

# The Language Flower

BNF-converter

April 8, 2011

This document was automatically generated by the *BNF-Converter*. It was generated together with the lexer, the parser, and the abstract syntax module, which guarantees that the document matches with the implementation of the language (provided no hand-hacking has taken place).

## The lexical structure of Flower

### Identifiers

Identifiers  $\langle Ident \rangle$  are unquoted strings beginning with a letter, followed by any combination of letters, digits, and the characters `_` `'`, reserved words excluded.

### Literals

String literals  $\langle String \rangle$  have the form `"x"`, where  $x$  is any sequence of any characters except `"` unless preceded by `\`.

CSci literals are recognized by the regular expression  $\langle digit \rangle \cdot \langle digit \rangle + \text{'e'}(\text{'+'} \mid \text{'-'})\langle digit \rangle +$

CFix literals are recognized by the regular expression  $\langle digit \rangle + \text{'.'}\langle digit \rangle + \mid \langle digit \rangle + \text{'.'} \mid \text{'.'}\langle digit \rangle +$

CHex literals are recognized by the regular expression `'0'x["0123456789abcdef"]+`

CDec literals are recognized by the regular expression `'0'd<digit>+`

COct literals are recognized by the regular expression `'0'o["01234567"]+`

CBin literals are recognized by the regular expression `'0'b["01"]+`

CInt literals are recognized by the regular expression  $\langle digit \rangle +$

### Reserved words and symbols

The set of reserved words is the set of terminals appearing in the grammar. Those reserved words that consist of non-letter characters are called sym-

bols, and they are treated in a different way from those that are similar to identifiers. The lexer follows rules familiar from languages like Haskell, C, and Java, including longest match and spacing conventions.

The reserved words used in Flower are the following:

```
else  end  for
fun   if   let
loop  then
```

The symbols used in Flower are the following:

```
:      =    ;;
->    (    )
*      ,
```

## Comments

Single-line comments begin with #.

Multiple-line comments are enclosed with (# and #).

## The syntactic structure of Flower

Non-terminals are enclosed between  $\langle$  and  $\rangle$ . The symbols  $::=$  (production),  $|$  (union) and  $\epsilon$  (empty rule) belong to the BNF notation. All other symbols are terminals.

$$\langle Program \rangle ::= \langle ListAbstractDeclaration \rangle$$

$$\langle AbstractDeclaration \rangle ::= \langle Quantifier \rangle \langle Declaration \rangle$$

$$| \langle Declaration \rangle$$

$$\langle ListAbstractDeclaration \rangle ::= \langle AbstractDeclaration \rangle$$

$$| \langle AbstractDeclaration \rangle \langle ListAbstractDeclaration \rangle$$

$$\langle Quantifier \rangle ::= \text{for } \langle ListBound \rangle$$

$$\langle Declaration \rangle ::= \text{let } \langle Ident \rangle : \langle Type \rangle = \langle Expr \rangle ; ;$$

$$\begin{aligned}
\langle \text{Expr1} \rangle & ::= \text{loop} \\
& \quad | \quad \text{fun } \langle \text{ListArg} \rangle \rightarrow \langle \text{Expr} \rangle \text{ end} \\
& \quad | \quad \text{if } \langle \text{Expr} \rangle \text{ then } \langle \text{Expr} \rangle \text{ else } \langle \text{Expr} \rangle \text{ end} \\
& \quad | \quad \langle \text{Ident} \rangle \\
& \quad | \quad \langle \text{Constant} \rangle \\
& \quad | \quad ( \langle \text{Expr} \rangle ) \\
\langle \text{Expr} \rangle & ::= \langle \text{Expr} \rangle \langle \text{Expr1} \rangle \\
& \quad | \quad \langle \text{Expr1} \rangle \\
\langle \text{Constant} \rangle & ::= \langle \text{CSci} \rangle \\
& \quad | \quad \langle \text{CFix} \rangle \\
& \quad | \quad \langle \text{CHex} \rangle \\
& \quad | \quad \langle \text{CDec} \rangle \\
& \quad | \quad \langle \text{COct} \rangle \\
& \quad | \quad \langle \text{CBin} \rangle \\
& \quad | \quad \langle \text{CInt} \rangle \\
& \quad | \quad \langle \text{String} \rangle \\
\langle \text{Arg} \rangle & ::= \langle \text{Ident} \rangle \\
\langle \text{ListArg} \rangle & ::= \langle \text{Arg} \rangle \\
& \quad | \quad \langle \text{Arg} \rangle \langle \text{ListArg} \rangle \\
\langle \text{Type} \rangle & ::= \langle \text{Type1} \rangle \rightarrow \langle \text{Type} \rangle \\
& \quad | \quad \langle \text{Type1} \rangle \\
\langle \text{Type1} \rangle & ::= \langle \text{Type1} \rangle \langle \text{Type2} \rangle \\
& \quad | \quad \langle \text{Type2} \rangle \\
\langle \text{Type2} \rangle & ::= \langle \text{Ident} \rangle \\
& \quad | \quad ( \langle \text{Type} \rangle ) \\
\langle \text{Kind} \rangle & ::= \langle \text{Kind1} \rangle \rightarrow \langle \text{Kind} \rangle \\
& \quad | \quad \langle \text{Kind1} \rangle \\
\langle \text{Kind1} \rangle & ::= * \\
& \quad | \quad ( \langle \text{Kind} \rangle ) \\
\langle \text{Bound} \rangle & ::= \langle \text{Ident} \rangle : \langle \text{Kind} \rangle \\
\langle \text{ListBound} \rangle & ::= \langle \text{Bound} \rangle \\
& \quad | \quad \langle \text{Bound} \rangle , \langle \text{ListBound} \rangle
\end{aligned}$$