

## Instantiating Fairness

The formulas  $\overline{\text{WF}_{\text{vars}_A}(\text{Consumer}_A)}$  and  $\text{WF}_{\text{vars}_A}(\overline{\text{Consumer}_A})$  need not be equivalent. This is because  $\text{WF}_{\text{vars}_A}(\text{Consumer}_A)$  is defined in terms of the enabling of action  $\langle \text{Consumer}_A \rangle_{\text{vars}_A}$ , and the following two assertions need not be equivalent

- $\overline{\langle \text{Consumer}_A \rangle_{\text{vars}_A}}$  is enabled
- $\overline{\langle \text{Consumer}_A \rangle_{\text{vars}_A}}$  (which equals  $\overline{\langle \text{Consumer}_A \rangle_{\text{vars}_A}})$  is enabled

More precisely,  $\text{WF}$  is defined in terms of the  $\text{ENABLED}$  operator, where  $\text{ENABLED } C$  is true in a state  $s$  iff action  $C$  is enabled in  $s$ . The formula  $\overline{\text{ENABLED } C}$  is true in state  $s$  iff action  $C$  is enabled in state  $\bar{s}$ . The formula  $\text{ENABLED } \bar{C}$  is true in state  $s$  iff  $\bar{C}$  is enabled in state  $s$ , where  $\bar{C}$  is the action that is true of the transition  $s \rightarrow t$  iff  $C$  is true of  $\bar{s} \rightarrow \bar{t}$ . The state predicates  $\overline{\text{ENABLED } C}$  and  $\text{ENABLED } \bar{C}$  are not necessarily equivalent. If  $\overline{\text{ENABLED } \langle \text{Consumer}_A \rangle_{\text{vars}_A}}$  is not equivalent to  $\text{ENABLED } \overline{\langle \text{Consumer}_A \rangle_{\text{vars}_A}}$ , then  $\overline{\text{WF}_{\text{vars}_A}(\text{Consumer}_A)}$  and  $\text{WF}_{\text{vars}_A}(\overline{\text{Consumer}_A})$  are not equivalent.

As an example, let  $C$  be the action  $x' \neq y'$ , where  $x$  and  $y$  are variables. Since for any state  $s$  there is a state  $t$  in which  $x \neq y$ , this action is always enabled. Hence,  $\text{ENABLED } C$  equals  $\text{TRUE}$ . Since the expression  $\text{TRUE}$  does depend on any variables,  $\overline{\text{TRUE}} = \text{TRUE}$  and hence  $\overline{\text{ENABLED } C} = \text{TRUE}$  for any refinement mapping. Now consider the refinement mapping defined by

$$\bar{x} = z \quad \bar{y} = z$$

We have

$$\begin{aligned} \bar{C} &= \overline{x' \neq y'} && \text{By definition of } C. \\ &= \bar{x'} \neq \bar{y'} && \text{By the meaning of } \bar{F} \text{ for a formula } F. \\ &= z' \neq z' && \text{By definition of } \bar{x} \text{ and } \bar{y}. \\ &= \text{FALSE} && \text{No value is unequal to itself.} \end{aligned}$$

Therefore,  $\overline{\text{ENABLED } C}$  equals  $\text{ENABLED FALSE}$ , which equals  $\text{FALSE}$ , so it is not equal to  $\overline{\text{ENABLED } C}$ , which equals  $\text{TRUE}$ .