# A Domain Specific Language for Usage Management

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# Outline

1 Introduction

2 Design

3 Implementation

4 Application



#### Introduction

What motivated us to do this DSL?

- Easier domain representation
- Internal v. External DSL

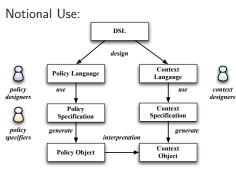
What is motivating our work?

- Applying policy-centric usage management dynamically, incorporating into network fabrics
- Providing attribution and query capabilities to policies and licensure
- Creating dynamic flexible policy environments

We think this DSL will help is in our longer term goals.



# Design — Notional Use

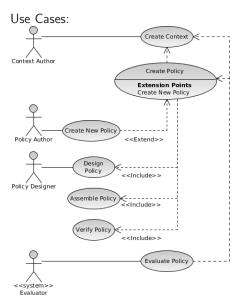


- DSL Domain specific language
- Policy Language Language elements specific to policy

- Context Language Language elements specific to context
- Policy Specification Actual specification of policy
- Context Specification Specification of context requirements
- Policy Object An object embodying policy created from the DSL
- Context Object An object containing context



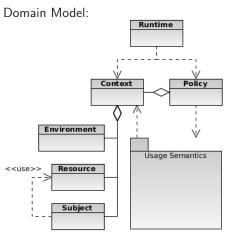
# Design — Use Cases



- Create Context Prior to creating a policy, the context in which that policy will be evaluated must be defined.
- Create Policy A designer creates a new type of policy, embodied by specific extension elements or semantic constraints over existing elements. An author will use these to create an instance of a policy.
- Evaluate Policy The policy is evaluated with a context.



# Design — Domain Model



The *Runtime* accesses and activates a *policy* and manages a *context* to which the policy is given a reference.

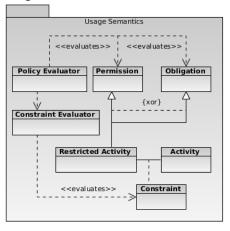
The *context* has access to information about the *environment*, *resource* managed, and the *subject* using the *resource*.

Interactions are described by specific usage semantics embodied in the policy.



# Design — Usage Semantics

#### Usage Semantics:



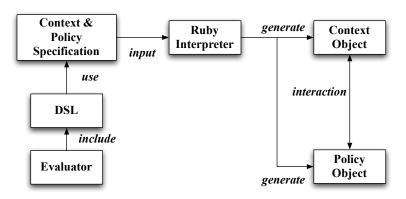
A policy evaluator examines and rectifies both permissions and Obligations.

A restricted activity is a specialization of either a permission or obligation, and is associated with a specific activity.

The association between an *activity* and a *restricted activity* is embodied by a *constraint*, which is evaluated by a *constraint evaluator*.

# Implementation — Lifecycle

### Typical DSL Lifecycle:



# Implementation — Attributes

Context			
Entity	Property (p)	Domain (D <sub>p</sub> )	Functions $(F_p)$
Environment (E)	OperatingSystem	{Windows, OSX, SELinux}	equatable
	Device	{Workstation, Handheld, Blackberry, Terminal}	equatable
	SecurityDomain	{ ABNet, SECNet, TELNet, OMNINet}	comparable
Subject (S)	SecurityClearance	{Top Secret, Secret, Confidential}	comparable
	Project	{Zebra, Yuma, Lion}	equatable
	Role	{Alpha, Beta, Delta}	equatable
Resource(R)	SecurityClassification	{ Top Secret, Secret, Confidential, Unclassified }	comparable

#### Environment (E):

Operating System  $\rightarrow$  {Windows, OSX, SELinux}  $\rightarrow$  equatable Device  $\rightarrow$  {Workstation, Handheld, Blackberry, Terminal}  $\rightarrow$  equatable Security Domain  $\rightarrow$  {ABNet, SECNet, TELNet, OMNINet}  $\rightarrow$  comparable Subject(S):

SecurityClearance  $\rightarrow$  {Top Secret, Secret, Confidential}  $\rightarrow$  comparable Project  $\rightarrow$  {Zebra, Yuma, Lion}  $\rightarrow$  equatable Role  $\rightarrow$  {Alpha, Beta, Delta}  $\rightarrow$  equatable

### Resource(S):

Classification  $\rightarrow$  {TopSecret, Secret, Confidential, Unclassified}  $\rightarrow$  comparable



# Implementation — Properties

```
property : OperatingSystem do
 values :windows, :osx, :selinux
 functions :set. :get. :equatable
end
property : device do
 values :workstation, :handheld, :blackberry, :terminal
 functions :set. :get. :equatable
end
property : project do
 values :zebra. :vuma. :lion
 functions :set, :get, :equatable
end
property : role do
 values :alpha, :beta, :delta
 functions :set, :get, :equatable
end
```



# Implementation — Properties

```
property : securitydomain do
 values :abnet, :secnet, :telnet, :omninet
 functions :set. :get. :comparable
 order :abnet. :secnet. :telnet. :omninet
end
property : securityclearance do
 values :topsecret. :secret. :confidential
 functions :set, :get, :comparable
 order :topsecret. :secret. :confidential
end
property : securityclassification do
 values :topsecret, :secret, :confidential,
    :unclassified
 functions :set, :get, :comparable
 order :topsecret. :secret. :confidential.
    ·unclassified
end
```



# Implementation — Entity, Context

```
entity :subject do
    contains :project, :role, :securityclearance
end

entity :environment do
    contains :device, :operatingsystem, :securitydomain
end

entity :resource do
    contains :securityclassification
end
```

```
context :multilevelsecurity do
contains :subject, :resource, :environment
end
```



# Implementation — Activities, Constraints

```
view = activity :view do
    # Some activity to enable viewing
end

c1 = constraint do
    securityclassification >= :secret
    && project == :yuma
    && securityclaerance >= :secret
    && device == :blackberry
    && securitydomain >= :secnet
end

restricted_view = restrict view do
    with c1
end
```

```
authorization = activity :project_authorization do is_authorized? :yuma end
```



# Implementation — Policies

```
pol = policy do
    policy_evaluators :standard
    constraint_evauators :propositional
    permit restricted_view do
    when authorization
    end
end
```

```
pol = policy do
   policy_evaluators :standard
   constraint_evauators :propositional
   permit restricted_view do
   when authorization
   count_limit restricted_view, 5
   end
end
```



# Implementation - Interface

- permissions?(). Returns the set of permissions for a given policy.
- obligations?(a). Returns the set of all obligations associated with a given permission.
- **remaining\_obligations(a)**. Returns the set of remaining obligations for a given permission.
- remaining\_count(a). Returns the set of remaining count for a given permission.
- allowed?(a, ctx). A boolean function that returns true/false whether a given activity can be carried out under a given context.
- reset(). Resets the policy by resetting its state.



# Application - CC REL

- The Creative Commons Rights Expression Language
- RDFa (Resource Description Framework in attributes) for HTML
   Web pages and resources referenced therein
- XMP (Extensible Metadata Platform) for stand-alone media
- http://wiki.creativecommons.org/CC\_REL



# Application - RDFa in HTML

Can simply associate web content with a CC license:

```
<div about="" instanceof="cc:Work" xmlns:cc="http://creativecommons.org/ns#"</pre>
     xmlns:dc="http://purl.org/dc/elements/1.1/" align="center">
    <a rel="license" href="http://creativecommons.org/licenses/by/3.0/">
    <img alt="Creative Commons License"</pre>
         src="http://i.creativecommons.org/l/by/3.0/us/88x31.png" /></a><br/>br />
    <span property="dc:title">The Lessig Blog</span>,
    a <span rel="dc:type" href="http://purl.org/dc/dcmitype/Text">
    collection of texts</span>
   by <a property="cc:attributionName" rel="cc:attributionURL"
        href="http://lessig.org/"> Lawrence Lessig </a>, <br />
    is licensed under a
    <a rel="license" href="http://creativecommons.org/licenses/by/3.0/">
    Creative Commons Attribution License</a>.<br />
    There are
   <a rel="cc:morePermissions" href="http://lessig.org/blog/other-license">
   alternative licensing options </a>
</div>
```

 But what are the semantics of http://creativecommons.org/licenses/by/3.0/?



# Application - License Deed Webpage

• Can a machine derive this page's semantics?





# Application - RDF Embedded

The previous webpage contains the following embedded RDF:

- From this, a machine can determine that this license:
  - · Permits:
    - #DerivativeWorks, #Distribution, #Reproduction
  - Requires:
    - #Attribution. #Notice
- However, what do these things mean? How are they implemented?



# Application - RDFa Embedded

- In addition to the RDF shown on previous slide, CC License Deeds also have embedded RDFa
- You can see that a machine can parse this data with something like the RDFa Distiller and Parser Tool:





# Application - RDFa Distiller

```
<?xml version="1.0" encoding="utf-8"?>
<rdf:RDF
  xmlns:cc="http://creativecommons.org/ns#"
  xmlns:dc="http://purl.org/dc/elements/1.1/"
  xmlns:dct="http://purl.org/dc/terms/"
  xmlns:rdf="http://www.w3.org/1999/02/22-rdf-svntax-ns#"
  xmlns:xml="http://www.w3.org/XML/1998/namespace"
  <rdf:Description rdf:about="http://creativecommons.org/">
    <dct:title xml:lang="en">Creative Commons</dct:title>
    <dc:title xml:lang="en">Creative Commons</dc:title>
  </rdf:Description>
  <cc:License rdf:about="http://creativecommons.org/licenses/bv/3.0/">
    <dct:creator rdf:resource="http://creativecommons.org/"/>
    <dc:creator rdf:resource="http://creativecommons.org/"/>
    <cc:requires rdf:resource="http://creativecommons.org/ns#Attribution"/>
    <cc:requires rdf:resource="http://creativecommons.org/ns#Notice"/>
    <cc:permits rdf:resource="http://creativecommons.org/ns#DerivativeWorks"/>
    <cc:permits rdf:resource="http://creativecommons.org/ns#Reproduction"/>
    <cc:permits rdf:resource="http://creativecommons.org/ns#Distribution"/>
    <dct:identifier xml:lang="en">CC BY 3.0</dct:identifier>
    <dct:title xml:lang="en">Attribution 3.0 Unported</dct:title>
    <dc:identifier xml:lang="en">CC BY 3.0</dc:identifier>
  </cc:License>
</rdf:RDF>
```

# Application - CC RDF Schema

- License RDF(a) references #DerivativeWorks, etc., in the CC namespace that's defined by a schema that's human-readable and machine-readable RDF.
- But... how immediately machine actionable is this schema?
- Partial screenshot below:



• We would like to investigate replacing or augmenting RDF(a) in the license deed with a license that's described with our DSI



# Application - DSL

- By investigating replacing the contents of a license like http://creativecommons.org/licenses/by/3.0/ with something that expresses the license in terms our DSL, we hope to:
  - Maintain equivalent license semantics
  - Express the semantics in a form that is easier for humans to read and write
  - Enable a machine to more directly execute the license and reason over it



## Conclusions

- Internal DSLs are convenient, but probably not appropriate for real systems
- Overall we like the DSL but could do without some of the Ruby cruft (e.g. **do...end**, etc.)
- Application and Optimization

