# Orthogonal Gradient Boosting for Interpretable Additive Rule Ensembles Supplementary Information

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#### 1 PROOF OF THEOREM 4.3

The condition of Theorem 4.3 states that:

Let  $Q \in \mathbb{R}^{n \times (t-1)}$  be the selected query matrix, g the corresponding gradient vector after a full weight correction, and  $q^*$  be a maximizer of the **orthogonal gradient boosting objective** function defined by

$$obj_{ogb}(q) = \frac{|\mathbf{g}^T \mathbf{q}|}{\|\mathbf{q}_\perp\| + \epsilon}$$
 (1)

where  $\mathbf{q}_{\perp}$  is the projection of q onto the orthogonal complement of range Q.

### 1.1 Property a

PROPOSITION 1.1. For  $\epsilon \to 0$ ,  $[q_1, \dots, q_{t-1}, q^*]$  is the best approximation to  $[q_1, \dots, q_{t-1}, g]$ .

Proof. If  $\epsilon \to 0$ , then  $\mathrm{obj}_{\mathrm{ogb}}(q) \to \frac{|g^Tq|}{||q_\perp||}$ . If  $\mathbf{q}^*$  is a maximizer of  $\mathrm{obj}_{\mathrm{ogb}}$ , then as shown in Lemma 4.1,  $\mathbf{q}^*$  minimises the minimum distance from all

$$f \in \text{span}\{q_1, \cdots, q_{t-1}, g\}$$

to the subspace of

$$\operatorname{span}\{\mathbf{q}_1,\cdots,\mathbf{q}_{t-1},\mathbf{q}^*\}.$$

Therefore, the subspace spanned by  $[q_1, \dots, q_{t-1}, q^*]$  is the best approximation to the subspace spanned by  $[q_1, \dots, q_{t-1}, g]$ .

#### 1.2 Property b

Proposition 1.2. For  $\epsilon \to \infty$ ,  $q^*$  is also a maximizer of  $obj_{gs}$  and any maximizer of  $obj_{gs}$  is also a maximizer of  $obj_{ogb}$ .

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Proof. Let  $q_1$  and  $q_2$  be any two queries and denote by  $\operatorname{obj}_{\operatorname{ogb}}^{(\epsilon)}(q)$  the  $\operatorname{obj}_{\operatorname{ogb}}$ -value of q for a specific  $\epsilon$ . Then

$$\begin{split} & \lim_{\epsilon \to \infty} \epsilon \left( \operatorname{obj}_{\operatorname{ogb}}^{(\epsilon)}(q_1) - \operatorname{obj}_{\operatorname{ogb}}^{(\epsilon)}(q_2) \right) \\ & = \lim_{\epsilon \to \infty} \epsilon \left( \frac{|g^T \mathbf{q}_1|}{\|\mathbf{q}_1^\perp\| + \epsilon} - \frac{|g^T \mathbf{q}_2|}{\|\mathbf{q}_2^\perp\| + \epsilon} \right) \\ & = \lim_{\epsilon \to \infty} \left( \frac{|g^T \mathbf{q}_1|}{\|\mathbf{q}_1^\perp\| / \epsilon + 1} - \frac{|g^T \mathbf{q}_2|}{\|\mathbf{q}_2^\perp\| / \epsilon + 1} \right) \\ & = |g^T \mathbf{q}_1| - |g^T \mathbf{q}_2| \\ & = \operatorname{obj}_{\operatorname{gs}}(q_1) - \operatorname{obj}_{\operatorname{gs}}(q_2) \end{split}$$

Thus for large enough  $\epsilon$ , the signs of  $\operatorname{obj}_{\operatorname{ogb}}^{(\epsilon)}(q_1) - \operatorname{obj}_{\operatorname{ogb}}^{(\epsilon)}(q_2)$  and  $\operatorname{obj}_{\operatorname{gs}}(q_1) - \operatorname{obj}_{\operatorname{gs}}(q_2)$  agree. Therefore, a query q is a  $\operatorname{obj}_{\operatorname{gs}}$ -maximizer, i.e.,  $\operatorname{obj}_{\operatorname{gs}}(q) \geq \operatorname{obj}_{\operatorname{gs}}(q')$  for all  $q' \in Q$ , if and only if q is a  $\operatorname{obj}_{\operatorname{ogb}}$ -maximizer, i.e.,  $\operatorname{obj}_{\operatorname{ogb}}(q) \geq \operatorname{obj}_{\operatorname{ogb}}(q')$  for all  $q' \in Q$ .

## 1.3 Property c

Proposition 1.3. For  $\epsilon = 0$  and  $\|\mathbf{q}_{\perp}\| > 0$ , the ratio  $(\text{obj}_{ogb}(q)/\text{obj}_{gb}(q))^2$  is equal to  $1 + (\|\mathbf{q}_{\parallel}\|/\|\mathbf{q}_{\perp}\|)^2$ .

PROOF. If  $\epsilon = 0$  and  $||q_{\perp}|| > 0$ , then

$$\begin{split} \left(\frac{\text{obj}_{\text{ogb}}(q)}{\text{obj}_{\text{gb}}(q)}\right)^2 &= \frac{\frac{|\mathbf{g}^T \mathbf{q}|^2}{\|\mathbf{q}_\perp\|^2}}{\frac{|\mathbf{g}^T \mathbf{q}|^2}{\|\mathbf{q}\|^2}} = \frac{\|\mathbf{q}\|^2}{\|\mathbf{q}_\perp\|^2} \\ &= \frac{\|\mathbf{q}\|^2 + \|\mathbf{q}_\perp\|^2}{\|\mathbf{q}_\perp\|^2} \\ &= 1 + \left(\frac{\|\mathbf{q}_\parallel\|}{\|\mathbf{q}_\perp\|}\right)^2 \end{split}$$

## 1.4 Property d

Proposition 1.4. The objective value  $obj_{ogb}(q)$  is upper bounded by  $\|\mathbf{g}\|/(1+\epsilon)$ .

Proof. If we divide the numerator and denominator of  $obj_{ogb}(\mathbf{q})$  with  $\|\mathbf{q}_b ot\|$ , then we can get

$$obj_{ogb}(\mathbf{q}) = \frac{|\mathbf{g}^T \mathbf{q}|}{\|\mathbf{q}_\perp\| + \epsilon}$$
$$= \frac{\frac{|\mathbf{g}^T \mathbf{q}_\perp|}{\|\mathbf{q}_\perp\|}}{1 + \frac{\epsilon}{\|\mathbf{q}_\perp\|}}$$

 $\begin{array}{l} \text{because } \frac{|g^Tq|}{\|q_\bot\|} = \frac{\|g\|\|q_\bot\||\cos \langle g,q_\bot\rangle\,|}{\|q_\bot\|} = \|g\||\cos \langle g,q_\bot\rangle\,| \text{ is always less than or equal to } \|g\|, \end{array}$ 

$$obj_{ogb}(q) \le \frac{\|g\|}{1 + \frac{\epsilon}{\|q_{\perp}\|}}$$

as  $\|\mathbf{q}_{\perp}\|$  is upper bounded by the number of data points n,

$$\operatorname{obj}_{\operatorname{ogb}}(\mathbf{q}) \le \frac{\|\mathbf{g}\|}{1 + \frac{\epsilon}{n}}.$$

Let 
$$\epsilon' = \frac{\epsilon}{n}$$
, then

$$obj_{ogb}(q) \le \frac{\|g\|}{1 + \epsilon'}$$

# 2 GREEDY APPROXIMATION TO BOUNDARY FUNCTION

Solving the prefix optimization problem for a query described in Section 4.3 results in an approximation to the optimal objective function value for the best objective value which can be obtained by adding the query into the model.

We test 2000 instances for different gradient values to see the difference between the greedy approach, the pre-sorting greedy approach, and the actual optimal objective values.

Figure 1 shows the number of instances of ratios between the best objective values and the true optimal objective value. It shows that more than 80% of the best objective function values of prefixes have the same value with the optimal objective value, and more than 95% of the instances have an objective value more than 95% of the optimal objective. The smallest ratio between the result of greedy approach and the optimal solution is about 75%.

#### 3 COMPARISON OF TESTING RISKS

Table 1 compares the testing empirical risks of the same benchmark datasets in Section 5. Table 1 has the same format with Table ?? except that it does not contain the comparison of running times.

As discussed in Section 5, the test performance of these algorithms follows similar trends to the training performance.

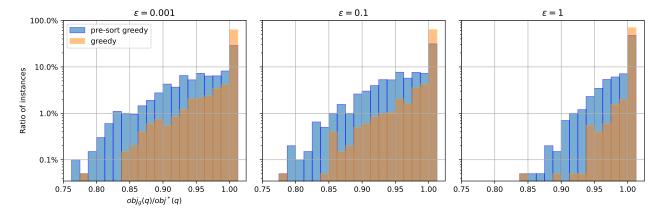


Figure 1: The number of instances of ratios between the best objective values obtained from the greedy search and the true optimal objective value.

Table 1: Comparison of Test Risks of Gradient Sum(S), Gradient boosting (G), XGBoost (X) and FCOGB (O) for benchmark datasets of classification (upper), regression (middle) and Poisson regression problems (lower).

DATASET	FEAT	ROW	$ar{ ilde{R}}_{ m O}$	FCOGB vs. GS					FCOGB vs. GB				FCOGB vs XGBoost					
				Z	$\Lambda_{ m SO}^{ m bc}$	$\bar{\Delta}_{SO}$	Δ	SO SO	Z	$\Lambda_{ m GO}^{ m bc}$	$\bar{\Delta}_{GO}$	Δ	wc GO	Δ	bc XO	$\bar{\Delta}_{XO}$	Δ	wc XO
TITANIC	7	1043	.712	.074	(17.6)	.035	.000	(2.4)	.147	(2.4)	.025	025	(28.2)	.147	(2.4)	.022	015	(4.4)
TIC-TAC-TOE	27	958	.751	.174	(23.6)	.101	058	(1.4)	.111	(6.2)	.060	.000	(2.4)	.089	(16.8)	.030	013	(28.4)
IRIS	4	150	.552	.141	(5.8)	089	180	(7.2)	.149	(2.4)	083	294	(5)	.149	(2.4)	058	147	(25.8)
BREAST	30	569	.352	.024	(7.8)	055	229	(7)	.009	(7.8)	011	026	(26)	.100	(7.8)	.031	006	(20.2)
WINE	13	178	.368	.314	(2.4)	.020	270	(5.8)	.409	(6)	.135	.003	(4)	.433	(6)	.090	043	(27)
IBM HR	32	1470	.217	.019	(6.2)	.003	005	(8.2)	.034	(15.4)	.004	.000	(1.4)	.065	(4)	.012	.004	(1.4)
TELCO CHURN	18	7043	.688	.058	(2.4)	.005	159	(1.4)	.037	(23.4)	.019	032	(8.2)	.017	(9.4)	.006	009	(12.4)
GENDER	20	3168	.999	.003	(2.4)	.001	.002	(3.2)	.000	(4.2)	.000	.000	(2.4)	.003	(11.4)	.001	.000	(2.4)
BANKNOTE	4	1372	.355	.142	(19.6)	.055	075	(8.2)	.133	(19.6)	.024	079	(9.4)	.120	(19.6)	.049	049	(9.4)
LIVER	6	345	.999	012	(3.8)	093	195	(29.8)	.057	(3.8)	024	067	(15.4)	.057	(3.8)	066	164	(29.4)
MAGIC	10	19020	.710	.056	(8.2)	.018	037	(4.2)	.018	(15.4)	.007	.000	(5.2)	.017	(18.4)	.007	003	(25.4)
ADULT	11	30162	.619	.146	(2.4)	.007	191	(1.4)	.059	(10.4)	.018	.000	(2.4)	.058	(4.2)	.011	.004	(20.4)
DIGITS5	64	3915	.381	.030	(4.2)	.014	008	(3.2)	.009	(4.2)	031	058	(19.4)	.070	(4.2)	004	034	(19.4)
INSURANCE	6	1338	.163	.104	(7.2)	.017	567	(1.4)	.172	(4.2)	.018	011	(5.2)	.172	(4.2)	.020	011	(5.2)
FRIEDMAN1	10	2000	.083	.013	(4)	002	012	(3.2)	.025	(4)	.006	.000	(1.4)	.025	(4)	.005	.000	(1.4)
FRIEDMAN2	4	10000	.149	.165	(3.2)	021	612	(1.4)	.084	(10.4)	.019	068	(5.2)	.080	(8.2)	.016	068	(5.2)
friedman3	4	5000	.060	.009	(4.6)	.003	012	(4)	.021	(4.6)	.002	.000	(1.4)	.021	(4.6)	.002	.000	(1.4)
WAGE	5	1379	.419	.017	(6.6)	021	048	(4.6)	.065	(6.6)	.000	027	(18.2)	.065	(6.6)	002	014	(21.4)
DEMOGRAPHICS	13	6876	.229	.011	(5.2)	.003	.000	(1.4)	.007	(3.2)	.002	.000	(1.4)	.007	(3.2)	.002	.000	(1.4)
GDP	1	35	.038	.003	(3.2)	.000	001	(5.2)	.003	(3.2)	.000	001	(5.2)	.003	(3.2)	.000	001	(5.2)
USED CARS	4	1770	.198	.178	(3.2)	019	549	(1.4)	.113	(14)	.055	.000	(3.2)	.089	(8.8)	.035	.000	(3.2)
DIABETES	10	442	.169	.026	(4.4)	004	033	(3.4)	.058	(4.4)	002	011	(29.8)	.058	(4.4)	.002	008	(29.4)
BOSTON	13	506	.097	.019	(4.4)	001	025	(3.8)	.044	(4.4)	.006	011	(8.6)	.044	(4.4)	.006	011	(9.2)
WORLD HAPPINESS	8	315	.051	.010	(5.2)	003	012	(3.8)	.013	(5.2)	.002	010	(4.4)	.023	(5.2)	.002	001	(23.6)
LIFE EXPECTANCY	21	1649	.041	.003	(4.2)	.000	001	(8.2)	.007	(4.2)	.001	.000	(20.4)	.007	(4.2)	.001	001	(26.4)
MOBILE PRICES	20	2000	.168	.168	(2.4)	008	648	(1.4)	.058	(3.2)	.002	004	(8.2)	.058	(3.2)	.004	.000	(4.2)
SUICIDE RATE	5	27820	.534	.081	(2.4)	.016	333	(1.4)	.018	(13.4)	.008	.000	(2.4)	.024	(11.4)	.009	.000	(2.4)
VIDEO GAMES	6	16327	.723	.000	(4.2)	.000	.000	(3.2)	.000	(10.4)	.000	.000	(3.2)	.000	(10.4)	.000	.000	(3.2)
RED WINE	11	1599	.048	.003	(4.2)	.000	002	(3.2)	.004	(4.2)	.001	.000	(1.4)	.004	(4.2)	.001	.000	(7.4)
COVID VIC	4	85	.185	.068	(3)	.030	062	(10)	.103	(10.6)	005	692	(2.8)	3.448	(3)	.418	.002	(25.8)
COVID	2	225	.515	.071	(6.8)	037	399	(2)	.279	(3.8)	.025	003	(27)	25.66	(6.8)	5.388	.290	(26.8)
BICYCLE	4	122	.505	.213	(7.8)	021	333	(3.4)	.054	(4.6)	035	337	(4)	.011	(7.8)	025	061	(22.2)
SHIPS	4	34	.338	.076	(3.2)	089	498	(2.4)	.160	(11)	.015	066	(23.2)	993.4	(3.2)	181.3	.605	(29.6)
SMOKING	2	36	.216	017	(16.6)	046	189	(6.6)	.075	(10.2)	.029	043	(4.2)	.950	(2.4)	.124	.035	(14.2)