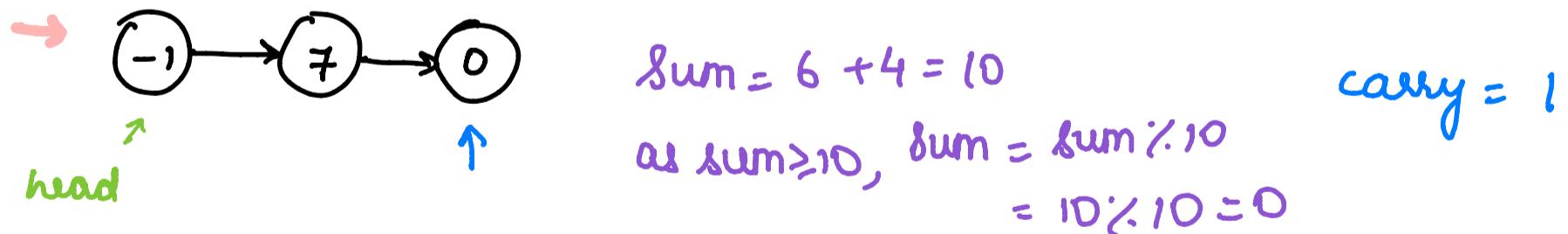
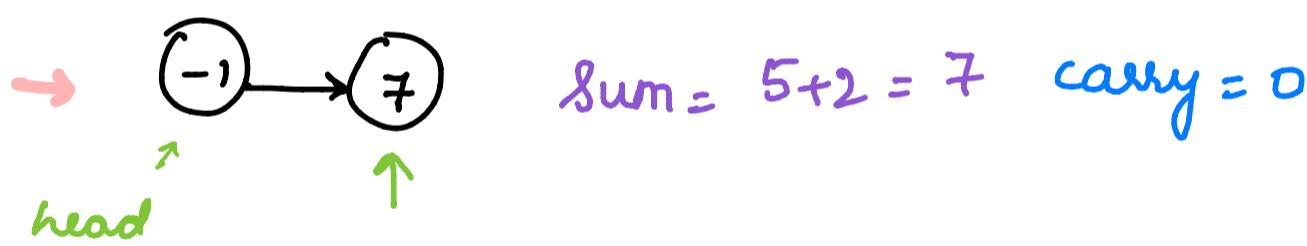
Initially,



1) traverse both list simultaneously  
if sum  $\geq 10$ , then set carry=1

2) add both values + carry

3) Create newNode with this value



## code

```
1  class Solution {
2  public:
3
4      ListNode* addTwoNumbers(ListNode* l1, ListNode* l2) {
5
6          ListNode* dummyNode;
7          ListNode* head;
8          dummyNode = head = new ListNode(-1);
9          if(!l1)
10              return l2;
11          if(!l2)
12              return l1;
13
14          int carry = 0;
15
16          while(l1 || l2){
17              int firstVal = l1 ? l1->val : 0;
18              int secondVal = l2 ? l2->val : 0;
19
20              int total = firstVal + secondVal + carry;
21              carry = total / 10;
22              total = total % 10;
23
24              ListNode* newNode = new ListNode(total);
25              dummyNode->next = newNode;
26
27              dummyNode = dummyNode->next;
28
29              l1 = l1 ? l1->next : l1;
30              l2 = l2 ? l2->next : l2;
31          }
32
33          if(carry)
34              dummyNode->next = new ListNode(1);
35
36          return head->next;
37      }
38  };
```

## ⑥ Add two numbers !!

→ Problem solving approach is same as previous problem

→ Points to note :

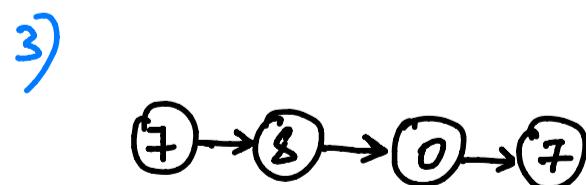
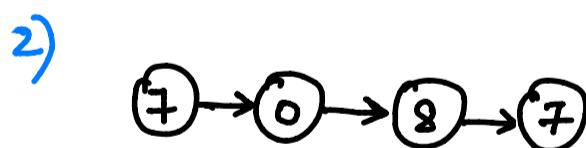
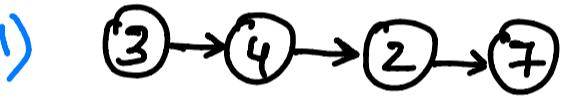
1. Reverse both the lists
2. Add them
3. Reverse the result

Eg →



Code →

```
1 class Solution {
2 public:
3     ListNode* reverseList(ListNode* head) {
4         ListNode* prev = NULL;
5         ListNode* curr = head;
6         ListNode* temp = NULL;
7         while(curr!=NULL)
8         {
9             temp = curr->next;
10            curr->next = prev;
11            prev = curr;
12            curr = temp;
13        }
14        return prev;
15    }
16
17    ListNode* addTwoNumbers(ListNode* l1, ListNode* l2) {
18        l1=reverseList(l1); // O(n)
19        l2=reverseList(l2); // O(n)
20        ListNode* dummyNode;
21        ListNode* head;
22        dummyNode = head = new ListNode(-1);
23        if(!l1)
24            return l2;
25        if(!l2)
26            return l1;
27
28        int carry = 0;
29
30        while(l1 || l2){
31            int firstVal = l1 ? l1->val : 0;
32            int secondVal = l2 ? l2->val : 0;
33
34            int total = firstVal + secondVal + carry;
35            carry = total / 10;
36            total = total % 10;
37
38            ListNode* newNode = new ListNode(total);
39            dummyNode->next = newNode;
40
41            dummyNode = dummyNode->next;
42
43            l1 = l1 ? l1->next : l1;
44            l2 = l2 ? l2->next : l2;
45        }
46
47        if(carry)
48            dummyNode->next = new ListNode(1);
49
50        return reverseList(head->next); // O(max(m,n))
51    }
52};
```



Result ↑

## ⑦ Linked List Cycle

$Tc \rightarrow O(n)$   
 $Sc \rightarrow O(1)$

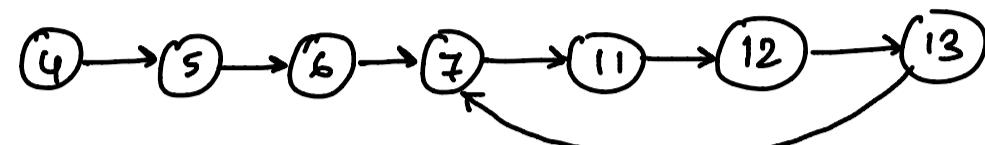
### Approach - 1

Create a set of nodes & insert every node into it,  
if already exist then return true else false.

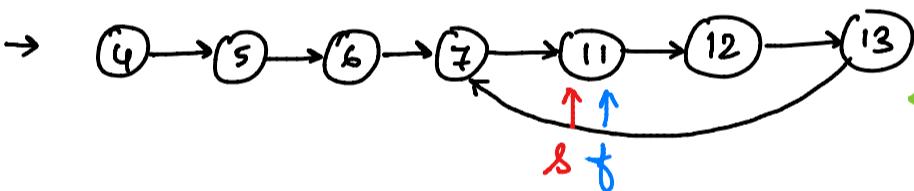
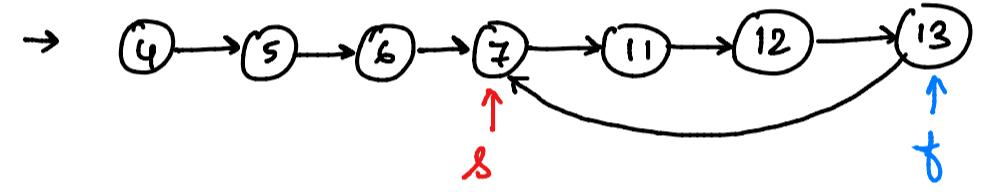
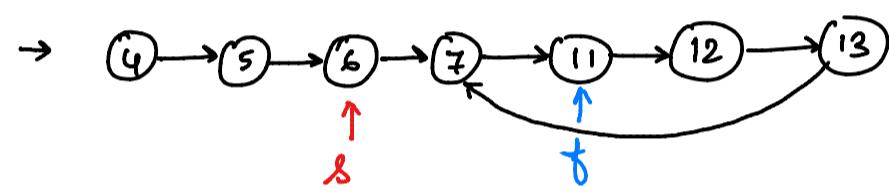
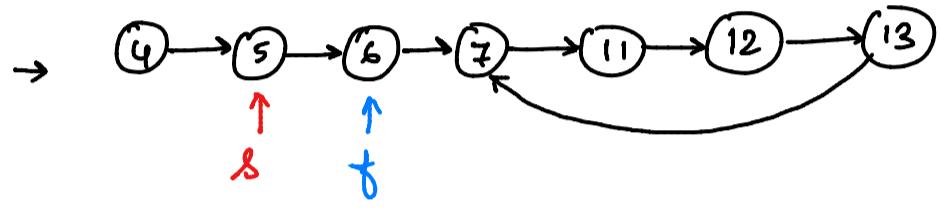
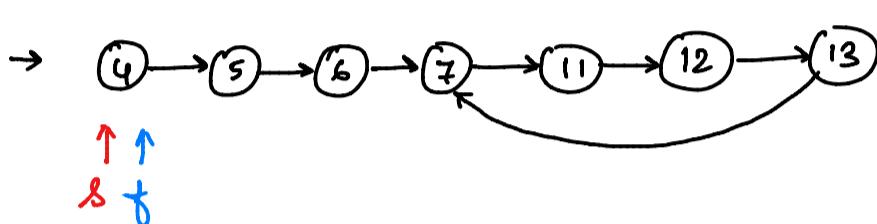
### Approach - 2

using fast & slow pointer  $Tc \rightarrow O(n)$   $Sc \rightarrow O(1)$

Eg



keep iterating till  
 $fast \rightarrow next \&& fast \rightarrow next \rightarrow next$  exist.

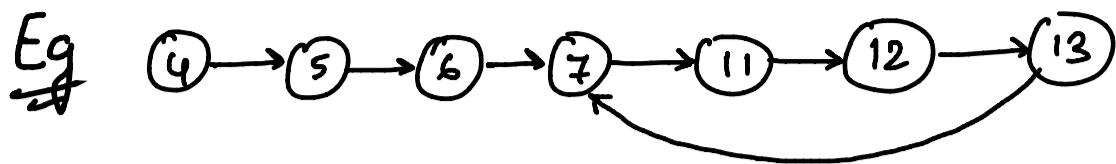


return true

Code →

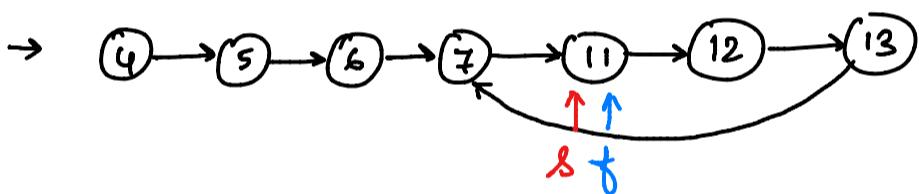
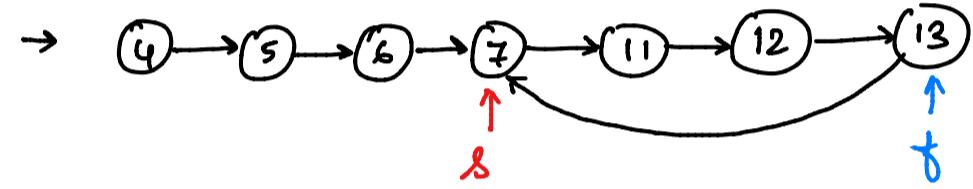
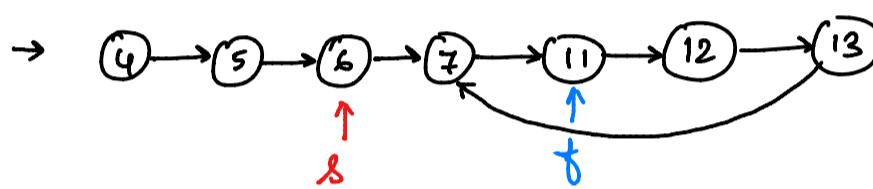
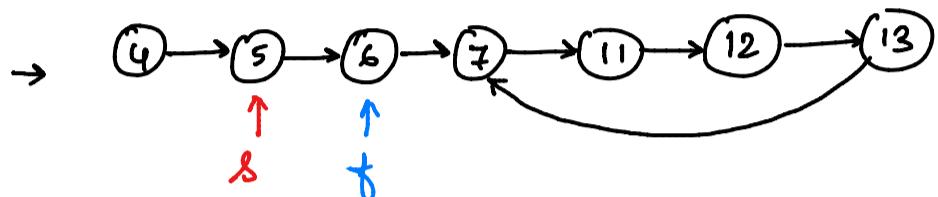
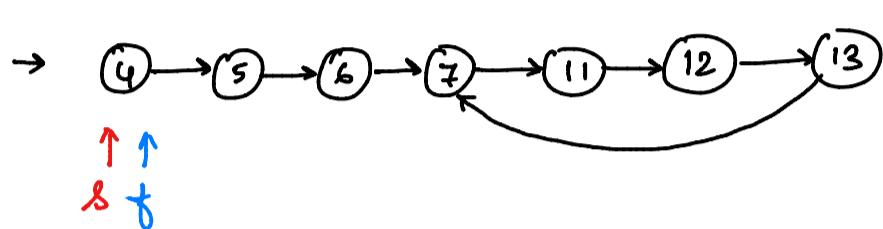
```
1 class Solution {
2 public:
3     bool hasCycle(ListNode *head) {
4         if(head==NULL) return false;
5         ListNode *fast = head, *slow = head;
6         while(fast->next!=NULL && fast->next->next!=NULL)
7         {
8             fast=fast->next->next;
9             slow=slow->next;
10            if(fast==slow) return true;
11        }
12        return false;
13    }
14};
```

⑧ Linked list cycle  $\Rightarrow$  returns the node where cycle begins.

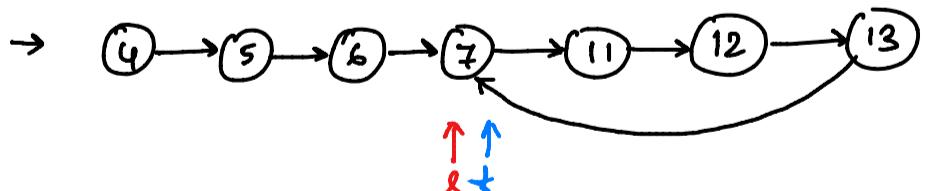
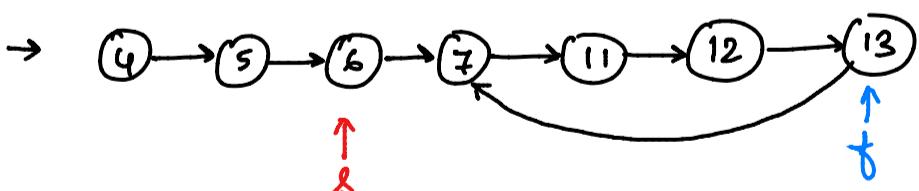
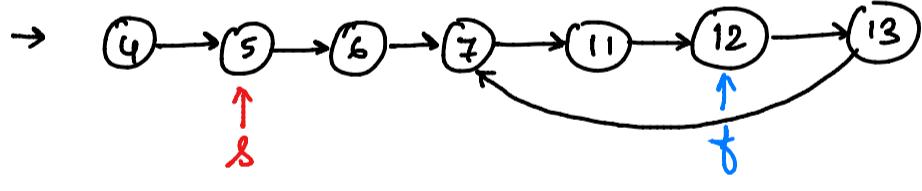
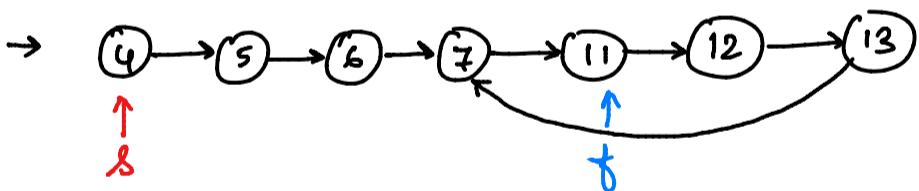


result = 7

keep iterating while  $\text{fast} \rightarrow \text{next}$  &  $\text{fast} \rightarrow \text{next} \rightarrow \text{next}$  exist.



← once  $\text{slow} == \text{fast}$ , then set  
 $\text{slow} = \text{head}$ , & move pointers  
by 1 unit



$\rightarrow$  when  $\text{slow} == \text{fast}$ , it denotes the node where cycle begins.

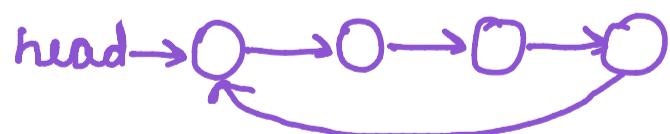
return slow;  $\Rightarrow$  7

code →

```
● ● ●  
1 class Solution {  
2 public:  
3     ListNode *detectCycle(ListNode *head) {  
4         if(head==NULL) return NULL;  
5         ListNode *fast = head, *slow = head;  
6         while(fast->next!=NULL && fast->next->next!=NULL)  
7         {  
8             fast = fast->next->next;  
9             slow = slow->next;  
10            if(fast == slow)  
11            {  
12                slow = head;  
13                while (slow != fast)  
14                {  
15                    slow = slow->next;  
16                    fast = fast->next;  
17                }  
18                return slow;  
19            }  
20        }  
21        return NULL;  
22    }  
23};
```

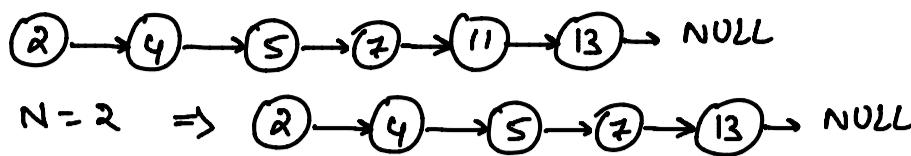
$$TC \rightarrow O(n) + O(n) = O(2n) = \underline{\underline{O(n)}}$$

\* worst case when its a loop



$$SC \rightarrow O(1)$$

## ⑨ Remove $N^{\text{th}}$ node from End of list



Approach - 1 → find length of list L, delete  $L-N+1^{\text{th}}$  node

Approach - 2 → a) Reverse b) Delete  $N^{\text{th}}$  node c) Reverse

Code →

$T_C \rightarrow O(n)$

$S_C \rightarrow O(1)$

```
1 class Solution {
2     public:
3         ListNode* reverseList(ListNode* head) {
4             ListNode *prev = NULL, *curr = head, *temp;
5             while(curr){
6                 temp = curr->next;
7                 curr->next = prev;
8                 prev = curr;
9                 curr = temp;
10            }
11            return prev;
12        }
13
14        ListNode* removeNthFromEnd(ListNode* head, int n) {
15            ListNode *dummy = new ListNode(-1);
16            dummy->next = reverseList(head);
17            head = dummy;
18            ListNode *curr = head;
19            ListNode *prev = NULL;
20            // Iteration
21            for(int i=0; i<n; i++)
22            {
23                prev = curr;
24                curr = curr->next;
25            }
26            // Deletion
27            prev->next = curr->next;
28            return reverseList(head->next);
29        }
30    };
```

## ⑩ Pallindrome Linked List

Approach - 1 → Create a copy of list & reverse it. Compare value by value.  
If all are equal true else false.

Approach - 2 → Reach middle node & return the remaining list as new list. Reverse the newList & compare its value by value.

Code →

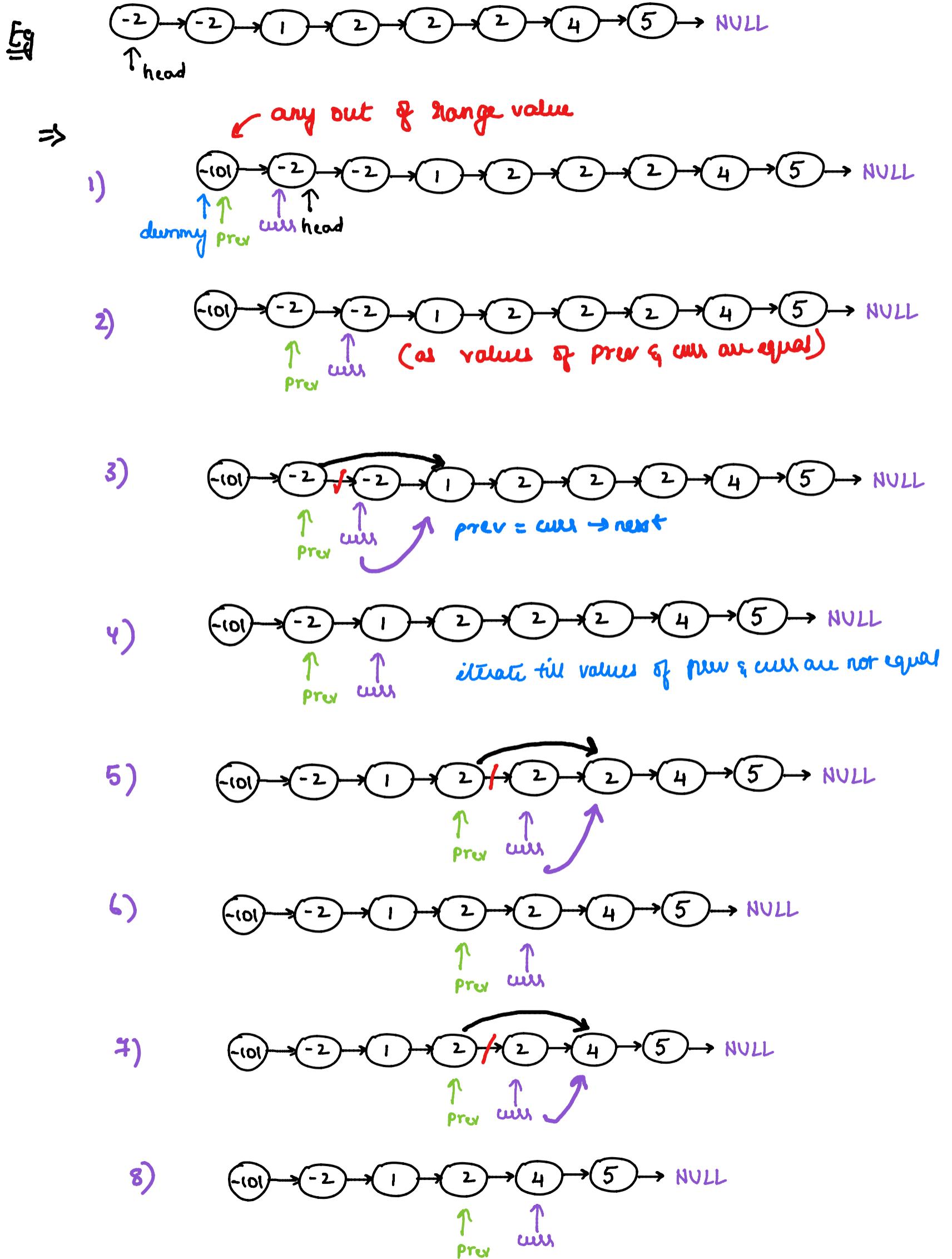
TC → O(n)

SC → O(1)

```
● ● ●
1 class Solution {
2 public:
3     ListNode* midNode(ListNode*head)
4     {
5         ListNode *fast = head, *slow = head;
6         while(fast->next!=NULL and fast->next->next!=NULL){
7             fast=fast->next->next;
8             slow = slow->next;
9         }
10        return slow;
11    }
12    ListNode* reverseList(ListNode* head) {
13        ListNode *prev = NULL, *curr = head, *temp;
14        while(curr!=NULL)
15        {
16            temp=curr->next;
17            curr->next = prev;
18            prev = curr;
19            curr = temp;
20        }
21        return prev;
22    }
23    bool compare(ListNode*l1,ListNode* l2)
24    {
25        while(l1!=NULL && l2!=NULL)
26        {
27            if(l1->val!=l2->val)    return false;
28            l1 = l1->next;
29            l2 = l2->next;
30        }
31        return true;
32    }
33    bool isPalindrome(ListNode* head) {
34        if(head==NULL) return false;
35        if(head->next == NULL) return true;
36        ListNode *mid = midNode(head);
37        ListNode *l2 = mid->next;
38        mid->next = NULL;
39        l2 = reverseList(l2);
40        return compare(head,l2);
41    }
42};
```

# (1) Remove duplicates from sorted list →

Given a linkedlist, return linkedlist without duplicates.



code →

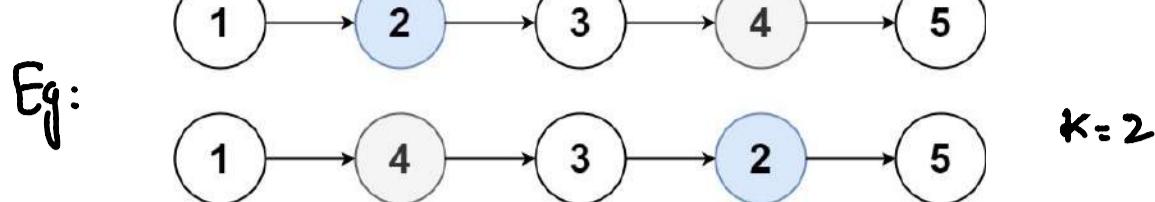
```
● ● ●

1 class Solution {
2 public:
3     ListNode* deleteDuplicates(ListNode* head) {
4         ListNode* dummy = new ListNode(101);
5         dummy->next = head;
6         ListNode* curr = head;
7         ListNode* prev = dummy;
8         while(curr!=NULL)
9         {
10             if(curr->val==prev->val){
11                 prev->next = curr->next;
12                 curr = curr->next;
13             } else {
14                 prev = curr;
15                 curr = curr->next;
16             }
17         }
18         return dummy->next;
19     }
20 };
21
22
23 // Another approach
24
25 ListNode* deleteDuplicates(ListNode* head) {
26     if(head==NULL || head->next==NULL) return head;
27     ListNode *curr = head;
28     while(curr->next!=NULL){
29         if(curr->val == curr->next->val){
30             curr->next = curr->next->next;
31         } else {
32             curr = curr->next;
33         }
34     }
35     return head;
36 }
```

$$T_C \rightarrow O(n)$$

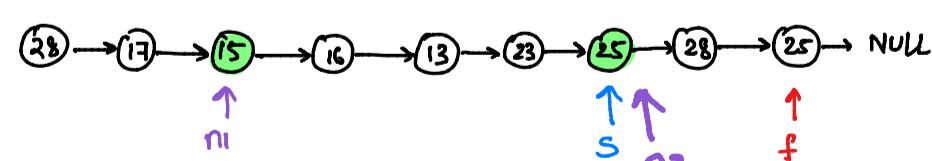
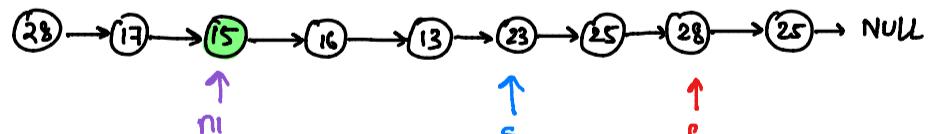
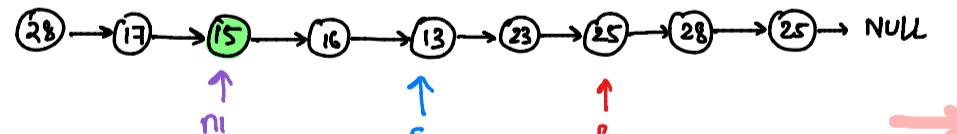
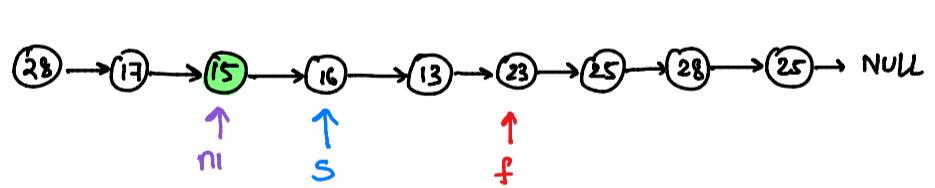
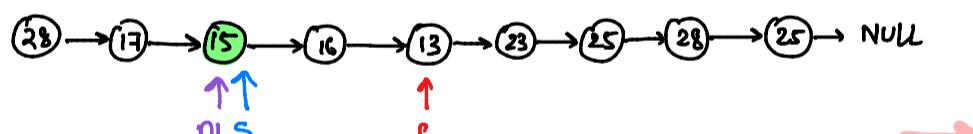
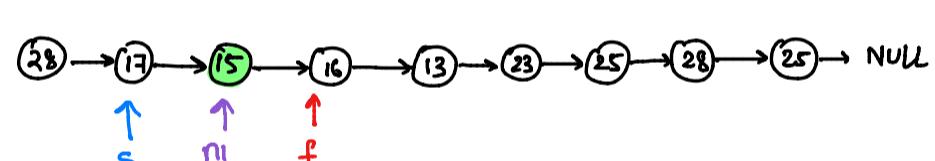
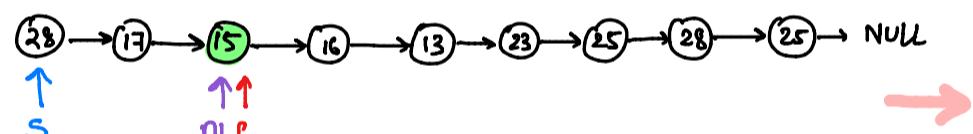
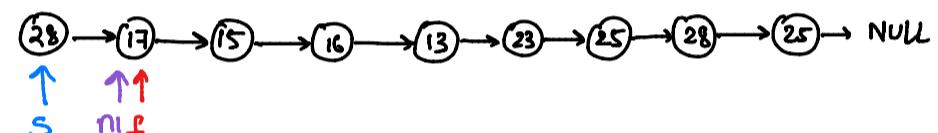
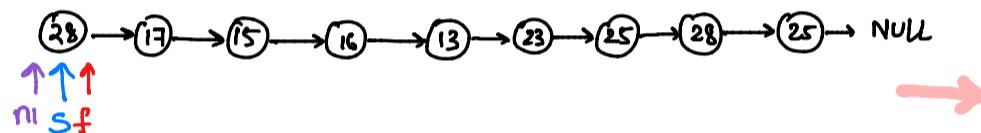
$$S_C \rightarrow O(1)$$

## (12) Swapping Nodes in Linked list →



given a linkedlist swap  $k^{th}$  node from both ends.

Eg \* for  $k-1$  times iterate f & mark n1, then iterate s & f  
(as it is 1 indexed) till f is not NULL, once null mark s as n2.  
swap(n1, n2)

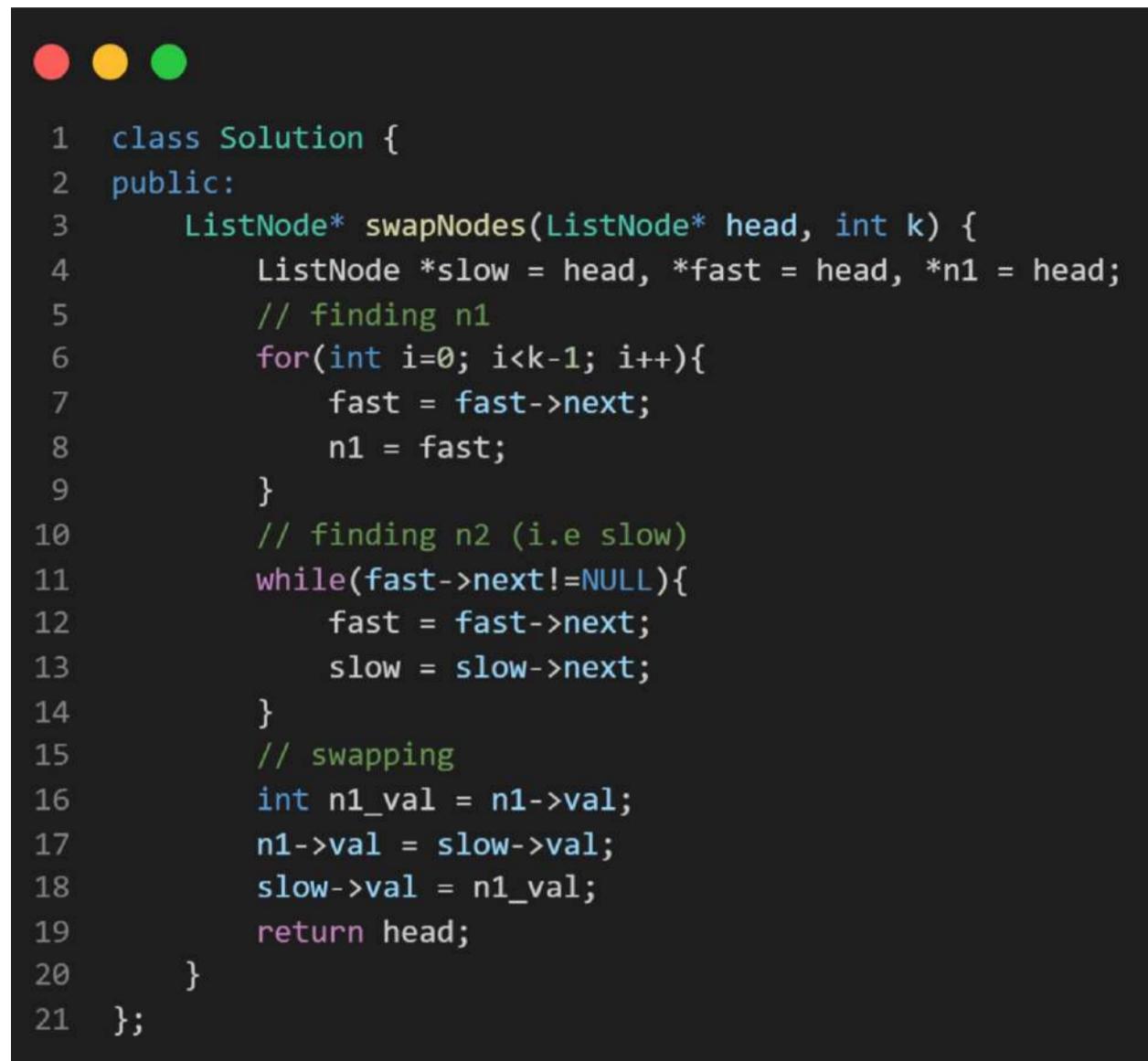


( $f \rightarrow next == NULL$ , so  $f$  is  $n_2$ )

swap(15, 25)

Result ⇒

code →



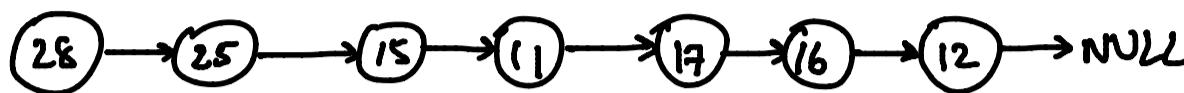
```
1 class Solution {
2 public:
3     ListNode* swapNodes(ListNode* head, int k) {
4         ListNode *slow = head, *fast = head, *n1 = head;
5         // finding n1
6         for(int i=0; i<k-1; i++){
7             fast = fast->next;
8             n1 = fast;
9         }
10        // finding n2 (i.e slow)
11        while(fast->next!=NULL){
12            fast = fast->next;
13            slow = slow->next;
14        }
15        // swapping
16        int n1_val = n1->val;
17        n1->val = slow->val;
18        slow->val = n1_val;
19        return head;
20    }
21 }
```

$T_C \rightarrow O(n)$

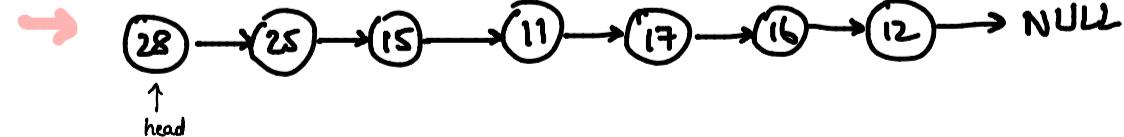
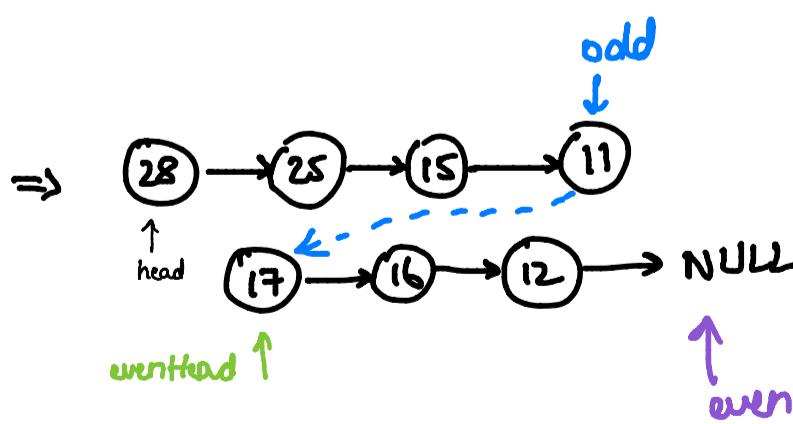
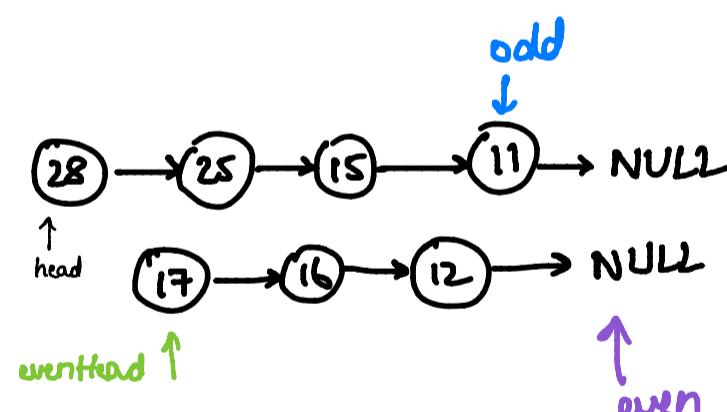
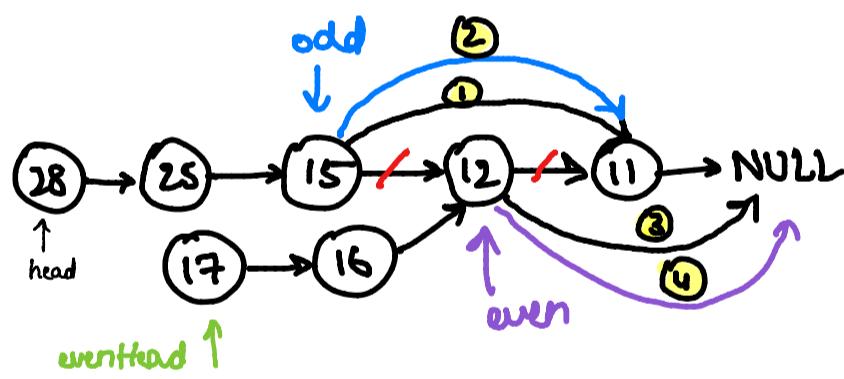
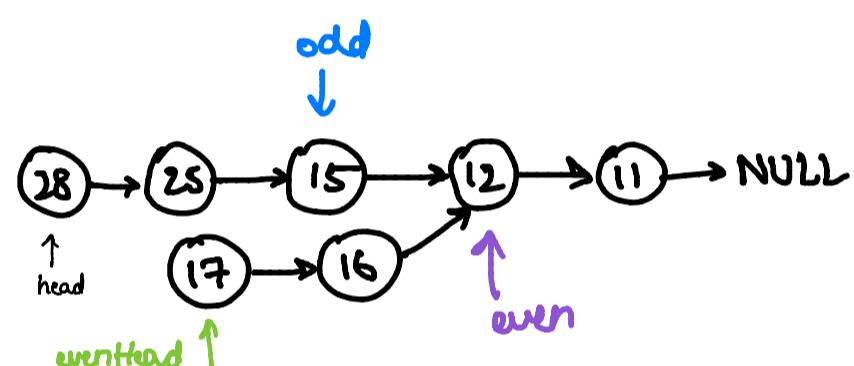
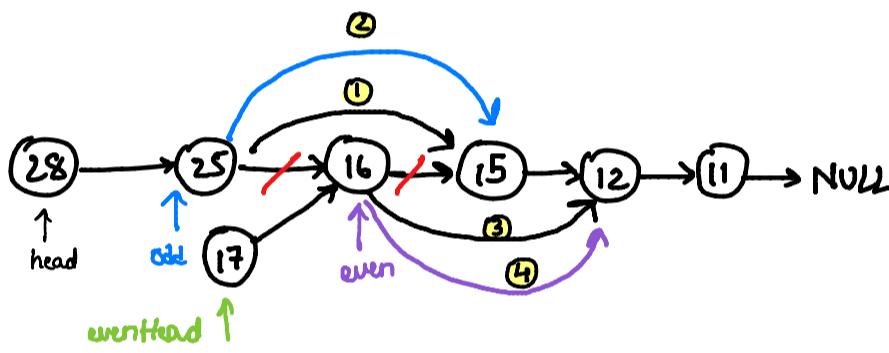
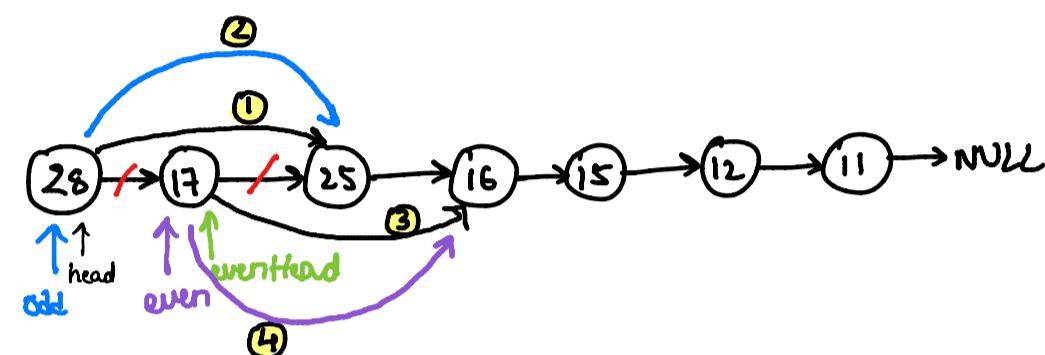
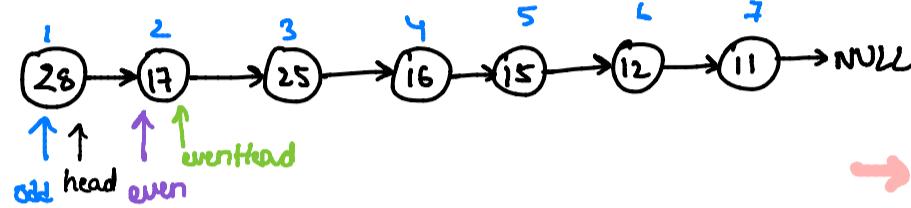
$S_C \rightarrow O(1)$

### 13 Odd Even Linked List →

gives a linkedlist group all odd indices nodes followed by even nodes



⇒



odd → next = evenhead

code →

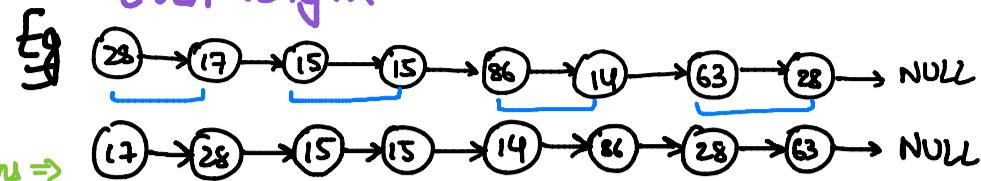


```
1 class Solution {
2 public:
3     ListNode* oddEvenList(ListNode* head) {
4         if(!head) return NULL;
5
6         ListNode *even = head->next;
7         ListNode *odd = head;
8         ListNode *evenHead = even;
9
10        while(even && even->next){
11            odd->next=even->next;
12            odd=odd->next;
13            even->next=odd->next;
14            even=even->next;
15        }
16
17        // like odd and even lists
18        odd->next = evenHead;
19        return head;
20    }
21};
```

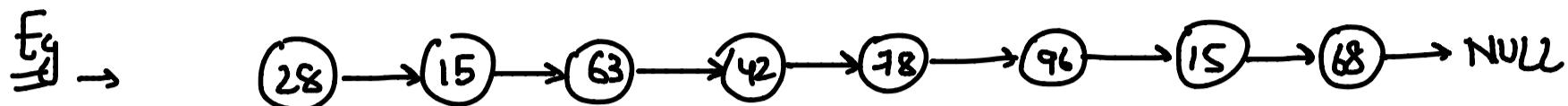
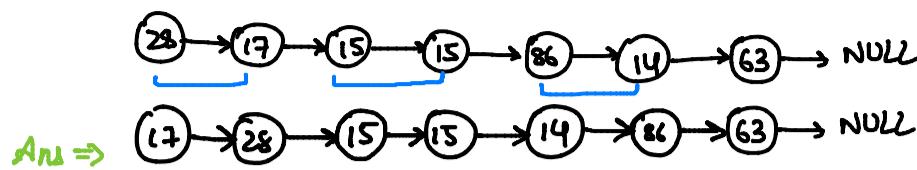
$Tc \rightarrow O(n)$   
 $Sc \rightarrow O(1)$

(14) Swap Nodes in Pairs → Given a linkedlist swap adjacent nodes.

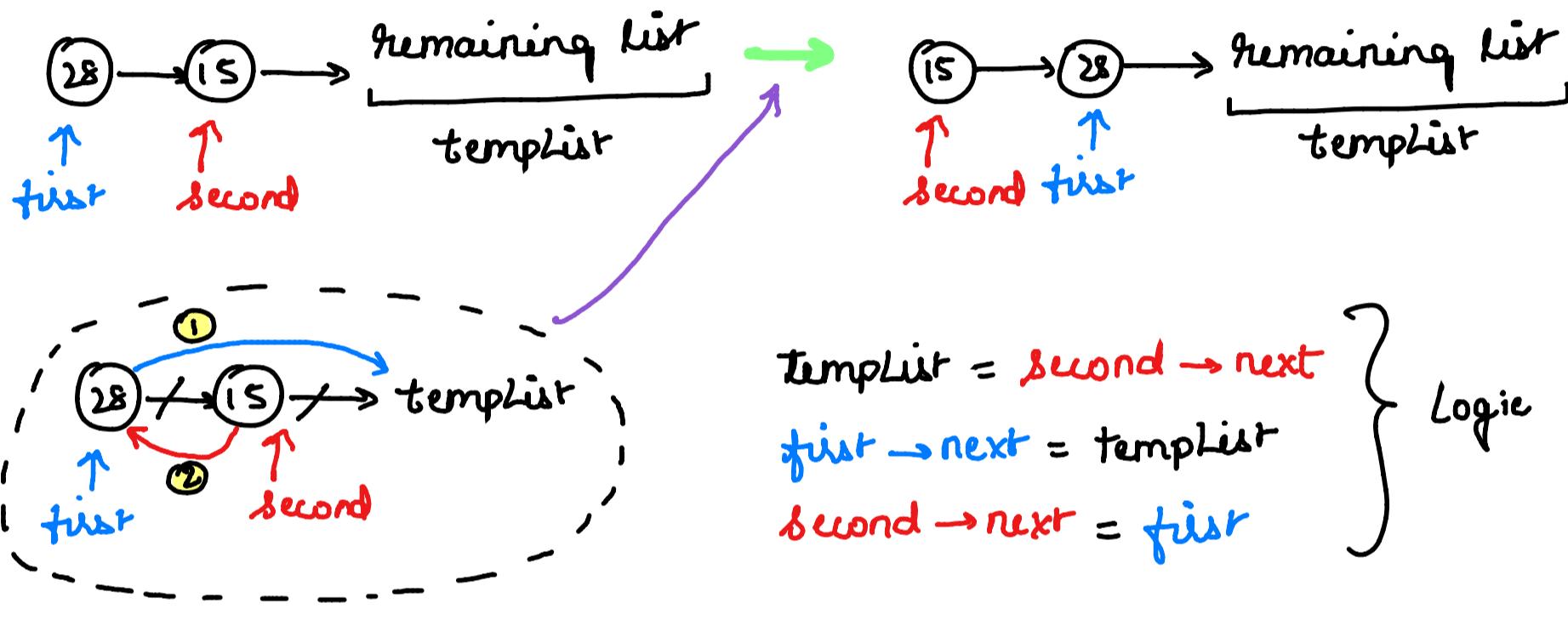
Even length



odd length



Consider for 1st pair,



Solve recursively for all pairs.

Code →

```
1 class Solution {
2 public:
3     ListNode* SwapAdjacentNodes(ListNode* head)
4     {
5         if(head==NULL || head->next==NULL) return head;
6         ListNode *first = head;
7         ListNode *second = head->next;
8         // start logic
9         ListNode *tempList = SwapAdjacentNodes(second->next);
10        first->next = temp;
11        second->next = first;
12        return second;
13    }
14    ListNode* swapPairs(ListNode* head) {
15        return SwapAdjacentNodes(head);
16    }
17};
```

TC → O(N)

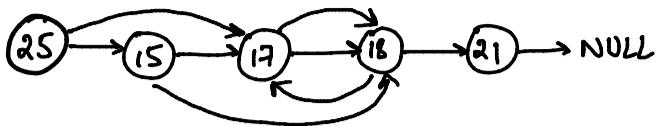
SC → O(1)

Recursive stack → O(N/2)  
 $\approx O(N)$

# 15 Copy list with random pointer

Given a list, clone & return.

Eg



mp

25	25
15	15
17	17
18	18
21	21
NULL	NULL

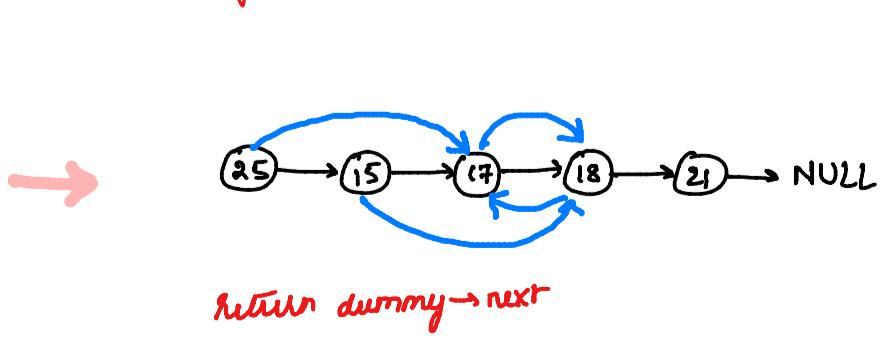
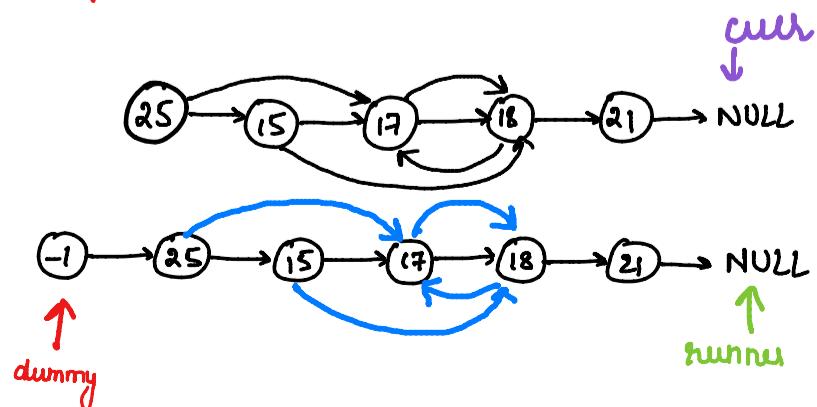
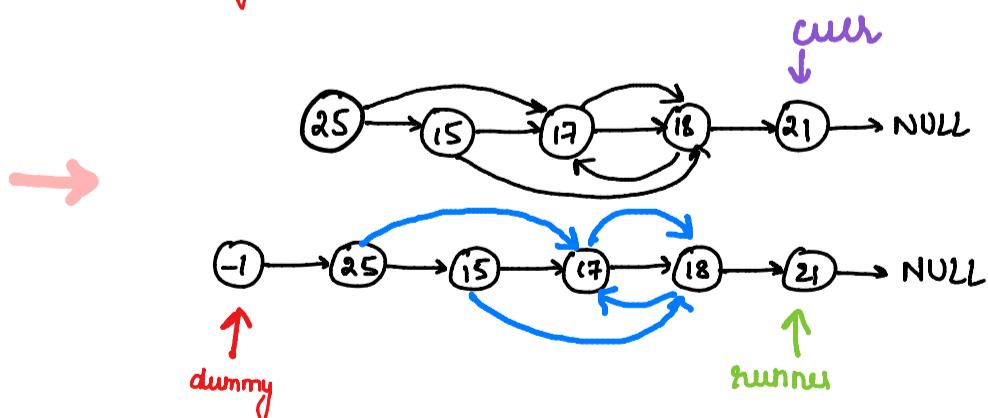
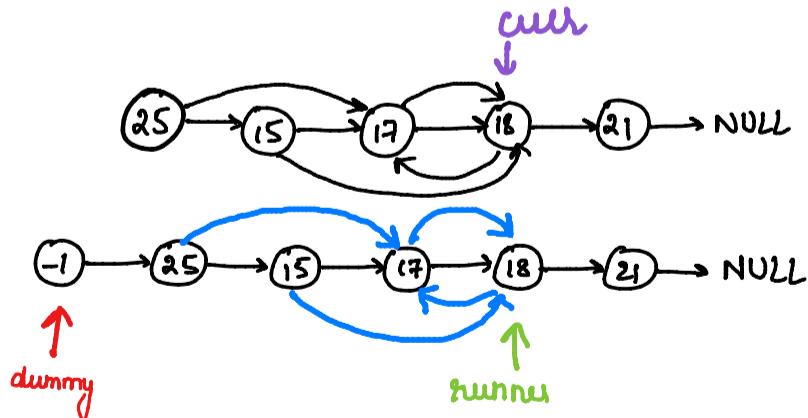
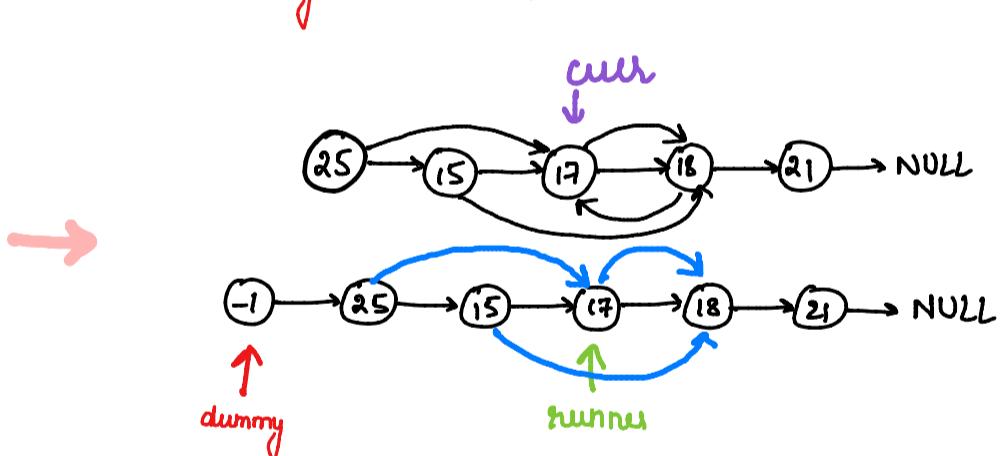
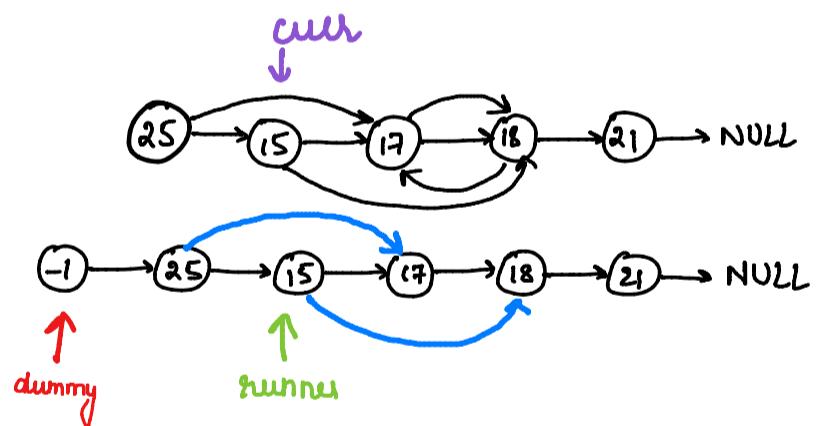
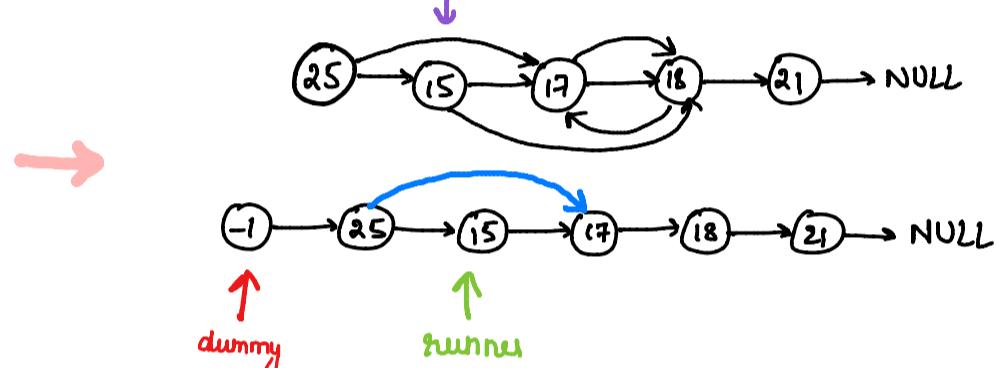
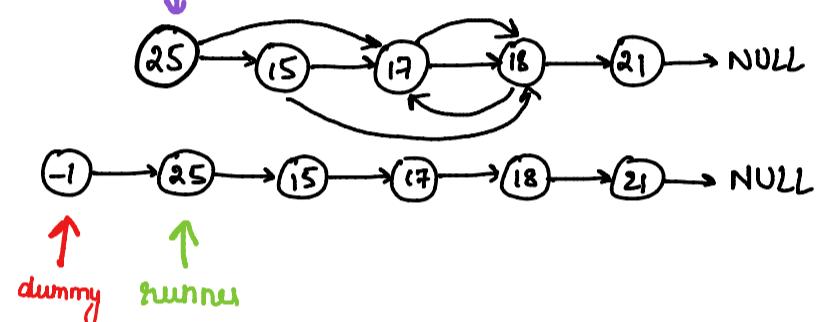
→ In 1st iteration create the list without random pointer & also maintain hashmap for mapping node pointed by random pointer

→ In 2nd iteration use map to link node pointed by random pointer

After 1st iteration →



→ curr (iterate till curr != null)



Code →

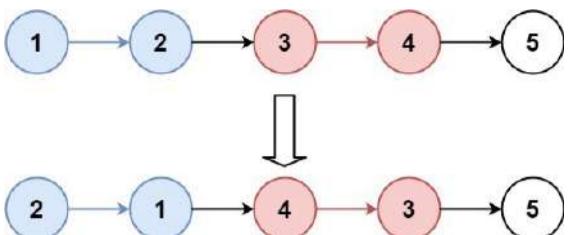
$T_C \rightarrow O(n)$

$SC \rightarrow O(n)$



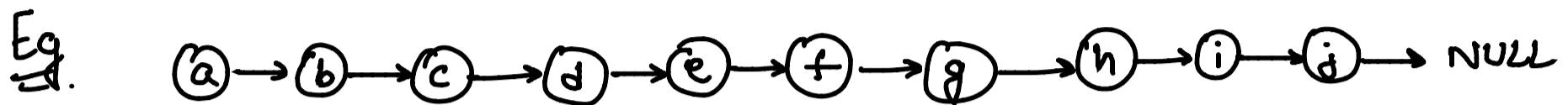
```
1 class Solution {
2 public:
3     Node* copyRandomList(Node* head) {
4
5         unordered_map<Node*, Node*> mp;
6         Node *dummy = new Node(100001);
7         Node *runner = dummy, *curr = head;
8
9         // initial iteration
10        while(curr != NULL){
11            Node *newNode = new Node(curr->val);
12            runner->next = newNode;
13            mp[curr] = newNode;
14            curr = curr->next;
15            runner = runner->next;
16        }
17
18        // setting starting points in both lists
19        curr = head;
20        runner = dummy->next;
21
22        // setting the random pointers
23        while(curr != NULL){
24            if(curr->random != NULL)
25                runner->random = mp[curr->random];
26            runner = runner->next;
27            curr = curr->next;
28        }
29
30        return dummy->next;
31    }
32}
```

## 16 Reverse Nodes in K-Group

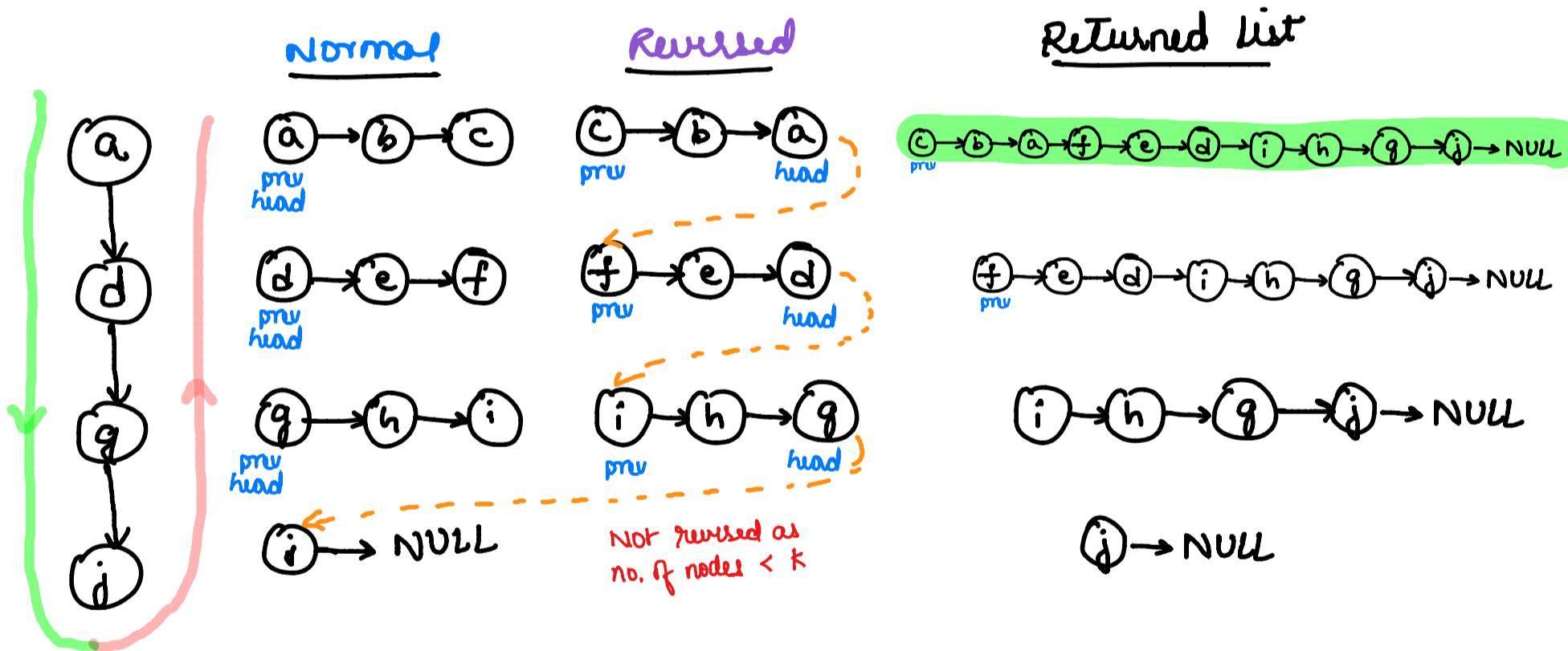
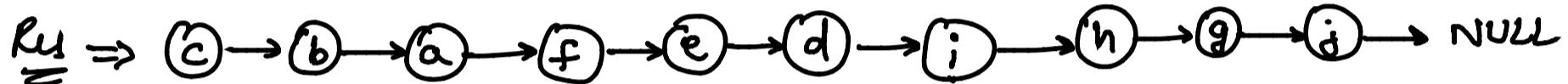


$\hookrightarrow k=2$

Given a linkedlist  $\in K$ , return a list with reversed nodes by  $K$ -groups.



$k = 3$



\* Consider the case if  $f \rightarrow e \rightarrow d$ , then  $\text{tempList} = i \rightarrow h \rightarrow g$ .  
 linking would happen as  $\text{head} \rightarrow \text{next} = \text{tempList}$  & this list would become  $\text{tempList}$  to  $c \rightarrow b \rightarrow a$

code →

$T_C \rightarrow O(n)$

$S_C \rightarrow O(1)$

```
1 class Solution {
2 public:
3     ListNode* reverseList(ListNode* head)
4     {
5         ListNode *prev = NULL, *curr = head, *temp;
6         while (curr!=NULL)
7         {
8             temp = curr->next;
9             curr->next = prev;
10            prev = curr;
11            curr = temp;
12        }
13        return prev;
14    }
15
16    ListNode* reverseInGroups(ListNode* head, int k)
17    {
18        ListNode *curr = head;
19        int currlen = 1;
20        if(head == NULL) return head;
21        while(curr->next!=NULL && currlen<k ){
22            curr=curr->next;
23            currlen+=1;
24        }
25        if(currlen<k) return head;
26        ListNode *tempNode = curr->next;
27        curr->next = NULL;
28
29        // start linking
30        ListNode *tempList = reverseInGroups(tempNode,k);
31        ListNode *prev = reverseList(head);
32        head->next = tempList;
33        return prev;
34    }
35
36    ListNode* reverseKGroup(ListNode* head, int k) {
37        return reverseInGroups(head,k);
38    }
39};
```

# 17 Design linked list → Implementation of Doubly Linked list

Code →

```

class Node{
public:
    int val;
    Node* prev;
    Node* next;
    Node(int val){
        this->val=val;
        prev = nullptr;
        next = nullptr;
    }
};

class MyLinkedList {
public:
    Node *head;
    Node *tail;
    MyLinkedList(){
        head = nullptr;
        tail = nullptr;
    }

    int get(int index){
        if(head == NULL)    return -1;
        Node *temp = head;
        int count = 0;
        while(temp!=NULL){
            temp=temp->next;
            count++;
        }
        if(index>=count)    return -1;
        temp = head;
        while(temp != NULL && index>0){
            temp=temp->next;
            index--;
        }
        return temp->val;
    }

    void addAtHead(int val){
        Node *newNode = new Node(val);
        if(head == NULL){
            head = newNode;
            tail = newNode;
        } else {
            newNode->next = head;
            head->prev = newNode;
            head = newNode;
        }
    }

    void addAtTail(int val){
        Node *temp = head;
        if(head == NULL){
            Node *newNode = new Node(val);
            head = newNode;
            tail = newNode;
            return;
        }
        while(temp->next!=NULL){
            temp = temp->next;
        }
        Node *newNode = new Node(val);
        temp->next = newNode;
        newNode->prev = temp;
        tail = newNode;
    }
}

```

```

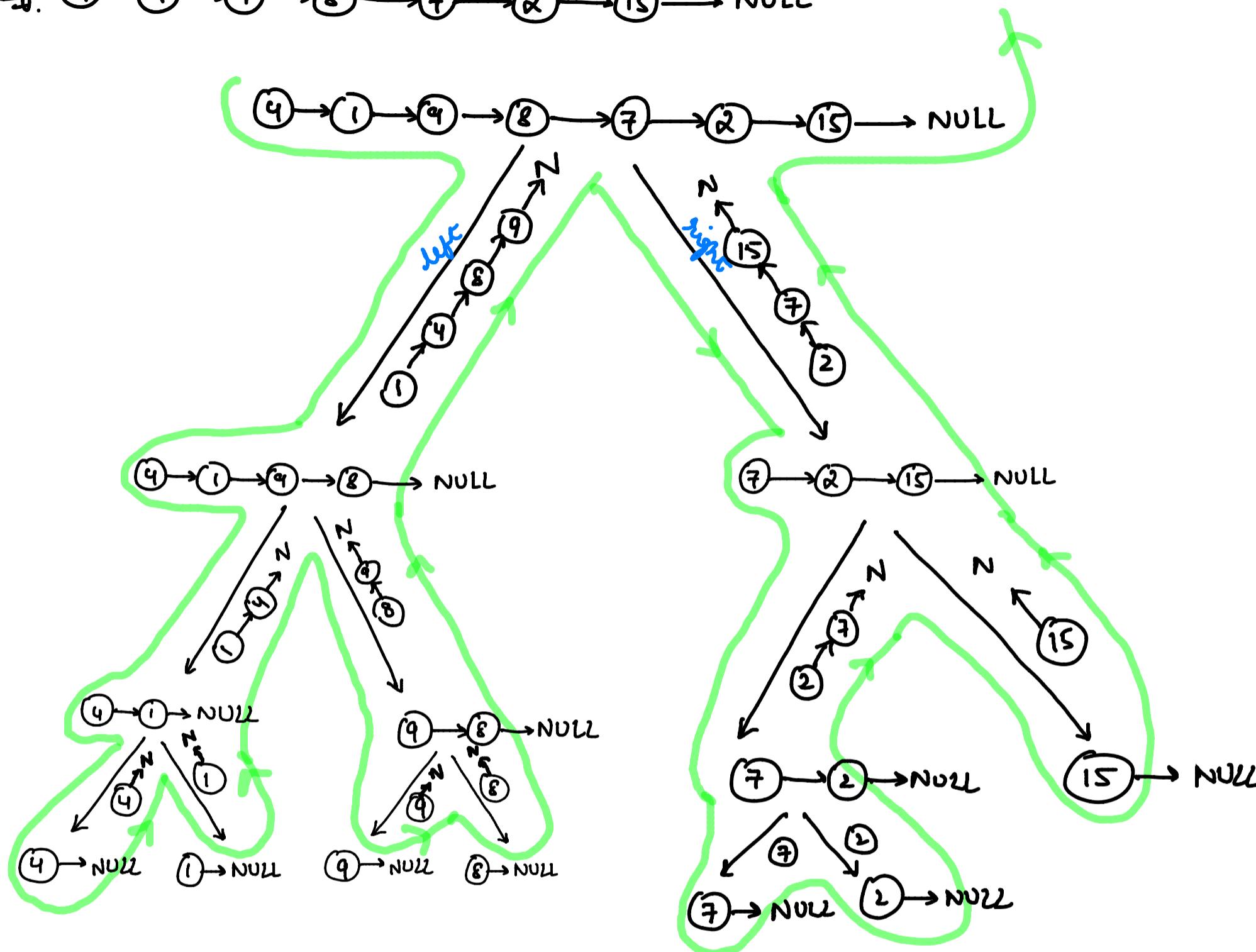
void addAtIndex(int index, int val){
    Node *temp = head;
    int count = 0;
    while(temp != NULL){
        temp = temp->next;
        count++;
    }
    if(index>count) return ;
    if(index==0){
        addAtHead(val);
        return;
    } else if(count == index){
        addAtTail(val);
        return;
    } else {
        temp = head;
        while(temp != NULL && index>0){
            temp = temp->next;
            index--;
        }
        Node* newNode = new Node(val);
        Node* temp2 = temp->prev;
        temp->prev->next = newNode;
        temp->prev = newNode;
        newNode->prev = temp2;
        newNode->next = temp;
    }
}

void deleteAtIndex(int index) {
    Node* temp = head;
    int count = 0;
    while(temp != NULL){
        temp=temp->next;
        count++;
    }
    if(index>=count) return;
    if(count==1 && index==0){
        head = NULL;
        return;
    } else if(count-1 == index){
        tail = tail->prev;
        tail->next = NULL;
        return;
    } else {
        if(index==0){
            head->next->prev = NULL;
            head = head->next;
            return;
        }
        temp=head;
        while(temp!=NULL && index>0){
            temp = temp->next;
            index--;
        }
        Node* temp2 = temp->next;
        temp->prev->next = temp2;
        temp->next->prev = temp->prev;
    }
}

```

18 Sort List → By following Merge Sort.

Eg.  $4 \rightarrow 1 \rightarrow 9 \rightarrow 8 \rightarrow 7 \rightarrow 2 \rightarrow 15 \rightarrow \text{NULL}$



In the last step while returning from both branches we have,

$\text{left} = 1 \rightarrow 4 \rightarrow 8 \rightarrow 9 \rightarrow \text{NULL}$  &  $\text{right} = 2 \rightarrow 7 \rightarrow 15 \rightarrow \text{NULL}$

so create dummy node & merge, i.e.  $(-1 \rightarrow 1 \rightarrow 2 \rightarrow 4 \rightarrow 7 \rightarrow 8 \rightarrow 9 \rightarrow 15 \rightarrow \text{NULL})$

return  $\text{dummy} \rightarrow \text{next}$ ,  $(1 \rightarrow 2 \rightarrow 4 \rightarrow 7 \rightarrow 8 \rightarrow 9 \rightarrow 15 \rightarrow \text{NULL})$

\* The same happens at every intermediate merge

code →

$Tc \rightarrow O(m+n)$

$Sc \rightarrow O(n)$



```
1 class Solution {
2 public:
3     ListNode* merge(ListNode* l1, ListNode* l2) {
4         ListNode *dummy = new ListNode(-1);
5         ListNode *curr = dummy;
6         while(l1 && l2){
7             if(l1->val < l2->val){
8                 curr->next = l1;
9                 l1 = l1->next;
10            } else {
11                curr->next = l2;
12                l2 = l2->next;
13            }
14            curr = curr->next;
15        }
16        if(l1) curr->next = l1;
17        if(l2) curr->next = l2;
18
19        return dummy->next;
20    }
21
22    ListNode* sortList(ListNode* head) {
23        if(!head || !head->next) return head;
24
25        ListNode *slow = head;
26        ListNode *fast = head->next;
27        while(fast && fast->next) {
28            slow = slow->next;
29            fast = fast->next->next;
30        }
31        // dividing the lists into 2 parts
32        fast = slow->next;
33        slow->next = NULL;
34
35        // sort & merge
36        head = sortList(head);
37        fast = sortList(fast)
38        return merge(head, fast);
39    }
40};
```

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