

[报告]J - I NEED A OFFER!

[source]

<http://202.114.18.202:8080/judge/contest/view.action?cid=6165#problem/B>

[Description]

Nim is a 2-player game featuring several piles of stones. Players alternate turns, and on his/her turn, a player's move consists of removing *one or more stones* from any single pile. Play ends when all the stones have been removed, at which point the last player to have moved is declared the winner. Given a position in Nim, your task is to determine how many winning moves there are in that position.

A position in Nim is called "losing" if the first player to move from that position would lose if both sides played perfectly. A "winning move," then, is a move that leaves the game in a losing position. There is a famous theorem that classifies all losing positions. Suppose a Nim position contains  $n$  piles having  $k_1, k_2, \dots, k_n$  stones respectively; in such a position, there are  $k_1 + k_2 + \dots + k_n$  possible moves. We write each  $k_i$  in binary (base 2). Then, the Nim position is losing if and only if, among all the  $k_i$ 's, there are an even number of 1's in each digit position. In other words, the Nim position is losing if and only if the *xor* of the  $k_i$ 's is 0.

Consider the position with three piles given by  $k_1 = 7$ ,  $k_2 = 11$ , and  $k_3 = 13$ . In binary, these values are as follows:

```
  111
 1011
 1101
```

There are an odd number of 1's among the rightmost digits, so this position is not losing. However, suppose  $k_3$  were changed to be 12. Then, there would be exactly two 1's in each digit position, and thus, the Nim position would become losing. Since a winning move is any move that leaves the game in a losing position, it follows that removing one stone from the third pile is a winning move when  $k_1 = 7$ ,  $k_2 = 11$ , and  $k_3 = 13$ . In fact, there are exactly three winning moves from this position: namely removing one stone from any of the three piles.

[Solution]

是 nim 的一个变化。令  $ans = a_1 \oplus a_2 \oplus \dots \oplus a_n$ , 第一步的方法数就是减少某一堆的值使  $ans = 0$  的方法数, 如果需要找出异或值为 0 的数, 所以对于某堆石子  $a_i$ , 现在对于  $ans \oplus a_i$ , 就是除了  $a_i$  以外其他的石子的异或值, 如果  $ans \oplus a_i \leq a_i$ , 那么对于  $a_i$  的话, 是可以减小到  $ans \oplus a_i$  的值, 然后使得所有数的异或值为 0。

[Code]

```
#include<stdio>
#include<stdlib>
int main()
{
    int a[10000];
    int n,t;
```

```

int ans=0;
while (scanf("%d",&n))
{
    if (!n) break;
    ans=0;
    t=0;
    for (int i=1;i<=n;i++)
    {
        scanf("%d",&a[i]);
        t=t^a[i];
    }
    if (t==0)
        printf("0\n");
    else
    {
        for (int i=1;i<=n;i++)
        {
            if ((t^a[i])<=a[i])
                ans++;
        }
        printf("%d\n",ans);
    }
}
return 0;
}

```