

## LOCAL OPTIMIZATIONS

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- works on single basic block

~~REGIONAL TECHNIQUES~~

### REGIONAL TECHNIQUES

- consider multiple blocks, but less than whole procedure
- single loop, loop nest, dominator region

### INTRAPROCEDURAL OR GLOBAL TECHNIQUES

- operate on entire procedure
- common unit of compilation

### INTERPROCEDURAL OR WHOLE PROGRAM TECHNIQUES

- operate on  $\geq 1$  procedure, upto whole program
- logistical issues related to accessing the code (link time?)

### LINK TIME OPTIMIZATIONS

## OPTIMIZATION

### LOCAL TECHNIQUES

- dependence graph (instruction scheduling)

### REGIONAL TECHNIQUES

- control-flow graph (natural loops)
- dominator tree

## INTRAPROCEDURAL OR GLOBAL TECHNIQUES

- control flow graph
- def-use chains, sparse evaluation graphs, SSA as graph.

## INTERPROCEDURAL OR WHOLE PROGRAM TECHNIQUES

- call (multi) graph

## ANALYSIS:

reasons about the code's behaviour

## TRANSFORMATION:

rewrites the code to change its behaviour

## VALUE NUMBERING

$$a^2 \leftarrow b^0 + c^1$$

$$b^4 \leftarrow a^2 - d^3$$

$$c^5 \leftarrow b^4 + c^1$$

$$d^4 \leftarrow a^2 - d^2$$

map: var  $\rightarrow$  valuen.

map: valuen.  $\rightarrow$  var

$$"2-3" \rightarrow (4)$$

$$(4) \rightarrow "b"$$

$$"0+1" \rightarrow (2)$$

$$(2) \rightarrow "a"$$

ORIGINAL

VALUEN.

CSE

$a = x + y$

$a^3 = x^1 + y^2$

$a^3 = x^1 + y^2$

$b = x + y$

$b^3 = x^1 + y^2$

$b^3 = x^1 + y^2$

$a = 17$

$a^4 = 17^4$

$a^4 = 17^4$

$c = x + y$

$c^3 = x^1 + y^2$

$c^3 = b^3$

exp  $\rightarrow$  vnvn  $\rightarrow$  varvar  $\rightarrow$  vn?

$1+2 \rightarrow (3)$

$(3) \rightarrow a$

$\rightarrow b$

$4 \rightarrow (4)$

$(4) \rightarrow a$

$(3) \rightarrow c$

VALUE NUMBERING

$a = x + y$

$a^3 = x^1 + y^2$

$a = 17$

$a^4 = 17^4$

$c = x + y$

$c^3 = x^1 + y^2$

issue: although  $c = x + y$  is redundant, its value (VN-3) is not available in any variable.

possible solution:

introduce temporary variables.

$a^3 = x^1 + y^2$

$t^3 = a^3$

$a^4 = 17^4$

$c^3 = t^3$

## SINGLE STATIC ASSIGNMENT FORM (SSA)

$$a = x + y$$

$$a_0 = x + y$$

$$b = x + y$$

$$b_0 = x + y$$

$$a = 17$$

$$a_1 = 17$$

$$c = x + y$$

$$c_0 = x + y$$

each definition to a variable creates a new version of a variable. a new name space could be created using this approach.

how to reconcile with the rest of the namespace in other BBs?

$$a_0 = x + y$$

$$a = a_1$$

$$b_0 = x + y$$

$$b = b_0$$

$$a_1 = 17$$

$$c = c_0$$

$$c_0 = x + y$$

ORIGINAL CODE

SSA FORM

SSA WITH VN

$$a = x + y$$

$$a_0 = x + y$$

$$a_0^3 = x' + y^2$$

$$a = 17$$

$$a_1 = 17$$

$$a_1^7 = 17^7$$

$$c = x + y$$

$$c_0 = x + y$$

$$c_0^3 = x' + y^2$$

OPTIMIZED CODE

$$a_0^3 = x' + y^2$$

$$a_1^7 = 17^7$$

$$c_0^3 = a_0^3$$



$$x^2 \quad 2x \quad nx^{n-1}$$

$$d(uv) = u \, dv + v \, du$$

$$\text{SSA FORM} \quad \frac{d}{dx} x + x \cdot dx$$

$$dx + x = 2x$$

a program is in SSA form when it meets two constraints:

- each definition has a distinct name
- each use refers to a single definition

### TRANSLATING TO SSA FORM

how to translate IR into SSA form?

- easy for straight line sequence code
- each assignment to variable is given unique name
- all of the uses reached by that assignment are renamed.

$$1: a = x + y$$

$$1: a_1 = x_0 + y_0$$

$$2: a = a + 3$$

$$2: a_2 = a_1 + 3$$

$$3: b = x + y$$

$$3: b_1 = x_0 + y_0$$