Topological k-means: a graph-based algorithm for data clustering using geodesic distances

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Abstract

Clustering is one of the most important tasks in machine learning and data science. Several clustering algorithms have been proposed to mitigate limitations from unsupervised learning in pattern recognition. However, clustering high dimensional data is still a challenge. In this paper, we propose topological k-means, a graph-based method for data clustering that uses the Dijkstra algorithm to compute geodesic distances between sample points in the discrete approximate data manifold. Moreover, the computational complexity of the proposed algorithm is reasonably low in comparison to modern clustering techniques, as it is linear in the number of samples and also in the number of edges of the k-NN graph. Computational experiments with real datasets show that the proposed method is capable of improving the quality of the obtained clusters in terms of external validity measures in comparison to regular k-means and a state-of-the-art approach, the HDBSCAN algorithm, defining a viable and promissing alternative to existing clustering algorithms.

Keywords: Clustering, k-means, HDBSCAN, high dimensional data, geodesic distance

1 Introduction

Machine learning has witnessed unprecedented growth in recent years, and within this expansive domain, data clustering stands out as a pivotal technique with profound implications for data analysis and pattern recognition [1–3]. Clustering, in the domain of machine learning, is a technique that involves grouping similar data points together, based on certain inherent characteristics, without the need for predefined labels. The fundamental premise underlying clustering is the identification of natural groupings or patterns within data, which facilitates a more nuanced understanding of the underlying structure [4]. This unsupervised learning approach has far-reaching applications across various domains, including image processing, pattern recognition, and data mining [5].

The most popular clustering method is k-means, a partitional algorithm, which aim to divide a dataset into k distinct, non-overlapping subgroups or clusters by minimizing the sum of squared distances between data points and the centroid of their assigned cluster [6, 7]. The main advantages of k-means are: 1) computational efficiency, as it is linear in the number of samples, making it suitable for large datasets; and 2) simplicity and versatility, as it is easy to implement and applicable in various domains, as it can hadle different types of data. However, k-means also has several limitations, among which we can cite: 1) as it relies on the Euclidean distance, it suffers from a kind of non-spherical blindness and high sensitivity to outliers in data; and 2) it struggles with clusters with varying densities and shapes.

Modern clustering algorithms are often capable of overcoming some limitations of k-means. Hierarchical Density-Based Spatial Clustering of Applications with Noise, or simply HDBSCAN, is an advanced clustering algorithm that extends the capabilities of traditional density-based methods, such as DBSCAN [8]. The key innovation lies in its ability to discover clusters of varying shapes and densities in a dataset while being robust to noise. HDBSCAN operates by constructing a hierarchical clustering tree and then extracting stable and significant clusters from this structure [9]. The main advantages of HDBSCAN are: 1) noise handling, as it explicitly identifies and labels noise points, providing a clear distinction between valid clusters and noise in the data; 2) flexibility in cluster shapes, as it is not constrained by assumptions about the shape or size of clusters, making it suitable for datasets with complex structures; and 3) adaptability to different densities, as it adapts well to clusters with varying densities, making it particularly effective in scenarios where clusters exhibit different levels of compactness. On the ohter hand, the main limitations of HDBSCAN are [10]: 1) computational complexity, as it is an intensive algorithm, especially for large datasets; 2) sensitivity to parameters, as the performance can be significantly influenced by tuning its parameters; and 3) limited for global structure, as it focuses on local density structures, and while it is excellent for identifying microclusters. Another issue with HDBSCAN is related to performance and approximation. For huge datasets, a random sample for k-means can get a good approximation of the overall result. But for HDBSCAN, it's not clear how to evaluate it in subsets of the data.

Despite many advances, one open problem in clustering is how to deal with high dimensional data, that is, situations in which the number of features m is order of magnitudes larger than the number of samples n [11–13]. To tackle this problem, we propose topological k-means, or simply top k-means, a graph-based algorithm that

computes geodesic distances between samples in the data manifold using Dijkstra's algorithm. Basically, the ideia is to approximate the data manifold by a k-NN graph and then use the length of the shortest paths as an approximation to the underlying geodesic distances. The main contributions of this paper are twofold: 1) after building a graph representation and replacing the Euclidean distances by the geodesic distances, topological k-means is able to detect clusters with different shapes, depending on the graph topology; and 2) topological k-means avoids the lack of discriminating power of the Euclidean distance in high dimensional spaces, producing meaningful clusters even when the number of features m is much larger than the number of samples n (the curse of the dimensionality).

The remaining of the paper is organized as follows: Section 2 presents an overview of the k-means clustering algorithm. Section 3 discusses Dijkstra's algorithm in detail, showing that is always returns the geodesic distances in graphs whose edge weights are positive. Section 4 describes the proposed topological k-means algorithm, explaining how it works and analyzing its computational complexity. Section 5 shows the computational experiments with regular k-means, topological k-means and HDBSCAN, presenting the obtained results in terms of three external cluster quality mesures: Rand index, normalized mutual information index and V-measure. Finally, Section 6 presents our conclusions and final remarks.

2 The k-means clustering algorithm

The k-means algorithm is an unsupervised method in the sense that learning is completely autonomous. Its first version was originally proposed in 1957 by Stuart P. Loyd in the context of pulse code modulation, but only published decades later, in 1982 [14]. The basic idea of k-means consists in grouping points according to a similarity measure, with the obtained results (clusters) strongly depending on the similarity measure chosen. In its regular version, k-means adopts the Euclidean distance, which makes the algorithm "blind" for non spherical clusters.

The problem formulation is as follows: given a data matrix $X^T = \{\vec{x}_1, \vec{x}_2, ..., \vec{x}_n\}$, where $\vec{x}_i \in R^d$, the objective is to partition X in k < n groups or clusters in order to minimize the intra-cluster scattering, that is, to find the optimal partition $S^* = \{s_1^*, s_2^*, ..., s_k^*\}$ that satisfies the following criterion, also known as within-cluster sum of the squares (WCSS) [15]:

$$S^* = \underset{S}{\operatorname{arg\,min}} \sum_{i=1}^k \sum_{\vec{x} \in s_i} ||\vec{x} - \vec{\mu}_i||^2$$
 (1)

where k denotes the number of partitions and $\vec{\mu}_i$ is the centroid of partition set s_i . Note that, intuitively, the objective function determines that we must make the centroid of each partition be as close as possible to the data points that define that partition.

It has been shown that finding the optimal solution to this problem is NP-Hard for any number of dimensions d (even two dimensions). The heuristic adopted by k-means to simplify the problem and hence lead to a polynomial time algorithm, consists of fixing the value of k. Therefore, the parameter k (number of clusters) must be chosen

before the algorithm is executed, what is unknown is some situations. In the following, we present two major results that help us better understand the k-means algorithm.

Theorem 1. Given a non-empty cluster s_k , its centroid or mean is the only choice of center which minimizes internal scattering

$$WS(\vec{c}^{(k)}) = \sum_{\vec{x} \in s_k} ||\vec{x} - \vec{c}^{(k)}||^2$$
 (2)

First, note that:

$$WS(\vec{c}^{(k)}) = \sum_{\vec{x} \in s_k} ||\vec{x} - \vec{c}^{(k)}||^2 = \sum_{\vec{x} \in s_k} (\vec{x} - \vec{c}^{(k)})^T (\vec{x} - \vec{c}^{(k)})$$
(3)

Applying the distributive, we have:

$$WS(\vec{c}^{(k)}) = \sum_{\vec{x} \in s_k} \left(\vec{x}^T \vec{x} - 2\vec{x}^T \vec{c}^{(k)} + \vec{c}^{(k)}^T \vec{c}^{(k)} \right)$$
(4)

The necessary condition for the minimization is:

$$\frac{\partial}{\partial \vec{c}^{(k)}} W S(\vec{c}^{(k)}) = 0 \tag{5}$$

which leads to:

$$2\sum_{\vec{x}\in s_k} \vec{x} - 2\sum_{\vec{x}\in s_k} \vec{c}^{(k)} = 0 \tag{6}$$

whose solution is:

$$\bar{c}^{(k)} = \frac{1}{n_k} \sum_{\vec{x} \in s_k} \vec{x} \tag{7}$$

where n_k is the number of samples in the k-th cluster. The following result shows that the nearest centroid rule, that is, assign each data point \vec{x} to the cluster whose center is the closest, produces an optimal solution to the k-means clustering problem.

Theorem 2. Given a set X of data points and a sequence of k centroids $\vec{c}^{(1)}, \vec{c}^{(2)}, ..., \vec{c}^{(k)}$ a partition into clusters minimize the objective function

$$\sum_{i=1}^{k} \sum_{\vec{x} \in s_i} ||\vec{x} - \vec{\mu}_i||^2 \tag{8}$$

if the algorithm assigns each sample \vec{x} to the cluster s_i with the closest centroid.

The proof is trivial, since each sample \vec{x} contributes exactly a single time to the sum defined by the previous objective function and when choosing to assign \vec{x} to the cluster whose centroid is the closest, we clearly minimize the contribution to the

objective function. Algorithm 1 illustrates the pseudo-code of the k-means clustering method.

Algorithm 1 k-means clustering algorithm

```
// Parameters:
//X: the n \times d data matrix (each row is a sample)
// n: the number of samples
// k: the number of clusters
function K-MEANS(X, n, k)
     \vec{s}_1, \vec{s}_2, ..., \vec{s}_k \leftarrow select\_random\_seeds(X, k)
                                                                                   ▷ Select k random centers
     for i \leftarrow 1; i \leq k; i + + do
          \vec{\mu}_i \leftarrow \vec{s}_i
     end for
     while not convergence do
          for i \leftarrow 1; i \leq k; i + + do
               \omega_i \leftarrow \{\}
                                                                          \triangleright Begin with k empty partitions
          end for
          for i \leftarrow 1; i \leq n; i + + \mathbf{do}
               j \leftarrow \arg\min_{l} \|\vec{\mu}_{l} - \vec{x}_{i}\|
                                                                                       ⊳ Find nearest centroid
              \omega_i \leftarrow \omega_j \cup \{\vec{x}_i\}
                                                                             ▶ Assign sample to the cluster
          end for
          for i \leftarrow 1; i \leq k; i + + do
              \vec{\mu}_i \leftarrow \frac{1}{|\omega_i|} \sum_{\vec{x} \in \omega_i} \vec{x}
                                                                                 ▶ Recalculate the centroids
          end for
     end while
     return \omega_1, \omega_2, ..., \omega_k
end function
```

The convergence of the k-means algorithm is usually measured in three different ways: 1) the centroids of the newly formed clusters do not move or their movement is less than a threshold; 2) the data points remain in the same clusters; and 3) a maximum number of iterations is reached.

2.1 Complexity analysis

Basically, the computational complexity of the k-means algorithm essentially depends on the following factors: the number of samples n, the number of centers k, the number of features d and the number of iterations for convergence t (which is often unknown).

First, note that the select_random_seeds(X, k) function and the initial FOR loop (outside the while) are accomplished in O(k). The critical part is executing the WHILE loop, which we will analyze next. To obtain the closest centroid, note that we must compare it against k centroids. Knowing that k Euclidean distances will be necessary for this and that, to calculate each distance, we have a cost of O(d). Thus, to decide which centroid is the closest to a sample, the computational cost is O(kd). As the process is repeated for all n data points, we have a cost of O(nkd). Finally, since this

process is repeated for an unknown number t of iterations, the total cost of the k-means algorithm is O(tnkd), which is nevertheless linear in the number of samples, leading to a quite efficient algorithm.

3 Geodesic distances in graphs

The problem of finding geodesic distances in graphs is solved by the computation of shortest paths in weighted graphs.

Definition 1. Given a weighted graph G = (V, E, w), where V is the set of vertices, E is the set of edges and $w : E \to R^+$ is a weighting function for the edges, the weight of a path $P = v_0v_1v_2...v_n$ is given by:

$$w(P) = \sum_{i=1}^{n} w(v_{i-1}, v_i)$$
(9)

The optimum path P^* from v_0 to v_n is a minimizer for w(P), that is:

$$P^* = \arg\min_{P} w(P) \tag{10}$$

provided there is a path between v_0 and v_n , otherwise the cost of the path is infinite. In the following, we present a nice property of optimum paths.

Theorem 3. Let G = (V, E, w) and $P = v_0 v_1 v_2 ... v_n$ the optimum path from v_0 to v_n . For $0 \le i < j \le n$, let $P' = v_i v_{i+1} ... v_j$ a subpath of P. Then, P' is the optimum path from v_i to v_j .

In other words, any subpath from an optimum path P is also optimum. This result shows that dynamic programming can be applied in the solution of shortest path problems. Dijkstra's algorithm combines dynamic programming and a greedy strategy to solve shortest path problems in a computationally efficient way. First, we define the main variables used by Dijkstra's algorithm.

- $\lambda(v)$: length of the shortest path from the root s to vertex v.
- $\pi(v)$: predecessor of vertex v in the shortest path tree.
- Q: priority queue to store the vertices (the smaller $\lambda(v)$, the higher the priority).

The priority queue Q has three basic primitives:

- Insert (Q, v): insert a vertex v at the end of Q.
- ExtractMin(Q): remove from Q the vertex having the smallest $\lambda(v)$.
- DecreaseKey $(Q, v, \lambda(v))$: update the priority of vertex v in Q.

Given the above, Dijkstra's method for building shortest paths in weighted graphs is presented in Algorithm 2 [16].

Algorithm 2 Dijkstra's algorithm

```
// Parameters:
//G: the input graph with n vertices and m edges
//w: the edge weights
// s: the root of the optimum path tree
function Dijkstra(G, w, s)
    for each v \in V do
                                                                            \triangleright Initialize \lambda(v) and \pi(v)
         \lambda(v) \leftarrow \infty
         \pi(v) \leftarrow nil
     end for
     \lambda(s) \leftarrow 0
                                                                         \triangleright Root starts with \lambda(s) = 0
    S \leftarrow \emptyset
    Q \leftarrow \emptyset
    for each v \in V do
                                                                        \triangleright Insert all the vertices in Q
         Insert(Q, v)
     end for
     while Q \neq \emptyset do
                                                                                              ▶ Main loop
         u \leftarrow ExtractMin(Q)
                                                                  ▷ Extract highest priority vertex
         S \leftarrow S \cup \{u\}
         for each v \in N(u) do
                                                                        \triangleright Process each neighbor of u
             \lambda(v) \leftarrow min\{\lambda(v), \lambda(v) + w(u, v)\}
                                                                   \triangleright Is it worth to reach v from u?
             if \lambda(v) was updated then
                                                           \triangleright Found a better entrance from s to u
                  \pi(v) \leftarrow u
                                                                             \triangleright Update v's predecessor
                  Decrease_Key(Q, v, \lambda(v))
                                                                               ▶ Update the priorities
             end if
         end for
     end while
end function
```

In summary, the algorithm works as follows: initially, all the vertices are initialized with $\lambda(v) = \infty$, because at this point, we do not know whether there will be a path from s to every other vertex or not. After that, all vertices are inserted into the priority queue and the root s is the only one to have its priority set to zero (this makes the root to be the first vertex to leave Q). At each iteration, the highest priority (lowest $\lambda(v)$) vertex u is removed from Q and inserted in S. An important property is that, when a vertex v is inserted in S, its $\lambda(v)$ does not change anymore, because at this point $\lambda(v) = d(s, v)$. After removing u from Q, we look at every neighbor v of u and check if it is a good idea to pass through u to reach v. This operation is known as the relaxation of the edge (u, v). If $\lambda(v)$ is reduced (u is a good route for v), the we update the predecessor of v and the priority of v in Q. This process is repeated until there is no more vertices in Q.

3.1 Complexity analysis

The implementation of Djikstra's algorithm can use static (adjacency matrix for G and a simple array for Q) or dynamic data structures (adjacency list for G and a binary heap for Q). In the following, we discuss the computational complexity of the algorithm when using dynamic strutures due to better memory management.

First of all, note that according to the algorithm, we have the following facts:

- The initialization of $\lambda(v)$ and $\pi(v)$ for every vertex in G is O(n).
- The insertion of the vertices in Q can be performed in O(n) in the binary heap, because initially, all vertices have the same priority.
- \bullet the main WHILE loop is executed exactly n times, one for each vertex.
- The primitive $\operatorname{ExtractMin}(Q)$ is $O(\log n)$, as it is proportional to a seach in a binary tree.
- The number of iterations of the FOR loop depends essentially on the degree of the vertex u, that is, d(u). Without loss of generality, suppose that the vertices are removed from Q in the following order: $v_1, v_2, ..., v_n$. Hence, the total number of iterations of the FOR loop is $d(v_1) + d(v_2) + ... + d(v_n)$. But, according to the Handshaking Lema in graph theory, the sum of the degrees of the n vertices of any graph is equal to two times the number of edges m, that is:

$$\sum_{i=1}^{n} d(v_i) = 2m \tag{11}$$

- The update of $\lambda(v)$ is O(1) (constant time).
- The primitive Descrease_Key is $O(\log n)$, as Q is a binary heap and we have to search for v.

Given the above, the function T(n) that measures the computational complexity of Djikstra's algorithm is given by:

$$T(n) = O(n) + O(n) + (O(1) + O(\log n)) * \sum_{i=1}^{n} d(v_i)$$

$$= O(n) + O(1) * O(m) + O(m) * O(\log n)$$

$$= O(m \log n)$$
(12)

Therefore, Dijkstra's algorithm is linear in the number of edges and logarithmic in the number of vertices, which is quite effcient for a large number of vertices n. In the proposed topological k-means algorithm, Dijkstra's algorithm is executed on k-NN graphs. Hence, the total number of edges is a constant (k << n) times the number of vertices n, that is, in the k-NN graph m is approximately O(n), which is even better.

3.2 Correctness of Dijkstra's algorithm

Before we move forward, it is crucial to show that whenever we apply Dijkstra's algorithm in a weighted graph G = (V, E, w) with positive weights, starting from a

root node s, it always builds a optimum path tree rooted in s. Moreover, the lengths of the optimum paths are the geodesic distance from s to all the remaining vertices.

Theorem 4. Given G = (V, E, w), with $w : E \to R^+$, Dijkstra's algorithm terminates with $\lambda(v) = d(s, v), \forall v \in V$, where d(s, v) is the geodesic distance from s to v.

The following steps show a proof by contradiction.

- 1. Note that along the entire execution of Dijkstra's algorithm $\lambda(v) \geq d(s, v), \forall v \in V$.
- 2. Suppose u is leaving Q and it is the first vertex to have $\lambda(u) \neq d(s, u)$ when it happens (when it enters S).
- 3. Then, $u \neq s$, because otherwise we would have $\lambda(s) = 0 = d(s, s)$ (u cannot be the root of the tree).
- 4. Hence, there is a path P_{su} , because otherwise $\lambda(u) = \infty = d(s, u)$. Therefore, there is a shortest path P_{su}^* .
- 5. Right before vertex u is removed from Q, the path has $s \in S$ and $u \in V S$.
- 6. Let y be the first vertex in P_{su}^* such that $y \in V S$ and let $x \in S$ be its predecessor.
- 7. As $x \in S$, $\lambda(x) = d(s, x)$. Moreover, by the time x was added in S, the edge (x, y) was relaxed, that is:

$$\lambda(y) = \lambda(x) + w(x, y) = d(s, x) + w(x, y) = d(s, y)$$
(13)

8. But y precedes u in P_{su}^* and as the edge weights are positive, have $d(s, y) \leq d(s, u)$. Then, by combining statements (7), (8) and (1), we have:

$$\lambda(y) = d(s, y) \le d(s, u) \le \lambda(u) \tag{14}$$

- 9. But, as both u and y belong to V-S, when u is chosen to leave Q, we must have $\lambda(u) \leq \lambda(y)$.
- 10. As $\lambda(u) \leq \lambda(y)$ and $\lambda(y) \leq \lambda(u)$, it follow that $\lambda(u) = \lambda(y)$, which implies:

$$\lambda(y) = d(s, y) = d(s, u) = \lambda(u) \tag{15}$$

yeilding a contradiction to our initial hypothesis. Therefore, $\nexists u \in V$ such that $\lambda(u) \neq d(s, u)$ when it leaves Q.

4 The topological k-means algorithm

The proposed topological k-means method is a graph-based algorithm for clustering that uses a k-NN graph to approximate the data manifold in the input space. The manifold hypothesis states that, often, many high dimensional datasets observed in real world problems lie in a low dimensional latent non-linear geometric structure inside the ambient space [17, 18]. Figure 1 shows a illustration of the k-NN graph for the digits dataset, with n=1797 samples, d=64 features and composed by c=10 classes for $nn=\sqrt{n}=42$ neighbors.

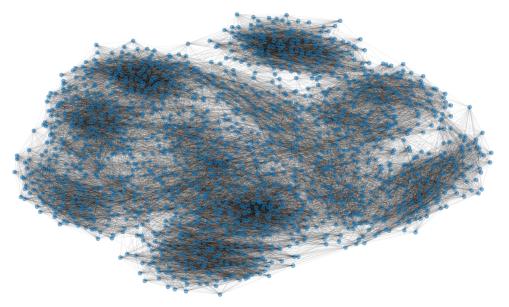


Fig. 1 k-NN graph of the digits dataset for nn = 42 neighbors.

The first question that arises in topological k-means is: how to build a graph from a multivariate dataset in order to create a discrete approximation to the data manifold? In graph-based machine learning, two popular methods are particularly interesting due to computational efficiency [19, 20]:

- 1. k-NN graph: for each data point $\vec{x}_i, i=1,2,...,n$, we must compute the Euclidean distance to every other sample $\vec{x}_j, j \neq i, j=1,2,...,n$. Then, we have to select only the k samples with smallest distance (nearest neighbors) and create an edge between them.
- 2. ϵ -neighborhood graph: we must define a radius ϵ and for each sample $\vec{x}_i, i = 1, 2, ..., n$ and every other sample $\vec{x}_j, j \neq i, j = 1, 2, ..., n$, compute the Euclidean distance between them. For every point inside the ball of radius ϵ centered at \vec{x}_i , an edge must be created.

A second question that arises is: how well the length of shortest paths in k-NN graphs can approximate the true underlying geodesic distances in the data manifold? The Asymptotic Convergence Theorem shows that, under certain regularity conditions, the length of shortest paths in k-NN graphs $d_G(\vec{x}_i, \vec{x}_j)$ converges to the geodesic distances $d_M(\vec{x}_i, \vec{x}_j)$ [21]. In summary, the authors show that the two distance metrics, $d_G(\vec{x}_i, \vec{x}_j)$ and $d_M(\vec{x}_i, \vec{x}_j)$ approximate each other arbitrarily closely, as the density of data points tends to infinity.

Theorem 5 (Asymtptic Convergence Theorem). Given $\lambda_1, \lambda_2, \mu > 0$, but as small as desired, then, for a sufficiently large density of points, the following inequality holds:

$$1 - \lambda_1 \le \frac{d_G(\vec{x}_i, \vec{x}_j)}{d_M(\vec{x}_i, \vec{x}_j)} \le 1 + \lambda_2 \tag{16}$$

with probability $1 - \mu$, where $d_G(\vec{x}_i, \vec{x}_j)$ is the recovered distance (length of shortest path) and $d_M(\vec{x}_i, \vec{x}_j)$ is the true geodesic distance in the manifold.

The mathematical details of the proof can be found in the work of Bernstein and colleagues [22]. In the following, we present the pseudo-code for the proposed topological k-means in Algorithm 3.

Algorithm 3 Topological k-means algorithm

```
// Parameters:
//X: the n \times d data matrix (each row is a sample)
// n: the number of samples
// k: the number of clusters
// nn: the number of neighbors in the k-NN graph
function Top-K-Means(X, n, k, nn)
    \vec{s}_1, \vec{s}_2, ..., \vec{s}_k \leftarrow select\_random\_seeds(X, k)

⊳ Select k random centers

    distances \leftarrow zeros(k, n)
                                                           ▶ Array to store the geodesic distances
    for i \leftarrow 1; i \leq k; i + + do
         \vec{\mu}_i \leftarrow \vec{s}_i
    end for
     while not convergence do
         G \leftarrow kNN\_Graph(X, nn)
                                                                               ▷ Build the k-NN graph
         for i \leftarrow 1; i \leq k; i + + \mathbf{do}
             \omega_i \leftarrow \{\}
                                                                    \triangleright Begin with k empty partitions
         end for
         for j \leftarrow 1; i \leq k; j + + \mathbf{do}
              distances[j] \leftarrow Dijkstra(G, \vec{\mu}_i)
                                                                                    ▶ Geodesic distances
         end for
         for i \leftarrow 1; i \leq n; i + + \mathbf{do}
             j \leftarrow distances[:,i].argmin()
                                                                         ▷ Select the nearest centroid
             \omega_j \leftarrow \omega_j \cup \{\vec{x}_i\}
                                                                       ▶ Assign sample to the cluster
         end for
         for i \leftarrow 1; i \leq k; i + + do
             \vec{\mu}_i \leftarrow \frac{1}{|\omega_i|} \sum_{\vec{x} \in \omega_i} \vec{x}
                                                                           ▶ Recalculate the centroids
         end for
     end while
     return \omega_1, \omega_2, ..., \omega_k
end function
```

At this point, some comments about the algorithm should be done. First, note that, in comparison to regular k-means, topological k-means require an additional parameter nn, which is the number of neighbors in the k-NN graph. Alternatively, one can choose to parametrize the algorithm in terms of ϵ , which is the radius that define the size of the neighborhood (ball centered in \vec{x}_i). Second, to correctly build the k-NN graph, at the end of each iteration, the new centers must be included in the data matrix X, because often they do not belong to the original data matrix. Finally, in this version of the algorithm, the center of the cluster ω_j is updated by computing the sample average of the samples assigned to it, but it is possible to compute an average over the geodesic distances, but that would require another execution of Dijkstra's algorithm inside the WHILE loop. Therefore, to reach a tradeoff between performance and computational cost, we chose to keep the sample average.

The main advantages of the proposed topological k-means over regular k-means can be summarized by two main points:

- Regular k-means uses the Euclidean distance, which means that it is "blind" for non-spherical clusters and quite sensitive to the presence of noise and outliers in data.
 With the incorporation of a graph-based geodesic distance, we intend to alleviate these limitations.
- In high dimensional spaces, the discriminating power of the Euclidan distance is poor due to the complex geometry of hyperspaces. Hence, the performance of k-means and other clustering algorithms can be severely degraded in datasets for which the number of features n is much larger than the number of samples m.

4.1 Complexity analysis

The computational complexity of the proposed topological k-means algorithm depends on the following variables: the number of samples n, the number of features d, the number of clusters k, the number of neighbors in the k-NN graph nn, the number of edges in the k-NN graph m and the number of iterations t, which is often unknown.

Note that, the definition of the initial centers can be done in O(k). The definition of the matrix of distances is performed in O(kn), and the initial FOR loop is also O(k). The analysis of the main WHILE loop reveals that:

- The k-NN graph construction can be done in $O(nd \log n)$ using a KD-Tree.
- Dijkstra's algorithm is executed k times per iteration, resulting in a total cost of $O(km \log n)$.
- The label assignment FOR loop has cost O(kn).
- The computation of the new centers has cost O(nd), as in the partition $\omega_1 + \omega_2 + \ldots + \omega_k = n$.

As the main WHILE loop is executed an unknown number of iterations t, the total cost of topological k-means is:

$$T(n) = O(kn) + O(tnd \log n) + O(tkm \log n) + O(tkn) + O(tnd)$$

$$= O(tnd \log n) + O(tkm \log n) = O(t(nd + km)\log n)$$
(17)

Therefore, topological k-means depends linearly and logarithmically in the number of samples n and linearly in the number of edges of the k-NN graph m, which in terms of computational complexity is considered a quite efficient algorithm.

5 Experiments and results

In order to test and evaluate the performance of the proposed method against regular k-means and HDBSCAN, a state-of-the-art hierarquical density based method for clustering that assumes the observed data is noisy. Three different cluster evaluation metrics (external indices) were selected to assess the performance of the algorithms: Rand index [23, 24], normalized mutual information score [25] and V-measure [26].

The first set of experiments was designed to compare the performance of the methods in real high dimensional datasets. All the selected datasets are freely available at the repository www.openml.org. Table 1 describes each dataset with names, number of samples n, number of features m and number of classes k. Note that the majority of the datasets contains microarray data from the project GEMLeR. This project provides a collection of gene expression datasets that can be used for benchmarking gene expression oriented machine learning algorithms. They can be used for estimation of different quality metrics (e.g. accuracy, precision, area under ROC curve, etc.) for classification, feature selection or clustering algorithms. Each gene expression sample in GEMLeR repository comes from a large publicly available expO (Expression Project For Oncology) repository by International Genomics Consortium [27].

Table 1 Number of samples, features and classes of the selected openML datasets for the first set of experiments.

Number	Datasets	# samples	# features	# classes
1	AP_Colon_Kidney	546	10935	2
2	AP_Breast_Kidney	604	10935	2
3	AP_Breast_Colon	630	10935	2
4	AP_Colon_Prostate	355	10935	2
5	AP_Prostate_Kidney	329	10935	2
6	AP_Uterus_Kidney	384	10935	2
7	$AP_Prostate_Uterus$	193	10935	2
8	AP_Lung_Uterus	250	10935	2
9	AP_Ovary_Kidney	458	10935	2
10	AP_Lung_Kidney	386	10935	2
11	${ m tr} 11.{ m wc}$	414	6429	9
12	${ m tr} 12.{ m wc}$	313	5804	8
13	${ m tr}45.{ m wc}$	690	8261	10
14	${ m tr}23.{ m wc}$	204	5832	6
15	${ m tr}31.{ m wc}$	927	10128	7
16	SRBCT	83	2308	4
17	pasture	36	22	2
18	$analcatdata_authorship$	841	70	2
19	stock	950	9	2
20	kc1-binary	145	94	2

It is possible to note that a common feature to most datasets is that the number of features m is much larger than the number of samples n. For example, in dataset number 7, AP_Prostate_Uterus, the number of features is more than 56 times the number of samples, a situation that represents a tough challenge for clustering algorithms.

The quantitative results for the clustering algorithms are presented in Table 2. Looking at the results, we see that for these datasets, the proposed topological k-means outperforms regular k-means in all three metrics. In all experiments, the parameter k (number of clusters) is set equal to the number of classes and the number of neighbors nn is set to $\lfloor \sqrt{n} \rfloor$. To generate the results, we computed the average performance after 30 random initializations of k-means and topological k-means. It is interesting to note that, using the standard parameter definitions, the NMI's and V-measures of HDBSCAN in the majority of these datasets are zero, because the algorithm classify all data points as noise. A possible explanation for this behavior is the sparsity of the input feature space. To check if the differences are significant, we performed a non-parametric Wilcoxon test. According to test, for a significance level $\alpha=0.05$, there are strong evidences that the proposed topological k-means performed better than regular k-means in terms of Rand index (p<0.001), normalized mutual information (p<0.001) and V-measure (p<0.001).

Table 2 Average cluster evaluation metrics obtained after 30 executions of regular k-means and topological k-means for twenty real high dimensional datasets from the openML repository.

	Regular k-means			Topological k-means		
Datasets	Rand	NMI	V-meas.	Rand	NMI	V-meas.
AP_Colon_Kidney	0.5715	0.0894	0.1296	0.6970	0.2412	0.3528
AP_Breast_Kidney	0.5131	0.0116	0.0170	0.5855	0.1079	0.1620
AP_Breast_Colon	0.5085	0.0080	0.0118	0.5341	0.0428	0.0666
AP_Colon_Prostate	0.5074	0.0248	0.0420	0.6705	0.1585	0.3075
AP_Prostate_Kidney	0.5908	0.1098	0.2002	0.6682	0.1677	0.3134
$\mathbf{AP}_{-}\mathbf{Uterus}_{-}\mathbf{Kidney}$	0.5102	0.0081	0.0122	0.6346	0.1382	0.2240
$AP_Prostate_Uterus$	0.5350	0.0461	0.0701	0.6578	0.2022	0.3096
$\mathbf{AP}_{-}\mathbf{Lung}_{-}\mathbf{Uterus}$	0.4982	0.0002	0.0003	0.5500	0.0677	0.1043
AP_Ovary_Kidney	0.5028	0.0007	0.0016	0.5826	0.0944	0.1463
AP_Lung_Kidney	0.5761	0.0968	0.1466	0.6356	0.1369	0.2234
${ m tr} 11.{ m wc}$	0.3571	0.1279	0.0992	0.5620	0.3875	0.2466
${ m tr} 12.{ m wc}$	0.3046	0.0994	0.0799	0.4602	0.2757	0.1895
${ m tr}45.{ m wc}$	0.3947	0.2498	0.1689	0.5305	0.3617	0.2206
${ m tr} {f 23.wc}$	0.4331	0.1524	0.1339	0.5313	0.2076	0.1609
${ m tr}31.{ m wc}$	0.3650	0.0899	0.0897	0.4991	0.2064	0.1617
\mathbf{SRBCT}	0.5886	0.2084	0.1702	0.6193	0.3303	0.2613
pasture	0.5254	0.1557	0.2487	0.6374	0.1941	0.3011
analcatdata_authorship	0.5469	0.1146	0.1780	$\boldsymbol{0.6002}$	0.1627	0.2475
stock	0.5360	0.0366	0.0530	0.5639	0.0947	0.1469
kc1-binary	0.5330	0.0326	0.0717	0.5992	0.1133	0.1723
Average	0.4949	0.0831	0.0962	0.5910	0.1846	0.2159
\mathbf{Median}	0.5117	0.0897	0.0848	0.5924	0.1652	0.2220
Maximum	0.5908	0.2498	0.2487	0.6970	0.3875	0.3528

The second set of experiments was conducted in a different list of datasets and the main goal is to compare the performances of topological k-means and HDBSCAN (with the standard parameters) in datasets with a large number of samples and classes. Table 3 describes each dataset with names, number of samples n, number of features m and number of classes k. The objective here is to analyze the performance of the proposed method in datasets with a large number of high density clusters. The parameters settings for topological k-means were the same as before. For HDBSCAN, as the algorithm has a large number of parameters, we used the default configuration, that is, we used the standard parameters defined in Python's scikit-learn library, with exception of the min_cluster_size parameter that was set to 10. The intuiton behind this choice is to defined that a valid cluster must have at least ten data points. The optimization of HDBSCAN parameters often requires metaheuristics methods [28], such as genetic algorithms and particle swarm optimization, leading to a significant increase in the overall computational cost.

Table 3 Number of samples, features and classes of the selected openML datasets for the second set of experiments.

Number	Datasets	# samples	# features	# classes
1	digits	1797	64	10
2	mfeat-fourier	2000	76	10
3	letter	20000	16	26
4	mfeat-karhunen	355	64	10
5	mfeat-factors	329	10935	10
6	optdigits	5620	64	10
7	mfeat-zernike	2000	47	10
8	abalone	4177	8	28
9	cnae-9	1080	856	9
10	satimage	6430	36	6
11	semeion	1593	256	10
12	vehicle	846	18	4
13	micro-mass	360	1300	10
14	har	10299	561	6
15	JapaneseVowels	9961	14	9
16	waveform- 5000	5000	40	3
17	texture	5000	40	11
18	mnist_784 (25%)	17500	784	10
19	audiology	226	69	24
20	yeast	1484	8	10

The quantitative results for the second set of experiments are presented in Table 4. The results indicate that, in general, the proposed topological k-means outperforms HDBSCAN in the selected datasets. To check if the differences in performance are statistically significant, we performed a non-parametric Wilcoxon test. Once again, according to test, for a significance level $\alpha=0.05$, there are strong evidences that the proposed topological k-means performed better than HDBSCAN with standard parameter settings in these datasets in terms of Rand index (p<0.001), normalized mutual information (p<0.001) and V-measure (p<0.001).

Table 4 Average cluster evaluation metrics obtained after 30 executions of topological k-means and HDBSCAN for twenty datasets from the openML repository.

	HDBSCAN			Topological Kmeans		
Datasets	Rand	NMI	V-meas.	Rand	NMI	V-meas.
digits	0.7890	1.3493	0.6415	0.8941	1.3662	0.6072
mfeat-fourier	0.7789	1.0675	0.5454	0.8682	1.1639	0.5226
letter	0.7001	1.3258	0.4467	0.9224	1.0481	0.3314
mfeat-karhunen	0.7637	1.0770	0.5399	0.8880	1.3092	0.5816
mfeat-factors	0.7178	0.9666	0.4949	0.8744	1.1990	0.5364
${f optdigits}$	0.8027	1.1547	0.5716	0.8920	1.3177	0.5849
yeast	0.2539	0.0508	0.0550	0.7171	0.3515	0.1873
mfeat-zernike	0.6347	0.8764	0.4613	0.8574	0.9607	0.4288
abalone	0.6450	0.1622	0.0903	0.8566	0.4462	0.1605
cnae-9	0.2060	0.0546	0.0441	0.6842	0.6045	0.3233
${f satimage}$	0.4038	0.2931	0.2651	0.8007	0.8147	0.4847
semeion	0.5023	0.6471	0.3815	0.8525	0.9288	0.4137
vehicle	0.2864	0.0380	0.0485	0.6252	0.1977	0.1514
micro-mass	0.5954	0.5311	0.3278	0.8295	1.2393	0.5813
har	0.7129	0.5993	0.4137	0.8011	0.9570	0.5744
JapaneseVowels	0.1919	0.0828	0.0679	0.7732	0.1245	0.0595
Waveform-5000	0.3413	0.0016	0.0027	0.625	0.2914	0.2813
texture	0.6017	0.8768	0.4995	0.8754	1.3268	0.5785
mnist_784 (25%)	0.3586	0.4261	0.2932	0.8264	0.7737	0.3516
audiology	0.5825	0.2383	0.1401	0.8341	1.0736	0.4067
Average	0.5434	0.5910	0.3165	0.8149	0.8519	0.4074
Median	0.5986	0.5652	0.3547	0.8433	0.9570	0.4213
Maximum	0.8027	1.3493	0.6415	0.9224	1.3662	0.6072

A few comments about the topological k-means should be addressed at this point. One important aspect about the proposed method is that we must ensure that the k-NN graph is connected, in order to allow the centers to move freely around the entire set of vertices. Otherwise, if the random initialization chooses all the centroids to be in a single connected component, the algorithm will never label the samples that belong to other connected components of G. There are two strategies to ensure that G is connected: 1) define the number of neighbors $nn = \lfloor \sqrt{n} \rfloor$ and while G is not connected, increment nn by one and build a novel k-NN graph; and 2) begin with a complete graph, extract the minimum spanning tree (MST) and add the MST edges to the k-NN graph built using $nn = \lfloor \sqrt{n} \rfloor$. Some authors argue that instead of the square root of the number of samples, a good estimative for the number of neighbors is the base two logarithm of the number of samples [29, 30]. However, in our experiments, $nn = log_2 n$ produced disconnected graphs for several high dimensional datasets, so we chose to consider the standard parameter configuration as $nn = \lfloor \sqrt{n} \rfloor$.

The main drawbacks of the proposed topological k-means are still related to regular k-means limitations. First, the random initialization scheme: the performance of the algorithm is directly related to the random selection of the initial centroids, something that is not under control. Better initialization strategies for k-means can be adapted to topological k-means, like k-means++ for instance [31]. Methods based in topological properties of the k-NN graph can also be employed to this purpose. Another issue

is related to the definition of the number of clusters k. In topological k-means it is possible to avoid the dependency on the parameter k by estimating its value through a community detection algorithm in the k-NN graph [32].

6 Conclusions and final remarks

Clustering algorithms are versatile tools in machine learning, contributing to various aspects of data analysis, exploration, and understanding. Their applications span across different domains, making them indispensable for researchers, data scientists, and analysts seeking to uncover patterns and structures within complex data. These algorithms play a crucial role in machine learning, particularly in unsupervised learning. Some examples of problems where clustering plays a major role are: pattern discovery and recognition, data exploration and understanding, anomaly detection, image and signal processing, dimensionality reduction, recommendation systems, document analysis, bioinformatics and in reducing labeling burden.

However, clustering high dimensional data still is a challenging task, especially in small sample size problems. Part of this complexity comes from the geometric properties of hyperspaces, that is, spaces with thousands of dimensions. High-dimensional data poses challenges due to the curse of dimensionality because traditional distance metrics, such as the Euclidean distance, may become less meaningful, negatively affecting the performance of clustering algorithms.

In this paper, we proposed a topological k-means algorithm that uses optimum path in graphs to approximate the true geodesic distances in the data manifold. In summary, the motivation for topological k-means was to improve the capacity of k-means to deal with high dimensional data. Overall, the main contributions of the proposed method are directly related to mitigate the the blindness of non-spherical clustering of regular k-means. Our results showed that topological k-means is capable of producing better clusters than regular k-means and even HDBSCAN, a state-of-the-art approach for data clustering. Despite being a viable and promissing algorithm, topological k-means is not perfect and has limitations. The most prominent is that the performance depends on a good initial choice of centroids, which most of the times is not under control.

In order to overcome topological k-means limitations, future works may include the development of better initialization strategies, such as the kmeans++ heuristic and a curvature-based method for sampling from data based on the shape operator. Another future improvement concerns the automatic determination of the number of clusters by means of community detection in the k-NN graph. Other strategies for graph construction can be adopted as a way to make the underlying cluster structure in data more evident, easing the work of topological k-means. Finally, the use of different metrics to weight the edges of the k-NN graph can be interesting to improve topological k-means capacity of learning different shapes of clusters.

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