Correctness of Bellman-Ford's Algorithm:

This proof is made by Contradiction:

show  $d[v] = \delta(s,v)$  after |V|-1 passes

If  $shortest\ distance[v] > shortest[u] + w\ then\ we\ can't\ update\ that\ shortest[v] = shortest[u] + w$ 

- Lemma:  $d[v] ≥ \delta(s,v)$  always Initially true Let v be first vertex for which  $d[v] < \delta(s,v)$
- Let u be the vertex that caused d[v] to change: d[v] = d[u] + w(u,v)
- If  $d[v] < \delta(s,v)$  then

$$\delta(s,v) \leq \delta(s,u) + w(u,v)$$

$$\delta(s,u) + w(u,v) \le d[u] + w(u,v)$$

Therefore, d[v] < d[u] + w(u,v). Contradiction.