

CS 33

Machine Programming (3)

Most of the slides in this lecture are either from or adapted from slides provided by the authors of the textbook “Computer Systems: A Programmer’s Perspective,” 2nd Edition and are provided from the website of Carnegie-Mellon University, course 15-213, taught by Randy Bryant and David O’Hallaron in Fall 2010. These slides are indicated “Supplied by CMU” in the notes section of the slides.

Today

- **Loops**
- **Switch statements**

Supplied by CMU.

“Do-While” Loop Example

C Code

```
int pcount_do(unsigned x)
{
    int result = 0;
    do {
        result += x & 0x1;
        x >>= 1;
    } while (x);
    return result;
}
```

Goto Version

```
int pcount_do(unsigned x)
{
    int result = 0;
loop:
    result += x & 0x1;
    x >>= 1;
    if (x)
        goto loop;
    return result;
}
```

- Count number of 1's in argument x (“popcount”)
- Use conditional branch either to continue looping or to exit loop

“Do-While” Loop Compilation

Goto Version

```
int pcount_do(unsigned x) {  
    int result = 0;  
loop:  
    result += x & 0x1;  
    x >>= 1;  
    if (x)  
        goto loop;  
    return result;  
}
```

Registers:

%edx	x
%eax	result

```
        movl    $0, %eax        # result = 0  
.L2:    # loop:  
        movl    %edx, %ecx  
        andl    $1, %ecx        # t = x & 1  
        addl    %ecx, %eax      # result += t  
        shrl    %edx            # x >>= 1  
        jne     .L2            # if !0, goto loop
```

Supplied by CMU.

Note that the condition codes are set as part of the execution of the shrl instruction.

General “Do-While” Translation

C Code

```
do  
    Body  
while (Test);
```

- **Body:**

```
{  
    Statement1;  
    Statement2;  
    ...  
    Statementn;  
}
```
- **Test returns integer**
 - = 0** interpreted as false
 - ≠ 0** interpreted as true

Goto Version

```
loop:  
    Body  
    if (Test)  
        goto loop
```

Supplied by CMU.

“While” Loop Example

C Code

```
int pcount_while(unsigned x) {  
    int result = 0;  
    while (x) {  
        result += x & 0x1;  
        x >>= 1;  
    }  
    return result;  
}
```

Goto Version

```
int pcount_do(unsigned x) {  
    int result = 0;  
    if (!x) goto done;  
loop:  
    result += x & 0x1;  
    x >>= 1;  
    if (x)  
        goto loop;  
done:  
    return result;  
}
```

- Is this code equivalent to the do-while version?
 - must jump out of loop if test fails

General “While” Translation

While version

```
while (Test)  
  Body
```



Do-While Version

```
if (!Test)  
  goto done;  
do  
  Body  
  while(Test) ;  
done:
```



Goto Version

```
if (!Test)  
  goto done;  
loop:  
  Body  
  if (Test)  
    goto loop;  
done:
```

Supplied by CMU.

“For” Loop Example

C Code

```
#define WSIZE 8*sizeof(int)
int pcount_for(unsigned x) {
    int i;
    int result = 0;
    for (i = 0; i < WSIZE; i++) {
        unsigned mask = 1 << i;
        result += (x & mask) != 0;
    }
    return result;
}
```

- Is this code equivalent to other versions?

“For” Loop Form

General Form

```
for (Init; Test; Update)  
    Body
```

```
for (i = 0; i < WSIZE; i++) {  
    unsigned mask = 1 << i;  
    result += (x & mask) != 0;  
}
```

Init

```
i = 0
```

Test

```
i < WSIZE
```

Update

```
i++
```

Body

```
{  
    unsigned mask = 1 << i;  
    result += (x & mask) != 0;  
}
```

“For” Loop → While Loop

For Version

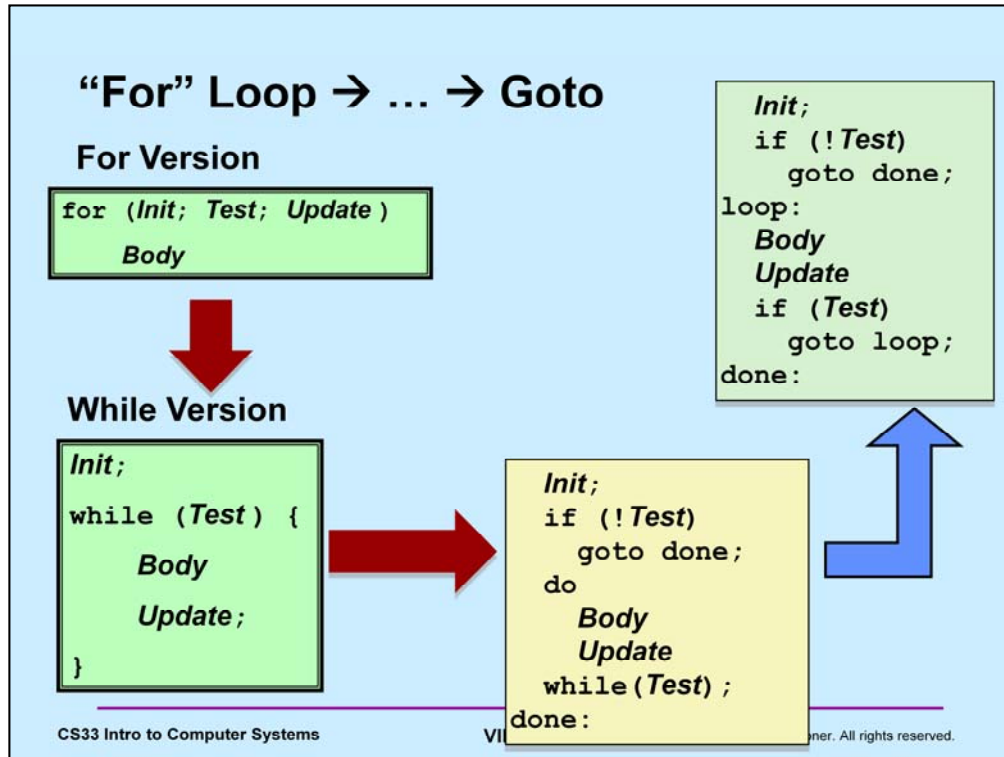
```
for (Init; Test; Update )  
    Body
```



While Version

```
Init;  
while (Test) {  
    Body  
    Update;  
}
```

Supplied by CMU.



Supplied by CMU.

“For” Loop Conversion Example

C Code

```
#define WSIZE 8*sizeof(int)
int pcount_for(unsigned x) {
    int i;
    int result = 0;
    for (i = 0; i < WSIZE; i++) {
        unsigned mask = 1 << i;
        result += (x & mask) != 0;
    }
    return result;
}
```

Initial test can be optimized away

Goto Version

```
int pcount_for_gt(unsigned x) {
    int i;
    int result = 0;
    i = 0;
if (!(i < WSIZE)) !Test
goto done;
loop:
    {
        unsigned mask = 1 << i;
        result += (x & mask) != 0;
    }
    i++;
    if (i < WSIZE)
        goto loop;
done:
    return result;
}
```

Supplied by CMU.

Today

- **Loops**
- **Switch statements**

Supplied by CMU.

Switch-Statement Example

```
long switch_eg
(long x, long y, long z)
{
    long w = 1;
    switch(x) {
        case 1:
            w = y*z;
            break;
        case 2:
            w = y/z;
            /* Fall Through */
        case 3:
            w += z;
            break;
        case 5:
        case 6:
            w -= z;
            break;
        default:
            w = 2;
    }
    return w;
}
```

- **Multiple case labels**
 - here: 5 & 6
- **Fall-through cases**
 - here: 2
- **Missing cases**
 - here: 4

Supplied by CMU.

Jump-Table Structure

Switch Form

```
switch(x) {  
  case val_0:  
    Block 0  
  case val_1:  
    Block 1  
    . . .  
  case val_n-1:  
    Block n-1  
}
```

Jump Table

jtab:	Targ0
	Targ1
	Targ2
	•
	•
	•
	Targn-1

Jump Targets

Targ0:	Code Block 0
Targ1:	Code Block 1
Targ2:	Code Block 2
	•
	•
	•
Targn-1:	Code Block n-1

Approximate Translation

```
target = JTab[x];  
goto *target;
```

Supplied by CMU.

Switch-Statement Example (IA32)

```
long switch_eg(long x, long y, long z)
{
    long w = 1;
    switch(x) {
        . . .
    }
    return w;
}
```

What range of values is covered by the default case?

Setup:

```
switch_eg:
    ...
    movl    8(%ebp), %ebx    # Setup
                                # %ebx = x
    movl    12(%ebp), %edx   # %edx = y
    movl    16(%ebp), %ecx   # %ecx = z
    cmpl    $6, %ebx        # Compare x:6
    ja      .L2              # If unsigned > goto default
    jmp     *.L7(, %ebx, 4)   # Goto *JTab[x] Note that w not initialized here
```

Supplied by CMU.


Switch-Statement Example (IA32)

```
long switch_eg(long x, long y, long z)
{
    long w = 1;
    switch(x) {
        . . .
    }
    return w;
}
```

Jump table

```
.section .rodata
.align 4
.L7:
.long .L2 # x = 0
.long .L3 # x = 1
.long .L4 # x = 2
.long .L5 # x = 3
.long .L2 # x = 4
.long .L6 # x = 5
.long .L6 # x = 6
```

Setup:

```
switch_eg:
    . . .                               # Setup
    movl 8(%ebp), %ebx                 # %ebx = x
    movl 12(%ebp), %edx                # %edx = y
    movl 16(%ebp), %ecx                # %ecx = z
    cmpl $6, %ebx                     # Compare x:6
    ja .L2                             # If unsigned > goto default
    Indirect
jump  jmp *.L7(,%ebx,4) # Goto *JTab[x]
```

Supplied by CMU.

Assembly-Setup Explanation

- Table structure

- each target requires 4 bytes
- base address at .L7

Jump table

```
.section .rodata
.align 4
.L7:
.long .L2 # x = 0
.long .L3 # x = 1
.long .L4 # x = 2
.long .L5 # x = 3
.long .L2 # x = 4
.long .L6 # x = 5
.long .L6 # x = 6
```

- Jumping

direct: `jmp .L2`

- jump target is denoted by label .L2

indirect: `jmp *.L7(, %ebx, 4)`

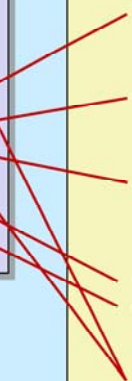
- start of jump table: .L7
- must scale by factor of 4 (labels have 32 bits = 4 Bytes on IA32)
- fetch target from effective address `.L7 + ebx*4`
 - » only for $0 \leq x \leq 6$

Jump Table

Jump table

```
.section .rodata
.align 4
.L7:
.long .L2 # x = 0
.long .L3 # x = 1
.long .L4 # x = 2
.long .L5 # x = 3
.long .L2 # x = 4
.long .L6 # x = 5
.long .L6 # x = 6
```

```
switch(x) {
case 1:      // .L3
    w = y*z;
    break;
case 2:      // .L4
    w = y/z;
    /* Fall Through */
case 3:      // .L5
    w += z;
    break;
case 5:
case 6:      // .L6
    w -= z;
    break;
default:    // .L2
    w = 2;
}
```



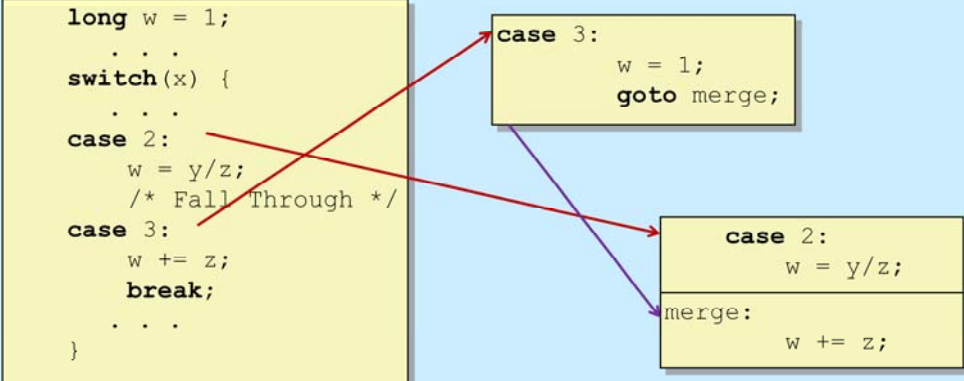
Supplied by CMU.

Handling Fall-Through

```
long w = 1;
...
switch(x) {
    ...
    case 2:
        w = y/z;
        /* Fall Through */
    case 3:
        w += z;
        break;
    ...
}
```

```
case 3:
    w = 1;
    goto merge;
```

```
case 2:
    w = y/z;
merge:
    w += z;
```



Supplied by CMU.

Code Blocks (Partial)

```
switch(x) {  
  case 1:      // .L3  
    w = y*z;  
    break;  
    . . .  
  case 3:      // .L5  
    w += z;  
    break;  
    . . .  
  default:    // .L2  
    w = 2;  
}
```

```
.L2:          # Default  
  movl $2, %eax # w = 2  
  jmp  .L8      # Goto done  
  
.L5:          # x == 3  
  movl $1, %eax # w = 1  
  jmp  .L9      # Goto merge  
  
.L3:          # x == 1  
  movl %ecx, %eax # z  
  imull %edx, %eax # w = y*z  
  jmp  .L8      # Goto done
```

Supplied by CMU.

Code Blocks (Rest)

```
switch(x) {  
    . . .  
    case 2: // .L4  
        w = y/z;  
        /* Fall Through */  
merge:    // .L9  
        w += z;  
        break;  
    case 5:  
    case 6: // .L6  
        w -= z;  
        break;  
}
```

```
.L4:                # x == 2  
    movl %edx, %eax  
    sarl $31, %edx  
    idivl %ecx      # w = y/z  
  
.L9:                # merge:  
    addl %ecx, %eax # w += z  
    jmp  .L8        # goto done  
  
.L6:                # x == 5, 6  
    movl $1, %eax  # w = 1  
    subl %ecx, %eax # w = 1-z
```

Supplied by CMU.

The code following the .L4 label requires some explanation. The `idivl` instruction is of a special form in that it takes a 64-bit dividend, which is implicitly assumed to reside in registers `edx` and `eax`. `y`, which we want to be the dividend, is in `edx`. It is copied to `eax` by the `movl` instruction. The `sarl` instruction propagates the sign bit of `edx` across the entire register. Thus, if one considers `edx` to contain the most-significant bits of the dividend and `eax` to contain the least-significant bits, the pair of registers now contains the 64-bit version of `y`. The `idivl` instruction computes the quotient of dividing this 64-bit value by the 32-bit value contained in register `ecx`, which is `z`. The quotient goes into register `eax` (implicitly) and the remainder goes into register `edx` (and is ignored).

Switch Code (Finish)

```
return w;
```

```
.L8:                # done:  
    popl %ebx  
    popl %ebp  
    ret
```

- **Noteworthy Features**

- jump table avoids sequencing through cases
 - » constant time, rather than linear
- use jump table to handle holes and duplicate tags
- use program sequencing to handle fall-through
- don't initialize w = 1 unless really need it

x86-64 Switch Implementation

- Same general idea, adapted to 64-bit code
- Table entries 64 bits (pointers)
- Cases use revised code

```
switch(x) {  
  case 1:      // .L3  
    w = y*z;  
    break;  
  . . .  
}
```

```
.L3:  
  movq    %rdx, %rax  
  imulq   %rsi, %rax  
  ret
```

Jump Table

```
.section .rodata  
.align 8  
.L7:  
.quad    .L2    # x = 0  
.quad    .L3    # x = 1  
.quad    .L4    # x = 2  
.quad    .L5    # x = 3  
.quad    .L2    # x = 4  
.quad    .L6    # x = 5  
.quad    .L6    # x = 6
```

Supplied by CMU.

IA32 Object Code

- Setup

- label .L2 becomes address 0x80483b9
- label .L7 becomes address 0x80484d0

Assembly code

```
switch_eg:
    . . .
    ja     .L2          # If unsigned > goto default
    jmp    *.L7(,%ebx,4) # Goto *JTab[x]
```

Disassembled object code

```
080483a0 <switch_eg>:
    . . .
    80483b0: 77 07                ja     80483b9 <switch_eg+0x19>
    80483b2: ff 24 9d d0 84 04 08 jmp    *0x80484d0(,%ebx,4)
```

Supplied by CMU.

IA32 Object Code (cont.)

- **Jump table**

- doesn't show up in disassembled code
- can inspect using gdb

`gdb switch`

`(gdb) x/7xw 0x80484d0`

- » examine 7 hexadecimal format "words" (4-bytes each)
- » use command "`help x`" to get format documentation

<code>0x80484d0:</code>	<code>0x080483b9</code>	<code>0x080483ca</code>	<code>0x080483d2</code>	<code>0x080483c0</code>
<code>0x80484e0:</code>	<code>0x080483b9</code>	<code>0x080483dd</code>	<code>0x080483dd</code>	

Supplied by CMU.

IA32 Object Code (cont.)

- Deciphering jump table

0x80484d0: 0x080483b9 0x080483ca 0x080483d2 0x080483c0
0x80484e0: 0x080483b9 0x080483dd 0x080483dd

Address	Value	x
0x80484d0	0x080483b9	0
0x80484d4	0x080483ca	1
0x80484d8	0x080483d2	2
0x80484dc	0x080483c0	3
0x80484e0	0x080483b9	4
0x80484e4	0x080483dd	5
0x80484e8	0x080483dd	6

Supplied by CMU.

Disassembled Targets

```
80483b9: b8 02 00 00 00      mov     $0x2,%eax
80483be: eb 24               jmp     80483e4 <switch_eg+0x44>
80483c0: b8 01 00 00 00      mov     $0x1,%eax
80483c5: 8d 76 00            lea     0x0(%esi),%esi # no-op
80483c8: eb 0f               jmp     80483d9 <switch_eg+0x39>
80483ca: 89 c8               mov     %ecx,%eax
80483cc: 0f af c2            imul    %edx,%eax
80483cf: 90                  nop     # no-op
80483d0: eb 12               jmp     80483e4 <switch_eg+0x44>
80483d2: 89 d0               mov     %edx,%eax
80483d4: c1 fa 1f            sar     $0x1f,%edx
80483d7: f7 f9              idiv    %ecx
80483d9: 01 c8               add     %ecx,%eax
80483db: eb 07               jmp     80483e4 <switch_eg+0x44>
80483dd: b8 01 00 00 00      mov     $0x1,%eax
80483e2: 29 c8               sub     %ecx,%eax
80483e4: 5b                  pop     %ebx
80483e5: 5d                  pop     %ebp
80483e6: c3                  ret
```

Matching Disassembled Targets

Value

0x80483b9

0x80483ca

0x80483d2

0x80483c0

0x80483b9

0x80483dd

0x80483dd

```

80483b9:  mov    $0x2,%eax
80483be:  jmp     80483e4 <switch_eg+0x44>
80483c0:  mov     $0x1,%eax
80483c5:  lea     0x0(%esi),%esi # no-op
80483c8:  jmp     80483d9 <switch_eg+0x39>
80483ca:  mov     %ecx,%eax
80483cc:  imul    %edx,%eax
80483cf:  nop                                     # no-op
80483d0:  jmp     80483e4 <switch_eg+0x44>
80483d2:  mov     %edx,%eax
80483d4:  sar     $0x1f,%edx
80483d7:  idiv    %ecx
80483d9:  add     %ecx,%eax
80483db:  jmp     80483e4 <switch_eg+0x44>
80483dd:  mov     $0x1,%eax
80483e2:  sub     %ecx,%eax
80483e4:  pop     %ebx
80483e5:  pop     %ebp
80483e6:  ret
    
```

Summarizing

- **C Control**
 - if-then-else
 - do-while
 - while, for
 - switch
- **Assembler Control**
 - conditional jump
 - conditional move
 - indirect jump
 - compiler generates code sequence to implement more complex control
- **Standard Techniques**
 - loops converted to do-while form
 - large switch statements use jump tables
 - sparse switch statements may use decision trees

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