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1  ┌────────────────────────── MODULE Cure ───────────────────────────┐
    │ See ICDCS2016: “Cure: Strong Semantics Meets High Availability and Low Latency”. │
5  └────────────────────────── EXTENDS Naturals, Sequences, FiniteSets ───────────────────────────┘
6  ┌──────────────────────────┐
7  │  $Max(a, b) \triangleq \text{IF } a < b \text{ THEN } b \text{ ELSE } a$  │
8  │  $Min(S) \triangleq \text{CHOOSE } a \in S : \forall b \in S : a \leq b$  │
9  │  $Injective(f) \triangleq \forall a, b \in \text{DOMAIN } f : (a \neq b) \Rightarrow (f[a] \neq f[b])$  │
10 └──────────────────────────┘
11 CONSTANTS
12   Key,           the set of keys, ranged over by  $k \in Key$ 
13   Value,         the set of values, ranged over by  $v \in Value$ 
14   Client,        the set of clients, ranged over by  $c \in Client$ 
15   Partition,     the set of partitions, ranged over by  $p \in Partition$ 
16   Datacenter,    the set of datacenters, ranged over by  $d \in Datacenter$ 
17   KeySharding,   the mapping from Key to Partition
18   ClientAttachment the mapping from Client to Datacenter
20  $NotVal \triangleq \text{CHOOSE } v : v \notin Value$ 
22 ASSUME
23    $\wedge KeySharding \in [Key \rightarrow Partition]$ 
24    $\wedge ClientAttachment \in [Client \rightarrow Datacenter]$ 
25 ┌──────────────────────────┐
26 VARIABLES
27   At the client side:
28   cvc, cvc[c]: the vector clock of client  $c \in Client$ 
29   At the server side (each for partition  $p \in Partition$  in  $d \in Datacenter$ ):
30   clock, clock[p][d]: the current clock
31   pvc, pvc[p][d]: the vector clock
32   css, css[p][d]: the stable snapshot
33   store, store[p][d]: the kv store
34   communication:
35   msgs, the set of messages in transit
36   incoming fifo[p][d]: incoming FIFO channel; for propagating updates and heartbeats
38  $cVars \triangleq \langle cvc \rangle$ 
39  $sVars \triangleq \langle clock, pvc, css, store \rangle$ 
40  $mVars \triangleq \langle msgs, incoming \rangle$ 
41  $vars \triangleq \langle cvc, clock, pvc, css, store, msgs, incoming \rangle$ 
42 ┌──────────────────────────┐
43 │  $Clock \triangleq Nat$  │
44 │  $VC \triangleq [Datacenter \rightarrow Clock]$  vector clock with an entry per datacenter  $d \in Datacenter$  │
45 │  $VCInit \triangleq [d \in Datacenter \mapsto 0]$  │
46 │  $Merge(vc1, vc2) \triangleq [d \in Datacenter \mapsto Max(vc1[d], vc2[d])]$  │
47 │  $KVTuple \triangleq [key : Key, val : Value \cup \{NotVal\}, vc : VC]$  │

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49  $DC \triangleq \text{Cardinality}(\text{Datacenter})$ 
50  $DCIndex \triangleq \text{CHOOSE } f \in [1 \dots DC \rightarrow \text{Datacenter}] : \text{Injective}(f)$ 
51  $LTE(vc1, vc2) \triangleq$  less-than-or-equal-to comparator for vector clocks
52   LET RECURSIVE  $LTEHelper(-, -, -)$ 
53      $LTEHelper(vc1h, vc2h, index) \triangleq$ 
54       IF  $index > DC$  THEN TRUE  $EQ$ 
55       ELSE LET  $d \triangleq DCIndex[index]$ 
56         IN CASE  $vc1h[d] < vc2h[d] \rightarrow$  TRUE  $LT$ 
57            $\square vc1h[d] > vc2h[d] \rightarrow$  FALSE  $GT$ 
58            $\square$  OTHER  $\rightarrow LTEHelper(vc1h, vc2h, index + 1)$ 
59   IN  $LTEHelper(vc1, vc2, 1)$ 

61  $Message \triangleq$ 
62    $[type : \{\text{"ReadRequest"}\}, key : Key, vc : VC, c : Client, p : Partition, d : Datacenter]$ 
63    $\cup [type : \{\text{"ReadReply"}\}, val : Value \cup \{\text{NotVal}\}, vc : VC, c : Client]$ 
64    $\cup [type : \{\text{"UpdateRequest"}\}, key : Key, val : Value, vc : VC, c : Client, p : Partition, d : Datacenter]$ 
65    $\cup [type : \{\text{"UpdateReply"}\}, ts : Clock, c : Client, d : Datacenter]$ 
66    $\cup [type : \{\text{"Replicate"}\}, d : Datacenter, kv : KVTuple]$ 
67    $\cup [type : \{\text{"Heartbeat"}\}, d : Datacenter, ts : Clock]$ 

69  $TypeOK \triangleq$ 
70    $\wedge cvc \in [Client \rightarrow VC]$ 
71    $\wedge clock \in [Partition \rightarrow [Datacenter \rightarrow Clock]]$ 
72    $\wedge pvc \in [Partition \rightarrow [Datacenter \rightarrow VC]]$ 
73    $\wedge css \in [Partition \rightarrow [Datacenter \rightarrow VC]]$ 
74    $\wedge store \in [Partition \rightarrow [Datacenter \rightarrow \text{SUBSET } KVTuple]]$ 
75    $\wedge msgs \subseteq Message$ 
76    $\wedge incoming \in [Partition \rightarrow [Datacenter \rightarrow Seq(Message)]]$ 
77 |-----|
78  $Init \triangleq$ 
79    $\wedge cvc = [c \in Client \mapsto VCInit]$ 
80    $\wedge clock = [p \in Partition \mapsto [d \in Datacenter \mapsto 0]]$ 
81    $\wedge pvc = [p \in Partition \mapsto [d \in Datacenter \mapsto VCInit]]$ 
82    $\wedge css = [p \in Partition \mapsto [d \in Datacenter \mapsto VCInit]]$ 
83    $\wedge store = [p \in Partition \mapsto [d \in Datacenter \mapsto$ 
84      $[key : \{k \in Key : KeySharding[k] = p\}, val : \{\text{NotVal}\}, vc : \{VCInit\}]]]$ 
85    $\wedge msgs = \{\}$ 
86    $\wedge incoming = [p \in Partition \mapsto [d \in Datacenter \mapsto \langle \rangle]]$ 
87 |-----|
88  $Send(m) \triangleq msgs' = msgs \cup \{m\}$ 
89  $SendAndDelete(sm, dm) \triangleq msgs' = (msgs \cup \{sm\}) \setminus \{dm\}$ 

91  $CanIssue(c) \triangleq \forall m \in msgs :$ 
92    $m.type \in \{\text{"ReadRequest"}, \text{"ReadReply"}, \text{"UpdateRequest"}, \text{"UpdateReply"}\} \Rightarrow m.c \neq c$ 
93 |-----|
94 Client operations at client  $c \in Client.$ 

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96  $Read(c, k) \triangleq$   $c \in Client$  reads from  $k \in Key$ 
97  $\wedge CanIssue(c)$ 
98  $\wedge Send([type \mapsto \text{"ReadRequest"}, key \mapsto k, vc \mapsto cvc[c],$ 
99  $c \mapsto c, p \mapsto KeySharding[k], d \mapsto ClientAttachment[c]])$ 
100  $\wedge UNCHANGED \langle cVars, sVars, incoming \rangle$ 

102  $ReadReply(c) \triangleq$   $c \in Client$  handles the reply to its read request
103  $\wedge \exists m \in msgs :$ 
104  $\wedge m.type = \text{"ReadReply"} \wedge m.c = c$  such  $m$  is unique due to well-formedness
105  $\wedge cvc' = [cvc \text{ EXCEPT } ![c] = Merge(m.vc, @)]$ 
106  $\wedge msgs' = msgs \setminus \{m\}$ 
107  $\wedge UNCHANGED \langle sVars, incoming \rangle$ 

109  $Update(c, k, v) \triangleq$   $c \in Client$  updates  $k \in Key$  with  $v \in Value$ 
110  $\wedge CanIssue(c)$ 
111  $\wedge Send([type \mapsto \text{"UpdateRequest"}, key \mapsto k, val \mapsto v,$ 
112  $vc \mapsto cvc[c], c \mapsto c, p \mapsto KeySharding[k], d \mapsto ClientAttachment[c]])$ 
113  $\wedge UNCHANGED \langle cVars, sVars, incoming \rangle$ 

115  $UpdateReply(c) \triangleq$   $c \in Client$  handles the reply to its update request
116  $\wedge \exists m \in msgs :$ 
117  $\wedge m.type = \text{"UpdateReply"} \wedge m.c = c$  such  $m$  is unique due to well-formedness
118  $\wedge cvc' = [cvc \text{ EXCEPT } ![c][m.d] = m.ts]$ 
119  $\wedge msgs' = msgs \setminus \{m\}$ 
120  $\wedge UNCHANGED \langle sVars, incoming \rangle$ 

121 |-----|
122  $Server\ operations\ at\ partition\ p \in Partition\ in\ datacenter\ d \in Datacenter.$ 

124  $ReadRequest(p, d) \triangleq$  handle a "ReadRequest"
125  $\wedge \exists m \in msgs :$ 
126  $\wedge m.type = \text{"ReadRequest"} \wedge m.p = p \wedge m.d = d$ 
127  $\wedge css' = [css \text{ EXCEPT } ![p][d] = Merge(m.vc, @)]$ 
128  $\wedge LET\ kvs \triangleq \{kv \in store[p][d] :$ 
129  $\wedge kv.key = m.key$ 
130  $\wedge \forall dc \in Datacenter \setminus \{d\} : kv.vc[dc] \leq css'[p][d][dc]\}$ 
131  $lkv \triangleq CHOOSE\ kv \in kvs : \forall akv \in kvs : LTE(akv.vc, kv.vc)$ 
132  $IN\ SendAndDelete([type \mapsto \text{"ReadReply"}, val \mapsto lkv.val, vc \mapsto lkv.vc, c \mapsto m.c], m)$ 
133  $\wedge UNCHANGED \langle cVars, clock, pvc, store, incoming \rangle$ 

135  $UpdateRequest(p, d) \triangleq$  handle a "UpdateRequest"
136  $\wedge \exists m \in msgs :$ 
137  $\wedge m.type = \text{"UpdateRequest"} \wedge m.p = p \wedge m.d = d$ 
138  $\wedge m.vc[d] < clock[p][d]$  waiting condition; (" $\leq$ " strengthened to " $<$ ")
139  $\wedge css' = [css \text{ EXCEPT } ![p][d] = Merge(m.vc, @)]$ 
140  $\wedge LET\ kv \triangleq [key \mapsto m.key, val \mapsto m.val,$ 
141  $vc \mapsto [m.vc \text{ EXCEPT } ![d] = clock[p][d]]]$ 

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142   IN     $\wedge store' = [store \text{ EXCEPT } ![p][d] = @ \cup \{kv\}]$ 
143          $\wedge SendAndDelete([type \mapsto \text{"UpdateReply"}, ts \mapsto clock[p][d], c \mapsto m.c, d \mapsto d], m)$ 
144          $\wedge incoming' = [incoming \text{ EXCEPT } ![p] = [dc \in Datacenter \mapsto$ 
145             IF  $dc = d$  THEN  $@[dc]$  ELSE  $Append(@[dc], [type \mapsto \text{"Replicate"}, d \mapsto d, kv \mapsto kv])]$ 
146          $\wedge \text{UNCHANGED } \langle cVars, clock, pvc \rangle$ 

148  $Replicate(p, d) \triangleq$  handle a "Replicate"
149      $\wedge incoming[p][d] \neq \langle \rangle$ 
150      $\wedge \text{LET } m \triangleq Head(incoming[p][d])$ 
151     IN     $\wedge m.type = \text{"Replicate"}$ 
152            $\wedge store' = [store \text{ EXCEPT } ![p][d] = @ \cup \{m.kv\}]$ 
153            $\wedge pvc' = [pvc \text{ EXCEPT } ![p][d][m.d] = m.kv.vc[m.d]]$ 
154            $\wedge incoming' = [incoming \text{ EXCEPT } ![p][d] = Tail(@)]$ 
155      $\wedge \text{UNCHANGED } \langle cVars, cvc, clock, css, msgs \rangle$ 

157  $Heartbeat(p, d) \triangleq$  handle a "Heartbeat"
158      $\wedge incoming[p][d] \neq \langle \rangle$ 
159      $\wedge \text{LET } m \triangleq Head(incoming[p][d])$ 
160     IN     $\wedge m.type = \text{"Heartbeat"}$ 
161            $\wedge pvc' = [pvc \text{ EXCEPT } ![p][d][m.d] = m.ts]$ 
162            $\wedge incoming' = [incoming \text{ EXCEPT } ![p][d] = Tail(@)]$ 
163      $\wedge \text{UNCHANGED } \langle cVars, cvc, clock, css, store, msgs \rangle$ 

164 |-----|
165 Clock management at partition  $p \in Partition$  in datacenter  $d \in Datacenter$ 
166  $Tick(p, d) \triangleq$   $clock[p][d]$  ticks
167      $\wedge clock' = [clock \text{ EXCEPT } ![p][d] = @ + 1]$ 
168      $\wedge pvc' = [pvc \text{ EXCEPT } ![p][d][d] = clock'[p][d]]$ 
169      $\wedge incoming' = [incoming \text{ EXCEPT } ![p] = [dc \in Datacenter \mapsto$ 
170         IF  $dc = d$  THEN  $@[dc]$  ELSE  $Append(@[dc], [type \mapsto \text{"Heartbeat"}, d \mapsto d, ts \mapsto pvc'[p][d][d]])]$ 
171      $\wedge \text{UNCHANGED } \langle cVars, cvc, css, store, msgs \rangle$ 

173  $UpdateCSS(p, d) \triangleq$  update  $css[p][d]$ 
174      $\wedge css' = [css \text{ EXCEPT } ![p][d] =$ 
175          $[dc \in Datacenter \mapsto Min(\{pvc[pp][d][dc] : pp \in Partition\})]$ 
176      $\wedge \text{UNCHANGED } \langle cVars, mVars, clock, pvc, store \rangle$ 

177 |-----|
178  $Next \triangleq$ 
179      $\vee \exists c \in Client, k \in Key : Read(c, k)$ 
180      $\vee \exists c \in Client, k \in Key, v \in Value : Update(c, k, v)$ 
181      $\vee \exists c \in Client : ReadReply(c) \vee UpdateReply(c)$ 
182      $\vee \exists p \in Partition, d \in Datacenter :$ 
183          $\vee ReadRequest(p, d)$ 
184          $\vee UpdateRequest(p, d)$ 
185          $\vee Replicate(p, d)$ 
186          $\vee Heartbeat(p, d)$ 
187          $\vee Tick(p, d)$ 

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188 $\vee \text{UpdateCSS}(p, d)$

190 $\text{Spec} \stackrel{\Delta}{=} \text{Init} \wedge \Box[\text{Next}]_{\text{vars}}$

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