

Accelerating Array Constraints in Symbolic Execution

David Perry^{*1}, Andrea Mattavelli^{*2}, Xiangyu Zhang¹, Cristian Cadar²

¹Department of Computer Science, Purdue University

²Department of Computing, Imperial College London

*The first two authors contributed equally to this paper



What is Symbolic Execution?

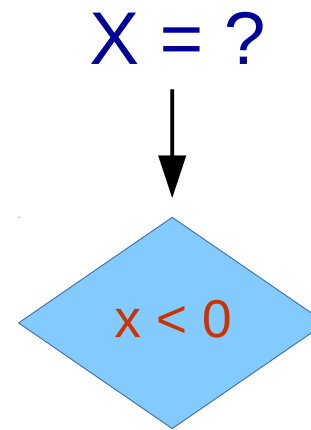
```
int bad_abs(int x) {  
    if(x < 0)  
        return -x;  
    if(x == 1234)  
        return -x;  
    return x;  
}
```

What is Symbolic Execution?

```
int bad_abs(int x) {                                X = ?  
    if(x < 0)  
        return -x;  
    if(x == 1234)  
        return -x;  
    return x;  
}
```

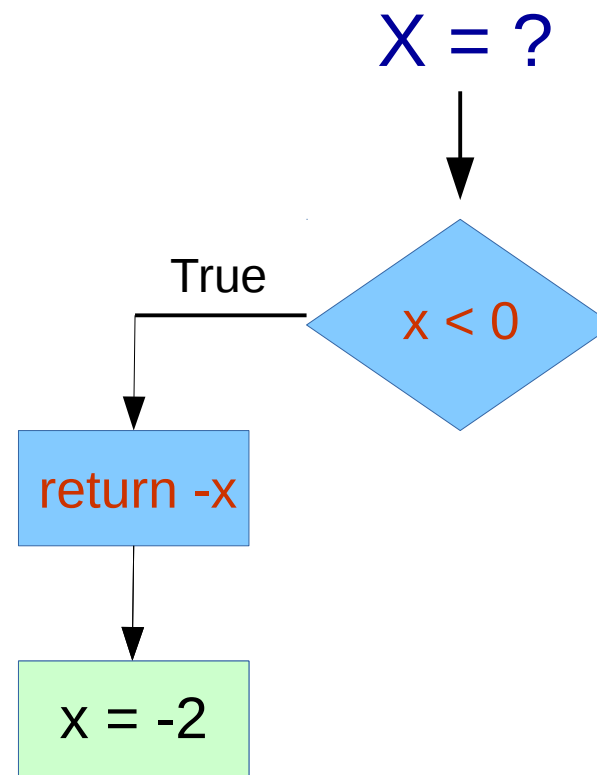
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What is Symbolic Execution?

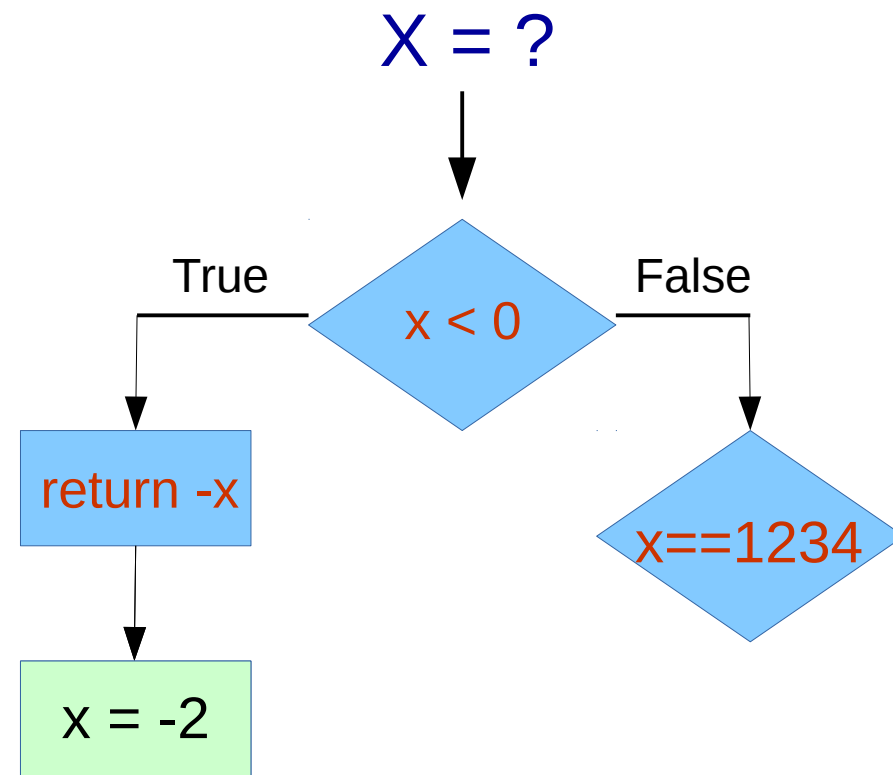
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}
```



Test1.out

What is Symbolic Execution?

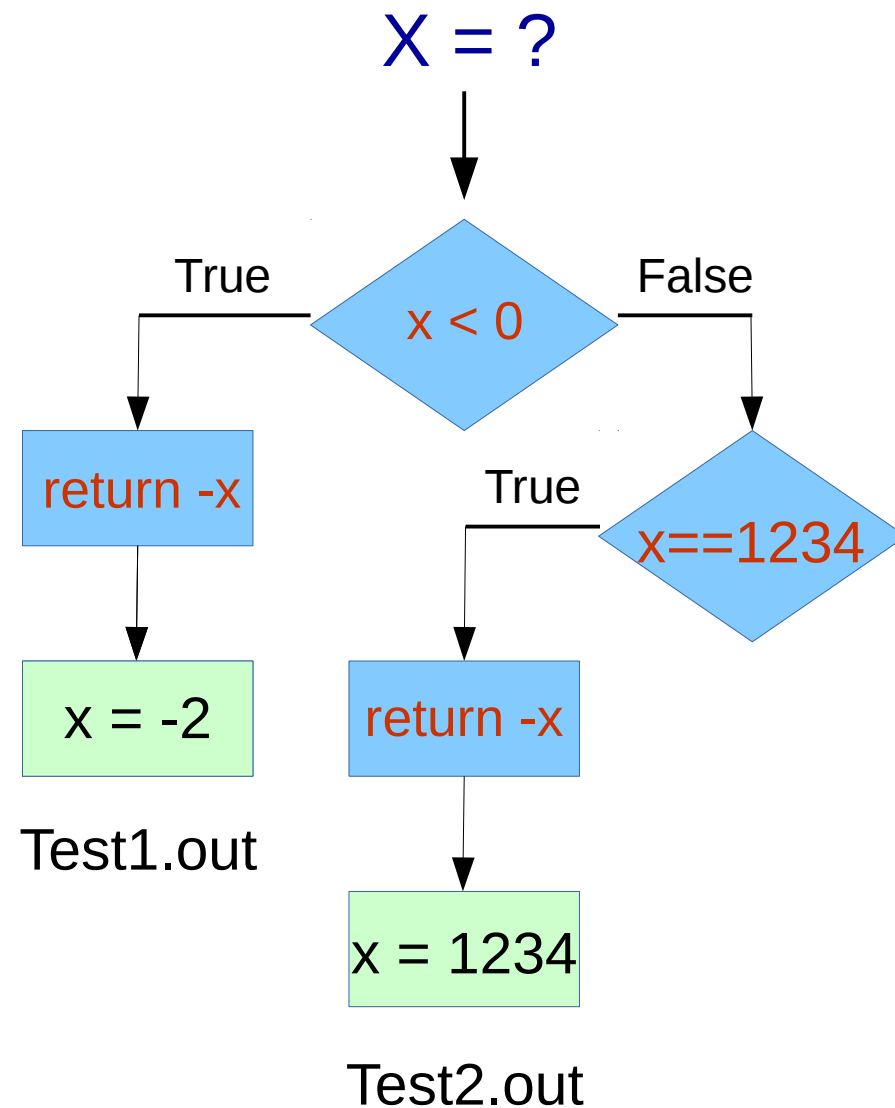
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    if(x == 1234)  
        return -x;  
    return x;  
}
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Test1.out

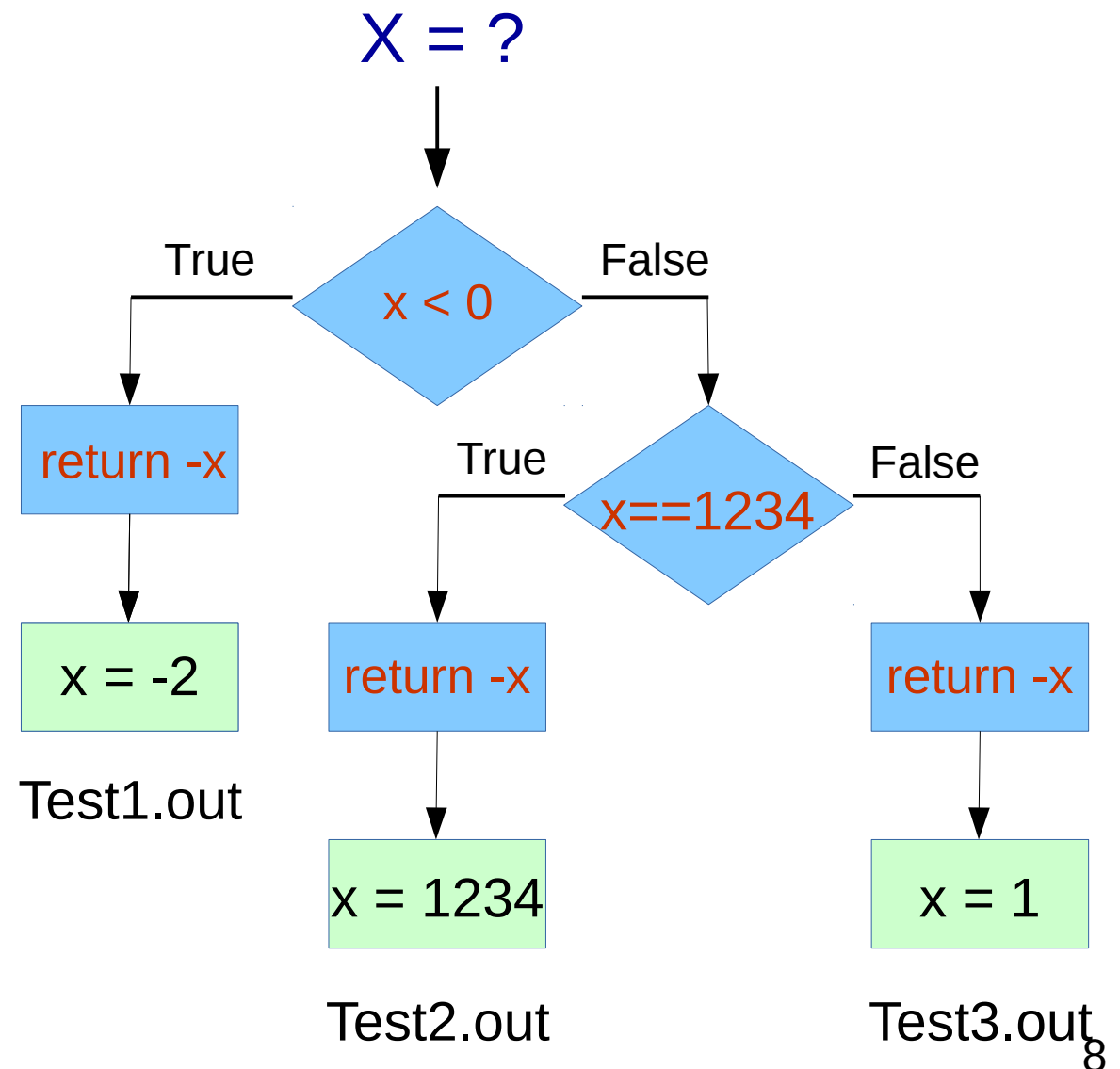
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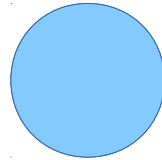
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    if(x < 0)  
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    if(x == 1234)  
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    return x;  
}
```

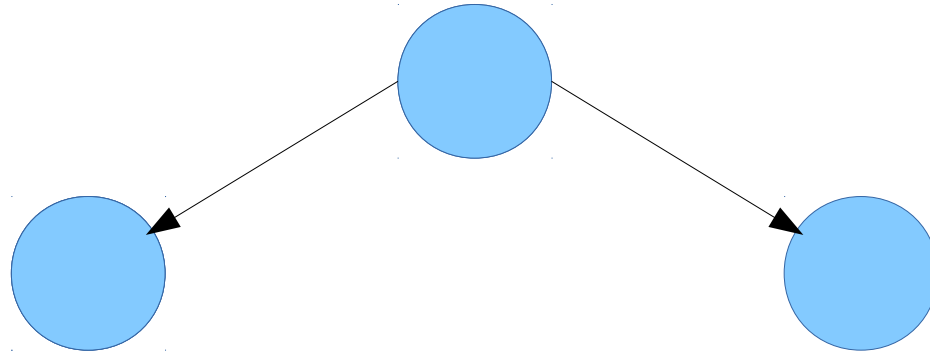


What are the technique's
limitations?

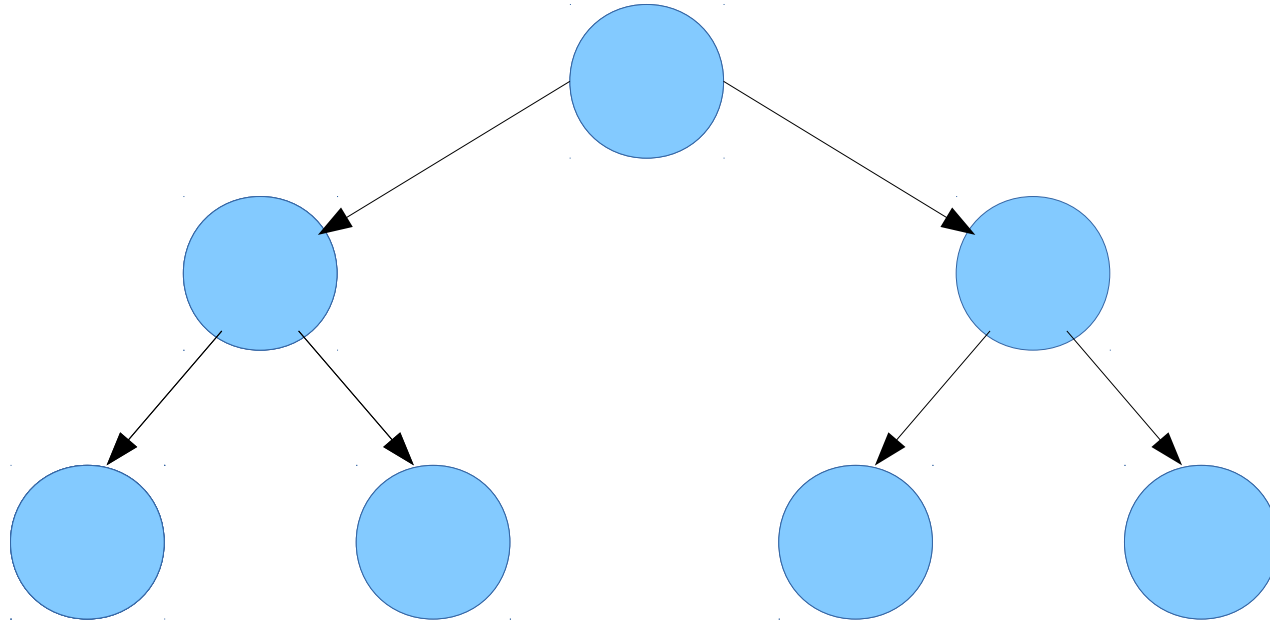
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What are the technique's limitations?

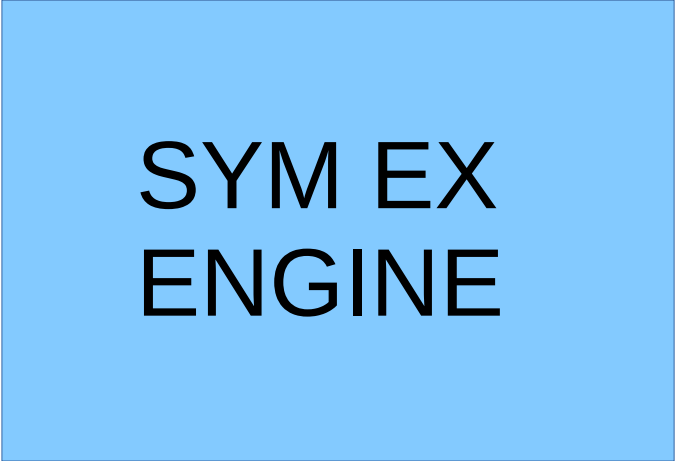


What are the technique's limitations?



PATH EXPLOSION!

What are the technique's limitations?



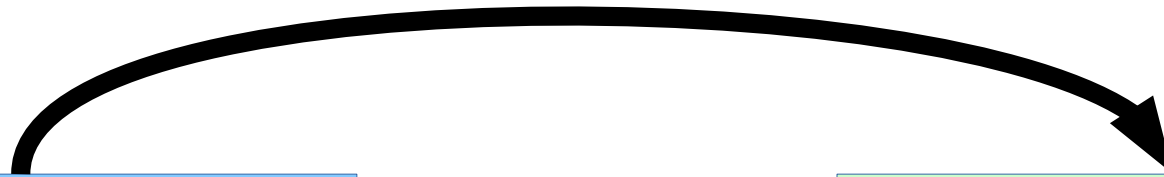
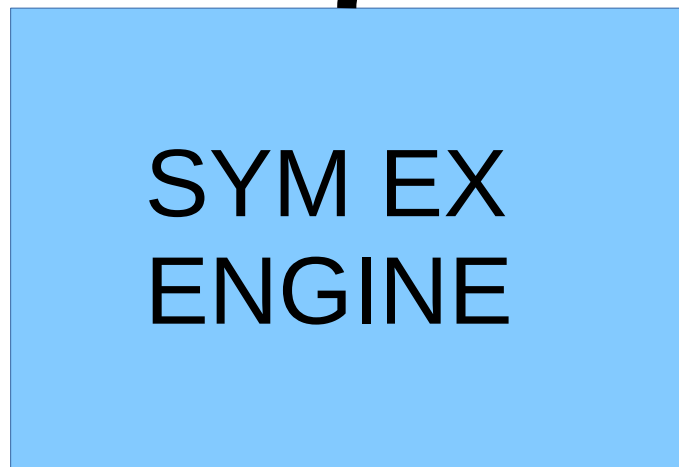
**SYM EX
ENGINE**



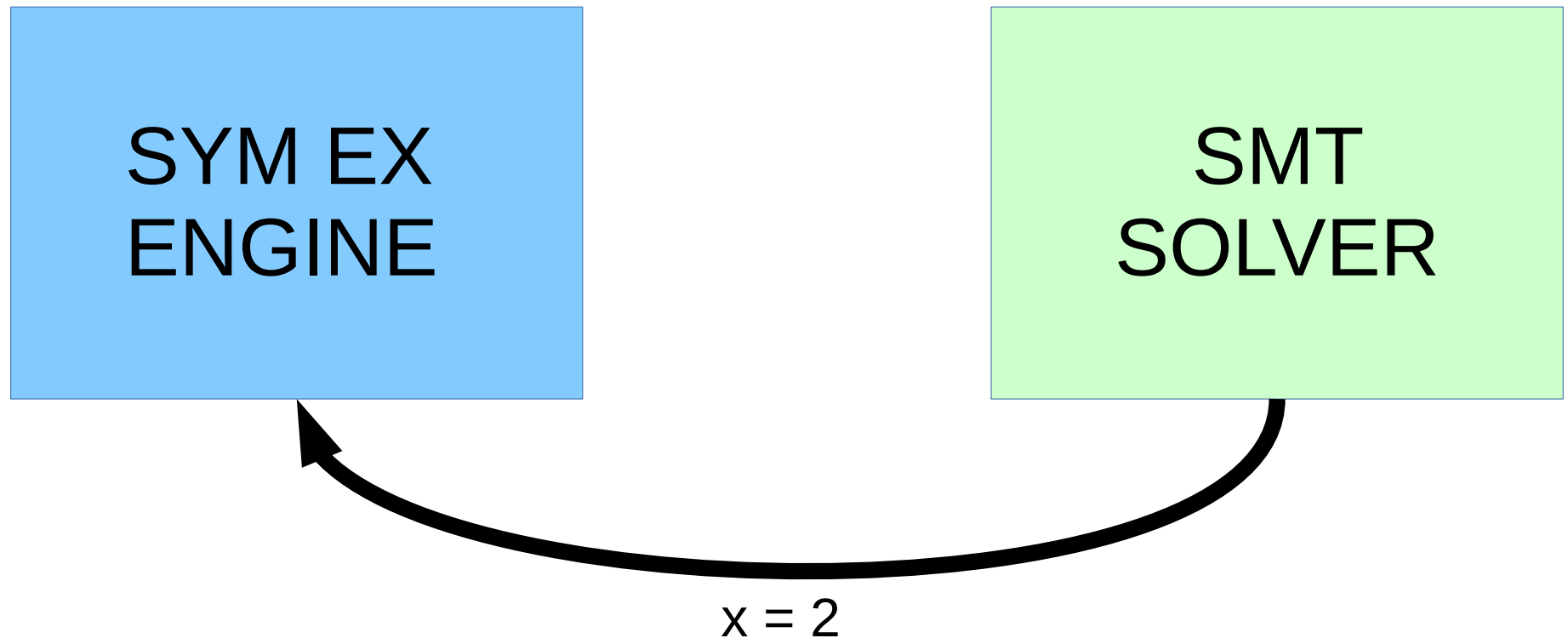
**SMT
SOLVER**

What are the technique's limitations?

$X > 0,$
 $X \neq 1234$

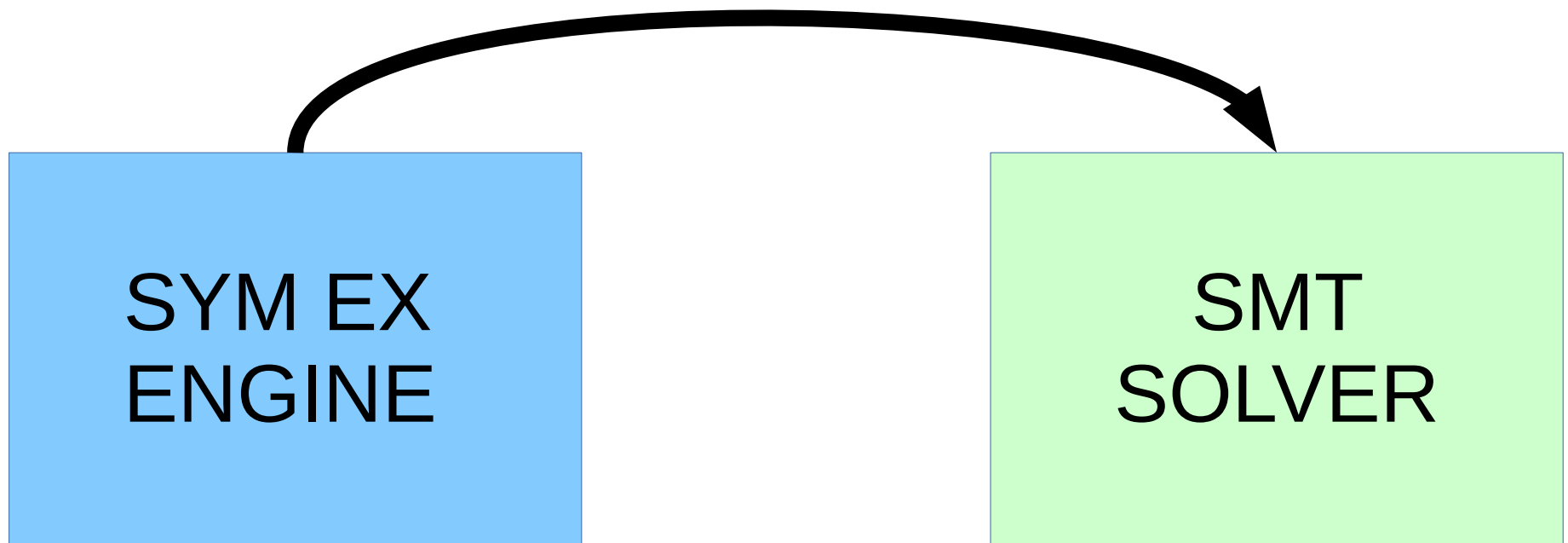


What are the technique's limitations?



What are the technique's limitations?

`largeArray[symIdx] != symVar`



What are the technique's limitations?

```
largeArray[sym_idx] != sym_var
```



SOLVER TIMEOUT!

A Challenging Example: BC

```
Void tokenMatch(char *input) {  
    Unsigned currState = 0;  
    Char charPtr = input;  
    Do {  
        Char currClass = equivClass[*charPtr];  
        if(accept[currState]) {  
            LastAcceptState = currState;  
            LastAcceptPos = charPtr;  
        }  
        while(check[base[currState] + currClass] != currState) {  
            CurrState = def[currState];  
            if(currState >= 298) {CurrClass = meta[currClass];}  
        }  
        CurrState = next[base[currState] + currClass]; charPtr++;  
    } while(base[currState] != 526);  
}
```

A Challenging Example: BC

```
Void tokenMatch(char *input) {  
    Unsigned currState = 0;  
    Char charPtr = input;  
    Do {  
        Char currClass = equivClass[*charPtr];  
        if(accept[currState]) {  
            LastAcceptState = currState;  
            LastAcceptPos = charPtr;  
        }  
        while(check[base[currState] + currClass] != currState)  
            currState = def[currState],  
            if(currState >= 298) {CurrClass = meta[currClass];}  
        }  
        CurrState = next[base[currState] + currClass]; charPtr++;  
    } while(base[currState] != 526);  
}
```

How does Symbolic Execution handle array reads?

```
int small[5] = {0,3,2,2,2};  
if(small[sym_idx] == 2) {...}
```

How does Symbolic Execution handle array reads?

```
int small[5] = {0,3,2,2,2};
```

```
if(small[sym_idx] == 2) {...}
```

- **Create a variable and assign its value for each offset**

- $\text{small}_0 = 0 \wedge \text{small}_1 = 3 \wedge \text{small}_2 = 2 \wedge \text{small}_3 = 3 \wedge \text{small}_4 = 2$

How does Symbolic Execution handle array reads?

```
int small[5] = {0,3,2,2,2};
```

```
if(small[sym_idx] == 2) {...}
```

- Create a variable and assign its value for each offset
 - $\text{small}_0 = 0 \wedge \text{small}_1 = 3 \wedge \text{small}_2 = 2 \wedge \text{small}_3 = 3 \wedge \text{small}_4 = 2$
- **Add a constraint for the conditional**
 - $\text{val} == \text{small}[\text{sym_idx}] == 2$

How does Symbolic Execution handle array reads?

```
int small[5] = {0,3,2,2,2};  
if(small[sym_idx] == 2) {...}
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- Create a variable and assign its value for each offset
 - $\text{small_0} = 0 \wedge \text{small_1} = 3 \wedge \text{small_2} = 2 \wedge \text{small_3} = 3 \wedge \text{small_4} = 2$
- Add a constraint for the conditional
 - $\text{val} == \text{small}[\text{sym_idx}] == 2$
- **Handle the read operation**
 - $\text{sym_idx} = 0 \rightarrow \text{val} = \text{small_0} \wedge \text{sym_idx} = 1 \rightarrow \text{val} = \text{small_1} \wedge$
 $\text{sym_idx} = 2 \rightarrow \text{val} = \text{small_2} \wedge \text{sym_idx} = 3 \rightarrow \text{val} = \text{small_3} \wedge$
 $\text{sym_idx} = 4 \rightarrow \text{val} = \text{small_4}$

How does Symbolic Execution handle array reads?

[illegible]

```
unsigned isBase64(unsigned k) {  
    if(k > 255)  
        return -1;  
    if(b64[k] >= 0)  
        return 1;  
    else return 0;  
}
```

How does Symbolic Execution handle array reads?

[illegible]

```
unsigned isBase64(unsigned k) {
    if(k > 255)
        return 1;
    if(b64[k] >= 0)
        return 1;
    else return 0;
}
```

How does Symbolic Execution handle array reads?

```
b64[0]=-1 ∧ b64[1]=-1 ∧ b64[2]=-1 ∧ b64[3]=-1 ∧ b64[4]=-1 ∧ b64[5]=-1 ∧ b64[6]=-1 ∧ b64[7]=-1 ∧ b64[8]=-1 ∧ b64[9]=-1 ∧ b64[10]=-1 ∧  
b64[11]=-1 ∧ b64[12]=-1 ∧ b64[13]=-1 ∧ b64[14]=-1 ∧ b64[15]=-1 ∧ b64[16]=-1 ∧ b64[17]=-1 ∧ b64[18]=-1 ∧ b64[19]=-1 ∧ b64[20]=-1 ∧ b64[21]=-  
1 ∧ b64[22]=-1 ∧ b64[23]=-1 ∧ b64[24]=-1 ∧ b64[25]=-1 ∧ b64[26]=-1 ∧ b64[27]=-1 ∧ b64[28]=-1 ∧ b64[29]=-1 ∧ b64[30]=-1 ∧ b64[31]=-1 ∧  
b64[32]=-1 ∧ b64[33]=-1 ∧ b64[34]=-1 ∧ b64[35]=-1 ∧ b64[36]=-1 ∧ b64[37]=-1 ∧ b64[38]=-1 ∧ b64[39]=-1 ∧ b64[40]=-1 ∧ b64[41]=-1 ∧ b64[42]=-  
1 ∧ b64[43]=62 ∧ b64[44]=-1 ∧ b64[45]=-1 ∧ b64[46]=-1 ∧ b64[47]=63 ∧ b64[48]=52 ∧ b64[49]=53 ∧ b64[50]=54 ∧ b64[51]=55 ∧ b64[52]=56 ∧  
b64[53]=57 ∧ b64[54]=58 ∧ b64[55]=59 ∧ b64[56]=60 ∧ b64[57]=61 ∧ b64[58]=-1 ∧ b64[59]=-1 ∧ b64[60]=-1 ∧ b64[61]=-1 ∧ b64[62]=-1 ∧  
b64[63]=-1 ∧ b64[64]=-1 ∧ b64[65]=0 ∧ b64[66]=1 ∧ b64[67]=2 ∧ b64[68]=3 ∧ b64[69]=4 ∧ b64[70]=5 ∧ b64[71]=6 ∧ b64[72]=7 ∧ b64[73]=8 ∧  
b64[74]=9 ∧ b64[75]=10 ∧ b64[76]=11 ∧ b64[77]=12 ∧ b64[78]=13 ∧ b64[79]=14 ∧ b64[80]=15 ∧ b64[81]=16 ∧ b64[82]=17 ∧ b64[83]=18 ∧  
b64[84]=19 ∧ b64[85]=20 ∧ b64[86]=21 ∧ b64[87]=22 ∧ b64[88]=23 ∧ b64[89]=24 ∧ b64[90]=25 ∧ b64[91]=-1 ∧ b64[92]=-1 ∧ b64[93]=-1 ∧  
b64[94]=-1 ∧ b64[95]=-1 ∧ b64[96]=-1 ∧ b64[97]=26 ∧ b64[98]=27 ∧ b64[99]=28 ∧ b64[100]=29 ∧ b64[101]=30 ∧ b64[102]=31 ∧ b64[103]=32 ∧  
b64[104]=33 ∧ b64[105]=34 ∧ b64[106]=35 ∧ b64[107]=36 ∧ b64[108]=37 ∧ b64[109]=38 ∧ b64[110]=39 ∧ b64[111]=40 ∧ b64[112]=41 ∧  
b64[113]=42 ∧ b64[114]=43 ∧ b64[115]=44 ∧ b64[116]=45 ∧ b64[117]=46 ∧ b64[118]=47 ∧ b64[119]=48 ∧ b64[120]=49 ∧ b64[121]=50 ∧  
b64[122]=51 ∧ b64[123]=-1 ∧ b64[124]=-1 ∧ b64[125]=-1 ∧ b64[126]=-1 ∧ b64[127]=-1 ∧ b64[128]=-1 ∧ b64[129]=-1 ∧ b64[130]=-1 ∧ b64[131]=-  
1 ∧ b64[132]=-1 ∧ b64[133]=-1 ∧ b64[134]=-1 ∧ b64[135]=-1 ∧ b64[136]=-1 ∧ b64[137]=-1 ∧ b64[138]=-1 ∧ b64[139]=-1 ∧ b64[140]=-1 ∧  
b64[141]=-1 ∧ b64[142]=-1 ∧ b64[143]=-1 ∧ b64[144]=-1 ∧ b64[145]=-1 ∧ b64[146]=-1 ∧ b64[147]=-1 ∧ b64[148]=-1 ∧ b64[149]=-1 ∧ b64[150]=-  
1 ∧ b64[151]=-1 ∧ b64[152]=-1 ∧ b64[153]=-1 ∧ b64[154]=-1 ∧ b64[155]=-1 ∧ b64[156]=-1 ∧ b64[157]=-1 ∧ b64[158]=-1 ∧ b64[159]=-1 ∧  
b64[160]=-1 ∧ b64[161]=-1 ∧ b64[162]=-1 ∧ b64[163]=-1 ∧ b64[164]=-1 ∧ b64[165]=-1 ∧ b64[166]=-1 ∧ b64[167]=-1 ∧ b64[168]=-1 ∧ b64[169]=-  
1 ∧ b64[170]=-1 ∧ b64[171]=-1 ∧ b64[172]=-1 ∧ b64[173]=-1 ∧ b64[174]=-1 ∧ b64[175]=-1 ∧ b64[176]=-1 ∧ b64[177]=-1 ∧ b64[178]=-1 ∧  
b64[179]=-1 ∧ b64[180]=-1 ∧ b64[181]=-1 ∧ b64[182]=-1 ∧ b64[183]=-1 ∧ b64[184]=-1 ∧ b64[185]=-1 ∧ b64[186]=-1 ∧ b64[187]=-1 ∧ b64[188]=-  
1 ∧ b64[189]=-1 ∧ b64[190]=-1 ∧ b64[191]=-1 ∧ b64[192]=-1 ∧ b64[193]=-1 ∧ b64[194]=-1 ∧ b64[195]=-1 ∧ b64[196]=-1 ∧ b64[197]=-1 ∧  
b64[198]=-1 ∧ b64[199]=-1 ∧ b64[200]=-1 ∧ b64[201]=-1 ∧ b64[202]=-1 ∧ b64[203]=-1 ∧ b64[204]=-1 ∧ b64[205]=-1 ∧ b64[206]=-1 ∧ b64[207]=-  
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1 ∧ b64[246]=-1 ∧ b64[247]=-1 ∧ b64[248]=-1 ∧ b64[249]=-1 ∧ b64[250]=-1 ∧ b64[251]=-1 ∧ b64[252]=-1 ∧ b64[253]=-1 ∧ b64[254]=-1 ∧  
b64[255]=-1 ∧
```

How does Symbolic Execution handle array reads?

$$B64k \geq 0 \wedge k \leq 255$$

How does Symbolic Execution handle array reads?

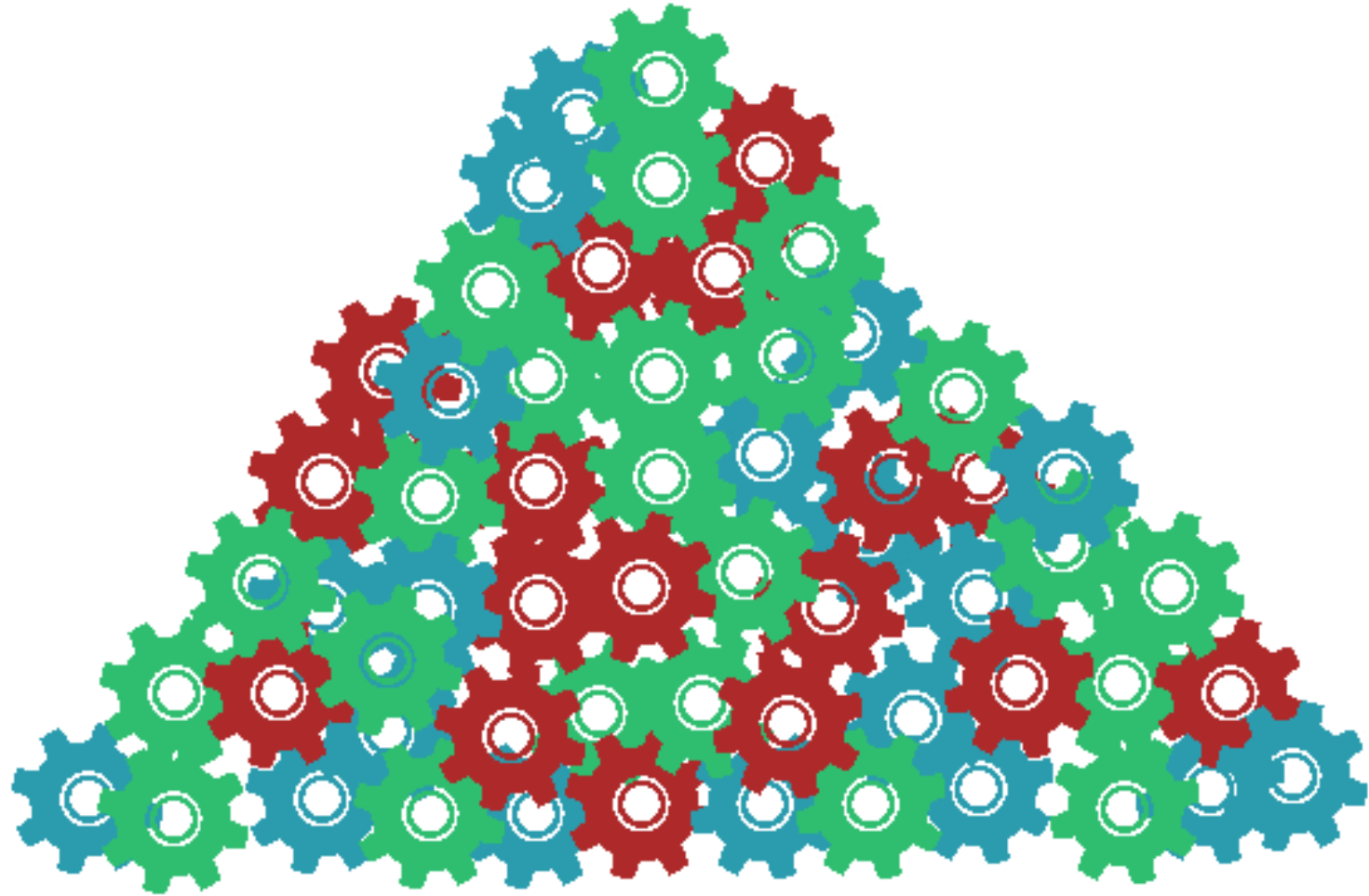
[illegible]

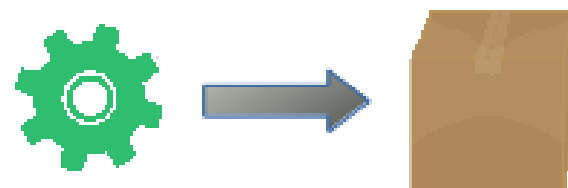
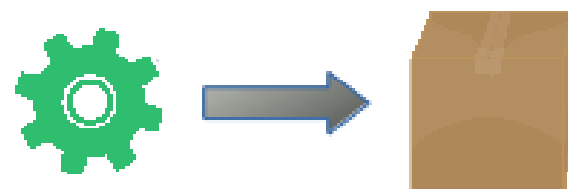
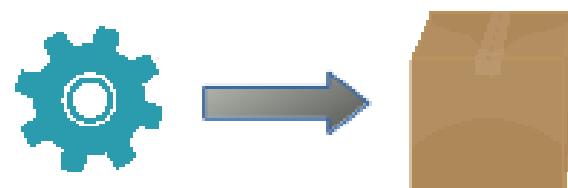
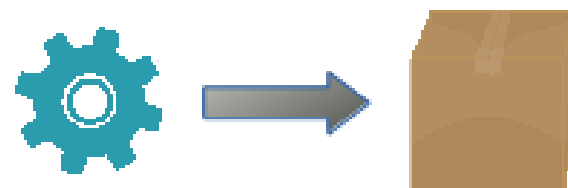
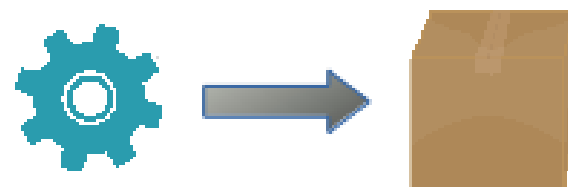
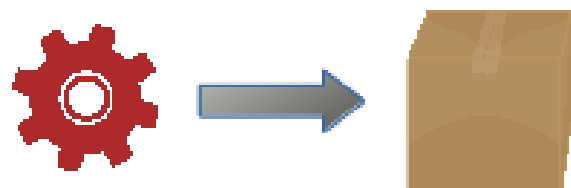
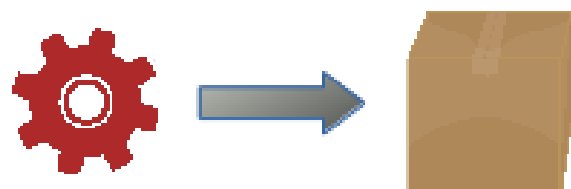
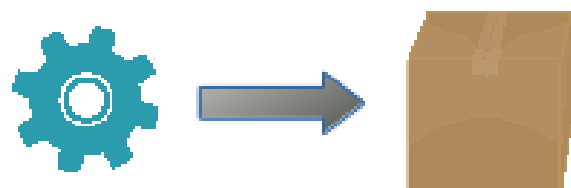
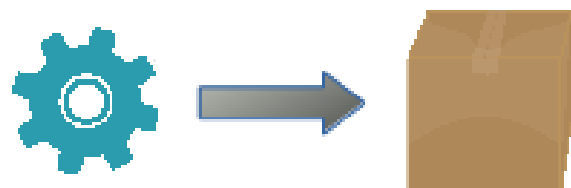
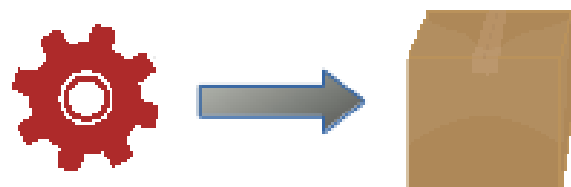
$$b_{64k} \geq 0 \wedge$$

30

Impact on the Solver

- This explosion of variables creates an incredibly large search space
- Solver wastes time exploring multiple solutions with different indexes that point to the same array value
- Complicated queries make the performance even worse as backtracking results in more redundant exploration







Our Technique

- Takes advantage of statically and dynamically available information about the target array
- Applies one of two optimizations depending on how the array read is used
- Yields a simplified query that results in dramatic performance improvement in programs that read from large arrays with symbolic indexes

Index-Based Transformation

- Only applicable to conditionals comparing array reads with statically known values
- Turns the comparison operation into a conjunctive formula that compares index ranges without the explicit array read
- Avoids the overhead of representing values in the array you don't need
- Creates the most concise encoding of either of our transformations

Index-Based Transformation

[illegible]

| Target Program | Original Constraints | Opt Constraints |
|---|----------------------|-----------------|
| <pre>isBase64(unsigned k) { if(k > 255) return -1; if(b64[k] >= 0) return 1; else return 0; }</pre> | | |

Index-Based Transformation

[illegible]

| Target Program | Original Constraints | Opt Constraints |
|---|--|-----------------|
| <pre>isBase64(unsigned k) { if(k > 255) return -1; if(b64[k] >= 0) return 1; else return 0; }</pre> | <pre>// Array Variables b64[0]=-1 ∧ b64[1]=-1 ∧ ... ∧ //Array Conditional b64k >= 0 ∧ // Read Operation k=0 → b64k=b64[0] ∧ k=1 → b64k=b64[1] ∧ ...</pre> | |

Index-Based Transformation

[illegible]

| Target Program | Original Constraints | Opt Constraints |
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| <pre>isBase64(unsigned k) { if(k > 255) return -1; if(b64[k] >= 0) return 1; else return 0; }</pre> | <pre>// Array Variables b64[0]=-1 ∧ b64[1]=-1 ∧ ... ∧ //Array Conditional b64k >= 0 ∧ // Read Operation k=0 → b64k=b64[0] ∧ k=1 → b64k=b64[1] ∧ ...</pre> | <div style="border: 2px solid red; padding: 5px; display: inline-block;"> k = 43 v </div> |

Index-Based Transformation

[illegible]

| Target Program | Original Constraints | Opt Constraints |
|---|--|---|
| <pre>isBase64(unsigned k) { if(k > 255) return -1; if(b64[k] >= 0) return 1; else return 0; }</pre> | <pre>// Array Variables b64[0]=-1 ∧ b64[1]=-1 ∧ ... ∧ //Array Conditional b64k >= 0 ∧ // Read Operation k=0 → b64k=b64[0] ∧ k=1 → b64k=b64[1] ∧ ...</pre> | <div style="border: 2px solid red; padding: 5px; margin-bottom: 10px;">k = 43 v</div> <div style="border: 2px solid green; padding: 5px;">47 <= k <= 57 v</div> |

Index-Based Transformation

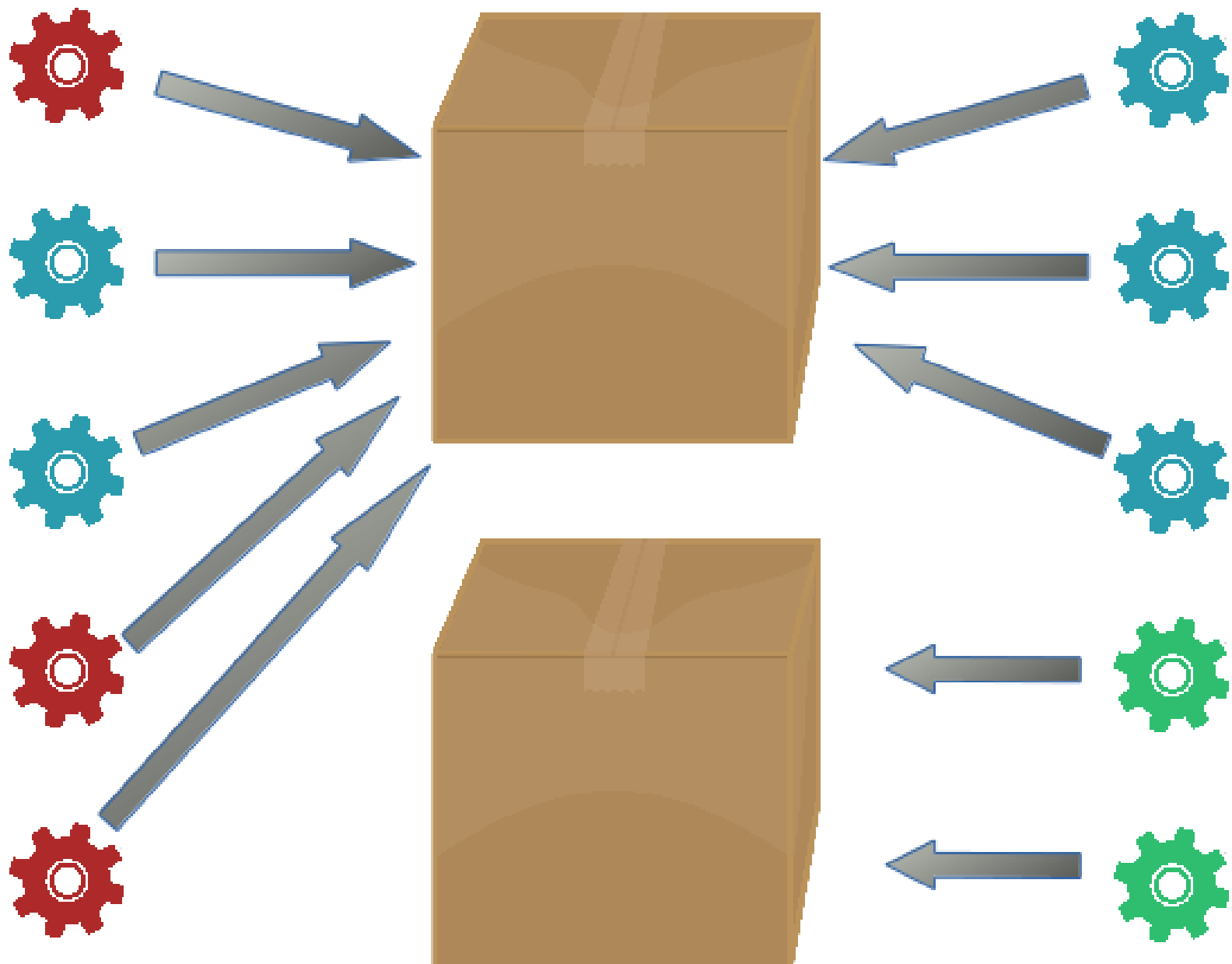
[illegible]

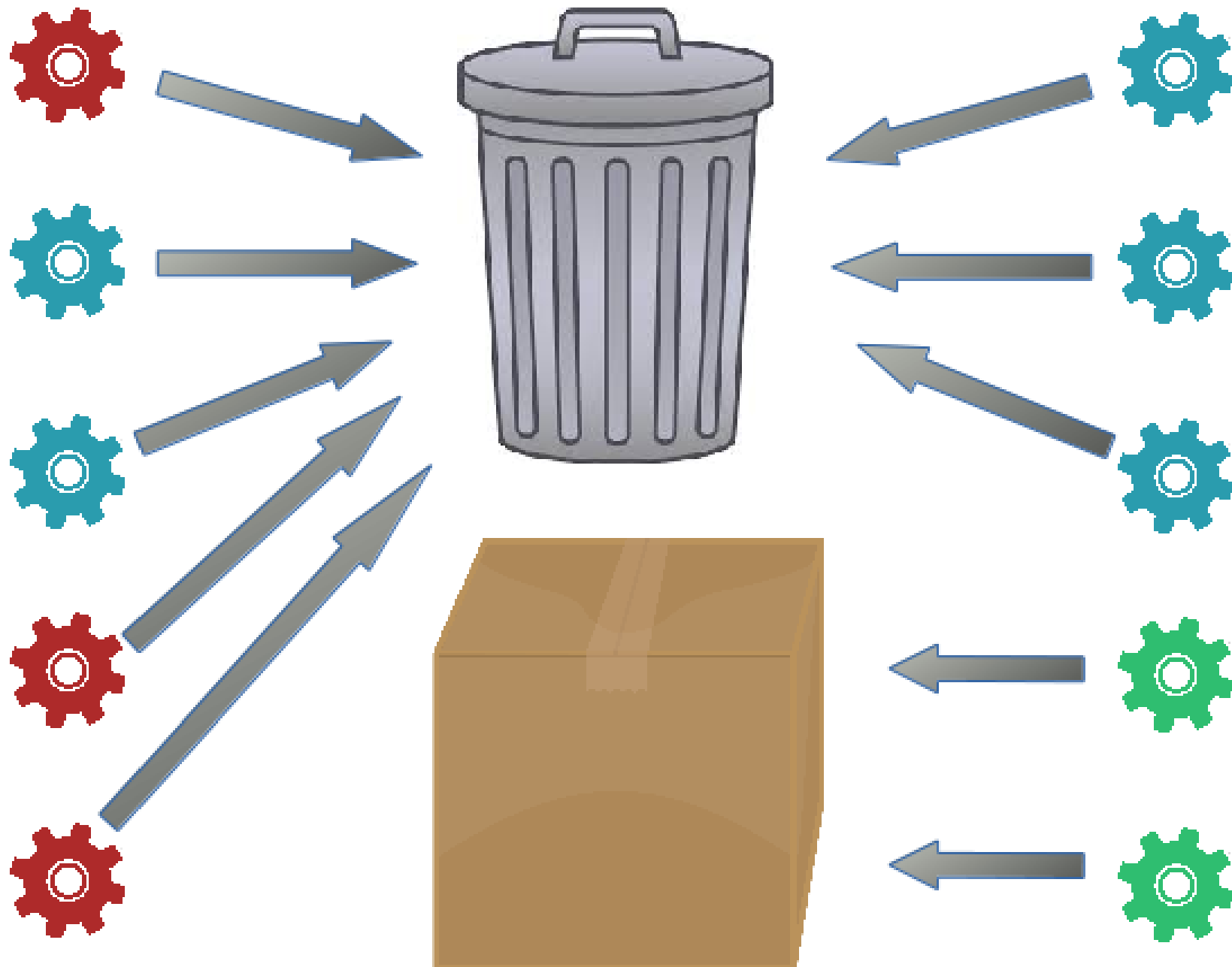
| Target Program | Original Constraints | Opt Constraints |
|---|--|---|
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Index-Based Transformation

[illegible]

| Target Program | Original Constraints | Opt Constraints |
|---|--|--|
| <pre>isBase64(unsigned k) { if(k > 255) return -1; if(b64[k] >= 0) return 1; else return 0; }</pre> | <pre>// Array Variables b64[0]=-1 ∧ b64[1]=-1 ∧ ... ∧ //Array Conditional b64k >= 0 ∧ // Read Operation k=0 → b64k=b64[0] ∧ k=1 → b64k=b64[1] ∧ ...</pre> | <pre>k = 43 v 47 <= k <= 57 v 65 <= k <= 90 97 <= k <= 122</pre> |





Value-Based Transformation

- Applicable to all array reads
- Turns the array read into a series of nested ITE's comparing the index to a range of offsets that correspond to the same value
- Removes the redundancy created by the traditional Theory of Arrays

Value-Based Transformation

[illegible]

| Target Program | Original Constraints | Opt Constraints |
|---|----------------------|-----------------|
| <pre>isBase64(unsigned k) { if(k > 255) return -1; if(b64[k] >= sym) return 1; else return 0; }</pre> | | |

Value-Based Transformation

[illegible]

| Target Program | Original Constraints | Opt Constraints |
|---|--|-----------------|
| <pre>isBase64(unsigned k) { if(k > 255) return -1; if(b64[k] >= sym) return 1; else return 0; }</pre> | <pre>// Array Variables b64[0]=-1 ∧ b64[1]=-1 ∧ ... ∧ //Array Conditional b64k >= sym ∧ // Read Operation k=0 → b64k=b64[0] ∧ k=1 → b64k=b64[1] ∧ ...</pre> | |

Value-Based Transformation

[illegible]

| Target Program | Original Constraints | Opt Constraints |
|---|--|--|
| <pre> isBase64(unsigned k) { if(k > 255) return -1; if(b64[k] >= sym) return 1; else return 0; } </pre> | <pre> // Array Variables b64[0]=-1 ∧ b64[1]=-1 ∧ ... ∧ //Array Conditional b64k >= sym ∧ // Read Operation k=0 → b64k=b64[0] ∧ k=1 → b64k=b64[1] ∧ ... </pre> | <pre> x = lte(0≤k≤42 44≤k≤46 ..., -1, </pre> |

Value-Based Transformation

[illegible]

| Target Program | Original Constraints | Opt Constraints |
|---|--|--|
| <pre> isBase64(unsigned k) { if(k > 255) return -1; if(b64[k] >= sym) return 1; else return 0; } </pre> | <pre> // Array Variables b64[0]=-1 ∧ b64[1]=-1 ∧ ... ∧ //Array Conditional b64k >= sym ∧ // Read Operation k=0 → b64k=b64[0] ∧ k=1 → b64k=b64[1] ∧ ... </pre> | <div style="border: 2px solid red; padding: 5px; margin-bottom: 10px;"> $x = \text{Ite}(0 \leq k \leq 42 \mid 44 \leq k \leq 46 \mid \dots, -1).$ </div> <div style="border: 2px solid green; padding: 5px;"> $\text{ite}(k == 43, 62,$ </div> |

Value-Based Transformation

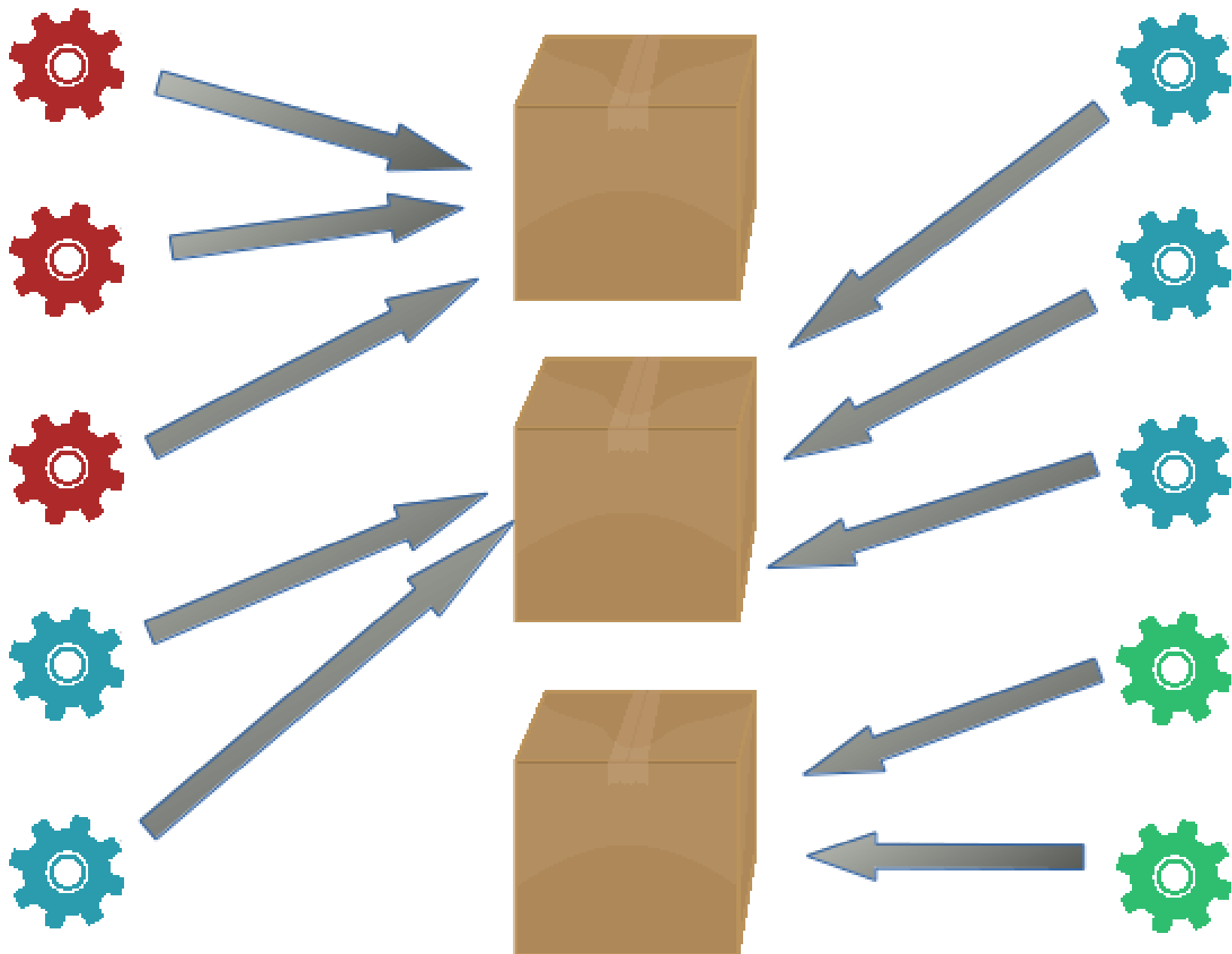
[illegible]

| Target Program | Original Constraints | Opt Constraints |
|---|--|---|
| <pre> isBase64(unsigned k) { if(k > 255) return -1; if(b64[k] >= sym) return 1; else return 0; } </pre> | <pre> // Array Variables b64[0]=-1 ∧ b64[1]=-1 ∧ ... ∧ //Array Conditional b64k >= sym ∧ // Read Operation k=0 → b64k=b64[0] ∧ k=1 → b64k=b64[1] ∧ ... </pre> | <div style="border: 2px solid red; padding: 5px; margin-bottom: 5px;"> $x = \text{Ite}(0 \leq k \leq 42 \mid 44 \leq k \leq 46 \mid \dots, -1).$ </div> <div style="border: 2px solid green; padding: 5px; margin-bottom: 5px;"> $\text{ite}(k == 43, 62,$ </div> <div style="border: 2px solid blue; padding: 5px;"> $\text{ite}(k == 47, 63,$ </div> |

Value-Based Transformation

[illegible]

| Target Program | Original Constraints | Opt Constraints |
|---|--|--|
| <pre> isBase64(unsigned k) { if(k > 255) return -1; if(b64[k] >= sym) return 1; else return 0; } </pre> | <pre> // Array Variables b64[0]=-1 ∧ b64[1]=-1 ∧ ... ∧ //Array Conditional b64k >= sym ∧ // Read Operation k=0 → b64k=b64[0] ∧ k=1 → b64k=b64[1] ∧ ... </pre> | <pre> x = lte(0≤k≤42 44≤k≤46 -1. ite(k == 43, 62, ite(k==47, 63, x >= sym </pre> |



What about partially symbolic arrays?

- Arrays can become partially symbolic due to program behavior or decisions made during analysis
- the Value-based transformation can be extended to handle them
- Add an individual case for each symbolic value in the array

Partially Symbolic Arrays

[illegible]

| Target Program | Original Constraints | Opt Constraints |
|---|---|-----------------|
| <pre>isBase64(unsigned k) { b64[0] = sym0; b64[1] = sym1; b64[2] = sym2; b64[3] = sym3; if(b64[k] >= sym) return 1; else return 0; }</pre> | <pre>// Array Variables b64[0]= ? ∧ b64[1]= ? ∧ ... ∧ //Array Conditional b64k >= sym ∧ // Read Operation k=0 → b64k=b64[0] ∧ k=1 → b64k=b64[1] ∧ ..</pre> | |

Partially Symbolic Arrays

[illegible]

| Target Program | Original Constraints | Opt Constraints |
|---|---|--|
| <pre>isBase64(unsigned k) { b64[0] = sym0; b64[1] = sym1; b64[2] = sym2; b64[3] = sym3; if(b64[k] >= sym) return 1; else return 0; }</pre> | <pre>// Array Variables b64[0]= ? ∧ b64[1]= ? ∧ ... ∧ //Array Conditional b64k >= sym ∧ // Read Operation k=0 → b64k=b64[0] ∧ k=1 → b64k=b64[1] ∧ ..</pre> | <pre>ite(k == 0, b64[0], ite(k==1, b64[1], ite(k==2, b64[2], ...</pre> |

How are the transformations implemented?



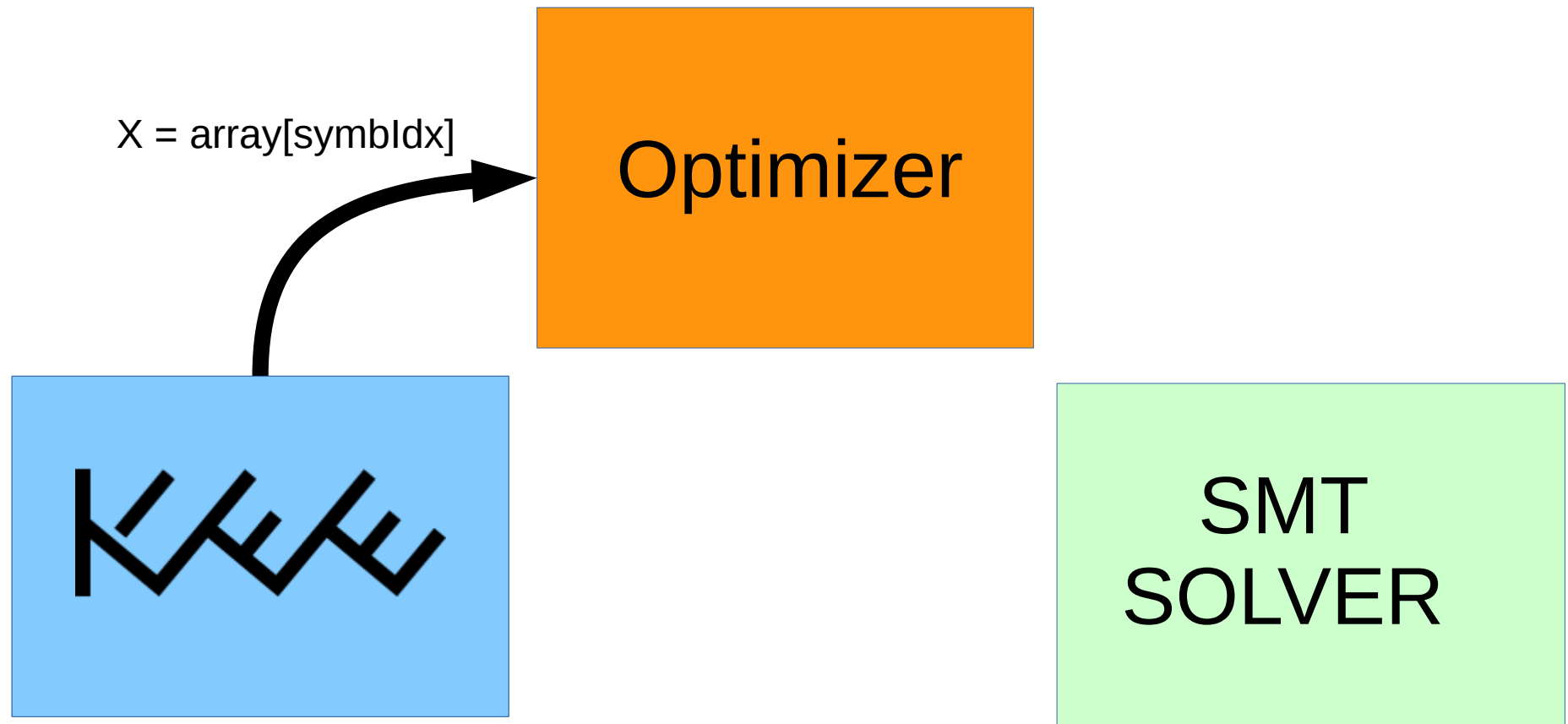
Optimizer



SMT
SOLVER

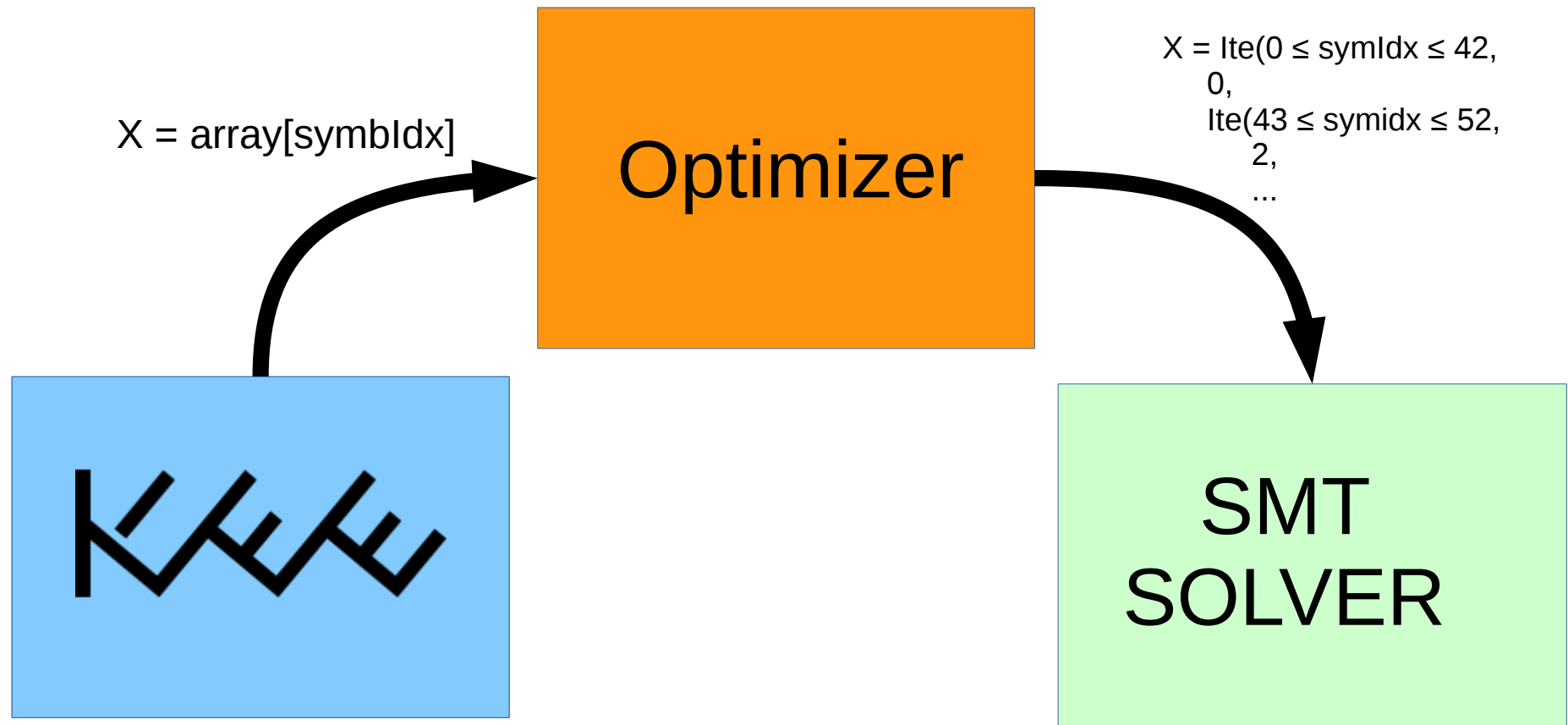
Implementation is available at: <https://srg.doc.ic.ac.uk/projects/klee-array/artifact.html>

How are the transformations implemented?



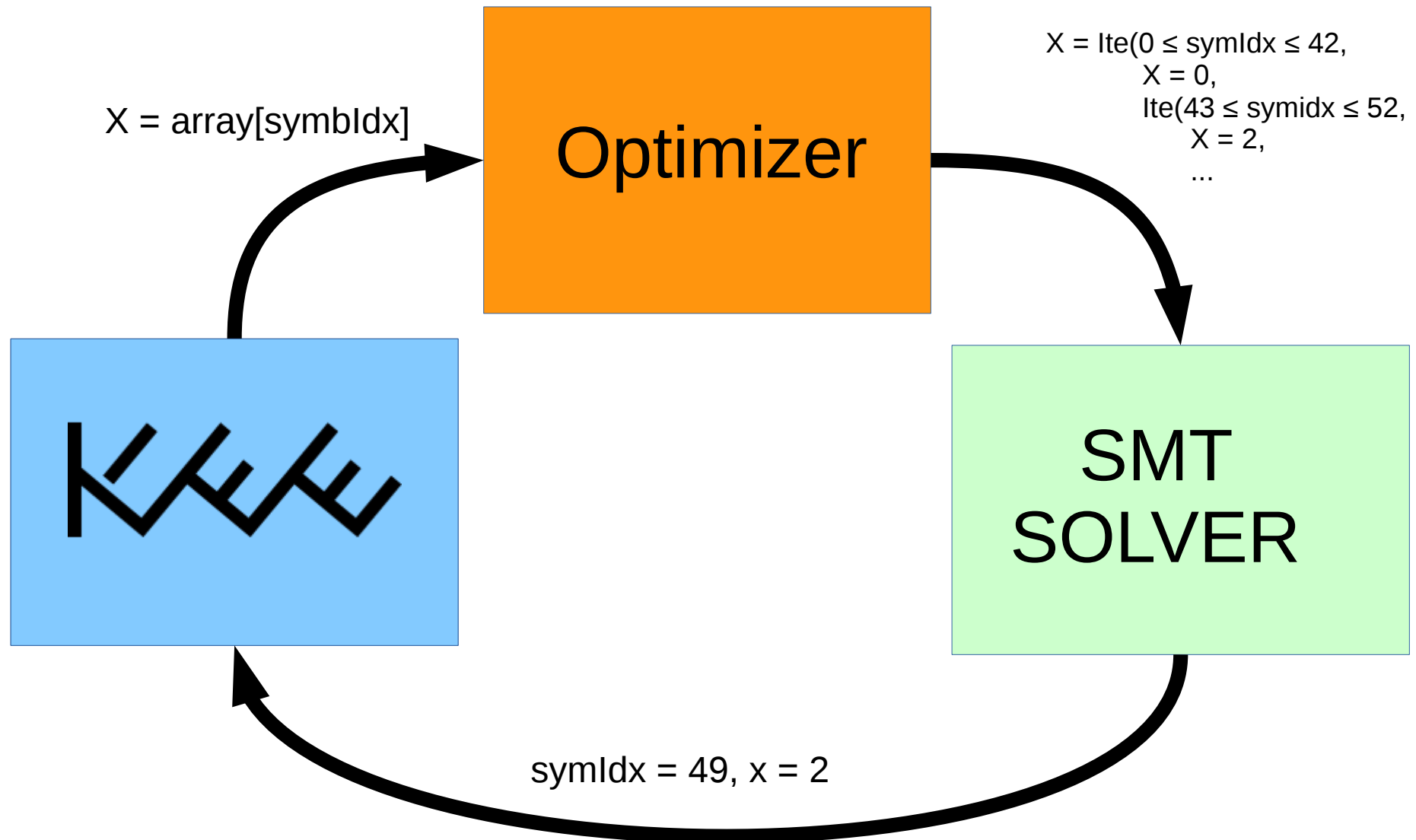
Implementation is available at: <https://srg.doc.ic.ac.uk/projects/klee-array/artifact.html>

How are the transformations implemented?



Implementation is available at: <https://srg.doc.ic.ac.uk/projects/klee-array/artifact.html>

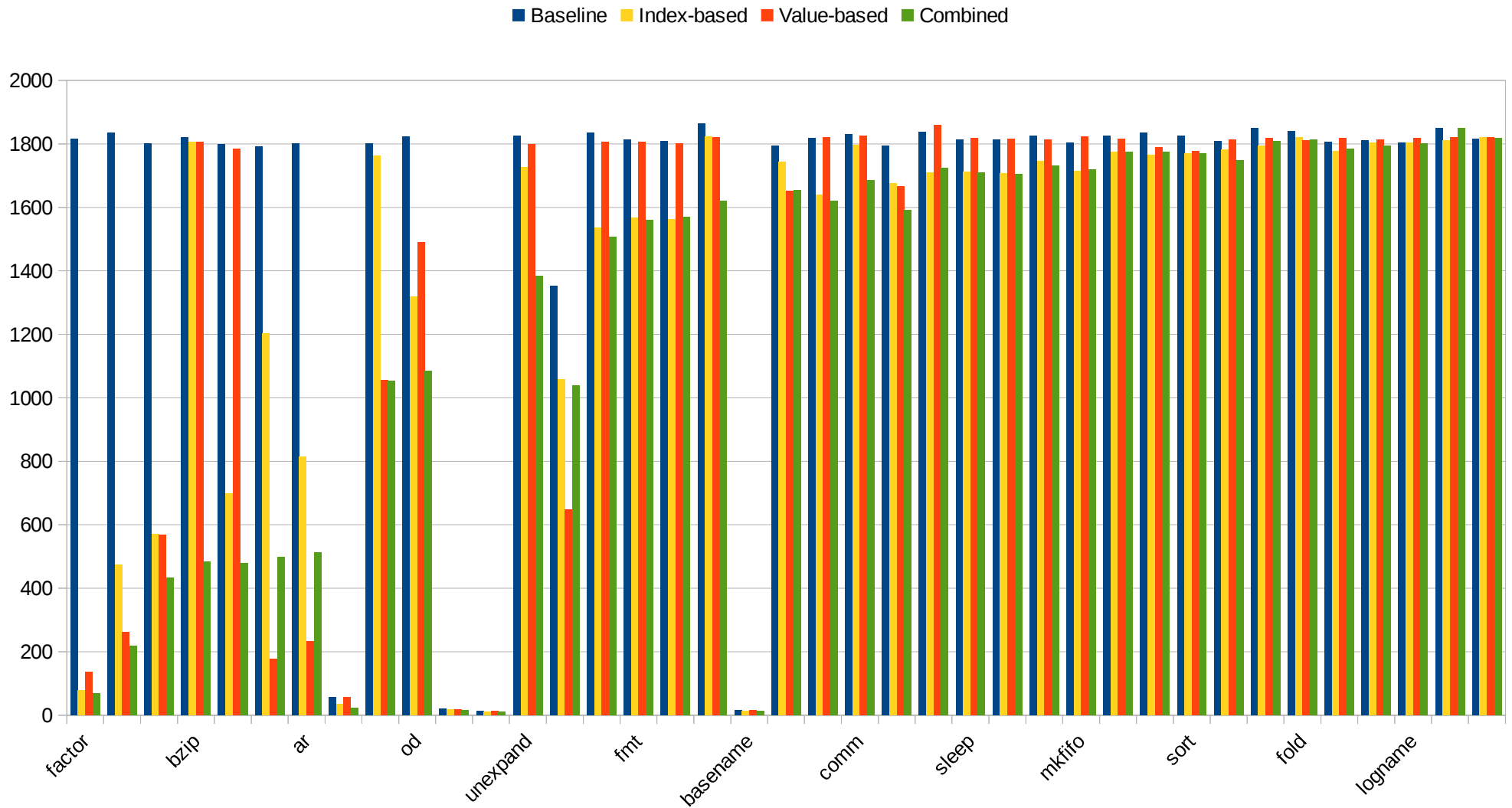
How are the transformations implemented?



Evaluation: Speed

- Performed on 104 programs from coreutils, binutils, and other open source repositories
- Baseline and Optimized runs execute the same number of instructions
- Instructions executed are logged to ensure fairness
- 7 programs with more than 3x speedup and none with slowdown

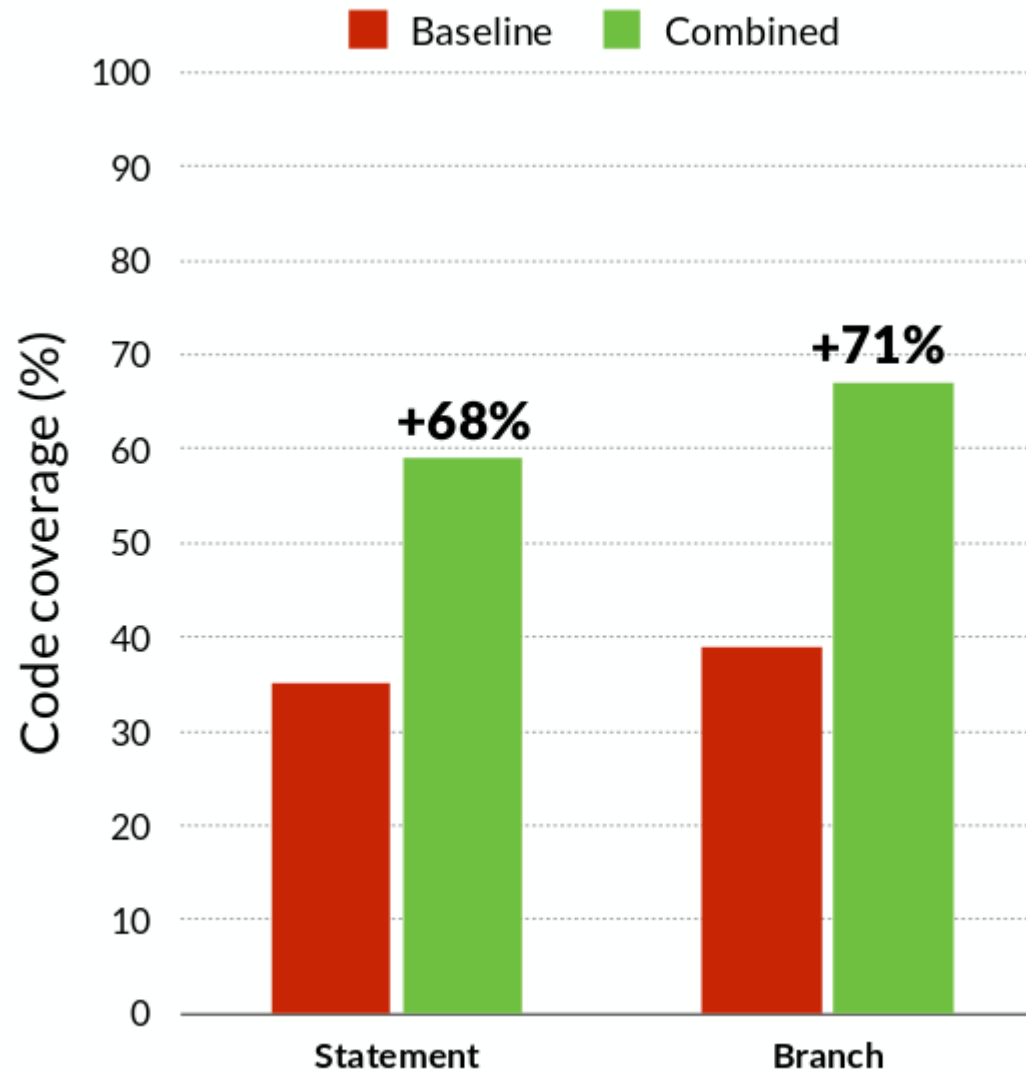
Evaluation: Speed



Evaluation: Code Coverage

- Performed on the program BC
- Baseline and optimized runs use a heuristic based search strategy designed to explore uncovered code
- Both programs are analyzed for 6 hours creating inputs with the most code coverage possible

Evaluation: Code Coverage



Related Work

- **Constraint Optimization**

- V. Ganesh and D. L. Dill. A decision procedure for bit-vectors and arrays. In Proc. of the 19th International Conference on Computer-Aided Verification (CAV'07) July 2007.
- Erete and A. Orso. Optimizing constraint solving to better support symbolic execution. In Proc. of the Workshop on Constraints in Software Testing, Verification, and Analysis (CSTVA'11), Mar. 2011.
- S. Dong, O. Olivo, L. Zhang, and S. Khurshid. Studying the influence of standard compiler optimizations on symbolic execution. In Proc. of the 26th International Symposium on Software Reliability Engineering (ISSRE'15), Nov. 2015.

Conclusion

- Our technique uses static and dynamic information to build specific encodings for individual arrays
- This encoding strategy yields considerable improvement in SMT solver performance
- This speedup allows symbolic execution to explore more code at a faster rate

Q&A

Thank You!

David Perry

PhD Student, Purdue University

Contact: perry74@purdue.edu