

# towards high assurance cryptographic software

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*Karthikeyan Bhargavan*



Joint work with many others at Inria, Cryspen, MSR, U. Porto, U. Aarhus, and elsewhere

# The Golden Age of Crypto?

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**New Constructions**

**New Protocols**

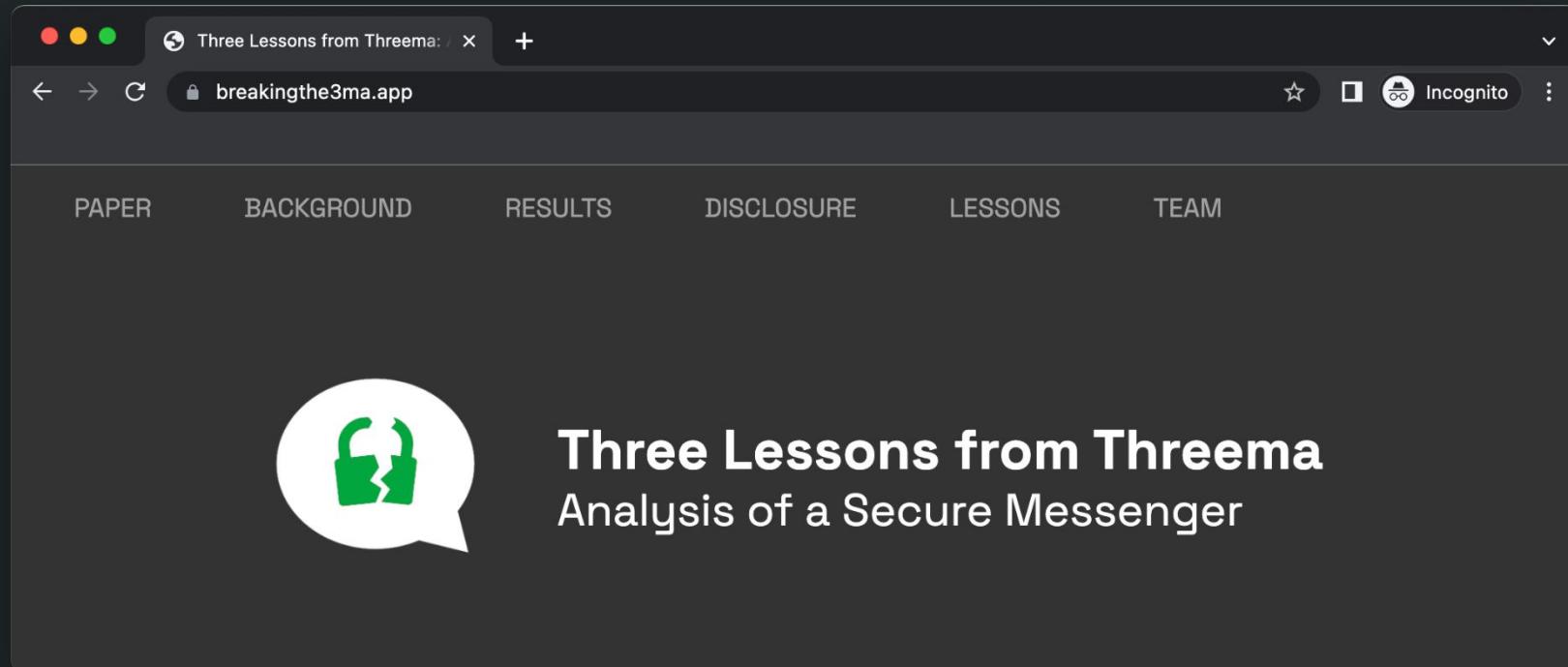
**New Applications**

Post-Quantum Crypto, Homomorphic Encryption, ...

Group Messaging, IoT Software Updates, ...

Blockchains, Privacy-Preserving Machine Learning, ...

# Designing Secure Protocols Is Still Hard



The screenshot shows a dark-themed web browser window. At the top, the address bar displays the URL [breakingthe3ma.app](https://breakingthe3ma.app). The main content area features a large white speech bubble icon containing a green padlock that is cracked open. To the right of the icon, the text "Three Lessons from Threema" is displayed in large, bold, white font, followed by "Analysis of a Secure Messenger" in a smaller white font. Below the main title, there is a navigation menu with links: PAPER, BACKGROUND, RESULTS, DISCLOSURE, LESSONS, and TEAM. The browser's interface includes standard controls like back, forward, and search, along with an Incognito mode button.

<https://breakingthe3ma.app>

# Three Lessons From Threema: Analysis of a Secure Messenger

Kenneth G. Paterson

*Applied Cryptography Group,  
ETH Zurich*

Matteo Scarlata

*Applied Cryptography Group,  
ETH Zurich*

Kien Tuong Truong

*Applied Cryptography Group,  
ETH Zurich*

Usenix 2023

- Don't roll your own crypto protocol
- Beware of cross-protocol interactions
- Formally analyze your protocol design  
using a (semi-automated) verification tool

# Formal Analysis during Protocol Design

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**Crypto Constructions**

HPKE, ...

**Crypto Protocols**

TLS 1.3, MLS, ...

# The TLS 1.3 experiment [2014-2018]

Multi-year effort to redesign IETF Transport Layer Security

- 4 years, 28 drafts, 12 IETF meetings

Major contributions from academic security researchers

- **Cryptographic analyses and proofs (of drafts 5,9,10)**  
[Dowling et al. CCS'15-J.Crypt 2021, Jager et al. CCS'15, Krawczyk et al. Euro S&P'16, ...]
- **Mechanized cryptographic proofs (of draft 18) with CryptoVerif**  
[Bhargavan et al. S&P'17]
- **Automated symbolic protocol analysis with Tamarin and ProVerif**  
[Cremers et al. Oakland'16 + CCS'17, Bhargavan et al. S&P'17 + CCS'22]
- **Verified implementation code in F\***  
[Bhargavan et al. S&P'17 and S&P'17]

# The TLS 1.3 experiment [2014-2018]

Multi-year effort to redesign IETF Transport Layer Security

- 4 years, 28 drafts, 12 IETF meetings

Major contributions from academic security researchers

ACM CCS 2023

## A Symbolic Analysis of Privacy for TLS 1.3 with Encrypted Client Hello

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# How to encrypt a message?



Encrypt a symmetric key using a peer's public key

- Standard PKE/KEM. Just use RSA, ECIES, PQ-KEM?

Long messages? A stream of messages?

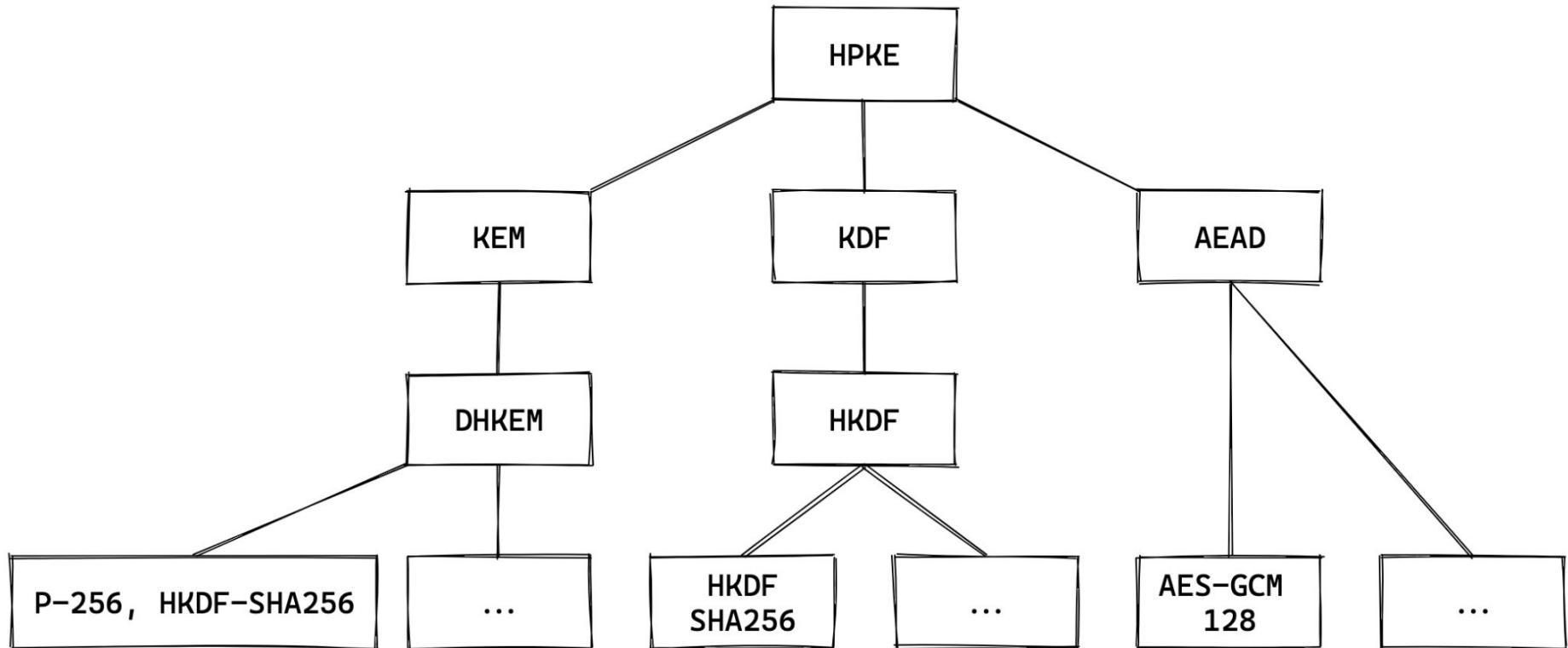
Sender authentication? PKE with Associated Data?

Stream: Internet Research Task Force (IRTF)  
RFC: [9180](#)  
Category: Informational  
Published: February 2022  
ISSN: 2070-1721  
Authors: R. Barnes K. Bhargavan B. Lipp C. Wood  
*Cisco* *Inria* *Inria* *Cloudflare*

# RFC 9180

# Hybrid Public Key Encryption

# HPKE: Agile, Modular Construction



# HPKE Proof using CryptoVerif

EuroCrypt 2022

## Analysing the HPKE Standard \*

Joël Alwen<sup>1</sup>, Bruno Blanchet<sup>2</sup>, Eduard Hauck<sup>3</sup> , Eike Kiltz<sup>3</sup> , Benjamin Lipp<sup>2</sup> , and  
Doreen Riepel<sup>3</sup> 

<sup>1</sup> Wickr

[jalwen@wickr.com](mailto:jalwen@wickr.com)

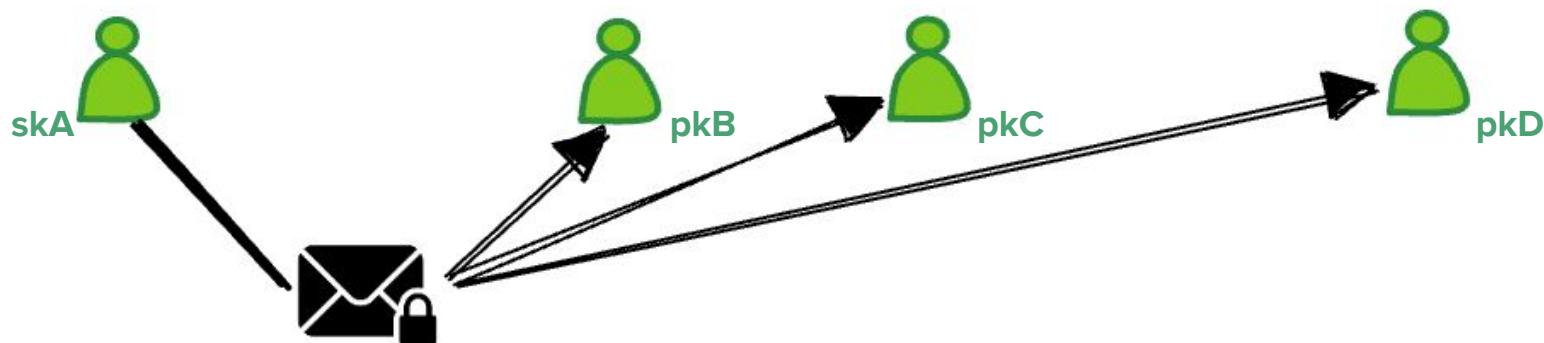
<sup>2</sup> Inria Paris

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<sup>3</sup> Ruhr-Universität Bochum

[{eduard.hauck,eike.kiltz,doreen.riepel}@rub.de](mailto:{eduard.hauck,eike.kiltz,doreen.riepel}@rub.de)

# How to encrypt group messages?

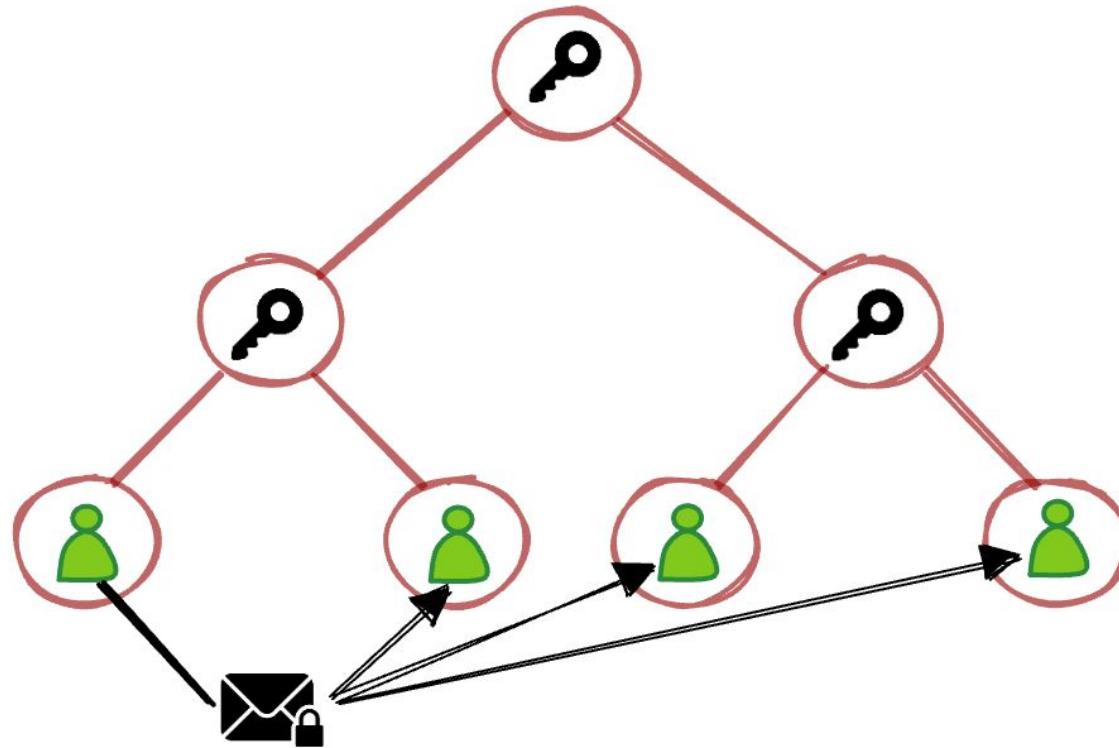


Use HPKE to encrypt message N times?

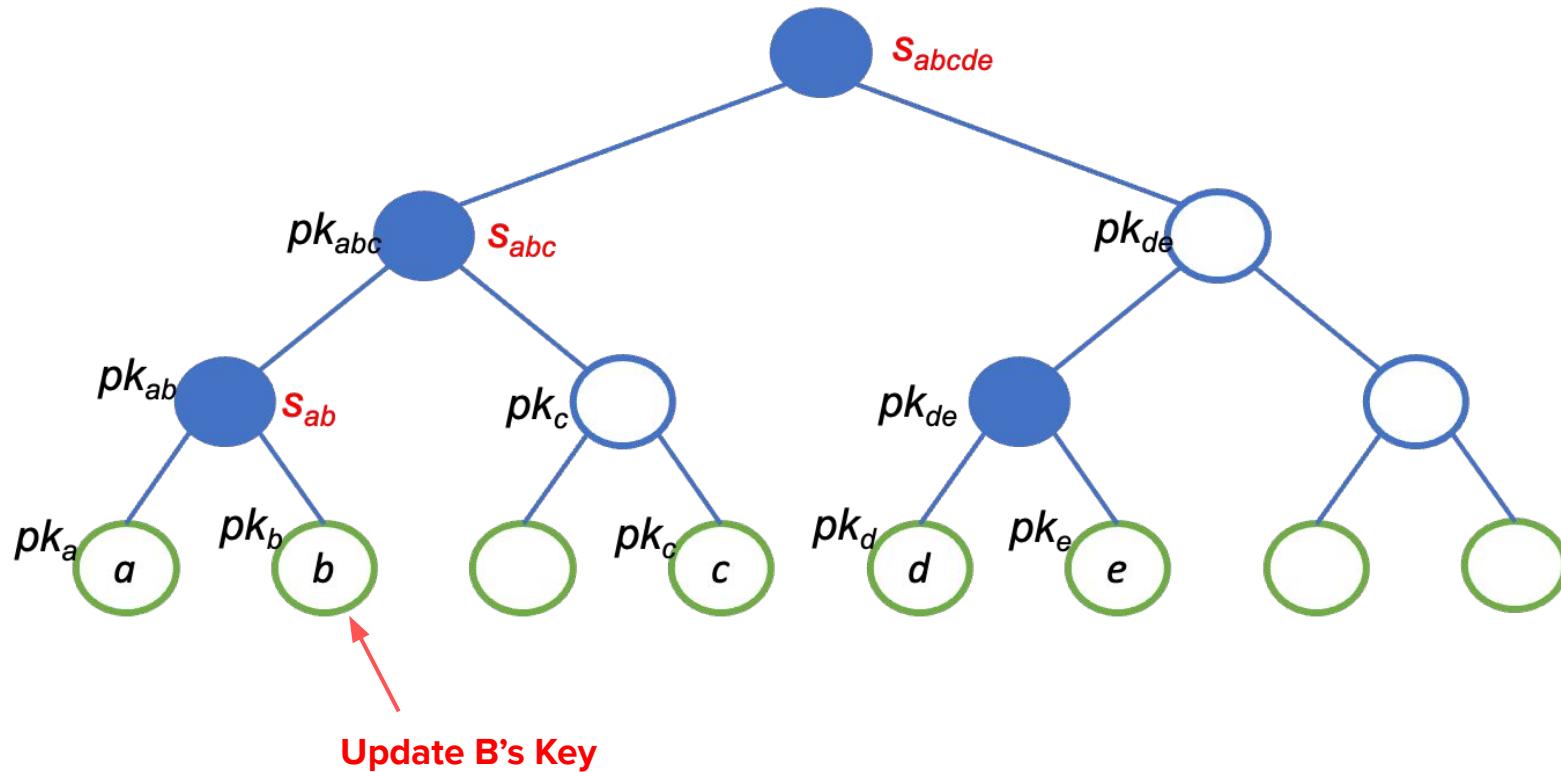
- **Sender Keys:** KEM key to N recipients, Sign every message
- $O(N)$  computation for key changes

# TreeKEM: keys for a tree of subgroups

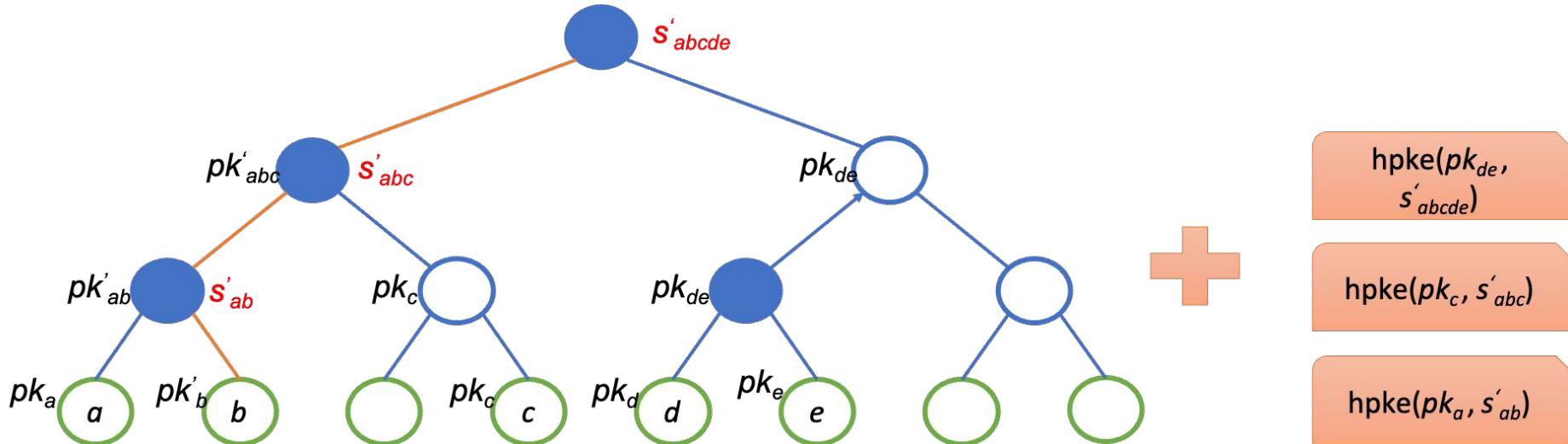
[Bhargavan, Barnes, Rescorla, 2018]



# TreeKEM: keys for a tree of subgroups



# TreeKEM: keys for a tree of subgroups



Maintain a tree of subgroup keys with efficient updates

- Add is  $O(1)$ , Remove and Update are  $O(\log n)$

draft-ietf-mls-protocol-20

← → C [datatracker.ietf.org/doc/html/draft-ietf-mls-protocol-20](https://datatracker.ietf.org/doc/html/draft-ietf-mls-protocol-20) ☆ □  Incognito :

draft-ietf-mls-protocol-20 **Internet-Draft**

Workgroup: Network Working Group  
Internet-Draft: draft-ietf-mls-protocol-20  
Published: 27 March 2023  
Intended Status: Standards Track  
Expires: 28 September 2023

R. Barnes  
Cisco  
B. Beurdouche  
Inria & Mozilla  
R. Robert  
Phoenix R&D  
J. Millican  
Meta Platforms  
E. Omara  
Google  
K. Cohn-Gordon  
University of Oxford

The Messaging Layer Security (MLS) Protocol

# Decomposing Messaging Layer Security

TreeSync      **synchronize membership and tree**



TreeKEM      **derive, encapsulate subgroup keys**



TreeDEM      **encrypt application messages**

# Decomposing Messaging Layer Security

TreeSync



**synchronize membership and tree**  
(authentication, integrity)

TREE HASH + SIG

TreeKEM



**derive, encapsulate subgroup keys**  
(forward secrecy,  
post-compromise security)

HPKE + KDF

TreeDEM

**encrypt application messages**  
(forward secrecy, sender auth)

KDF + AEAD + SIG

# Decomposing Messaging Layer Security

TreeSync

synchronize membership and tree  
(authentication)

TREE HASH + SIG

## TreeSync: Authenticated Group Management for Messaging Layer Security

Théophile Wallez  
*Inria Paris*

Jonathan Protzenko  
*Microsoft Research*

Benjamin Beurdouche  
*Mozilla*

Karthikeyan Bhargavan  
*Inria Paris*

# Many Ongoing Analyses of IETF MLS

	MLS Version	Part Analyzed	Adversarial Model	Considers Group Splits	Framework
[25]	Draft 1 (ART)	CGKA in static groups	active	yes	part game-based, part symbolic
[6]	Draft 6	CGKA	passive	no	game-based
[18]	Draft 7	Messaging	insider	yes	symbolic
[7]	Draft 11	Messaging	semi-active	yes	game-based
[20]	Draft 11	Key derivation	insider	n/a	game-based
[26]	Draft 11	Multi-group messaging	n/a	n/a	n/a
this work	Draft 12	CGKA	insider	yes	UC

From: On The Insider Security of MLS, Alwen et al. CRYPTO 2022

# Verifying Crypto Implementations

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**Verified Crypto Code**

**Verified Protocol Code**

HACL\*, Vale, libjade, libcrux, ...

miTLS, LibSignal\*, Noise\*, MLS\*, ...



**NSS**

**BoringSSL**

**Web**

**CRYPTO  
LIBRARIES**

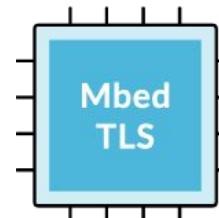
**OS**



**Lang**



**IoT**





12

NSS

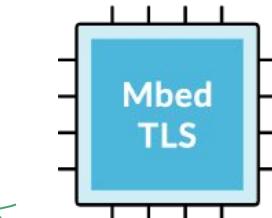
BoringSSL

Web

TRUSTED (?)  
COMPUTING  
BASE



OS



Mbed  
TLS



Lang



6

IoT



10

# Many Bugs in Classic Crypto Code

## CVE-2022-21449: Psychic Signatures in Java

Neil Madden

19 April, 2022

cryptography,  
Security

API security,  
cryptography, Java,  
jose, jwt, web-  
security

The long-running BBC sci-fi show [Doctor Who](#) has a recurring plot device where the Doctor manages to get out of trouble by showing an identity card which is actually completely blank. Of course, this being Doctor Who, the card is really made out of a special “[psychic paper](#)”, which causes the person looking at it to see whatever the Doctor wants them to see: a security pass, a warrant, or whatever.



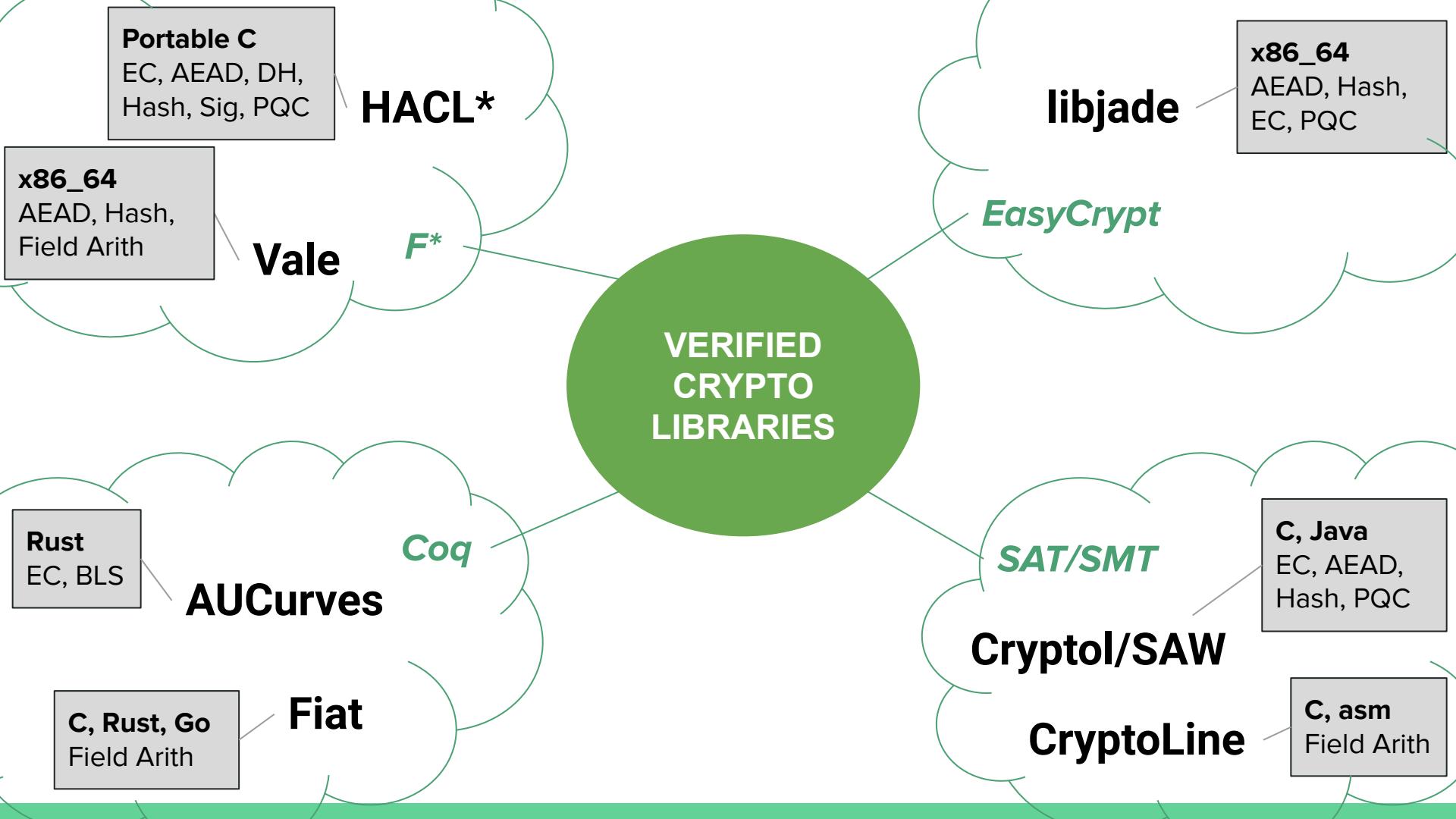
<https://neilmadden.blog/2022/04/19/psychic-signatures-in-java/>

# Many Bugs in Classic Crypto Code

This is why the very first check in the ECDSA verification algorithm is to ensure that  $r$  and  $s$  are both  $\geq 1$ .

Guess which check Java forgot?

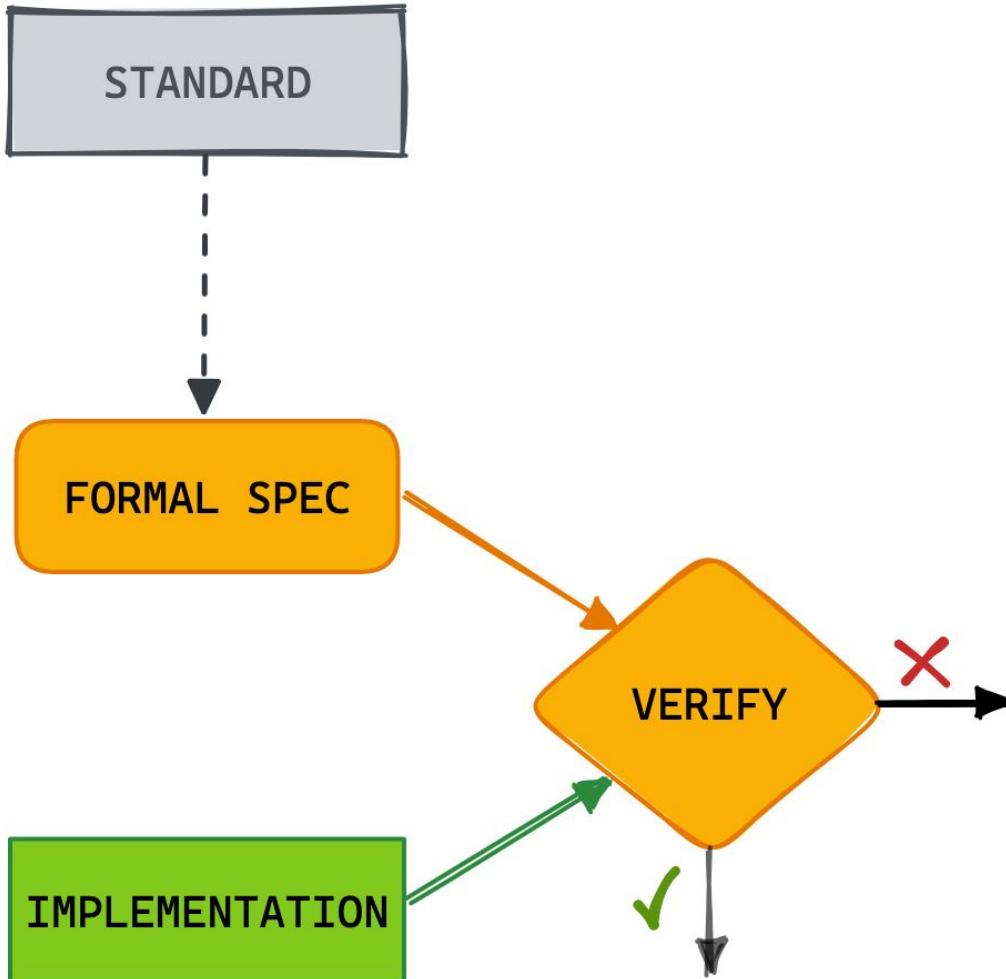
That's right. Java's implementation of ECDSA signature verification didn't check if  $r$  or  $s$  were zero, so you could produce a signature value in which they are both 0 ([appropriately encoded](#)) and Java would accept it as a valid signature for any message and for any public key. The digital equivalent of a blank ID card.



**Good news:** For any modern crypto algorithm,  
there is probably a verified implementation.

---

**But...** research code with low-level APIs, and  
specs written in unfamiliar formal languages.



# Verified Cryptography Workflow

# STANDARD

Internet Research Task Force (IRTF)  
Request for Comments: 8439  
Obsoletes: [7539](#)  
Category: Informational  
ISSN: 2070-1721

Y. Nir  
Dell EMC  
A. Langley  
Google, Inc.  
June 2018



# FORMAL SPEC

## ChaCha20 and Poly1305 for IETF Protocols

### Abstract

This document defines the ChaCha20 stream cipher and the Poly1305 authenticator, both as stand-alone components and in a "combined mode", or Authenticated Encryption.

**IETF RFC or  
NIST Standard**

### 2.1. The ChaCha Quarter Round

The basic operation of the ChaCha algorithm is the quarter round. It operates on four 32-bit unsigned integers, denoted a, b, c, and d. The operation is as follows (in C-like notation):

```
a += b; d ^= a; d <<= 16;  
c += d; b ^= c; b <<= 12;  
a += b; d ^= a; d <<= 8;  
c += d; b ^= c; b <<= 7;
```

**In English +  
Pseudocode**

### 2.1.1. Test Vector for the ChaCha Quarter Round

For a test vector, we will use the same numbers as in the example, adding something random for c.

```
a = 0x11111111  
b = 0x01020304  
c = 0x9b8d6f43  
d = 0x01234567
```

**+ Test Vectors**

# IMPLEMENTATION

## STANDARD



## FORMAL SPEC

## IMPLEMENTATION

```
let line (a:idx) (b:idx) (d:idx) (s:rotval U32) (m:state) : Tot state =
  let m = m.[a] ← (m.[a] +. m.[b]) in
  let m = m.[d] ← ((m.[d] ^. m.[a]) <<<. s) in m
```

```
let quarter_round a b c d : Tot shuffle =
  line a b d (size 16) @
  line c d b (size 12) @
  line a b d (size 8) @
  line c d b (size 7)
```

F\* Spec  
(HACL\*)

```
proc chacha20_line(a : int, b : int, d : int, s : int, st : State) = {
  var state;
  state <- st;
  state.[a] <- ((state).[a]) + ((state).[b]);
  state.[d] <- ((state).[d]) `^` ((state).[a]);
  state.[d] <- rotate_left ((state).[d]) (s);
  return state;
}

proc chacha20_quarter_round(a : int, b : int, c : int, d : int, st : State) = {
  var state;
  state <@ chacha20_line (a, b, d, 16, st);
  state <@ chacha20_line (c, d, b, 12, state);
  state <@ chacha20_line (a, b, d, 8, state);
  state <@ chacha20_line (c, d, b, 7, state);
  return state;
}
```

EasyCrypt Spec  
(libjade)

STANDARD

F\* Implementation

FORMAL SPEC

Translate

```
let line st a b d r =
  let sta = st.(a) in
  let stb = st.(b) in
  let std = st.(d) in
  let sta = sta +. stb in
  let std = std ^. sta in
  let std = rotate_left std r in
  st.(a) ← sta;
  st.(d) ← std

let quarter_round st a b c d =
  line st a b d (size 16);
  line st c d b (size 12);
  line st a b d (size 8);
  line st c d b (size 7)
```

IMPLEMENTATION

Portable C Code

```
static inline void quarter_round(uint32_t *st, uint32_t a, uint32_t b, uint32_t c, uint32_t d)
{
    uint32_t sta = st[a];
    uint32_t stb0 = st[b];
    uint32_t std0 = st[d];
    uint32_t sta10 = sta + stb0;
    uint32_t std10 = std0 ^ sta10;
    uint32_t std2 = std10 << (uint32_t)16U | std10 >> (uint32_t)16U;
    st[a] = sta10;
    st[d] = std2;
    ...
}
```

STANDARD

FORMAL SPEC

IMPLEMENTATION

```
inline fn __line_ref(reg u32[16] k,  
                     inline int a b c r)  
    -> reg u32[16]  
{  
    k[a] += k[b];  
    k[c] ^= k[a];  
    _, _, k[c] = #ROL_32(k[c], r);  
    return k;  
  
inline fn __quarter_round_ref(reg u32[16] k,  
                             inline int a b c d)  
    -> reg u32[16]  
{  
    k = __line_ref(k, a, b, d, 16);  
    k = __line_ref(k, c, d, b, 12);  
    k = __line_ref(k, a, b, d, 8);  
    k = __line_ref(k, c, d, b, 7);  
    return k;  
}
```

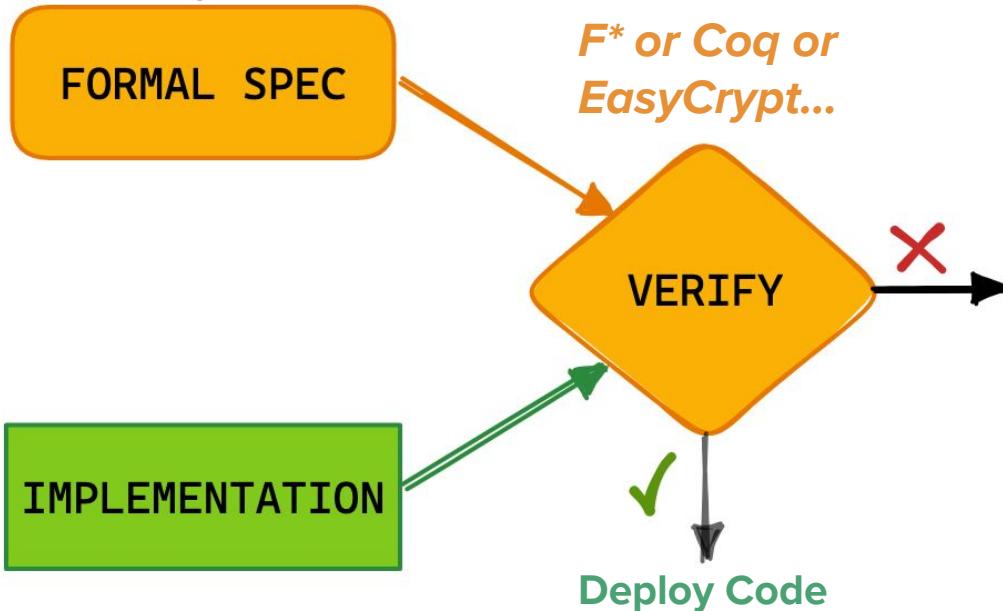
Jasmin  
Implementation

Translate

Intel AVX2  
Assembly

```
vpaddd %ymm4, %ymm0, %ymm0  
vpxor %ymm0, %ymm12, %ymm12  
(%rsp), %ymm12, %ymm12  
vpaddd %ymm12, %ymm8, %ymm8  
vpaddd %ymm6, %ymm2, %ymm2  
vpxor %ymm8, %ymm4, %ymm4  
vpxor %ymm2, %ymm14, %ymm14  
vpslld $12, %ymm4, %ymm15  
vpsrlld $20, %ymm4, %ymm4  
vpxor %ymm15, %ymm4, %ymm4  
vpshufb (%rsp), %ymm14, %ymm14  
vpaddd %ymm4, %ymm0, %ymm0  
vpaddd %ymm14, %ymm10, %ymm10  
vpxor %ymm0, %ymm12, %ymm12  
vpxor %ymm10, %ymm6, %ymm6  
vpshufb 32(%rsp), %ymm12, %ymm12  
vpslld $12, %ymm6, %ymm15  
vpsrlld $20, %ymm6, %ymm6  
...
```

STANDARD



# Verified Cryptography Workflow

## Potential Implementation Bug

- Memory Safety Violation
- Functional Correctness Flaw
- Side Channel Vulnerability

Fix and re-verify

**Good news:** For any modern crypto algorithm,  
there is probably a verified implementation

- 
- You don't have to sacrifice **performance**
  - **Mechanized proofs** that you can run and re-run yourself
  - You (mostly) don't have to read or understand the proofs

But... not always easy to use, extend, or  
combine code from different libraries

---

- You do need to carefully audit the formal specs, written in tool-specific spec languages like F\*, Coq, EasyCrypt
- You do need to safely use their low-level APIs, which often embed subtle pre-conditions

# hacspec: a tool-independent spec language

## Design Goals

- **Easy to use** for crypto developers
- **Familiar** language and tools
- **Succinct** specs, like pseudocode
- **Strongly typed** to avoid spec errors
- **Executable** for spec debugging
- **Testable** against RFC test vectors
- **Translations** to formal languages like  
**F\*, Coq, EasyCrypt, ...**

# hacspeс: a tool-independent spec language

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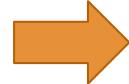
## A purely functional subset of Rust

- Safe Rust without external side-effects
- No mutable borrows
- All values are copyable
- Rust tools & development environment
- A library of common abstractions
  - Arbitrary-precision Integers
  - Secret-independent Machine Ints
  - Vectors, Matrices, Polynomials,...

# hacspe: purely functional crypto code in Rust

```
inner_block (state):  
    Qround(state, 0, 4, 8, 12)  
    Qround(state, 1, 5, 9, 13)  
    Qround(state, 2, 6, 10, 14)  
    Qround(state, 3, 7, 11, 15)  
    Qround(state, 0, 5, 10, 15)  
    Qround(state, 1, 6, 11, 12)  
    Qround(state, 2, 7, 8, 13)  
    Qround(state, 3, 4, 9, 14)  
end
```

ChaCha20 RFC



```
fn inner_block(st: State) -> State {  
    let mut state = st;  
    state = chacha20_quarter_round(0, 4, 8, 12, state);  
    state = chacha20_quarter_round(1, 5, 9, 13, state);  
    state = chacha20_quarter_round(2, 6, 10, 14, state);  
    state = chacha20_quarter_round(3, 7, 11, 15, state);  
    state = chacha20_quarter_round(0, 5, 10, 15, state);  
    state = chacha20_quarter_round(1, 6, 11, 12, state);  
    state = chacha20_quarter_round(2, 7, 8, 13, state);  
    chacha20_quarter_round(3, 4, 9, 14, state)  
}
```

State-passing style

ChaCha20 in  
hacspe

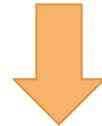
Call-by-value

# hacspe: abstract integers for field arithmetic

```
n = le_bytes_to_num(msg[((i-1)*16)..(i*16)] | [0x01])
a += n
a = (r * a) % p
```

Poly1305 RFC  
(update\_block)

Modular 130-bit Prime Field Arithmetic



```
pub fn poly1305_encode_block(b: &PolyBlock) -> FieldElement {
    let n = U128_from_le_bytes(U128Word::from_seq(b));
    let f = FieldElement::from_secret_literal(n);
    f + FieldElement::pow2(128)
}

pub fn poly1305_update_block(b: &PolyBlock, (acc,r,s): PolyState) -> PolyState {
    ((poly1305_encode_block(b) + acc) * r, r, s)
}
```

Poly1305 in  
hacspe

Modular Arithmetic over User-Defined Field

# hacspe: secret integers for “constant-time” specs

## Separate Secret and Public Values

- New types: U8, U32, U64, U128
- Can do arithmetic: +, \*, -
- Can do bitwise ops: ^, |, &
- Cannot do division: /, %
- Cannot do comparison: ==, !=, <, ...
- Cannot use as array indexes: x[u]

## Enforces secret independence

- A “constant-time” discipline
- Important for some crypto specs

```
fn chacha20_line(a: StateIdx, b: StateIdx, d: StateIdx,  
                  s: usize, mut state: State) -> State {  
    state[a] = state[a] + state[b];  
    state[d] = state[d] ^ state[a];  
    state[d] = state[d].rotate_left(s);  
    state  
}
```

ChaCha20 in  
hacspe

```
fn sub_bytes(state: Block) -> Block {  
    let mut st = state;  
    for i in 0..BLOCKSIZE {  
        st[i] = SBOX[U8::declassify(state[i])];  
    }  
    st  
}
```

AES in  
hacspe

# hacspe: translation to formal languages

```
pub fn chacha20_quarter_round(  
    a: StateIdx,  
    b: StateIdx,  
    c: StateIdx,  
    d: StateIdx,  
    mut state: State,  
) -> State {  
    state = chacha20_line(a, b, d, 16, state);  
    state = chacha20_line(c, d, b, 12, state);  
    state = chacha20_line(a, b, d, 8, state);  
    chacha20_line(c, d, b, 7, state)  
}
```

ChaCha20 in  
hacspe

```
let chacha20_quarter_round (a b c d: state_idx_t) (state: state_t) : state_t =  
let state:state_t = chacha20_line a b d 16 state in  
let state:state_t = chacha20_line c d b 12 state in  
let state:state_t = chacha20_line a b d 8 state in  
chacha20_line c d b 7 state
```

F\* Spec

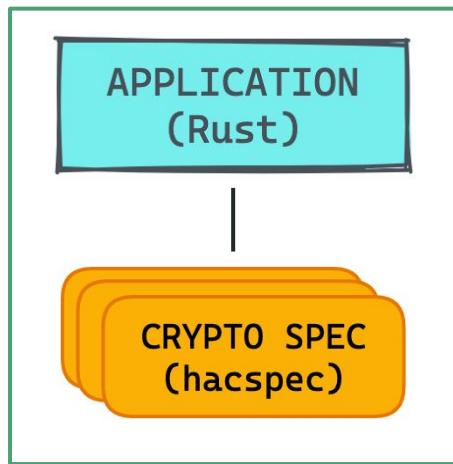
```
Definition chacha20_quarter_round (a : int32) (b : int32) (c : int32)  
                                (d : int32) (state : State) : State :=  
let state := chacha20_line a b d 16 state : State in  
let state := chacha20_line c d b 12 state : State in  
let state := chacha20_line a b d 8 state : State in  
chacha20_line c d b 7 state.
```

Coq Spec

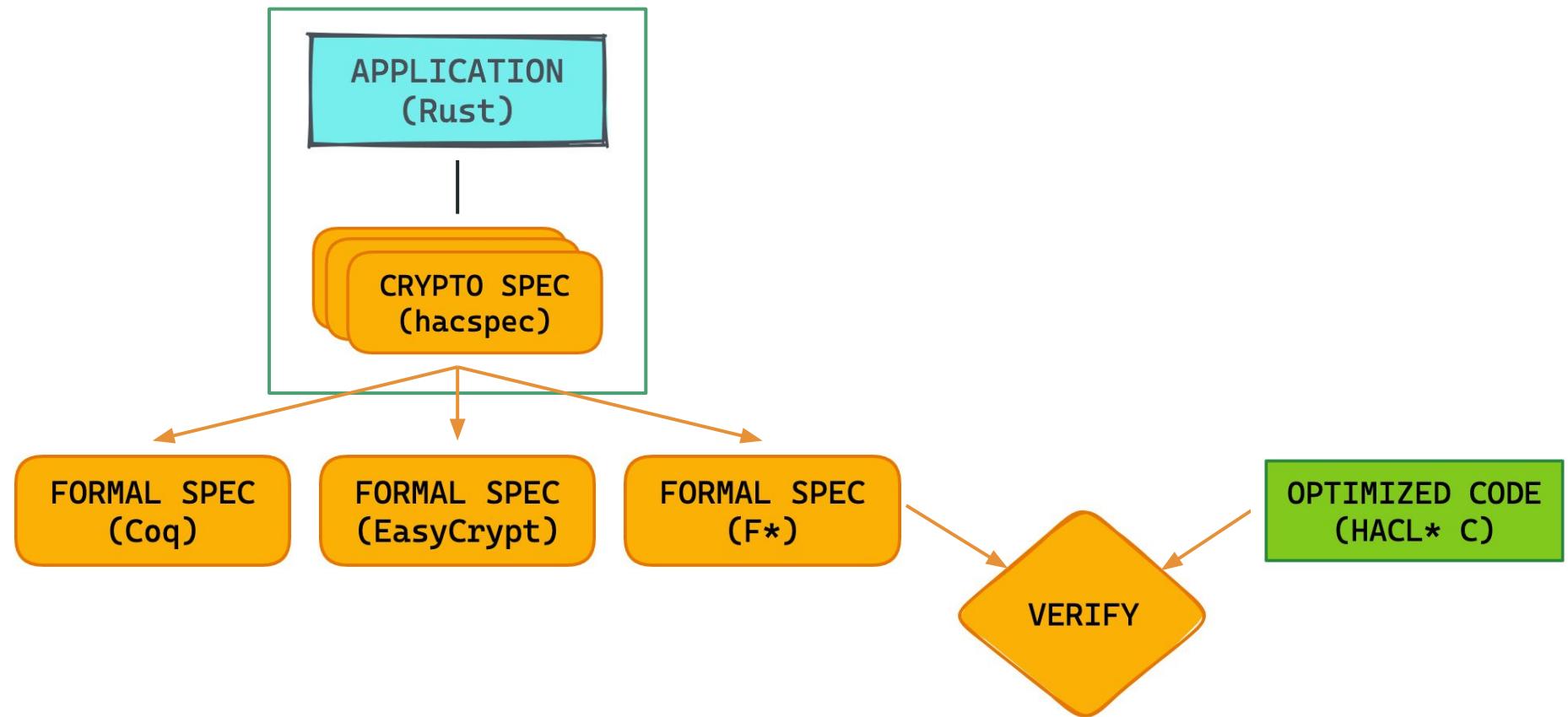
```
proc chacha20_quarter_round(a : int, b : int, c : int, d : int,  
                           state : State) = {  
    var _res;  
    state <@ chacha20_line (a, b, d, 16, state);  
    state <@ chacha20_line (c, d, b, 12, state);  
    state <@ chacha20_line (a, b, d, 8, state);  
    _res <@ chacha20_line (c, d, b, 7, state);  
    return _res;  
}
```

EasyCrypt Spec

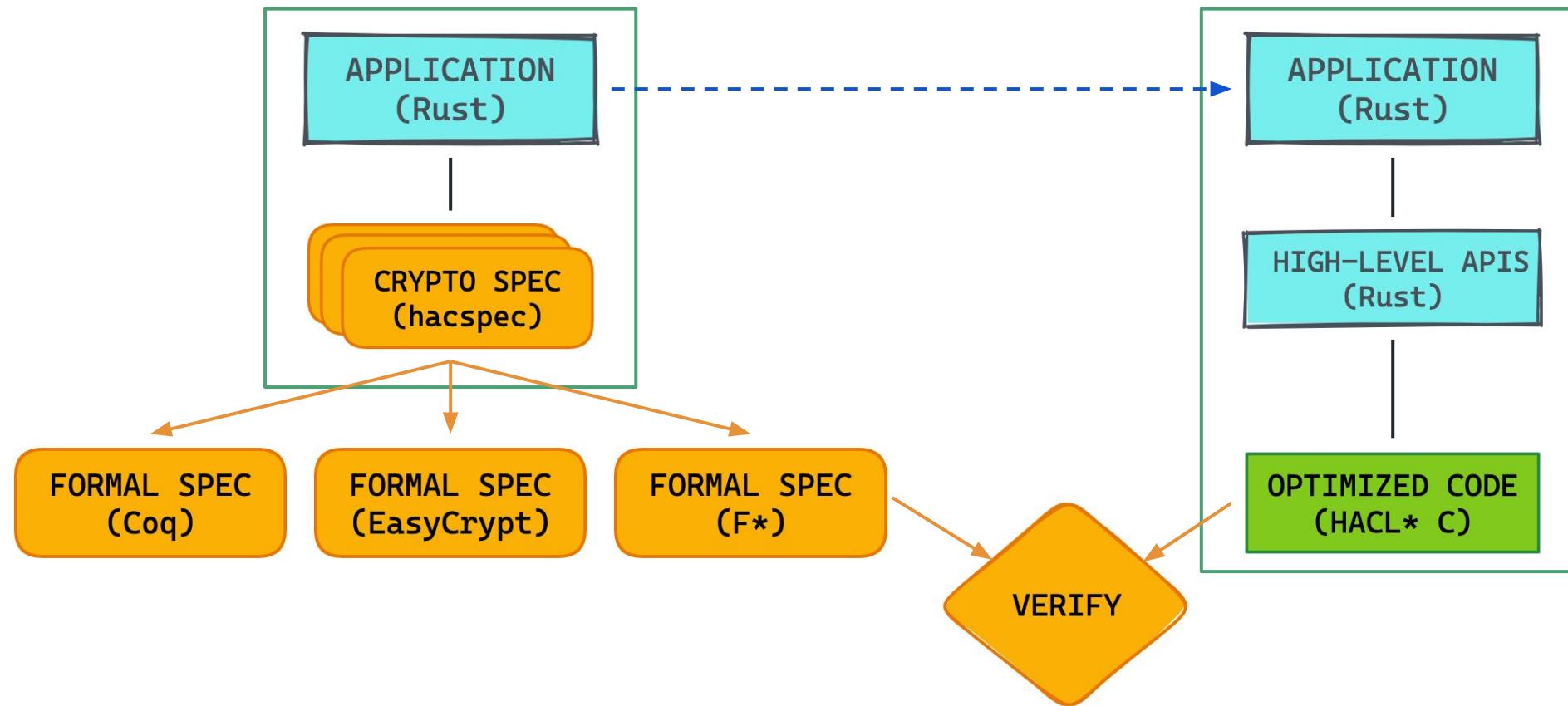
# hacspe: towards high-assurance crypto software



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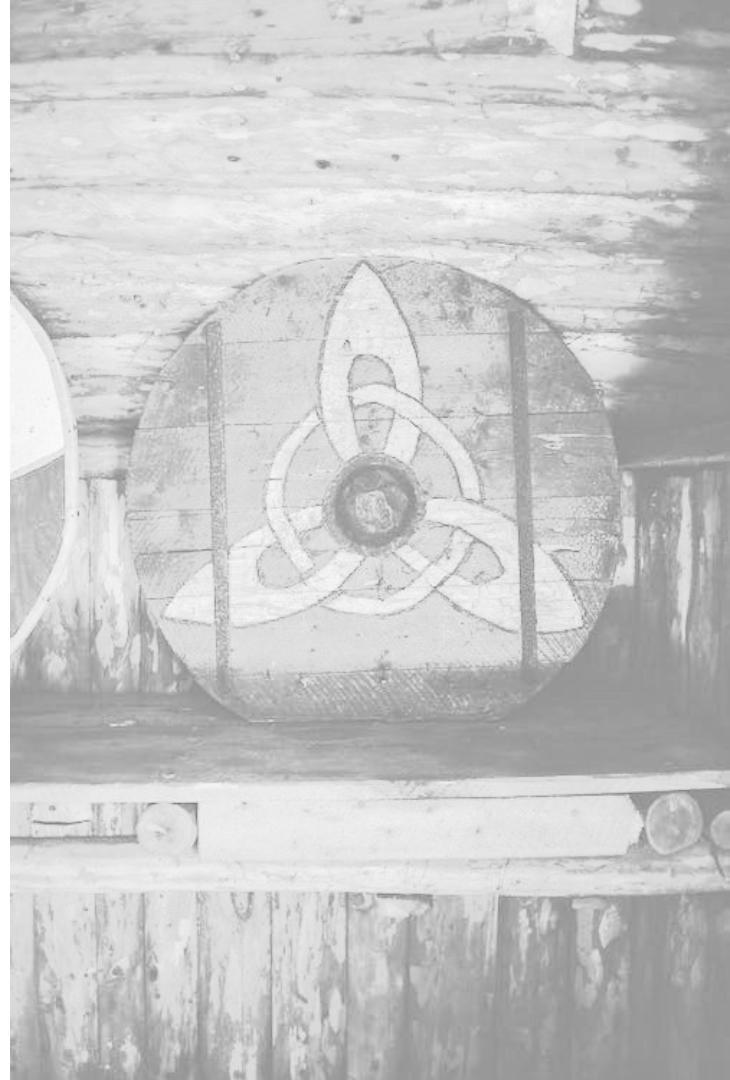
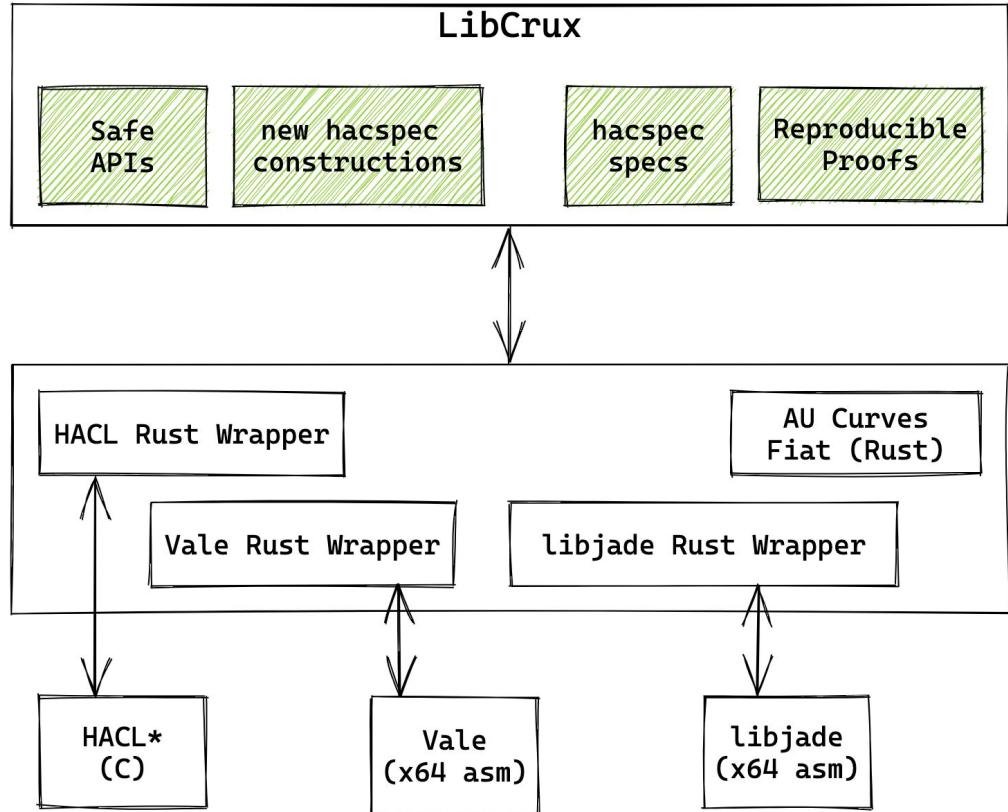
# hacspe: towards high-assurance crypto software



# libcrux: a library of verified cryptography

---

# libcrux: architecture



# Unsafe APIs: Array Constraints

```
void  
Hacl_Chacha20Poly1305_32_aead_encrypt(  
    uint8_t *k, ←  
    uint8_t *n, ←  
    uint32_t aadlen,  
    uint8_t *aad,  
    uint32_t mlen,  
    uint8_t *m,  
    uint8_t *cipher, ←  
    uint8_t *mac ←  
);
```

Fixed Length

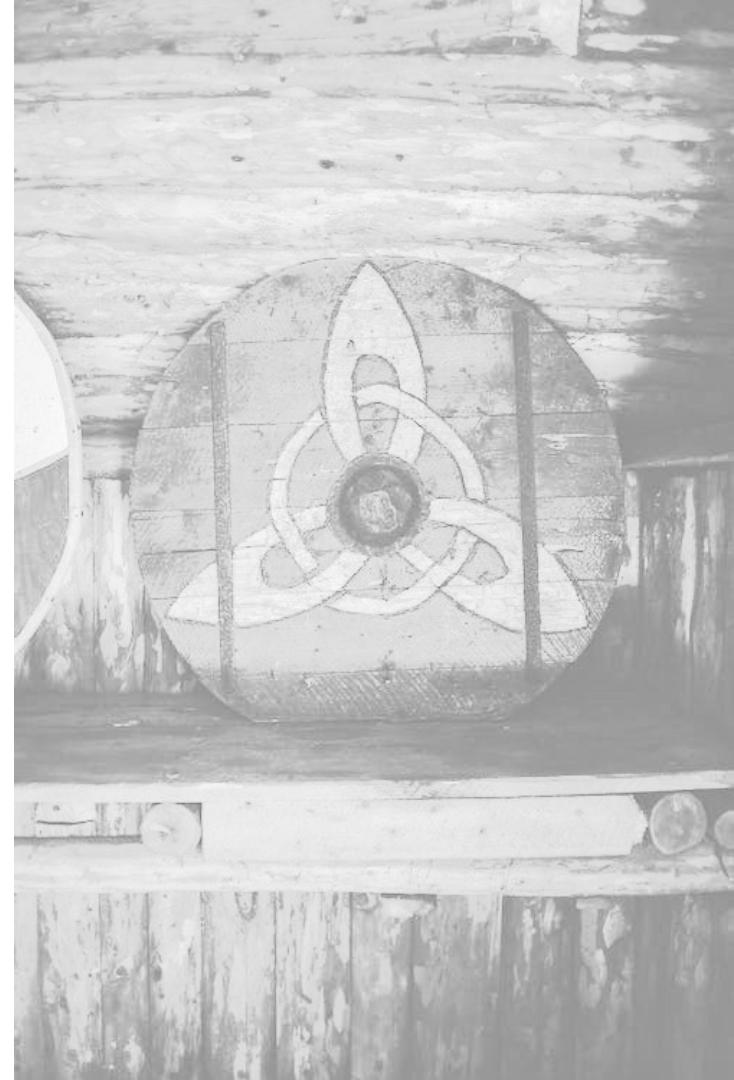
Disjoint



# Verified F\* API: Preconditions

```
let aead_encrypt_st (w:field_spec) =
  key:lbuffer uint8 32ul
  -> nonce:lbuffer uint8 12ul
  -> alen:size_t
  -> aad:lbuffer uint8 alen
  -> len:size_t
  -> input:lbuffer uint8 len
  -> output:lbuffer uint8 len
  -> tag:lbuffer uint8 16ul ->
Stack unit
(requires fun h ->
  live h key /\ live h nonce /\ live h aad /\
  live h input /\ live h output /\ live h tag /\
  disjoint key output /\ disjoint nonce output /\
  disjoint key tag /\ disjoint nonce tag /\
  disjoint output tag /\ eq_or_disjoint input output /\
  disjoint aad output)
```

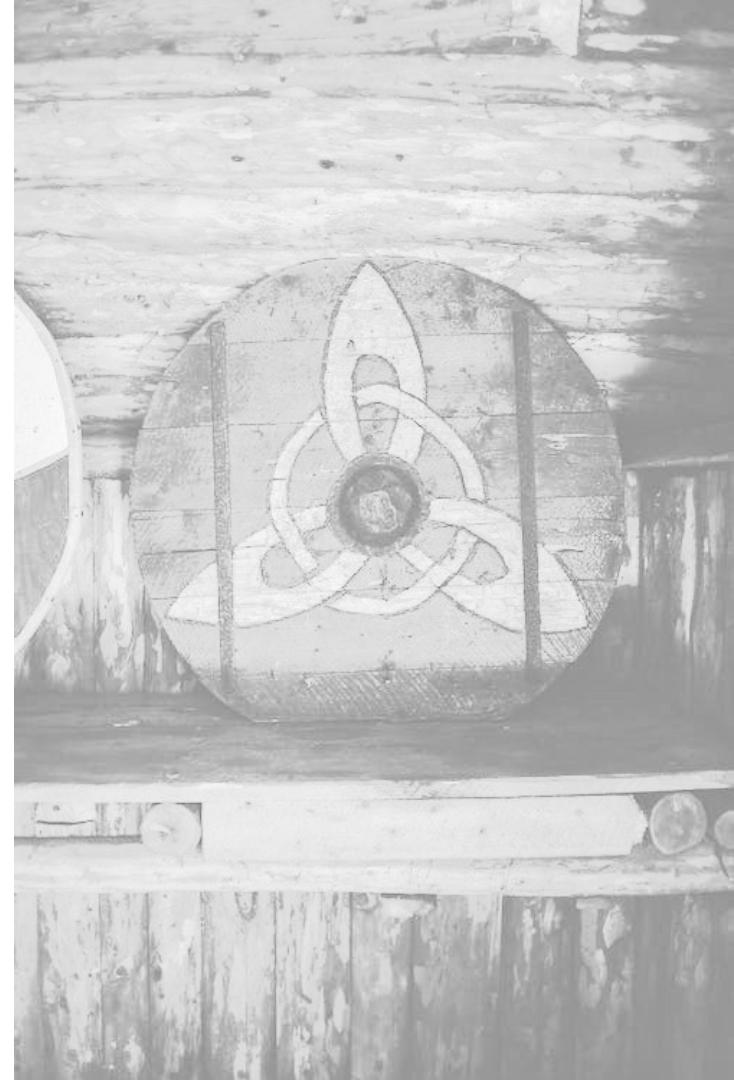
Length Constraints



# Verified F\* API: Preconditions

```
let aead_encrypt_st (w:field_spec) =
  key:lbuffer uint8 32ul
  -> nonce:lbuffer uint8 12ul
  -> alen:size_t
  -> aad:lbuffer uint8 alen
  -> len:size_t
  -> input:lbuffer uint8 len
  -> output:lbuffer uint8 len
  -> tag:lbuffer uint8 16ul ->
Stack unit
(requires fun h ->
  live h key /\ live h nonce /\ live h aad /\ 
  live h input /\ live h output /\ live h tag /\ 
  disjoint key output /\ disjoint nonce output /\ 
  disjoint key tag /\ disjoint nonce tag /\ 
  disjoint output tag /\ eq_or_disjoint input output /\ 
  disjoint aad output)
```

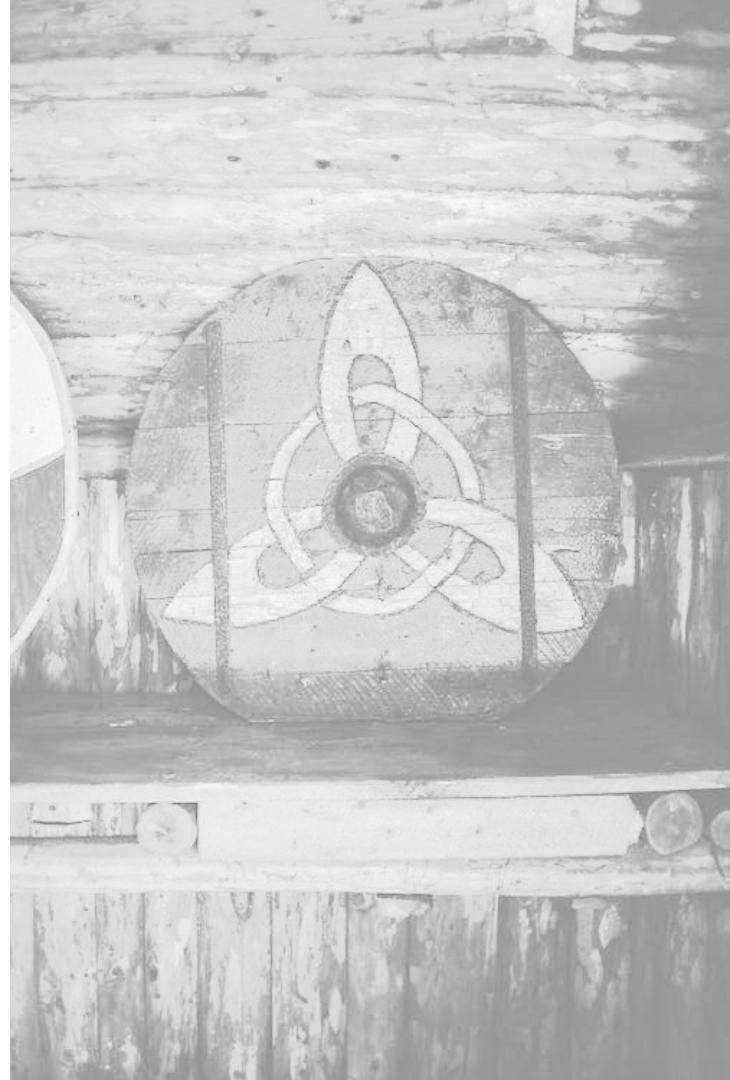
Disjointness Constraints



# libcrux: Typed Rust APIs

```
type Chacha20Key = [u8; 32];
type Nonce = [u8; 12];
type Tag = [u8; 16];

fn encrypt(
    key: &Chacha20Key,
    msg_ctxt: &mut [u8],
    nonce: Nonce,
    aad: &[u8]
) -> Tag
```

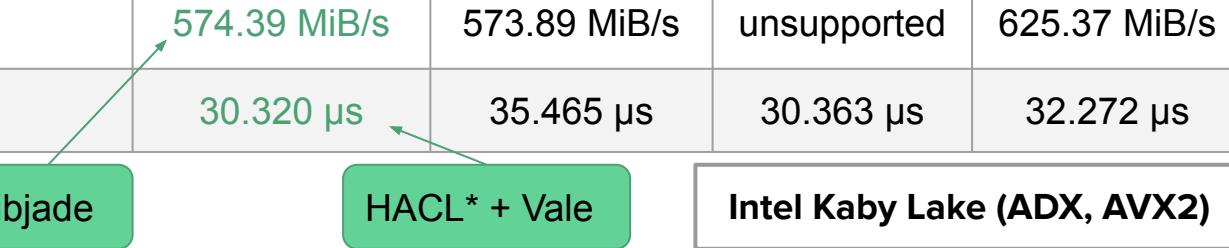


Crypto Standard	Platforms	Specs	Implementations
<b>ECDH</b> <ul style="list-style-type: none"><li>• x25519</li><li>• P256</li></ul>	Portable + Intel ADX Portable	hacspec, F* hacspec, F*	HACL*, Vale HACL*
<b>AEAD</b> <ul style="list-style-type: none"><li>• Chacha20Poly1305</li><li>• AES-GCM</li></ul>	Portable + Intel/ARM SIMD Intel AES-NI	hacspec, F*, EasyCrypt hacspec, F*	HACL*, libjade Vale
<b>Signature</b> <ul style="list-style-type: none"><li>• Ed25519</li><li>• ECDSA P256</li><li>• BLS12-381</li></ul>	Portable Portable Portable	hacspec, F* hacspec, F* hacspec, Coq	HACL* HACL* AUCurves
<b>Hash</b> <ul style="list-style-type: none"><li>• Blake2</li><li>• SHA2</li><li>• SHA3</li></ul>	Portable + Intel/ARM SIMD Portable Portable + Intel SIMD	hacspec, F* hacspec, F* hacspec, F*, EasyCrypt	HACL* HACL* HACL*, libjade
<b>HKDF, HMAC</b>	Portable	hacspec, F*	HACL*
<b>HPKE</b>	Portable	hacspec	hacspec

# libcrux: performance

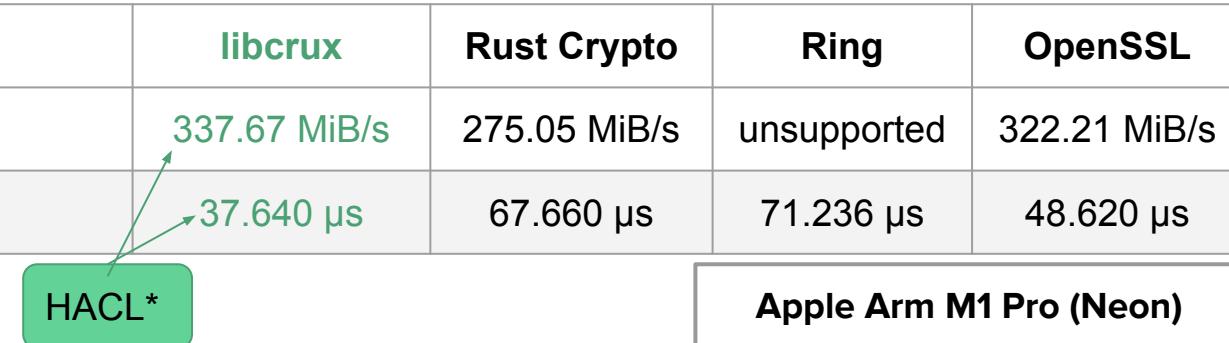
	<b>libcrux</b>	<b>Rust Crypto</b>	<b>Ring</b>	<b>OpenSSL</b>
Sha3 256	574.39 MiB/s	573.89 MiB/s	unsupported	625.37 MiB/s
x25519	30.320 µs	35.465 µs	30.363 µs	32.272 µs

**libjade**      **HACL\* + Vale**      **Intel Kaby Lake (ADX, AVX2)**

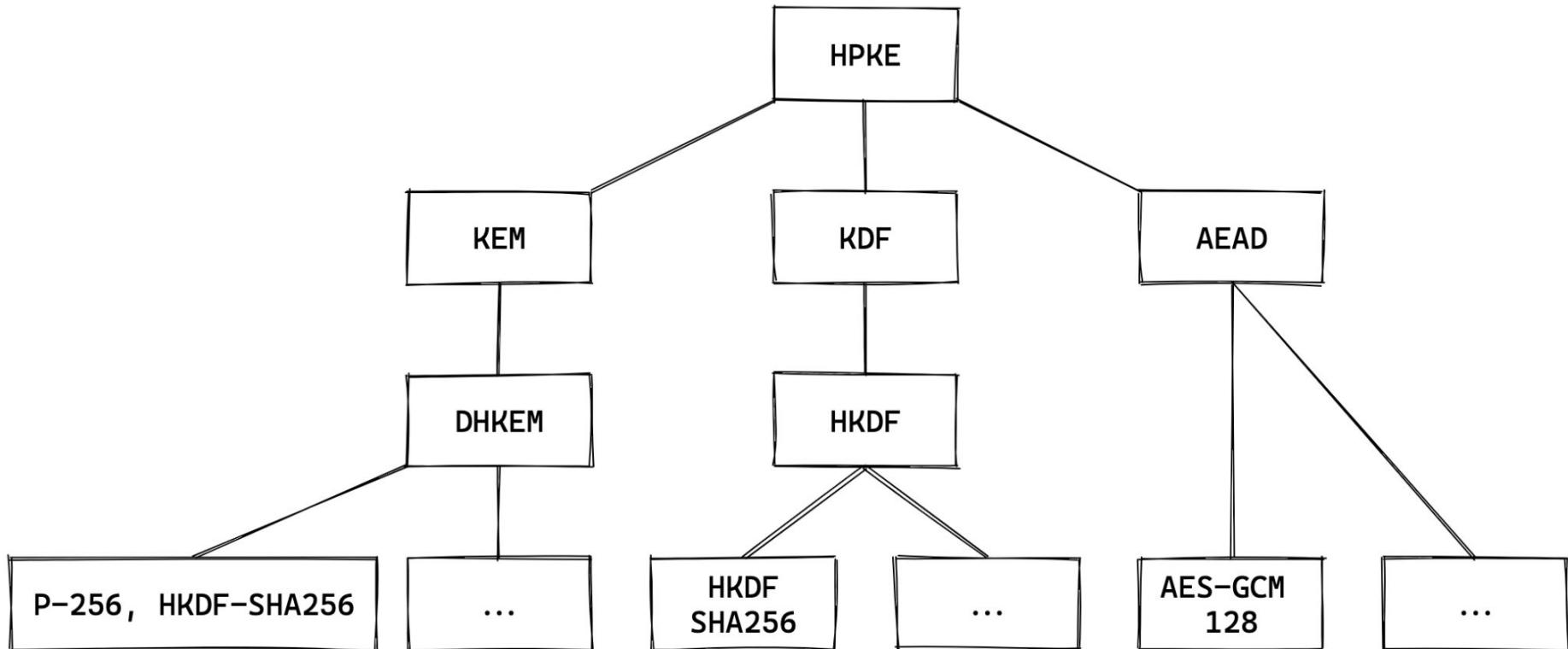


	<b>libcrux</b>	<b>Rust Crypto</b>	<b>Ring</b>	<b>OpenSSL</b>
Sha3 256	337.67 MiB/s	275.05 MiB/s	unsupported	322.21 MiB/s
x25519	37.640 µs	67.660 µs	71.236 µs	48.620 µs

**HACL\***      **Apple Arm M1 Pro (Neon)**



# Building HPKE over libcrux



# Ongoing Work: more proof backends for hacspec

## Security Analysis Tools

- **SSProve**: modular crypto proofs
- **EasyCrypt**: verified constructions
  
- **ProVerif**: symbolic protocol proofs
- **CryptoVerif**: verified protocols
- **Squirrel**: protocol verifier

## Program Verification Tools

- **QuickCheck**: logical spec testing
- **Creusot**: verifying spec contracts
- **Aeneas**: verifying Rust code
  
- **LEAN**: verification framework
- <Your favourite prover here>

# Conclusions

- **Protocol verification** tools are available for analyzing real-world protocols
  - Symbolic analyzers (ProVerif, Tamarin, DY\*)
  - Computational provers (CryptoVerif, EasyCrypt, Squirrel, SSProve)
  - Many case studies (HPKE, MLS, TLS 1.3, Noise, Signal)
- **Fast verified code** is available for most modern crypto algorithms
  - Portable C (HACL\*, Fiat-Crypto), **Assembly** (Vale, libjade, CryptoLine)
  - Ongoing work: PQC, ZKP, FHE, MPC, ...
- **hacspec** is a common spec language for multiple verification tools
  - Try it: [hacspec.org](https://hacspec.org)
- **libcrux** provides safe Rust APIs to multiple verified crypto libraries
  - Try it: [libcrux.org](https://libcrux.org)

# Thanks!

- HACL\*: <https://github.com/hacl-star/hacl-star>
- Vale: <https://github.com/project-everest/vale>
- libjade: <https://github.com/formosa-crypto/libjade>
- AUCurves: <https://github.com/AU-COBRA/AUCurves>
- hacspec: <https://github.com/hacspec/hacspec>
- libcrux: <https://github.com/cryspen/libcrux>

We are hiring R&D crypto/proof engineers at Inria and Cryspen. Get in touch!

