

CYBERSECURITY IN QUANTUM ERA

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- So before we continue, we provide some necessary definitions.
- *Cybersecurity* is the protection of computer systems and networks from information disclosure, theft or damage from malicious attack (cyberattacks).
- Also, we provide some definitions concerning cryptography.

ALICE, BOB AND EVE

- Usually Alice and Bob want to have a secure communication.

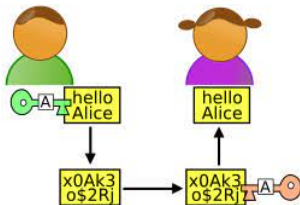


FIGURE: Bob encrypts the message *Hello Alice* and he sends it to Alice.

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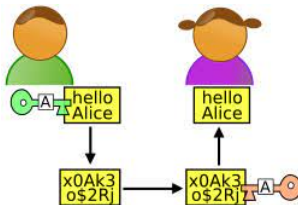


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- Although, their **friend** Eve 🧐, sometimes, tries to intercept their communication and if she is lucky or smart enough she may decrypt the message. Usually, Eve has enough resources, sophisticated algorithms, very powerful computers and she is very skilled in social engineering.

THREE CRYPTOGRAPHIC PRIMITIVES

- There two basic schemes. **Symmetric cryptosystems** (Scs) and **Public key Cryptosystems** (PkCs). We use both of them to build security protocols, such as SSL/TLS or IPsec or ssh. Also, sometimes instead of a PkC we use a **key agreement protocol**, such as Diffie-Hellman.

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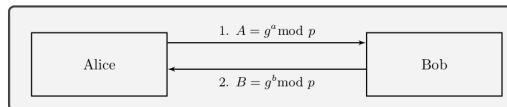


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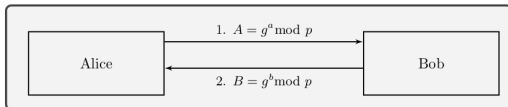


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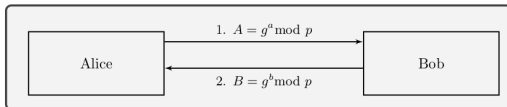


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- Here Eve, if somehow discovers a or b she will compromise all the communication.

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- Public Key Cryptography, as well as Key agreement protocols, as we use them today, are based on the following number theoretic problems : **Factorization** and **Discrete Logarithm Problem**.



FIGURE: Whit Diffie and Martin Hellman. ACM Turing Medal 2015.


WHAT ABOUT QUANTUM COMPUTERS?

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- The key word here is : *Shor's Algorithm*.
- It runs only in quantum computers (since it exploits quantum properties) and solves the problem of factorization and DL problem in polynomial time.

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- If a quantum computer with large memory ever constructed, then the most well known public key cryptosystem, RSA (and also Diffie-Hellman), shall break and all the current security in Internet will be compromised.

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- Also, Eve needs a quantum computer with enough memory. We measure the memory of quantum computers in qubits.
- For instance, today IBM-Q, the quantum computer of IBM, has 65 qubits memory. Although, IBM promises 1000-qubit quantum computer by 2023.
- *Bristlecone*, the google's quantum computer, it has 72 qubits

CYBERSECURITY IN QUANTUM ERA?

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- The current quantum computers are noisy
- How large is a 2048-bit integer?

HOW LARGE IS A 2048 OR 4096 BITS INTEGER?

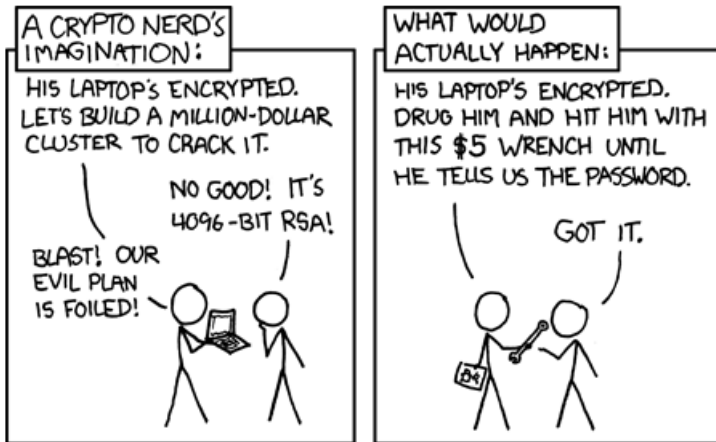


FIGURE: Licence:xkcd.com, CC BY-NC 2.

HOW LARGE IS A 2048 BITS INTEGER?

- Here is the number you have to factor in order to break the security of <https://www.sfhmmy.gr/>

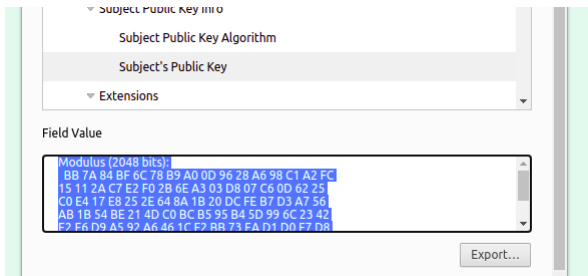


FIGURE: RSA public key

HOW LARGE IS A 2048 BITS INTEGER?

- ...and if we represent it as an integer

```
236669791753739784024992868689634815865603851656209245508903792881110423\  
238404152066731350154873684923058959342095561097416509926668052067194395\  
159648229493748798615332050052858896009784504620170814783692610398045900\  
140167222626905623214902579061673663159365906153364449644240912381048000\  
238183883103063286643646262270394166332736638459211300770684353422869195\  
996527266327813527690931019987991992600942095272228931181847018205192648\  
687960611937525233108111461550726182358572994215168369025892791967296674\  
131323898777664515932960356538839891068969078068777021646419769454794321\  
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- A quantum computer capable of implementing Shor's algorithms will return the two prime factors.

HOW LARGE IS A 2048 BITS INTEGER?

- The current record for factorization RSA modulus with 829 bits and it was factored in Feb. 2020 using CADO-NFS software.

RSA-250 = 214032465020744961264423072839333563008614715144755017797754920881418023447
 1401366433455190958046796109928518724709145876873962619215573630474547705208
 0511905649310668769159001975940569345745223058932597669747168173806936489469
 9871578499497937497937

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- The computation involved tens of thousands of machines and was finished in a few months.
- The (heuristic) time complexity of GNFS is $2^{O(\sqrt[3]{\log_2(N)})}$.

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- Although, we don't know or predict when a large memory quantum computer has been built. But we have such computers with low memory.

MAYBE IT'S TIME TO PANIC

- *Dystopia Scenario* : Say, all the encrypted data you have exchanged the previous year with a plethora of servers, were collected and kept in some data center. They will be kept and decrypted when a large memory quantum computer will be built.

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- When we are talking about PQC we (usually) mean cryptography which is **not** based on Factorization and Discrete Logarithm Problem (DLP).

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- Supersingular Isogeny Problem 2006, (SIDH) Key exchange



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- What is more, there are eight alternate candidate algorithms.
- For more information <https://csrc.nist.gov/projects/post-quantum-cryptography>

SO WHY WE DO NOT USE THEM AND STOP WORRYING ABOUT QUANTUM COMPUTERS?

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- Well, Post quantum cryptographic algorithms are not very efficient...yet
- We must have good evidences that they are indeed quantum resistant. So, years of research need to be done.

LATTICE BASED PQC

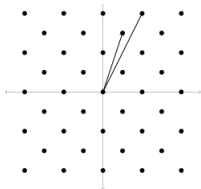


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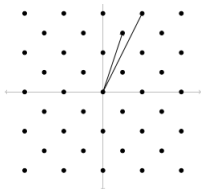


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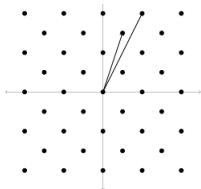


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- The problem of finding a shortest vector (which always exists) is called Shortest Vector Problem (SVP).

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- Another example is the Short Integer Problem (SIS).

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- The best result until now, provides an algorithm with complexity $2^{0.268n+o(n)}$ instead of $2^{0.298n+o(n)}$ in classic computers.

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$$\{p_1 = p_2 = \dots = p_m = 0 \pmod{p}\}$$

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- Let $p_1(x_1, \dots, x_k), \dots, p_m(x_1, \dots, x_k)$ some quadratic polynomials over a finite field. For simplicity consider that we reduce all the coefficients of p_i modulo a prime number p .
- The MQ problem is the problem of finding the solutions of

$$\{p_1 = p_2 = \dots = p_m = 0 \pmod{p}\}$$

- It seems easy, but it was proved to be NP-hard.

Thank you!