#### Article Title

LogiKEy Workbench: Deontic Logics, Logic Combinations and Expressive Ethical and Legal Reasoning (Isabelle/HOL Dataset)

#### Authors

Christoph Benzmüller<sup>2,1</sup>, Ali Farjami<sup>1</sup>, David Fuenmayor<sup>2</sup>, Paul Meder<sup>1</sup>, Xavier Parent<sup>1</sup>, Alexander Steen<sup>1</sup>, Leendert van der Torre<sup>1,3</sup>, Valeria Zahoransky<sup>4</sup>

### **Affiliations**

<sup>1</sup>University of Luxembourg, Esch sur Alzette, Luxembourg

## Corresponding author(s)

Christoph Benzmüller (c.benzmueller@fu-berlin.de)
Xavier Parent (x.parent.xaviert@gmail.com)

#### Abstract

The LogiKEy workbench for ethical and legal reasoning is presented. This workbench simultaneously supports development, experimentation, assessment and deployment of formal logics and ethical and legal theories at different conceptual layers. More concretely, it comprises, in form of a data set (Isabelle/HOL theory files), formal encodings of multiple deontic logics, logic combinations, deontic paradoxes and normative theories in the higher-order proof assistant system Isabelle/HOL. The data was acquired through application of the LogiKEy methodology, which supports experimentation with different normative theories, in different application scenarios, and which is not tied to specific logics or logic combinations. Our workbench consolidates related research contributions of the authors and it may serve as a starting point for further studies and experiments in flexible and expressive ethical and legal reasoning. It may also support hand-on teaching of non-trivial logic formalisms in lecture courses and tutorials.

### Keywords

Trustworthy and responsible AI; Knowledge representation and reasoning; Automated theorem proving; Model finding; Normative reasoning; Normative systems; Semantical embedding; Higher-order logic

### Specifications Table

Subject	Computer Science	
Specific subject area	Artificial intelligence; Knowledge representation and reasoning; Normative reasoning	
Type of data	formal theories (.thy files) encoded in Isabelle/HOL syntax, readable (.png or .pdf) views of this data	
How data was acquired	The data was acquired through manual encoding of various deontic logics, logic combinations, examples of contrary-to-duty paradoxes, excerpts of legal texts and exemplary ethical theories utilizing the LogiKEy methodology [8], which is itself based on shallow semantical embeddings (SSEs) [1] of logics and theories in classical higher-order logic. The concrete encodings were conducted in the higher-order proof assistant system Isabelle/HOL [27]; however, they are conceptually transferable to other expressive reasoning systems.	

<sup>&</sup>lt;sup>2</sup>Freie Universität Berlin, Berlin, Germany

<sup>&</sup>lt;sup>3</sup>Zhejiang University, Hangzhou, China

<sup>&</sup>lt;sup>4</sup>University of Oxford, Oxford, UK

Data format	Raw, processed, analyzed and cleaned data. The data is provided in the syntax format of the Isabelle/HOL proof assistant, which has been used to process, analyze and verify it; the data files were also annotated by hand. Isabelle/HOL is freely available: https://isabelle.in.tum.de	
Parameters for data collection	One objective was to empirically assess the expressiveness and proof automation capabilities of Isabelle/HOL and its integrated tools in normative reasoning when utilizing the LogiKEy methodology and the SSE approach. Another objective was to provide a reusable foundation for further experiments in expressive ethical and legal reasoning.	
Description of data collection	The data were written by hand. As part of the data collection process it has been demonstrated that non-trivial, normative reasoning is supported in the provided framework. This in particular included studies of paradoxes in normative reasoning [11] and whether and how they can eventually be avoided. An integral aspect of the data collection process also has been to provide evidence for the practical normative reasoning performance of the various reasoning tools integrated with Isabelle/HOL when utilizing the LogiKEy approach. Useful comments were added to the data files. The practical performance of the logic encodings can be independently assessed by users in combination with the Isabelle/HOL system. It has also been demonstrated how deontic logics can be flexible combined with other logic formalisms within the LogiKEy approach.	
Data source location	The data is hosted on github.com.	
Data accessibility	The data is accessible via logikey.org, which redirects to the repository https://github.com/cbenzmueller/LogiKEy on github.com, where the data is hosted and maintained. The two subdirectories 2020-DataInBrief-Article and 2020-DataInBrief-Data are associated with this article; the latter contains the data set.	
C. Benzmüller, X. Parent, and L. van der Torre. Designing nor theories of ethical and legal reasoning: LogiKEy framework, met ogy, and tool support. Artificial Intelligence (to appear), 2020. Pr https://arxiv.org/abs/1903.10187.  Further related research articles include: [4, 5, 7, 14, 15, 16]		

#### Value of the Data

- The provided data can be reused, independent of the related research article(s), as a starting point for further studies and experiments in expressive ethical and legal reasoning. Moreover, it can be reused, extended and adapted to support also other various other application directions, including, e.g., the study of deontic modality and quantifiers in linguistics.
- The data collection is beneficial for research and application in a range of areas, including but not limited to: machine ethics (ethico-legal governor systems), explainable and trustworthy AI, regulatory technologies, argumentation, natural language semantics. To that end the data includes reusable SSEs of a portfolio of deontic logics, logic combinations, paradoxes in normative reasoning and ethical theories in classical higher-order logic (HOL), aka Church's type theory [3], interpretable in the Isabelle/HOL proof assistant system [27]. The data set may also be used to support the teaching of expressive, classical and non-classical logic formalisms and their combinations in lecture courses and tutorials.
- To reuse the data interested researchers, students and practitioners only need to download the provided data files, include them in their formalization projects and suitably extend or adapt them. For example, the contributed data includes a sample encoding of selected statements from the GDPR (General Data Protection Regulation) and an encoding of Gewirth's ethical argument and principle, known as the Principle of Generic Consistency (PGC) [18], in a suitable extension of higher-order deontic logic. These are two examples in the area of knowledge representation and reasoning with an emphasis on regulatory and ethical aspects. They can be reused as a starting point for the encoding and automated solution of similar ethico-legal theories.

- The data set advances the state of the art in deontic logic [17] as follows. Sixty years after Von Wright's invention of deontic logic, the question has always been how deontic logics and normative theories can be used in computer science applications. The LogiKEy workbench and associated methodology addresses this challenge; it has the potential to revolutionize the area of deontic logic itself.
- The data set is useful also for stimulation of cross-fertilization effects between different research communities including the deontic logics and normative reasoning communities, the area of higher-order logics, and the area of interactive and automated theorem proving with its various sub-communities targeting very different logic formalisms.
- The presented encodings put a particular emphasis on the modeling of (regulative) norms. We agree with, e.g., Jones and Sergot [20] that deontic logic is needed when it is necessary to make explicit, and then reason about, the distinction between what ought to be the case and what is the case. Furthermore, the adequate handling of the deontic paradoxes (like in particular Chisholm's paradox of contrary-to-duty obligation, which deals with norm violation) posed a core challenge for knowledge representation frameworks. This problem motivated the design of deontic logics (and logic combinations) more sophisticated and finergrained than the traditional ones, like modal logic. These frameworks are automatized for the first time. It is demonstrated that a computer or a machine can reason about norm violation during run-time.

### **Data Description**

The data is provided in form of Isabelle/HOL source files, which are hosted at logikey.org. The individual data files belong to different categories.

Contributed data files in category I are listed in Table 1. They provide encodings of SSEs, and associated tests, of various deontic logics in meta-logic HOL. A category I example file is displayed in Figs. 1–2; this data file contains (an extension of) the SSE of dyadic deontic logic (DDL) by Carmo and Jones [10] in HOL and studies, resp. verifies, its properties.

Contributed data files in category II are listed in Table 2. They study paradoxes and smaller examples of normative reasoning. An example is displayed in Fig. 3, which presents an analysis of Chisholm's paradox of contrary-to-duty obligation [13] in Standard Deontic Logic (SDL).

Contributed data files in category III are listed in Table 3. They provide an encoding of (excerpt of) legal and ethical theories and arguments formalized using the deontic logics as provided in category I files and further examined in the category II files.

In addition to data listed in Tables 1—3 the data set provided at logikey.org also includes the following:

- Subdirectory 2020-DataInBrief-Data/Course-Material-1 contains Isabelle/HOL data files stemming from a lecture course on deontic logic at University of Luxembourg based on [28].
- Subdirectory 2020-DataInBrief-Data/Climate-Engineering contains Isabelle/HOL data files related to the formalization and assessment [16] of selected arguments in climate engineering.
- Subdirectory 2020-DataInBrief-Data/US-Constitution-Loophole contains Isabelle/HOL data files related to a formalization and assessment [31] of Kurt Gödel's claim that the US Constitution contains a loophole for establishing a dictatorship.
- Subdirectory 2020-DataInBrief-Data/WiseMenPuzzle contains Isabelle/HOL data files related to a formalization and study [1] of the well known Wise Men Puzzle; this data set, which has been published before [2], is included here to make it better available for logikey.org users.

Further related data sets, including selected formalisations in computational metaphysics (cf. [1, 23] and the references therein), will be added to logikey.org as we think fit.

# Experimental Design, Materials, and Methods

The data was acquired through manual encodings of logics, theories and arguments in the Isabelle/HOL [27] proof assistant system. The modeling process was following the LogiKEy methodology depicted in Fig. 4. This methodology supports formalization projects in the area of ethical and legal reasoning at different layers of abstraction. This methodology is briefly explained below at hand of selected examples from our contributed data set; we address all three different layers and discuss examples.<sup>1</sup>

Layer L1 example development (files CJ\_DDL.thy and CJ\_DDL\_Tests.thy): File CJ\_DDL.thy contains the encoding (of a quantified extension) of the DDL of Carmo and Jones in HOL. This encoding of DDL in HOL is exemplary for *Layer L1* developments in LogiKEy. First, the desired object *logic* was selected (Step 1);

<sup>&</sup>lt;sup>1</sup>For a general description of the LogiKEy framework, methodology and tool support see the related research article [8].

Table 1: Category I data files — deontic logics, extensions of deontic logics and logic combinations
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File	Dependency	Reading	Description
SDL.thy	Main.thy	[9, 12]	Provides a SSE of standard deontic logic (SDL) in HOL. An unary deontic operator is defined. The D axiom is postulated and correspondence to seriality of the accessibility is proved. The added first-order and higher-order quantifiers are constant domain (possibilist notion of quantification). This is verified by proving the Barcan formula and its converse.
CJ_DDL.thy	Main.thy	[5]	Provides a SSE of a dyadic deontic logic (DDL) by Carmo and Jones [11] in HOL. Different modal operators are introduced: dyadic deontic obligation, monadic deontic operator for actual obligation, monadic deontic operator for primary obligation, and further alethic modalities.  Moreover, constant domain first-order and higher-order quantifiers are added.
$CJ_{DDL}_{Tests.thy}$	$CJ_DDL.thy$	[5]	Contains soundness and proof automation tests for the embedding of DDL in HOL given in CJ_DDL.thy. For example, the monadic modal operators $\Box$ , $\Box$ <sub>p</sub> and $\Box$ <sub>a</sub> are identified as S5, KT and KD modalities, respectively. Relevant lemmata from the original work of Carmo and Jones are automated.
E.thy	Main.thy	[7],[8, Fig.6]	Provides a SSE of a quantified extension of Aqvist's System E in HOL. The file also runs a number of reasoning tasks (validity checking, refutation, correspondence theory).
$Lewis\_DDL.thy$	Main.thy	[24]	Provides a SSE of Lewis's DDL. The file runs a number of reasoning tasks (validity checking, refutation, correspondence theory). The relationship with Åqvist's dyadic deontic operator is also studied.
IOL_out2.thy	Main.thy	[4]	Provides a SSE of a quantified extension of IO logic $(out_2)$ [25, 28]. The file also contains an analysis of a benchmark example discussed in the literature on moral luck.
IO_out2_STIT.thy	Main.thy	[4, 26]	Provides a SSE of a quantified extension of IO logic (out <sub>2</sub> ) [25, 28] and elements of STIT logics [19] in HOL. The file also contains proof automation tests and soundness checks.
$CJ_{oldsymbol{-}}DDLplus.thy$	Main.thy	[14, 15]	A modification of the SSE developed in file CJ_DDL.thy is presented; see Figs. 1 and 2. This theory provides the starting point for an extension of a higher-order variant of DDL into a two-dimensional semantics as originally presented by Kaplan for his logic of demonstratives [21, 22]. The logic extension is completed in file Extended_CJ_DDL.thy. The displayed lines in Fig. 2 show automations of various lemmata from the original paper of Carmo and Jones [10], where they were proved manually with pen and paper.
Extended_CJ_DDL.thy	$CJ_{DDLplus}$ .thy	[14, 15]	Contains a further extension and combination of the higher-order DDL encoded in file CJ_DDLplus.thy with relevant parts (for the work presented in the related research article [8]) of Kaplan's logic of demonstratives (LD) [21, 22].

```
Isabelle2019/HOL - CJ_DDLplus.thy
28 subsubsection < Set-Theoretic Conditions for DDL>
29 consts
       av::"w \Rightarrow wo" (**set of worlds that are open alternatives (aka. actual versions) of w*)
 30
31 pv::"w\(\pi\)wo" (**set of worlds that are possible alternatives (aka. potential versions) of w*)
32 ob:: "wo wo bool" (**set of propositions which are obligatory in a given context (of type wo) *)
33
34
      axiomatization where
       sem 3a: "\forall w. I(av w)" and (** av is serial: in every situation there is always an open alternative*)
 35
       sem_4a: "\forallw. av w \sqsubseteq pv w" and (** open alternatives are possible alternatives*)
       sem 4b: "\forall w. pv w w" and (** pv is reflexive: every situation is a possible alternative to itself*)
 37
       sem 5a: "\forall X. \neg (ob \ X \perp)" and (** contradictions cannot be obligatory*)
 38
       sem 5b: "\forall X Y Z. (X \sqcap Y) =_s (X \sqcap Z) \longrightarrow (ob X Y \longleftrightarrow ob X Z)" and
       sem 5c: "\forall X \ Y \ Z. \mathcal{I}(X \sqcap Y \sqcap Z) \land ob \ X \ Y \land ob \ X \ Z \longrightarrow ob \ X \ (Y \sqcap Z)" and
       sem\_5d \textbf{:} \ "\forall X \ Y \ Z. \ (Y \sqsubseteq X \wedge ob \ X \ Y \wedge X \sqsubseteq Z) \longrightarrow ob \ Z \ ((Z \sqcap (\sim \! X)) \sqcup Y)" \ \textbf{and}
       sem\_5e: "\forall X \ Y \ Z. \ Y \sqsubseteq X \land ob \ X \ Z \land \mathcal{I}(Y \sqcap Z) \longrightarrow ob \ Y \ Z"
 42
 43
 44
       lemma True nitpick[satisfy] oops (**model found: axioms are consistent*)
 45
46 subsubsection (Verifying Semantic Conditions)
       lemma sem 5b1: "ob XY \longrightarrow \text{ob } X (Y \cap X)" by (metis (no types, lifting) sem 5b)
 47
       lemma sem 5b2: "(ob X (Y \cap X) \longrightarrow ob X Y)" by (metis (no types, lifting) sem 5b)
      lemma sem 5ab: "ob XY \longrightarrow \mathcal{I}(X \cap Y)" by (metis (full_types) sem_5a sem_5b)
 49
      lemma sem_5bd1: "Y \sqsubseteq X \land ob \ X \ Y \land X \sqsubseteq Z \longrightarrow ob \ Z \ ((\sim X) \sqcup Y)" using sem_5b sem_5d by smt
      lemma sem_5bd2: "ob X Y \land X \sqsubseteq Z \longrightarrow ob Z ((Z \sqcap (\simX)) \sqcup Y)" using sem_5b sem_5d by (smt sem_5b1) lemma sem_5bd3: "ob X Y \land X \sqsubseteq Z \longrightarrow ob Z ((\simX) \sqcup Y)" by (smt sem_5bd2 sem_5b) lemma sem_5bd4: "ob X Y \land X \sqsubseteq Z \longrightarrow ob Z ((\simX) \sqcup (X \sqcap Y))" using sem_5bd3 by auto
 53
 54 lemma sem_5bcd: "(ob X Z \land ob Y Z) \longrightarrow ob (X \sqcup Y) Z" using sem_5b sem_5c sem_5d oops
       (** 5e and 5ab justify redefinition of @{text "O\langle \varphi | \sigma \rangle"} as (ob A B)*)
      lemma "ob A B \longleftrightarrow (\mathcal{I}(A \sqcap B) \land (\forall X. X \sqsubseteq A \land \mathcal{I}(X \sqcap B) \longrightarrow ob X B))" using sem_5e sem_5ab by blast
 56
 57
 58 subsection <(Shallow) Semantic Embedding of DDL>
 59
      subsubsection (Basic Propositional Logic)
 60
 61 abbreviation pand::"m\Rightarrowm\Rightarrowm" (infixr"\wedge" 51) where "\varphi \wedge \psi \equiv \lambda c w. (\varphi c w)\wedge (\psi c w)"
       abbreviation por::"m\Rightarrowm\Rightarrowm" (infixr"\lor" 50) where "\varphi \lor \psi \equiv \lambda c w. (\varphi c w)\lor(\psi c w)" abbreviation pimp::"m\Rightarrowm\Rightarrowm" (infix"\rightarrow" 49) where "\varphi \rightarrow \psi \equiv \lambda c w. (\varphi c w)\longrightarrow(\psi c w)"
       abbreviation pequ::"m\Rightarrowm\Rightarrowm" (infix"\leftrightarrow" 48) where "\varphi \leftrightarrow \psi \equiv \lambda c \text{ w. } (\varphi \text{ c w}) \longleftrightarrow (\psi \text{ c w})"
       abbreviation pnot::"m\Rightarrowm" ("¬_" [52]53) where "¬\varphi \equiv \lambda c w. ¬(\varphi c w)"
 66
      subsubsection (Modal Operators)
 67
68 abbreviation cjboxa :: "m \Rightarrow m" ("\square_a_" [52]53) where "\square_a \varphi \equiv \lambda c w. \forall v. (av w) v \longrightarrow (\varphi c v)"
69 abbreviation cjdiaa :: "m \Rightarrow m" ("\lozenge_a" [52]53) where "\lozenge_a \varphi \equiv \lambda c w. \exists v. (av w) v \land (\varphi c v)"
70 abbreviation cjboxp :: "m \Rightarrow m" ("\square_p_" [52]53) where "\square_p \varphi \equiv \lambda c w. \forall v. (pv w) v \longrightarrow (\varphi c v)"
71 abbreviation cjdiap :: "m \Rightarrow m" ("\lozenge_p" [52]53) where "\lozenge_p \varphi \equiv \lambda c w. \exists v. (pv w) v \land (\varphi c v)"
       abbreviation cjtaut :: "m" ("\top") where "\top \equiv \lambda c w. True"
 72
       abbreviation cjcontr :: "m" ("\bot") where "\bot \equiv \lambda c w. False"
 73
 74
      subsubsection (Deontic Operators)
 75
       abbreviation cjod :: "m\Rightarrowm\Rightarrowm" ("O(_|_)"54) where "O(\varphi|\sigma) \equiv \lambdac w. ob (\sigma c) (\varphi c)"
 76
       abbreviation cjoa :: "m\Rightarrowm" ("O<sub>a</sub>" [53]54) where "O<sub>a</sub>\varphi \equiv \lambda c w. (ob (av w)) (\varphi c) \wedge (\existsx. (av w) x \wedge \neg (\varphi c x))" abbreviation cjop :: "m\Rightarrowm" ("O<sub>i</sub>" [53]54) where "O<sub>i</sub>\varphi \equiv \lambda c w. (ob (pv w)) (\varphi c) \wedge (\existsx. (pv w) x \wedge \neg (\varphi c x))"
 77
 78
 79
      subsubsection < Logical Validity (Classical) >
 80
       abbreviation modvalidctx :: "m\Rightarrowc\Rightarrowbool" ("[_]M") where "[\varphi]M \equiv \lambdac. \forallw. \varphi c w" (**context-dep. modal validity*)
82 abbreviation modinvalidctx :: "m\Rightarrowc\Rightarrowbool" ("  [ ] ^{Mn} ]  where "  [ \varphi ] ^{M} \equiv \lambda c. \forall w. \neg \varphi c w " (**ctxt-dep. modal invalidity*) 
       abbreviation modvalid :: "m\Rightarrowbool" ("[ ]") where "[\varphi] \equiv \forall c. [\varphi]^M c" (**general modal validity*)
       abbreviation modinvalid :: "m\Rightarrowbool" ("[]") where "[\varphi] \equiv \forall c. [\varphi]^M c" (**general modal invalidity*)
 84
```

Figure 1: Data file CJ DDLplus.thy; in lines 29-85 the SSE of the DDL by Carmo and Jones [10] in HOL is presented

Carmo and Jones's DDL in the given case. A *semantics* (Step 2) for this object logic was sought and found in the original papers by Carmo and Jones [11]; such a mathematical description of a semantics, a neighborhood semantics in the given case, constitutes the ideal starting point for the definition of a SSE of the object logic in HOL, which in turn enables its *automation* (Step 3) with off-the-shelf reasoning tools for HOL. The automation of DDL was subsequently assessed (Step 4) with automated theorem provers and model finders integrated with Isabelle/HOL. Then, by pen and paper means on a theoretical level, the *faithfulness* (Step 4) of the embedding of DDL in HOL was studied and proved; this proof has been

```
Isabelle2019/HOL - CJ_DDLplus.thy
     85
     86 subsection < Verifying the Embedding>
          subsubsection < Avoiding Modal Collapse>
     87
           lemma "|P \rightarrow O_a P|" nitpick oops (**(actual) deontic modal collapse is countersatisfiable*)
     88
           lemma "[P \rightarrow O_i P]" nitpick oops (**(ideal) <u>deontic</u> modal collapse is <u>countersatisfiable</u>*)
     89
           \textbf{lemma} \ "[P \rightarrow \square_a P]" \ \textbf{nitpick oops} \ (**\underline{alethic} \ modal \ collapse \ is \ \underline{countersatisf}. \ (implies \ other \ necessity \ operators)*)
     91
          subsubsection (Necessitation Rule)
     92
           lemma NecDDLa: "|A| \Longrightarrow |\Box_a A|" by simp (** Valid only using classical (not LD) validity*)
     93
     94 lemma NecDDLp: "[A] \Longrightarrow [\Box_pA]" by simp (** Valid only using classical (not \overline{LD}) validity*)
     96 subsubsection «Lemmas for Semantic Conditions» (* extracted from Benzmüller et al. paper*)
           abbreviation mboxS5 :: "m\Rightarrowm" ("\squareS5 " [52]53) where "\squareS5\varphi \equiv \lambda c w. \forall v. \varphi c v"
     97
           abbreviation mdiaS5 :: "m\Rightarrowm" ("\diamondsuitS5" [52]53) where "\diamondsuitS5\varphi \equiv \lambda c w. \exists v. \varphi c v"
     98
           lemma C_2: "[O\langle A \mid B\rangle \rightarrow \diamondsuit^{S5}(B \land A)]" by (simp add: sem_5ab)
           lemma C_3: "((\diamondsuit^{55}(A \land B \land C)) \land O(B|A) \land O(C|A)) \rightarrow O((B \land C)|A)" by (simp add: sem_5c)
   100
   101 lemma C_4: "\lfloor (\Box^{S5}(A \to B) \land \diamondsuit^{S5}(A \land C) \land O(C|B) \rightarrow O(C|A) \rfloor" using sem_5e by blast
           lemma C_5: "[\squareS5(A \leftrightarrow B) \rightarrow (O(C|A) \rightarrow O(C|B))]" using C_2 sem_5e by blast lemma C_6: "[\squareS5(C \rightarrow (A \leftrightarrow B)) \rightarrow (O(A|C) \leftrightarrow O(B|C))]" by (metis sem_5b)
   102
   103
           lemma C 7: "| \mathbf{O} \langle \mathbf{B} | \mathbf{A} \rangle \rightarrow \square^{S5} \mathbf{O} \langle \mathbf{B} | \mathbf{A} \rangle |" by blast
   104
           lemma C_8: "[O(B|A) \rightarrow O(A \rightarrow B|T)]" using sem_5bd4 by presburger
   105
   106
   107 subsubsection (Verifying Axiomatic Characterisation)
   108 (**The following theorems have been taken from the original Carmo and Jones' paper.*)
           lemma CJ_3: "|\Box_p A \rightarrow \Box_a A|" by (simp add: sem_4a)
   109
           lemma CJ_4: "[\neg O(\bot|A)]" by (simp add: sem_5a)
           lemma CJ_5: "[(\mathbf{O}\langle B|A) \land \mathbf{O}\langle C|A\rangle) \rightarrow \mathbf{O}\langle B \land C|A\rangle]" nitpick oops (**countermodel found*)
   111
           lemma CJ_5_minus: "\lfloor \diamondsuit^{55}(A \land B \land C) \land (O(B|A) \land O(C|A)) \rightarrow O(B \land C|A) \rfloor" by (simp add: sem_5c)
           lemma CJ 6: "| \mathbf{O} \langle \mathbf{B} | \mathbf{A} \rangle \rightarrow \mathbf{O} \langle \mathbf{B} | \mathbf{A} \wedge \mathbf{B} \rangle |" by (smt C 2 C 4)
   113
           \textbf{lemma} \ \textbf{CJ\_7:} \ " [\textbf{A} \leftrightarrow \textbf{B}] \longrightarrow [\textbf{O} \langle \textbf{C} | \textbf{A} \rangle \leftrightarrow \textbf{O} \langle \textbf{C} | \textbf{B} \rangle] " \ \textbf{using} \ \text{sem\_5ab sem\_5e} \ \textbf{by} \ \text{blast}
   114
           lemma CJ 8: "[C \rightarrow (A \leftrightarrow B)] \longrightarrow [O(A|C) \leftrightarrow O(B|C)]" using C 6 by simp
   115
   116 lemma CJ_9a: "[\diamondsuit_p \mathbf{O}\langle B|A\rangle \to \Box_p \mathbf{O}\langle B|A\rangle]" by simp lemma CJ_9p: "[\diamondsuit_a \mathbf{O}\langle B|A\rangle \to \Box_a \mathbf{O}\langle B|A\rangle]" by simp
   118 lemma CJ 9 var a: "| \mathbf{O} \langle \mathbf{B} | \mathbf{A} \rangle \rightarrow \square_{\mathbf{a}} \mathbf{O} \langle \mathbf{B} | \mathbf{A} \rangle |" by simp
   119 lemma CJ_9_{var_b}: "[O(B|A) \rightarrow \square_p O(B|A)]" by simp
           lemma CJ_10: "[\diamondsuit_p(A \land B \land C) \land O(C|B) \rightarrow O(C|A \land B)]" by (smt C_4) lemma CJ_11a: "[(O_aA \land O_aB) \rightarrow O_a(A \land B)]" nitpick oops (** countermodel found*)
   122 lemma CJ 11a var: "|\diamondsuit_a(A \land B) \land (O_aA \land O_aB) \rightarrow O_a(A \land B)|" using sem 5c by auto
          \textbf{lemma CJ\_11p: "[(O_iA \land O_iB) \rightarrow O_i(A \land B)]" nitpick oops (** \underline{countermodel} \ found*)}
   123
           lemma CJ_11p_var: "[\diamondsuit_p(A \land B) \land (O_iA \land O_iB) \rightarrow O_i(A \land B)]" using sem_5c by auto
   124
           lemma CJ_12a: "[\Box_a A \rightarrow (\neg O_a A \land \neg O_a (\neg A))]" using sem_5ab by blast (*using C_2 by blast *)
   125
   126 lemma CJ_12p: "[\Box_p A \rightarrow (\neg O_i A \land \neg O_i (\neg A))]" using sem_5ab by blast (*using C_2 by blast*)
           lemma CJ_13a: "[□a(A ↔ B) → (OaA ↔ OaB)]" using sem_5b by metis (*using C_6 by blast *) lemma CJ_13p: "[□p(A ↔ B) → (OiA ↔ OiB)]" using sem_5b by metis (*using C_6 by blast *)
   127
   129 lemma CJ O O: "| \mathbf{O} \langle \mathbf{B} | \mathbf{A} \rangle \rightarrow \mathbf{O} \langle \mathbf{A} \rightarrow \mathbf{B} | \mathbf{T} \rangle |" using sem 5bd4 by presburger
           (**An ideal obligation which is actually possible both to fulfill and to violate entails an actual obligation.*)
           lemma CJ Oi Oa: "|(\mathbf{O}_{i}A \wedge \diamond_{a}A \wedge \diamond_{a}(\neg A)) \rightarrow \mathbf{O}_{a}A|" using sem 5e sem 4a by blast
   131
           (**Bridge relations between conditional obligations and actual/ideal obligations:*)
   132
   133 lemma CJ_14a: "[O(B|A) \land \Box_a A \land \Diamond_a B \land \Diamond_a \neg B \rightarrow O_a B]" using sem_5e by blast
 134 lemma CJ_14p: "[\mathbf{O} \backslash B | A \rangle \wedge \Box_p A \wedge \Diamond_p B \wedge \Diamond_a B \rightarrow \mathbf{O}_i B]" using som_se by blast lemma CJ_15a: "[(\mathbf{O} \backslash B | A \rangle \wedge \Diamond_a (A \wedge B) \wedge \Diamond_a (A \wedge B)] \rightarrow \mathbf{O}_a (A \rightarrow B)]" using CJ_O_O sem_5e by fastforce lemma CJ_15p: "[(\mathbf{O} \backslash B | A \rangle \wedge \Diamond_p (A \wedge B) \wedge \Diamond_p (A \wedge B)] \rightarrow \mathbf{O}_i (A \rightarrow B)]" using CJ_O_O sem_5e by fastforce
○137 end
```

Figure 2: Data file CJ\_DDLplus.thy; in lines 87–137 lemmata from Carmo and Jones's paper [10] are proved

published [5, 6]. Furthermore, *implications* of the embedding of DDL in HOL were studied (Step 5); see for example the additional theorems in file CJ\_DDL\_Tests.thy and the analysis of contrary-to-duty scenarios conducted in files Chisholm\_DDL\_Monadic.thy and Chisholm\_DDL\_Dyadic.thy. Since the DDL of Carmo and Jones has not been automated before with other systems or approaches, there are no *benchmarks* (Step 7) available that we could use to properly assess and compare the competitiveness of our solution. The publication of this data set can be seen as a first step towards the built-up and *contribution* (Step 8) of such a benchmark suite to the community; future work includes the conversion of our Isabelle/HOL encodings into TPTP THF format [30], so that they can be used as benchmark problems in the yearly CASC world-championship competitions [29].

Layer L2 example development (file GDPR\_CJ\_DDL.thy): In file GDPR\_CJ\_DDL.thy we selected (Step 1) statements from the General Data Protection Regulation (GDPR) for formalisation. The analysis (Step 2) of

Table 2: Category II data files: paradoxes and examples of normative reasoning

File	Dependency	Reading	Description
Chisholm_SDL.thy	SDL.thy	_	Contains a detailed analysis of Chisholm's contrary-to-duty paradox [13] in SDL, including an independence analysis for all wide and narrow scoping options that arise in the axiomatization of the paradox; see Fig. 3.
Chisholm_CJ_DDL_Monadic.thy	$CJ_DDL.thy$	_	Contains a study analogous to Chisholm_SDL.thy for monadic obligation in DDL.
Chisholm_CJ_DDL_Dyadic.thy	$CJ_DDL.thy$	_	Contains a study analogous to Chisholm_SDL.thy for dyadic obligation in DDL.
Chisholm_E.thy	CJ_DDL.thy	_	Contains a study analogous to Chisholm_SDL.thy for deontic logic E.
IO_Experiments	IO_out2_STIT	[26]	Contains a study of different paradoxes from the literature in IO logic (out <sub>2</sub> ); the file imports IO_out2_STIT.

Table 3: Category III data files: (excerpts of) legal and ethical theories and arguments

File	Dependency	Reading	Description
GDPR_SDL.thy	SDL.thy	[8, Fig. 7]	Contains a modeling of selected statements from the GDPR in SDL. It is confirmed by automated means that SDL cannot handle contrary-to-duty scenarios.
GDPR_CJ_DDL.thy	$CJ_DDL.thy$	_	Contains a modeling of selected statements from the GDPR in DDL. It is confirmed by automated means that the logic can handle contrary-to-duty scenarios. The problems identified in GDPR_SDL.thy, i.e., inconsistency and explosion, are avoided. The reasoner return "intuitive" answers to queries.
GDPR_E.thy	E.thy	[8, Fig. 8]	Contains a modeling of selected statements from the GDPR in logic E. It is confirmed by automated means that the logic can handle contrary-to-duty scenarios, and does not face the problems identified in GDPR_SDL.thy, i.e., inconsistency and explosion.
GewirthArgument.thy	$Extended_CJ_DDL.thy.$	[15, 14], [8, Fig. 10]	Contains a formalization and partial automation of Gewirth's supporting argument for his <i>Principle of Generic Consistency</i> . This principle constitutes, loosely speaking, an emendation of the <i>Golden Rule</i> , i.e., the principle of treating others as one's self would wish to be treated. Gewirth's argument and theory is assessed, emended (minor corrections) and verified.

these statements revealed that obligation aspects in the context of data processing needed to be addressed. Natural language phrases in the studied parts of the GDPR indeed contains occurrences of the deontic modalities. This motivated the choice of a suitable deontic logic (Step 3), such as DDL, for the formal encoding of these challenging aspects. In the given case it also became apparent that a propositional encoding would hardly suffice in practical applications, so the selected deontic logic DDL needed to be

```
Isabelle2019/HOL - Chisholm_SDL.thy
    1 theory Chisholm SDL imports SDL
                                                                             (*Christoph Benzmüller & Xavier Parent, 2019*)
   2 begin (*Unimportant*) nitpick_params [user axioms,expect=genuine,show all,format=2]
       (*** Chisholm Example ***)
        consts go::\sigma tell::\sigma kill::\sigma
        abbreviation "D1 \equiv O<go>" (*It ought to be that Jones goes to assist his neighbors.*)
         abbreviation "D2w ≡ O<go → tell>" (*It ought to be that if Jones goes, then he tells them he is coming.*)
         abbreviation "D2n \equiv go \rightarrow O<tell>"
         abbreviation "D3w ≡ O<¬go → ¬tell>" (*If Jones doesn't go, then he ought not tell them he is coming.*)
         abbreviation "D3n \equiv \neg go \rightarrow \bigcirc < \neg tell >"
  10
         abbreviation "D4 ≡ ¬go" (*Jones doesn't go. (This is encoded as a locally valid statement.)*)
  11
  12
       (**** Chisholm_A: All-wide scoping is leading to an inadequate, dependent set of the axioms.****)
  13
       lemma "|(D1 ∧ D2w ∧ D3w) → D4|" nitpick oops (*countermodel*)
  14
       lemma "|(D1 \land D2w \land D4) \rightarrow D3w|" by blast (*proof*)
   15
       lemma "[(D1 \land D3w \land D4) \rightarrow D2w]" nitpick oops (*countermodel*)
       lemma "[(D2w ∧ D3w ∧ D4)] → D1]" nitpick oops (*countermodel*)
  17
        (* Consistency *)
  18
       lemma "[(D1 ∧ D2w ∧ D3w)] ∧ [D4]<sub>1</sub>" nitpick [satisfy] oops (*Consistent? Yes*)
  19
        (* Oueries *)
  20
       lemma assumes "\lfloor (D1 \land D2w \land D3w) \rfloor \land \lfloor D4 \rfloor_1" shows "\lfloor \bigcirc < \neg \text{tell} > \rfloor_1" nitpick oops (*Should James not tell? No*) lemma assumes "\lfloor (D1 \land D2w \land D3w) \rfloor \land \lfloor D4 \rfloor_1" shows "\lfloor \bigcirc < \text{tell} > \rfloor_1" using assms by blast (*Should J. tell? Yes*) lemma assumes "\lfloor (D1 \land D2w \land D3w) \rfloor \land \lfloor D4 \rfloor_1" shows "\lfloor \bigcirc < \text{kill} > \rfloor_1" nitpick oops (*Should James kill? No*)
  21
  22
  23
  24
        (**** Chisholm_B: All-narrow scoping is leading to a inadequate, dependent set of the axioms.*)
  25
       lemma "|(D1 \land D2n \land D3n) \rightarrow D4|" nitpick oops (*countermodel*)
  26
       lemma "[(D1 ∧ D2n ∧ D4) → D3n]" nitpick oops (*countermodel*)
       lemma "[(D1 ∧ D3n ∧ D4) \rightarrow D2n]" by blast (*proof*)
lemma "[(D2n ∧ D3n ∧ D4) \rightarrow D1]" nitpick oops (*countermodel*)
  28
  29
        (* Consistency *)
  30
       lemma "[(D1 ∧ D2n ∧ D3n)] ∧ [D4]<sub>1</sub>" nitpick [satisfy] oops (*Consistent? Yes*)
  31
        (* Oueries *)
  32
       lemma assumes "|(D1 \land D2n \land D3n)| \land |D4|_1" shows "|O<\neg tell>_1" using assms by smt (*Should J. not tell? Yes*)
  33
       \textbf{lemma assumes} \ "\lfloor (D1 \land D2n \land D3n) \rfloor \land \lfloor D4 \rfloor_1 " \ \textbf{shows} \ "\lfloor \bigcirc <\textbf{tell} > \rfloor_1 " \ \textbf{nitpick pops} \ (*Should James tell? No*) \}
        lemma assumes "|(D1 ∧ D2n ∧ D3n)| ∧ |D4|<sub>1</sub>" shows "|O<kill>|<sub>1</sub>" nitpick oops (*Should James kill? No*)
  35
  36
        (*** Chisholm C: Wide-narrow scoping is leading to an adequate, independence of the axioms.*)
  37
       lemma "[(D1 ∧ D2w ∧ D3n) \rightarrow D4]" nitpick oops (*countermodel*) lemma "[(D1 ∧ D2w ∧ D4) \rightarrow D3n]" nitpick oops (*countermodel*)
  38
  39
        lemma "[(D1 ∧ D3n ∧ D4) → D2w]" nitpick oops (*countermodel*)
  40
       lemma "[(D2w \land D3n \land D4) \rightarrow D1]" nitpick oops (*countermodel*)
        (* Consistency *)
  42
       lemma "\lfloor (D1 \land D2w \land D3n) \rfloor \land \lfloor D4 \rfloor_1" nitpick [satisfy] oops (*Consistent? No*)
  43
  44
        \begin{array}{l} \textbf{lemma assumes} \ "[(D1 \land D2w \land D3n)] \land [D4]_1" \ \textbf{shows} \ "[\bigcirc < \neg tell > ]_1" \ \textbf{using} \ D \ assms \ \textbf{by} \ smt \ (*\underline{Shld} \ J. \ not \ tell? \ Yes*) \\ \textbf{lemma assumes} \ "[(D1 \land D2w \land D3n)] \land [D4]_1" \ \textbf{shows} \ "[\bigcirc < tell > ]_1" \ \textbf{using} \ assms \ \textbf{by} \ blast \ (*\underline{Should} \ J. \ tell? \ Yes*) \\ \textbf{lemma assumes} \ "[(D1 \land D2w \land D3n)] \land [D4]_1" \ \textbf{shows} \ "[\bigcirc < tell > ]_1" \ \textbf{using} \ D \ assms \ \textbf{by} \ blast \ (*\underline{Should} \ J. \ tell? \ Yes*) \\ \end{array} 
  45
  46
  47
  48
  49
        (*** Chisholm D: Narrow-wide scoping is leading to a inadequate, dependent set of the axioms.*)
        lemma "[(D1 \land D2n \land D3w) \rightarrow D4]" nitpick oops (*countermodel*)
  50
       lemma "|(D1 \land D2n \land D4) \rightarrow D3w|" by blast (*proof*)
       lemma "[(D1 ∧ D3w ∧ D4) \rightarrow D2n]" by blast (*proof*)
lemma "[(D2n ∧ D3w ∧ D4) \rightarrow D1]" nitpick oops (*countermodel*)
  52
  53
        (* Consistency *)
  54
       lemma "\lfloor (D1 \land D2n \land D3w) \rfloor \land \lfloor D4 \rfloor_1" nitpick [satisfy] oops (*Consistent? Yes*)
  56
        (* Queries *)
       \textbf{lemma assumes} \ "\lfloor (D1 \land D2n \land D3w) \rfloor \land \lfloor D4 \rfloor_1 " \ \textbf{shows} \ "\lfloor \bigcirc < \neg tell > \rfloor_1 " \ \textbf{nitpick oops} \ (*Should James not tell? No*) \}
  57
  58 lemma assumes "|(D1 \land D2n \land D3w)| \land |D4||" shows "|\bigcirc <tell>||" nitpick oops (*Should James tell? No*)
  59 lemma assumes "[(D1 \land D2n \land D3w)] \land [D4]_1" shows "[O < kill > ]_1" nitpick cops (*Should James kill? No*)
⊖60 end
```

Figure 3: Data file Chisholm SDL.thy studies Chisholm's paradox in combination with wide-narrow scoping issues

combined with, respectively extended by, a notion of quantification, which led to the addition of quantifiers to the file CJ\_DDL.thy. Subsequently the two GDPR articles were formalized (Step 4) using logical connectives as provided in the imported file CJ\_DDL.thy, and then some exploration (Step 5) and assessment studies were conducted. This included the analysis of the contrary-to-duty scenario as reported in related research articles [8, 5]. With our data set we contribute (Step 6) this work to the wider research community and enable its reuse.

Layer L3 example development: Layer L3 example developments have just started. The idea is to populate regulatory governor architectures [8] with ethical and legal theories from Layer L2, so that reasoning with the theories can be utilized to explain and control the behaviour of (autonomous) AI systems. To realize such applications it is required to select (Step 1) some ethical and/or legal theory from Layer L2, to devise and implement (or reuse) a respective governor architecture (Step 2), to populate (Step 3) this governor system with the selected ethical and/or legal theory, and to assess (Step 4) the well-functioning of this system in empirical studies.

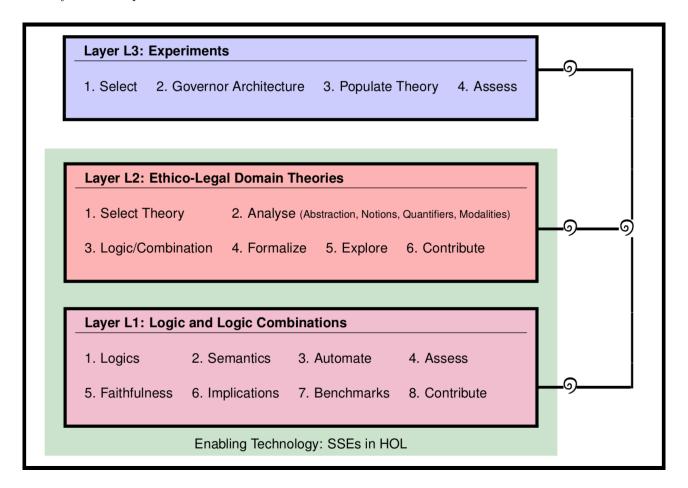


Figure 4: The LogiKEy logic and knowledge development methodology

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### Competing Interests

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