

## OVERVIEW OF COMPILATION

7, 1					
	PROGRAMMING LANG	UACIE			
	,				
ahn	the programming lan	guage is a f	ormal language with		
	mathematical properties and well-defined meanings, as				
			ith evolved properties and		
	' '		properties and		
	ambiguities.				
			expressiveness		
	specify computations		conciseness		
			clarity:		
16n	ompiler:				
<u></u>	a specialized program which translates a source program				
	into a set of operations that are defined on the target				
	computer.		J		
	(or poter)	anna - Mara			
		compiler			
	program from	ntend optimizer	/ Program		
	\\	< infrastructur			
den	compilers that programming languages rather than the instruction				
	set of a computer are often alled source-to-source translators.				
dan	chterpreten				
	an interpreter takes as an input an executable specification				
	and produces as output the result of executing the specification				
	produces as output the result of executing				
-	Daviese		acalt		
	Pregram	c'nter proter	results		



Control of the Contro						
_	a good compiler contains a microcosm of computer science.					
	it makes practical applications of:					
	- greedy algorithms (register allocation) - heuristic search techniques (list scheduling)					
	- graph algorithms (dead code elimination),					
	- dynamic programming Cinstruction selection),					
	- finite automata & push down automata (scanning & parring)					
	- fixed-point algorithms (data-flow analysis)					
	V					
	it deals with problems, such as!					
	- dynamic allocation - synchronization					
	- naming					
	- locality					
	- memory hierarchy management					
	- pipeline scheduling					
	formal language theory has led to tools that automate					
	the production of scanners and parsers.					
	ad hoc methods - apportimate - efficiency.					
	Theretage.					

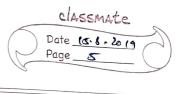
the first principle is that a compiler most observe is inviolable.

THE FUNDAMENTAL PRINCIPLES OF COMPILATION

the compiler most preserve the meaning of the program being compiled (social contract)

*	the compiler must improve the input program in some				
	discernable way.				
	COMPILER STRUCTURE 1955				
	Source , brontend IR, back end , target				
	Program Program				
	(ompiler				
	two-phase compiler				
	the compiler can make multiple passes over the 1R form of the code before emitting the target program.				
	3 0 1 0				
	two-phase structure: may simplify the process of retargeting.				
	introducing IR makes it possible to add more phases to compilation.				
	source, Grant and is optimizen is back and target,				
	program				
	1				
- (1	three-phase compiler optimizer				
	- foster target prog.				
	- Smaller tourget proof-				
	- bewer page faults				
	- less power.				

	bront end:	optimizer:	backend:		
	- sman	- ept1	- select		
	- parse	- opt 2	- Schedule		
	- CSA	- opt n	- allomte		
		$\sim$			
		analysis			
	ch frastructure;				
	- symbol tables				
	- trees				
	- graphs				
	- sets				
	- grammars				
	HIGH LEVEL VIEW OF TRANSLATION				
	TEN OF TRANSIA (ON				
	the form	/ syntax	meaning / emantics		
			J		
db	the source	language is usea	by a set, usually infinite,		
	ob strings o	defined by some	finite set of rules, called a		
	grammar.	0	partie del ef toles, carros se		
	sentence -> subject work of the				
	sentence -> subject verb object endmark				



din. the process of discovering words in a string of characters and classifying them according to their parts of speech is called scanning.

derivation in some set of grammatical rules is called parsing.

a compiler implements the abstractions defined by the

semantic analysis / context sensitive analysis.

Source language.

e.g. symbolic names.

procedures

parameters

lexical scopes

dan the process of analyzing code to discover facts from context and using that knowledge to improve the code is often called code optimization.

about the bloco of values at runtime.

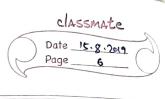
System of smultaneous egms,

don dependence analysis uses number theoretic tests to reason
about the values that can be assumed by subarript expressions

data flow analysis involves reasoning, at compile-time,

used to disambiguate reference to array elements

control-flow ops.



Varp > activation record pointer.

the problem of allocating an arbitrary set of values to a bounded set of machine reguisters in a way that minimizes loads and stores as NP-complete.

cinstruction scheduling is a hard-problem (NP-complete).

desirable properties of a compiler,

- Space - beedback (correct every syntax error)

- compile-time officiency.

-debugging

loop-invariant computation,

original FORTRAN: moltipass system

- scanner parser register allocator

eptimizations

Challenges for compilers;

- multiple functional units - long memory latencies - parallel code generation.

HETRIPICY !