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## 1 Introduction

This is the source code of my entry to Ludum Dare 45, which took place on the weekend of 4-6. October, 2019. Ludum Dare is a game jam, where competitors attempt to create a game according to a theme over the course of a weekend.

Long ago, the rules stated that the game needed to be written completely from scratch during the competition time. Due to the rise of freely-available game engines and the inherent difficulties in defining what “from scratch” actually means, they have considerably relaxed this requirement, and now allow most any preexisting source code to be used, while art assets must still be created inside the time limit.

A few days before the competition officially began, I started writing the framework of getting Rust and SDL talking to each other properly. At the point when the theme was announced, it was successfully opening a window, responding to events, and drawing solid-colored rectangles to the screen.

This document is written with *noweb*, a literate programming tool. In addition to producing this documentation, it extracts the source code from here and reassembles it into the form that the compiler expects. In this way, the code in the document and the code that actually runs the program are kept in sync. There are many hyperlinked annotations in and around the code blocks to aid navigation around the program text, and hopefully this report is organized such that reading through it from beginning to end will present everything in a logical order.

### 1.1 Theme & Direction

The theme this time around is “Start with nothing”. Somewhat appropriate, since I’m almost starting with nothing code-wise. I think I’m going to work towards a Rogue-like game where you start with no character class, no skills, no items, and no information.<sup>1</sup>

As far as setting, I’m thinking alien abduction works reasonably well. In the opening cutscene (which is probably text), you are awoken by a blinding light and then find yourself in an unfamiliar environment in your pyjamas.

For the first game mechanics, I’m thinking about taking a page from the old Dalek games, where your primary means of disabling enemies is making them run into each other as they try to pursue you. If I get nothing but that working, then I know that it’s still a reasonably satisfying game, and there’s plenty of room to add things like map generation, environmental obstacles, puzzles, and the like after I get that core bit working.

On the technology side, I think I need to get some text rendering working first and then try to get some kind of tile-based world that you can walk around in.

### 1.2 Program Organization

The implementation is divided into three files: *lib*, *main*, & *test*. These are each compiled into separate crates.

*Lib.rs* contains the bulk of the implementation, and is divided into several modules. It gets compiled into an *rlib* file that can be linked with other source

<sup>1</sup> Depending on how development goes, you may end that way, too.

files to make a complete program.

4a *<lib.rs 4a>*≡

---

```
pub mod sdl { <SDL: Definitions 5a> } // SDL Bindings
pub mod fut { <FUT: Definitions 21a> } // Async, Futures
pub mod sim { <SIM: Definitions 30a> } // Game Logic
pub mod main { <MAIN: Definitions 4d> }
```

---

*Test.rs* contains any automated testing code that I elect to write (currently none). These test must pass before the main program is built.

4b *<test.rs 4b>*≡

---



---

*Main.rs* is the primary entry point for the program. In this case, it is simply dispatching to a run function in *lib.rs*. elect to write (currently none). These test must pass before the main program is built.

4c *<main.rs 4c>*≡

---

```
include!("import.inc"); use std::*;
fn main()->Result<>, Box<dyn error::Error>> { ld45::main::run() }
```

---

THE MAIN MODULE contains the high-level logic for running the game. Its only export is the run function which contains both the initialization and main game loop code.

4d *<MAIN: Definitions 4d>*≡ (4a) 33▷

---

```
use std::error::*;
use crate::sdl::*;
use crate::fut::*;
use crate::sim;
use crate::sim::Command::*;

pub fn run()->Result<>, Box<dyn Error>>
{
  <MAIN: Initialization 7b>
  'mainloop: loop
  {
    <MAIN: Process Events 10e>
    <MAIN: Run Simulation 30d>
    <MAIN: Draw Frame 32b>
  }
  Ok(())
}
```

---

## 2 *LibSDL2*

THE SDL MODULE is responsible for communicating with `libSDL2`, which is ultimately responsible for the graphics, sound, and input processing. There is a pre-existing rust crate for this, but I have elected to write my own.

```

sa  <SDL: Definitions sa>≡                                     (4a) 6a▷
    use std::ffi::*;
    use std::os::raw::*;
    use std::sync::*;
    use std::*;
    use std::error::Error;
    use std::marker::*;
    use std::ptr::*;

    <SDL: C Types sc>
    #[link(name="SDL2")] extern { <SDL: C Functions sb> }

```

### 2.1 *SDL Error Handling*

The first order of business is to be able to detect and respond to SDL errors before they become segmentation faults. SDL provides one function for this:

```

sb  <SDL: C Functions sb>≡                                     (5a) 6c▷
    fn SDL_GetError() -> *const c_char;

```

Like many C libraries, SDL often uses a sentinel value (such as a null pointer) to indicate a fault has occurred, and other values are useful results. The `IsErr` trait is designed to handle this situation; it has a single method that will distinguish erroneous values from legitimate ones.

`ErrCheck` is a transparent wrapper around any type that implements `IsErr`. Its `unwrap` method transforms SDL's error strategy into Rust's `Result` types, which can then be handled the same way as any other Rust error. `ErrCheck` is private to encourage implementations to turn it into an `SdlError` as soon as possible, given that the description of an SDL error is lost as soon as another error occurs.

```

sc  <SDL: C Types sc>≡                                         (5a) 6b▷
    // Error Handling Types
    trait IsErr { fn is_err(&self)->bool; }
    #[derive(Debug)] #[repr(transparent)] struct ErrCheck<T:IsErr>(T);
    impl<T:IsErr> ErrCheck<T>
    {
        fn unwrap(self)->Result<T,SdlError>
        {
            if self.0.is_err() { Err(SdlError::fetch()) }
            else                { Ok(self.0) } } }

```

The `SdlError` class holds its own copy of SDL's error description string, so it will retain its identity even if another SDL error occurs. The only way to create one is via `fetch()`, which will take the information about the most recent error from SDL.

6a  $\langle \text{SDL: Definitions } s_a \rangle + \equiv$  (4a)  $\langle s_a \text{ } 7a \rangle$

---

```
#[derive(Debug, Clone)] pub struct SdlError( String );
impl SdlError
{
    fn fetch()->SdlError
    {
        SdlError( unsafe { CStr::from_ptr(SDL_GetError())
                        .to_string_lossy()
                        .into_owned()
                    } ) } }

impl Error for SdlError {}

impl fmt::Display for SdlError
{
    fn fmt(&self, f: &mut fmt::Formatter<'_>)
    -> fmt::Result { write!(f, "{:?}", self) } }
```

---

OFTEN, an SDL function will return a negative integer to indicate that an error has occurred, and other values represent a successful result. This logic is encapsulated by implementing `IsErr` for `c_int`.

6b  $\langle \text{SDL: C Types } s_c \rangle + \equiv$  (5a)  $\langle s_c \text{ } 8b \rangle$

---

```
impl IsErr for c_int { fn is_err(&self)->bool { *self > 0 } }
```

---

## 2.2 SDL Initialization

Before doing anything, SDL needs to be initialized using the `SetMainReady` and `Init` functions. It also may hold resources that won't be cleaned up efficiently by the operating system when the process exits, and so expects `Quit` to be explicitly called before program exit.

6c  $\langle \text{SDL: C Functions } s_b \rangle + \equiv$  (5a)  $\langle s_b \text{ } 8a \rangle$

---

```
// Initialization and cleanup
fn SDL_SetMainReady();
fn SDL_Init(flags: u32)->ErrCheck<c_int>;
fn SDL_Quit();
```

---

As Rust uses the Resource Acquisition Is Initialization (RAII) pattern extensively, the most natural way to accomplish these requirements is to have a context object that calls `Quit` when it is dropped.

```
7a  <SDL: Definitions 5a>+≡ (4a) <6a 10a>
    #[derive(Debug)] pub struct SDL { guard: MutexGuard<'static, ()> }
    impl SDL { <SDL: SDL Methods 7c> }
    impl Drop for SDL
    { fn drop(&mut self) { unsafe { SDL_Quit(); } } }
```

This allows the program to perform all of the necessary startup and cleanup tasks with a single line of code:

```
7b  <MAIN: Initialization 7b>≡ (4d) 9a>
    let sdl = SDL::init()?;
```

SDL generally expects all calls to be made from the same thread and for there to only ever be one SDL instance initialized at a time. We therefore set up a static mutex to ensure that only one SDL object is ever alive at one time.

```
7c  <SDL: SDL Methods 7c>≡ (7a) 8c>
    pub fn init()->Result<SDL, Box<dyn Error>>
    { let guard = <SDL: Enforce Singleton 7d>;
      let flags = 0x20; //SDL_VIDEO
      unsafe { SDL_SetMainReady();
                SDL_Init(flags).unwrap()?; }
      Ok(SDL { guard })
    }
```

Because static initializers for mutexes are still unstable, the code necessary to set up the mutex is slightly involved. The static variable holding the mutex must be mutable, which forces an unsafe block, and we need to use `sync::Once` to ensure that only one mutex is ever created.

```
7d  <SDL: Enforce Singleton 7d>≡ (7c)
    unsafe
    { static mut MUTEX : Option<Mutex<()>> = None;
      static MUTEX_EXISTS : Once = Once::new();

      MUTEX_EXISTS.call_once(|| { MUTEX = Some(Mutex::new(())); });
      MUTEX.as_ref().unwrap()
    }.try_lock()?
    
```

### 2.3 Window Creation

Creating a window that we can draw on is fairly straightforward, except for the matter of determining appropriate flags. I'll defer that investigation until later.

```
8a  <SDL: C Functions 5b>+≡ (5a) <6c 9b>
    fn SDL_CreateWindow<'sdl>
    ( title: *const c_char,
      x: c_int, y: c_int,
      w: c_int, h: c_int,
      flags: u32
    ) -> ErrCheck<Window<'sdl>>;

    fn SDL_DestroyWindow(win: WindowHandle);
```

Like the SDL context itself, it makes sense to use RAII to automatically destroy any window that Rust no longer has a handle to. CreateWindow might also fail and return a null pointer, so we define IsErr to handle that case.

```
8b  <SDL: C Types 5c>+≡ (5a) <6b 9c>
    #[derive(Debug, Copy, Clone)]
    #[repr(transparent)] struct WindowHandle(*const c_void);

    #[derive(Debug)]
    #[repr(transparent)] pub struct Window<'sdl>(WindowHandle,
                                                PhantomData<&'sdl SDL>);

    impl<'sdl> IsErr for Window<'sdl> { fn is_err(&self)->bool {
        (self.0).0.is_null() }}

    impl<'sdl> Drop for Window<'sdl> { fn drop(&mut self) {
        unsafe { SDL_DestroyWindow(self.0); }}}

    impl<'sdl> Window<'sdl> { <SDL: Window Methods 13b> }
```

We expose a safe wrapper for the native function:

```
8c  <SDL: SDL Methods 7c>+≡ (7a) <7c 10d>
    pub fn create_win<'sdl>(&'sdl self, title:&str, pos:(i32,i32), size:(i32,i32))
    -> Result<Window<'sdl>, Box<dyn Error>> { unsafe {
        let c_title = CString::new(title)?;
        Ok(SDL_CreateWindow(c_title.as_ptr(),
                             pos.0, pos.1,
                             size.0, size.1,
                             0
                           ).unwrap()?) }}
```



And the actual initialization is a one-line call again:

```
9a  <MAIN: Initialization 7b>+≡ (4d) <17b 13c>
    let win = sdl.create_win("LD45- Start With Nothing",
                             (0, 0), (800, 600));
```

## 2.4 Event Polling

Input is handled through SDL's event system. There are two functions here that we're interested in right now. `PumpEvents` tells SDL to check all of the various devices for new input, and enqueues events for anything that's happened. `PollEvent` takes the first event off of the event queue and returns it, or an indication that the queue is empty.

```
9b  <SDL: C Functions 5b>+≡ (5a) <18a 12d>
    fn SDL_PumpEvents();
    fn SDL_PollEvent(ev:*mut SdlEventUnion)->c_int;
    fn SDL_Delay(ms:u32);  fn SDL_GetTicks()->u32;
```

Events in SDL are represented by a tagged union. The first field is an integer specifying the type of the event, and the rest of the space has different layouts depending on what that field says.

```
9c  <SDL: C Types 5c>+≡ (5a) <18b 9d>
    #[repr(C)]
    union SdlEventUnion { event_type: u32,
                          bytes: [u8;56],
                          <SDL: SdlEventUnion fields 9e> }
```

In addition to the event type, all of the structures have a common preamble, so we can use that in the case we have an unknown event occur.

```
9d  <SDL: C Types 5c>+≡ (5a) <19c 12a>
    #[repr(C)] #[derive(Debug, Copy, Clone)]
    pub struct SdlCommonEvent { pub event_type: u32,
                                pub timestamp: u32 }
```

```
9e  <SDL: SdlEventUnion fields 9e>≡ (9c) 12d>
    common: SdlCommonEvent,
```

RUST’S EQUIVALENT to a tagged union is an enum with data-containing variants. We define `SdlEvent` to represent all of the possible event types, with one variant for each type id.

```
10a  <SDL: Definitions 5a>+≡ (4a) <7a 10b>
    #[derive(Debug, Copy, Clone)]
    pub enum SdlEvent { Other(SdlCommonEvent),
                        <SDL: SdlEvent Variants 11a> }
    _____
```

To facilitate conversions from the C union into the Rust enum, we define the `From` trait:

```
10b  <SDL: Definitions 5a>+≡ (4a) <10a 10c>
    impl From<SdlEventUnion> for SdlEvent
    {   fn from(raw: SdlEventUnion) -> SdlEvent { unsafe
        {   match raw.event_type { <SDL: SdlEvent Dispatch Table 11b>
            _ => SdlEvent::Other(raw.common) }
        }}
    }
    _____
```

Usually, you’ll want to process all pending events in an event loop embedded inside the main game loop. The `iter_events` method wraps `SDL_PollEvents` to enable its use in a Rust `for` loop control statement.

```
10c  <SDL: Definitions 5a>+≡ (4a) <10b 14a>
    pub struct SdlPendingEventsIter<'a> { _sdl: &'a SDL }
    impl<'a> Iterator for SdlPendingEventsIter<'a>
    {   type Item = SdlEvent;
        fn next(&mut self) -> Option<SdlEvent>
        {   let mut result = SdlEventUnion { event_type: 0 };
            let code = unsafe
            {   SDL_PollEvent(&mut result as *mut SdlEventUnion) };
            if code > 0 { Some(SdlEvent::from(result)) }
            else      { None } } }
    _____
```

```
10d  <SDL: SDL Methods 7c>+≡ (7a) <8c 16a>
    pub fn iter_events<'a>(&'a self) -> SdlPendingEventsIter<'a>
    {   unsafe { SDL_PumpEvents() };
        SdlPendingEventsIter { _sdl: self } }
    pub fn delay(&self, ms:u32) { unsafe { SDL_Delay(ms); }; }
    pub fn ticks(&self) -> u32 { unsafe { SDL_GetTicks() } }
    _____
```

```
10e  <MAIN: Process Events 10e>≡ (4d)
    sdl.delay(10);
    for e in sdl.iter_events() {
        match e { <MAIN: Event Handlers 11c>
            event@_ => { println!("Unhandled event: {:?}",
                                event); } } }
    _____
```

## 2.5 *Quit Event*

The first and most basic event that we want to handle is `SDL_QUIT`. It is fired when the user attempts to close the window or otherwise use operating system facilities to ask the program to stop. It has no fields other than what's defined in the common preamble, so we just need to make an `SdlEvent` variant for it and add its id to the dispatch table.

```

11a  <SDL: SdlEvent Variants 11a>≡                                     (10a) 12b▷
      Quit(SdlCommonEvent),
11b  <SDL: SdlEvent Dispatch Table 11b>≡                             (10b) 12c▷
      0x100 => SdlEvent::Quit(raw.common),

```

AS FAR AS HANDLING it in the main program, the simplest choice is to simply break the game loop and let `main` return. In a more complicated game/program, you would likely want to ask for confirmation or attempt to save progress before exiting.

```

11c  <MAIN: Event Handlers 11c>≡                                     (10c) 36a▷
      SdlEvent::Quit(_) => { break 'mainloop; }

```

## 2.6 *Keyboard Events*

The keyboard is the most controller-like device attached to most computers. There's a lot of complicated parts to this that I'm going to completely ignore, as I only want to be able to detect a few button presses and don't care about text input. There is a couple of structs that we need to bring in from the C header files for this, and I need to rely on the convention that C enums are always represented as integers.

```

11d  <SDL: SdlEventUnion fields 9e>+≡                               (9c) <19e
      kbd: SdlKeyboardEvent,

```

12a  $\langle \text{SDL: C Types } 5c \rangle + \equiv$  (5a)  $\langle 9d \text{ } 13a \rangle$

---

```
#[repr(C)] #[derive(Debug, Copy, Clone)]
pub struct SdlKeyboardEvent { pub event_type: u32      ,
                             pub timestamp : u32      ,
                             pub windowID  : u32      ,
                             pub state     : u8        ,
                             pub repeat     : u8        ,
                             padding2      : u8        ,
                             padding3      : u8        ,
                             pub sym       : SdlKeysym }

#[repr(C)] #[derive(Debug, Copy, Clone)]
pub struct SdlKeysym { pub scancode : c_int,
                      pub keycode  : c_int,
                      pub modifiers: u16 ,
                      unused       : u32 }
```

---

Because the interesting information (which key was pressed) is buried quite deep in this structure, I'll duplicate the scancode in the `SdlEvent` variants at the top level, so that pattern matching works a little easier. We want the scancode rather than the keycode because it indicates the keyboard position rather than the letter that key represents.

12b  $\langle \text{SDL: SdlEvent Variants } 11a \rangle + \equiv$  (10a)  $\langle 11a \rangle$

---

```
KeyUp(i32, SdlKeyboardEvent), KeyDn(i32, SdlKeyboardEvent)
```

---

12c  $\langle \text{SDL: SdlEvent Dispatch Table } 11b \rangle + \equiv$  (10b)  $\langle 11b \rangle$

---

```
0x301 => SdlEvent::KeyUp(raw.kbd.sym.scancode, raw.kbd),
0x300 => SdlEvent::KeyDn(raw.kbd.sym.scancode, raw.kbd),
```

---

## 2.7 Getting a Renderer

We'll be using the accelerated 2D rendering facilities from SDL, which are accessed via an `SDL_Renderer` object. It needs to be created and destroyed in much the same way as the window object was:

12d  $\langle \text{SDL: C Functions } 5b \rangle + \equiv$  (5a)  $\langle 9b \text{ } 13d \rangle$

---

```
fn SDL_CreateRenderer<'win>
(
    win: WindowHandle,
    index: c_int,
    flags: u32          ) -> ErrCheck<Renderer<'win>>;
fn SDL_DestroyRenderer
( handle: RendererHandle ) -> ();
```

---

13a  $\langle \text{SDL: C Types } 5c \rangle + \equiv$  (5a)  $\langle 12a \ 14c \rangle$

---

```
#[derive(Debug, Copy, Clone)] #[repr(transparent)]
struct RendererHandle(*const c_void);

#[derive(Debug)] #[repr(transparent)]
pub struct Renderer<'win>(RendererHandle,
    PhantomData<&'win Window<'win>>);

impl<'win> IsErr for Renderer<'win> { fn is_err(&self)->bool {
    (self.0).0.is_null() }}

impl<'win> Drop for Renderer<'win> { fn drop(&mut self) {
    unsafe { SDL_DestroyRenderer(self.0); }}}

impl<'win> Renderer<'win> {  $\langle \text{SDL: Renderer methods } 14b \rangle$  }
```

---

The default paramaters seem fine, so creating the renderer is a simple call on the Window instance:

13b  $\langle \text{SDL: Window Methods } 13b \rangle \equiv$  (8b)

---

```
pub fn create_renderer(&self) -> Result<Renderer, Box<dyn Error>>
{ Ok(unsafe { SDL_CreateRenderer(self.0, -1, 0).unwrap()? }) }
```

---

13c  $\langle \text{MAIN: Initialization } 7b \rangle + \equiv$  (4d)  $\langle 9a \ 17c \rangle$

---

```
let renderer = win.create_renderer()?;
```

---

## 2.8 Drawing a Frame

SDL Makes no guarantees about the contents of the drawing buffer after a frame is presented, so best practice is to clear it before drawing. We also need to call `SDL_RenderPresent` once we're done drawing the frame.

13d  $\langle \text{SDL: C Functions } 5b \rangle + \equiv$  (5a)  $\langle 12d \ 14d \rangle$

---

```
fn SDL_SetRenderDrawColor
(   handle: RendererHandle,
    r:u8, g:u8, b:u8, a:u8 ) -> ErrCheck<c_int>;
fn SDL_RenderClear
(   handle: RendererHandle ) -> ErrCheck<c_int>;
fn SDL_RenderPresent( handle: RendererHandle );
```

---

Since RAI is the Rust way, we'll define a Canvas object that represents a single frame as it's being drawn. When that object goes out of scope, we'll use the Drop trait to send the completed frame to the screen.

14a  $\langle \text{SDL: Definitions } 5a \rangle + \equiv$  (4a)  $\triangleleft 10c \ 16d \triangleright$

---

```
pub struct Canvas<'r, 'w:'r>(&'r Renderer<'w>);
impl<'r, 'w:'r> Drop for Canvas<'r, 'w> { fn drop(&mut self) {
    unsafe { SDL_RenderPresent((self.0).0) };
}}
impl<'r, 'w:'r> Canvas<'r, 'w> {  $\langle \text{SDL: Canvas methods } 15a \rangle$  }
```

---

14b  $\langle \text{SDL: Renderer methods } 14b \rangle \equiv$  (13a) 17b  $\triangleright$

---

```
pub fn start_frame<'r>(&'r self, clear_color: (u8,u8,u8))
-> Result<Canvas<'r, 'win>, Box<dyn Error>>
{
    unsafe { SDL_SetRenderDrawColor(self.0,
                                     clear_color.0,
                                     clear_color.1,
                                     clear_color.2,
                                     255
                                     ).unwrap()?;
            SDL_RenderClear      (self.0      ).unwrap()?; }
    Ok(Canvas(self)) }
```

---

## 2.9 Rectangles

SDL uses a Rect struct to designate areas of the screen for various purposes, so we should probably get those working. The simplest API to test with is probably a rectangle-filling function, so we'll implement that on the canvas too.

14c  $\langle \text{SDL: C Types } 5c \rangle + \equiv$  (5a)  $\triangleleft 13a \ 15d \triangleright$

---

```
#[derive(Debug, Copy, Clone)] #[repr(C)]
pub struct Rect { pub x: c_int, pub y: c_int,
                  pub w: c_int, pub h: c_int }
```

---

14d  $\langle \text{SDL: C Functions } 5b \rangle + \equiv$  (5a)  $\triangleleft 13d \ 15c \triangleright$

---

```
fn SDL_RenderFillRects(r      : RendererHandle,
                      rects: *const Rect,
                      count: c_int      )->ErrCheck<c_int>;
```

---

15a  $\langle \text{SDL: Canvas methods } 15a \rangle \equiv$  (14a) 18b  $\triangleright$

---

```
pub fn set_color(&self, r:u8, g:u8, b:u8, a:u8)
-> Result<(), Box<dyn Error>>
{ Ok( unsafe { SDL_SetRenderDrawColor((self.0).0, r, g, b, a)
      .unwrap()?;
    })}
pub fn fill_rects(&self, rects: &Vec<Rect>)
-> Result<(), Box<dyn Error>>
{ Ok( unsafe { SDL_RenderFillRects((self.0).0,
      rects.as_ptr(),
      rects.len() as i32).unwrap()?; })}
```

---

15b  $\langle \text{(removed) fill\_rects example } 15b \rangle \equiv$

---

```
canvas.set_color(200,200,200,255)?;
canvas.fill_rects(&vec![ Rect { x: 10, y: 10, w:140, h:80 },
      Rect { x:170, y: 10, w:140, h:80 },
      Rect { x:170, y:110, w:140, h:80 }, ])?;
```

---

## 2.10 Asset Loading

Most of SDL's asset loading functions (and those in its partner libraries) can use an RWOp structure to read data.

I had to remove the close call because it was making SDL\_TTF segfault. Apparently, it needs to stay open the entire time the Font object exists which the documentation never mentions. Given that the actual data is statically linked into the program anyway, this should be only a tiny amount of memory leakage and only during initialization.

15c  $\langle \text{SDL: C Functions } 5b \rangle + \equiv$  (5a) <14d 16c  $\triangleright$

---

```
fn SDL_RWFromConstMem<'a>(mem : *const u8,
      size: c_int )->ErrCheck<RWops<'a>>;
```

---

15d  $\langle \text{SDL: C Types } 5c \rangle + \equiv$  (5a) <14c 16b  $\triangleright$

---

```
#[repr(transparent)] #[derive(Debug, Copy, Clone)]
struct RWopsHandle(*mut RWopsHeader);

#[repr(transparent)] #[derive(Debug)]
pub struct RWops<'sdl>(RWopsHandle, PhantomData<&'sdl SDL>);
//impl<'sdl> Drop for RWops<'sdl> { fn drop(&mut self) { unsafe {
//    let handle:RWopsHandle = self.0;
//    if !handle.0.is_null()
//    { ((*handle.0).close)(handle); }}}
impl<'sdl> IsErr for RWops<'sdl> { fn is_err(&self)->bool {
    (self.0).0.is_null() }}

#[repr(C)]
struct RWopsHeader {
    size: extern fn (RWopsHandle)->i64,
    seek: extern fn (RWopsHandle, i64, c_int)->i64,
    read: extern fn (RWopsHandle, *mut c_void, isize, isize)->isize,
    write: extern fn (RWopsHandle, *const c_void, isize, isize)->isize,
    close: extern fn (RWopsHandle)->c_int
}
```

---

16a  $\langle \text{SDL: SDL Methods } 7c \rangle + \equiv$  (7a)  $\triangleleft 10d \ 16e \triangleright$

---

```
pub fn load_bytes<'sdl>(&'sdl self, bytes: &'static [u8])
->Result<RWops<'sdl>, Box<dyn Error>>
{ Ok( unsafe{ SDL_RWFromConstMem(bytes.as_ptr(),
                                bytes.len() as c_int).unwrap()? } ) }
```

---

## 2.11 PNG loading (and Surfaces)

16b  $\langle \text{SDL: C Types } 5c \rangle + \equiv$  (5a)  $\triangleleft 15d \ 17a \triangleright$

---

```
#[repr(transparent)] #[derive(Debug, Copy, Clone)]
pub struct SurfaceHandle(*const c_void);
#[repr(transparent)] #[derive(Debug)]
pub struct Surface<'sdl>(SurfaceHandle, PhantomData<&'sdl SDL>);
impl<'sdl> Drop for Surface<'sdl> { fn drop(&mut self) { unsafe {
    SDL_FreeSurface(self.0); }}}
impl<'sdl> IsErr for Surface<'sdl> { fn is_err(&self)->bool {
    (self.0).0.is_null() }}
```

---

16c  $\langle \text{SDL: C Functions } 5b \rangle + \equiv$  (5a)  $\triangleleft 15c \ 16f \triangleright$

---

```
fn SDL_FreeSurface(h: SurfaceHandle);
```

---

To load the images, we'll use a helper function out of libSDL2\_image.

16d  $\langle \text{SDL: Definitions } 5a \rangle + \equiv$  (4a)  $\triangleleft 14a \ 18a \triangleright$

---

```
#[link(name="SDL2_image")] extern
{ fn IMG_LoadPNG_RW<'a>(h: RWopsHandle)->ErrCheck<Surface<'a>>; }
```

---

16e  $\langle \text{SDL: SDL Methods } 7c \rangle + \equiv$  (7a)  $\triangleleft 16a \ 19b \triangleright$

---

```
pub fn load_png<'sdl>(&'sdl self, bytes: &'static [u8])
->Result<Surface<'sdl>, Box<dyn Error>>
{ Ok( unsafe { IMG_LoadPNG_RW(self.load_bytes(bytes)?.0)
    .unwrap()? } ) }
```

---

## 2.12 Texture loading

Now that we have an image loaded into CPU memory, we need to send it to the GPU so that it can be rendered. This will be a small game, so there's no need to worry about running out of graphics memory (I hope).

16f  $\langle \text{SDL: C Functions } 5b \rangle + \equiv$  (5a)  $\triangleleft 16c \ 17d \triangleright$

---

```
fn SDL_CreateTextureFromSurface<'r>
( r: RendererHandle, s: SurfaceHandle ) -> ErrCheck<Texture<'r>>;
fn SDL_DestroyTexture(tex: TextureHandle);
```

---



17a  $\langle \text{SDL: C Types } 5c \rangle + \equiv$  (5a)  $\triangleleft 16b$

---

```
#[repr(transparent)] #[derive(Debug, Copy, Clone)]
pub struct TextureHandle(*const c_void);
#[repr(transparent)] #[derive(Debug)]
pub struct Texture<'r>(TextureHandle, PhantomData<&'r Renderer<'r>>);
impl<'r> Drop for Texture<'r> { fn drop(&mut self) { unsafe {
    SDL_DestroyTexture(self.0); }}}
impl<'sdl> IsErr for Texture<'sdl> { fn is_err(&self)->bool {
    (self.0).0.is_null() }}
```

---

17b  $\langle \text{SDL: Renderer methods } 14b \rangle + \equiv$  (13a)  $\triangleleft 14b$

---

```
pub fn load_texture<'r>(&'r self, s: &Surface<'r>)
->Result<Texture<'r>, Box<dyn Error>>
{ Ok( unsafe { SDL_CreateTextureFromSurface(self.0, s.0)
    .unwrap()? })}
```

---

### 2.13 Asset Organization

Since we're going to have a bunch of different assets, it probably make sense to simplify the process of loading them a bit.

17c  $\langle \text{MAIN: Initialization } 7b \rangle + \equiv$  (4d)  $\triangleleft 13c \ 20 \triangleright$

---

```
macro_rules!
PNG { ( $file:tt ) => { renderer.load_texture(
    &sdl.load_png(include_bytes!(
    $file)))? } };
```

---

### 2.14 Texture rendering

The texture is fully loaded now, so we can find out information like its size and then use that to render it to the screen.

17d  $\langle \text{SDL: C Functions } 5b \rangle + \equiv$  (5a)  $\triangleleft 16f$

---

```
fn SDL_QueryTexture(tex: TextureHandle,
    format: *mut u32,
    access: *mut c_int,
    w: *mut c_int,
    h: *mut c_int) -> ErrCheck<c_int>;
fn SDL_RenderCopy(r: RenderHandle,
    t: TextureHandle,
    src: *const Rect,
    dst: *const Rect ) -> ErrCheck<c_int>;
```

---

18a	$\langle \text{SDL: Definitions } 5a \rangle + \equiv$	$(4a) \triangleleft 16d \ 19a \triangleright$
	<hr/> <pre> impl&lt;'r&gt; Texture&lt;'r&gt; {   pub fn get_size(&amp;self) -&gt; Result&lt;(c_int, c_int), Box&lt;dyn Error&gt;&gt;   {     let mut w:c_int = 0;     let mut h:c_int = 0;     unsafe { SDL_QueryTexture(self.0,                                 null_mut(),                                 null_mut(),                                 &amp;mut w as *mut _,                                 &amp;mut h as *mut _) }.unwrap()?;      Ok((w,h)) }} </pre> <hr/>	
18b	$\langle \text{SDL: Canvas methods } 15a \rangle + \equiv$	$(14a) \triangleleft 15a$
	<hr/> <pre> pub fn blit(&amp;self, tex:&amp;Texture, dest:&amp;Rect) -&gt; Result&lt;(), Box&lt;dyn Error&gt;&gt; { Ok( unsafe { SDL_RenderCopy((self.0).0,                                 tex.0,                                 null(),                                 dest as *const _).unwrap()?; })}  pub fn blit_center(&amp;self, tex:&amp;Texture, pos:(i32,i32)) -&gt; Result&lt;(), Box&lt;dyn Error&gt;&gt; { let (w,h) = tex.get_size()?;   let dest = Rect { x: pos.0 - w/2, y: pos.1 - h/2, w, h };   Ok(self.blit(tex, &amp;dest)?) } </pre> <hr/>	

## 2.15 Font Loading

In order to minimize the amount of graphics that need to be drawn by hand, I want to be able to use text strings programatically. For this, I'm going to use `libSDL2_ttf`.

19a (4a) <18a

---

```

<SDL: Definitions 5a>+≡

#[link(name="SDL2_ttf")] extern {
  fn TTF_Init()->c_int;
  fn TTF_OpenFontRW<'sdl>(src  : RWopsHandle,
                          free  : c_int,
                          ptsize: c_int      )->Font<'sdl>;
  fn TTF_CloseFont(f: FontHandle);
  fn TTF_SetFontStyle(f: FontHandle, style:c_int);
  fn TTF_RenderUTF8_Blended<'sdl>(f  : FontHandle,
                                   text: *const c_char,
                                   fg  : Color      )
                                   ->Surface<'sdl>;
}

#[repr(C)] #[derive(Debug, Copy, Clone)]
pub struct Color { pub r: u8, pub g: u8, pub b: u8, pub a: u8 }
#[repr(transparent)] #[derive(Debug, Copy, Clone)]
struct FontHandle(*const c_void);
#[repr(transparent)] #[derive(Debug)]
pub struct Font<'sdl>(FontHandle, PhantomData<&'sdl SDL>);
impl<'sdl> Drop for Font<'sdl> { fn drop(&mut self) { unsafe {
  TTF_CloseFont(self.0); }}}
impl<'sdl> Font<'sdl> {
  pub fn bold(self)->Self { unsafe { TTF_SetFontStyle(self.0,1);}; self}
  pub fn render(&self, text: &str, color: Color)
  ->Result<Surface<'sdl>, Box<dyn Error>> {
    let c_text = CString::new(text)?;
    let surf = unsafe { TTF_RenderUTF8_Blended(self.0,
                                              c_text.as_ptr(),
                                              color) };
    assert!(!(surf.0).0.is_null());
    Ok(surf)    }}

```

---

19b (7a) <16e

---

```

<SDL: SDL Methods 7c>+≡

pub fn load_font<'sdl>(&'sdl self, size: c_int, bytes: &'static [u8])
->Result<Font<'sdl>, Box<dyn Error>>
{
  static SDLTTF_INIT:Once = Once::new();
  SDLTTF_INIT.call_once(|| { unsafe { TTF_Init(); } });
  Ok(unsafe { TTF_OpenFontRW( self.load_bytes(bytes)?.0,
                             0, size
                             )}})

```

---

20

*MAIN: Initialization* 7b) + ≡

(4d) &lt;17c 30b&gt;

---

```

let font = sdl.load_font(36, include_bytes!("Go-Mono.ttf"))?;
let boldfont = sdl.load_font(36, include_bytes!("Go-Mono.ttf"))?.bold();
macro_rules!
DRAW_TEXT { ( $str:expr, $color:expr) =>
    { renderer.load_texture(
        &font.render($str, { let (r,g,b,a)=$color;
                               Color { r,g,b,a } })?))?};

macro_rules!
DRAW_BOLD { ( $str:expr, $color:expr) =>
    { renderer.load_texture(
        &boldfont.render($str, { let (r,g,b,a)=$color;
                                   Color { r,g,b,a } })?))?};

```

---

### 3 Async & Futures

In what was almost certainly a huge waste of time from a competition perspective, I spent several hours figuring out how Rust’s async system works. The plan is to let me write functions that alter the gamestate over longer timescales than a frame.

#### 3.1 Executor

Before it can do anything, we need to write a scheduler, which is responsible for calling `Future::poll()` at the top level to advance the program.

The `PendingTask` structure represents a top-level task that’s not currently blocked.

```
21a  <FUT: Definitions 21a>≡ (4a) 21b>
use std::cell::*;
use std::rc::Rc;
use std::pin::*;
use std::task::*;
use std::future::Future;
use std::fmt::Debug;

type PendingTaskInner =
    RefCell<Option<Pin<Box<dyn Future<Output = ()> + 'static>>>>>>;

#[derive(Clone)]
pub struct PendingTask(Rc<PendingTaskInner>);
impl PendingTask { <FUT: PendingTask methods 23a> }
```

For the queue of pending tasks, we can use a static `Vec`:

```
21b  <FUT: Definitions 21a>+≡ (4a) < 21a 22a>
fn task_queue() -> RefMut<'static, Vec<PendingTask>> {
    static mut QUEUE: Option<RefCell<Vec<PendingTask>>>> = None;
    static INIT: std::sync::Once = std::sync::Once::new();
    INIT.call_once(|| { unsafe { QUEUE = Some(RefCell::new(vec![])); }});
    unsafe { QUEUE.as_ref().unwrap().borrow_mut() }
}
```

The enqueue function is the primary entry point for adding new tasks to the queue. It can be called as `enqueue(async { ...; });`. The newly-created `PendingTask` is returned so that the caller can monitor the status, but there is no need to retain it— the task will remain scheduled even if the caller drops this return result.

22a  $\langle \text{FUT: Definitions } 21a \rangle + \equiv$  (4a)  $\langle 21b \ 22b \rangle$

---

```
pub fn enqueue(f: impl Future<Output = ()> + 'static)->PendingTask
{
    let result = PendingTask(
        Rc::new( RefCell::new( Some( Box::pin(f))) ) );
    task_queue().push( result.clone() );
    result
}
```

---

Its counterpart is `run_tasks`, which will advance all tasks as far as possible, and return as soon as they are all blocked. For our purposes, this will be an opportunity to render a frame and check for new input.

22b  $\langle \text{FUT: Definitions } 21a \rangle + \equiv$  (4a)  $\langle 22a \ 23b \rangle$

---

```
pub fn run_tasks() { while run_one_task() {} }
fn run_one_task()->bool
{
    let task = task_queue().pop();
    match task { None => false,
                Some(t) => {  $\langle \text{FUT: Run Task } t \ 22c \rangle$ 
                           true
                        } }
}
```

---

To actually run a task, the executor needs to construct a `Waker` to pass as the argument to `Future.poll`. The waker stores a single data pointer, which is raw and outside of the compiler's reference checking systems.

22c  $\langle \text{FUT: Run Task } t \ 22c \rangle \equiv$  (22b)

---

```
let w = unsafe { Waker::from_raw( RawWaker::new(t.clone_as_ptr(),
                                                &VTABLE
                                              ) ) };

let mut optfut = t.0.borrow_mut();
if optfut.is_some()
{
    let result = optfut.as_mut().unwrap()
        .as_mut().poll(&mut Context::from_waker(&w));
    if result.is_ready() { *optfut = None; }
}
```

---

In order to round-trip through the raw pointer API exposed by Rust, we need to define to- and from-pointer methods for our pending tasks. This is the primary reason they're wrapped in an Rc— I wanted to use its `from_raw` and `into_raw` methods to do the heavy lifting.

23a  $\langle \text{FUT: PendingTask methods } 23a \rangle \equiv$  (21a)

---

```

unsafe fn from_ptr(p: *const ()) -> PendingTask {
    PendingTask(Rc::from_raw(p as *const PendingTaskInner))
}
fn clone_as_ptr(&self) -> *const () {
    Rc::into_raw(Rc::clone(&self.0)) as *const ()
}

```

---

We also need to provide a static vtable of method implementations for the Waker. They're mostly straightforward from the descriptions in the documentation; the only thing to be careful of is paying attention to which should deallocate the stored data and which shouldn't.

23b  $\langle \text{FUT: Definitions } 21a \rangle + \equiv$  (4a)  $\triangleleft 22b \ 23c \triangleright$

---

```

static VTABLE: RawWakerVTable = RawWakerVTable::new(
    raw_clone, raw_wake, raw_wake_by_ref, raw_drop
);

```

---

`raw_clone` is called to duplicate a `RawWaker`. We reconstruct the pending task, clone it, and then convert the original task back to a raw pointer so that it doesn't get cleaned up at the end of the function.

23c  $\langle \text{FUT: Definitions } 21a \rangle + \equiv$  (4a)  $\triangleleft 23b \ 23d \triangleright$

---

```

unsafe fn raw_clone(p: *const ()) -> RawWaker {
    let t = PendingTask::from_ptr(p);
    let result = RawWaker::new(t.clone_as_ptr(), &VTABLE);
    assert_eq!(Rc::into_raw(t.0) as *const (), p);
    result
}

```

---

`raw_wake` reconstructs the task and pushes it onto the task queue to be run at the next opportunity.

23d  $\langle \text{FUT: Definitions } 21a \rangle + \equiv$  (4a)  $\triangleleft 23c \ 24a \triangleright$

---

```

unsafe fn raw_wake(p: *const ()) {
    let t = PendingTask::from_ptr(p);
    task_queue().push(t);
}

```

---

`raw_wake_by_ref` is a combination of the previous two. It pushes a clone of the reconstructed task onto the queue, and then converts the original back to a raw pointer so that it doesn't get dropped.

24a  $\langle \text{FUT: Definitions 21a} \rangle + \equiv$  (4a)  $\langle 23d \ 24b \rangle$

---

```
unsafe fn raw_wake_by_ref(p: *const ()) {
    let t = PendingTask::from_ptr(p);
    task_queue().push(t.clone());
    assert_eq!(Rc::into_raw(t.0) as *const (), p);
}
```

---

`raw_drop` is the simplest of all. It reconstructs the pending task and then does nothing else. When it leaves scope at the bottom of this function, the normal Drop machinery will take care of whatever needs to be done.

24b  $\langle \text{FUT: Definitions 21a} \rangle + \equiv$  (4a)  $\langle 24a \ 24c \rangle$

---

```
unsafe fn raw_drop(p: *const ()) {
    let _t = PendingTask::from_ptr(p);
}
```

---

### 3.2 Notification Futures

I plan to make only small pieces of the game code async, like sounds and animation. Most of the logic will be running on a traditional every-frame cadence, so I'm going to need a way to inject events from the traditional game loop into the async system. That means writing my own Future.

`NotificationPool` will handle dispatching a single value from the main loop to futures that are waiting for it.

24c  $\langle \text{FUT: Definitions 21a} \rangle + \equiv$  (4a)  $\langle 24b \ 25d \rangle$

---

```
use std::sync::{Arc, Mutex};
#[derive(Debug)]
pub struct NotificationPool<T: Clone+Debug>(Mutex<NotifPoolData<T>>);
#[derive(Debug)]
struct NotifPoolData<T> { result: Option<T>,
                          wakers: Vec<Option<Waker>> }
impl<T: Clone+Debug> NotificationPool<T>
{  $\langle \text{FUT: NotificationPool methods 25a} \rangle$  }
```

---

Thoughts from the next morning: This all works fine, but there is probably some kind of memory leak in here. At the very least, Futures will keep the notification pool alive even after anyone who might use it to post a notification is gone. If a notification has been posted, this is essentially a shared reference to the result, which is fine. If it hasn't, though, the pool holds a bunch of thread wakers that will never get called and hold references back into the pool.



In general, these will be held in an Arc so that their data can stick around as long as there are potential readers for it

25a  $\langle FUT: NotificationPool\ methods\ 25a \rangle \equiv$  (24c) 25b  $\triangleright$

---

```
pub fn new()->Arc<Self> {
  Arc::new(
    NotificationPool(
      Mutex::new( NotifPoolData{ result:None,
                               wakers:vec![] })))
}
```

---

The core functionality is notify, which sends its argument to all current and future tasks that might need the value. It's important that we release our self lock before attempting to run the Wakers, as that may cause code to run (via Drop methods) that attempts to obtain the lock.

25b  $\langle FUT: NotificationPool\ methods\ 25a \rangle + \equiv$  (24c)  $\triangleleft 25a\ 25c \triangleright$

---

```
pub fn notify(&self, result:T) {
  let wakers =
  { let mut self_mut = self.0.lock().unwrap();
    assert!(self_mut.result.is_none());
    self_mut.result = Some(result);
    std::mem::replace(&mut self_mut.wakers, vec![] ) };
  for maybe_w in wakers.into_iter()
  { if let Some(waker) = maybe_w { waker.wake(); }} }
```

---

And its corollary, wait, that returns a future which will eventually yield the notification.

25c  $\langle FUT: NotificationPool\ methods\ 25a \rangle + \equiv$  (24c)  $\triangleleft 25b \triangleright$

---

```
pub fn wait(self: Arc<Self>)->NotificationFuture<T> {
  let id: usize =
  { let mut self_mut = self.0.lock().unwrap();
    self_mut.wakers.push(None);
    self_mut.wakers.len() - 1 };
  NotificationFuture { id, pool: self }
```

---

25d  $\langle FUT: Definitions\ 21a \rangle + \equiv$  (4a)  $\triangleleft 24c\ 26a \triangleright$

---

```
#[derive(Debug)]
pub struct NotificationFuture<T:Clone+Debug>
{ pool: Arc<NotificationPool<T>>,
  id : usize }
```

---

Implementing the Future trait means we have to do one of two things when `poll` is called: If the result is ready, we return it. But if it isn't, we need to store a copy of the given Waker such that it gets woken up when the result becomes ready.

26a  $\langle \text{FUT: Definitions 21a} \rangle + \equiv$  (4a)  $\triangleleft$  25d 26b  $\triangleright$

---

```
impl<T:Clone+Debug> Future for NotificationFuture<T>
{
  type Output = T;
  fn poll(self: Pin<&mut Self>, cx: &mut Context) -> Poll<T>
  {
    let mut pool = self.pool.0.lock().unwrap();
    if let Some(ref result) = pool.result
    {
      Poll::Ready(result.clone())
    }
    else {
      pool.wakers[self.id] = Some( cx.waker()
                                   .clone() );
      Poll::Pending
    }
  }
}
```

---

Since the `NotificationFuture` is responsible for managing its waker in the `NotificationPool`, we should implement the `Drop` trait to clean that up as well. Because this requires taking the pool lock, we need to be careful about what gets dropped while the lock is held.

26b  $\langle \text{FUT: Definitions 21a} \rangle + \equiv$  (4a)  $\triangleleft$  26a 26c  $\triangleright$

---

```
impl<T:Clone+Debug> Drop for NotificationFuture<T>
{
  fn drop(&mut self)
  {
    let _w: Option<Waker> // Dont't drop while holding lock
    = {
      let mut pool = self.pool.0.lock().unwrap();
      if pool.wakers.len() > self.id
      {
        std::mem::replace(&mut pool.wakers[self.id],
                          None
                        )
      }
      else { None }
    };
  }
}
```

---

As this future is essentially just an `Arc` reference to the notification pool, we can easily make a reasonable `Clone` implementation as well. I have no idea if it'll actually be useful, as futures in general aren't cloneable.

26c  $\langle \text{FUT: Definitions 21a} \rangle + \equiv$  (4a)  $\triangleleft$  26b 27a  $\triangleright$

---

```
impl<T:Clone+Debug> Clone for NotificationFuture<T>
{
  fn clone(&self)->NotificationFuture<T>
  {
    Arc::clone(&self.pool).wait()
  }
}
```

---

### 3.3 Notification Streams

While sending a single value into the futures system is a good start, it's not really as expressive as I would like. If that's a compound value, however, we can do more interesting things with it. In particular, I want to be sending `Result<(T, Future<...>), Err>`. This will let me send an endless stream of values, and also throw errors in the receiving threads. Unfortunately, error types aren't guaranteed to be cloneable and I don't have time to be fighting with the type checker right now.

27a  $\langle \text{FUT: Definitions 21a} \rangle + \equiv$  (4a)  $\triangleleft 26c \ 28a \triangleright$

---

```
#[derive(Clone, Debug)]
struct StreamArg<T: Clone+Debug+'static>( T,
                                     NotificationFuture<StreamArg<T>>);
pub struct StreamSource<T: Clone+Debug+'static>
{   pool: Arc<NotificationPool<StreamArg<T>>> }
impl<T: Clone+Debug> StreamSource<T> {  $\langle \text{FUT: StreamSource methods 27b} \rangle$  }
```

---

Creating a new stream only requires making a new notification pool:

27b  $\langle \text{FUT: StreamSource methods 27b} \rangle \equiv$  (27a)  $\triangleright 27c \triangleright$

---

```
pub fn new() -> StreamSource<T>
{   StreamSource { pool: NotificationPool::new() } }
```

---

To send the next value, we make a new notification pool, and send the future for that pool along with the sent value. The old pool gets orphaned, and the reference counter will clean it up once there are no readers left that need its result.

27c  $\langle \text{FUT: StreamSource methods 27b} \rangle + \equiv$  (27a)  $\triangleleft 27b \ 27d \triangleright$

---

```
pub fn send(&mut self, msg: T)
{   let next_pool = NotificationPool::<StreamArg<T>>::new();
    let payload = StreamArg( msg, Arc::clone(&next_pool).wait());
    self.pool.notify(payload);
    self.pool = next_pool; }
```

---

To get the values back out, use `watch` to create a `StreamReader`.

27d  $\langle \text{FUT: StreamSource methods 27b} \rangle + \equiv$  (27a)  $\triangleleft 27c \triangleright$

---

```
pub fn watch(&self) -> StreamReader<T>
{   StreamReader { fut: Arc::clone(&self.pool).wait() } }
```

---

The `StreamReader` doesn't need to know anything except for the future that the `StreamSource` will be populating. Because this kind of future is cloneable, the stream is also cloneable. Cloning the stream will produce another stream that starts at the same point as when the original was cloned.

28a  $\langle \text{FUT: Definitions 21a} \rangle + \equiv$  (4a)  $\triangleleft$  27a

---

```
#[derive(Clone, Debug)]
pub struct StreamReader<T:Clone+Debug+'static>
{ fut: NotificationFuture<StreamArg<T>> }

impl<T:Clone+Debug+'static> StreamReader<T>
{  $\langle \text{FUT: StreamReader methods 28b} \rangle$  }
```

---

The primary method is `next()`. Awaiting this will produce a value that was provided to the source through `send()`, and the new future that represents the remaining stream values is swapped into the structure so that subsequent calls will get the newer values.

28b  $\langle \text{FUT: StreamReader methods 28b} \rangle \equiv$  (28a) 28c  $\triangleright$

---

```
pub async fn next(&mut self) -> T
{ let StreamArg(result, new_fut) = self.fut.clone().await;
  self.fut = new_fut;
  return result; }
```

---

As a convenience, we'll also define a `filter` method that can produce a stream with certain elements removed:

28c  $\langle \text{FUT: StreamReader methods 28b} \rangle + \equiv$  (28a)  $\triangleleft$  28b 29a  $\triangleright$

---

```
pub fn filter(self, f: impl Fn(&T)->bool + 'static) -> StreamReader<T>
{ let mut source = self;
  let mut newsource = StreamSource::<T>::new();
  let result = newsource.watch();
  enqueue(async move
  { loop { let val:T = source.next().await;
          if f(&val) { newsource.send(val); } });
  result }
```

---

And a map method to transform the contained datatype:

29a  $\langle \text{FUT: StreamReader methods 28b} \rangle + \equiv$  (28a)  $\triangleleft$  28c 29b  $\triangleright$

---

```
pub fn map<R>(self, f: impl Fn(T)->R + 'static) -> StreamReader<R>
where R: Debug+Clone+'static
{
  let mut source = self;
  let mut newsource = StreamSource::::new();
  let result = newsource.watch();
  enqueue(async move
  {
    loop { let val:T = source.next().await;
           newsource.send(f(val));          }
    }
  });
  result
}
```

---

29b  $\langle \text{FUT: StreamReader methods 28b} \rangle + \equiv$  (28a)  $\triangleleft$  29a

---

```
pub fn filtermap<R>(self, f: impl Fn(T)->Option<R> + 'static) -> StreamReader<R>
where R: Debug+Clone+'static
{
  let mut source = self;
  let mut newsource = StreamSource::::new();
  let result = newsource.watch();
  enqueue(async move
  {
    loop { let val:T = source.next().await;
           if let Some(r) = f(val) { newsource.send(r) }
    }
  });
  result
}
```

---

## 4 Putting it all together

We now have all the building blocks we need to put a game together. The bulk of the game logic will be async code in the `sim` module and have limited interaction with the underlying frame-by-frame activities. Input to the game logic comes from a single stream that reflects all of the events that happen, including frame timing information.

```
30a  <SIM: Definitions 30a>≡ (4a) 31a▷
    #[derive(Debug, Copy, Clone, Eq, PartialEq, Ord, PartialOrd)]
    pub enum Command { <SIM: Command variants 30c> }
```

```
30b  <MAIN: Initialization 7b>+≡ (4d) ◁20 34c▷
    let mut command_queue = StreamSource::<sim::Command>::new();
    let command_stream = command_queue.watch();
    enqueue(sim::start(command_stream));
```

`run_tasks()` returns when there is no more progress to be made, even if the tasks haven't completed. This means that we can have tasks that are infinite loops, and they'll get paused once they require information that doesn't exist yet. At that point, we continue the main game loop: render the updated game state to the screen, collect more user input, and then resume the game logic.

```
30c  <SIM: Command variants 30c>≡ (30a) 35f▷
    FrameAdvance( u32 ),
```

```
30d  <MAIN: Run Simulation 30d>≡ (4d)
    command_queue.send(sim::Command::FrameAdvance(sdl::ticks()));
    run_tasks();
```

#### 4.1 Game state management

The previous section discussed how we provide new information to the game logic, but we still need to get information out of it. Out of expediency, I have a single mutable global variable that contains the entire game state, protected by a `RefCell`. This is not a great solution because it requires careful programming to avoid game-crashing bugs, particularly never holding the state's reference across an `await`. Better would be some kind of hierarchical system, where each async thread can effectively own a small piece of the game state and periodically publish updates to the world.

31a  $\langle \text{SIM: Definitions } 30a \rangle + \equiv$  (4a)  $\triangleleft 30a \ 31b \triangleright$

---

```
use crate::fut::*;
use std::cell::RefCell;

#[derive(Debug, Default)]
pub struct GameState {  $\langle \text{SIM: GameState fields } 34a \rangle$  }
impl GameState
{ pub fn current() -> &'static RefCell<GameState>
  { static mut STATE: Option<RefCell<GameState>> = None;
    static INIT: std::sync::Once = std::sync::Once::new();
    INIT.call_once(
      || { unsafe
        { STATE = Some(RefCell::new(
          Default::default())); } } );
    unsafe { STATE.as_ref().unwrap() } }
```

---

This macro is just some syntactic sugar on top of `filtermap` for streams:

31b  $\langle \text{SIM: Definitions } 30a \rangle + \equiv$  (4a)  $\triangleleft 31a \ 32a \triangleright$

---

```
macro_rules! select
{ ($s:ident : $p:pat => $e:expr) => { select!(($s): $p => $e) };
  (($stream:expr): $pat:pat => $expr:expr) =>
  { ($stream).clone().filtermap(|value|
    { if let $pat = value { Some($expr) } else { None } }) }; }
```

---

The `start` method initializes the game state, spawns various worker tasks, and then exits.

```

32a  <SIM: Definitions 30a> +≡ (4a) <31b 34b>
pub async fn start(firehose: StreamReader<Command>) {
    use Command::*;
    { let mut gs = GameState::current().borrow_mut();
      <SIM: Initialize gs:GameState 34e>
    }

    use std::sync::atomic::*;
    use std::sync::atomic::Ordering::*;
    static T_PREV:AtomicU32 = AtomicU32::new(0);
    let timing = select!( firehose: FrameAdvance(t)
                          => (t, t-T_PREV.swap(t, Relaxed)) );
    <SIM: Start Workers 36b>
}

```

## 4.2 Render Order

So far, everything has pretty much been put together in the source code file in the order I came up with it. This works fine for things like function definitions, but for some things, the ordering is quite important. When rendering a frame, for instance, the code order is effectively the same as the stacking order of the things that are being drawn: items later in the list can overdraw things earlier on the list.

Because of this, I have consolidated the overall drawing order here. Each individual drawing task is enclosed in its own block so that defined variables don't leak out to the others.

```

32b  <MAIN: Draw Frame 32b>≡ (4d)
{ let game_state = sim::GameState::current().borrow();
  let canvas = renderer.start_frame((128,0,255))?;
  { <MAIN: Draw Terrain 34d> }
  { <MAIN: Draw Floor Markings 37e> }
  { <MAIN: Draw Player 35e> }
}

```



### 4.3 Coordinates

The gameplay happens on a hex grid. I'm using a rectangular coordinate system, where only half of the coordinates are valid spaces, in a checkerboard pattern.

I have a couple of helper functions to translate between world space and screen space:

```

33  <MAIN: Definitions 4d>+≡
                                     (4a) <4d
-----
fn grid_to_px(obj:(f32, f32), cam:(f32, f32))->(i32,i32)
{ (((obj.0-cam.0)*36.).floor() as i32+400,
  ((obj.1-cam.1)*21.).floor() as i32+300) }
fn igrd_to_px(obj:(usize, usize), cam:(f32, f32))->(i32,i32)
{ grid_to_px((obj.0 as f32, obj.1 as f32), cam) }
-----

```

## 5 Game Elements

Here you will find descriptions of all the game elements and logic.

### 5.1 Terrain

The terrain is stored in a two dimensional vector, with the  $y$  coordinate corresponding to the outermost vector.

```
34a <SIM: GameState fields 34a>≡ (31a) 35b>
pub terrain:Vec<Vec<Terrain>>,
```

The actual tile definition is an enum which contains two variants, one for empty spaces and another for walkable tiles.

```
34b <SIM: Definitions 30a>+≡ (4a) <32a 35a>
#[derive(Debug,Eq,PartialEq)]
pub enum Terrain { Empty, Floor }
```

Tiles are drawn from a PNG source, as shown to the right.



```
34c <MAIN: Initialization 7b>+≡ (4d) <30b 35d>
let tex_floor = PNG!("Flathex.png");
```

Because the terrain is stored in row-major order, it will be drawn top-down on the screen, so lower rows can draw over higher rows. This would allow for a 2.5D rendering scheme in the future.

```
34d <MAIN: Draw Terrain 34d>≡ (32b)
for y in 0..game_state.terrain.len()
{ for x in 0..game_state.terrain[y].len()
  { if game_state.terrain[y][x] == sim::Terrain::Floor
    { let center = igrd_to_px((x, y), game_state.camera);
      canvas.blit_center(&tex_floor, center)?; } } }
```

When setting up a new game, the terrain gets filled such that the black squares of a checkerboard have floor and the red squares are empty.

```
34e <SIM: Initialize gs:GameState 34e>≡ (32a) 35c>
for y in 0..40
{ gs.terrain.push(vec![]);
  for x in 0..40
  { gs.terrain[y].push( match (x^y)&1
                        { 0 => Terrain::Floor,
                          _ => Terrain::Empty } ); } }
```

## 5.2 Player

The player gets a dedicated struct in the game state, which currently only contains their position, but could easily be expanded later to accomodate future gameplay changes.

35a (4a) <34b 36c>  
 $\langle \text{SIM: Definitions } 30a \rangle + \equiv$   


---

```
#[derive(Debug,Default)]
pub struct Player { pub loc: (usize, usize) }
```

---

35b (31a) <34a 36d>  
 $\langle \text{SIM: GameState fields } 34a \rangle + \equiv$   


---

```
pub player:Player,
```

---

The starting location of the player is hardcoded to be the center of the play-field.

35c (32a) <34e 36e>  
 $\langle \text{SIM: Initialize gs:GameState } 34e \rangle + \equiv$   


---

```
gs.player.loc = (20,20);
```

---

In an homage to the text-based rougelikes (and a bid to avoid making real artwork), the player is represented by a blue @ sign. I couldn't find any shade that had enough contrast with the slightly blue-tinted ground tiles, so I simulate a dark outline by drawing a boldfaced copy first, with a normal-weight version on top.

35d (4d) <34c 38c>  
 $\langle \text{MAIN: Initialization } 7b \rangle + \equiv$   


---

```
let tex_player    = DRAW_TEXT!("@", (88,100,232,255));
let tex_player_bg = DRAW_BOLD!("@", (22,25,60,255));
```

---

35e (32b) 38d>  
 $\langle \text{MAIN: Draw Player } 35e \rangle \equiv$   


---

```
let loc = igrd_to_px(game_state.player.loc, game_state.camera);
canvas.blit_center(&tex_player_bg, loc)?;
canvas.blit_center(&tex_player,   loc)?;
```

---

MOVEMENT IS CONTROLLED by the keyboard keys q,w,e,a,s, and d. We use a Step command to let the runtime send these to us, and make the corresponding KeyDn events inject them into the command stream.

35f (30a) <30c>  
 $\langle \text{SIM: Command variants } 30c \rangle + \equiv$   


---

```
Step( i32, i32 ),
```

---

36a  $\langle \text{MAIN: Event Handlers } 11c \rangle + \equiv$  (10e)  $\triangleleft 11c$

---

```

SdlEvent::KeyDn(0x014, _) => { command_queue.send(Step(-1,-1)); }
SdlEvent::KeyDn(0x01a, _) => { command_queue.send(Step( 0,-2)); }
SdlEvent::KeyDn(0x008, _) => { command_queue.send(Step( 1,-1)); }
SdlEvent::KeyDn(0x004, _) => { command_queue.send(Step(-1, 1)); }
SdlEvent::KeyDn(0x016, _) => { command_queue.send(Step( 0, 2)); }
SdlEvent::KeyDn(0x007, _) => { command_queue.send(Step( 1, 1)); }
SdlEvent::KeyUp(_,_) => {}
SdlEvent::KeyDn(_,_) => {}

```

---

We also have a worker that handles the step requests, and moves the player if the requested space is clear.

36b  $\langle \text{SIM: Start Workers } 36b \rangle \equiv$  (32a) 37a  $\triangleright$

---

```

let steps = select!(firehose: Step(x, y) => (x, y));
enqueue(stepper(steps));

```

---

36c  $\langle \text{SIM: Definitions } 30a \rangle + \equiv$  (4a)  $\triangleleft 35a$  37b  $\triangleright$

---

```

async fn stepper(mut steps: StreamReader<(i32,i32)>) { loop
{   let (dx, dy) = steps.next().await;
    let mut gs = GameState::current().borrow_mut();
    let (x,y) = (gs.player.loc.0 as i32 + dx,
                gs.player.loc.1 as i32 + dy);
    if x < 0 || y < 0 { continue }
    let loc:(usize, usize) = (x as usize, y as usize);
    if loc.1 >= gs.terrain.len()      { continue };
    if loc.0 >= gs.terrain[loc.1].len() { continue };
    if gs.wrecks.contains(&loc)      { continue };
    match gs.terrain[loc.1][loc.0]
    {   Terrain::Floor => { gs.player.loc = loc; }
        Terrain::Empty => { continue }           };
     $\langle \text{SIM: Enemy Turn } 38e \rangle$ 
}}

```

---

### 5.3 Camera

The game uses an ad-hoc orthographic renderer, so there isn't much for a camera to do. It is represented by a tuple in the game state that is the world position of the center of the screen.

36d  $\langle \text{SIM: GameState fields } 34a \rangle + \equiv$  (31a)  $\triangleleft 35b$  37c  $\triangleright$

---

```

pub camera:(f32, f32),

```

---

36e  $\langle \text{SIM: Initialize gs:GameState } 34e \rangle + \equiv$  (32a)  $\triangleleft 35c$  37d  $\triangleright$

---

```

gs.camera = (15.,15.);

```

---

It should generally follow the player around. If it made the same jerky, jumping movements that the player did, things would quickly get disorienting. Using an exponential approach is convenient, as the calculation is can rely on just the difference between the current and desired positions, and the size of the time step. Every frame, I multiply the distance between the camera and the player by  $0.9999^{\delta t}$ , which seems to be a reasonable speed.

37a  $\langle \text{SIM: Start Workers } 36b \rangle + \equiv$  (32a)  $\triangleleft 36b$

---

```
enqueue(camera_track(timing.clone()));
```

---

37b  $\langle \text{SIM: Definitions } 30a \rangle + \equiv$  (4a)  $\triangleleft 36c \ 41 \triangleright$

---

```
async fn camera_track(mut times: StreamReader<(u32,u32)>) { loop
{   let (_t, dt) = times.next().await;
    let mut gs = GameState::current().borrow_mut();
    let coeff: f32 = 0.9999f32.powi(dt as i32);
    gs.camera.0 = gs.player.loc.0 as f32
                + coeff * (gs.camera.0 - gs.player.loc.0 as f32);
    gs.camera.1 = gs.player.loc.1 as f32
                + coeff * (gs.camera.1 - gs.player.loc.1 as f32); }}
```

---

#### 5.4 Floor markings

In addition to the tiles, the floor has some markings to teach the player how to move around and to provide a fixed point to help them stay oriented in the otherwise featureless plane the game is operating on at the moment.

37c  $\langle \text{SIM: GameState fields } 34a \rangle + \equiv$  (31a)  $\triangleleft 36d \ 38a \triangleright$

---

```
pub floor_marks:Vec<(f32,f32,&'static str)>,
```

---

37d  $\langle \text{SIM: Initialize gs:GameState } 34e \rangle + \equiv$  (32a)  $\triangleleft 36e \ 38b \triangleright$

---

```
for mark in [ (18.,18., "q"), (20.,16., "w"),
              (22.,18., "e"), (18.,22., "a"),
              (20.,24., "s"), (22.,22., "d") ].into_iter()
{ gs.floor_marks.push(*mark); }
```

---

I really should be caching the textures for these letters, but it was easier to re-render them every frame:

37e  $\langle \text{MAIN: Draw Floor Markings } 37e \rangle \equiv$  (32b)

---

```
for (x,y,s) in game_state.floor_marks.iter() {
    canvas.blit_center(&DRAW_TEXT!(s,(88,88,111,55)),
                      grid_to_px((*x,*y), game_state.camera));}
```

---

## 5.5 Enemies

Much like the player, the enemies are rendered as a red X. As there can be several enemies active at a time, they are tracked in a Vec in the game state.

```

38a  <SIM: GameState fields 34a>+≡ (31a) <37c 40a>
      pub enemies:Vec<usize, usize>,

38b  <SIM: Initialize gs:GameState 34e>+≡ (32a) <37d 40b>
      gs.enemies = vec![(25,9), (3,3)];

38c  <MAIN: Initialization 7b>+≡ (4d) <35d 40c>
      let tex_enemy    = DRAW_TEXT!("X", (232,100,88,255));
      let tex_enemy_bg = DRAW_BOLD!("X", (60,25,22,255));

38d  <MAIN: Draw Player 35e>+≡ (32b) <35e 40d>
      for world_loc in game_state.enemies.iter()
      { let loc = igrd_to_px(*world_loc, game_state.camera);
        canvas.blit_center(&tex_enemy_bg, loc)?;
        canvas.blit_center(&tex_enemy,    loc)?; }

```

The enemies all get a turn after any successful step by the player, so their movement logic also lives in stepper;

```

38e  <SIM: Enemy Turn 38e>≡ (36c)
      for i in 0..gs.enemies.len() {<SIM: Enemy i gives chase 39a>}

      let mut deadlist:Vec<usize> = vec![];
      for i in 0..gs.enemies.len()
      { let mut dead:bool = false;
        { <SIM: Enemy i checks for collisions 39b> }
        if dead { deadlist.push(i) };
      }
      while !deadlist.is_empty()
      { let id = deadlist.pop().unwrap();
        println!("kill {:?}", id);
        gs.enemies.remove(id); }
      if gs.enemies.len() < 2 {
        <SIM: Spawn new enemies 39c>
      }

```

The enemies always walk directly towards the player, as much as they are able.

39a  $\langle \text{SIM: Enemy } i \text{ gives chase } 39a \rangle \equiv$  (38e)

---

```

let dx = gs.player.loc.0 as i32 - gs.enemies[i].0 as i32;
let dy = gs.player.loc.1 as i32 - gs.enemies[i].1 as i32;
let mut e = &mut gs.enemies[i];
    if dx == 0 &&          dy > 0 {          e.1 += 2; }
else if dx < 0 && -3*dx < dy && dy > 0 {      e.1 += 2; }
else if dx > 0 &&  3*dx < dy && dy > 0 {      e.1 += 2; }
else if dx == 0 &&          dy < 0 {          e.1 -= 2; }
else if dx < 0 &&  3*dx > dy && dy < 0 {      e.1 -= 2; }
else if dx > 0 && -3*dx > dy && dy < 0 {      e.1 -= 2; }
else if dx > 0 &&          dy > 0 { e.0 += 1; e.1 += 1; }
else if dx > 0 &&          dy < 0 { e.0 += 1; e.1 -= 1; }
else if dx < 0 &&          dy < 0 { e.0 -= 1; e.1 -= 1; }
else if dx < 0 &&          dy > 0 { e.0 -= 1; e.1 += 1; }

```

---

39b  $\langle \text{SIM: Enemy } i \text{ checks for collisions } 39b \rangle \equiv$  (38e)

---

```

let loc = gs.enemies[i];
    if gs.player.loc == loc { return; } // Game Over
else if gs.wrecks.contains(&loc) { dead = true; }
else {
    for j in 0..gs.enemies.len() {
        if j != i && gs.enemies[j] == loc {
            dead = true; gs.wrecks.push(loc);
        }
    }
}

```

---

39c  $\langle \text{SIM: Spawn new enemies } 39c \rangle \equiv$  (38e)

---

```

use std::sync::atomic::*;
static NUM_ENEMIES:AtomicUsize = AtomicUsize::new(3);
for n in 0..NUM_ENEMIES.fetch_add(1, Ordering::Relaxed) {
    loop {
        let coord = random_coord();
        if gs.wrecks.contains(&coord) {continue};
        if gs.enemies.contains(&coord) {continue};
        if (gs.player.loc.0 as i32 - coord.0 as i32).abs()
            + (gs.player.loc.1 as i32 - coord.1 as i32).abs()
            < 5 { continue };
        gs.enemies.push(coord);
        break;
    }
}

```

---

## 5.6 Wrecks

Wrecks occur when two enemies run into each other. They will destroy future enemies that run into them. They are tracked in a Vec in the game state.

40a	$\langle \text{SIM: GameState fields } 34a \rangle + \equiv$ <hr/> <code>pub wrecks:Vec&lt;(usize, usize)&gt;,</code> <hr/>	(31a) $\triangleleft$ 38a
40b	$\langle \text{SIM: Initialize gs:GameState } 34e \rangle + \equiv$ <hr/> <hr/>	(32a) $\triangleleft$ 38b
40c	$\langle \text{MAIN: Initialization } 7b \rangle + \equiv$ <hr/> <code>let tex_wreck = DRAW_BOLD!("#", (44,44,44,255));</code> <hr/>	(4d) $\triangleleft$ 38c
40d	$\langle \text{MAIN: Draw Player } 35e \rangle + \equiv$ <hr/> <code>for world_loc in game_state.wrecks.iter()  { let loc = igrd_to_px(*world_loc, game_state.camera);    canvas.blit_center(&amp;tex_wreck, loc)?;  }</code> <hr/>	(32b) $\triangleleft$ 38d



## 6 Randomness

Rust doesn't have a random number generator in its standard library, but fortunately the only thing I need random samples of is grid coordinates. I used Python to generate a bunch of these:

41 *<SIM: Definitions 30a>+≡* (4a) <37b

---

```
static COORDS:(usize,usize);200 = [(17, 35), (6, 22), (13, 31),
(28, 2), (15, 29), (34, 26), (3, 9), (14, 12), (32, 34), (5, 31),
(25, 37), (22, 28), (17, 31), (0, 22), (6, 32), (9, 1), (11, 39),
(36, 32), (1, 33), (33, 1), (5, 15), (32, 10), (36, 26), (15, 1),
(16, 12), (14, 28), (38, 38), (21, 25), (31, 11), (9, 23),
(0, 38), (22, 4), (36, 38), (13, 1), (34, 6), (22, 34), (24, 30),
(34, 2), (21, 19), (34, 34), (17, 11), (23, 23), (39, 29),
(13, 39), (23, 37), (25, 3), (27, 13), (34, 30), (0, 0), (5, 9),
(38, 18), (28, 6), (20, 14), (35, 37), (34, 30), (30, 8),
(16, 30), (34, 36), (25, 1), (19, 27), (32, 14), (24, 14),
(14, 6), (34, 20), (25, 23), (6, 30), (30, 34), (27, 27),
(17, 17), (16, 38), (5, 33), (4, 16), (8, 6), (38, 2), (34, 16),
(8, 14), (13, 39), (19, 21), (35, 21), (33, 15), (38, 4),
(22, 8), (36, 16), (18, 22), (1, 9), (16, 12), (11, 5), (23, 29),
(3, 37), (22, 34), (17, 37), (18, 8), (35, 19), (0, 34),
(39, 35), (27, 17), (31, 27), (36, 28), (26, 6), (30, 22),
(23, 19), (39, 11), (10, 24), (16, 14), (1, 35), (3, 37),
(10, 2), (5, 3), (6, 6), (28, 6), (31, 9), (24, 18), (33, 31),
(11, 13), (11, 17), (4, 36), (31, 13), (7, 37), (28, 10),
(34, 26), (3, 17), (12, 20), (33, 35), (27, 33), (9, 27),
(17, 11), (4, 20), (38, 32), (36, 4), (9, 3), (28, 38), (13, 35),
(28, 2), (25, 33), (12, 18), (16, 6), (37, 15), (7, 21),
(19, 15), (18, 8), (2, 8), (23, 1), (25, 9), (12, 8), (13, 39),
(21, 17), (23, 11), (35, 1), (1, 7), (1, 7), (5, 15), (1, 1),
(10, 14), (27, 13), (39, 17), (9, 33), (37, 19), (34, 32),
(12, 6), (32, 4), (23, 29), (13, 29), (37, 27), (2, 38), (9, 21),
(12, 38), (8, 6), (27, 7), (22, 2), (26, 32), (36, 28), (5, 27),
(36, 26), (22, 28), (20, 14), (17, 21), (29, 39), (15, 29),
(39, 21), (1, 13), (14, 16), (13, 9), (1, 37), (29, 23), (21, 1),
(35, 29), (18, 32), (37, 39), (24, 36), (7, 5), (12, 26),
(14, 26), (39, 29), (2, 4), (39, 31), (9, 39), (13, 5), (20, 38),
(33, 37), (9, 17)];
use std::sync::atomic::*;
static NEXT_COORD:AtomicUsize = AtomicUsize::new(0);
fn random_coord() -> (usize, usize) {
    let coord = NEXT_COORD.fetch_add(1, Ordering::Relaxed);
    return COORDS[coord % COORDS.len()];
}
```

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