1) 
$$\lim_{n\to\infty} \frac{T_1(n)}{T_2(n)} = \lim_{n\to\infty} \frac{3\log_n + 3}{4\log(\log_n)} = \lim_{n\to\infty} \frac{3}{n} = \lim_{n\to\infty} \frac{3}{n} = \infty$$

So,  $T_2(n) \in O(T_1(n))$ 

2)  $\lim_{n\to\infty} \frac{T_1(n)}{T_2(n)} = \lim_{n\to\infty} \frac{3\log_n + 3}{n\log_n n} = \lim_{n\to\infty} \frac{3}{n\log_n n} = \lim_{n\to\infty} \frac{3}{n$ 

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b)

L) 
$$\lim_{n\to\infty} \frac{T_{2}(n)}{T_{3}(n)} = \lim_{n\to\infty} \frac{4\log(\log n)}{n^{3} + 8n^{4}} = \lim_{n\to\infty} \frac{1}{n \cdot \ln n} \cdot \frac{1}{5n^{4} + 32n^{3}} = 0$$

So,  $T_{2}(n) \in O(T_{3}(n))$ 

2)  $\lim_{n\to\infty} \frac{T_{2}(n)}{T_{4}(n)} = \lim_{n\to\infty} \frac{4\log(\log n)}{2000n + 1} = \lim_{n\to\infty} \frac{1}{n \cdot \ln n} \cdot \frac{1}{2000} = 0$ 

So,  $T_{2}(n) \in O(T_{4}(n))$ 

3)  $\lim_{n\to\infty} \frac{T_{2}(n)}{T_{3}(n)} = \lim_{n\to\infty} \frac{4\log(\log n)}{(n^{3} + n^{2} + n^$ 

3) 
$$\lim_{n\to\infty} \frac{T_3(n)}{T_6(n)} = \lim_{n\to\infty} \frac{n^6 + 3n^4}{3^n + n^2} = \lim_{n\to\infty} \frac{5n^4 + 32n^3}{3^n \cdot 1n^3 + 2n} \Rightarrow \lim_{n\to\infty} \frac{12n}{3^n \cdot (1n^3)^4} = 0$$

So,  $T_3(n) \in O(T_6(n))$ 

4)  $\lim_{n\to\infty} \frac{T_3(n)}{T_7(n)} = \lim_{n\to\infty} \frac{n^5 + 2n^4}{n^2 + 1000n} = \frac{5n^4 + 32n^3}{n^3 \cdot (1n^3)^4 + 1000n} \Rightarrow \frac{120}{n^6 \cdot (1n^4)^4} = 0$ 

So,  $T_2(n) \in O(T_3(n))$ 

5)  $\lim_{n\to\infty} \frac{T_2(n)}{T_7(n)} = \lim_{n\to\infty} \frac{n^5 + 2n^4}{n^2 + 1000} = \lim_{n\to\infty} \frac{9n^4 + 22n^3}{n^3 \cdot (1n^2 + 2n^2)} \Rightarrow \lim_{n\to\infty} \frac{120}{2^n \cdot (1n^2)^4} = 0$ 

So,  $T_3(n) \in O(T_3(n))$ 

4)  $\lim_{n\to\infty} \frac{T_4(n)}{T_7(n)} = \lim_{n\to\infty} \frac{2000n+1}{n^6 + 1000n} = \lim_{n\to\infty} \frac{2000}{2^n \cdot (1n^2 + 2n^2)} = 0$ 

So,  $T_4(n) \in O(T_5(n))$ 

2)  $\lim_{n\to\infty} \frac{T_4(n)}{T_7(n)} = \lim_{n\to\infty} \frac{2000n+1}{n^6 + 1000n} = \lim_{n\to\infty} \frac{2000}{2^n \cdot (1n^2 + 2n^2)} = 0$ 

So,  $T_4(n) \in O(T_7(n))$ 

3)  $\lim_{n\to\infty} \frac{T_4(n)}{T_7(n)} = \lim_{n\to\infty} \frac{2000n+1}{n^6 + 1000n} = \lim_{n\to\infty} \frac{2000}{2^n \cdot (1n^2 + 2n^2)} = 0$ 

So,  $T_4(n) \in O(T_7(n))$ 

4)  $\lim_{n\to\infty} \frac{T_7(n)}{T_7(n)} = \lim_{n\to\infty} \frac{2000n+1}{n^6 + 1000n} = \lim_{n\to\infty} \frac{2000}{2^n \cdot (1n^2 + 2n^2)} = 0$ 

So,  $T_4(n) \in O(T_7(n))$ 

2)  $\lim_{n\to\infty} \frac{T_7(n)}{T_7(n)} = \lim_{n\to\infty} \frac{2000n+1}{n^6 + 1000n} = \lim_{n\to\infty} \frac{2000}{2^n \cdot (1n^2 + 2n^2)} = 0$ 

So,  $T_4(n) \in O(T_7(n))$ 

2)  $\lim_{n\to\infty} \frac{T_7(n)}{T_7(n)} = \lim_{n\to\infty} \frac{2000n+1}{n^6 + 1000n} = \lim_{n\to\infty} \frac{2000}{2^n \cdot (1n^2 + 2n^2)} = 0$ 

So,  $T_4(n) \in O(T_7(n))$ 

2)  $\lim_{n\to\infty} \frac{T_7(n)}{T_7(n)} = \lim_{n\to\infty} \frac{2000n+1}{n^6 + 1000n} = \lim_{n\to\infty} \frac{2n^6 - 1000n}{2^n \cdot (1n^2 + 2n^2)} = 0$ 

So,  $T_4(n) \in O(T_7(n))$ 

2)  $\lim_{n\to\infty} \frac{T_7(n)}{T_7(n)} = \lim_{n\to\infty} \frac{2000n+1}{n^6 + 1000n} = \lim_{n\to\infty} \frac{2n^6 - 1000n}{2^n \cdot (1n^2 + 2n^2)} = 0$ 

2)  $\lim_{n\to\infty} \frac{T_7(n)}{T_7(n)} = \lim_{n\to\infty} \frac{2000n+1}{n^6 + 1000n} = \lim_{n\to\infty} \frac{2n^6 - 1000n}{2^n \cdot (1n^2 + 2n^2)} = 0$ 

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2)  $\lim_{n\to\infty} \frac{T_7(n)}{T_7(n)} = \lim_{n\to\infty} \frac{2000n+1}{n^6 + 1000n} = \lim_{n\to\infty} \frac{2n^6 - 1000n}{2^n \cdot (1n^2 + 2n^2)} = 0$ 

2)  $\lim_{n\to\infty} \frac{T_7(n)}{T_7(n)} = \lim_{n\to\infty} \frac{2n^6 - 1000n}{2^n \cdot (1n^2 + 2n^2)} = 0$ 

2)  $\lim_{n\to\infty} \frac{T_7(n)}{T_$ 

- 3) d) Gets an array and size of the array as parameters. Counts how many same elements there are. If the number of an element is bigger than the half of the size of the array, returns the element. If there are not identical elements as many as the half of the size of the array, returns -L.
  - b) Worst case occurs when the number of identical elements is less than the half of the size of the array. The time complexity of worst case is  $O(n^2)$ . Best case occurs when there are identical elements as many as the half of the size of the array to first element of array. The time complexity is O(n).
- 9) Gets an array and size of the array as parameters and pends the max number in the nums array. Creates a new array with size of one more of max number in nums array. Apter that, counts how many save elements are in nums array. For example, nums[0] is 2, then increases map[2] by L. Finally, wheck every elements in map array is one of the elements is bigger than half of size of the nums array, returns the repeating value. Otherwise, returns -1.
- b) The best and worst case time complexities of this function is O(n). Because first and second loop iterate n times regardless of best case and worst case. So, the general time complexity is S(n):
- 5) Time complexity of second purction is better than parst purction. Because in worst case, second purction has O(n) time complexity, whereas, first purction has  $O(n^2)$ . But their best cases are almost equal. In best case, second purction takes O(n) at most. In terms of space complexity, first function better than second purction. First function uses O(n) memory whereas second purction's space complexity can be anything according to the max number in array.

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```
Note: Works when there is no regative element.
I. maxA= ALOJ
                            Best and worst case time complexity is depends on
2. MaxB= BLO]
3. For i=1 ich
                         mak of n and m. So time complexity is
4. IF MaxA CALI]
                          Q(max(n,m)),
5. maxA = A[i]
6. For i=1 icm
7. IF MaxB < B []
3. maxB = B[i]
 9. Return max A * max B
1. Declare on array X size of 1+m integers
2. For i=0 icn+m
                                   Bust and worst case time complority
3. Max= ACO]
4. For J=L J <n
                               is depends on (men), made of mand n.
S. IF max < ALJ]
6.
                               So time complexity is Q((m+n), max(m,n))
         max = ASJJ
7-
        index=J
8.
           First = true
        For k=0 kkm
10.
            If make BLEJ
11.
          max= B[k]
12.
          index=k
13.
         first=false
14.
       IF pirst == true
15.
       Alindex] = -L
16.
       Bcindex] = -1
      X[i]= max
1. For i=n izposition
                              Best case occurs when the element is added
                            to the end of the array. It takes O(L). Worst
2. ACI] = ACI-L]
                            case occurs when the element is added to the beginning
3. Assposition ] = element
                            OF the array- It takes O(n). Avarage is O(n).
                             Best case occurs when the element is deleted
L. For i=position is a
                           from the end of the array. It takes O(4). Worst
      ACI] = ACI+L]
                           case occurs when the element is deleted from the
3. Asposition] = null
                           beginning of the array. It takes O(n). Avarage
                           is O(n).
                                                   Gaga Gayar
                                                    1901042629
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