Article Title

LogiKEy Workbench: Deontic Logics, Logic Combinations and Expressive Legal and Ethical Reasoning (Isabelle/HOL Dataset)

Authors

Christoph Benzmüller^{2,1}, Ali Farjami¹, David Fuenmayor², Paul Meder¹, Xavier Parent¹, Alexander Steen¹, Leon van der Torre^{1,3}

Affiliations

¹University of Luxembourg, Esch sur Alzette, Luxembourg

Corresponding author(s)

Christoph Benzmüller (c.benzmueller@fu-berlin.de)
Xavier Parent (x.parent.xaviert@gmail.com)

Abstract

The LogiKEy workbench for legal and ethical reasoning is presented. This workbench simultaneously supports development, experimentation, assessment and deployment of formal logics and legal and ethical theories at different conceptual layers. More concretely, it comprises, in form of a data set (Isabelle/HOL theory files), formal encodings of multiple deontic logics, logic combinations, deontic paradoxes and normative theories in the higher-order proof assistent system Isabelle/HOL. The data was acquired through application of the LogiKEy methodology, which supports experimentation with different normative theories, in different application scenarios, and which is not tied to specific logics or specific logic combinations. Our workbench consolidates related research contributions of the authors and it may serve as a starting point for further studies and experiments in flexible and expressive legal and ethical reasoning. It may also support hand-on teaching of non-trivial logic formalisms in lecture courses and tutorials.

${\bf Keywords}$

Thrustworthy and responsible AI; Knowledge representation and reasoning; Automated theorem proving; Model finding; Normative reasoning; Normative systems; Philosophical and ethical issues; Semantical embedding; Higher-order logic

Specifications Table

Subject	Computer Science	
Specific subject area	Artificial intelligence; Knowledge representation and reasoning; Normative reasoning	
Type of data	.thy files (formal theories encoded in Isabelle/HOL syntax)	
How data were acquired	The data was acquired through manual encoding of various deontic logics, logic combinations, examples of contrary-to-duty paradoxes, excerpts of legal texts and exemplary ethical theories utilizing the LogiKEy methodology [7], which is itself based on shallow semantical embeddings (SSEs) [1] of logics and theories in classical higher-order logic. The concrete encodings were conducted in the higher-order proof assistant system Isabelle/HOL [22]; however, they are conceptually transferable to many other expressive reasoning systems.	
Data format	Raw, processed, analyzed and cleaned data. The data is provided in the syntax format of the Isabelle/HOL proof assistant, which has been used to process, analyze and verify it; Isabelle/HOL is freely available at https://isabelle.in.tum.de	

²Freie Universität Berlin, Berlin, Germany

³Freie Zhejiang University, Hangzhou, China

Parameters for data collection	One objective was to empirically assess the expressivity and proof automation capabilities of Isabelle/HOL and its integrated tools in normative reasoning when utilising the LogiKEy methodology and the SSE approach. Another objective was to provide a reusable foundation for further experiments in expressive legal and ethical reasoning.	
Description of data collection	The data was manually constructed and curated. As part of the data collection process it has been demonstrated that non-trivial, contrary-to-duty compliant normative reasoning is supported in the provide framework. An integral aspect of the data collection process also has been to provide evidence for the practical normative reasoning performance of the various reasoning tools integrated with Isabelle/HOL when utilizing the LogiKEy approach. Useful comments were added to the data files. The practical performance of the logic encodings can be independently assessed by users in combination with the Isabelle/HOL system. It has also been demonstrated how deontic logics can be flexible combined with other logic formalisms within the LogiKEy approach.	
Data source location	The data is hosted on github.com.	
Data accessibility	The data is accessible via https://logikey.org which redirects to the repository https://github.com/cbenzmueller/LogikEy on github.com, where the data is hosted and maintained. The two subdirectorie 2020-DataInBrief-Article and 2020-DataInBrief-Data are associated with this article; the latter contains the data set.	
Related research article	C. Benzmüller, X. Parent, and L. van der Torre. Designing norma theories of legal and ethical reasoning: LogiKEy framework, metho ogy, and tool support. Artificial Intelligence (to appear), 2020. Prepr. https://arxiv.org/abs/1903.10187. Further related research articles: [3, 4, 6, 13, 14, ?]	

Value of the Data

- The provided data can be reused, independent of the related research article(s), as a starting point for further studies and experiments in flexible and expressive legal and ethical reasoning. Moreover, it can be reused, extended and adapted to support also other various other application directions, including e.g. deontic modality and quantifiers in linguistics.
- The data collection is beneficial for research and application in a range of areas, including but not limited to: machine ethics (ethico-legal governor systems), explainably and trustworthy AI, regulatory technologies. To that end the data includes reusable SSEs of a portfolio of deontic logics, logic combinations, paradoxes in normative reasoning and ethical theories in classical higher-order logic (HOL), aka Church's type theory [2], interpretable in the Isabelle/HOL proof assistant system [22]. The data set may also be used to support the teaching of non-trivial logic formalisms in lecture courses.
- To reuse the data interested researchers, students and practitioners only need to download the provided data files, include them in their formalization projects and suitable extend or adapt them. For example, the contributed data includes a sample encoding of selected statements from the GDPR (General Data Protection Regulation) and an encoding of Gewirth's ethical argument and principle, known as the Principle of Generic Consistency (PGC), in a suitable extension of higher-order deontic logic. These are two examples in the area of knowledge representation and reasoning with an emphasis on regulatory and ethical aspects. They can be reused as a starting point for the encoding and automated solution of similar ethico-legal theories.
- Fourty years after Von Wright's invention of deontic logic, the question has always been how deontic logics and normative theories can be used in computer science applications. The LogiKEy workbench and associated methodology addresses this challenge; it has the potential to revolutionize the area of deontic logic itself.
- The data set is useful also for stimulation of cross-fertilization effects between different research communities including the deontic logics and normative reasoning communities, the area of higher-order logics,

and the area of interactive and automated theorem proving with its various sub-communities targeting very different logic formalisms.

• The presented encodings put a particular emphasis on the modeling of (regulative) norms. We agree with, e.g., Jones and Sergot [16] that deontic logic is needed when it is necessary to make explicit, and then reason about, the distinction between what ought to be the case and what is the case. Furthermore, the adequate handling of norm violation (also called contrary-to-duty situations) has posed a core challenge for knowledge representation frameworks. This problem motivated the design of deontic logics (and logic combinations) more sophisticated than traditional ones, like modal logic. These frameworks are automatized for the first time.

Data Description

The data is provided in form of Isabelle/HOL source files, which are bundled together in a single zip-file. This zip-file is comprised of the following individual files, which belong to different categories. An example file is displayed in Fig. 1; this file contains (an extension of) the SSE of dyadic deontic logic (DDL) by Carmo and Jones [10] in HOL.

Contributed data files in category I are listed in Table 1. They provide encodings of SSEs, and associated tests, of various deontic logics in meta-logic HOL.

Contributed data files in category II are listed in Table 2. They study paradoxes and smaller examples of normative reasoning.

Contributed data files in category III are listed in Table 3. They provide encodings of (excerpts of) legal and ethical theories and arguments formalized using the deontic logics as provided in category I files and further examined in the category II files.

In addition to the files listed in Tables 1—3 our data set contains further material:

• Subdirectory 2020-DataInBrief-Data/CourseMaterial1 2020-DataInBrief-Data/CourseMaterial1

There are related works in normative reasoning (see e.g. [26, 15, ?, ?]) and in computational metaphysic (see e.g. [19] and the references therein) that have employed the LogiKEy methodology, but which, for various reasons, are not included in our data set.

Experimental Design, Materials, and Methods

The data was acquired through manual encoding of the problem in the Isabelle/HOL [22] proof assistant system. The modeling process was following the LogiKEy methodology depicted in Fig. 4. This methodology supports formalization projects in the area of legal and ethical reasoning at different layers of abstraction. This methodology is explained here at hand of selected examples from our contributed data set, and we address all three different layers. For a general description of the LogiKEy framework, methodology and tool support see the related research article [7].

Layer L1 example development (files CJ-DDL.thy and CJ-DDL Tests.thy): File CJ-DDL.thy contains the encoding (of a quantified extension) of the dyadic deontic logic of Carmo and Jones in HOL. This encoding of DDL in HOL is exemplary for a Layer L1 development in the LogiKEy methodology. First, the desired object logic was selected (Step 1); DDL in the given case. A semantics (Step 2) for this object logic was sought and found in the original papers by Carmo and Jones [9, 10]; such a mathematical description of a semantics, a neighborhood semantics in the given case, constitutes the ideal starting point for the definition of a SSE of the object logic in HOL, which in turn enables its automation (Step 3) with off-the-shelf reasoning tools for HOL. The automation of DDL was subsequently assessed (Step 4) with automated theorem provers and model finders integrated with Isabelle/HOL. Then, by pen and paper means on a theoretical level, the faithfulness (Step 4) of the embedding of DDL in HOL was studied and proved; this proof has been published [4, 5]. Furthermore, implications of the embedding of DDL in HOL were studied (Step 5); see for example the additional theorems in file CJ-DDL_Tests.thy and the contrary-to-duty studies conducted with this DDL logic in files Chisholm DDL Monadic.thy and Chisholm DDL Dyadic.thy. Since the DDL of Carmo and Jones has not been automated before with other systems or approaches, there are no benchmarks (Step 7) available that we could use to properly assess and compare the competitiveness of our solution. The publication of this data set can be seen as first step towards the built-up and contribution (Step 8) of such a benchmark suite to the community; future work includes the conversion of our Isabelle/HOL encodings into TPTP THF format [25] so that they can be used as benchmarks in the yearly CASC world-championship competitions [24].

Layer L2 example development (file GDPR_CJ-DDL.thy): In file GDPR_CJ-DDL.thy we selected (Step 1) two statements from the General Data Protection Regulation (GDPR) for formalisation. The analysis (Step

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Table 1: Category I data files —	deomine rogics.	evicusions or	deomine roste	s and logic combinations

File	Dependency	Reading	Description
SDL.thy	Main.thy	[8, 11]	Provides a consistent SSE of SDL in HOL. An unary deontic operator is defined. The D axiom is postulated and correspondence to seriality of the accessibility is proved. The added first-order and higher-order quantifiers are constant domain (possibilist notion of quantification). This is verified by proving the Barcan formula and its converse.
CJ-DDL.thy	Main.thy	[4]	Provides a consistent SSE of a dyadic deontic logic (DDL) by Carmo and Jones [10] in HOL. Different modal operators are introduced: dyadic deontic obligation, monadic deontic operator for actual obligation, monadic deontic operator for primary obligation, and further alethic modalities. Moreover, constant domain first-order and higher-order quantifiers are added.
$CJ ext{-}DDL_-Tests.thy$	CJ-DDL.thy	[4]	Contains soundness and proof automation tests for the embedding of DDL in HOL given in CJ-DDL.thy. For example, the monadic modal operators \Box , \Box _p and \Box _a are identified as S5, KT and KD modalities, respectively. Relevant lemmata from the original work of Carmo and Jones [10] are automated.
E.thy	Main.thy	[6],[7, Fig.6]	Provides a consistent SSE of a quantified extension of Aqvist's System E in HOL. The file also runs a number of reasoning tasks (validity checking, refutation, correspondance theory).
Lewis-DDL.thy	Main.thy	[20]	Provides a consistent SSE of Lewis's DDL. The file also runs a number of reasoning tasks (validity checking, refutation, correspondence theory). The relationship with Åqvist's dyadic deontic operator is also studied.
IO_out2_STIT.thy	Main.thy	[3]	Provides a consistent SSE of a quantified extension of IO logic (out2) [21, 23] and elements of STIT logics [?] in HOL. The file also contains proof automation tests and soundness checks.
CJ-DDLplus.thy	Main.thy	[13, 14]	A modification of the SSE developed in file CJ-DDL.thy is presented; see Figs. 1 and 2. This theory provides the starting point for an extension of a higher-order variant of DDL into a two-dimensional semantics as originally presented by David Kaplan for his logic of demonstratives [17, 18]. The logic extension is completed in file ExtendedCJ-DDL.thy. The displayed lines in Fig. 2 show automations of various lemmata from the original paper of Carmo and Jones [10], where they were proved manually with pen and paper.
ExtendedCJ-DDL.thy	CJ-DDLplus.thy	[13, 14]	Contains a further extension and combination of the higher-order DDL encoded in file CJDDLplus.thy with relevant parts (for the work presented in the related research article [7]) of Kaplan's logic of demonstratives (LD) [17, 18].

²⁾ of these statements of the GDPR revealed that obligation aspects in the context of data processing needed to be addressed and that natural language phrases in the studied parts of the GDPR indeed contains challenge deontic modalities. This motivated the choice of a suitable deontic logic (Step 3),

```
Isabelle2019/HOL - CJDDLplus.thy
28
      subsubsection <Set-Theoretic Conditions for DDL>
29
30 consts
       av:: "w \Rightarrow wo" (**set of worlds that are open alternatives (aka. actual versions) of w*)
 31
32 pv::"w \Rightarrow wo" (**set of worlds that are possible alternatives (aka. potential versions) of w*)
33 ob:: "wo \Rightarrow wo \Rightarrow bool" (**set of propositions which are obligatory in a given context (of type wo) *)
34
35 axiomatization where
36 sem 3a: "\forall w. I(av w)" and (** av is serial: in every situation there is always an open alternative*)
       sem 4a: "∀w. av w □ pv w" and (** open alternatives are possible alternatives*)
37
       sem 4b: "∀w. pv w w" and (** pv is reflexive: every situation is a possible alternative to itself*)
       sem 5a: "\forall X. \neg(ob X \bot)" and (** contradictions cannot be obligatory*)
       sem 5b: "\forall X \ Y \ Z. \ (X \sqcap Y) =_s (X \sqcap Z) \longrightarrow (ob \ X \ Y \longleftrightarrow ob \ X \ Z)" and
       sem 5c: "\forall X \ Y \ Z. \mathcal{I}(X \sqcap Y \sqcap Z) \land ob \ X \ Y \land ob \ X \ Z \longrightarrow ob \ X \ (Y \sqcap Z)" and
       sem 5d: "\forall X \ Y \ Z. \ (Y \sqsubseteq X \land ob \ X \ Y \land X \sqsubseteq Z) \longrightarrow ob \ Z \ ((Z \sqcap (\sim X)) \sqcup Y)" and
 42
       sem\_5e \hbox{:} "\forall X \ Y \ Z. \ Y \sqsubseteq X \land ob \ X \ Z \land \mathcal{I}(Y \sqcap Z) \longrightarrow ob \ Y \ Z"
 43
44
       lemma True nitpick[satisfy] oops (**model found: axioms are consistent*)
 45
46
      subsubsection < Verifying Semantic Conditions>
 47
       lemma sem_5b1: "ob X Y -
                                                       \rightarrow ob X (Y \cap X)" by (metis (no_types, lifting) sem_5b)
       lemma sem_5b2: "(ob X (Y \cap X) \longrightarrow ob X Y)" by (metis (no_types, lifting) sem_5b)
       lemma sem_5ab: "ob X Y \longrightarrow \mathcal{I}(X \sqcap Y)" by (metis (full_types) sem_5a sem_5b)
        lemma sem 5bd1: "Y \sqsubseteq X \land ob \ X \ Y \land X \sqsubseteq Z \longrightarrow ob \ Z \ ((\sim X) \sqcup Y)" using sem 5b sem 5d by smt
 51
        lemma \ sem\_5bd2: "ob \ X \ Y \land X \sqsubseteq Z \longrightarrow ob \ Z \ ((Z \sqcap (\sim X)) \sqcup Y)" \ using \ sem\_5b \ sem\_5d \ by \ (smt \ sem\_5b1) 
       lemma sem_5bd3: "ob X Y \wedge X \sqsubseteq Z \longrightarrow ob Z ((\simX) \sqcup Y)" by (smt sem_5bd2 sem_5b)
       lemma sem 5bd4: "ob X Y \wedge X \sqsubseteq Z \longrightarrow ob Z ((\sim X) \sqcup (X \sqcap Y))" using sem 5bd3 by auto
       lemma sem_5bcd: "(ob X Z \land ob Y Z) \longrightarrow 5 (X \sqcup Y) Z" using sem_5b sem_5c sem_5d oops
       (** 5e and 5ab justify redefinition of @{text "O\langle \varphi | \sigma \rangle"} as (ob A B)*)
 56
       \textbf{lemma} \ \texttt{"ob A B} \longleftrightarrow \ (\mathcal{I}(A \sqcap B) \land (\forall X.\ X \sqsubseteq A \land \mathcal{I}(X \sqcap B) \longrightarrow \texttt{ob X B})) \texttt{"} \ \textbf{using } sem\_5e \ sem\_5ab \ \textbf{by} \ blast
 57
 58
      subsection <(Shallow) Semantic Embedding of DDL>
 59
      subsubsection < Basic Propositional Logic>
 61
       abbreviation pand::"m\Rightarrowm\Rightarrowm" (infixr"\wedge" 51) where "\varphi \wedge \psi \equiv \lambda c w. (\varphi c w)\wedge (\psi c w)"
       abbreviation por::"m\Rightarrowm\Rightarrowm" (infixr"\lor" 50) where "\varphi \lor \psi \equiv \lambda c \text{ w. } (\varphi \text{ c w}) \lor (\psi \text{ c w})"
       abbreviation pimp::"m\Rightarrowm\Rightarrowm" (infix"\rightarrow" 49) where "\varphi \rightarrow \psi \equiv \lambda c \text{ w. } (\varphi \text{ c w}) \longrightarrow (\psi \text{ c w})"
abbreviation pequ::"m\Rightarrowm\Rightarrowm" (infix"\leftrightarrow" 48) where "\varphi \leftrightarrow \psi \equiv \lambda c \text{ w. } (\varphi \text{ c w}) \longleftrightarrow (\psi \text{ c w})"
       abbreviation pnot::"m\Rightarrowm" ("¬_" [52]53) where "¬\varphi \equiv \lambdac w. ¬(\varphi \in w)"
66
 67
      subsubsection (Modal Operators)
 68
       abbreviation cjboxa :: "m\Rightarrowm" ("\square_a" [52]53) where "\square_a\varphi \equiv \lambda c w. \forall v. (av w) v \longrightarrow (\varphi c v)" abbreviation cjdiaa :: "m\Rightarrowm" ("\lozenge_a" [52]53) where "\lozenge_a\varphi \equiv \lambda c w. \exists v. (av w) v \wedge (\varphi c v)" abbreviation cjboxp :: "m\Rightarrowm" ("\square_p" [52]53) where "\square_p\varphi \equiv \lambda c w. \forall v. (pv w) v \longrightarrow (\varphi c v)" abbreviation cjdiap :: "m\Rightarrowm" ("\lozenge_p" [52]53) where "\lozenge_p\varphi \equiv \lambda c w. \exists v. (pv w) v \wedge (\varphi c v)"
 70
 71
       abbreviation cjtaut :: "m" ("T") where "T \equiv \lambda c w. True"
 73
       abbreviation cjcontr :: "m" ("\bot") where "\bot \equiv \lambda c w. False"
 74
 76 subsubsection (Deontic Operators)
       abbreviation cjod :: "m\Rightarrowm\Rightarrowm" ("O(_|_)"54) where "O(\varphi|\sigma) \equiv \lambdac w. ob (\sigma c) (\varphi c)"
 77
       abbreviation cjoa :: "m\Rightarrowm" ("\mathbf{O}_a" [53]54) where "\mathbf{O}_a\varphi \equiv \lambda c w. (ob (av w)) (\varphi c) \wedge (\existsx. (av w) x \wedge \neg (\varphi c x))" abbreviation cjop :: "m\Rightarrowm" ("\mathbf{O}_i" [53]54) where "\mathbf{O}_i\varphi \equiv \lambda c w. (ob (pv w)) (\varphi c) \wedge (\existsx. (pv w) x \wedge \neg (\varphi c x))"
 78
 79
 81 subsubsection < Logical Validity (Classical) >
      abbreviation modvalidctx :: "m\Rightarrowc\Rightarrowbool" ("[\_]^{M}") where "[\varphi]^{M} \equiv \lambda c. \forall w. \varphi \in w" (**context-dep. modal validity*) abbreviation modinvalidctx :: "m\Rightarrowc\Rightarrowbool" ("[\_]^{M}") where "[\varphi]^{M} \equiv \lambda c. \forall w. \neg \varphi \in w" (**ctxt-dep. modal invalidity*)
       abbreviation modvalid :: "m\Rightarrowbool" ("[_]") where "[\varphi] \equiv \forall c. [\varphi] ^{M} ^{M} c" (**general modal validity*) abbreviation modinvalid :: "m\Rightarrowbool" ("[_]") where "[\varphi] \equiv \forall c. [\varphi] ^{M} ^{M} ^{M} c" (**general modal invalidity*)
```

Figure 1: Data file CJDDLplus.thy; in lines 29-85 the SSE of DDL in HOL is presented

such as DDL, for the formal encoding of these challenges aspects. In the given case it became apparent that a propositional encoding would hardly suffice in practical applications, so the selected deontic logic DDL, needed to be *combined with*, respectively extended by, a notion of quantification, which led to the addition of quantifiers to the file CJ-DDL.thy. Subsequently the two GDPR articles were *formalized* (Step 4) using logical connectives as provided in the imported file CJ-DDL.thy, and then some *exploration* (Step 5) and assessment studies were conducted. This included the contrary-to-duty studies as reported in related research articles [7, 4]. With our data set we *contribute* (Step 6) this work to the wider research community and enable its reuse.

```
Isabelle2019/HOL - CJDDLplus.thy
     86
           subsection (Verifying the Embedding)
     87
     88 subsubsection < Avoiding Modal Collapse>
           lemma "[P \to O_a P]" nitpick oops (**(actual) <u>deontic</u> modal collapse is <u>countersatisfiable</u>*) lemma "[P \to O_i P]" nitpick oops (**(ideal) <u>deontic</u> modal collapse is <u>countersatisfiable</u>*)
     90
     91 lemma "[P \to \square_a P]" nitpick oops (**alethic modal collapse is countersatisf. (implies other necessity operators)*)
     92
           subsubsection < Necessitation Rule >
     93
     94 lemma NecDDLa: "|A| \implies |\Box_a A|" by simp (** Valid only using classical (not LD) validity*)
            lemma NecDDLp: "[A] \Longrightarrow [\Box_p A]" by simp (** Valid only using classical (not \overline{LD}) validity*)
           subsubsection (Lemmas for Semantic Conditions) (* extracted from Benzmüller et al. paper*)
     97
    abbreviation mboxS5 :: "m\Rightarrowm" ("\squareS5_" [52]53) where "\squareS5\varphi \equiv \lambda c w. \forall v. \varphi c v" 99 abbreviation mdiaS5 :: "m\Rightarrowm" ("\lozengeS5_" [52]53) where "\lozengeS5\varphi \equiv \lambda c w. \exists v. \varphi c v"
           lemma C_2: "| \mathbf{O} \langle \mathbf{A} | \mathbf{B} \rangle \rightarrow \diamond^{S5} (\mathbf{B} \wedge \mathbf{A}) |" by (simp add: sem_5ab)
   101 lemma C_3: "[((\diamonds5(A \wedge B \wedge C)) \wedge O(B|A) \wedge O(C|A)) \rightarrow O((B \wedge C)|A)|" by (simp add: sem_5c)
            lemma C_4: "[(□S5(A → B) ∧ ♦S5(A ∧ C) ∧ \mathbf{O}\langle C|B\rangle)) → \mathbf{O}\langle C|A\rangle]" using sem_5e by blast lemma C_5: "[□S5(A ↔ B) → (\mathbf{O}\langle C|A\rangle) using C_2 sem_5e by blast
   104 lemma C_6: "\squareS5(C \rightarrow (A \leftrightarrow B)) \rightarrow (O(A|C) \leftrightarrow O(B|C))]" by (metis sem_5b)
            lemma C_7: "[\mathbf{O}\langle B|A\rangle \to \square^{S5}\mathbf{O}\langle B|A\rangle]" by blast
   106 lemma C_8: "[O(B|A) \rightarrow O(A \rightarrow B|T)]" using sem_5bd4 by presburger
   107
  108 subsubsection (Verifying Axiomatic Characterisation)
            (**The following theorems have been taken from the original Carmo and Jones' paper.*)
  109
            lemma CJ_3: "[\Box_p A \rightarrow \Box_a A]" by (simp add: sem_4a)
            lemma CJ 4: "|\neg O(\bot |A)|" by (simp add: sem 5a)
  112 lemma CJ_5: "[(\mathbf{O}\langle B|A\rangle \land \mathbf{O}\langle C|A\rangle) \rightarrow \mathbf{O}\langle B \land C|A\rangle]" nitpick oops (**countermodel found*)
           lemma CJ_5_minus: "[\diamondsuit55(A \land B \land C) \land (O\langleB|A\rangle \land O\langleC|A\rangle) \rightarrow O\langleB\landC|A\rangle]" by (simp add: sem_5c) lemma CJ_6: "[O\langleB|A\rangle \rightarrow O\langleB|A\landB\rangle]" by (smt C_2 C_4)
   113
            \textbf{lemma} \ CJ\_7 \textbf{:} \ "[A \leftrightarrow B] \longrightarrow [\textbf{O}\langle C|A\rangle \leftrightarrow \textbf{O}\langle C|B\rangle] " \textbf{using} \ sem\_5 ab \ sem\_5 e \ \textbf{by} \ blast
           lemma CJ_8: "[C \rightarrow (A \leftrightarrow B)] \longrightarrow [O\langle A|C\rangle \leftrightarrow O\langle B|C\rangle]" using C_6 by simp lemma CJ_9a: "[\diamondsuit_pO\langle B|A\rangle \rightarrow \Box_pO\langle B|A\rangle]" by simp
  116
  118 lemma CJ 9p: "| \diamondsuit_a O \langle B | A \rangle \rightarrow \square_a O \langle B | A \rangle |" by simp
   119 lemma CJ_9_{var_a}: "[O(B|A) \rightarrow \square_aO(B|A)]" by simp
            lemma CJ_9var_b: "[O(B|A) \rightarrow \Box_p O(B|A)]" by simp
            lemma CJ_10: "[\diamondsuit_p(A \land B \land C) \land O(C|B) \rightarrow O(C|A \land B)]" by (set C_4)
            lemma CJ 11a: "|(O<sub>a</sub>A ∧ O<sub>a</sub>B) → O<sub>a</sub>(A ∧ B)|" nitpick oops (** countermodel found*)
            lemma CJ_11a_var: "[ \diamondsuit_a(A \land B) \land (O_aA \land O_aB) \rightarrow O_a(A \land B)]" using sem_5c by auto
   123
            \begin{array}{l} \textbf{lemma} \ CJ\_11p: \ "[(\textbf{O}_iA \ \land \ \textbf{O}_iB) \rightarrow \textbf{O}_i(A \ \land B)] \ " \ \textbf{nitpick oops} \ (** \ \underline{\textbf{countermodel}} \ found*) \\ \textbf{lemma} \ CJ\_11p\_var: \ "[\ \diamondsuit_p(A \ \land B) \ \land \ (\textbf{O}_iA \ \land \ \textbf{O}_iB) \rightarrow \ \textbf{O}_i(A \ \land B)] \ " \ \textbf{using } \ sem\_5c \ \textbf{by} \ auto \end{array}
   126 lemma CJ_12a: "\square_a A \rightarrow (\neg O_a A \land \neg O_a (\neg A))" using sem_5ab by blast (*using C_2 by blast *)
           \begin{array}{l} \textbf{lemma} \ CJ\_12p: \ "[\ \square_p A \to (\neg \mathbf{O}_i A \ \land \ \neg \mathbf{O}_i (\neg A))]" \ \textbf{using} \ sem\_5ab \ \textbf{by} \ blast \ (*using \ C\_2 \ by \ blast*) \\ \textbf{lemma} \ CJ\_13a: \ "[\ \square_a (A \leftrightarrow B) \to (\mathbf{O}_a A \leftrightarrow \mathbf{O}_a B)]" \ \textbf{using} \ sem\_5b \ \textbf{by} \ metis \ (*using \ C\_6 \ by \ blast*) \end{array}
  129 lemma CJ 13p: "\squarep(A \leftrightarrow B) \rightarrow (O<sub>i</sub>A \leftrightarrow O<sub>i</sub>B) |" using sem 5b by metis (*using C 6 by blast *)
   130 lemma CJ_O_O: "[O(B|A) \rightarrow O(A \rightarrow B|T)]" using sem_5bd4 by presburger
            (**An ideal obligation which is actually possible both to fulfill and to violate entails an actual obligation.*)
  132 lemma CJ_Oi_Oa: "[(O_iA \land \diamondsuit_aA \land \diamondsuit_a(\neg A)) \rightarrow O_aA]" using sem 5e sem 4a by blast
            (**Bridge relations between conditional obligations and actual/ideal obligations:*)
            \begin{array}{l} \textbf{lemma} \ CJ\_14a\text{: } "\lfloor \mathbf{O} \langle B | A \rangle \ \land \ \Box_a A \ \land \ \diamondsuit_a B \ \land \ \diamondsuit_a \neg B \ \rightarrow \ \mathbf{O}_a B \rfloor " \ \textbf{using} \ sem\_5e \ \textbf{by} \ blast \\ \textbf{lemma} \ CJ\_14p\text{: } "\lfloor \mathbf{O} \langle B | A \rangle \ \land \ \Box_p A \ \land \ \diamondsuit_p B \ \land \ \diamondsuit_p \neg B \ \rightarrow \ \mathbf{O}_i B \rfloor " \ \textbf{using} \ sem\_5e \ \textbf{by} \ blast \end{array}
  136 lemma CJ 15a: "(O(B|A) \land \diamondsuit_a(A \land B) \land \diamondsuit_a(A \land B)) \rightarrow O_a(A \rightarrow B)" using CJ O O sem 5e by fastforce
  137 lemma CJ_15p: "[(\mathbf{O}\langle B|A\rangle \land \diamondsuit_p(A \land B) \land \diamondsuit_p(A \land \neg B)) \rightarrow \mathbf{O}_i(A \rightarrow B)]" using CJ_O_O sem_5e by fastforce
⊝138 end
```

Figure 2: Data file CJDDLplus.thy; in lines 87–137 lemmata from Carmo and Jones paper [10] are proved

Layer L3 example development: Layer L3 example developments have just started. The idea is to populate ethico-legal governor architectures [7] with Layer L2 legal and ethical theories, so that reasoning with the theories can be utilized to explain and control the behaviour of (autonomous) AI systems. To realize such applications it is required to select (Step 1) some legal and/or ethical theory from Layer L2, to devise and implement a respective governor architecture (Step 2), to populate (Step 3) this governor system with the selected ethico-legal theory, and to assess (Step 4) the well-functioning of this system in empirical studies.

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Table 2: Category II data files: paradoxes and examples of normative reasoning

File	Dependency	Reading	Description
${\sf Chisholm_SDL.thy}$	SDL.thy	_	Contains a detailed analysis of Chisholm's contrary-to-duty paradox [12] in SDL, including an independence analysis for all wide and narrow scoping options that arise in the axiomatization of the paradox; see Fig. 3.
Chisholm_CJ-DDL_Monadic.thy	CJ-DDL.thy	-	Contains a study analogous to Chisholm_SDL.thy for monadic obligation in DDL.
Chisholm_CJ-DDL_Dyadic.thy	CJ-DDL.thy	-	Contains a study analogous to Chisholm_SDL.thy for dyadic obligation in DDL.
Chisholm_E.thy	CJ-DDL.thy	_	Contains a study analogous to Chisholm_SDL.thy for deontic logic E.
IO_Experiments	IO_out2_STIT	-	Contains a study of different paradoxes from the literature in IO logic (out2).

Table 3: Category III data files: (excerpts of) legal and ethical theories and arguments

File	Dependency	Reading	Description
GDPR_SDL.thy	SDL.thy	[7, Fig. 7]	Contains a modeling of selected statements from the GDPR in SDL. It is demonstrated that this modeling leads to contrary-to-duty issues, i.e. inconsistency and explosion.
GDPR_CJ-DDL.thy	CJ-DDL.thy	todo	Contains a modeling of selected statements from the GDPR in DDL. It is demonstrated that this modeling is stable against the contrary-to-duty issues identified in GDPR_SDL.thy, i.e. inconsistency and explosion is avoided and inferences are supported as expected.
$GDPR_E.thy$	E.thy	[7, Fig. 8]	Contains a modeling of selected statements from the GDPR in logic E. It is demonstrated that this modeling is stable against the contrary-to-duty issues identified in GDPR_SDL.thy, i.e. inconsistency and explosion is avoided and inferences are supported as expected.
GewirthArgument.thy	ExtendedCJ-DDL.thy.	[14, 13],[7, Fig. 10]	Contains a formalization and partial automation of Gewirth's supporting argument for his <i>Principle of Generic Consistency</i> . This principle constitutes, loosely speaking, an emendation of the <i>Golden Rule</i> , i.e., the principle of treating others as one's self would wish to be treated. Gewirth's argument and theory is assessed, emended (minor corrections) and verified.

Competing Interests

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```
Chisholm_SDL ~
                                               Isabelle2019/HOL - Chisholm_SDL.thy
                                                                 (*Christoph Benzmüller & Xavier Parent, 2019*)
   1 theory Chisholm SDL imports SDL
2 begin (*Unimportant*) nitpick_params [user_axioms,expect=genuine,show_all,format=2]
   3
      (*** Chisholm Example ***)
   4
       consts go::\sigma tell::\sigma kill::\sigma
       abbreviation "D1 ≡ O<go>" (*It ought to be that Iones goes to assist his neighbors.*)
   6
       abbreviation "D2w \equiv O<go \rightarrow tell>" (*It ought to be that if Jones goes, then he tells them he is coming.*)
       abbreviation "D2n \equiv go \rightarrow O<tell>"
       abbreviation "D3w \equiv O<\neggo \rightarrow \negtell>" (*If Jones doesn't go, then he ought not tell them he is coming.*)
       abbreviation "D3n \equiv \neg go \rightarrow \bigcirc < \neg tell > 0
  10
       abbreviation "D4 \equiv \neg go" (*Jones doesn't go. (This is encoded as a locally valid statement.)*)
  11
  12
      (*** Chisholm_A: All-wide scoping is not leading to an inadequate, dependent set of the axioms.***)
  13
      lemma "[(D1 ∧ D2w ∧ D3w) → D4]" nitpick oops (*countermodel*)
  14
      lemma "[(D1 \land D2w \land D4) \rightarrow D3w|" by blast (*proof*)
  15
      lemma "[(D1 \land D3w \land D4) \rightarrow D2w]" nitpick oops (*countermodel*)
      lemma "[(D2w ∧ D3w ∧ D4) → D1]" nitpick oops (*countermodel*)
  17
      (* Consistency *
  18
  19
      lemma "[(D1 ∧ D2w ∧ D3w)] ∧ [D4]<sub>1</sub>" nitpick [satisfy] oops (*Consistent? Yes*)
  20
      (* Oueries *)
      lemma assumes "|(D1 \land D2w \land D3w)| \land |D4||_1" shows "|\bigcirc < \neg tell > |_1" nitpick oops (*Should James not tell? No*)
  21
      lemma assumes "(D1 ∧ D2w ∧ D3w) | ∧ |D4|<sub>1</sub>" shows "|O<tell>|<sub>1</sub>" using assms by blast (*Should J. tell? Yes*)
  22
      lemma assumes "[(D1 ∧ D2w ∧ D3w)] ∧ [D4] | shows "[O<kill>| | nitpick oops (*Should James kill? No*)
  23
  24
      (*** Chisholm B: All-narrow scoping is leading to a inadequate, dependent set of the axioms.*)
  25
      lemma "\lfloor (D1 \land D2n \land D3n) \rightarrow D4 \rfloor" nitpick oops (*countermodel*)
  26
      lemma "[(D1 ∧ D2n ∧ D4) \rightarrow D3n]" nitpick oops (*countermodel*)
lemma "[(D1 ∧ D3n ∧ D4) \rightarrow D2n]" by blast (*proof*)
  27
      lemma "[(D2n \land D3n \land D4) \rightarrow D1]" nitpick oops (*countermodel*)
      (* Consistency *)
  30
  31
      lemma "|(D1 ∧ D2n ∧ D3n)| ∧ |D4|<sub>1</sub>" nitpick [satisfy] oops (*Consistent? Yes*)
      (* Oueries *)
  32
      35
  36
      (*** Chisholm_C: Wide-narrow scoping is leading to an adequate, independence of the axioms.*)
  37
      lemma "[(D1 \land D2w \land D3n) \rightarrow D4]" nitpick oops (*countermodel*) lemma "[(D1 \land D2w \land D4) \rightarrow D3n]" nitpick oops (*countermodel*)
  38
  39
      lemma "[(D1 \land D3n \land D4) \rightarrow D2w]" nitpick oops (*countermodel*)
      lemma "[(D2w ∧ D3n ∧ D4) → D1]" nitpick oops (*countermodel*)
  41
      (* Consistency *)
  42
      lemma "|(D1 ∧ D2w ∧ D3n)| ∧ |D4|<sub>1</sub>" nitpick [satisfy] oops (*Consistent? No*)
  43
      (* Queries *)
       \begin{array}{l} \textbf{lemma assumes} \ " [ (D1 \ \land D2w \ \land D3n) ] \ \land \ [D4]_1" \ \textbf{shows} \ " [ \bigcirc <\neg tell >]_1" \ \textbf{using} \ D \ assms \ \textbf{by} \ smt \ (*\underline{Shld} \ J. \ not \ tell? \ Yes*) \\ \textbf{lemma assumes} \ " [ (D1 \ \land D2w \ \land D3n) ] \ \land \ [D4]_1" \ \textbf{shows} \ " [ \bigcirc <\neg tell >]_1" \ \textbf{using} \ assms \ \textbf{by} \ blast \ (*\underline{Should} \ J. \ tell? \ Yes*) \\ \end{array} 
  45
  46
      lemma assumes "[(D1 ∧ D2w ∧ D3n)] ∧ [D4]<sub>1</sub>" shows "[O<kill>]<sub>1</sub>" using D assms by blast (*Should J. kill? Yes*)
  48
      (*** Chisholm_D: Narrow-wide scoping is leading to a inadequate, dependent set of the axioms.*)
  49
  50 lemma "|(D1 ∧ D2n ∧ D3w) → D4|" nitpick oops (*countermodel*)
  51 lemma "[(D1 \land D2n \land D4) \rightarrow D3w]" by blast (*proof*)
      lemma "[(D1 \wedge D3w \wedge D4) \rightarrow D2n]" by blast (*proof*)
  52
      lemma "[(D2n \land D3w \land D4) \rightarrow D1]" nitpick oops (*countermodel*)
  53
      (* Consistency *)
      lemma "|(D1 ∧ D2n ∧ D3w)| ∧ |D4|<sub>1</sub>" nitpick [satisfy] oops (*Consistent? Yes*)
  55
      lemma assumes "|(D1 \land D2n \land D3w)| \land |D4||" shows "|\bigcirc < \neg tell > |" nitpick oops (*Should James not tell? No*)
  lemma assumes "[(D1 \land D2n \land D3w)] \land [D4]<sub>1</sub>" shows "[O<tell>]<sub>1</sub>" nitpick oops (*Should James tell? No*) lemma assumes "[(D1 \land D2n \land D3w)] \land [D4]<sub>1</sub>" shows "[O<kill>]<sub>1</sub>" nitpick oops (*Should James kill? No*)
\bigcirc 60 end
```

Figure 3: Data file Chisholm SDL.thy studies Chisholm's paradox in combination with wide-narrow scoping issues

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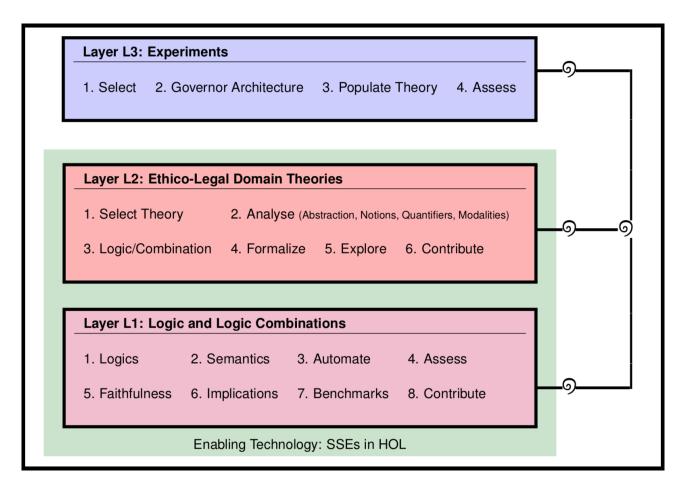


Figure 4: The LogiKEy logic and knowledge development methodology

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