

Week 05 Studio

Orders of Growth,

Data Abstraction

CS1101S AY21/22 Semester 1
Studio 05E

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Admin

- Contact tracing (QR code + class photo)
- “RA1 is so bloody hard!”
 - NO BELL CURVE!!! If you meet criteria for A, you will get the A
 - Prof Martin is “mean” and wanted to challenge everyone
 - 6% only, please focus on mid-terms
- Mission Beyond the First Dimension due 6th September 2359
- Mastery check 1
 - Can PM me and try to book before all of us get busy with mid-terms :(((

Recap

Orders of Growth

What Is It?

- Limited resources for computational processes
 - Time and space
- Asymptotic analysis and notation
 - Provides abstraction / rough measures of resources used with respect to problem size
 - Big Omega (Ω) / Big Theta (Θ) / Big O (O)

Orders of Growth

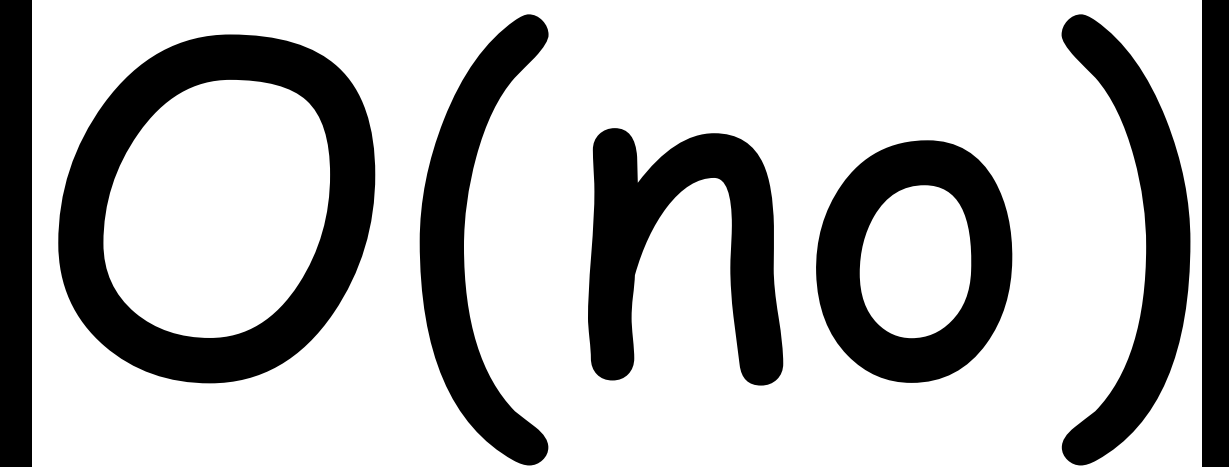
Textbook Definition

- Let n be a parameter that measures the size of the problem, and $R(n)$ be the amount of resources the process requires
- For sufficiently large value of n :
 - $\Omega(g(n))$ if $\forall n > n_0, \exists k \in \mathbb{R}^+$ such that $R(n) \geq k \cdot g(n)$
 - $\Theta(g(n))$ if $\forall n > n_0, \exists k_1, k_2 \in \mathbb{R}^+$ such that $k_1 \cdot g(n) \leq R(n) \leq k_2 \cdot g(n)$
 - $O(g(n))$ if $\forall n > n_0, \exists k \in \mathbb{R}^+$ such that $R(n) \leq k \cdot g(n)$
- Huh?

Orders of Growth

Summary

- Ω : lower bound (best case?)
- Θ : tight bound
- O : upper bound (worst case?)
- Every algorithm technically has $\Omega(1)$ and $O(\infty)$
 - But that does not help with analysis!
 - Often interested in Big O notation
- Will be brought up again in CS2040S, CS3230, CS3233



$O(\text{no})$

Orders of Growth

Things to Take Note

- Constants / Coefficients: disregard
 - $O(n) \equiv O(0.5n)$, $O(100000) \equiv O(1)$
- Addition: drop minor terms
 - $O(n^2 + 3n + 5) \equiv O(n^2)$
- Multiplication: take product
 - $O(n \cdot \log(n) \cdot n) \equiv O(n^2 \cdot \log n)$

When you watch the video
at 2x speed to save time
but it is still $O(n)$

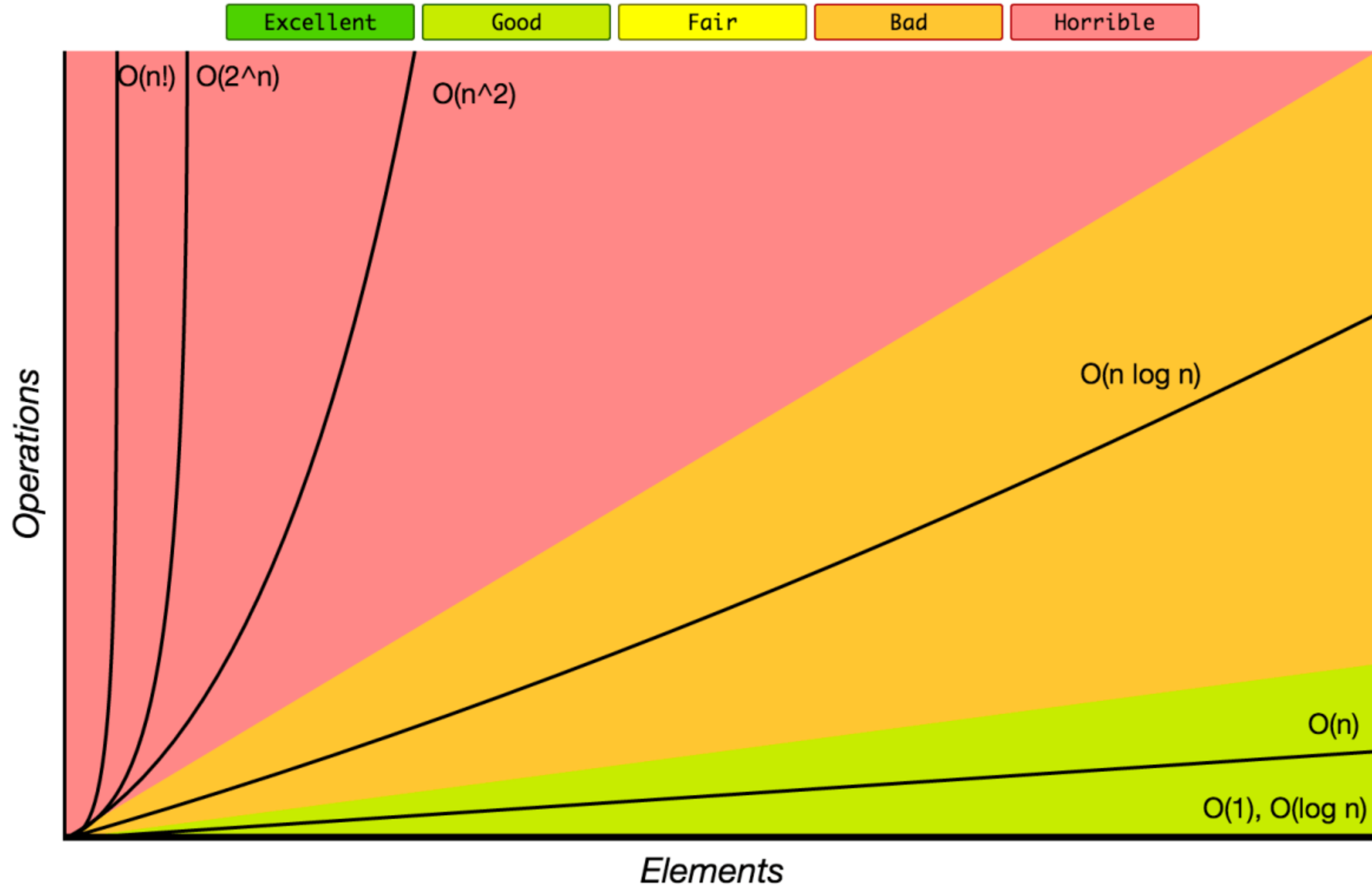


Orders of Growth

Common Functions

Notation	$O(1)$	$O(\log n)$	$O(n)$	$O(n \log n)$	$O(n^2)$	$O(n^c)$	$O(c^n)$	$O(n!)$
Name	Constant	Logarithmic	Linear	Log-linear	Quadratic	Polynomial	Exponential	Factorial
e.g.	Look-up table	fast_expt, binary search	expt	Comparison-based sorting	Permutation, combination	Permutation, combination	Travelling salesman problem with DP	Travelling salesman problem with brute force

Big-O Complexity Chart



Data Abstraction - Pairs

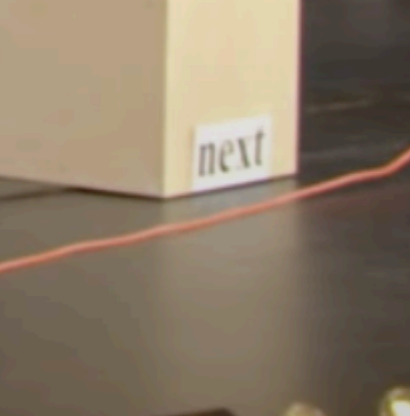
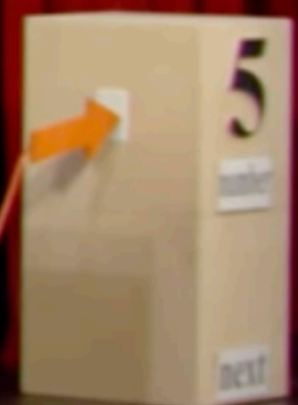
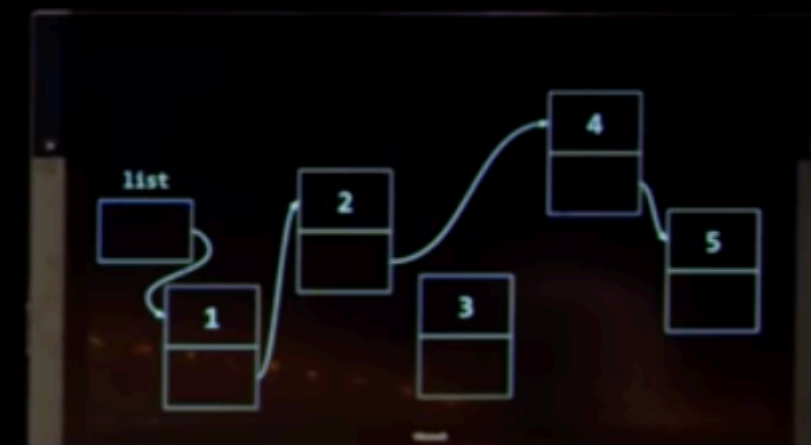
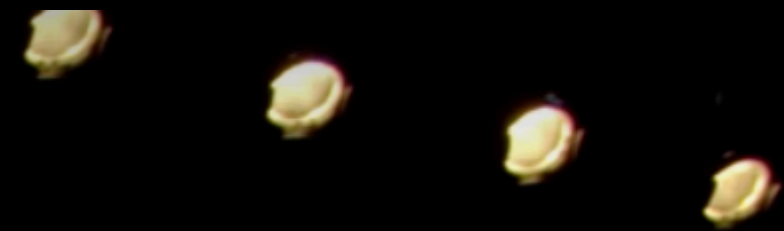
What Is It?

- A kind of data structure
- Does this seem familiar?
 - `function pair(x, y) { return f => f(x, y); }`
 - `function head(p) { return p((x, y) => x); }`
 - `function tail(p) { p((x, y) => y); }`
- Rational number

Data Abstraction - Lists

What Is It?

- Another kind of data structure
- “List” of things (number, string, boolean, pair, or even list itself)
 - `list(123, "hello", true, list(1, 2, 3),
pair("head", "tail"));`
- Linked list in C
- Python list, Java Array/ArrayList are not exactly a list



Data Abstraction - Lists

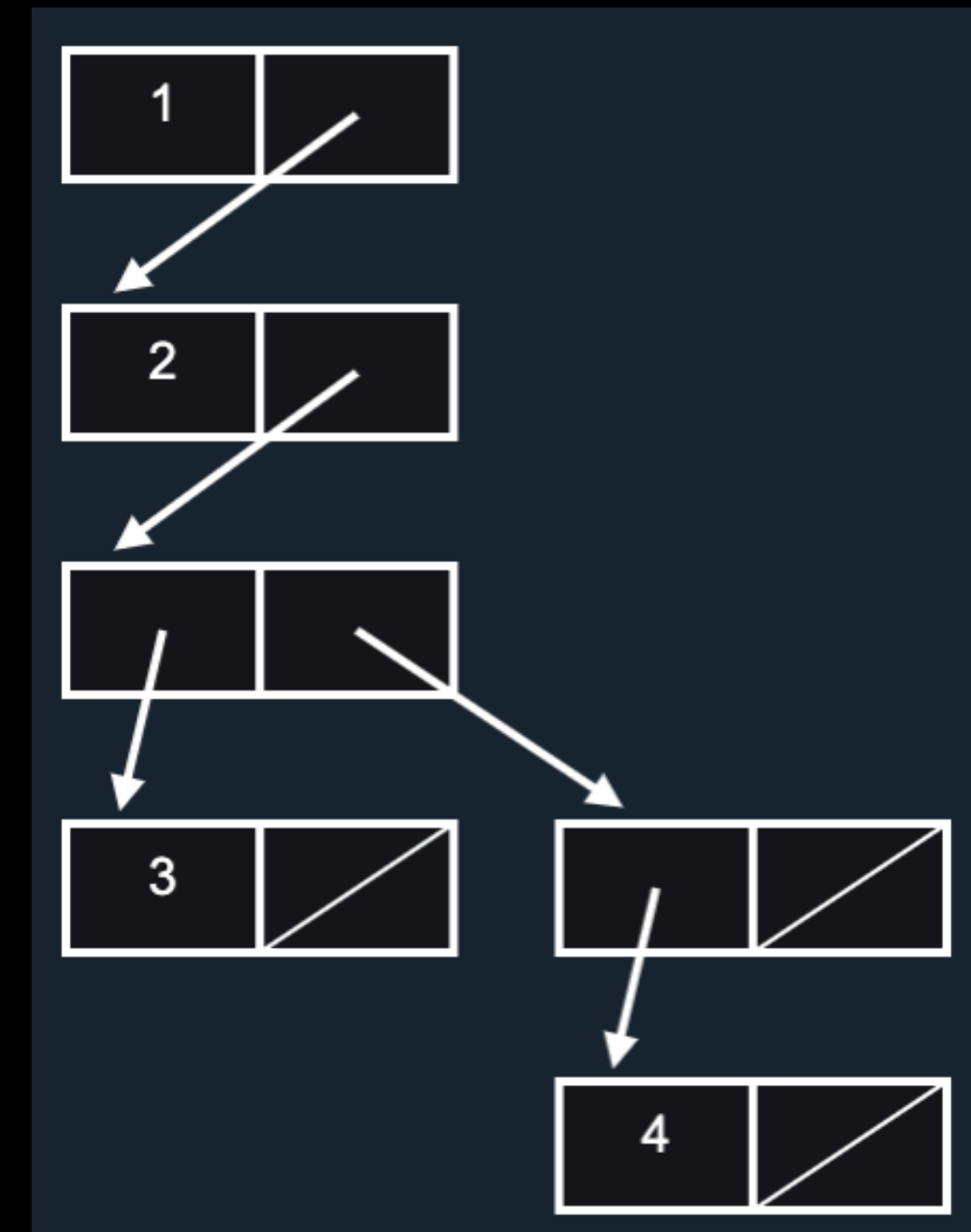
Definition

- Recursive in nature
 - A list is either null or a pair whose tail is a list
 - A list is either null or a pair whose tail is (either null or a pair whose tail is a list)
 - A list is either null or a pair whose tail is (either null or a pair whose tail is (either null or a pair whose tail is a list))
 - ... you get it!

Data Abstraction - Lists

Notations

- `list(1, 2, list(3), pair(4, null));`
 - Box notation: `[1, [2, [[3, null], [[4, null], null]]]]`
 - `list_to_string`
 - List notation: `list(1, 2, list(3), list(4))`
 - `display_list` (display lists nicely if detected)
 - Box-and-pointer diagram: slash for `null`
 - `draw_data`
- `[null, null]` a pair or list?



Data Abstraction - Lists

Length

- Count the number of items in the list
- Nested list?
- Recursive / iterative?
- Head over to Source Academy!

Any Questions?