# **ArosTemplate**

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# **InSkill**

#### Ubuntu 下 CodeBlocks 更改调试终端

在环境设置里进行如下设置:把 Terminal to launch console programs 那个选项改成 "gnome-terminal -t \$TITLE -x",原来是"xterm -T \$TITLE -e"。

#### HDU 上的 DFS 爆栈问题的简易解决方法

在文件 gui 头处加上这么一句"#pragma comment(linker, "/STACK:1024000000,1024000000")" 后面两个数字随便写,你觉得能过就好,另外不要超了栈内存的真正上限。基于 VC++的编译预处理命令,这个代码必须拿 C++来提交,所以 C++会出现的那种 long long 和\_\_int64 的问题也要注意到。

### 通过内嵌汇编把堆空间作为栈空间使用\_hdu\_4118

```
#include<cstdio>
#include<cstring>
#include<algorithm>
using namespace std;
const int MAXN = 100000+5, MAXM = 200000+5;
int T, N, X, Y, Z;
int e, head[MAXN], next[MAXM], v[MAXM];
int cnt[MAXN];
long long w[MAXM], ans;
void addedge(int x, int y, int z)
   v[e] = y; w[e] = z;
   next[e] = head[x]; head[x] = e++;
void dfs(int u, int fa = 0)
   cnt[u] = 1;
   for (int i = head[u]; i != -1; i = next[i]) if (v[i] != fa)
       dfs(v[i], u);
       ans += min(cnt[v[i]], N-cnt[v[i]])*2*w[i];
       cnt[u] += cnt[v[i]];
```

```
void call dfs()
    const int STACK SIZE = 1<<23;
    static char stack[STACK_SIZE];
    int bak;
    _asm__volatile_
        "movl %%esp, %0\n"
        "movl %1, %%esp\n":
        "=q"(bak):
        "g"(stack+STACK_SIZE-1):
    );
    dfs(1);
    __asm___volatile__
        "mov1 %0, %%esp\n":
        "q"(bak):
    );
}
int main()
    scanf("%d", &T);
    for (int cas = 1; cas <= T; cas++)
        e = 0;
       memset(head, -1, sizeof(head));
        scanf("%d", &N);
       for (int i = 1; i < N; i++)
           scanf("%d%d%d", &X, &Y, &Z);
           addedge(X, Y, Z);
           addedge(Y, X, Z);
       ans = 0:
        call_dfs();
       printf("Case #%d: %I64d\n", cas, ans);
    return 0;
```

# Graph

spfa

```
#include<cstdio>
#include<cstring>
#include<queue>
using namespace std;
const int MAXN = 1000+5, MAXM = 1000+5;
const int INF = 0x3f3f3f3f;
int n, m, e, s;
int v[MAXM], next[MAXM], head[MAXN];
int w[MAXM], d[MAXN];
int ing cnt[MAXN]; //存在负权回路时需要
bool ing[MAXN];
queue<int> Q;
void addedge(int x, int y, int z)
   v[e] = y; w[e] = z;
   next[e] = head[x]; head[x] = e;
   e++;
bool spfa()
   for (int i = 1; i <= n; i++)
      d[i] = (i == s ? 0 : INF);
   memset(inq, 0, sizeof(inq));
   memset(ing cnt, 0, sizeof(ing cnt));
   while (!Q.empty()) Q.pop();
   Q.push(s);
   inq[s] = 1;
   ing cnt[s]++;
   while (!Q.empty())
      int u = Q.front(); Q.pop();
      inq[u] = 0;
      for(int e = head[u]; e != -1; e = next[e])
          if(d[v[e]] > d[u]+w[e])
             d[v[e]] = d[u]+w[e];
             if(!inq[v[e]])
                 Q.push(v[e]);
```

```
inq[v[e]] = 1;
    inq_cnt[v[e]]++;
    if (inq_cnt[v[e]] > n)
        return 0;
    }
}
return 1;

int main()

{
// freopen("input.txt", "r", stdin);
// freopen("output.txt", "w", stdout);
    memset(head, -1, sizeof(head));
    e = 0;
return 0;
}
```

#### 二维最短路 hdu 4396

```
题意: 求至少经过 K 条边, 到达终点的最短路 (K<=50)。
思路:因为 K<=500,所以每个节点最多扩展成 50 个节点,最后一个节点表示到达该节点时
经过的边数(收集到的木材/10)已经满足 K 值对应的要求。然后 spfa,每个节点表示为(编
号,经过的边数)。
*/
#include<cstdio>
#include<cstring>
#include<algorithm>
#include<queue>
using namespace std;
const int MAXN = 5000+5, MAXM = 200000+5, MAXK = 50+5;
const int INF = 0x3f3f3f3f3f;
int N, M, A, B, C, S, T, K, mk;
int e, head[MAXN], next[MAXM], v[MAXM];
int d[MAXN][MAXK], w[MAXM];
bool inq[MAXN][MAXK];
queue<pair<int, int> > Q;
void addedge(int x, int y, int z)
   v[e] = y; w[e] = z;
   next[e] = head[x]; head[x] = e++;
```

```
void spfa(int s)
   for (int i = 1; i <= N; i++)
       for (int j = 1; j <= mk; j++)
           d[i][j] = INF;
   Q.push(make_pair(s, 0));
   while (!Q.empty())
       int u = Q.front().first, k = Q.front().second;
       Q.pop();
       inq[u][k] = 0;
       for (int i = head[u]; i != -1; i = next[i])
           int 1 = k + (k < mk ? 1 : 0);
           if (d[u][k]+w[i] < d[v[i]][l])
               d[v[i]][1] = d[u][k]+w[i];
               if (!ing[v[i]][1])
                   Q.push(make pair(v[i], 1));
                   inq[v[i]][1] = 1;
           }
void init()
   e = 0;
   memset(head, -1, sizeof(head));
int main()
   while (scanf("%d%d", &N, &M) != EOF)
       init();
       for (int i = 0; i < M; i++)
           scanf("%d%d%d", &A, &B, &C);
           addedge(A, B, C);
           addedge(B, A, C);
       scanf("%d%d%d", &S, &T, &K);
       mk = (K-1)/10+1;
       spfa(S);
```

```
printf("%d\n", d[T][mk] < INF ? d[T][mk] : -1);
}
return 0;
}</pre>
```

#### 找环\_hdu\_4337

```
#include<cstdio>
#include<cstring>
#include<algorithm>
#include<vector>
using namespace std;
const int MAXN = 150+5, MAXM = 22500+5, MAXP = 50+5;
int N, M, a, b;
int e, head[MAXN], next[MAXM], v[MAXM];
int mark[MAXN];
vector<int> vec;
void Init()
   e = 0;
   memset(head, -1, sizeof(head));
   memset(mark, -1, sizeof(mark));
   vec.clear();
void addedge(int x, int y)
   v[e] = y;
   next[e] = head[x]; head[x] = e++;
bool dfs(int u, int step = 0)
   mark[u] = step;
   vec.push back(u);
   for (int i = head[u]; i != -1; i = next[i])
      if (mark[v[i]] == -1)
          if (dfs(v[i], step+1))
             return 1;
       else if (step-mark[v[i]]+1 == N)
          return 1:
   mark[u] = -1;
```

```
vec.pop back();
   return 0;
int main()
   while (scanf("%d%d", &N, &M) != EOF)
      Init();
      for (int i = 0; i < M; i++)
          scanf("%d%d", &a, &b);
          addedge(a, b);
          addedge(b, a);
      if (dfs(1))
          for (int i = 0; i < (int)vec.size(); i++)
             if (i)
                 printf(" ");
             printf("%d", vec[i]);
          printf("\n");
      else
          printf("no solution\n");
   return 0;
```

#### 最小生成树的最佳替换边 hdu 4126

/\*

题意:给定一个图 G,有 q 次询问(相互独立),每次询问(u,v,w),表示将<u,v>这条边的边权更改为 w,求此时的最小生成树的值。

算法: 先求得最小生成树, 对于每次询问(u,v,w), 分两种情况讨论:

- 1、若<u,v>是非最小生成树上的边,那么不用考虑,最小生成树仍是原来的值。
- 2、若<u,v>是最小生成树上的边,那么我们就需要在这条边所导致的两个集合中分别选出一个点 i,j 并且 g[i][j]最小来替代那条被增加的边,那么反过来考虑,对于一条非最小生成树上的边<u,v>,它可以替代哪些边呢?就是 u->x1->x2->...->xk->v 这条路径(因为是树,所以这条路径唯一)上的边。那么现在要求的就是对于每条树上的边<u,v>,得到一个best[u][v]表示去掉它,最小的替代边的权值。这个可以用 dfs 来做,以每个点为起点做dfs,遍历整个最小生成树,得到 best[i][j],这样复杂度就是  $O(N^2)$ 的,然后对于每次询问就可以 O(1) 回答了。这个 dfs 的写法还是有点技巧的,具体就见代码吧。

```
#include<cstdio>
#include<cstring>
#include<algorithm>
using namespace std;
const int MAXN = 3000+5, MAXM = 6000+5;
const int INF = 0x3f3f3f3f;
int N, M, Q, X, Y, C;
int g[MAXN][MAXN], best[MAXN][MAXN], dis[MAXN], pre[MAXN];
int e, head[MAXN], next[MAXM], v[MAXM];
bool vis[MAXN];
void addedge(int x, int y)
   v[e] = y;
   next[e] = head[x]; head[x] = e++;
void init()
    e = 0;
    memset(head, -1, sizeof(head));
   for (int i = 0; i < N; i++)
       for (int j = 0; j < i; j++)
            g[i][j] = g[j][i] = best[i][j] = best[j][i] = INF;
int prim()
    for (int i = 0; i < N; i++)
       vis[i] = 0, pre[i] = -1, dis[i] = INF;
    int res = 0;
    dis[0] = 0;
    for (int j = 0; j < N; j++)
       int u = -1;
       for (int i = 0; i < N; i++)
            if (!vis[i] && (u == -1 || dis[i] < dis[u]))
                u = i;
       vis[u] = 1;
       res += dis[u];
        if (pre[u] != -1)
            addedge(u, pre[u]);
            addedge(pre[u], u);
        for (int i = 0; i < N; i++)
            if (!vis[i] && g[u][i] < dis[i])</pre>
```

```
dis[i] = g[u][i];
               pre[i] = u;
   return res;
int dfs(int st, int u, int fa)
   int mini = INF;
   for (int i = head[u]; i != -1; i = next[i]) if (v[i] != fa)
       int cur = dfs(st, v[i], u);
       mini = min(mini, cur);
       best[u][v[i]] = best[v[i]][u] = min(best[u][v[i]], cur);
   if (st != fa)
       mini = min(mini, g[st][u]);
   return mini;
int main()
   while (scanf("%d%d", &N, &M))
       if (!N && !M)
           break:
       init();
       for (int i = 0; i < M; i++)
           scanf("%d%d%d", &X, &Y, &C);
           g[X][Y] = g[Y][X] = C;
       int mst = prim();
       for (int i = 0; i < N; i++)
           dfs(i, i, -1);
       scanf("%d", &Q);
       double ans = 0;
       for (int i = 0; i < Q; i++)
           scanf("%d%d%d", &X, &Y, &C);
           if (pre[X] == Y || pre[Y] == X)
               ans += mst-g[X][Y]+min(C, best[X][Y]);
               ans += mst;
```

```
ans /= Q;
    printf("%.4f\n", ans);
}
return 0;
}
```

#### 最小树形图 hdu 4009

/\*

题意:有 n 个地方需要供水,每个地方都可以选择是自己挖井,还是从别的地方引水,根据方法不同和每个地方的坐标不同,花费也不同,现在给出每个地方的坐标,花费的计算方法,以及每个地方可以给哪些地方供水(即对方可以从这里引水),求给所有地方供水的最小花费。 思路:显然对于每个地方,只有一种供水方式就足够了,这样也能保证花费最小,而每个地方都可以自己挖井,所以是不可能出现无解的情况的,为了方便思考,我们引入一个虚拟点,把所有自己挖井的都连到这个点,边权为挖井的花费,而如果 i 能从 j 处引水,则从 j 向 i 连边,边权为引水的花费,然后对这个有向图,以虚拟点为根,求最小树形图即可(最小树形图即为有向图的最小生成树)。

```
*/
#include<cstdio>
#include<cstring>
#include<cmath>
#include<algorithm>
using namespace std;
const int MAXN = 1000+5, MAXM = 1001000+5;
const int INF = 0x3f3f3f3f;
int N, X, Y, Z, K, x[MAXN], y[MAXN], z[MAXN];
int e, u[MAXM], v[MAXM], w[MAXM];
int pre[MAXN], id[MAXN], vis[MAXN];
int in[MAXN]:
int Directed MST(int root, int NV, int NE) //number vertices from zero!!!
   int res = 0;
   for (;;)
       //1.找最小入边
       for (int i = 0; i < NV; i++)
           in[i] = INF, id[i] = -1, vis[i] = -1;
       for (int i = 0; i < NE; i++)
           int s = u[i], t = v[i];
           if (w[i] < in[t] && s != t)
               pre[t] = s;
               in[t] = w[i];
```

```
for (int i = 0; i < NV; i++)
           if (i == root)
              continue;
           if (in[i] == INF)
               return -1;//除了跟以外有点没有入边,则根无法到达它
       //2.找环
       int cntnode = 0;
       in[root] = 0;
       for (int i = 0; i < NV; i++)
       {//标记每个环
          res += in[i];
          int t = i;
          for (; vis[t] != i && id[t] == -1 && t != root; t = pre[t])
               vis[t] = i;
           if (t != root && id[t] == -1)
               for (int s = pre[t]; s != t; s = pre[s])
                  id[s] = cntnode;
              id[t] = cntnode++;
       if (!cntnode)
           break://无环
       for (int i = 0; i < NV; i++)
           if (id[i] == -1)
              id[i] = cntnode++;
       //3.缩点,重新标记
       for (int i = 0; i < NE; i++)
          int t = v[i];
           u[i] = id[u[i]];
          v[i] = id[v[i]];
           if (u[i] != v[i])
               w[i] -= in[t];
       NV = cntnode;
       root = id[root];
   return res;
void addedge(int x, int y, int z)
```

```
u[e] = x; v[e] = y; w[e] = z;
    e++;
int main()
   while (scanf("%d%d%d%d", &N, &X, &Y, &Z))
       if (!N && !X && !Y && !Z)
           break:
       e = 0;
       int root = 0;
       for (int i = 1; i <= N; i++)
           scanf("%d%d%d", &x[i], &y[i], &z[i]);
           addedge(root, i, z[i]*X);
       for (int i = 1; i <= N; i++)
           scanf("%d", &K);
           for (int k = 1, j; k \le K; k++)
               scanf("%d", &j);
               if(z[i] < z[j])
                   addedge(i, j,
(abs(x[i]-x[j])+abs(y[i]-y[j])+abs(z[i]-z[j]))*Y+Z);
                   addedge(i, j,
(abs(x[i]-x[j])+abs(y[i]-y[j])+abs(z[i]-z[j]))*Y);
       printf("%d\n", Directed_MST(root, N+1, e));
   return 0;
```

# **Network**

最大流 ISAP\_hdu\_3879

```
#include<cstdio>
#include<cstring>
```

```
#include<algorithm>
using namespace std;
const int MAXN = 55000+5, MAXM = 155000*2+5;
const int INF = 0x3f3f3f3f;
int N, M;
int n, s, t;
int e, v[MAXM], next[MAXM], head[MAXN];
int cap[MAXM];
int h[MAXN], gap[MAXN];
void init()
   e = 0;
  memset(head, -1, sizeof(head));
  memset(gap, 0, sizeof(gap));
   memset(h, 0, sizeof(h));
void addedge(int x, int y, int c)
   v[e] = y; cap[e] = c;
   next[e] = head[x]; head[x] = e++;
   v[e] = x; cap[e] = 0;
   next[e] = head[y]; head[y] = e++;
int sap(int u, int f)
   if (u == t)
     return f;
   int minh = n-1, rf = f;
   for (int i = head[u]; i != -1; i = next[i]) if (cap[i])
      if (h[v[i]]+1 == h[u])
          int cf = sap(v[i], min(cap[i], rf));
          cap[i] -= cf;
          cap[i^1] += cf;
          rf -= cf;
          if (h[s] >= n)
             return f-rf;
          if (!rf)
             break;
      minh = min(minh, h[v[i]]);
   }
   if (rf == f)
```

```
gap[h[u]]--;
      if (!gap[h[u]])
          h[s] = n;
      h[u] = minh+1;
      gap[h[u]]++;
   return f-rf;
int maxflow()
   int res = 0;
   gap[0] = n;
   while (h[s] < n)
     res += sap(s, INF);
   return res;
int main()
   freopen("input.txt", "r", stdin);
// freopen("output.txt", "w", stdout);
   while (scanf("%d%d", &N, &M) != EOF)
      init();
      n = N+M+2; s = N+M+1; t = s+1;
      for (int i = 1; i <= N; i++)
         int P;
          scanf("%d", &P);
          addedge(i, t, P);
      int tp = 0;
      for (int i = 1; i <= M; i++)
          int x, y, z;
          scanf("%d%d%d", &x, &y, &z);
          tp += z;
          addedge(s, N+i, z);
          addedge(N+i, x, INF);
          addedge(N+i, y, INF);
      int f = maxflow();
      printf("%d\n", tp-f);
   return 0;
```

#### 最大流-邻接表

```
#include<cstdio>
#include<cstring>
#include<algorithm>
#include<queue>
using namespace std;
const int MAXN = 1000+5, MAXM = 1000+5;
const int INF = 0x3f3f3f3f;
int e, s, t, n;
int v[MAXM], next[MAXM], head[MAXN];
int cap[MAXM], a[MAXN], f;
int pv[MAXN], pe[MAXN];
queue<int> Q;
void addedge(int u_, int v_, int c_)
  v[e] = v_{;} cap[e] = c_{;}
  next[e] = head[u_]; head[u_] = e;
   e++;
  v[e] = u_{;} cap[e] = 0;
   next[e] = head[v ]; head[v ] = e;
   e++;
void maxflow()
   f = 0;
   for (;;)
      memset(a, 0, sizeof(a));
      a[s] = INF;
      Q.push(s);
      while (!Q.empty())
          int u = Q.front(); Q.pop();
          for (int e = head[u]; e != -1; e = next[e])
             if(!a[v[e]] && cap[e])
                 Q.push(v[e]);
                 a[v[e]] = min(a[u], cap[e]);
                 pv[v[e]] = u; pe[v[e]] = e;
      if (!a[t]) break;
      for (int v = t; v != s; v = pv[v])
```

#### 最大流-邻接矩阵

```
#include<cstdio>
#include<cstring>
#include<algorithm>
#include<queue>
using namespace std;
const int MAXN = 1000+5;
const int INF = 0x3f3f3f3f;
int s, t, n;
int p[MAXN];
int cap[MAXN][MAXN], flow[MAXN][MAXN], a[MAXN], f;
queue<int> 0;
void addedge(int u , int v , int c )
   cap[u_][v_] = c_;
void maxflow()
   f = 0;
   memset(flow, 0, sizeof(flow));
   for(;;)
      memset(a, 0, sizeof(a));
      a[s] = INF;
      Q.push(s);
      while(!Q.empty())
```

```
int u = Q.front(); Q.pop();
         for(int v = 1; v \le n; v++)
             if(!a[v] \&\& cap[u][v] > flow[u][v])
                p[v] = u; Q.push(v);
                a[v] = min(a[u], cap[u][v]-flow[u][v]);
      if(a[t] == 0) break;
      for(int v = t; v != s; v = p[v])
         flow[p[v]][v] += a[t];
         flow[v][p[v]] -= a[t];
      f += a[t];
int main()
// freopen("input.txt", "r", stdin);
// freopen("output.txt", "w", stdout);
   memset(cap, 0, sizeof(cap));
   return 0;
```

## 最小费用最大流-邻接表

```
#include<cstdio>
#include<cstring>
#include<algorithm>
#include<queue>
using namespace std;
const int MAXN = 1000+5, MAXM = 1000+5;
const int INF = 0x3f3f3f3f;
int e, s, t, n;
int v[MAXM], next[MAXM], head[MAXN];
int cap[MAXM], f;
int cost[MAXM], d[MAXN], c;
int pv[MAXN], pe[MAXN], c;
int pv[MAXN], pe[MAXN];
bool inq[MAXN];
queue<int> Q;
void addedge(int u_, int v_, int c_, int w_)
```

```
v[e] = v_; cap[e] = c_; cost[e] = w_;
   next[e] = head[u_]; head[u_] = e;
   e++;
   v[e] = u_{;} cap[e] = 0; cost[e] = -w_{;}
   next[e] = head[v_]; head[v_] = e;
   e++;
void mincostflow()
   f = 0; c = 0;
   for (;;)
      memset(ing, 0, sizeof(ing));
      for (int i = 1; i <= n; i++)
          d[i] = (i == s ? 0 : INF);
      Q.push(s); inq[s] = 1;
      while (!Q.empty())
          int u = Q.front(); Q.pop();
          ing[u] = 0;
          for (int e = head[u]; e != -1; e = next[e])
             if(cap[e] && d[v[e]] > d[u]+cost[e])
                 d[v[e]] = d[u] + cost[e];
                if (!ing[v[e]])
                    Q.push(v[e]), inq[v[e]] = 1;
                pv[v[e]] = u; pe[v[e]] = e;
      if (d[t] == INF) break;
      int a = INF;
      for (int v = t; v != s; v = pv[v])
          a = min(a, cap[pe[v]]);
      for (int v = t; v != s; v = pv[v])
          cap[pe[v]] -= a;
          cap[pe[v]^1] += a;
      f += a;
      c += d[t]*a;
int main()
```

```
// freopen("input.txt", "r", stdin);
// freopen("output.txt", "w", stdout);
  memset(cap, 0, sizeof(cap));
  memset(cost, 0, sizeof(cost));
  memset(head, -1, sizeof(head));
  e = 0;
  return 0;
}
```

#### 最小费用最大流-邻接矩阵

```
#include<cstdio>
#include<cstring>
#include<algorithm>
#include<queue>
using namespace std;
const int MAXN = 1000+5;
const int INF = 0x3f3f3f3f;
int s, t, n;
int cost[MAXN][MAXN], d[MAXN], c;
int cap[MAXN][MAXN], flow[MAXN][MAXN], f;
int p[MAXN];
bool ing[MAXN];
queue<int> Q;
void addedge(int u , int v , int c , int w )
   cap[u_][v_] = c_;
   cost[u_][v_] = w_; cost[v_][u_] = -w_;
void mincostflow()
   f = 0, c = 0;
   memset(flow, 0, sizeof(flow));
   for(;;)
       for(int i = 1; i <= n; i++)
           d[i] = (i == s ? 0 : INF);
      memset(inq, 0, sizeof(inq));
      Q.push(s); inq[s] = 1;
      while(!Q.empty())
          int u = Q.front(); Q.pop();
          inq[u] = 0;
```

```
for(int v = 1; v \le n; v++)
             if(cap[u][v] > flow[u][v] && d[v] > d[u]+cost[u][v])
                d[v] = d[u] + cost[u][v];
                if(!ing[v])
                    Q.push(v), inq[v] = 1;
                p[v] = u;
      if (d[t] == INF) break;
      int a = INF;
      for(int v = t; v != s; v = p[v])
          a = min(a, cap[p[v]][v]-flow[p[v]][v]);
      for(int v = t; v != s; v = p[v])
          flow[p[v]][v] += a;
          flow[v][p[v]] -= a;
      c += d[t]*a;
      f += a;
int main()
// freopen("input.txt", "r", stdin);
// freopen("output.txt", "w", stdout);
   memset(cap, 0, sizeof(cap));
   memset(cost, 0, sizeof(cost));
   return 0;
```

# Number

## 组合数 C(N, R)

```
int com(int n, int r)
{// return C(n, r)
  if (n-r > r) r = n-r; // C(n, r) = C(n, n-r)
  int s = 1;
  for (int i = 0, j = 1; i < r; i++)
  {</pre>
```

# **Structure**

AC 自动机

AC 自动机\_hdu\_2222

```
网络流上流传最广的 AC 自动机模板题,问你目标串中出现了几个模式串
如果一个结点是单词末尾的话 out 标记为 true,在 search 的时候对于每个结点都向 fail
指针找,找到 out 为 true 的就将其标记为 false,且 ans+=out
*/
#include<cstdio>
#include<cstring>
#include<algorithm>
#include<queue>
using namespace std;
const int MAXN = 1000000+5, MAXM = 50+5;
const int MAX_NODE = 500000+5, MAX_CHD = 26;
int T, N;
int chd[MAX NODE][MAX CHD], fail[MAX NODE], out[MAX NODE];
int ID[1<<8], nv;
char key[MAXM], des[MAXN];
queue<int> Q;
namespace AC_Automaton
   void Initialize()
       fail[0] = 0;
       for (int i = 0; i < MAX_CHD; i++)
          ID[i+'a'] = i;
   void Reset()
```

```
memset(chd[0], 0, sizeof(chd[0]));
       nv = 1:
    void Insert(char *pat)
       int u = 0;
       for (int i = 0; pat[i]; i++)
           int c = ID[pat[i]];
           if (!chd[u][c])
               memset(chd[nv], 0, sizeof(chd[nv]));
               out[nv] = 0;
               chd[u][c] = nv++;
           u = chd[u][c];
        out[u]++;
   void Construct()
       for (int i = 0; i < MAX_CHD; i++)
           if (chd[0][i])
               fail[chd[0][i]] = 0;
                Q.push(chd[0][i]);
       while (!Q.empty())
           int u = Q.front(); Q.pop();
           for (int i = 0; i < MAX_CHD; i++)
               int v = chd[u][i];
               if (v)
                    Q.push(v);
                   fail[v] = chd[fail[u]][i];
               else
                    chd[u][i] = chd[fail[u]][i];
int main()
```

```
AC_Automaton::Initialize();
scanf("%d", &T);
while (T--)
    scanf("%d", &N);
   AC_Automaton::Reset();
    for (int i = 0; i < N; i++)
       scanf("%s", key);
       AC_Automaton::Insert(key);
    AC_Automaton::Construct();
    scanf("%s", des);
   int ans = 0;
    for (int i = 0, u = 0; des[i]; i++)
       u = chd[u][ID[des[i]]];
       for (int t = u; t; )
           ans += out[t];
           out[t] = 0;
           t = fail[t];
    printf("%d\n", ans);
return 0;
```

#### AC 自动机+DP hdu 2825

```
/*
求长度为 n 的字符串中包含至少 k 个给出的关键字的字符串的个数,结果模 MOD。
*/
#include<cstdio>
#include<algorithm>
#include<algorithm>
#include<queue>
using namespace std;

//MAX_NODE = StringNumber*StringLength
const int MAX_NODE = 100+5;
//字符集大小,一般字符形式的题 26 个
```

```
const int MAX CHD = 26;
//每个节点的儿子,即当前节点的状态转移
int chd[MAX NODE][MAX CHD];
//记录题目给的关键数据(点的权值)
int out[MAX NODE];
//传说中的 fail 指针
int fail(MAX NODE);
//字母对应的 ID
int ID[1<<8];
//已使用节点个数
int nv:
//队列,用于广度优先计算 fail 指针
queue<int> Q;
//特定题目需要
const int MAXN = 25+5;
const int MOD = 20090717;
int N, M, K, d[2][MAX_NODE][1<<10];
namespace AC_Automaton
   //初始化,计算字母对应的儿子 ID,如:'a'->0 ... 'z'->25
   void Initialize()
       fail[0] = 0;
       for (int i = 0; i < MAX_CHD; i++)
          ID[i+'a'] = i;
   //重新建树需先 Reset
   void Reset()
       memset(chd[0], 0, sizeof(chd[0]));
       nv = 1:
   //将权值为 key 的字符串 a 插入到 trie 中
   void Insert(char *pat, int key)
       int u = 0:
       for (int i = 0; pat[i]; i++)
          int c = ID[pat[i]];
          if (!chd[u][c])
              memset(chd[nv], 0, sizeof(chd[nv]));
              out[nv] = 0;
```

```
chd[u][c] = nv++;
          }
           u = chd[u][c];
       out[u] = key;
   //建立 AC 自动机,确定每个节点的权值以及状态转移
   void Construct()
       for (int i = 0; i < MAX_CHD; i++)
           if (chd[0][i])
              fail[chd[0][i]] = 0;
               Q.push(chd[0][i]);
          }
       while (!Q.empty())
           int u = Q.front(); Q.pop();
           for (int i = 0; i < MAX_CHD; i++)
              int &v = chd[u][i];
              if (v)
                  Q.push(v);
                  fail[v] = chd[fail[u]][i];
                  //以下一行代码要根据题目所给 out 的含义来写
                  out[v] |= out[fail[v]];
               else
                  v = chd[fail[u]][i];
//解题
int solve()
   int tot = (1 << M)-1, ans = 0, s = 0, t = 1;
   memset(d[t], 0, sizeof(d[t]));
   d[t][0][0] = 1;
   for (int i = 0; i < N; i++)
       swap(s, t);
       memset(d[t], 0, sizeof(d[t]));
```

```
for (int u = 0; u < nv; u++)
           for (int a = 0; a <= tot; a++) if (d[s][u][a])
               for (int k = 0; k < MAX_CHD; k++)
                   int v = chd[u][k], b = (a|out[v]);
                   d[t][v][b] = (d[t][v][b]+d[s][u][a])%MOD;
   for (int a = 0; a <= tot; a++)
       int cnt = 0;
       for (int i = 0; i < M; i++)
           if (a&(1<<i))
               cnt++;
       if (cnt >= K)
           for (int u = 0; u < nv; u++)
               ans = (ans+d[t][u][a])%MOD;
   return ans;
int main()
   AC_Automaton::Initialize();
   while (scanf("%d%d%d", &N, &M, &K) != EOF)
       if (!N && !M && !K)
           break;
       AC_Automaton::Reset();
       for (int i = 0; i < M; i++)
           char temp[11];
           scanf("%s", temp);
           AC_Automaton::Insert(temp, 1<<i);</pre>
       AC_Automaton::Construct();
       printf("%d\n", solve());
   return 0;
```

#### AC 自动机+概率 DP\_hdu\_3689

```
字符集中有一些字符,给出每个字符的出现概率(它们的和保证为1),再给出一个串 S,问任
给一个长度为 N 的字符串 A (只能包含字符集中的字符), 使得 S 是 A 的子串的概率。
#include<cstdio>
#include<cstring>
#include<algorithm>
#include<queue>
using namespace std;
const int MAXN = 1000+5, MAXM = 10+5;
const int INF = 0x3f3f3f3f;
const int MAX_NODE = MAXN, MAX_CHD = 26;
int chd[MAX_NODE][MAX_CHD], fail[MAX_NODE], out[MAX_NODE];
int ID[1<<8], nv;
double P[MAX_CHD], d[MAXN][MAX_NODE];
char ch[5], word[MAXM];
queue<int> 0;
namespace AC Automaton
   void Initialize()
       fail[0] = 0;
       for (int i = 0; i < MAX_CHD; i++)
           ID[i+'a'] = i;
   void Reset()
       memset(chd[0], 0, sizeof(chd[0]));
       nv = 1;
   void Insert(char *pat)
       int u = 0;
       for (int i = 0; pat[i]; i++)
           int c = ID[pat[i]];
           if (!chd[u][c])
              memset(chd[nv], 0, sizeof(chd[nv]));
              out[nv] = 0;
               chd[u][c] = nv++;
```

```
u = chd[u][c];
        out[u]++;
    void Construct()
       for (int i = 0; i < MAX_CHD; i++)
            if (chd[0][i])
               fail[chd[0][i]] = 0;
               Q.push(chd[0][i]);
       while (!Q.empty())
           int u = Q.front(); Q.pop();
           for (int i = 0; i < MAX_CHD; i++)
               int &v = chd[u][i];
               if (v)
                    Q.push(v);
                    fail[v] = chd[fail[u]][i];
                else
                   v = chd[fail[u]][i];
int main()
   AC_Automaton::Initialize();
    while (scanf("%d%d", &N, &M))
       if (!N && !M)
           break;
       memset(P, 0, sizeof(P));
       memset(d, 0, sizeof(d));
       AC_Automaton::Reset();
       for (int i = 0; i < N; i++)
           scanf("%s", ch);
            scanf("%lf", &P[ID[ch[0]]]);
        scanf("%s", word);
```

#### AC 自动机+矩阵\_poj\_2778

```
问你长度为 N 的串中不包含模式串的串有几个
n属于1 ~ 2000000000 看到这个数据范围我们就应该敏感的想到这是矩阵~
最多 100 个结点, 先建好所有结点(不包括模式串结尾的和 fail 指向结尾的结点, 所以其实
最多只有 90 个有效结点)之间的转化关系,然后二分矩阵乘法,复杂度
O(100<sup>3</sup>*log(200000000))
*/
#include<cstdio>
#include<cstring>
#include<algorithm>
#include<queue>
using namespace std;
const int MAXM = 10+5;
const int MAX NODE = 100+5, MAX CHD = 4;
const long long MOD = 100000;
typedef long long MAT[MAX_NODE][MAX_NODE];
MAT g, G;
int M, N;
int chd[MAX NODE][MAX CHD], fail[MAX NODE], ID[1<<8], nv;
bool out[MAX_NODE];
char DNA[MAXM];
queue<int> Q;
namespace AC_Automaton
   void Initialize()
```

```
fail[0] = 0;
   ID['A'] = 0; ID['C'] = 1; ID['T'] = 2; ID['G'] = 3;
void Reset()
   memset(chd[0], 0, sizeof(chd[0]));
   nv = 1;
void Insert(char *pat)
   int u = 0;
   for (int i = 0; pat[i]; i++)
       int c = ID[pat[i]];
       if (!chd[u][c])
           memset(chd[nv], 0, sizeof(chd[nv]));
           out[nv] = 0;
           chd[u][c] = nv++;
       u = chd[u][c];
   out[u] = 1;
void Construct()
   for (int i = 0; i < MAX_CHD; i++)
       if (chd[0][i])
           fail[chd[0][i]] = 0;
           Q.push(chd[0][i]);
   while (!Q.empty())
       int u = Q.front(); Q.pop();
       for (int i = 0; i < MAX_CHD; i++)
           int &v = chd[u][i];
           if (v)
               Q.push(v);
               fail[v] = chd[fail[u]][i];
               out[v] |= out[fail[v]];
           }
           else
```

```
v = chd(fail(u))(i);
           }
namespace Matrix
   void Copy(int size, MAT x, MAT y)
       for (int i = 0; i < size; i++)
           for (int j = 0; j < size; j++)
               y[i][j] = x[i][j];
   void Mutiply(int size, MAT x, MAT y, MAT z)
       MAT tx, ty;
       Copy(size, x, tx);
       Copy(size, y, ty);
       for (int i = 0; i < size; i++)
           for (int j = 0; j < size; j++)
               z[i][i] = 0;
               for (int k = 0; k < size; ++k)
                   z[i][j] = (z[i][j]+tx[i][k]*ty[k][j])%MOD;
   void Power(int size, MAT x, int n, MAT y)
       MAT tx, r;
       Copy(size, x, tx);
       for (int i = 0; i < size; i++)
           for (int j = 0; j < size; j++)
               r[i][j] = (i == j ? 1 : 0);
       while (n)
           if (n&1)
               Mutiply(size, r, tx, r);
           n >>= 1:
           if (!n)
               break:
           Mutiply(size, tx, tx, tx);
        Copy(size, r, y);
```

```
int main()
   AC_Automaton::Initialize();
   memset(g, 0, sizeof(g));
   AC_Automaton::Reset();
   scanf("%d%d", &M, &N);
   for (int i = 0; i < M; i++)
       scanf("%s", DNA);
       AC_Automaton::Insert(DNA);
   AC_Automaton::Construct();
   for (int u = 0; u < nv; u++) if (!out[u])
       for (int k = 0; k < MAX_CHD; k++) if (!out[chd[u][k]])
           q[u][chd[u][k]]++;
   Matrix::Power(nv, g, N, G);
   long long ans = 0;
   for (int i = 0; i < nv; i++)
       ans = (ans+G[0][i])%MOD;
   printf("%lld\n", ans);
   return 0;
```

#### DP

#### 离散 DP hdu 4028

using namespace std;

const int MAX = 40, MAXN = MAX+5;

题意:给你n个钟的指针,第i个指针转一圈的时间是i单位,问你从n个钟任选一些指针使得,全部指针第一次回到原来的位置是经过的时间大于等于m,求又多少种选法。思路:显然时间是你选的指针的最小公倍数,但是好大,dp无从下手。看完神牛的题解才知道有一种 dp 叫做离散 dp,就是直接保存有用的状态就好了,其他的不用,这样空间就可以满足了,因为其实状态数很少。状态设定很简单:dp[i][j]:i表示以i指针结尾,最小公倍数(lcm)为j的方案数。转移也很简单就是dp[i][j]=dp[i][j]+dp[i-1][j];离散用了map,STL太强了,只能这么感慨,map要注意lcm的转移;还有初始状态为dp[i][i]=1;要在更新这个状态的时候加进去:\*/
#include<cstdio>
#include<cstdio>
#include<algorithm>
#include<map>

```
const int INF = 0x3f3f3f3f;
int T, N;
long long M;
struct cmp
   bool operator()(const long long a, const long long b)
       return a > b;
};
map<long long, long long, cmp> d[MAXN];
long long gcd(long long x, long long y)
   return !y ? x : gcd(y, x%y);
long long lcm(long long x, long long y)
   return x/gcd(x, y)*y;
int main()
   scanf("%d", &T);
   for (int i = 1; i <= MAX; i++)
       d[i] = d[i-1];
       d[i][i]++;
       map<long long, long long, cmp>::iterator p = d[i-1].begin();
       for (; p != d[i-1].end(); p++)
           d[i][lcm(p->first, i)] += p->second;
   for (int cas = 1; cas <= T; cas++)
       scanf("%d%I64d", &N, &M);
       long long ans = 0;
       map<long long, long long, cmp>::iterator p = d[N].begin();
       for (; p != d[N].end() && p->first >= M; p++)
           ans += p->second;
       printf("Case #%d: %I64d\n", cas, ans);
   return 0;
```

#### 区间 DP\_hdu\_4293\_1

```
题意:每个区间有权值,给若干区间,求最大收益。
思路: d[i]表示长度为i且包含以I结尾的区间时最大的人数。
#include<cstdio>
#include<cstring>
#include<algorithm>
using namespace std;
const int MAXN = 500+5;
const int INF = 0x3f3f3f3f;
int N, a, b, A[MAXN], B[MAXN], r[MAXN];
int mp[MAXN][MAXN], num[MAXN], d[MAXN];
bool cmp(const int a, const int b)
   return B[a] < B[b];
int main()
   while (scanf("%d", &N) != EOF)
       memset(mp, 0, sizeof(mp));
       memset(num, 0, sizeof(num));
       memset(d, 0, sizeof(d));
       int n = 0, ans = 0;
       for (int i = 1; i <= N; i++)
           scanf("%d%d", &a, &b);
           if (a+b >= N)
               continue;
           int &m = mp[a+1][N-b];
           if (!m)
               m = ++n;
               A[n] = a+1;
               B[n] = N-b;
               r[n] = n;
           num[m] = min(num[m]+1, N-a-b);
       sort(r+1, r+1+n, cmp);
       for (int i = 1; i <= n; i++)
           for (int j = 0; j < A[r[i]]; j++)
               d[B[r[i]]] = max(d[B[r[i]]], d[i]+num[r[i]]);
```

#### 树形背包 DP\_hdu\_4276

```
题意: 一个有 N 个节点的树形的地图,知道了每条变经过所需要的时间,现在给出时间 T,问
能不能在T时间内从1号节点到N节点。每个节点都有相对应的价值,而且每个价值只能被
取一次,问如果可以从1号节点走到n号节点的话,最多可以取到的最大价值为多少。
分析: 先求出从 1 号节点到 n 号节点的最短路,如果花费大于时间 T,则直接输出不符合,将
最短路上的权值全部赋值为 0, 在总时间 T 上减去最短路的长度, 表示最短路已经走过, 对其
它点进行树形背包求解,需要注意的是如果不是最短路上的边都要走两次,即走过去还要再走
回来, 状态转移方程: dp[i][j]=max(dp[i][j],dp[i][k]+dp[i][j-2*val-k])
*/
#include<cstdio>
#include<cstring>
#include<algorithm>
using namespace std;
const int MAXN = 100+5, MAXM = 500+5;
int N, T, a, b, t, A[MAXN];
int e, head[MAXN], next[MAXM], v[MAXM], w[MAXM];
int fa[MAXN], d[MAXN][MAXM];
void addedge(int x, int y, int z)
   v[e] = y; w[e] = z;
   next[e] = head[x]; head[x] = e++;
void mark(int u)
   for (int i = head[u]; i != -1; i = next[i]) if (v[i] != fa[u])
         fa[v[i]] = u;
         mark(v[i]);
void dfs(int u, int C)
   fill(d[u], d[u]+1+C, A[u]);
   for (int i = head[u]; i != -1; i = next[i]) if (v[i] != fa[u])
```

```
int cost = w[i]*2;
        if (cost <= C)
            dfs(v[i], C-cost);
            for (int j = C; j >= 0; j--)
                for (int k = 0; k \le j-cost; k++)
                    d[u][j] = max(d[u][j], d[u][j-k-cost]+d[v[i]][k]);
int main()
    while (scanf("%d%d", &N, &T) != EOF)
        e = 0;
       memset(head, -1, sizeof(head));
        for (int i = 1; i < N; i++)
            scanf("%d%d%d", &a, &b, &t);
            addedge(a, b, t);
            addedge(b, a, t);
        for (int i = 1; i <= N; i++)
            scanf("%d", &A[i]);
       int ans = 0;
       mark(1);
        for (int u = N; ;)
            ans += A[u];
           A[u] = 0;
            if (u == 1)
                break:
            for (int i = head[u]; i != -1; i = next[i]) if (v[i] == fa[u])
                u = v[i];
                T -= w[i];
                w[i] = 0;
                w[i^1] = 0;
                break:
            printf("Human beings die in pursuit of wealth, and birds die
in pursuit of food!\n");
        else
```

```
{
    dfs(1, T);
    ans += d[1][T];
    printf("%d\n", ans);
}
return 0;
}
```

## **KMP**

扩展 KMP\_hdu\_4300

```
这道题问的就是将 1 个串如何变为 stringA+stringB 的形式,使得 stringA 是 stringB
经过映射得到相同的串。映射那步其实没有什么价值,假设 str 为原串 s 经过映射后得到的
串,我们可以以 str 为模式串,以 s 为原串做一次扩展 KMP,得到 extend 数组, extend[i]
表示原串以第:开始与模式串的前缀的最长匹配。经过O(n)的枚举,我们可以得到,若
extend[i]+i=len 且 i>=extend[i]时,表示 stringB 即为该点之前的串, stringA 即
为该点之前的 str 串,最后输出即可。
#include<cstdio>
#include<cstring>
#include<algorithm>
using namespace std;
const int MAXN = 100000+5, MAXM = 50+5;
const int INF = 0x3f3f3f3f;
int T, extend[MAXN], next[MAXN];
char S[MAXM], tex1[MAXN], tex2[MAXN], match[1<<8];</pre>
void get next(char *pat)
   int len2 = strlen(pat), k = 0;
   next[0] = len2;
   while (k+1 < len2 && pat[k] == pat[k+1])
      k++;
   next[1] = k;
   for(int id = 1, i = 2; i < len2; i++)
       int u = i-id;
      if (next[u]+i >= next[id]+id)
          int j = next[id]+id-i;
          if (j < 0)
```

```
i = 0;
            while (j+i < len2 && pat[j] == pat[j+i])
                j++;
            next[i] = i;
            id = i;
        else
            next[i] = next[u];
void ext_kmp(char *str, char *pat)
    get_next(pat);
    int len1 = strlen(str), len2 = strlen(pat), k = 0;
    while (k < len1 && k < len2 && str[k] == pat[k])
        k++;
    extend[0] = k;
    for (int id = 0, i = 1; i < len1; i++)
        int u = i-id;
        if (i+next[u] < extend[id]+id)</pre>
            extend[i] = next[u];
        else
            int j = extend[id]+id-i;
            if (i < 0)
                i = 0;
            while (j+i < len1 && str[j+i] == pat[j])</pre>
                j++;
            extend[i] = j;
            id = i;
int main()
    scanf("%d", &T);
    while (T--)
        scanf("%s%s", S, tex1);
        int lenS = strlen(S);
        for (int i = 0; i < lenS; i++)
            match[(int)S[i]] = 'a'+i;
        int len = strlen(tex1);
        for (int i = 0; i < len; i++)
```

```
tex2[i] = match[(int)tex1[i]];
tex2[len] = 0;
ext_kmp(tex1, tex2);
for (int i = 0; i <= len; i++)
{
    if ((i+extend[i] == len && i*2 >= len) || i == len)
    {
        for (int j = 0; j < i; j++)
            printf("%c", tex1[j]);
        for (int j = 0; j < i; j++)
            printf("%c", tex2[j]);
        printf("\n");
        break;
    }
}
return 0;
}</pre>
```

#### 扩展 KMP\_hdu\_4333

```
扩展 KMP 能求出一个串所有后缀串(即 s[i...len])和模式串的最长公共前缀。于是只要将
这个串复制一遍,求出该串每个后缀与其本身的最长公共前缀即可,当公共前缀>=len时,
显然相等,否则只要比较下一位就能确定这个串与原串的大小关系。
至于重复串的问题,只有当这个串有循环节的时候才会产生重复串,用 KMP 的 next 数组求出
最小循环节。
*/
#include<cstdio>
#include<cstring>
#include<algorithm>
using namespace std;
const int MAXN = 100000+5, MAXM = 200000+5;
int T;
int extend[MAXM], next[MAXN], fail[MAXN];
char a[MAXN], aa[MAXM];
void get next(char *pat)
   next[0] = strlen(pat);
   int k = 0;
   while (pat[k+1] \&\& pat[k] == pat[k+1])
      k++;
   next[1] = k;
   for(int id = 1, i = 2; pat[i]; i++)
```

```
int u = i - id;
        if (next[u]+i >= next[id]+id)
            int i = next[id]+id-i;
            if (i < 0)
                j = 0;
            while (pat[j+i] && pat[j] == pat[j+i])
                j++;
            next[i] = j;
            id = i;
        else
            next[i] = next[u];
void ext_kmp(char *str, char *pat)
    get_next(pat);
    int k = 0;
    while (str[k] && pat[k] && str[k] == pat[k])
        k++;
    extend[0] = k;
    for (int id = 0, i = 1; str[i]; i++)
        int u = i-id;
        if (i+next[u] < extend[id]+id)</pre>
            extend[i] = next[u];
        else
            int j = extend[id]+id-i;
            if (j < 0)
             j = 0;
            while (str[j+i] && str[j+i] == pat[j])
             j++;
            extend[i] = j;
            id = i;
void get_fail(char *pat)
    fail[0] = -1;
    for (int i = 1, j = -1; pat[i]; i++)
```

```
while (j != -1 && pat[j+1] != pat[i])
           j = fail[j];
       if (pat[j+1] == pat[i])
           j++;
       fail[i] = j;
int main()
   scanf("%d", &T);
   for (int cas = 1; cas <= T; cas++)
       scanf("%s", a);
       int len = strlen(a);
       strcpy(aa, a);
       strcpy(aa+len, a);
       ext_kmp(aa, a);
       get_fail(a);
       int cir = len-fail[len-1]-1, cnt = 0;
       //求出循环节长度 cir, 原串循环不一定完整;
       if (len%cir)
           cir = len;
       for (int i = 0; i < cir; i++)
           if (extend[i] < len && aa[i+extend[i]] < a[extend[i]])</pre>
               cnt++;
       printf("Case %d: %d %d %d\n", cas, cnt, 1, cir-cnt-1);
   return 0;
```

# 大数

#### bign-bint

```
//比较高效的大数
#include<cstdio>
#include<cstring>
using namespace std;
const int base = 10000; // (base^2) fit into int
const int width = 4; // width = log base
const int maxn = 1000; // n*width: 可表示的最大位数
struct bint
{
```

```
int len, s[maxn];
   bint (int r = 0)
   { // r应该是字符串!
      for (len = 0; r > 0; r /= base)
          s[len++] = r*base;
   bint &operator = (const bint &r)
      memcpy(this, &r, (r.len+1)*sizeof(int));// !
      return *this;
};
bool operator < (const bint &a, const bint &b)
   int i;
   if (a.len != b.len) return a.len < b.len;
   for (i = a.len-1; i >= 0 && a.s[i] == b.s[i]; i--);
   return i < 0 ? 0 : a.s[i] < b.s[i];
bool operator <= (const bint &a, const bint &b)</pre>
   return !(b < a);
bint operator + (const bint &a, const bint &b)
   bint res; int i, cv = 0;
   for (i = 0; i < a.len || i < b.len || cy > 0; i++)
      if (i < a.len)
          cy += a.s[i];
      if (i < b.len)
          cy += b.s[i];
      res.s[i] = cy%base; cy /= base;
   res.len = i;
   return res;
bint operator - (const bint &a, const bint &b)
   bint res; int i, cy = 0;
   for (res.len = a.len, i = 0; i < res.len; i++)
      res.s[i] = a.s[i]-cy;
      if (i < b.len)
          res.s[i] -= b.s[i];
```

```
if (res.s[i] < 0)
          cy = 1, res.s[i] += base;
          cy = 0;
   while (res.len > 0 && res.s[res.len-1] == 0)
      res.len--:
   return res;
bint operator * (const bint &a, const bint &b)
   bint res; res.len = 0;
   if (0 == b.len)
      res.s[0] = 0;
      return res;
   int i, j, cy;
   for (i = 0; i < a.len; i++)
      for (j=cy=0; j < b.len | | cy > 0; j++, cy/= base)
          if (j < b.len)
             cy += a.s[i]*b.s[j];
          if (i+j < res.len)</pre>
             cy += res.s[i+j];
          if (i+j >= res.len)
             res.s[res.len++] = cy%base;
          else
             res.s[i+j] = cy%base;
   return res;
bint operator / (const bint &a, const bint &b)
{ // ! b != 0}
   bint tmp, mod, res;
   int i, lf, rg, mid;
   mod.s[0] = mod.len = 0;
   for (i = a.len-1; i >= 0; i--)
      mod = mod*base+a.s[i];
      for (lf = 0, rg = base-1; lf < rg; )
          mid = (lf+rg+1)/2;
```

```
if (b*mid <= mod)</pre>
             lf = mid;
          else
             rq = mid-1;
      res.s[i] = lf;
      mod = mod-b*lf;
   res.len = a.len;
   while (res.len > 0 \&\& res.s[res.len-1] == 0)
      res.len--;
   return res; // return mod 就是%运算
int digits(bint &a) // 返回位数
   if (a.len == 0) return 0;
   int l = (a.len-1)*4;
   for (int t = a.s[a.len-1]; t; ++1, t/=10);
   return 1;
bool read(bint &b, char buf[]) // 读取失败返回 0
   if (1 != scanf("%s", buf)) return 0;
   int w, u, len = strlen(buf);
   memset(&b, 0, sizeof(bint));
   if ('0' == buf[0] && 0 == buf[1]) return 1;
   for (w = 1, u = 0; len; )
      u += (buf[--len]-'0')*w;
      if (w*10 == base)
         b.s[b.len++] = u;
         u = 0;
          w = 1;
       else
          w *= 10;
   if (w != 1)
      b.s[b.len++] = u;
   return 1;
void write(const bint &v)
   int i;
```

#### bign-lrj

```
#include<cstdio>
#include<iostream>
using namespace std;
const int maxn = 200;
struct bign{
 int len, s[maxn];
 bign() {
  memset(s, 0, sizeof(s));
  len = 1;
 }
 bign(int num) {
  *this = num;
 }
 bign(const char* num) {
  *this = num;
 }
```

```
bign operator = (int num) {
 char s[maxn];
 sprintf(s, "%d", num);
 *this = s;
 return *this;
bign operator = (const char* num) {
 len = strlen(num);
 for(int i = 0; i < len; i++) s[i] = num[len-i-1] - '0';
 return *this;
string str() const {
 string res = "";
 for(int i = 0; i < len; i++) res = (char)(s[i] + '0') + res;
 if(res == "") res = "0";
 return res;
}
bign operator + (const bign& b) const{
 bign c:
 c.len = 0;
 for(int i = 0, g = 0; g \mid | i < max(len, b.len); i++) {
  int x = g;
   if(i < len) x += s[i];
   if(i < b.len) x += b.s[i];
   c.s[c.len++] = x % 10;
   g = x / 10;
 return c;
void clean() {
 while(len > 1 && !s[len-1]) len--;
bign operator * (const bign& b) {
 bign c; c.len = len + b.len;
 for(int i = 0; i < len; i++)
  for(int j = 0; j < b.len; j++)
     c.s[i+j] += s[i] * b.s[j];
 for(int i = 0; i < c.len-1; i++){
   c.s[i+1] += c.s[i] / 10;
   c.s[i] %= 10;
```

```
c.clean();
   return c;
  bign operator - (const bign& b) {
   bign c; c.len = 0;
   for(int i = 0, g = 0; i < len; i++) {
     int x = s[i] - g;
     if(i < b.len) x -= b.s[i];
     if(x >= 0) g = 0;
     else {
      g = 1;
      x += 10;
     c.s[c.len++] = x;
   c.clean();
   return c;
  bool operator < (const bign& b) const{</pre>
   if(len != b.len) return len < b.len;</pre>
   for(int i = len-1; i >= 0; i--)
     if(s[i] != b.s[i]) return s[i] < b.s[i];
   return false;
  bool operator > (const bign& b) const{
   return b < *this;
  bool operator <= (const bign& b) {</pre>
   return !(b > *this);
  bool operator == (const bign& b) {
   return !(b < *this) && !(*this < b);
  bign operator += (const bign& b) {
  *this = *this + b;
   return *this:
};
```

```
istream& operator >> (istream &in, bign& x) {
    string s;
    in >> s;
    x = s.c_str();
    return in;
}

ostream& operator << (ostream &out, const bign& x) {
    out << x.str();
    return out;
}

int main() {
    bign a;
    cin >> a;
    a += "123456789123456789000000000";
    cout << a*2 << endl;
    return 0;
}</pre>
```

#### bign-str

```
#include<cstdio>
#include<cstring>
using namespace std;
const int MAXSIZE = 200;
void Add(char *str1, char *str2, char *str3);
void Minus(char *str1, char *str2, char *str3);
void Mul(char *str1, char *str2, char *str3);
void Div(char *str1, char *str2, char *str3);
int main(void)
   char str1[MAXSIZE], str2[MAXSIZE], str3[MAXSIZE];
   while (scanf("%s %s", str1, str2) == 2)
      if (strcmp(str1, "0"))
          memset(str3, '0', sizeof(str3)); // !!!!!
          Add(str1, str2, str3);
          printf("%s\n", str3);
          memset(str3, '0', sizeof(str3));
          Minus(str1, str2, str3);
```

```
printf("%s\n", str3);
          memset(str3, '0', sizeof(str3));
          Mul(str1, str2, str3);
          printf("%s\n", str3);
          memset(str3, '0', sizeof(str3));
          Div(str1, str2, str3);
          printf("%s\n", str3);
      else
          if (strcmp(str2, "0"))
             printf("%s\n-%s\n0\n0\n", str2, str2);
          else
             printf("0\n0\n0\n0\n");
   return 0;
void Add(char *str1, char *str2, char *str3)
{// str3 = str1 + str2;
   int i, j, i1, i2, tmp, carry;
   int len1 = strlen(str1), len2 = strlen(str2);
   char ch:
   i1 = len1-1; i2 = len2-1;
   j = carry = 0;
   for (; i1 \ge 0 \&\& i2 \ge 0; ++i, --i1, --i2)
      tmp = str1[i1]-'0'+str2[i2]-'0'+carry;
      carry = tmp/10;
      str3[j] = tmp%10+'0';
   while (i1 >= 0)
      tmp = str1[i1--]-'0'+carry;
      carry = tmp/10;
      str3[j++] = tmp%10+'0';
   while (i2 >= 0)
      tmp = str2[i2--]-'0'+carry;
      carry = tmp/10;
      str3[j++] = tmp%10+'0';
   }
   if (carry)
      str3[j++] = carry+'0';
```

```
str3[i] = '\0';
   for (i = 0, --j; i < j; ++i, --j)
      ch = str3[i]; str3[i] = str3[i]; str3[i] = ch;
void Minus(char *str1, char *str2, char *str3)
{// str3 = str1-str2 (str1 > str2)}
   int i, j, i1, i2, tmp, carry;
   int len1 = strlen(str1), len2 = strlen(str2);
   char ch;
   i1 = len1-1; i2 = len2-1;
   j = carry = 0;
   while (i2 >= 0)
       tmp = str1[i1]-str2[i2]-carry;
      if (tmp < 0)
          str3[j] = tmp+10+'0'; carry = 1;
      else
          str3[j] = tmp+'0'; carry = 0;
       --i1; --i2; ++j;
   while (i1 >= 0)
      tmp = str1[i1]-'0'-carry;
      if (tmp < 0)
          str3[j] = tmp+10+'0'; carry = 1;
      else
          str3[j] = tmp+'0'; carry = 0;
      --i1; ++j;
   while (str3[j] == '0' \&\& j > 0)
      --j;
   str3[++j] = '\0';
   for (i=0, --j; i < j; ++i, --j)
```

```
ch = str3[i]; str3[i] = str3[j]; str3[j] = ch;
  }
void Mul(char *str1, char *str2, char *str3)
   int i, j, i1, i2, tmp, carry, jj;
   int len1 = strlen(str1), len2 = strlen(str2);
   char ch;
   jj = carry = 0;
   for (i1=len1-1; i1 >= 0; --i1)
      j = jj;
      for (i2=len2-1; i2 >= 0; --i2, ++j)
          tmp = (str3[j]-'0')+(str1[i1]-'0')*(str2[i2]-'0')+carry;
          if (tmp > 9)
             carry = tmp/10; str3[j] = tmp%10+'0';
          }
          else
             str3[j] = tmp+'0'; carry = 0;
      if (carry)
          str3[j] = carry+'0'; carry = 0; ++j;
      ++jj;
   --j;
   while (str3[j] == '0' \&\& j > 0)
      --j;
   str3[++j] = '\0';
   for (i=0, --j; i < j; ++i, --j)
      ch = str3[i]; str3[i] = str3[j]; str3[j] = ch;
void Div(char *str1, char *str2, char *str3)
   int i1, i2, i, j, jj, tag, carry, cf, c[MAXSIZE];
   int len1 = strlen(str1), len2 = strlen(str2), lend;
   char d[MAXSIZE];
   memset(c, 0, sizeof(c));
```

```
memcpy(d, str1, len2);
lend = len2; j = 0;
for (i1=len2-1; i1 < len1; ++i1)
   if (lend < len2)
      d[lend] = str1[i1+1]; c[j] = 0;
       ++j; ++lend;
   else if (lend == len2)
       jj = 1;
      for (i=0; i < lend; ++i)
          if (d[i] > str2[i]) break;
          else if (d[i] < str2[i])</pre>
              jj = 0; break;
       if(ii == 0)
          d[lend] = str1[i1+1]; c[j] = 0;
          ++j; ++lend;
          continue;
   if (jj==1 || lend > len2)
       cf = jj=0;
      while (d[jj] <= '0' && jj < lend)
          ++ii;
      if (lend-jj > len2)
          cf = 1;
       else if (lend-jj < len2)</pre>
          cf = 0;
       else
          i2 = 0; cf = 1;
          for (i = jj; i < lend; ++i)
             if (d[i] < str2[i2])</pre>
                 cf = 0; break;
```

```
else if (d[i] > str2[i2])
          break;
      ++i2;
}//else
while (cf)
   i2 = len2-1; cf = 0;
   for (i = lend-1; i >= lend-len2; --i)
      d[i] = d[i]-str2[i2]+'0';
      if (d[i] < '0')
          d[i] = d[i]+10; carry = 1;
          --d[i-1];
       else
          carry = 0;
       --i2;
   ++c[j]; jj=0;
   while (d[jj] <= '0' && jj < lend)
      ++jj;
   if (lend-jj > len2)
      cf = 1;
   else if (lend-jj < len2)</pre>
       cf = 0:
   else
      i2 = 0; cf = 1;
      for (i = jj; i < lend; ++i)
          if (d[i] < str2[i2])</pre>
             cf = 0; break;
          else if (d[i] > str2[i2])
             break;
          ++i2;
   }//else
```

```
}//while
       jj = 0;
       while (d[jj] \le '0' \&\& jj < lend)
          ++jj;
      for (i = 0; i < lend-jj; ++i)
          d[i] = d[i+jj];
      d[i] = str1[i1+1]; lend = i+1;
      ++j;
   }//else
}//for
i = tag = 0;
while (c[i] == 0)
   ++i;
for (; i < j; ++i, ++tag)
   str3[tag] = c[i]+'0';
str3[tag] = '\0';
```

# 后缀数组

# 第 K 个子串\_hdu\_3553

```
#include<cstdio>
#include<cstring>
#include<algorithm>
#include<set>
using namespace std;
const int MAXN = 100000+5;
int T;
int sa[MAXN], height[MAXN], rank[MAXN], tmp[MAXN], top[MAXN];
int Tr[MAXN<<2];</pre>
long long K, sumlen[MAXN];
char S[MAXN];
namespace SuffixArray
   void makesa(char *s, int n)
       int lena = n < 256 ? 256 : n;
       memset(top, 0, lena*sizeof(int));
       for (int i = 0; i < n; i++)
            top[rank[i] = s[i]&(-1)]++;
       for (int i = 1; i < lena; i++)
            top[i] += top[i-1];
```

```
for (int i = 0; i < n; i++)
           sa[--top[rank[i]]] = i;
       for (int k = 1; k < n; k <<= 1)
           for (int i = 0; i < n; i++)
               int j = sa[i]-k;
               if (j < 0)
                   j += n;
               tmp[top[rank[j]]++] = j;
           int j = sa[tmp[0]] = top[0] = 0;
           for (int i = 1; i < n; i++)
               if (rank[tmp[i]] != rank[tmp[i-1]] || rank[tmp[i]+k] !=
rank[tmp[i-1]+k])
                   top[++j] = i;
               sa[tmp[i]] = j;
           memcpy(rank, sa , n*sizeof(int));
           memcpy(sa , tmp, n*sizeof(int));
           if (j+1 >= n)
               break:
   void lcp(char *s, int n)
       height[0] = 0;
       for (int i = 0, k = 0, j = rank[0]; i+1 < n; i++, k++)
           while (k >= 0 \&\& s[i] != s[sa[j-1]+k])
               height[j] = k--;
               j = rank[sa[j]+1];
namespace SegTr
   void Build(int idx, int L, int R)
       if (L == R)
           Tr[idx] = R;
           return;
```

```
int mid = (L+R)>>1, left = idx<<1, right = idx<<1|1;
        Build(left, L, mid);
       Build(right, mid+1, R);
        Tr[idx] = (height[Tr[left]] <= height[Tr[right]] ? Tr[left] :</pre>
Tr[right]);
    int Query(int idx, int L, int R, int l, int r)
       if (1 <= L && R <= r)
            return Tr[idx];
       int mid = (L+R)>>1, left = idx<<1, right = idx<<1|1;
       int ql = 0, qr = 0;
       if (1 <= mid)</pre>
            ql = Query(left, L, mid, l, r);
       if (mid < r)
            qr = Query(right, mid+1, R, 1, r);
       if (ql && !qr)
            return ql;
        else if (!ql && qr)
            return qr;
        else
            return (height[ql] <= height[qr] ? ql : qr);</pre>
void solve(int len, int &rk, int &rl)
    int h = 0;
    long long a = 1, b = len;
    while (a < b)
       int q = SegTr::Query(1, 1, len, a+1, b);
       if (K <= (height[q]-h)*(b-a+1))</pre>
            rk = a; rl = h+1+(K-1)/(b-a+1);
            return;
       K = (height[q]-h)*(b-a+1);
        if (K <= sumlen[q-1]-sumlen[a-1]-height[q]*(q-a))</pre>
            b = q-1; h = height[q];
            continue;
       K -= sumlen[q-1]-sumlen[a-1]-height[q]*(q-a);
        a = q;
       h = height[q];
```

```
rk = a; rl = h+K;
int main()
   scanf("%d", &T);
   for (int cas = 1; cas <= T; cas++)
       scanf("%s%I64d", S, &K);
       int len = strlen(S);
       SuffixArray::makesa(S, len+1);
       SuffixArray::lcp(S, len+1);
       for (int i = 1; i <= len; i++)
           sumlen[i] = sumlen[i-1]+len-sa[i];
       SegTr::Build(1, 1, len);
       int rk, rl;
       solve(len, rk, rl);
       printf("Case %d: ", cas);
       for (int i = 0; i < rl; i++)
           printf("%c", S[sa[rk]+i]);
       printf("\n");
   return 0:
```

# 多串子串并集\_后缀数组\_hdu\_4416

```
/*
求多串的子串并集元素的个数,先用没出现过的不同的字符把多个串拼接,用后缀数组求这个串的不同子串的个数,再减去含有拼接字符的子串的个数。用上述方法求「A、B1、.....、BN」中不同子串的个数 sumB,答案就是 sumAB-sumB。
*/
#include<cstdio>
#include<algorithm>
using namespace std;
const int MAXN = 300000+5, MAXM = 100000+5;
int T, N, L[MAXM];
int len, sa[MAXN], height[MAXN], rank[MAXN], tmp[MAXN], top[MAXN];
int a[MAXN];
char A[MAXM];
void makesa(int *s, int n)
{
    int lena = n < 256 ? 256 : n;
```

```
memset(top, 0, lena*sizeof(int));
   for (int i = 0; i < n; i++)
        top[rank[i] = s[i]&(-1)]++;
    for (int i = 1; i < lena; i++)
        top[i] += top[i-1];
   for (int i = 0; i < n; i++)
        sa[--top[rank[i]]] = i;
   for (int k = 1; k < n; k <<= 1)
       for (int i = 0; i < n; i++)
           int j = sa[i]-k;
           if (j < 0)
                j += n;
           tmp[top[rank[j]]++] = j;
       int j = sa[tmp[0]] = top[0] = 0;
       for (int i = 1; i < n; i++)
            if (rank[tmp[i]] != rank[tmp[i-1]] || rank[tmp[i]+k] !=
rank[tmp[i-1]+k])
                top[++i] = i;
           sa[tmp[i]] = j;
       memcpy(rank, sa , n*sizeof(int));
       memcpy(sa , tmp, n*sizeof(int));
       if (i+1 >= n)
           break:
void lcp(int *s, int n)
   height[0] = 0;
   for (int i = 0, k = 0, j = rank[0]; i+1 < n; i++, k++)
       while (k \ge 0 \&\& s[i] != s[sa[j-1]+k])
           height[j] = k--;
           i = rank[sa[i]+1];
int main()
    scanf("%d", &T);
    for (int cas = 1; cas <= T; cas++)
```

```
scanf("%d%s", &N, A);
    len = 0;
   for (L[0] = 0; A[L[0]]; L[0]++)
       a[len++] = A[L[0]]-'a'+1;
    for (int i = 1; i <= N; i++)
       a[len++] = 26+i;
       scanf("%s", A);
       for (L[i] = 0; A[L[i]]; L[i]++)
            a[len++] = A[L[i]]-'a'+1;
    a[len] = 0;
    long long sumAB = 0, sumB = 0;
   makesa(a, len+1);
   lcp(a, len+1);
    for (int i = 1; i <= len; i++)
        sumAB += len-sa[i]-height[i];
    long long l = len;
    for (int i = 0; i < N; i++)
       1 -= L[i];
       sumAB -= (L[i]+1)*1;
       1--;
    len -= L[0]+1;
   makesa(a+L[0]+1, len+1);
   lcp(a+L[0]+1, len+1);
    for (int i = 1; i <= len; i++)
        sumB += len-sa[i]-height[i];
   l = len;
    for (int i = 1; i < N; i++)
       1 -= L[i];
       sumB -= (L[i]+1)*1;
       1--;
    printf("Case %d: %I64d\n", cas, sumAB-sumB);
return 0;
```

## 最长重复不重叠子串\_后缀数组+按 height 分组+二分\_poj\_1743

```
#include<cstdio>
```

```
#include<cstring>
#include<algorithm>
using namespace std;
const int MAXN = 20000+5;
const int INF = 0x3f3f3f3f;
int N, a[MAXN], s[MAXN];
int sa[MAXN], height[MAXN], rank[MAXN], tmp[MAXN], top[MAXN];
void makesa(int *s, int n)
    int lena = n < 256 ? 256 : n;
   memset(top, 0, lena*sizeof(int));
   for (int i = 0; i < n; i++)
        top[rank[i] = s[i]&(-1)]++;
   for (int i = 1; i < lena; i++)
        top[i] += top[i-1];
   for (int i = 0; i < n; i++)
        sa[--top[rank[i]]] = i;
   for (int k = 1; k < n; k <<= 1)
       for (int i = 0; i < n; i++)
           int j = sa[i]-k;
           if (j < 0)
               j += n;
            tmp[top[rank[j]]++] = j;
       int j = sa[tmp[0]] = top[0] = 0;
       for (int i = 1; i < n; i++)
            if (rank[tmp[i]] != rank[tmp[i-1]] || rank[tmp[i]+k] !=
rank[tmp[i-1]+k])
               top[++j] = i;
           sa[tmp[i]] = j;
       memcpy(rank, sa , n*sizeof(int));
       memcpy(sa , tmp, n*sizeof(int));
       if (j+1 >= n)
           break:
void lcp(int *s, int n)
   height[0] = 0;
   for (int i = 0, k = 0, j = rank[0]; i+1 < n; i++, k++)
        while (k >= 0 \&\& s[i] != s[sa[j-1]+k])
```

```
height[j] = k--;
           j = rank[sa[j]+1];
int main()
   while (scanf("%d", &N) && N)
       int len = 0;
       for (int i = 0; i < N; i++)
           scanf("%d", &a[i]);
           if (i)
               s[len++] = a[i]-a[i-1]+88;
       s[len] = 0;
       makesa(s, len+1);
       lcp(s, len+1);
       int l = 4, r = max(l+1, N/2), ans = -1;
       while (1 < r)
           int mid = (1+r)>>1, t = 0, mini = sa[0], maxi = sa[0];
           for (int i = 1; i <= len; i++)
               if (height[i] >= mid)
                   mini = min(mini, sa[i]);
                   maxi = max(maxi, sa[i]);
               else
                   t = max(t, maxi-mini);
                   mini = maxi = sa[i];
           t = max(t, maxi-mini);
           if (t > mid)
               ans = mid;
               l = mid+1;
           else
               r = mid;
```

```
printf("%d\n", ans+1);
}
return 0;
}
```

# 线段树

矩形并面积\_离散化+扫描线+线段树\_hdu\_4419

```
#include<cstdio>
#include<cstring>
#include<algorithm>
#include<map>
#define left(x) x<<1</pre>
#define right(x) x<<1|1</pre>
using namespace std;
const int MAXN = 10000+5, MAXM = 20000+5;
const int ALL = 1<<3;</pre>
int T, N, clr[MAXN], X[MAXM], Y[MAXM], y[MAXM], r[MAXM];
int ID[1<<8];
int Tsum[ALL][MAXM<<2], Tcov[ALL][MAXM<<2];</pre>
char C[5];
bool cmp(const int a, const int b)
    return X[a] < X[b];
//void Build(int idx, int L, int R)
//{
// for (int k = 1; k < ALL; k++)</pre>
        Tsum[k][idx] = Tcov[k][idx] = 0;
// if (R-L == 1)
//
        return;
// int mid = (L+R)>>1;
// Build(left(idx), L, mid);
// Build(right(idx), mid, R);
//}
void Update(int tr, int idx, int L, int R, int l, int r, int c)
    if (1 <= L && R <= r)
        Tcov[tr][idx] += c;
    else
        int mid = (L+R)>>1;
```

```
if (1 < mid)
           Update(tr, left(idx), L, mid, l, r, c);
       if (mid < r)
           Update(tr, right(idx), mid, R, l, r, c);
   if (Tcov[tr][idx])
       Tsum[tr][idx] = y[R-1]-y[L-1];
   else if (R-L == 1)
       Tsum[tr][idx] = 0;
   else
       Tsum[tr][idx] = Tsum[tr][left(idx)]+Tsum[tr][right(idx)];
int main()
   ID['R'] = 1 << 0; ID['G'] = 1 << 1; ID['B'] = 1 << 2;
   scanf("%d", &T);
   for (int cas = 1; cas <= T; cas++)
       scanf("%d", &N);
       for (int i = 0; i < N; i++)
           scanf("%s%d%d%d%d", C, &X[left(i)], &Y[left(i)],
&X[right(i)], &Y[right(i)]);
           clr[i] = ID[C[0]];
           y[left(i)] = Y[left(i)];
           y[right(i)] = Y[right(i)];
           r[left(i)] = left(i);
           r[right(i)] = right(i);
       int n = N << 1;
       sort(r, r+n, cmp);
       sort(y, y+n);
       map<int, int> dp;
       for (int i = 1; i <= n; i++)
           dp[y[i-1]] = i;
11
       Build(1, 1, n);
       long long area[ALL] = {};
       for (int i = 0; i < n; i++)
           for (int k = 1; k < ALL; k++)
               if (k&clr[r[i]>>1])
                   Update(k, 1, 1, n, dp[Y[left(r[i]>>1)]],
dp[Y[right(r[i]>>1)]], (r[i]&1 ? -1 : 1));
               if (i+1 < n)
                   area[k] += (long
```

# 线段树求矩形并周长\_hdu\_1828

```
思路:扫描线+线段树。记录完全覆盖住当前区间的线段条数,区间左右端点被几条线段覆盖。
叶节点表示长度为 1 的区间。用一个查询函数求一共有多少孤立线段。
*/
#include<cstdio>
#include<cstring>
#include<algorithm>
using namespace std;
const int MAXN = 20000+5, MAXM = 10000+5;
const int A = 10000, Len = 20000;
int N, x[MAXM], y[MAXM], r[MAXM];
int Tr[MAXN<<2], Tcov[MAXN<<2], cov1[MAXN<<2], covr[MAXN<<2],</pre>
mark[MAXN<<2];</pre>
bool cmpx(const int a, const int b)
   return x[a] < x[b];
bool cmpy(const int a, const int b)
   return y[a] < y[b];
//void Init(int idx, int L, int R)
//{
// if (L == R)
// {
//
       Tr[idx] = 0;
       covl[idx] = covr[idx] = 0;
```

```
// mark[idx] = 0;
//
      return;
// }
// Tr[idx] = 0;
// covl[idx] = covr[idx] = 0;
// mark[idx] = 0;
//}
void PushDown(int idx, int L, int R)
   int left = 2*idx, right = 2*idx+1;
   Tcov[left] += mark[idx];
   Tr[left] = Tcov[left] ? 1 : 0;
   covl[left] += mark[idx];
   covr[left] += mark[idx];
   mark[left] += mark[idx];
   Tcov[right] += mark[idx];
   Tr[right] = Tcov[right] ? 1 : 0;
   covl[right] += mark[idx];
   covr[right] += mark[idx];
   mark[right] += mark[idx];
   mark[idx] = 0;
void Update(int idx, int L, int R, int l, int r, int c)
   if (1 <= L && R <= r)
       Tcov[idx] += c;
       covl[idx] += c;
       covr[idx] += c;
       if (Tcov[idx] | R-L == 1)
           Tr[idx] = Tcov[idx] ? 1 : 0;
           mark[idx] += c;
           return;
   if (mark[idx])
       PushDown(idx, L, R);
   int mid = (L+R)/2, left = 2*idx, right = 2*idx+1;
   if (1 < mid)
       Update(left, L, mid, l, r, c);
   if (mid < r)
       Update(right, mid, R, l, r, c);
   covl[idx] = covl[left];
   covr[idx] = covr[right];
```

```
Tr[idx] = Tr[left]+Tr[right]-(covr[left] && covl[right] ? 1 : 0);
}
int main()
    while (scanf("%d", &N) != EOF)
        for (int i = 0, j; i < N; i++)
            j = 2*i;
            scanf("%d%d", &x[j], &y[j]);
            x[j] += A; y[j] += A;
            r[j] = j;
            j = 2*i+1;
            scanf("%d%d", &x[j], &y[j]);
            x[j] += A; y[j] += A;
            r[j] = j;
        int ans = 0;
        sort(r, r+2*N, cmpx);
        for (int i = 0; i < 2*N; )
            bool flag = 1;
            for (; (flag | | x[r[i]] == x[r[i-1]]) && i < 2*N; i++)
                flag = 0;
                int k = r[i];
                if (!(k%2))
                    Update(1, 0, Len, y[k], y[k^1], 1);
                else
                    Update(1, 0, Len, y[k^1], y[k], -1);
            if (i < 2*N)
                ans += (x[r[i]]-x[r[i-1]])*Tr[1]*2;
        sort(r, r+2*N, cmpy);
        for (int i = 0; i < 2*N; )
            bool flag = 1;
            for (; (flag | | y[r[i]] == y[r[i-1]]) && i < 2*N; i++)
                flag = 0;
                int k = r[i];
                if (!(k%2))
                    Update(1, 0, Len, x[k], x[k^1], 1);
                else
```

```
Update(1, 0, Len, x[k^1], x[k], -1);
}
if (i < 2*N)
    ans += (y[r[i]]-y[r[i-1]])*Tr[1]*2;
}
printf("%d\n", ans);
}
return 0;
}</pre>
```

## 线段树求体积并\_hdu\_3642

```
题意: 就是给你一些长方体, 求这些长方体相交至少 3 次的体积和。
思路: 对 z 轴扫描线,每次在 xy 平面对 x 轴扫描线、对 y 轴离散化用线段树求面积并,再把
分段求得的体积加和。
*/
#include<cstdio>
#include<cstring>
#include<algorithm>
using namespace std;
const int MAXN = 2000+5, MAXM = 2000+5, MAXP = 2000000+5;
int T, N, X[MAXM], Y[MAXM], Z[MAXM], rx[MAXM], ry[MAXM], rz[MAXM];
int Tr[MAXN<<2], Tcov[MAXN<<2], mark[MAXN<<2];</pre>
int match[MAXP], toy[MAXN];
bool cmpz(const int a, const int b)
   return Z[a] < Z[b];
bool cmpx(const int a, const int b)
   return X[a] < X[b];
bool cmpy(const int a, const int b)
   return Y[a] < Y[b];
//void Init(int idx, int L, int R)
//{
// if (R-L == 1)
// {
11
       Tr[idx] = 0;
11
       Tcov[idx] = 0;
//
   mark[idx] = 0;
```

```
return;
// }
// int mid = (L+R)/2, left = idx*2, right = idx*2+1;
// Init(left, L, mid);
// Init(right, mid, R);
// Tr[idx] = 0;
// Tcov[idx] = 0;
// mark[idx] = 0;
//}
void PushDown(int idx, int L, int R)
    int mid = (L+R)/2, left = idx*2, right = idx*2+1;
    Tcov[left] += mark[idx];
    Tr[left] = Tcov[left] > 2 ? toy[mid]-toy[L] : 0;
    mark[left] += mark[idx];
    Tcov[right] += mark[idx];
    Tr[right] = Tcov[right] > 2 ? toy[R]-toy[mid] : 0;
    mark[right] += mark[idx];
    mark[idx] = 0;
void Update(int idx, int L, int R, int l, int r, int c)
    if (1 <= L && R <= r)
       Tcov[idx] += c;
       if (Tcov[idx] > 2 | R-L == 1)
           mark[idx] += c;
           Tr[idx] = Tcov[idx] > 2 ? toy[R]-toy[L] : 0;
           return;
       }
    if (mark[idx])
       PushDown(idx, L, R);
    int mid = (L+R)/2, left = idx*2, right = idx*2+1;
    if (1 < mid)
       Update(left, L, mid, l, r, c);
    if (mid < r)
       Update(right, mid, R, l, r, c);
    Tr[idx] = Tr[left]+Tr[right];
int main()
    scanf("%d", &T);
    for (int cas = 1; cas <= T; cas++)
```

```
memset(match, 0, sizeof(match));
       scanf("%d", &N);
       for (int i = 0; i < N; i++)
           for (int j = 0; j < 2; j++)
               int k = 2*i+j;
               scanf("%d%d%d", &X[k], &Y[k], &Z[k]);
               Y[k] += 1000000;
               rx[k] = ry[k] = rz[k] = k;
           }
       sort(rx, rx+2*N, cmpx);
       sort(ry, ry+2*N, cmpy);
       sort(rz, rz+2*N, cmpz);
       int cnt = 0;
       for (int i = 0; i < 2*N; i++)
           if (!match[Y[ry[i]]])
               match[Y[ry[i]]] = ++cnt;
               toy[cnt] = Y[ry[i]];
           }
       long long ans = 0;
       for (int i = 0; i < 2*N; )
           long long area = 0;
           for (int j = 0; j < 2*N; )
               int curX = X[rx[j]];
               for (; curX == X[rx[j]] && j < 2*N; j++)
                   int k = rx[j]/2;
                   if (Z[2*k] \le Z[rz[i]] \&\& Z[rz[i]] < Z[2*k+1])
                       Update(1, 1, cnt, match[Y[2*k]],
match[Y[2*k+1]], (rx[j]&1 ? -1 : 1));
               if (j < 2*N)
                   area += (long long)(X[rx[j]]-X[rx[j-1]])*Tr[1];
           int curZ = Z[rz[i]];
           for (; curZ == Z[rz[i]] && i < 2*N; i++);
           if (i < 2*N)
               ans += (Z[rz[i]]-Z[rz[i-1]])*area;
       printf("Case %d: %I64d\n", cas, ans);
```

```
return 0;
}
```

#### 线段树区间修改单点查询 220B

```
题意: N个数, M个询问,每次问 Ai 到 Aj 里有多少个数 x 出现了 x 次。
思路: 离线+线段树区间修改、单点查询。按右端点将查询区间排序。扫描数列,假设当前数
a 第 x 次出现,那么当 x>=a 时,区间[pos[a][x-a]+1,pos[a][x-a+1]]上所有点+1;
当 x>a 时,区间[pos[a][x-a-1]+1,pos[a][x-a]]上所有点-1, pos[a][x]表示数 a
第 x 次出现的位置,为了方便,设所有数第一次出现的位置为 0。若当前扫描到的位置有查询
区间的右端点,则在线段树上查询左端点处的值,即为该次查询的答案。
*/
#include<cstdio>
#include<cstring>
#include<algorithm>
#include<vector>
using namespace std;
const int MAXN = 100000+5;
int N, M, a[MAXN], s[MAXN], t[MAXN], r[MAXN], ans[MAXN];
int Tr[MAXN<<2], mark[MAXN<<2];</pre>
vector<int> pos[MAXN];
bool cmp(const int a, const int b)
   return t[a] < t[b];
void PushDown(int idx)
   int left = idx <<1, right = (idx <<1)^1;
   Tr[left] += mark[idx];
   mark[left] += mark[idx];
   Tr[right] += mark[idx];
   mark[right] += mark[idx];
   mark[idx] = 0;
void Update(int idx, int L, int R, int l, int r, int c)
   if (1 <= L && R <= r)
       Tr[idx] += c;
       mark[idx] += c;
       return;
   if (mark[idx])
```

```
PushDown(idx);
   int mid = (L+R)>>1, left = idx<<1, right = (idx<<1)^1;
   if (1 <= mid)
        Update(left, L, mid, l, r, c);
   if (mid < r)
        Update(right, mid+1, R, 1, r, c);
int Query(int idx, int L, int R, int x)
   if (x == L \& R == x)
       return Tr[idx];
   if (mark[idx])
       PushDown(idx);
   int mid = (L+R)>>1, left = idx<<1, right = (idx<<1)^1;
   if (x \le mid)
       return Query(left, L, mid, x);
   else
       return Query(right, mid+1, R, x);
int main()
   scanf("%d%d", &N, &M);
   for (int i = 1; i <= N; i++)
       scanf("%d", &a[i]);
       if (a[i] <= N && !pos[a[i]].size())</pre>
           pos[a[i]].push_back(0);
   for (int i = 0; i < M; i++)
       scanf("%d%d", &s[i], &t[i]);
       r[i] = i;
   sort(r, r+M, cmp);
   for (int i = 1, j = 0; i \le N && j < M; i++)
       if (a[i] <= N)
           pos[a[i]].push_back(i);
           if (pos[a[i]].size() > a[i])
                Update(1, 1, N, pos[a[i]][pos[a[i]].size()-a[i]-1]+1,
pos[a[i]][pos[a[i]].size()-a[i]], 1);
           if (pos[a[i]].size() > a[i]+1)
                Update(1, 1, N, pos[a[i]][pos[a[i]].size()-a[i]-2]+1,
pos[a[i]][pos[a[i]].size()-a[i]-1], -1);
```

# 最长上升子序列

# 二维 LIS+方案输出\_sgu\_521

```
正向、反向分别求 LIS, 再枚举每个点.....
#include<cstdio>
#include<cstring>
#include<algorithm>
#include<vector>
using namespace std;
const int MAXN = 100000+5;
int N, x[MAXN], y[MAXN], id[MAXN];
int Y[MAXN], f[MAXN], d[2][MAXN], cnt[MAXN];
bool mark[MAXN];
bool cmp (const int &a, const int &b)
   if(x[a] != x[b])
       return x[a] < x[b];
    else
       return y[a] > y[b];
int LIS(int x)
    int maxi = 0;
    for (int i = 1; i \le N; i++)
       int j = lower bound(f+1, f+1+maxi, Y[i])-f;
       maxi = max(maxi, j);
       f[i] = Y[i];
       d[x][i] = i;
    return maxi;
```

```
int main()
   while (scanf("%d", &N) != EOF)
       memset(cnt, 0, sizeof(cnt));
       for (int i = 1; i <= N; i++)
           scanf("%d%d", &x[i], &y[i]);
           id[i] = i;
       sort(id+1, id+1+N, cmp);
       for (int i = 1; i <= N; i++)
           Y[i] = y[id[i]];
       int maxlen = LIS(0);
       for (int i = 1; i <= N; i++)
           Y[i] = -y[id[N-i+1]];
       LIS(1);
       vector<int> ans[2];
       for (int i = 1; i <= N; i++)
           mark[i] = (d[0][i]+d[1][N-i+1] == maxlen+1);
           if (mark[i])
               cnt[d[0][i]]++;
               ans[0].push_back(id[i]);
       for (int i = 1; i <= N; i++)
           if (mark[i] && cnt[d[0][i]] == 1)
               ans[1].push_back(id[i]);
       for (int i = 0; i < 2; i++)
           sort(ans[i].begin(), ans[i].end());
           printf("%u", ans[i].size());
           for (vector<int>::iterator it = ans[i].begin(); it !=
ans[i].end(); it++)
               printf(" %d", *it);
           printf("\n");
   return 0;
```

#### 某矩形的 LIS bupt 394

```
离线读入所有点(左下、右上),在左下点查询,右上点更新。
*/
#include<cstdio>
#include<cstring>
#include<algorithm>
using namespace std;
const int MAXN = 100000+5, MAXM = 200000+5;
int T, N, x[MAXM], y[MAXM], id[MAXM];
int f[MAXN], g[MAXN];
bool cmp (const int &a, const int &b)
    if (x[a] != x[b])
       return x[a] < x[b];
    else
       return y[a] > y[b];
int LIS(int n)
    int maxi = 0;
    for (int i = 0; i < n; i++)
       if (!(id[i]&1))
           g[id[i]>>1] = lower_bound(f+1, f+1+maxi, y[id[i]])-f;
       else
           if (g[id[i]>>1] > maxi)
               f[++maxi] = y[id[i]];
           else
               f[g[id[i]>>1]] = min(f[g[id[i]>>1]], y[id[i]]);
       }
    return maxi;
int main()
    scanf("%d", &T);
    while (T--)
       scanf("%d", &N);
       for (int i = 0; i < N; i++)
           scanf("%d%d%d%d", &x[i<<1], &y[i<<1], &x[i<<1|1],
```

## 最长上升子序列\_poj\_3903

```
#include<cstdio>
#include<cstring>
#include<algorithm>
using namespace std;
const int MAXN = 100000+5;
const int INF = 0x7ffffffff;
int N, a[MAXN], f[MAXN];
//int d[MAXN];
int main()
   while (scanf("%d", &N) != EOF)
       int maxi = 0;
       for (int i = 1; i <= N; i++)
           scanf("%d", &a[i]);
           int x = lower_bound(f+1, f+1+maxi, a[i])-f;
           maxi = max(maxi, x);
           f(x) = a[i];
11
           d[i] = x;
       printf("%d\n", maxi);
   return 0;
```

#### Mahjong\_hdu\_4431

```
#include<cstdio>
#include<cstring>
#include<algorithm>
```

```
#include<vector>
using namespace std;
const int MAX = 34;
const char *mahiong[] = {
    "1m", "2m", "3m", "4m", "5m", "6m", "7m", "8m", "9m",
    "1s", "2s", "3s", "4s", "5s", "6s", "7s", "8s", "9s",
    "1p", "2p", "3p", "4p", "5p", "6p", "7p", "8p", "9p",
    "1c", "2c", "3c", "4c", "5c", "6c", "7c"
};
int T, cnt[MAX];
char tile[10];
int id(char *s)
    if (s[1] == 'm')
        return s[0]-'1';
    else if (s[1] == 's')
        return 9+s[0]-'1';
    else if (s[1] == 'p')
        return 18+s[0]-'1';
    else
        return 27+s[0]-'1';
//bool check_standard_dfs(int dep)
//{
// if (dep == 5)
        return 1;
// bool res = 0;
// if (!dep)
// {
11
        for (int i = 0; i < MAX && !res; i++) if (cnt[i] >= 2)
11
11
            cnt[i] -= 2;
11
            res = check_standard_dfs(dep+1);
11
            cnt[i] += 2;
11
// }
// else
// {
11
        for (int i = 0; i < MAX && !res; i++)
11
11
            if (cnt[i] >= 3)
11
11
                cnt[i] -= 3;
//
                res = check_standard_dfs(dep+1);
//
                cnt[i] += 3;
```

```
11
           if (i < 27 && i%9 <= 6 && cnt[i] >= 1 && cnt[i+1] >= 1 &&
cnt[i+2] >= 1)
//
11
               for (int j = 0; j < 3; j++)
11
                   cnt[i+j]--;
               res = check_standard_dfs(dep+1);
11
11
               for (int j = 0; j < 3; j++)
11
                   cnt[i+j]++;
11
           }
11
// }
// return res;
//}
bool check_standard()
   bool res = 0;
   for (int i = 0; i < MAX && !res; i++) if (cnt[i] >= 2)
       int tmp[MAX], num = 0;
       memcpy(tmp, cnt, sizeof(cnt));
        tmp[i] -= 2;
       for (int j = 0; j < MAX; j++)
           if (tmp[j] >= 3)
                tmp[j] -= 3;
               num++;
           if (j < 27 \&\& j\%9 < 7)
               while (tmp[j] >= 1 \&\& tmp[j+1] >= 1 \&\& tmp[j+2] >= 1)
                   for (int k = 0; k < 3; k++)
                       tmp[j+k]--;
                   num++;
        res = (num == 4);
   return res;
bool check_ChiiToitsu()
```

```
for (int i = 0; i < MAX; i++)
       if (cnt[i] && cnt[i] != 2)
           return 0;
    return 1:
bool check_KokushiMuso()
    int res = 0;
    for (int i = 0; i < 3; i++)
       if (cnt[i*9+0] >= 1 && cnt[i*9+8] >= 1)
           res += cnt[i*9+0]+cnt[i*9+8];
       else
           return 0;
    for (int i = 27; i < MAX; i++)
        if (cnt[i] >= 1)
           res += cnt[i];
       else
           return 0;
    return (res == 14);
int main()
    scanf("%d", &T);
    while (T--)
       memset(cnt, 0, sizeof(cnt));
       for (int i = 0; i < 13; i++)
           scanf("%s", tile);
           cnt[id(tile)]++;
       vector<int> ans;
        for (int i = 0; i < MAX; i++) if (cnt[i] < 4)
           cnt[i]++;
           if (check_KokushiMuso() | check_ChiiToitsu() | 
check_standard())
               ans.push_back(i);
           cnt[i]--;
       if (ans.size())
```

```
{
    printf("%d", (int)ans.size());
    for (int i = 0; i < (int)ans.size(); i++)
        printf(" %s", mahjong[ans[i]]);
    printf("\n");
}
else
    printf("Nooten\n");
}
return 0;
}</pre>
```

#### **RMQ-ST**

```
#include<cstdio>
#include<cstring>
#include<cmath>
#include<algorithm>
using namespace std;
const int MAXN = 50000+5, MAXM = 16;
int N, Q;
int a[MAXN], st[MAXN][MAXM];
int pow2[MAXM];
inline int Most(const int &a, const int &b)
  return a > b ? a : b;
void InitRMQ(const int &n)
  pow2[0] = 1;
   for (int i = 1; i <= MAXM; i++)
      pow2[i] = pow2[i-1]<<1; //预处理2的i次方,最大次幂要大于MAXN
   for (int i = 1; i <= n; i++)
      stmax[i][0] = a[i];
   int k = int(log(double(n))/log(2.0))+1;
   for (int j = 1; j < k; j++)
      for (int i = 1; i <= n; i++)
         if (i+pow2[j-1]-1 <= n)
             stmax[i][j] = Most(stmax[i][j-1],
stmax[i+pow2[j-1]][j-1]);
         else
             break; // st[i][j] = st[i][j-1];
```

```
}
int Query(int x, int y) // x, y均为下标:1...n

{
    int k = int(log(double(y-x+1))/log(2.0));
    return Most(stmax[x][k], stmax[y-pow2[k]+1][k]);
}
int main()

{
    scanf("%d%d", &N, &Q);
    for (int i = 1; i <= N; i++)
        scanf("%d", &a[i]);
    InitRMQ(N);
    while (Q--)
    {
        int A, B;
        scanf("%d%d", &A, &B);
        int ans = Query(A, B);
    }
    return 0;
}
```

#### Trie 树\_编辑距离阈值匹配\_UVALive\_4769

```
求字典中存在前缀与查询串编辑距离小于阈值的词的个数
*/
#include<cstdio>
#include<cstring>
#include<algorithm>
#include <iostream>
using namespace std;
const int MAXM = 10+5;
const int MAX NODE = 3000000+5, MAX CHD = 26;
int N, M, edth;
int nv, chd[MAX_NODE][MAX_CHD], out[MAX_NODE], ID[1<<8];</pre>
int vis[MAX_NODE], mark[MAX_NODE];
char word[MAXM];
namespace Trie
   void Initialize()
       for (int k = 0; k < MAX CHD; k++)
           ID[k+'a'] = k;
```

```
void Reset()
       memset(chd[0], 0, sizeof(chd[0]));
       nv = 1:
   void Insert(char *pat)
       int u = 0;
       for (int i = 0; pat[i]; i++)
           int c = ID[pat[i]];
           if (!chd[u][c])
               memset(chd[nv], 0, sizeof(chd[nv]));
               out[nv] = 0;
               chd[u][c] = nv++;
           u = chd[u][c];
           out[u]++;
void dfs(int u, char *p, int d, int c)
   vis[u] = c;
   if (!(*p))
       mark[u] = c;
   if (mark[u] == c)
       return;
   if (chd[u][ID[*p]])
       dfs(chd[u][ID[*p]], p+1, d, c);
   if (d)
       for (int i = 0; i < MAX_CHD; i++) if (chd[u][i])
           dfs(chd[u][i], p, d-1, c);
       for (int i = 0; i < MAX_CHD; i++) if (chd[u][i])
           dfs(chd[u][i], p+1, d-1, c);
       dfs(u, p+1, d-1, c);
int calc(int u, int c)
   if (vis[u] != c)
       return 0;
   if (mark[u] == c)
```

```
return out[u];
    int res = 0;
    for (int i = 0; i < MAX_CHD; i++) if (chd[u][i])
       res += calc(chd[u][i], c);
    return res;
int main()
    scanf("%d", &N);
    Trie::Initialize();
    Trie::Reset();
    for (int i = 1; i <= N; i++)
        scanf("%s", word);
       Trie::Insert(word);
    scanf("%d", &M);
    for (int i = 1; i <= M; i++)
        scanf("%s%d", word, &edth);
       dfs(0, word, edth, i);
       printf("%d\n", calc(0, i));
    return 0;
```

# 编辑距离+BK 树\_hdu\_4323

```
1.dp 求编辑距离
2.bk 树找相差d的单词
*/
#include<cstdio>
#include<iostream>
#include<algorithm>
#include<queue>
using namespace std;
const int MAXN = 1500+5, MAXM = 10+5, MAXP = 400+5;
const int INF = 0x3f3f3f3f3;
int T, n, m, t, cnt;
int d[MAXM][MAXM], next[MAXN][MAXM];
char strl[MAXN][MAXM], str2[MAXM];
int Distance(char *s1, char *s2)
```

```
int 11 = strlen(s1), 12 = strlen(s2);
   for (int i = 0; i <= 11; i++)
      for (int j = 0; j \le 12; j++)
          if (!(i*j))
             d[i][j] = i+j;
          else
             d[i][j] = min(d[i-1][j]+1, d[i][j-1]+1);
             if (s1[i-1] == s2[j-1])
                d[i][j] = min(d[i][j], d[i-1][j-1]);
             else
                 d[i][j] = min(d[i][j], d[i-1][j-1]+1);
11
            printf("%d,%d:%d\n", i, j, d[i][j]);
   return d[11][12];
void dfs(int u)
   int dis = Distance(str1[u], str2);
   if (u && dis <= t)
      cnt++;
   for (int k = dis-t; k \le dis+t; k++)
      if (k \ge 0 \& next[u][k])
          dfs(next[u][k]);
int main()
   scanf("%d", &T);
   for (int cas = 1; cas <= T; cas++)
      memset(next, 0, sizeof(next));
      scanf("%d%d", &n, &m);
      strcpy(str1[0], "");
      for (int i = 1; i <= n; i++)
          scanf("%s", str1[i]);
          for (int j = 0; ;)
             int dis = Distance(str1[i], str1[j]);
             if (!next[j][dis])
                next[j][dis] = i;
```

#### 后缀自动机 SPOI LCS2

```
#include<cstdio>
#include<cstring>
#include<algorithm>
using namespace std;
const int MAXN = 100000+5, MAXM = 10+5;
char s[MAXN];
//MAX NODE = StringLength*2
const int MAX NODE = 500000+5;
//字符集大小,一般字符形式的题 26 个
const int MAX CHD = 26;
//已使用节点个数
int nv;
//每个节点的儿子,即当前节点的状态转移
int chd[MAX_NODE][MAX_CHD];
//此节点代表最长串的长度
int ml[MAX_NODE];
//父亲/失败指针
int fa[MAX_NODE];
//字母对应的 id
int id[1<<8];
//特定题目需要
int mml[MAX_NODE][MAXM], r[MAX_NODE];
```

```
namespace Suffix Automaton
   //初始化,计算字母对应的儿子 id,如:'a'->0 ... 'z'->25
   void Initialize()
       for (int i = 0; i < MAX_CHD; i++)
           id['a'+i] = i;
   //增加一个节点
   void Add(int u, int _ml, int _fa, int v = -1)
       ml[u] = _ml; fa[u] = _fa;
       if (v == -1)
           memset(chd[u], -1, sizeof(chd[u]));
       else
           memcpy(chd[u], chd[v], sizeof(chd[v]));
   //建立后缀自动机
   void Construct(char *str)
       nv = 1; Add(0, 0, -1);
       int cur = 0;
       for (int i = 0; str[i]; i++)
           int c = id[str[i]], p = cur;
           cur = nv++; Add(cur, i+1, -1);
           for (; p != -1 && chd[p][c] == -1; p = fa[p])
               chd[p][c] = cur;
           if (p == -1)
               fa[cur] = 0;
           else
               int q = chd[p][c];
               if (ml[q] == ml[p]+1)
                   fa[cur] = q;
               else
                   int r = nv++; Add(r, ml[q], fa[q], q);
                   ml[r] = ml[p]+1; fa[q] = fa[cur] = r;
                   for (; p != -1 \&\& chd[p][c] == q; p = fa[p])
                       chd[p][c] = r;
```

```
bool cmp(const int &a, const int &b)
    return ml[a] > ml[b];
int main()
    Suffix_Automaton::Initialize();
    scanf("%s", s);
    Suffix_Automaton::Construct(s);
    for (int i = 0; i < nv; i++)
       r[i] = i;
    sort(r, r+nv, cmp);
    memset(mml, 0, sizeof(mml));
    int cnt = 0;
    for (int i = 1; scanf("%s", s) != EOF; i++, cnt++)
       int 1 = 0, u = 0;
       for (int j = 0; s[j]; j++)
            int c = id[s[i]];
            if (chd[u][c] != -1)
               l++, u = chd[u][c];
            else
                while (u != -1 \&\& chd[u][c] == -1)
                   u = fa[u];
                if (u != -1)
                   1 = ml[u]+1, u = chd[u][c];
                else
                   1 = 0, u = 0;
            mml[u][i] = max(mml[u][i], 1);
    int ans = 0;
    for (int i = 0; i < nv; i++)
       int mini = ml[r[i]];
        for (int j = 1; j <= cnt; j++)
            mini = min(mini, mml[r[i]][j]);
            mml[fa[r[i]]][j] = max(mml[fa[r[i]]][j], mml[r[i]][j]);
```

```
ans = max(ans, mini);
}
printf("%d\n", ans);
return 0;
}
```

#### 斯坦纳树\_hdu\_4085

```
/*
斯坦纳树
最后的答案可能是一个森林,所以我们要先求出斯坦纳树后进行 DP。转移的时候要注意一点,
只有人的个数和房子的个数相等的时候才算合法状态,所以我们要加一个 check()函数进行
检查。
*/
#include<cstdio>
#include<cstring>
#include<algorithm>
#include<queue>
using namespace std;
const int MAXN = 50+5, MAXM = 2000+5;
const int MAX = 10;
const int INF = 0x3f3f3f3f;
int T, N, M, K, X, Y, Z;
int bit[MAXN], head[MAXN], e, next[MAXM], v[MAXM], w[MAXM];
int inq[MAXN][1<<MAX], d[MAXN][1<<MAX], dp[1<<MAX];
queue<int> Q;
void addedge(int x, int y, int z)
   v[e] = y; w[e] = z;
   next[e] = head[x]; head[x] = e++;
void init()
   e = 0;
   memset(head, -1, sizeof(head));
   memset(d, 0x3f, sizeof(d));
   memset(bit, 0, sizeof(bit));
   memset(ing, 0, sizeof(ing));
   memset(dp, 0x3f, sizeof(dp));
void spfa()
   while (!Q.empty())
```

```
int u = Q.front()&((1<<MAX)-1), st = Q.front()>>MAX;
        Q.pop();
        inq[u][st] = 0;
        for (int i = head[u]; i != -1; i = next[i])
            int nst = st|bit[v[i]];
            if (d[u][st]+w[i] < d[v[i]][nst])
                d[v[i]][nst] = d[u][st]+w[i];
                if (nst == st && !inq[v[i]][nst])
                    Q.push(nst<<MAX|v[i]);
                    inq[v[i]][nst] = 1;
            }
bool check(int st)
    int res = 0;
    for (int i = 0; i < K; i++)
       if (st&(1<<i))
            res++;
       if (st&(1<<(K+i)))
            res--;
    return !res;
int main()
    freopen("put.in", "r", stdin);
    scanf("%d", &T);
    while (T--)
       init();
        scanf("%d%d%d", &N, &M, &K);
        for(int i = 0; i < M; i++)
            scanf("%d%d%d", &X, &Y, &Z);
            addedge(X, Y, Z);
            addedge(Y, X, Z);
        int tot = (1 << (K << 1))-1;
```

```
for (int i = 1; i <= K; i++)
           bit[i] = 1 << (i-1);
           d[i][bit[i]] = 0;
           bit[N-K+i] = 1 << (K+i-1);
           d[N-K+i][bit[N-K+i]] = 0;
        for (int i = 0; i <= tot; i++)
           for (int j = 1; j <= N; j++)
               for (int k = (i-1)&i; k; k = (k-1)&i) //枚举i的所有子
                   d[j][i] = min(d[j][i],
d[j][k|bit[j]]+d[j][(i-k)|bit[j]]);
               if (d[j][i] < INF)</pre>
                   Q.push(i<<MAX|j);
                   inq[j][i] = 1;
           }
            spfa();
        for (int i = 0; i <= tot; i++)
           for (int j = 1; j <= N; j++)
                dp[i] = min(dp[i], d[j][i]);
       for (int i = 0; i \le tot; i++) if (check(i))
           for (int j = (i-1)&i; j; j = (j-1)&i) if (check(j))
               dp[i] = min(dp[i], dp[j]+dp[i-j]);
        if (dp[tot] < INF)</pre>
           printf("%d\n", dp[tot]);
        else
           printf("No solution\n");
   return 0;
```

# 最大非空连续和+方案\_hdu\_1003

```
#include<cstdio>
#include<cstring>
#include<algorithm>
using namespace std;
const int MAXN = 100000+5;
```

```
const int INF = 0x3f3f3f3f;
int T, N, a, s, t;
int main()
    scanf("%d", &T);
    for (int cas = 1; cas <= T; cas++)
        scanf("%d", &N);
       int sum = 0, mini = 0, maxi = -INF, p = 1;
        for (int i = 1; i \le N; i++)
            scanf("%d", &a);
            sum += a;
            if (sum-mini > maxi)
               maxi = sum-mini;
                s = p;
                t = i;
            if (sum < mini)</pre>
                mini = sum;
               p = i+1;
       if (cas > 1)
            printf("\n");
       printf("Case %d:\n", cas);
       printf("%d %d %d\n", maxi, s, t);
   return 0;
```