

# Computation of Persistent Homology Using the Delaunay-Rips Complex

An efficient family of simplicial complexes for topological data analysis

Amish Mishra and Francis Motta

## Abstract

Topological Data Analysis is an emerging area rooted in theories from Algebraic Topology, which enables researchers to extract discriminating geometric and topological features from data. We give an overview of some of the popular methods of extracting features from point-cloud data, which first requires one to construct a 1-parameter family of spaces on the data using the geometry of the cloud. We demonstrate their benefits and shortcomings and introduce a new, more efficient construction that we name the Delaunay-Rips Complex. We justify conditions on the data that guarantee stability of our method when computing persistent homology. Aided by intuitive examples, we also provide an empirical run-time comparison of the two existing methods with our new algorithm on the computation of the persistence diagrams of some synthetic data sets.

## 1 Introduction

Welcome to our paper!

## 2 Background

### 2.1 Topological Data Analysis

We are living in an information age where data-driven decision making is a huge area of interest. With so much data at our hands, many questions naturally arise.

- How do we extract relevant information from the data?
- How do we even know what is relevant and what is not?
- If we are unable to visualize large quantities of data, especially data in high dimensions, then how do we know what sort of data set we are inspecting?
- Further, how can we compare information extracted from one data set with another?

These are the sorts of difficult and fascinating questions tackled in the field of Topological Data Analysis (TDA). From the name itself, TDA hints at leveraging ideas borrowed from topology with data analysis techniques to measure and quantify qualitative features of data. At a more nuanced level, TDA appears as the child of Algebraic Topology, Computer Science, Statistics, Data Analysis, and Computational Geometry. Results from each field have found beautiful applications in TDA and have shed new light on applicability of theoretical results from mathematics.

### 2.2 Simplicial Complexes

The idea behind extracting topological features from a given point-cloud requires there to be a method of assigning some sort of “shape” to the data. Only then are we able to study its associated topological properties. The scenario we face is that we have a finite dimensional metric space out of the point-cloud that we need to assign a shape to. With the human eye, we may be able to make out some appropriate surface our data set can live on. However, this situation pokes at a fundamental question that arises when

computing: how can we get a computer to do what a human can do? We need to introduce the idea of a *simplicial complex*, a computational efficient way for a computer to build a surface onto a data set.

**Definition 2.1.** A simplicial complex is a collection  $K$  of non-empty subsets of a set  $K_0$  such that  $\{v\} \in K$  for all  $v \in K_0$ , and if  $\sigma \in K$  then  $\tau \in K$  for all  $\tau \subseteq \sigma$ . The elements of  $K_0$  are called vertices of  $K$ , and the elements of  $K$  are called simplices. Additionally, we say that a simplex has dimension  $p$  or is a  $p$ -simplex if it has cardinality of  $p + 1$ . We use  $K_p$  to denote the collection of  $p$ -simplices. The  $k$ -skeleton of  $K$  is the union of the sets  $K_p$  for all  $p \in \{0, 1, \dots, k\}$ . If  $\tau$  and  $\sigma$  are simplices such that  $\tau \subset \sigma$ , then we call  $\tau$  a face of  $\sigma$ , and we say that  $\tau$  is a face of  $\sigma$  of codimension  $k'$  if the dimensions of  $\tau$  and  $\sigma$  differ by  $k'$ . The dimension of  $K$  is defined as the maximum of the dimensions of its simplices. A map of simplicial complexes,  $f : K \rightarrow L$ , is a map  $f : K_0 \rightarrow L_0$  such that  $f(\sigma) \in L$  for all  $\sigma \in K$ . [EH10]

## 2.3 Simplicial Homology

## 2.4 Vietoris-Rips Complex

One of the simplest ways to build a complex on a data set  $X$  is by considering the pairwise distance between the points. The approach described here is an algorithmic, bottom-up approach that adds higher and higher dimensional simplices to the complex for a fixed scale. For a given scale  $\varepsilon > 0$ , if  $d(x, x') \leq 2\varepsilon$  for  $x, x' \in X$ , then we add the edge between  $x$  and  $x'$  into our complex. Once all of the edges are added, we add the higher dimensional simplices if their faces are already in the complex. That is, we add the  $k$ -simplex  $\sigma = \{x_0, x_1, \dots, x_k\}$  to the complex if every subset  $u \subset \sigma$  is already in the complex. Formally, we define the Vietoris-Rips complex [Ott+17] for scale  $\varepsilon > 0$

$$VR_\varepsilon(X) = \{\sigma \subseteq X \mid d(x, x') \leq 2\varepsilon, \forall x, x' \in \sigma\}.$$

## 2.5 Delaunay Triangulation

Although the Vietoris-Rips complex is simple to implement, constructing it on data sets with large numbers of points results in computation drawback. As the scale increases, we see that adding certain simplices does not affect the homology of the point cloud. We need some way to “weed” out these extraneous simplices as we construct our complex to increase computational efficiency. Turning to a tool of Computational Geometry, we incorporate the Delaunay Triangulation in our construction. Our definition is adapted from “A roadmap for the computation of persistent homology” [Ott+17]. Assume our data  $X$  lives in the space  $\mathbb{R}^n$ . Let  $x \in X$ . We define

$$V_x = \{p \in \mathbb{R}^d \mid d(p, x) \leq d(p, x') \forall x' \in X\}.$$

Each  $V_x$  is called a Voronoi cell. Note that  $\{V_x\}_{x \in X}$  forms a cover of  $\mathbb{R}^n$ . This cover is known as the Voronoi decomposition of  $\mathbb{R}^n$  with respect to  $X$ . To construct the Delaunay triangulation from this cover, we connect  $x, x' \in X$  with an edge if  $V_x$  and  $V_{x'}$  are neighbors (that is, the Voronoi cells share a wall). When the points in  $X$  are in general position, this gives us a graph (1-skeleton) on  $X$  that is known as the Delaunay Triangulation. Formally, we define [EH10]

$$Del(X) = \{\sigma \subset X \mid \bigcap_{u \in \sigma} V_u \neq \emptyset\}.$$

We will use  $Del(X)$  as the underlying graph structure when defining the Delaunay-Rips complex in section 3.1.

As an alternate definition to the Delaunay Triangulation that we will employ in section 3.5, we have the following

**Definition 2.2.** For a set  $X$  of points in  $\mathbb{R}^d$ , a Delaunay triangulation is a triangulation  $DT(X)$  such that no point in  $X$  is inside the circum-hypersphere of any simplex in  $DT(X)$ .

## 2.6 Persistence

# 3 Delaunay-Rips Complex

## 3.1 Definition and Construction

The Delaunay-Rips complex is our alternative method of building a complex on a data set  $X$ . The idea is similar to the construction of the Delaunay-Čech complex defined in [BE16]. Delaunay-Rips utilizes the conceptual simplicity of the Vietoris-Rips complex while cutting down on the number of high dimensional and extraneous simplices. This computational speed-up is by virtue of using the Delaunay Triangulation as the “backbone” of building the Vietoris-Rips complex on  $X$ . The idea is that we build the Vietoris-Rips complex on  $X$  but only add edges if the edges occur in the Delaunay 1-skeleton of the point cloud. The higher dimensional  $k$ -simplices are then added the traditional way they are in section 2.4. Formally, we define the Delaunay-Rips Complex for a given scale  $\varepsilon > 0$

$$DR_\varepsilon(X) = \{\sigma \subseteq Del(X) \mid d(x, x') \leq 2\varepsilon, \forall x, x' \in \sigma\}.$$

## 3.2 Example Data Set

1. Demonstrate construction on a small data-set (5-8 point data-set).

## 3.3 Run-time Analysis Comparison

In this section, we immediately show some empirical results of the performance of computing persistence using the Delaunay-Rips filtration on varying datasets. The runtimes are measured as the time between the data set being inputted into the corresponding algorithm and the persistence diagram being produced.

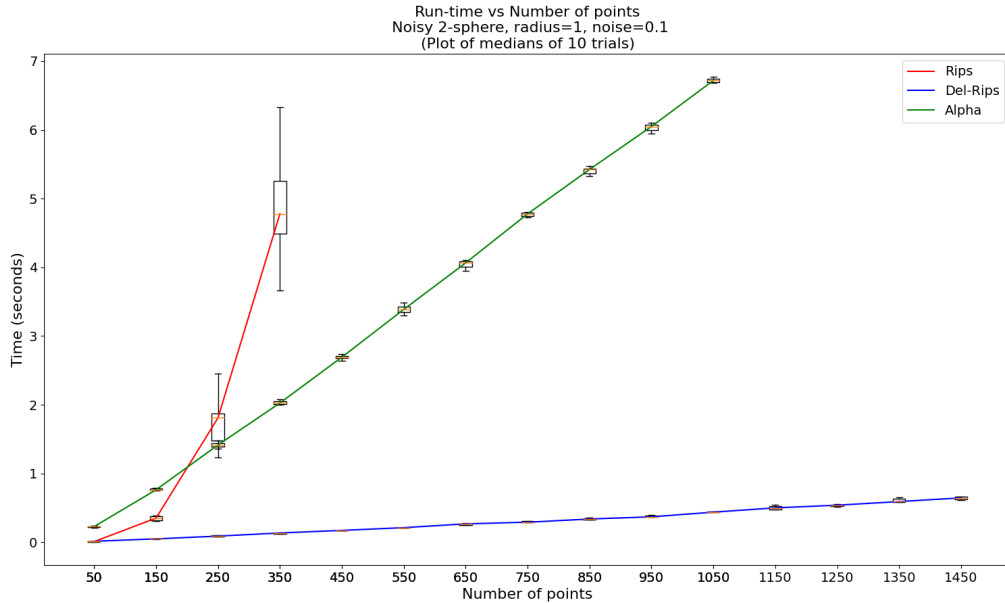


Figure 1: Runtime comparison of Rips, Alpha, and Delaunay-Rips as number of points are increased.

## 3.4 Persistence Diagram Instability

The Delaunay-Rips construction gains computational efficiency at the cost of stability. We demonstrate a simple, yet clear example of how this instability can arise. In figure 3, we can visually see a particular

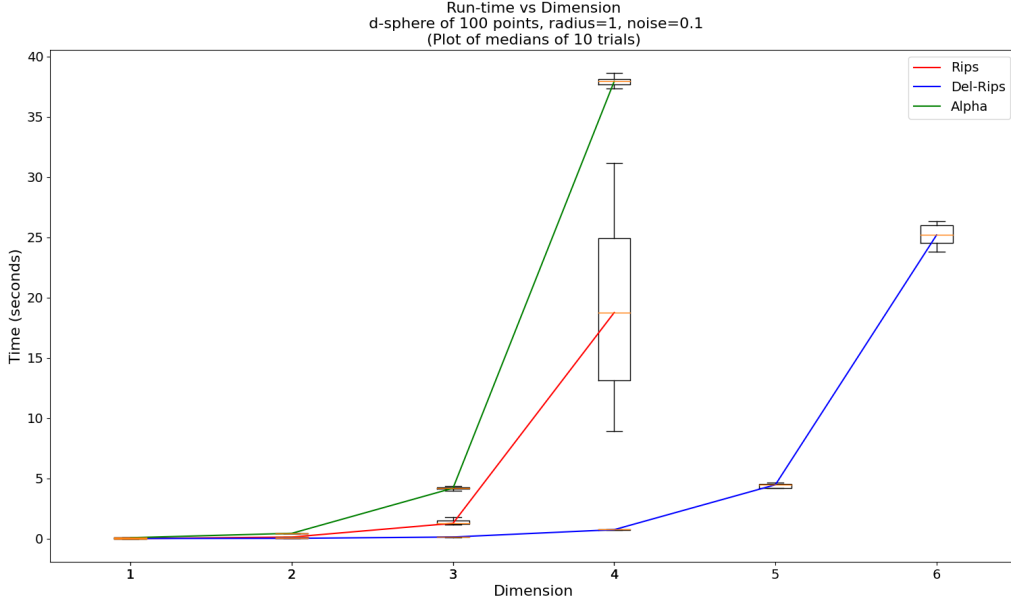


Figure 2: Runtime comparison of Rips, Alpha, and Delaunay-Rips as dimension of data increases.

configuration of four points giving a radically different Persistence Diagram as we move the right-most point across the circumcircle of the other 3 points. In the figure, we have marked the Delaunay Triangulation of the points to show case the point at which an edge flip occurs (namely when all four points lie on the same circle). We now proceed to formally prove the instability of this particular configuration of points.

**Definition 3.1.** Let  $X$  and  $Y$  be two non-empty subsets of a metric space  $(M, d)$ . We define their **Hausdorff distance**  $d_H(X, Y)$  by

$$d_H(X, Y) := \max \left\{ \sup_{x \in X} d(x, Y), \sup_{y \in Y} d(X, y) \right\}.$$

**Definition 3.2.** If  $X$  and  $Y$  are two compact metric spaces, then the **Gromov-Hausdorff distance**  $d_{GH}(X, Y)$  is defined to be the infimum of all numbers  $d_H(f(X), g(Y))$  for all metric spaces  $M$  and all isometric embeddings  $f : X \rightarrow M$  and  $g : Y \rightarrow M$ .

**Definition 3.3.** Let  $X$  and  $Y$  be two persistence diagrams. Let  $\eta : X \rightarrow Y$  be a bijection and  $\|x - y\|_\infty := \max\{|x_1 - y_1|, |x_2 - y_2|\}$  for  $x = (x_1, x_2), y = (y_1, y_2)$ . We define the **bottleneck distance** between the diagrams as

$$W_\infty(X, Y) = \inf_{\eta : X \rightarrow Y} \sup_{x \in X} \|x - \eta(x)\|_\infty.$$

Let  $(\mathcal{P}, d_{GH})$  be the space of point clouds equipped with the Gromov-Hausdorff metric and let  $(\mathcal{D}, W_\infty)$  be the space of Persistence Diagrams equipped with the bottle neck metric. Let

$$Pers_1 : \mathcal{P} \rightarrow \mathcal{D}$$

where  $Pers_1(P)$  is the persistence diagram showing the  $H_1$  classes of the point cloud  $P$  constructed using the Delaunay-Rips complex. Our example comes from 4 points taken in  $\mathbb{R}^2$  where the instability is demonstrated as the discontinuity of  $Pers_1$ .

Let  $P \in \mathcal{P}$  as  $P = \{(-1, 0), (\frac{1}{2}, \frac{\sqrt{3}}{2}), (\frac{1}{2}, -\frac{\sqrt{3}}{2}), (1, 0)\}$ . Note that the points all lie on the unit circle, so the Delaunay 1-skeleton has an edge between every pair of points (See figure). Thus,  $Pers_1(P)$  has no  $H_1$  class with non-zero persistence using the Delaunay-Rips filtration, as can be verified by the reader.

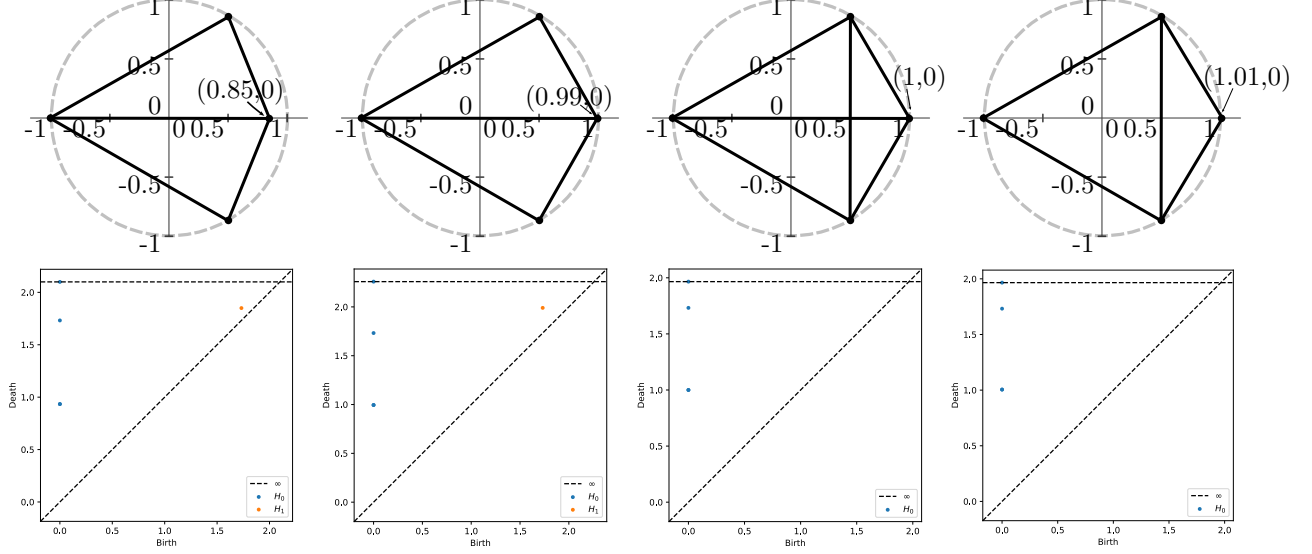


Figure 3: Persistence Diagrams of 4 point example

Fix  $\varepsilon = 0.1$ . We now show that for any  $\delta > 0$ , there exists  $P' \in \mathcal{P}$  such that  $d_{GH}(P, P') < \delta$ , but  $W_\infty(Pers_1(P), Pers_1(P')) \geq \varepsilon$ . Take  $P' = \{(-1, 0), (\frac{1}{2}, \frac{\sqrt{3}}{2}), (\frac{1}{2}, -\frac{\sqrt{3}}{2}), (1-x, 0)\}$  with  $0 < x < \delta < \frac{2-\sqrt{3}}{2}$ . This is a small perturbation of  $P$  by pushing the point  $(1, 0)$  inside the unit circle thereby putting the points in general position. We only work with  $\delta < 2 - \sqrt{3}$  so that  $Pers_1(P')$  maintains an  $H_1$  class with non-zero persistence; for this example to work, we further need  $\delta < \frac{2-\sqrt{3}}{2}$ . We just compute the Hausdorff distance  $d_H$  between  $P$  and  $P'$  in the plane taking the isometric embedding of  $P$  to be the map that sends each of its points to itself in  $\mathbb{R}^2$  and the same embedding for  $P'$ . Since the Gromov-Hausdorff distance is the infimum of  $d_H(f(P), g(P'))$  over all isometric embeddings  $f : P \rightarrow X$  and  $g : P' \rightarrow X$  into any metric space  $X$ ,  $d_H(P, P')$  serves as an upper bound for  $d_{GH}(P, P')$ . We find that

$$d_{GH}(P, P') \leq d_H(P, P') = x < \delta.$$

Recall that  $Pers_1(P)$  has no  $H_1$  class with non-zero persistence. Thus, to compute  $W_\infty(Pers_1(P), Pers_1(P'))$ , we must match the  $H_1$  class of  $Pers_1(P')$  with the diagonal. The  $H_1$  class of  $Pers_1(P')$  has birth  $\sqrt{3}$  and death  $2-x$  as calculated in the Appendix, section A.1. Using the max norm, we find

$$d := W_\infty(Pers_1(P), Pers_1(P')) = 2 - x - \sqrt{3} \geq 2 - \frac{2-\sqrt{3}}{2} - \sqrt{3} \geq 0.1 = \varepsilon.$$

Hence, our map  $Pers_1$  is discontinuous at  $P$ . This gives us insight into when the Delaunay-Rips construction of the Persistence Diagram experiences instability—namely when points are not in general position. We now have motivation to ask if we have stability of the PD when the underlying Delaunay-Rips complex does not change under a perturbation of the point cloud.

### 3.5 Stability in a Neighborhood

1. What is the best our method can do? Use knowledge on stability of Delaunay Triangulation.

Let  $P$  denote our point cloud and  $P'$  denote the perturbed point cloud. We use  $|xy|$  notation to denote the Euclidean distance between points  $x, y \in \mathbb{R}^n$ . Let  $V_x$  denote the closed Voronoi region for the point  $x \in P$ . Here is a lemma that should work in  $\mathbb{R}^n$ :

**Lemma 3.4.** *Two points  $p_i, p_j \in P$  are strong Voronoi neighbors if and only if there exists an  $m$  such that*

$$\max\{|mp_i|, |mp_j|\} < \min_{k \neq i, j} |mp_k|.$$

We claim

**Theorem 3.5.** *Let  $p_i, p_j \in P$  be strong Voronoi neighbors and  $m \in V_{p_i} \cap V_{p_j}$ . For  $\varepsilon = \min\{\frac{1}{4}(\min_{k \neq i,j} |mp_k| - |mp_j|), \frac{1}{2} \min_{i \neq j} |p_i p_j|\}$ , an  $\varepsilon$ -perturbation of  $P$  leaves  $p_i$  and  $p_j$  as strong Voronoi neighbors.*

*Proof.* Let  $p_i, p_j \in P$  satisfy Lemma 3.4 with  $m \in V_{p_i} \cap V_{p_j}$ . Note that  $0 < \varepsilon < \frac{1}{2} \min_{i \neq j} |p_i p_j|$  to ensure a unique correspondence between the points of  $P$  and the points of  $P'$ . We begin with the conclusion of Lemma 3.4:

$$\begin{aligned}
\max\{|mp_i|, |mp_j|\} &< \min_{k \neq i,j} |mp_k| \\
2 \max\{|mp_i|, |mp_j|\} &< 2 \min_{k \neq i,j} |mp_k| \\
\min_{k \neq i,j} |mp_k| + 3 \max\{|mp_i|, |mp_j|\} &< 3 \min_{k \neq i,j} |mp_k| + \max\{|mp_i|, |mp_j|\} \\
\frac{1}{4} \min_{k \neq i,j} |mp_k| + \frac{3}{4} \max\{|mp_i|, |mp_j|\} &< \frac{3}{4} \min_{k \neq i,j} |mp_k| + \frac{1}{4} \max\{|mp_i|, |mp_j|\} \\
\max\{|mp_i|, |mp_j|\} + \frac{1}{4} (\min_{k \neq i,j} |mp_k| - \max\{|mp_i|, |mp_j|\}) &< \min_{k \neq i,j} |mp_k| - \frac{1}{4} (\min_{k \neq i,j} |mp_k| - \max\{|mp_i|, |mp_j|\}) \\
\max\{|mp_i|, |mp_j|\} + \varepsilon &\leq \max\{|mp_i|, |mp_j|\} + \frac{1}{4} (\min_{k \neq i,j} |mp_k| - \max\{|mp_i|, |mp_j|\}) \\
&< \min_{k \neq i,j} |mp_k| - \frac{1}{4} (\min_{k \neq i,j} |mp_k| - \max\{|mp_i|, |mp_j|\}) \leq \min_{k \neq i,j} |mp_k| - \varepsilon.
\end{aligned} \tag{1}$$

Now, without loss of generality, let

$$|mp'_i| = \max\{|mp'_i|, |mp'_j|\}.$$

We note by the triangle inequality that

$$\max\{|mp'_i|, |mp'_j|\} \leq |mp_i| + |p_i p'_i| \leq \max\{|mp_i|, |mp_j|\} + \varepsilon. \tag{2}$$

Similarly, we have by the triangle inequality

$$\min_{k \neq i,j} |mp_k| - \varepsilon = \min_{k \neq i,j} (|mp_k| - \varepsilon) \leq \min_{k \neq i,j} (|mp_k| - |p_k p'_k|) \leq \min_{k \neq i,j} |mp'_k|. \tag{3}$$

Putting together equations 1, 2, 3, we have

$$\max\{|mp'_i|, |mp'_j|\} \leq \max\{|mp_i|, |mp_j|\} + \varepsilon < \min_{k \neq i,j} |mp_k| - \varepsilon \leq \min_{k \neq i,j} |mp'_k|.$$

Thus, we now apply Lemma 3.4 and have that  $p'_i$  and  $p'_j$  remain strong Voronoi neighbors.  $\square$

**Lemma 3.6.** *Given a subset  $S = \{p_1, p_2, \dots, p_{d+1}\} \subset P$ , there exists  $\varepsilon > 0$  such that any  $\varepsilon$ -perturbation  $S'$  of  $S$  has  $\{p'_1, p'_2, \dots, p'_{d+1}\} \subset P'$  as linearly independent.*

*Proof.* Let  $\mathcal{S}$  be the unique  $(d-1)$ -sphere containing  $S$  with center  $O$ . We proceed by linear algebra, treating each  $p_i$  as a vector  $\mathbf{p}_i \in \mathbb{R}^d$ . We locate the center  $O \in \mathbb{R}^d$  by finding the intersection of the perpendicular bisecting planes to the line segments  $\overline{\mathbf{p}_2 \mathbf{p}_1}, \overline{\mathbf{p}_3 \mathbf{p}_1}, \dots, \overline{\mathbf{p}_{d+1} \mathbf{p}_1}$ . Thus, we will have  $d$  hyperplanes each in  $d$  variables: this is a system of linear equations that has a solution since our subset  $S$  is in general position (that is, no two perpendicular bisecting hyperplanes are parallel). The vector equation for the perpendicular bisecting hyperplane of  $\overline{\mathbf{p}_i \mathbf{p}_1}$  can be found by determining the normal vector and a point on the hyperplane. A point on the plane we can use is simply the midpoint of  $\overline{\mathbf{p}_i \mathbf{p}_1}$ :

$$\mathbf{b}_i := \frac{\mathbf{p}_i + \mathbf{p}_1}{2}.$$

The normal vector to the plane is the directional vector based at the midpoint and pointing in the direction of one end of the line segment  $\overline{\mathbf{p}_i \mathbf{p}_1}$ . We find the normal vector to be

$$\mathbf{n}_i := \mathbf{p}_i - \mathbf{b}_i = \frac{\mathbf{p}_i - \mathbf{p}_1}{2}.$$

Thus, the vector equation of the perpendicular bisecting hyperplane of  $\overline{\mathbf{p}_i \mathbf{p}_1}$  is

$$\mathbf{n}_i \cdot \mathbf{r} = \mathbf{n}_i \cdot \mathbf{b}_i$$

where  $\mathbf{r}$  is the position vector for an arbitrary point on the hyperplane.

We are now ready to construct the following maps whose composition will output the circumcenter  $O \in \mathbb{R}^d$  as a function of the  $d+1$  points defining the sphere  $\mathcal{S}$ . This construction is instrumental to obtaining  $\varepsilon$  balls around each point that ensures the circumcenter  $O$  remains bounded. Define

$$f : \underbrace{\mathbb{R}^d \times \cdots \times \mathbb{R}^d}_{d+1} \rightarrow \mathcal{M}_d(\mathbb{R}) \times \mathbb{R}^d$$

$$f : (\mathbf{p}_1, \mathbf{p}_2, \dots, \mathbf{p}_{d+1}) \mapsto \left( \begin{bmatrix} \mathbf{n}_2 \\ \mathbf{n}_3 \\ \vdots \\ \mathbf{n}_{d+1} \end{bmatrix}, \begin{bmatrix} \mathbf{n}_2 \cdot \mathbf{b}_2 \\ \mathbf{n}_3 \cdot \mathbf{b}_3 \\ \vdots \\ \mathbf{n}_{d+1} \cdot \mathbf{b}_{d+1} \end{bmatrix} \right)$$

$$g : \text{GL}_d(\mathbb{R}) \times \mathbb{R}^d \rightarrow \mathbb{R}^d$$

$$g : (A, b) \mapsto A^{-1}b.$$

We prove that the map  $f$  is continuous. Let the domain of  $f$  have the standard topology of  $\mathbb{R}^{d(d+1)}$  (i.e.

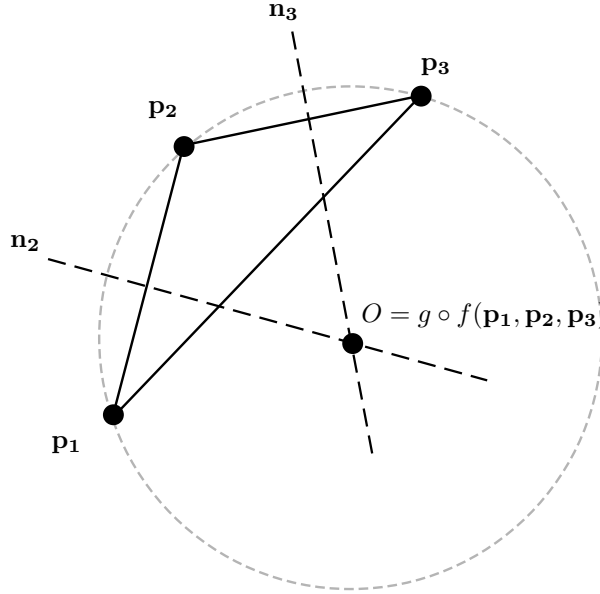


Figure 4: Locating the circumcenter of 3 points in  $\mathbb{R}^2$  as a function of the 3 points.

the open balls in  $\mathbb{R}^{d(d+1)}$  form a basis for the topology). Observing that  $\mathcal{M}_d \cong \mathbb{R}^{d^2}$ , we have  $\mathcal{M}_d(\mathbb{R}) \times \mathbb{R}^d \cong \mathbb{R}^{d^2} \times \mathbb{R}^d \cong \mathbb{R}^{d(d+1)}$ . We let this space have the standard topology of  $\mathbb{R}^{d(d+1)}$ . Note that  $f$  is a polynomial in each component of its output by how it is defined. Thus,  $f$  itself is continuous.

We carefully prove that the map  $g$  is continuous. Let  $\mathbb{R}^d$  have the standard topology and  $\text{GL}_d(\mathbb{R}) \times \mathbb{R}^d$  have the subspace topology of  $\mathcal{M}_d(\mathbb{R}) \times \mathbb{R}^d$ . Note that taking the inverse of a matrix is a continuous function in the topology induced by the Frobenius norm (proven in detail in Section A.2). Mapping  $b \mapsto b$  is a continuous function in the standard topology of  $\mathbb{R}^d$ . Since these two functions are defined on spaces with the same topology, recall that the product of two continuous functions is continuous. Hence,  $g(A, b) = A^{-1}b$  is a continuous function.

Observe that  $g \circ f$  is the circumcenter of the sphere through the input points. This composition is continuous since  $g$  and  $f$  are continuous. Let  $p = (\mathbf{p}_1, \mathbf{p}_2, \dots, \mathbf{p}_{d+1}) \in \mathbb{R}^{d(d+1)}$  and let  $O = g \circ f(p)$ . Thus, for any  $\varepsilon_1 > 0$ , there exists a  $\delta > 0$  such that when  $p' = (\mathbf{p}_1', \mathbf{p}_2', \dots, \mathbf{p}_{d+1}')$  satisfies  $\|p - p'\| < \delta$ , then

$$\|O - O'\| < \varepsilon_1$$

where  $O' = g \circ f(p')$ . Thus, we have that  $B_\delta(p) \subset \mathbb{R}^{d(d+1)}$  is an open ball of radius  $\delta$  centered at  $p$ . Let  $\pi_i$  be the projection map that projects  $p$  into its  $i$ th coordinate:

$$\pi_i : \underbrace{\mathbb{R}^d \times \dots \times \mathbb{R}^d}_{d+1} \rightarrow \mathbb{R}^d$$

$$\pi_i : (\mathbf{p}_1, \mathbf{p}_2, \dots, \mathbf{p}_{d+1}) \mapsto \mathbf{p}_i.$$

This map is an open map, thus  $\pi_i(B_\delta(p))$  is an open set in  $\mathbb{R}^d$ . Since we have  $\mathbf{p}_i \in \pi_i(B_\delta(p))$  by definition, we can find an open ball  $B_{\delta_i}(\mathbf{p}_i) \subset \mathbb{R}^d$  for each  $1 \leq i \leq d+1$ . Letting  $\varepsilon = \min\{\delta_i\}_{i=1}^{d+1}$ , we have an  $\varepsilon$  neighborhood around each point  $\mathbf{p}_i$  defining  $\mathcal{S}$  that ensures the circumcenter  $O$  stays in an  $\varepsilon_1$  ball.  $\square$

**Lemma 3.7.** *For a Delaunay simplex  $D = \{p_1, p_2, \dots, p_{d+1}\} \subset P$ , there exists  $\varepsilon > 0$  such that any  $\varepsilon$ -perturbation  $D'$  of  $D$  has  $\{p'_1, p'_2, \dots, p'_{d+1}\} \subset P'$  as a Delaunay simplex.*

*Proof.* Let  $\mathcal{S}$  be the unique  $(d-1)$ -sphere containing  $D$  with center  $O$ . Since  $D$  is a Delaunay simplex, for all  $p_k \in P \setminus D$ ,  $d(O, p_1) < d(O, p_k)$ . Let  $t(p_k) := (d(O, p_k) - d(O, p_1))/4$  and select  $p_{\min} \in P \setminus D$  that minimizes  $t$ . It follows that,

$$d(O, p_1) + 2t = d(O, p_{\min}) - 2t.$$

Since  $P$  is in general position and finite,  $t > 0$ . Let  $\varepsilon_1 = t$ , then we obtain from Lemma 3.6 that there is  $\delta > 0$  such that each  $d(p_i, p'_i) < \delta$  for  $i = 1, 2, \dots, d+1$  implies  $d(O, O') < \varepsilon_1$ . Now, we take  $\varepsilon = \min(\delta, \varepsilon_1)$  and  $D'$  is the  $\varepsilon$ -perturbation of  $D$ . Let  $\mathcal{S}'$  be the unique  $(d-1)$ -sphere containing  $D'$  with center  $O'$ . Using the triangle inequality, we have the following chains of inequalities:

$$\begin{aligned} d(O', p'_1) &\leq d(O', O) + d(O, p_1) + d(p_1, p'_1) \\ &< \varepsilon_1 + d(O, p_1) + \delta \leq t + d(O, p_1) + t \\ &= d(O, p_1) + 2t = d(O, p_{\min}) - 2t. \end{aligned}$$

Since  $d(O, p_{\min})$  minimized  $t$ , it must be that  $d(O, p_{\min}) \leq d(O, p_k)$  for any  $p_k \in P \setminus D$ . Thus, we have

$$\begin{aligned} d(O, p_{\min}) - 2t &\leq d(O, p_k) - t - t < d(O, p_k) - \varepsilon_1 - \delta \\ &\leq d(O, p_k) - d(O, O') - d(p_k, p'_k) \leq d(O', p'_k). \end{aligned}$$

We have concluded that

$$d(O', p'_1) < d(O', p'_k)$$

for any  $p'_k \in P' \setminus D'$ . Hence,  $D'$  is a Delaunay simplex.  $\square$

**Lemma 3.8.** *For a non-Delaunay simplex  $D = \{p_1, p_2, \dots, p_{d+1}\} \subset P$ , there exists  $\varepsilon > 0$  such that any  $\varepsilon$ -perturbation  $D'$  of  $D$  has  $\{p'_1, p'_2, \dots, p'_{d+1}\} \subset P'$  as a non-Delaunay simplex.*

*Proof.* Let  $\mathcal{S}$  be the unique  $(d-1)$ -sphere containing  $D$  with center  $O$ . Since  $D$  is a non-Delaunay simplex, there exists a  $p_k \in P \setminus D$  such that  $d(O, p_k) < d(O, p_1)$ . Let  $t := (d(O, p_1) - d(O, p_k))/4$  so that

$$d(O, p_k) + 2t = d(O, p_1) - 2t.$$

Since  $P$  is in general position,  $t > 0$ . Let  $\varepsilon_1 = t$ , then we obtain from Lemma 3.6 that there is  $\delta > 0$  such that each  $d(p_i, p'_i) < \delta$  for  $i = 1, 2, \dots, d+1$  implies  $d(O, O') < \varepsilon_1$ . Now, we take  $\varepsilon = \min(\delta, \varepsilon_1)$  and  $D'$  is the



$\varepsilon$ -perturbation of  $D$ . Let  $\mathcal{S}'$  be the unique  $(d-1)$ -sphere containing  $D'$  with center  $O'$ . Using the triangle inequality, we have the following chains of inequalities:

$$\begin{aligned}
d(O', p'_k) &\leq d(O', O) + d(O, p_k) + d(p_k, p'_k) \\
&< \varepsilon_1 + d(O, p_k) + \delta \leq t + d(O, p_k) + t \\
&= d(O, p_k) + 2t = d(O, p_1) - 2t \\
&= d(O, p_1) - t - t \leq d(O, p_1) - \varepsilon_1 - \delta \\
&< d(O, p_1) - d(O, O') - d(p_1, p'_1) \leq d(O', p'_1).
\end{aligned}$$

Since  $d(O', p'_1)$  is the radius of  $\mathcal{S}'$ , we have that  $p'_k \in \text{int}(\mathcal{S}')$ . Hence,  $D'$  is a non-Delaunay simplex.  $\square$

**Theorem 3.9.** *Let  $P \subset \mathbb{R}^d$  be in general position, that is no  $d+2$  points lie on the same  $(d-1)$ -sphere and no  $d+1$  points are affinely dependent. Then, there exists  $\varepsilon > 0$  such that any  $\varepsilon$ -perturbation  $P'$  of  $P$  has the same Delaunay Triangulation as  $P$ . That is,*

$$\text{Del}(P) = \text{Del}(P').$$

*Proof.* Let  $D$  be a subset of  $P$  of size  $d+1$ . Depending if  $D$  is a Delaunay simplex in  $P$ , obtain  $\varepsilon_D$  from Lemma 3.7 or from Lemma 3.8. For any  $\varepsilon_D$ -perturbation  $D'$  of  $D$ , we have  $\text{Del}(D) = \text{Del}(D')$ . Now, find

$$\varepsilon = \min\{\varepsilon_D \mid \text{for every } D \subset P\}.$$

Hence, for any  $\varepsilon$ -perturbation  $P'$  of  $P$ ,  $P'$  has the same Delaunay Triangulation as  $P$ .  $\square$

Now, we need a theorem that gives us stability of PDs in PD space equipped with the bottle-neck distance. So,

## 4 Application of Delaunay-Rips

1. Demonstrate value by talking about as dimensions change and number of points change.
2. Particular examples of how using special data sets affect the run-time of Rips/Alpha drastically but maybe not Del-Rips.
3. Performance: accuracy in ML algorithm, or classification. Instability may cause performance to go down even though run-time is unaffected.

### 4.1 Synthetic Data

Focus is two things: Accuracy and Speed. Demonstrating accuracy (by comparing with Rips and Alpha) builds trust in Del-Rips. Demonstrating speed encourages using Del-Rips over Rips or Alpha.

Comparing Ripser, Cechmate's Alpha, and our Del-Rips implementation.

1. Data 1: Noisy 1000 pt 1-sphere
2. Data 2: Noisy 100 pt. 2-sphere
3. Data 3: Noisy 100 pt. Torus in 3 dimensions
4. Data 4: Noisy 100 pt. Klein bottle in 4 dimensions

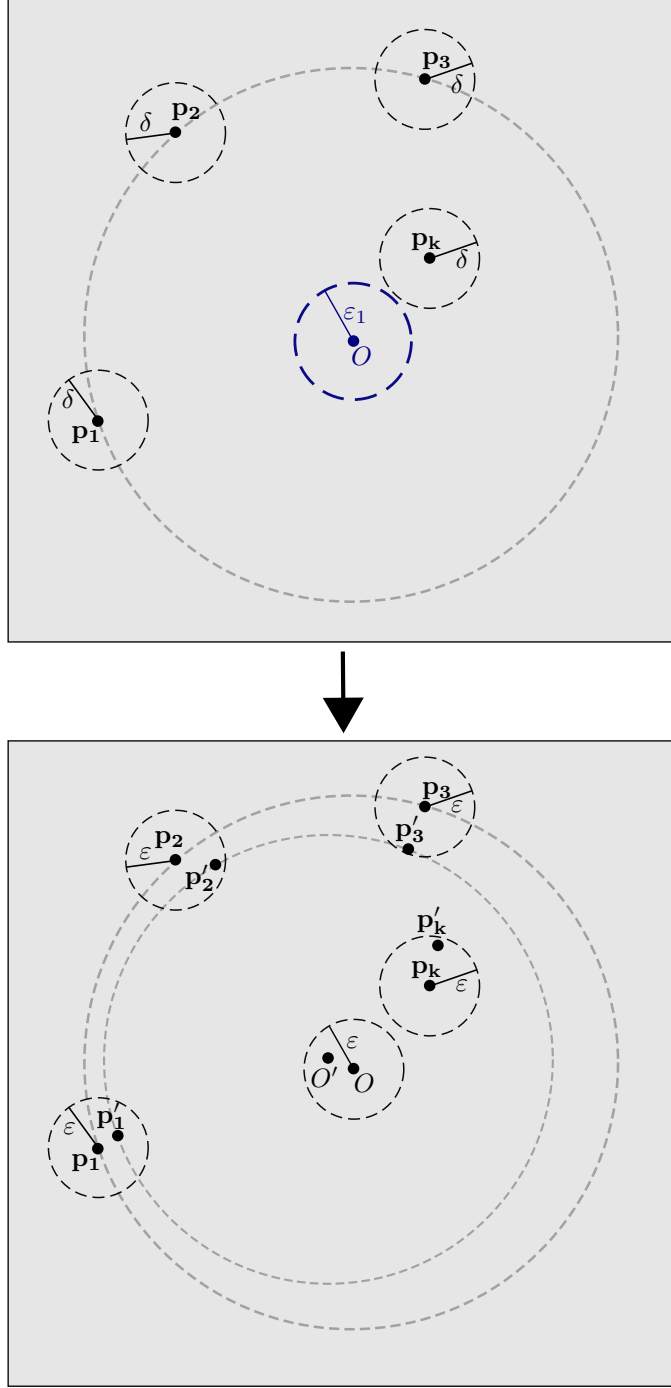


Figure 5: Maintaing a non-Delaunay triangle in  $\mathbb{R}^2$ . We take  $\varepsilon = \min(\delta, \varepsilon_1)$ . We illustrate in this figure how  $p'_k$  remains within the circumcircle through  $p'_1, p'_2$ , and  $p'_3$ .

Section with tables. Make a table with a column labeled “Rips” and the other column labeled “Alpha”. Each row of the table is a data set. In each cell of the table, put the PD for the corresponding column method alongside the PD made by using Del-Rips. Place underneath the bottleneck distance between the two PDs. Do this for each data set (row).

Next table will have 3 columns (Rips, Alpha, Del-Rips). Each row is (again) a data set. Each cell will be the corresponding run-time (in seconds).

**Need a better measurement of accuracy: maybe look at a classification problem (see how the quality of prediction is affected when using del-rips)**

Section with graphs. First graph will be of a 2-sphere with points increasing from 100-10,000 with runtime on the y-axis. Second graph will be a  $d$ -sphere of 50 points with dimension increasing from 2-7 with the runtime on the y-axis. Each graph will compare Rips, Alpha, and Del-Rips. **plot the standard deviation along each curve for each data point: or better, a box-whisker plot at each point for this non-gaussian distribution**

Note that we are comparing runtime of packages, not the run time of the filtration. So the theoretical runtime analysis will not compare with the actual runtime in this section. Or, we could compare rips (cechmate), alpha (cechmate), del-rips because they all use PHAT to reduce the boundary matrix and then the only comparison to make would be the filtration step.

Highlight separately the speed up in the filtration step and then the PD step (since you have less simplices).

## 4.2 Real Data

# 5 Conclusion

## 5.1 Further Questions

## A Appendix: Math

### A.1 Boundary Matrix Calculation for Instability

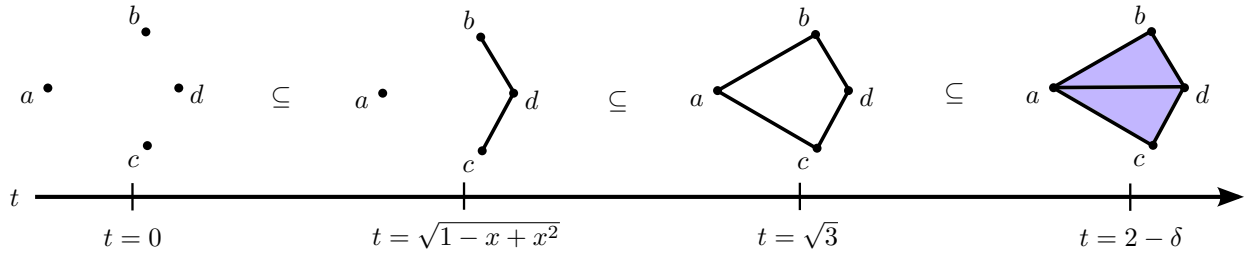


Figure 6: filtration

We have  $P' = \{(-1, 0), (\frac{1}{2}, \frac{\sqrt{3}}{2}), (\frac{1}{2}, -\frac{\sqrt{3}}{2}), (1-x, 0)\}$  with  $0 < x < \delta < 2 - \sqrt{3}$ . Our filtration has 4 key scale values,  $t = 0 \leq \sqrt{1-x+x^2} < \sqrt{3} < 2 - \delta$  as shown in Figure 6. We construct our boundary matrix  $B$  and reduce it to  $\bar{B}$  using the standard algorithm:

$$B = \begin{matrix} & \begin{matrix} a & b & c & d & bd & cd & ab & ac & ad & abd & acd \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \\ bd \\ cd \\ ab \\ ac \\ ad \\ abd \\ acd \end{matrix} & \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix} \end{pmatrix}$$

$$\overline{B} = \begin{matrix} & a & b & c & d & bd & cd & ab & ac & ad & abd & acd \\ \begin{matrix} a \\ b \\ c \\ d \\ bd \\ cd \\ ab \\ ac \\ ad \\ abd \\ acd \end{matrix} & \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix} \end{matrix}.$$

The persistence pairs for the  $H_0$  class with their persistence diagram coordinate (birth/death pair) come out as follows:

$$\begin{aligned} (a, N/A) &: (0, \infty) \\ (b, ab) &: (0, \sqrt{3}) \\ (c, cd) &: (0, \sqrt{1-x+x^2}) \\ (d, bd) &: (0, \sqrt{1-x+x^2}). \end{aligned}$$

The  $H_1$  classes come out as

$$\begin{aligned} (ad, abd) &: (2-x, 2-x) \\ (ac, acd) &: (\sqrt{3}, 2-x). \end{aligned}$$

The only point with non-zero persistence is  $(\sqrt{3}, 2-x)$ .

## A.2 Continuity of Matrix Inverse

Let  $V = (\mathcal{M}_n(\mathbb{R}), \|\cdot\|)$  be the normed vector space of  $n \times n$  matrices over  $\mathbb{R}$  with the Frobenius norm defined as

$$\|A\|_F = \sqrt{\sum_{i=1}^n \sum_{j=1}^n a_{ij}^2}$$

where  $a_{ij}$  is the entry in the  $i$ th row and  $j$ th column of  $A$ . Recall the definition of the vector 2-norm

$$\|v\|_2 = \sqrt{\sum_{i=1}^n |v_i|^2}$$

for  $v \in \mathbb{R}^n$ . We use the vector 2-norm with the Frobenius matrix norm because of their compatibility: that is,

$$\|Mv\|_2 \leq \|M\|_F \|v\|_2.$$

We proceed by dropping the subscripts and using  $\|\cdot\|$  for matrix and vector norms.

The heavy machinery used to establish the continuity of the matrix inverse was done in [Wil94] pages 92-93. However, their proof did not explicitly compute inequalities involving the Frobenius norm (that is,  $\|I\| = \sqrt{n}$  for  $I \in V$  being the identity). For completeness, we provide derivations of similar results using the Frobenius norm.

**Lemma A.1.** *Given  $A, E \in V$  with  $A$  invertible, the matrix  $A + E$  is invertible if*

$$\|A^{-1}E\| < 1.$$

*Proof.* We mimic the proof found in [Wil94] page 92. We write

$$A + E = A(I + A^{-1}E) \quad (4)$$

to see that  $A + E$  is invertible if  $I + A^{-1}E$  is invertible. From [Wil94] page 82, if  $\lambda$  is an eigenvalue of  $M$ , then

$$||M|| \geq |\lambda|.$$

Thus, if  $\lambda_i$  is an eigenvalue of  $A^{-1}E$ , then we have

$$||A^{-1}E|| \geq |\lambda_i|.$$

Since the eigenvalues of  $I + A^{-1}E$  are of the form  $1 + \lambda_i$ , to ensure  $I + A^{-1}E$  is invertible, we need  $\lambda_i \neq -1$  for any  $i$ . One way to ensure this is to require  $||A^{-1}E|| < 1$  so that

$$|\lambda_i| \leq ||A^{-1}E|| < 1.$$

This condition on  $||A^{-1}E||$  guarantees the invertibility of  $A + E$ .  $\square$

**Lemma A.2.** *Given  $A, E \in V$  with  $A$  invertible, let  $x, b \in \mathbb{R}^n$  such that  $Ax = b$  and  $||A^{-1}|| ||E|| < 1$ . Then, there exists  $h \in \mathbb{R}^n$  such that  $(A + E)(x + h) = b$  and*

$$||(A + E)^{-1} - A^{-1}|| \leq \frac{\sqrt{n}||A^{-1}||^2 ||E||}{1 - ||A^{-1}|| ||E||}.$$

*Proof.* We mimic the proof found in [Wil94] pages 92-93. Begin by defining

$$F := A^{-1}E$$

$$G := (I + F)^{-1}.$$

Note that since  $||A^{-1}E|| \leq ||A^{-1}|| ||E|| < 1$ , we have that  $(I + F)$  is invertible by the proof of Lemma A.1. Hence,  $G$  indeed is well-defined. Multiplying each side by  $(I + F)$  results in

$$I = G + FG$$

and using the reverse triangle inequality yields

$$||I|| = ||G + FG|| = ||G - (-FG)|| \geq ||G|| - ||-FG|| = ||G|| - ||FG|| \geq ||G|| - ||F|| ||G||$$

$$\sqrt{n} \geq ||G|| - ||F|| ||G||$$

since  $||I|| = \sqrt{n}$ . Since  $||F|| = ||A^{-1}E|| < 1$ , we have

$$||G|| \leq \frac{\sqrt{n}}{1 - ||F||}.$$

We will use this inequality in a bit. Now, we rewrite  $(A + E)(x + h) = b$  as

$$(A + E)h = -Ex.$$

Applying Lemma A.1, we have  $A + E$  is invertible and using equation 4, we have the existence of  $h$ :

$$h = -(A + E)^{-1}Ex = -(I + F)^{-1}A^{-1}Ex = -GA^{-1}Ex.$$

Further,

$$||h|| \leq ||G|| ||A^{-1}E|| ||x|| \leq \frac{\sqrt{n}||A^{-1}E|| ||x||}{1 - ||F||} = \frac{\sqrt{n}||A^{-1}E|| ||x||}{1 - ||A^{-1}E||} \leq \frac{\sqrt{n}||A^{-1}|| ||E|| ||x||}{1 - ||A^{-1}|| ||E||}$$

using the hypothesis that  $\|A^{-1}\| \|E\| < 1$ . Division by  $\|x\|$  results in

$$\frac{\|h\|}{\|x\|} \leq \frac{\sqrt{n}\|A^{-1}\| \|E\|}{1 - \|A^{-1}\| \|E\|}.$$

Using the substitutions

$$x = A^{-1}b,$$

$$h = (A + E)^{-1}b - x = (A + E)^{-1}b - A^{-1}b = ((A + E)^{-1} - A^{-1})b$$

we have

$$\begin{aligned} \frac{\|((A + E)^{-1} - A^{-1})b\|}{\|A^{-1}\| \|b\|} &\leq \frac{\|h\|}{\|x\|} \leq \frac{\sqrt{n}\|A^{-1}\| \|E\|}{1 - \|A^{-1}\| \|E\|} \\ \frac{\|((A + E)^{-1} - A^{-1})b\|}{\|b\|} &\leq \frac{\sqrt{n}\|A^{-1}\|^2 \|E\|}{1 - \|A^{-1}\| \|E\|}. \end{aligned}$$

Since this inequality is true for any  $b$ , we have that

$$\sup \left\{ \frac{\|((A + E)^{-1} - A^{-1})b\|}{\|b\|} : b \neq 0 \right\} \leq \frac{\sqrt{n}\|A^{-1}\|^2 \|E\|}{1 - \|A^{-1}\| \|E\|}.$$

It follows by the definition of the norm that

$$\|(A + E)^{-1} - A^{-1}\| \leq \frac{\sqrt{n}\|A^{-1}\|^2 \|E\|}{1 - \|A^{-1}\| \|E\|}.$$

□

**Theorem A.3.** *Define the inverse of a matrix as the map*

$$\iota : V \rightarrow V$$

$$\iota(A) = A^{-1}.$$

*This map is continuous in the topology on  $V$  induced by the Frobenius norm.*

*Proof.* Let  $\varepsilon > 0$  and  $A \in V$  be arbitrary and invertible and choose  $\delta = \min(\varepsilon/(2\sqrt{n}\|A^{-1}\|^2), 1/(2\|A^{-1}\|))$ . We find  $E \in V$  so that

$$\|(A + E) - A\| < \delta.$$

Hence, we have

$$\|(A + E) - A\| = \|E\| < \delta \leq 1/(2\|A^{-1}\|)$$

so that

$$\|A^{-1}\| \|E\| < 1/2 < 1.$$

Thus, Lemma A.2 applies and we have

$$\begin{aligned} \|\iota(A + E) - \iota(A)\| &= \|(A + E)^{-1} - A^{-1}\| \leq \frac{\sqrt{n}\|E\| \|A^{-1}\|^2}{1 - \|E\| \|A^{-1}\|} \leq \frac{\sqrt{n}\|A^{-1}\|^2}{1 - 1/2} \|E\| \\ &= 2\sqrt{n}\|A^{-1}\|^2 \|E\| < 2\sqrt{n}\|A^{-1}\|^2 \delta \leq \varepsilon. \end{aligned}$$

Since  $A$  was arbitrary, we have that  $\iota$  (the matrix inverse) is a continuous function in the topology induced by the Frobenius norm. □

## B Appendix: Tech

### B.1 Pseudo-code Implementation

### B.2 Github Repo of Actual, Clean Code

1. We want to compare the best implementation of Del-Rips with Ripser and Cechmate's Alpha.

---

**Algorithm 1:** An algorithm to compute Delaunay-Rips Filtration

---

**Input:**  $P = (\mathbf{p}_1, \dots, \mathbf{p}_n)$ ,  $dim$  = maximum homology dimension to compute  
**Output:** Delaunay Rips Filtration  
triangulation  $\leftarrow$  Delaunay( $P$ )  
filtration  $\leftarrow [([i], 0)]$  for  $1 \leq i \leq n$  /\* Add the 0-simplices into the filtration \*/  
**for** each simplex  $\in$  triangulation **do**  
    **for**  $d \leftarrow 1$  **to**  $dim + 1$  **do**  
        faces  $\leftarrow$  all  $d$ -simplex subsets of current simplex  
        **for** each face  $\in$  faces **do**  
            **if** face  $\notin$  filtration and  $d == 1$  **then**  
                value  $\leftarrow$  distance(face[0], face[1]) /\* Calculate the euclidean distance between the two points \*/  
                filtration.append((face, value)) /\* Append 1-simplices \*/  
            **end**  
            **if** face  $\notin$  filtration and  $dim > 1$  **then**  
                find subface of face with greatest value  
                filtration.append((subface, value)) /\* Add higher order simplices \*/  
            **end**  
        **end**  
    **end**  
**end**

---

### B.3 Machine Specs

1. Eluktronics laptop

## Whole bibliography

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