


More Dynamic Programming: Matrix Chain Multiplication

Matrix Chain Multiplication

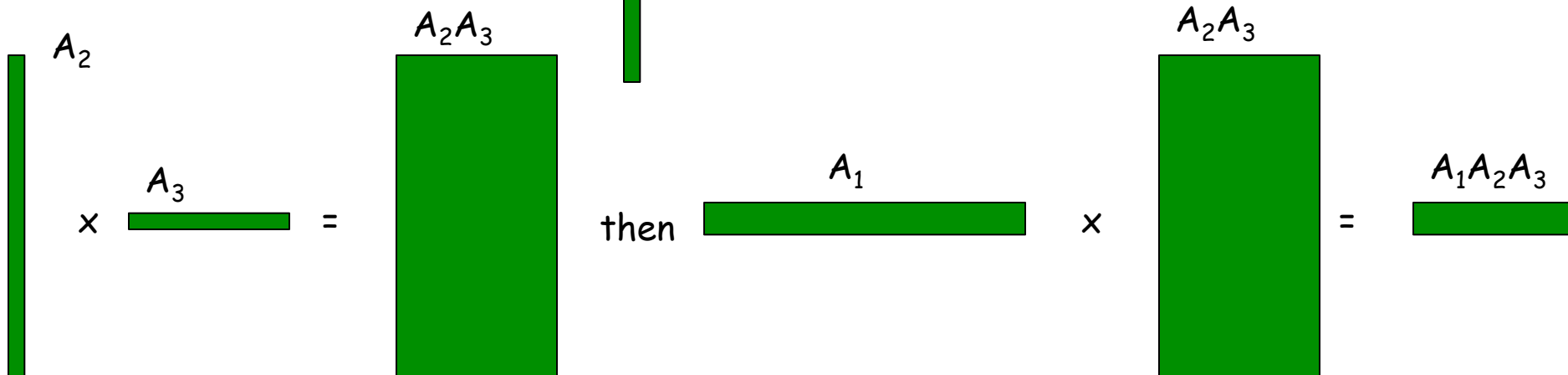
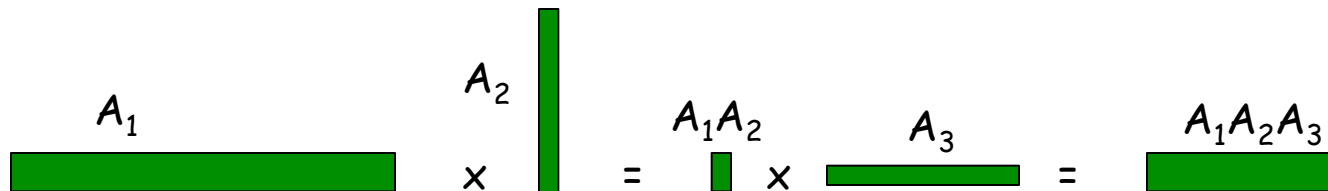
Problem Statement

- Given a chain of matrices to be multiplied together, determine a parenthesizing of the chain that minimizes the total number of steps required to complete the multiplication
- Matrix Multiplication
 - Given $A_{p \times q}$ and $B_{q \times r}$
 - AB requires pqr total element-wise multiplications
 - (and a similar number of additions, but we will ignore these)
 - Example: A_1 10×100 , A_2 100×5 , A_3 5×50
 -  $(A_1 A_2) A_3$ requires $10 \times 100 \times 5 + 10 \times 5 \times 50 = 7500$
 -  $A_1 (A_2 A_3)$ requires $100 \times 5 \times 50 + 10 \times 100 \times 50 = 75000$

Matrix Chain Multiplication

Example

- A_1 : 10×100
- A_2 : 100×5
- A_3 : 5×50
 - $(A_1 A_2) A_3$ requires $10 \times 100 \times 5 + 10 \times 5 \times 50 = 7500$
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Matrix Chain Multiplication

How Many Ways $P(n)$ of Parenthesizing are Possible for $A_1A_2A_3...A_n$?




- For $n = 1$, $P(n) = 1$
- For $n > 1$, the final product is the product of two fully parenthesized matrix subproducts, where the split can occur anywhere after matrix $1, 2, \dots, n-1$

$$P(n) = \sum_{k=1}^{n-1} P(k)P(n-k)$$

- $P(n) = \Omega(4^n / n^{3/2})$
 - Exponential time to check all possible parentheses placements

Matrix Chain Multiplication

Dynamic Programming Solution

- Optimal substructure
 - Total number of multiply operations needed for a given split:
 -  Multiplications needed for optimal solution to subproduct 1
 -  Multiplications needed for optimal solution to subproduct 2
 -  Multiplications needed to combine subproducts 1 and 2
 - The optimal solution will be the best choice among all possible splits

Matrix Chain Multiplication

Dynamic Programming Solution

- Suppose we have n matrices: $A_1 \times A_2 \times \dots \times A_n$
 - The dimensions of all the matrices can be completely specified with $a_0, a_1, a_2, \dots, a_n$
 - A_i has dimensions $a_{i-1} \times a_i$
 - Interior dimensions of neighboring matrices are the same

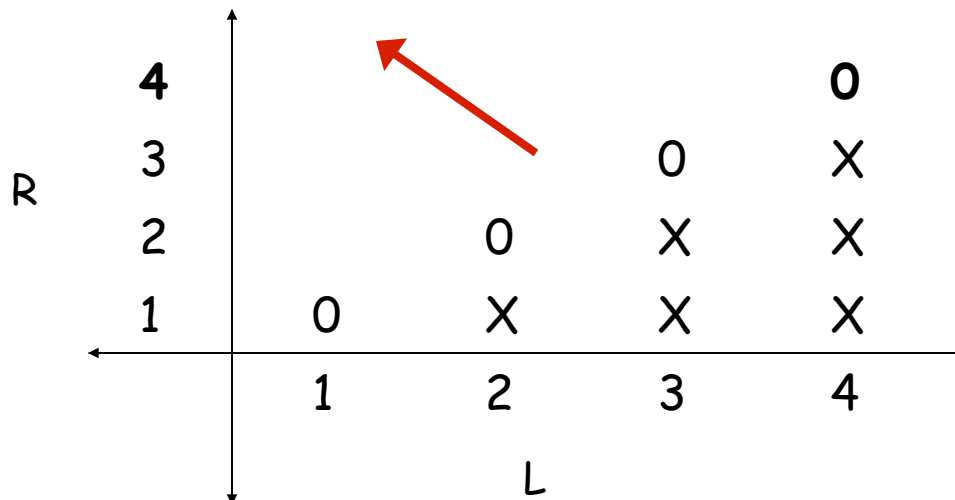
First Try

- $S[n] :=$ minimum number of multiplications needed to combine the first n matrices
 - $S[n] = \min_{k=1,2,\dots,n-1} (a_0 a_k a_n + S[k] + S[?])$
 - Need a subproblem corresponding to the optimal solution for multiplying $A_{k+1} \times \dots \times A_n$

Matrix Chain Multiplication

Dynamic Programming Solution: Second Try

- $S[L][R]$:= minimum number of multiplications needed to combine matrices A_L through A_R
 - $S[L][R] = \min_{k=L, L+1, \dots, R-1} (a_{L-1}a_k a_R + S[L][k] + S[k+1][R])$ for $1 \leq L < R \leq n$
 - $S[L][L] = 0$
- How is the matrix filled in?
 - From diagonal outward, using increasing size of $[L,R]$ interval



Matrix Chain Multiplication

Dynamic Programming Solution:

- Complexity: $O(n^3)$
 - Three nested loops each iterating over at most n items

MATRIX-CHAIN-MULTIPLICATION (a_0, \dots, a_n)

```
1. for L=1 to n do S[L][L] = 0
2. for d=1 to n-1 do
3.     for L=1 to n-d do
4.         R = L+d
5.         S[L][R] = ∞
6.         for k=L to R-1
7.             tmp = S[L][k] + S[k+1][R] + aL-1 · ak · aR
8.             if S[L][R] > tmp then S[L][R] = tmp
9. return S[1][n]
```

d loops over the size of interval
L loops over possible left endpoints