

VE281

Data Structures and Algorithms

Queues

Review

- Generic programming
 - Data-type independent way of programming
- Iterators for containers
- Stacks
 - Methods: **push**, **pop**, **size**, etc.
 - Implementations: arrays versus linked lists
 - Applications

Outline

- Queues

Queues

- A “line” of items in which the **first** item inserted into the queue is the **first** one out.
- FIFO access: first in, first out
- Restricted form of a linear list: insert at one end and remove from the other.

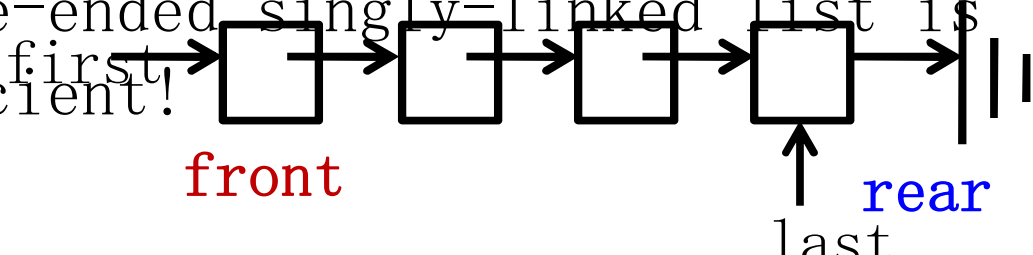


Methods of Queue

- **size()**: number of elements in the queue.
- **isEmpty()**: check if queue has no elements.
- **enqueue(Object o)**: add object **o** to the **rear** of the queue.
- **dequeue()**: remove the **front** object of the queue if not empty; otherwise, throw **queueEmpty**.
- **Object &front()**: return a reference to the front element of the queue.
- **Object &rear()**: return a reference to the rear element of the queue.

Queues Using Linked Lists

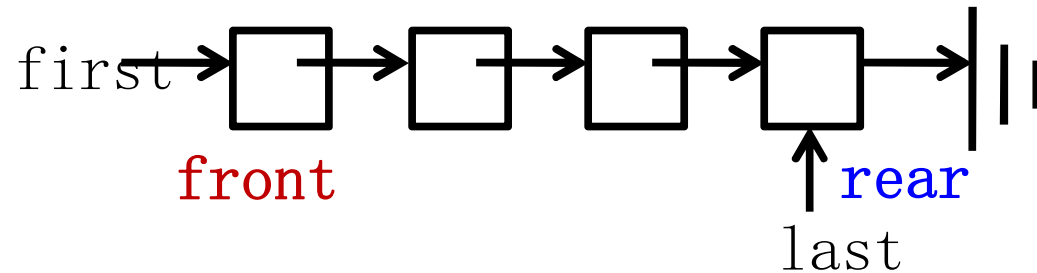
- Which type of linked list should we choose?
 - We need fast **enqueue** and **dequeue** operations.

- Double-ended singly-linked list is sufficient!
- 
- front rear last

- **enqueue** (Object o) : append object at the end

```
LinkedList::append(Object o) ;
```

Queues Using Linked Lists



- **size()**: return size;
- **isEmpty()**: return (size == 0);
- **Object &front()**: return a reference to the object stored in the first node.
- **Object &rear()**: return a reference to the object stored in the last node.

Queues Using Arrays

Array[MAXSIZE] :

2	3	1	4			
---	---	---	---	--	--	--

 front rear

- If we stick to the requirement that the n elements of a queue are the beginning n elements of the array,
 - what is the complexity of **enqueue**?
 - what is the complexity of **dequeue**?
- A better way is to let the elements “drift” within the array

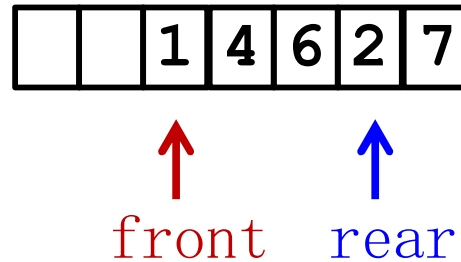
enqueue(6);

dequeue();

dequeue();

2	3	1	4	6		
---	---	---	---	---	--	--

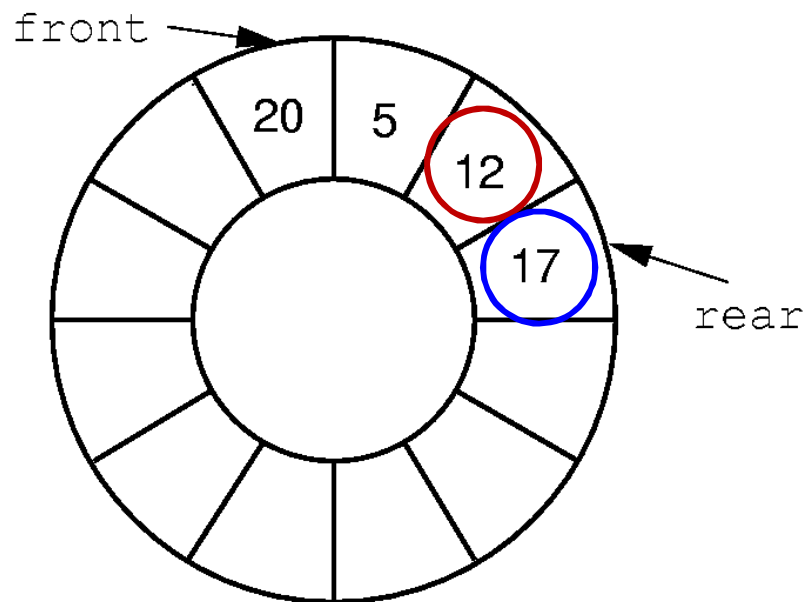
Queues Using Arrays



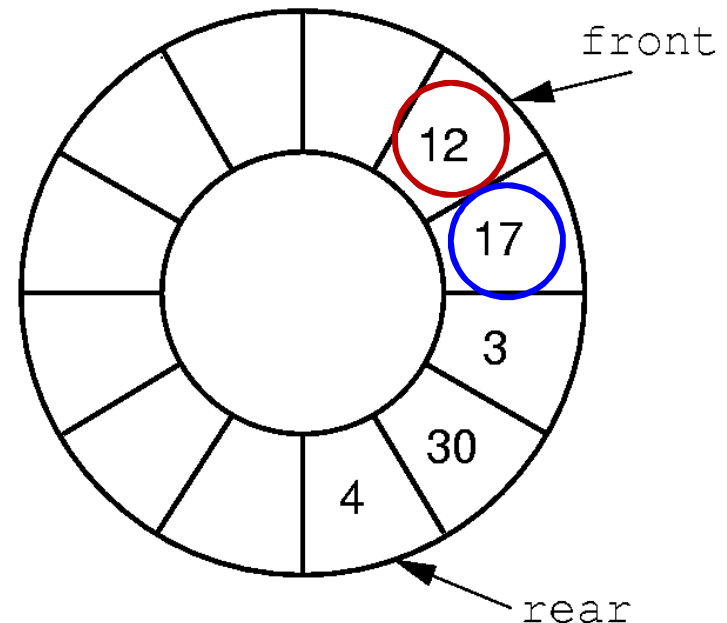
- We maintain two integers to indicate the front and the rear of the queue.
- However, as items are added and removed, the queue “drifts” toward the end.
 - Eventually, there will be no space to the right of the queue, even though there is space in the array.

Queues Using Arrays

- To solve the problem of memory waste, we use a **circular array**.



(a)



(b)

Circular Arrays

- We can implement a circular array using a plain linear array:
 - When front/rear equals the **last** index (i.e., MAXSIZE-1), increment of front/rear gives the **first** index (i.e., 0).



front

rear

enqueue (5)



rear

front

Circular Arrays

- To realize the “circular” increment, we can use modulo operation:

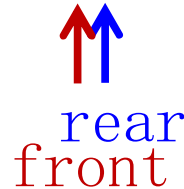
front = (front+1) % MAXSIZE;

If **front == MAXSIZE-1**, the statement sets **front** to 0.

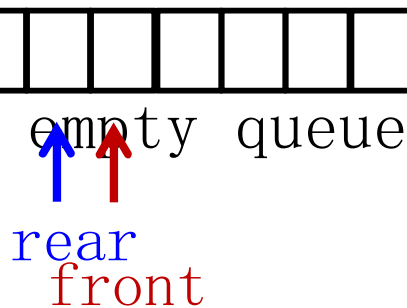
Boundary Conditions

- Suppose that **front** points to the first element in the queue and that **rear** points to the last element in the queue.

- What will a queue with one element look like?

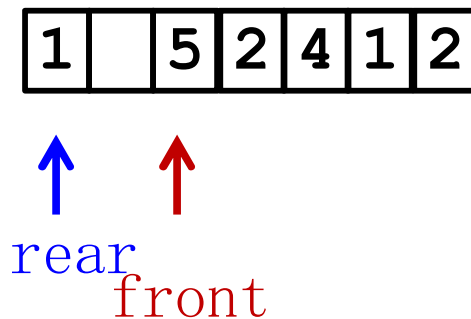


- What will an empty queue look like?

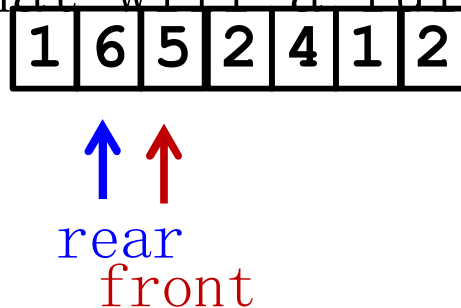


Boundary Conditions

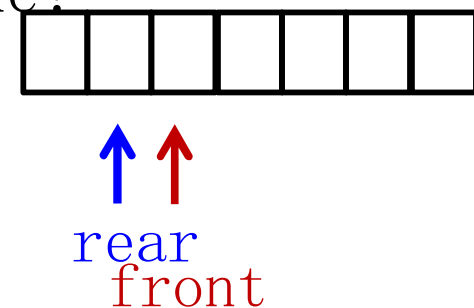
- What will a queue with one empty slot look like?



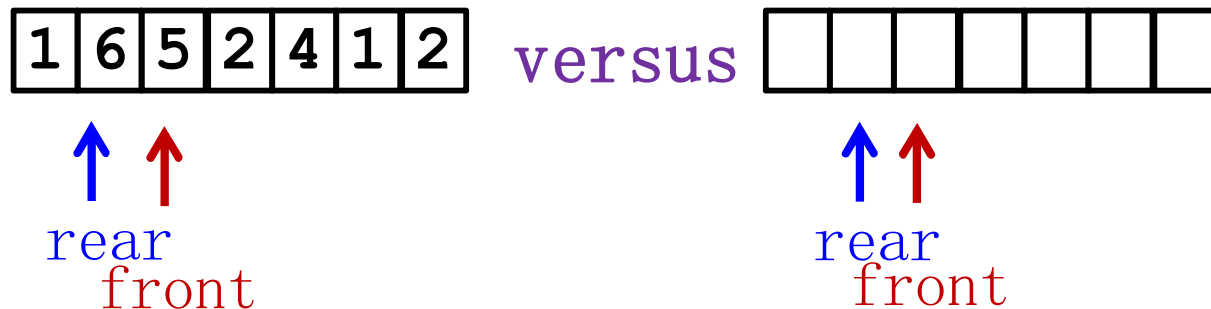
- What will a full queue look like?



versus an
empty
queue



Boundary Conditions



- To distinguish between the full array and the empty array, we need a flag indicating **empty** or **full**, or a **count** on the number of elements in the queue.

- Question: what will a queue created by the default constructor look like?
- Diagram illustrating the default constructor state for a queue. It shows an array with 7 cells.
- Array:

--	--	--	--	--	--	--
- A red arrow labeled "front" points to the first cell (index 0), and a blue arrow labeled "rear" points to the last cell (index 6).

Queues Using Arrays

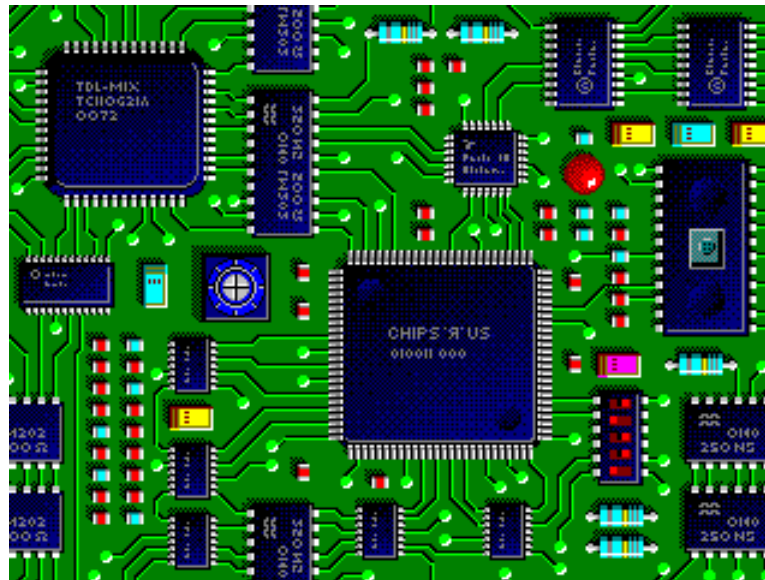
- **enqueue(Object o)**: increment **rear**, wrapping to the beginning of the array if the end of the array is reached; if **rear** becomes **front**, reallocate arrays.
- **dequeue()**: increment **front**, wrapping to the beginning of the array if the end of the array is reached; if empty, throw **queueEmpty**.
- **isEmpty()**: **return (count == 0);**
- **size()**: **return count;**

Application of Queues

- Request queue of a web server
 - Each user can send a request.
 - The arriving requests are stored in a **queue** and processed by the computer in a **first-come-first-serve** way.

Application of Queue: Wire Routing

- Select paths to connect all pairs of pins that need to be connected together.
- An important problem in **electronic design automation**.



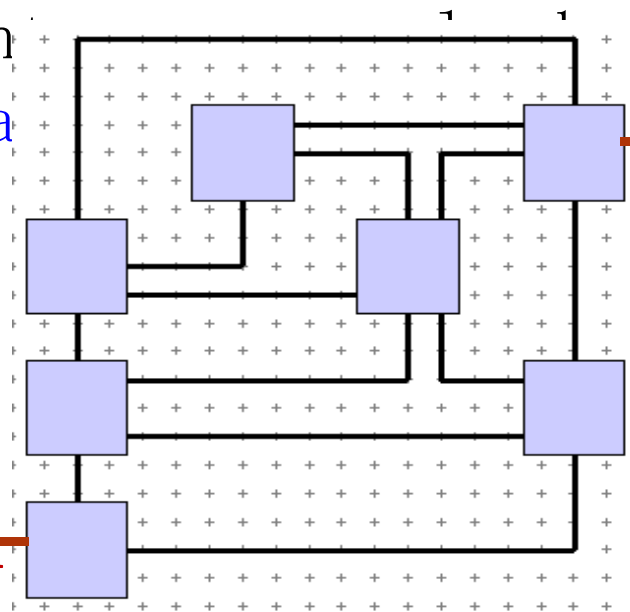
A Simplified Problem

- Condition: We have all blocks laid on the chip. We also have some of the wires routed.
- Problem: We want to connect the next pair of pins.
- Constraint: w wires

horizontal

StartPin

End Pin

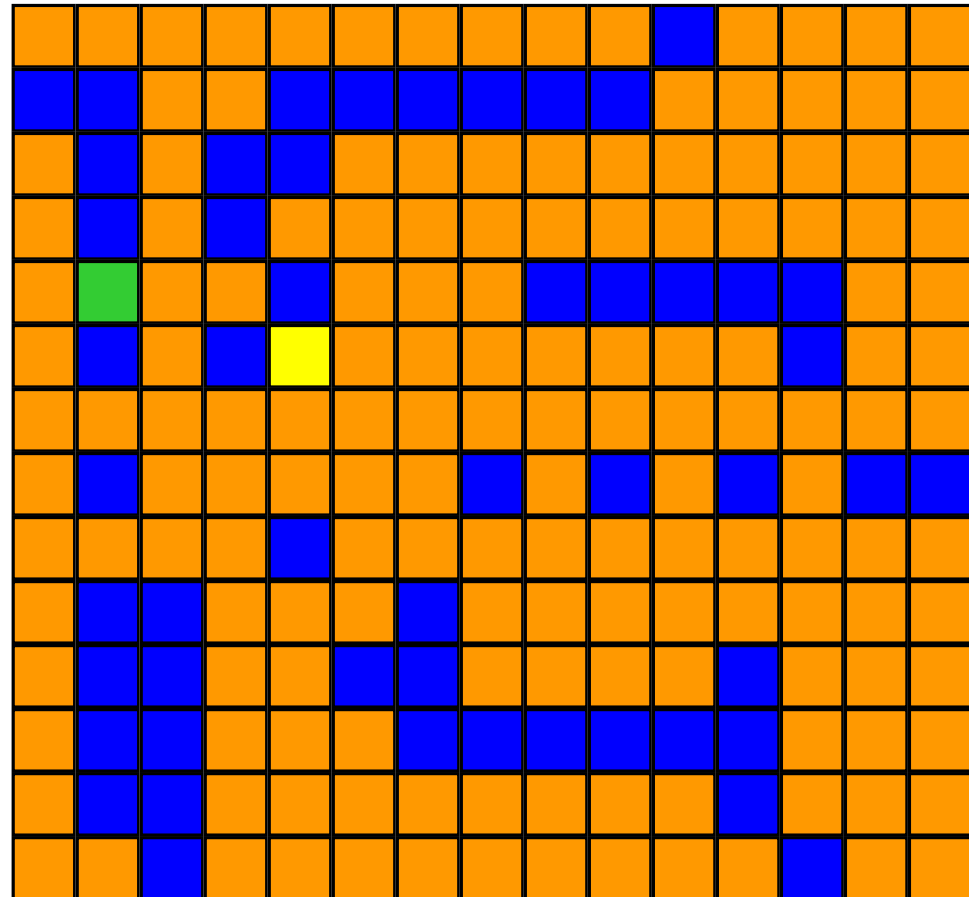


Modeling as a Grid

■ Start
Pin
■ End
Pin

- Blue squares are **blocked** squares.
- Orange squares are **available** to route a wire

How to find a path from the start pin to the end pin?



Wire Routing: Lee's Algorithm

- A **queue** of reachable squares from the start pin is used.
- The queue is initially empty, and the square of the start pin is the **examine cell**.
 - This cell has a distance value of 0.
- All **unreached unblocked** squares adjacent to the **examine cell** are marked with their distance (this is 1 more than the distance value of the **examine cell**) and **added to the queue**.
- Then a cell is removed from the **queue** and made the new **examine cell**.
- This process is repeated until the end pin is reached (path found) or the queue becomes

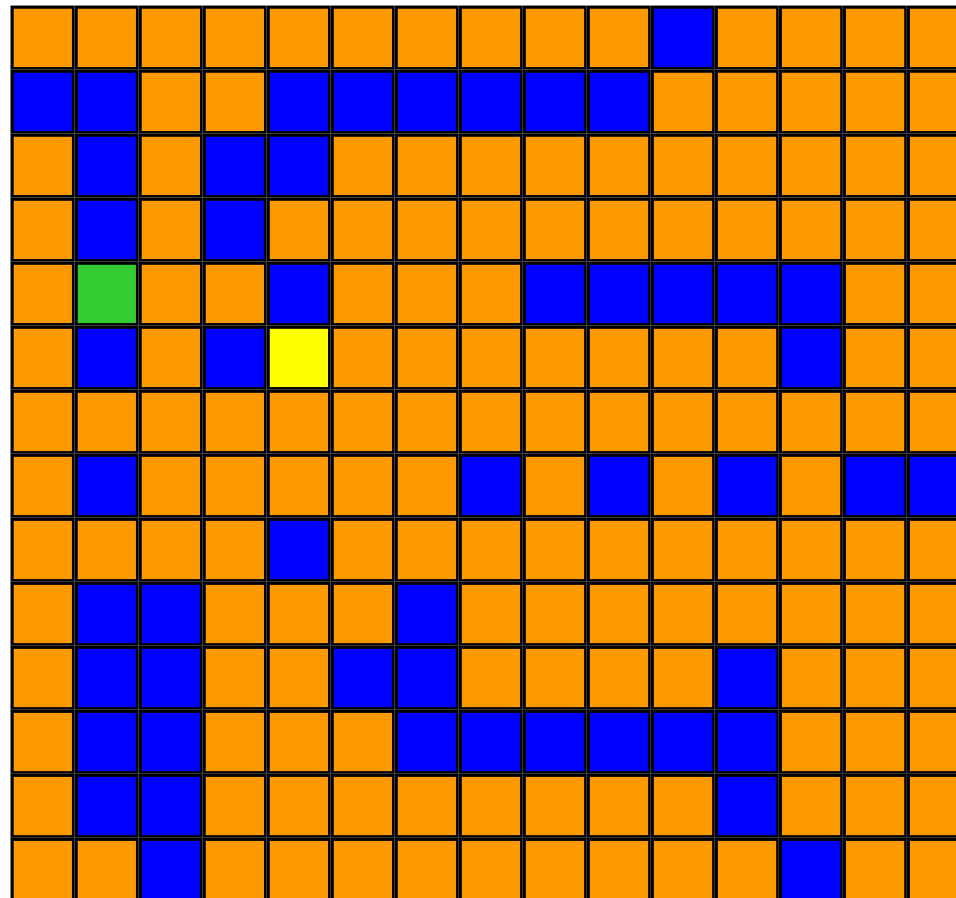
Loop



Illustration of Lee's Algorithm

 start pin

 end pin

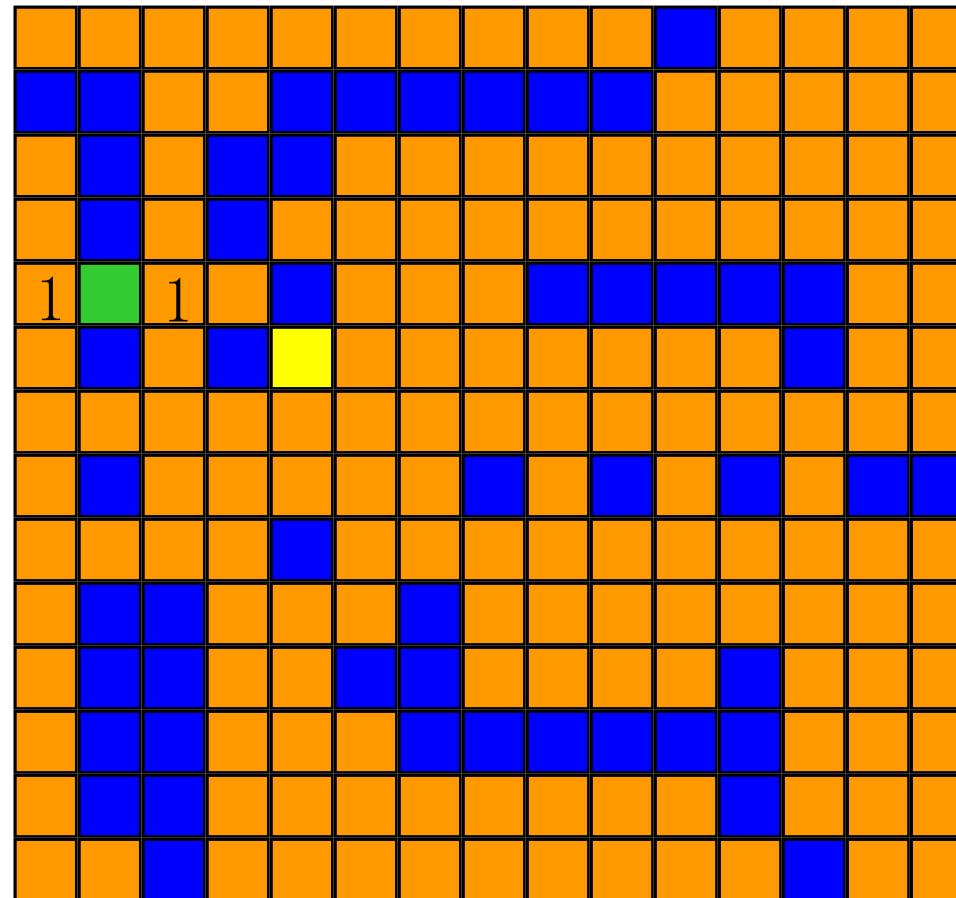


Label all reachable squares 1
unit from start.

Illustration of Lee's Algorithm

 start pin

 end pin

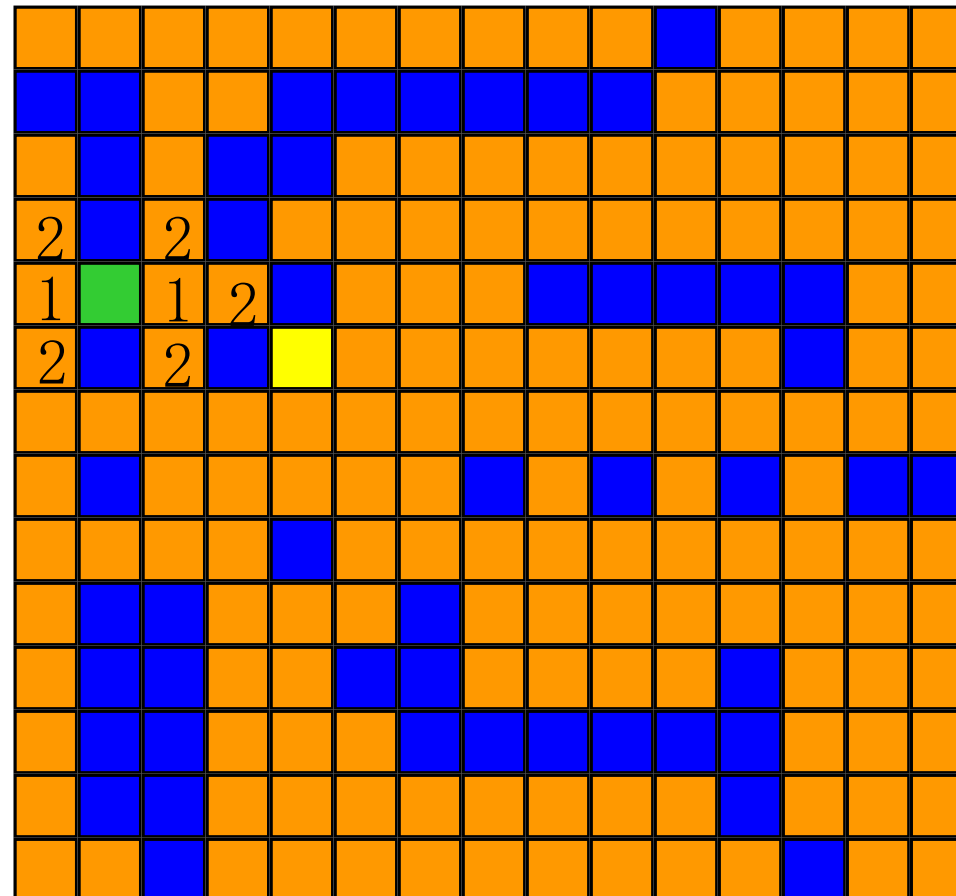


Label all reachable squares **2**
unit from start.

Illustration of Lee's Algorithm

 start pin

 end pin

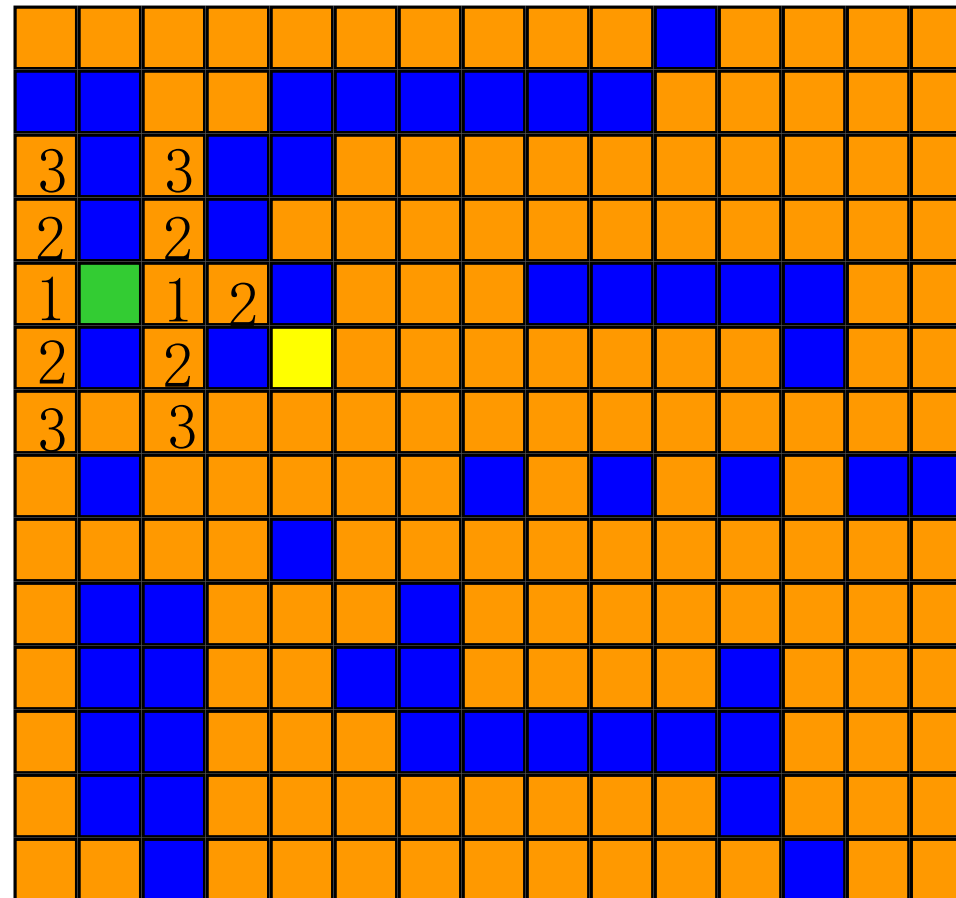


Label all reachable squares 3
unit from start.

Illustration of Lee's Algorithm

 start pin

 end pin

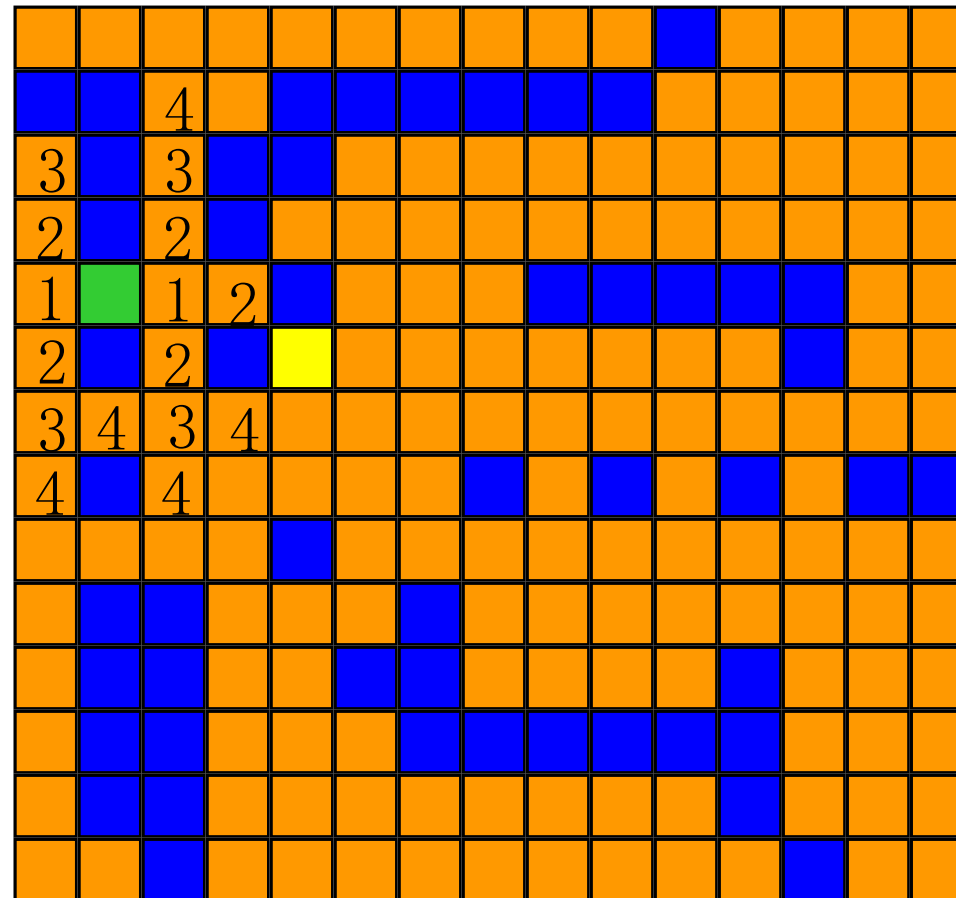


Label all reachable squares 4
unit from start.

Illustration of Lee's Algorithm

 start pin

 end pin

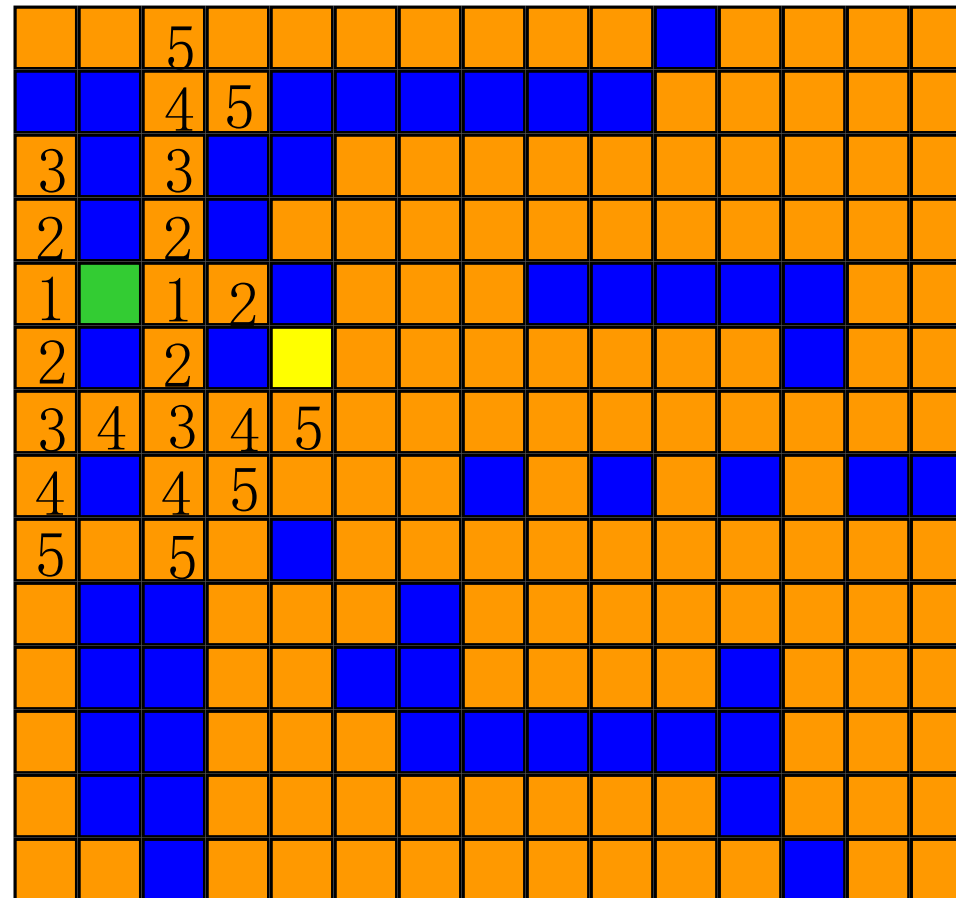


Label all reachable squares 5
unit from start.

Illustration of Lee's Algorithm

 start pin

 end pin

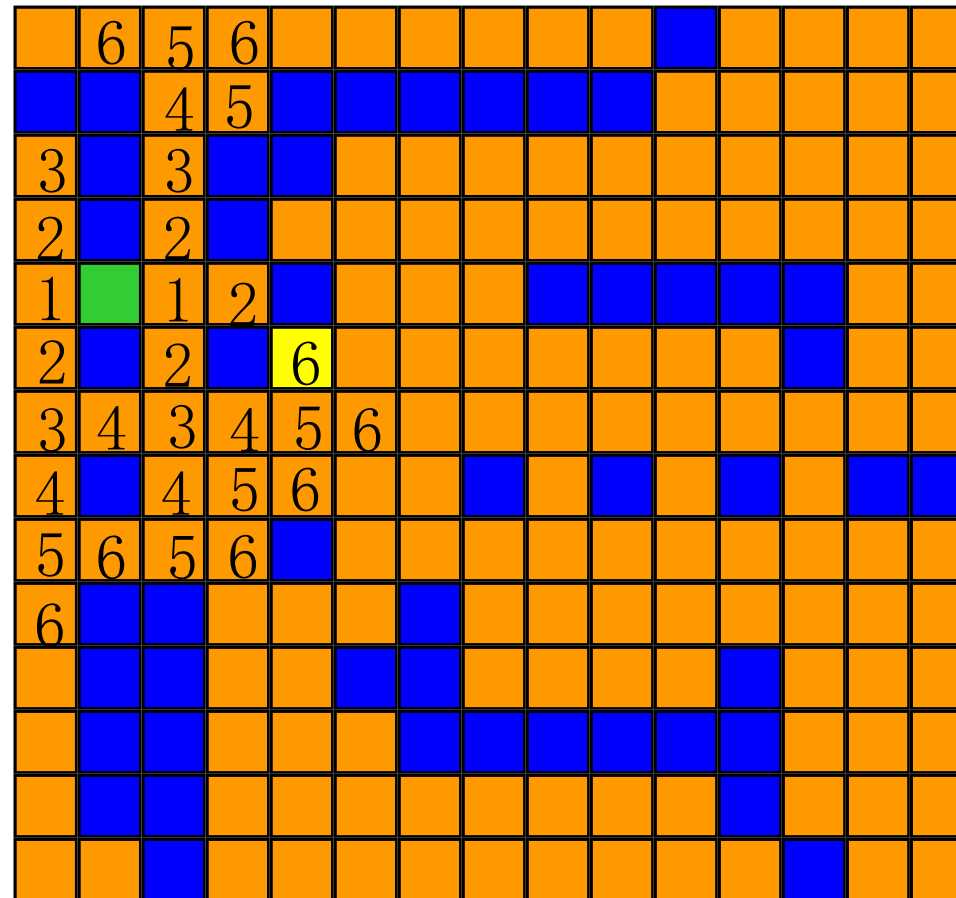


Label all reachable squares **6**
unit from start.

Illustration of Lee's Algorithm

 start pin



 end pin

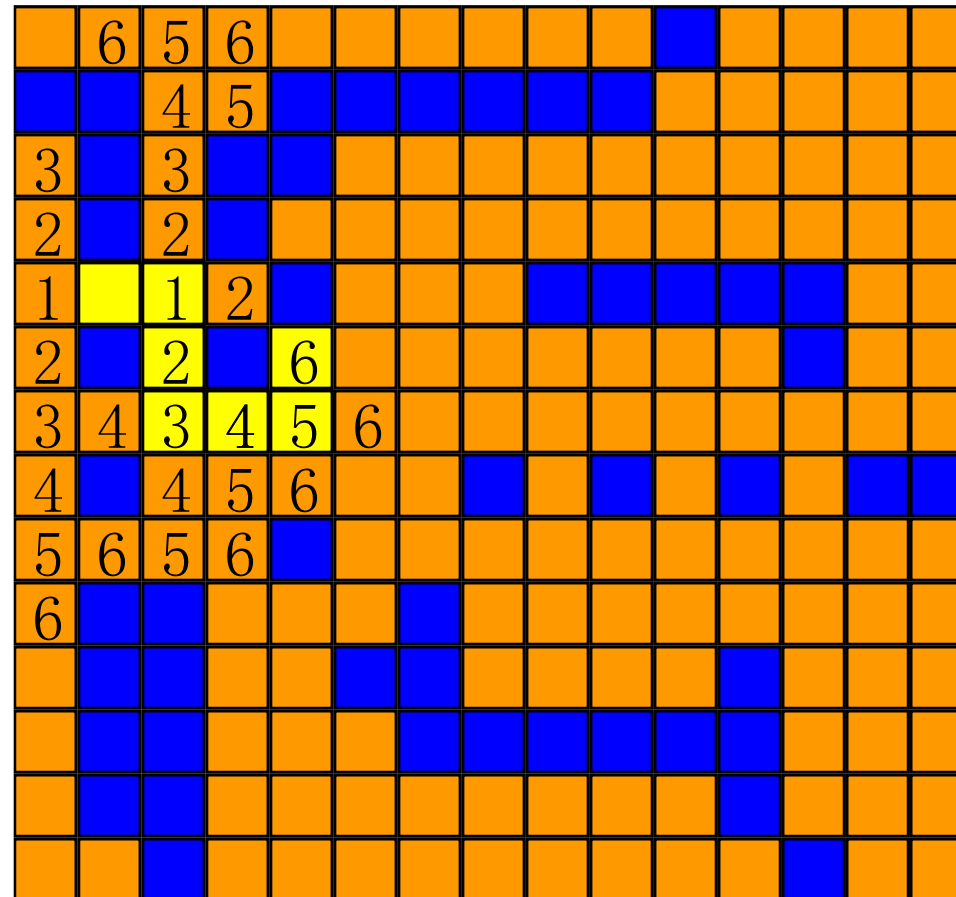


End pin reached.

Trace back.

Illustration of Lee's Algorithm

 start pin
 end pin



Deque

- Not a proper English word, pronounced as “deck” .
- A combination of stack and queue.
 - Items can be inserted and removed from **both ends** the list.
- Modification methods supported
 - **push_front(Object o)**
 - **push_back(Object o)**
 - **pop_front()**
 - **pop_back()**

Deque Implementation

- Linked list
 - Which type of linked list will you choose to support fast insertion and removal?
 - Double-ended doubly-linked list
- Circular array
 - front and rear not only need to be incremented (**push_rear**, **pop_front**), but also need to be decremented (**push_front**, **pop_rear**).