

CS 325 Analysis of Algorithms

Implementation Assignment 1: Divide and Conquer

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Pseudo Code:

The code uses an object `distObj` that has a list starting with smallest distance and the following pairs sharing that smallest distance.

Brute Force:

```
brute_force(list of points)
```

```
    point1 = first point in list
```

```
    point2 = second point in list
```

```
    min = the distance of these first pair of points
```

```
    ln = length of list
```

```
    if ln is equal to 2
```

```
        if min is less than current objects distance
```

```
            set min as the objects current smallest distance
```

```
            add pair of points to list of smallest distance points
```

```
            return object
```

```
    else
```

```
        for i in range of ln - 1
```

```
            for j in range of i+1 to ln
```

```
                dist = get distance between point i and point j
```

```
                if dist is less than distObj
```

```
                    distObj = set dist as new min distance
```

```
                    clear list of points and add new points to list
```

```
                elif dist is equal to current min distance
```

```
                    add points to list
```

```
    return distObj
```

Enhanced Divide and Conquer:

closest_pair(list sorted by x, object list of smallest distance and points)

length = length of xlist

index = 0

if length \leq 3 #base case

for point in xlist

if index + 1 is greater then length of xlist

then add the two points

increase index

return distObj

mid = middle of list

leftArrayX = split the left half

rightArrayX = split the right half

distObj1 = recursive call closest_pair(leftArrayX, distObj)

distObj2 = recursive call closest_pair(rightArrayX, distObj)

merge distObj1 and distObj2

get delta by getting distance from combined lists

min_cross = closest_cross_pair(xlist, delta) #find min distance on boundary

return list of merging distObj1 and min_cross

closest_cross_pair(list, delta)

distObj = get current smallest distance

length = length of list

mid = length of list divided by 2

if length is less than 7

strip = list

else

strip = list from range of [mid - delta, mid +delta]

```

sort strip by y coordinates
index = 0
for point in strip
    for i in range of 1 through 8
        if index + i < length of strip
            then add those points to distObj
        increase index by 1
return distObj

```

Asymptotic Analysis:

Brute Force Algorithm:

In the brute force algorithm, we have a nested for loop. Making this algorithm have a runtime of $O(n^2)$. There are constant operations as well but the nested for loop overtakes them.

Enhanced Divide and Conquer:

In the enhanced Divide and Conquer, the elements are pre-sorted when entering the function. In each recursive call we split the array in two halves. Then when we call the `closest_cross_pair` function, it runs a for loop adding points that have the minimal distance. The for loop has an $O(n)$ runtime, and the recursive call has a $O(\log n)$ runtime because each call halves the problem. So the total runtime will be $O(n \log n)$.

Plotting and Analysis:

plots and explanation go here