# **DUALITY**

### Lagrangian

**standard form problem** (not necessarily convex)

minimize 
$$f_0(x)$$
  
subject to  $f_i(x) \leq 0, \quad i = 1, \dots, m$   
 $h_i(x) = 0, \quad i = 1, \dots, p$ 

variable  $x \in \mathbf{R}^n$ , domain  $\mathcal{D}$ , optimal value  $p^*$ 

The basic idea: take the constraints into account by augmenting the objective function with a weighted sum of the constraint functions.

### Lagrangian

**Lagrangian:**  $L: \mathbb{R}^n \times \mathbb{R}^m \times \mathbb{R}^p \to \mathbb{R}$ , with  $\operatorname{dom} L = \mathcal{D} \times \mathbb{R}^m \times \mathbb{R}^p$ ,

$$L(x, \lambda, \nu) = f_0(x) + \sum_{i=1}^{m} \lambda_i f_i(x) + \sum_{i=1}^{p} \nu_i h_i(x)$$

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$$L(x, \lambda, \nu) = f_0(x) + \sum_{i=1}^{m} \lambda_i f_i(x) + \sum_{i=1}^{p} \nu_i h_i(x)$$

- weighted sum of objective and constraint functions
- $\lambda_i$  is Lagrange multiplier associated with  $f_i(x) \leq 0$
- $\nu_i$  is Lagrange multiplier associated with  $h_i(x) = 0$

### Lagrange dual function

Lagrange dual function:  $g: \mathbb{R}^m \times \mathbb{R}^p \to \mathbb{R}$ ,

$$g(\lambda, \nu) = \inf_{x \in \mathcal{D}} L(x, \lambda, \nu) = \inf_{x \in \mathcal{D}} \left( f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x) \right)$$

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concave

**lower bound property:** if  $\lambda \succeq 0$ , then  $g(\lambda, \nu) \leq p^*$ 

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concave

**lower bound property:** if  $\lambda \succeq 0$ , then  $g(\lambda, \nu) \leq p^*$ 

proof: if  $\tilde{x}$  is feasible and  $\lambda \succeq 0$ , then

$$f_0(\tilde{x}) \ge L(\tilde{x}, \lambda, \nu) \ge \inf_{x \in \mathcal{D}} L(x, \lambda, \nu) = g(\lambda, \nu)$$

minimizing over all feasible  $\tilde{x}$  gives  $p^* \geq g(\lambda, \nu)$ 

#### Least-squares solution of linear equations

minimize 
$$x^T x$$
  
subject to  $Ax = b$ ,

• Lagrangian is  $L(x,\nu) = x^T x + \nu^T (Ax - b)$ 

#### Least-squares solution of linear equations

minimize  $x^T x$ subject to Ax = b,

Convex quadratic function optimality condition

• Lagrangian is  $L(x, \nu) = x^T x + \nu^T (Ax - b)$ 

$$\nabla_x L(x, \nu) = 2x + A^T \nu = 0$$



$$x = -(1/2)A^T \nu$$

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$$\nabla_x L(x,\nu) = 2x + A^T \nu = 0$$
 
$$x = -(1/2)A^T \nu$$
 
$$g(\nu) = L((-1/2)A^T \nu, \nu) = -\frac{1}{4} \nu^T A A^T \nu - b^T \nu$$
 concave

lower bound property:  $p^{\star} \geq -(1/4)\nu^T A A^T \nu - b^T \nu$  for all  $\nu$ 

#### Standard form LP

minimize 
$$c^T x$$
  
subject to  $Ax = b$   
 $x \succ 0$ ,

$$L(x, \lambda, \nu) = c^T x - \sum_{i=1}^n \lambda_i x_i + \nu^T (Ax - b) = -b^T \nu + (c + A^T \nu - \lambda)^T x.$$

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$$g(\lambda,\nu) = \inf_{x} L(x,\lambda,\nu) = -b^{T}\nu + \inf_{x} (c + A^{T}\nu - \lambda)^{T}x,$$

$$g(\lambda, \nu) = \begin{cases} -b^T \nu & A^T \nu - \lambda + c = 0 \\ -\infty & \text{otherwise.} \end{cases}$$



### Dual problem

For each pair  $(\lambda, \nu)$  with  $\lambda \succeq 0$ , the Lagrange dual function gives us a lower bound on the optimal value  $p^*$  of the optimization problem

What is the best lower bound that can be obtained from the Lagrange dual function?

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#### Lagrange dual problem

maximize  $g(\lambda, \nu)$  subject to  $\lambda \succeq 0$ .

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subject to  $f_i(x) \leq 0, \quad i = 1, ..., m$   
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- $\lambda$ ,  $\nu$  are dual feasible if  $\lambda \succeq 0$ ,  $(\lambda, \nu) \in \operatorname{dom} g$
- finds best lower bound on  $p^*$ , obtained from Lagrange dual function
- a convex optimization problem; optimal value denoted  $d^*$

#### Lagrange dual of standard form LP

minimize 
$$c^T x$$
  
subject to  $Ax = b$   
 $x \succeq 0$ 

$$g(\lambda, \nu) = \begin{cases} -b^T \nu & A^T \nu - \lambda + c = 0 \\ -\infty & \text{otherwise.} \end{cases}$$

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maximize 
$$g(\lambda, \nu) = \begin{cases} -b^T \nu & A^T \nu - \lambda + c = 0 \\ -\infty & \text{otherwise} \end{cases}$$
 maximize  $-b^T \nu$  subject to  $A^T \nu - \lambda + c = 0$  subject to  $\lambda \succeq 0$ .

maximize 
$$-b^T \nu$$
  
subject to  $A^T \nu - \lambda + c = 0$   
 $\lambda \succeq 0$ .

$$\begin{array}{ll} \text{maximize} & -b^T \nu \\ \text{subject to} & A^T \nu + c \succeq 0, \end{array}$$



#### Lagrange dual of inequality form LP

minimize 
$$c^T x$$
  $\to L(x,\lambda) = c^T x + \lambda^T (Ax - b) = -b^T \lambda + (A^T \lambda + c)^T x$ , subject to  $Ax \leq b$ .

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  $L(x,\lambda) = c^T x + \lambda^T (Ax - b) = -b^T \lambda + (A^T \lambda + c)^T x$ , subject to  $Ax \leq b$ .

$$g(\lambda) = \inf_{x} L(x, \lambda) = -b^{T} \lambda + \inf_{x} (A^{T} \lambda + c)^{T} x.$$

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$$g(\lambda) = \begin{cases} -b^T \lambda & A^T \lambda + c = 0 \\ -\infty & \text{otherwise.} \end{cases}$$

maximize 
$$-b^T \lambda$$
  
subject to  $A^T \lambda + c = 0$   
 $\lambda \succeq 0$ ,

# Weak Duality

maximize  $g(\lambda, \nu)$  subject to  $\lambda \succeq 0$ .

minimize 
$$f_0(x)$$
  
subject to  $f_i(x) \leq 0, \quad i = 1, \dots, m$   
 $h_i(x) = 0, \quad i = 1, \dots, p,$ 

weak duality:  $d^{\star} \leq p^{\star}$ 

always holds (for convex and nonconvex problems)

$$p^{\star} - d^{\star} \longrightarrow$$
 optimal duality gap

• can be used to find nontrivial lower bounds for difficult problems

# Strong Duality

strong duality:  $d^{\star} = p^{\star}$ 

- does not hold in general
- (usually) holds for convex problems
- conditions that guarantee strong duality in convex problems are called constraint qualifications

# Slater's constraint qualification

#### Convex problem

minimize 
$$f_0(x)$$
  
subject to  $f_i(x) \le 0$ ,  $i = 1, ..., m$ ,  
 $Ax = b$ ,

if it is strictly feasible, i.e.,

$$\exists x \in \mathbf{int}\, \mathcal{D}: \qquad f_i(x) < 0, \quad i=1,\dots,m, \qquad Ax = b$$
 strong duality holds

• also guarantees that the dual optimum is attained (if  $p^* > -\infty$ )

# Slater's constraint qualification

Slater's condition can be refined when some of the inequality constraint functions are affine.

If the first k constraint functions  $f_1, \ldots, f_k$  are affine,

$$\exists x \in \mathbf{int} \, \mathcal{D} : f_i(x) \leq 0, \quad i = 1, \dots, k, \qquad f_i(x) < 0, \quad i = k+1, \dots, m, \qquad Ax = b.$$

refined Slater condition reduces to feasibility when the constraints are all linear equalities and inequalities

Inequality form LP

primal problem

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax \preceq b \end{array}$$

#### Inequality form LP

#### primal problem

 $\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax \leq b \end{array}$ 

dual function

$$g(\lambda) = \inf_{x} \left( (c + A^T \lambda)^T x - b^T \lambda \right) = \begin{cases} -b^T \lambda & A^T \lambda + c = 0 \\ -\infty & \text{otherwise} \end{cases}$$

dual problem

$$\begin{array}{ll} \text{maximize} & -b^T \lambda \\ \text{subject to} & A^T \lambda + c = 0, \quad \lambda \succeq 0 \end{array}$$

#### Quadratic program

```
primal problem (assume P \in \mathbf{S}^n_{++})  \begin{aligned} & \text{minimize} & & x^T P x \\ & \text{subject to} & & A x \preceq b \end{aligned}
```

#### Quadratic program

### **primal problem** (assume $P \in \mathbf{S}_{++}^n$ )

minimize 
$$x^T P x$$
 subject to  $Ax \leq b$ 

#### dual function

$$g(\lambda) = \inf_{x} \left( x^T P x + \lambda^T (Ax - b) \right) = -\frac{1}{4} \lambda^T A P^{-1} A^T \lambda - b^T \lambda$$

#### dual problem

$$\begin{array}{ll} \text{maximize} & -(1/4)\lambda^TAP^{-1}A^T\lambda - b^T\lambda \\ \text{subject to} & \lambda \succeq 0 \end{array}$$

• in fact,  $p^* = d^*$  always

assume strong duality holds,  $x^*$  is primal optimal,  $(\lambda^*, \nu^*)$  is dual optimal

$$f_0(x^*) = g(\lambda^*, \nu^*) = \inf_x \left( f_0(x) + \sum_{i=1}^m \lambda_i^* f_i(x) + \sum_{i=1}^p \nu_i^* h_i(x) \right)$$

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$$\leq f_0(x^*) + \sum_{i=1}^m \lambda_i^* f_i(x^*) + \sum_{i=1}^p \nu_i^* h_i(x^*)$$

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$$\leq f_0(x^*)$$

assume strong duality holds,  $x^*$  is primal optimal,  $(\lambda^*, \nu^*)$  is dual optimal

$$f_0(x^\star) = g(\lambda^\star, \nu^\star) = \inf_x \left( f_0(x) + \sum_{i=1}^m \lambda_i^\star f_i(x) + \sum_{i=1}^p \nu_i^\star h_i(x) \right)$$
 
$$\leq f_0(x^\star) + \sum_{i=1}^m \lambda_i^\star f_i(x^\star) + \sum_{i=1}^p \nu_i^\star h_i(x^\star)$$
 
$$\leq f_0(x^\star)$$
 hence, the two inequalities hold with equality

•  $x^*$  minimizes  $L(x, \lambda^*, \nu^*)$ 

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$$\leq f_0(x^*)$$

hence, the two inequalities hold with equality

- $x^*$  minimizes  $L(x, \lambda^*, \nu^*)$
- $\lambda_i^{\star} f_i(x^{\star}) = 0$  for  $i = 1, \dots, m$  (known as complementary slackness):

$$\lambda_i^{\star} > 0 \Longrightarrow f_i(x^{\star}) = 0, \qquad f_i(x^{\star}) < 0 \Longrightarrow \lambda_i^{\star} = 0$$

# Karush-Kuhn-Tucker(KKT) conditions

if strong duality holds and x,  $\lambda$ ,  $\nu$  are optimal, then they must satisfy the KKT conditions

the following four conditions are called KKT conditions (for a problem with differentiable  $f_i$ ,  $h_i$ ):

- 1. primal constraints:  $f_i(x) \leq 0$ ,  $i = 1, \ldots, m$ ,  $h_i(x) = 0$ ,  $i = 1, \ldots, p$
- 2. dual constraints:  $\lambda \succeq 0$
- 3. complementary slackness:  $\lambda_i f_i(x) = 0$ ,  $i = 1, \ldots, m$
- 4. gradient of Lagrangian with respect to x vanishes:

$$\nabla f_0(x) + \sum_{i=1}^m \lambda_i \nabla f_i(x) + \sum_{i=1}^p \nu_i \nabla h_i(x) = 0$$

### KKT conditions for convex problems

When the primal problem is convex, the KKT conditions are also sufficient for the points to be primal and dual optimal.

if  $\tilde{x}$ ,  $\tilde{\lambda}$ ,  $\tilde{\nu}$  satisfy KKT for a convex problem, then they are optimal:

• from 4th condition (and convexity):  $g(\tilde{\lambda}, \tilde{\nu}) = L(\tilde{x}, \tilde{\lambda}, \tilde{\nu})$ 

$$g(\tilde{\lambda}, \tilde{\nu}) = L(\tilde{x}, \tilde{\lambda}, \tilde{\nu})$$

$$= f_0(\tilde{x}) + \sum_{i=1}^m \tilde{\lambda}_i f_i(\tilde{x}) + \sum_{i=1}^p \tilde{\nu}_i h_i(\tilde{x})$$

$$= f_0(\tilde{x}),$$

• from complementary slackness:  $f_0(\tilde{x}) = L(\tilde{x}, \tilde{\lambda}, \tilde{\nu})$  hence,  $f_0(\tilde{x}) = g(\tilde{\lambda}, \tilde{\nu})$ 

for any convex optimization problem with differentiable objective and constraint functions, any points that satisfy the KKT conditions are primal and dual optimal, and have zero duality gap.

# Example

#### **Equality constrained convex quadratic minimization**

minimize 
$$(1/2)x^T P x + q^T x + r$$
  
subject to  $Ax = b$ ,

$$P \in \mathbf{S}^n_+$$

KKT conditions for this problem

# Example

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minimize 
$$(1/2)x^T P x + q^T x + r$$
  
subject to  $Ax = b$ ,

$$P \in \mathbf{S}^n_+$$

#### KKT conditions for this problem

$$Ax^* = b$$
$$Px^* + q + A^T \nu^* = 0,$$

$$\left[\begin{array}{cc} P & A^T \\ A & 0 \end{array}\right] \left[\begin{array}{c} x^{\star} \\ \nu^{\star} \end{array}\right] = \left[\begin{array}{c} -q \\ b \end{array}\right]$$

## Duality and problem reformulations

- ✓ equivalent formulations of a problem can lead to very different duals
- ✓ reformulating the primal problem can be useful when the dual is difficult to derive, or uninteresting

### The conjugate function

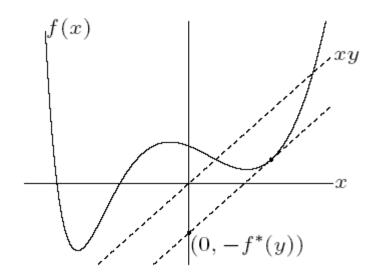
the conjugate of a function f is

$$f^*(y) = \sup_{x \in \mathbf{dom}\, f} (y^T x - f(x))$$

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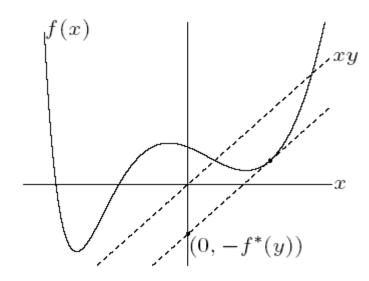


•  $f^*$  is convex (even if f is not)

## The conjugate function

#### the conjugate of a function f is

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•  $f^*$  is convex (even if f is not)

f is convex and differentiable



maximum gap occurs in  $y = \nabla f(x^*)$ 

minimize  $f_0(Ax+b)$ 

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- dual function is constant:  $g = \inf_x L(x) = \inf_x f_0(Ax + b) = p^*$
- we have strong duality, but dual is quite useless

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#### reformulated problem and its dual

dual function follows from

$$g(\nu) = \inf_{x,y} (f_0(y) - \nu^T y + \nu^T A x + b^T \nu)$$

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#### reformulated problem and its dual

dual function follows from

$$g(\nu) = \inf_{x,y} (f_0(y) - \nu^T y + \nu^T A x + b^T \nu) = \begin{cases} -f_0^*(\nu) + b^T \nu & A^T \nu = 0 \\ -\infty & \text{otherwise} \end{cases}$$

$$b^{T}\nu + \inf_{y}(f_{0}(y) - \nu^{T}y) = b^{T}\nu - f_{0}^{*}(\nu),$$

maximize 
$$b^T \nu - f_0^*(\nu)$$
  
subject to  $A^T \nu = 0$ .

### **Dual Norm**

Let  $\|\cdot\|$  be a norm on  $\mathbb{R}^n$ . The associated dual norm, denoted  $\|\cdot\|_*$ , is defined as

$$||z||_* = \sup\{z^T x \mid ||x|| \le 1\}$$

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The dual of the Euclidean norm is the Euclidean norm

The dual of the  $\ell_1$ -norm is the  $\ell_{\infty}$ -norm dual of the  $\ell_{\infty}$ -norm is the  $\ell_1$ -norm

the dual of the  $\ell_p$ -norm is the  $\ell_q$ -norm, where q satisfies 1/p + 1/q = 1.

# Example

Norm

$$f(x) = ||x||$$

$$f^*(y) = \begin{cases} 0 & ||y||_* \le 1 \\ \infty & \text{otherwise,} \end{cases} ||z||_* = \sup\{z^T x \mid ||x|| \le 1\}$$

the dual of the  $\ell_p$ -norm is the  $\ell_q$ -norm, where q satisfies 1/p+1/q=1.

$$||y||_* \le 1 \longrightarrow \text{ for all } x, \ y^T x - ||x|| \le 0 \longrightarrow \text{ maximum value } 0$$

$$z \in \mathbb{R}^n \text{ with } ||z|| \le 1 \text{ and}$$

$$||y||_* > 1 \longrightarrow y^T z > 1. \text{ Taking } x = tz \text{ and letting } t \to \infty, \longrightarrow$$

$$y^T x - ||x|| = t(y^T z - ||z||) \to \infty,$$

### Implicit constraints

#### LP with box constraints:

## Implicit constraints

#### LP with box constraints:

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax = b \\ & -\mathbf{1} \preceq x \preceq \mathbf{1} \end{array}$$

reformulation with box constraints made implicit

$$\begin{array}{ll} \text{minimize} & f_0(x) = \left\{ \begin{array}{ll} c^T x & -1 \preceq x \preceq 1 \\ \infty & \text{otherwise} \end{array} \right. \\ \text{subject to} & Ax = b \end{array}$$

### Implicit constraints

#### LP with box constraints:

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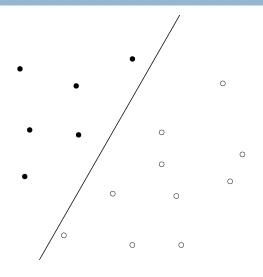
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$$f_0(x) = \begin{cases} c^T x & -1 \leq x \leq 1 \\ \infty & \text{otherwise} \end{cases} \qquad g(\nu) = \inf_{-1 \leq x \leq 1} (c^T x + \nu^T (Ax - b))$$
$$= -b^T \nu - ||A^T \nu + c||_1$$

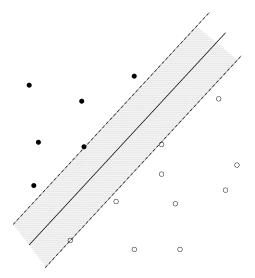
$$\text{maximize } -b^T \nu - \|A^T \nu + c\|_1$$

# CLASSIFICATION



$$g(x) = w^T x + w_0$$

$$\{x_i, i = 1, ..., n\}$$
 two classes,  $\omega_1$  and  $\omega_2$ ,



OBJECTIVE: Maximize distance between two hyperplane  $\frac{2}{||W||^2}$ 

$$\begin{pmatrix} \min & \frac{1}{2} w^T w \\ w^T x_i + w_0 \ge +1 \\ w^T x_i + w_0 \le -1 \end{pmatrix} \quad x \in \begin{cases} \omega_1 \\ \omega_2 \end{cases}$$

OBJECTIVE: Maximize distance between two hyperplane  $\frac{Z}{||W||^2}$ 

$$\begin{pmatrix} \min & \frac{1}{2} w^T w \\ w^T x_i + w_0 \ge +1 \\ w^T x_i + w_0 \le -1 \end{pmatrix} \quad x \in \begin{cases} \omega_1 \\ \omega_2 \end{cases}$$

 $\omega_1$  with corresponding numeric value,  $y_i = +1$ 

 $\omega_2$  with corresponding numeric value,  $y_i = -1$ 

$$w^T x_i + w_0 \ge +1$$
 for  $y_i = +1$   
 $w^T x_i + w_0 \le -1$  for  $y_i = -1$ 

$$y_i(w^T x_i + w_0) \ge 1$$
  $i = 1, ..., n$ 

$$\min \quad \frac{1}{2} w^T w$$

$$y_i(w^T x_i + w_0) \ge 1 \quad i = 1, \dots, n$$

Lagrangian

$$L_p = \frac{1}{2} w^T w - \sum_{i=1}^n \alpha_i (y_i (w^T x_i + w_0) - 1)$$

$$\min \quad \frac{1}{2} w^T w$$

$$y_i(w^T x_i + w_0) \ge 1 \quad i = 1, \dots, n$$

Lagrangian

$$L_p = \frac{1}{2} w^T w - \sum_{i=1}^n \alpha_i (y_i (w^T x_i + w_0) - 1)$$

Differentiating  $L_p$  with respect to  $w_0$  and w

$$\sum_{i=1}^{n} \alpha_i y_i = 0$$

$$w = \sum_{i=1}^{n} \alpha_i y_i x_i$$

Dual Function g 
$$L_D = \sum_{i=1}^n \alpha_i - \frac{1}{2} \sum_{i=1}^n \sum_{j=1}^n \alpha_i \alpha_j y_i y_j x_i^T x_j$$

Dual problem

Max 
$$L_D = \sum_{i=1}^n \alpha_i - \frac{1}{2} \sum_{i=1}^n \sum_{j=1}^n \alpha_i \alpha_j y_i y_j x_i^T x_j$$
  $\alpha_i \geq 0$   $\sum_{i=1}^n \alpha_i y_i = 0$ 

Dual Function g

$$L_D = \sum_{i=1}^{n} \alpha_i - \frac{1}{2} \sum_{i=1}^{n} \sum_{j=1}^{n} \alpha_i \alpha_j y_i y_j x_i^T x_j$$

Dual problem

Max 
$$L_D = \sum_{i=1}^n lpha_i - rac{1}{2} \sum_{i=1}^n \sum_{j=1}^n lpha_i lpha_j y_i y_j x_i^T x_j$$
  $lpha_i \geq 0$   $\sum_{i=1}^n lpha_i y_i = 0$ 

inner products of patterns, *xi* 

nonlinear support vector machines

### KKT conditions

4<sup>th</sup> condition

$$w = \sum_{i=1}^{n} \alpha_i y_i x_i$$

### KKT conditions

4<sup>th</sup> condition

$$w = \sum_{i=1}^{n} \alpha_i y_i x_i$$

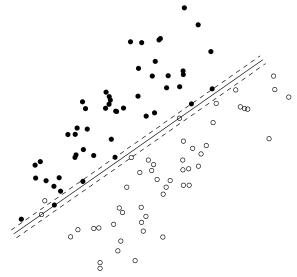
**Complementary Slackness** 

$$\alpha_i(y_i(x_i^Tw+w_0)-1)=0$$

Nonzero dual optimal

$$y_i(x_i^T w + w_0) - 1) = 0$$

**Support Vectors** 



- √two sets of points cannot be linearly separated
- ✓ seek an affine function that approximately classifies the points
- √ relax the constraints by introducing nonnegative variables

$$w^T x_i + w_0 \ge +1 - \xi_i$$
 for  $y_i = +1$   
 $w^T x_i + w_0 \le -1 + \xi_i$  for  $y_i = -1$   
 $\xi_i \ge 0$   $i = 1, ..., n$ 

#### 'regularisation' parameter

$$\frac{1}{2}\mathbf{w}^T\mathbf{w} + C\sum_{i} \xi_{i}$$

$$y_i(\mathbf{w}^T \mathbf{x}_i + w_0) \ge 1 - \xi_i \quad i = 1, ..., n$$
  
 $\xi_i \ge 0 \quad i = 1, ..., n$ 

#### 'regularisation' parameter

$$\frac{1}{2}w^Tw + C\sum_i \xi_i$$

$$\frac{1}{2} \mathbf{w}^T \mathbf{w} + C \sum_{i} \xi_i$$

$$y_i (\mathbf{w}^T \mathbf{x}_i + w_0) \ge 1 - \xi_i \quad i = 1, \dots, n$$

$$\xi_i \ge 0 \qquad \qquad i = 1, \dots, n$$

Lagrangian

$$L_p = \frac{1}{2} \mathbf{w}^T \mathbf{w} + C \sum_{i} \xi_i - \sum_{i=1}^n \alpha_i (y_i (\mathbf{w}^T \mathbf{x}_i + \mathbf{w}_0) - 1 + \xi_i) - \sum_{i=1}^n r_i \xi_i$$

Differentiating  $L_p$  with respect to  $w_0$ , w and zeta

$$\mathbf{w} = \sum_{i=1}^{n} \alpha_i y_i \mathbf{x}_i \qquad \sum_{i=1}^{n} \alpha_i y_i = 0 \qquad C - \alpha_i - r_i = 0$$

Dual Function g

$$L_D = \sum_{i=1}^{n} \alpha_i - \frac{1}{2} \sum_{i=1}^{n} \sum_{j=1}^{n} \alpha_i \alpha_j y_i y_j x_i^T x_j$$

Dual problem

$$L_D = \sum_{i=1}^n \alpha_i - \frac{1}{2} \sum_{i=1}^n \sum_{j=1}^n \alpha_i \alpha_j y_i y_j x_i^T x_j$$
$$\sum_{i=1}^n \alpha_i y_i = 0 \qquad 0 \le \alpha_i \le C$$