Numerical errors in computer simulations

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Enter:

Example 1

A3

Start your spreadsheet program (Excel, ...)

=(A2*10)-0.9

A4-A30 =(A?*10)-0.9

Enter: Cell Value A1 0.1 A2 =(A1*10)-0.9

Cell Value
A1 2
A2-A30 =(A?*10)-18

Give this a thought!

What's happening?



General

In this course we will outline different numerical errors that may appear in computer simulations, and how these errors can affect the simulation results.

- · Errors in the mathematical model (physics)
- Errors in the entered parameters
- · Errors in the program (implementation)
- Roundoff- and truncation errors
- · Break errors

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Significant digits

A numerical result \tilde{x} is an approximation of the real value x.

Absolute error

$$\delta = |\bar{x} - x|, x \neq 0$$

Relative error

$$\frac{\delta}{\tilde{x}} = |\frac{\tilde{x} - x}{\tilde{x}}|$$

· Error margin

$$\tilde{x} - \delta < x < \tilde{x} + \delta$$

 $x = \tilde{x} + \delta$

Today's outline

- Roundoff and truncation errors

Significant digits

• x has m significant digits if the absolute error in x is smaller or equal to 5 at the (m + 1)-th position:

$$10^{q-1} \le |\tilde{x}| \le 10^q$$

 $|x - \tilde{x}| = 0.5 \times 10^{q-m}$

· For example:

$$x = \frac{1}{3}, \tilde{x} = 0.333 \Rightarrow \delta = 0.00033333...$$

3 significant digits

Representation of numbers

- . Computers represent a number with a finite number of digits: each number is therefore an approximation due to roundoff and truncation errors
- In the decimal system, a digit c at position n has a value of $c \times 10^{n-1}$

$$(4521)_{10} = 4 \times 10^3 + 5 \times 10^2 + 2 \times 10^1 + 1 \times 10^0$$

Representation of numbers

You could use another basis computers often use the basis 2.

$$\begin{aligned} (4521)_{10} &= 1 \times 2^{12} + 0 \times 2^{11} + 0 \times 2^{10} + 0 \times 2^9 + 1 \times 2^8 + \dots \\ &\dots 1 \times 2^7 + 0 \times 2^9 + 1 \times 2^5 + 0 \times 2^4 + \dots \\ &\dots 1 \times 2^3 + 0 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 \\ &= (1000110101001)_2 \end{aligned}$$

· In general:

$$(c_m \dots c_1 c_0)_q = c_0 q^0 + c_1 q^1 + \dots + c_m q^m, c \in \{0, 1, 2, \dots, q-1\}$$

Excercise

. Convert the following decimal number to base-2: 214

- Excel:
- Decimal: =DEC2BIN(214)
 - Octal: =DEC20CT(214)
 - Hexadecimal: =DEC2HEX (214)
- Matlab:
 - Decimal: dec2bin(214)
 - · Other base: dec2base(214, <base>)

Representation of numbers

- . Numbers are stored in binary in the memory of a computer, in segments of a specific length (called a word).
- · We distinguish multiple types of numbers:
 - Integers: -301, -1, 0, 1, 96, 2293, . . .
 - Floating points: -301.01.0.01.3.14159265.14498.2
- · A binary integer representation looks like the following bit sequence:

$$z = \sigma \left(c_0 2^0 + c_1 2^1 + \ldots + c_{\lambda-1} 2^{\lambda-1} \right)$$

 σ is the sign of z (+ or -), and λ is the length of the word

. Endianness: the order of bits stored by a computer

Arithmetic operations with binary numbers

Addition

0 + 0 = 0													
0 + 1 = 1		1	4	5		1	0	0	1	0	0	0	1
1 + 0 = 1	+		2	3	+	0	0	0	1	0	1	1	1
1 + 1 = 0		1	6	8		1	0	1	0	1	0	0	0
(carry one)				_									

Subtraction

 Multiplication and division are more expensive, and more elaborate

Roundoff and truncation errors Break errors Loss of digits (Un)stable methods Symbolic math Summ

Excercise

Try the following commands in Matlab:

Command	Result
intmin	-2147483648
intmax	2147483647
i = int16(intmax)	i = 32767
whos i	int16 information
i = i + 100	i = 32767
realmax	1.7977e+308
f = 0.1	
whos f	double information
format long e	
realmax	1.797693134862316e+308
f	
fprintf("%0.16f",f)	0.1000000000000000
fprintf("%0.20f",f)	0.10000000000000000555

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Representation of real (floating point) numbers

 Formally, a real number is represented by the following bit sequence

$$x = \sigma \left(2^{-1} + c_2 2^{-2} + \ldots + c_m 2^{-m}\right) 2^{e-1023}$$

Here, σ is the sign of x and e is an integer value.

 A floating point number hence contains sections that contain the sign, the exponent and the mantissa

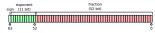


Image: Wikimedia Commons CC by-SA

Roundoff and trun

Representation of integer numbers

- In Matlab, integers of the type int32 are represented by 32-bit words (λ = 31).
- The set of numbers that an int32 z can represent is:

$$-2^{31} \le z \le 2^{31} - 1 \approx 2 \times 10^9$$

- If, during a calculation, an integer number becomes larger than 2^λ – 1, the computer reports an overflow¹
- · How can a computer identify an overflow?

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Representation of real (floating point) numbers

• Example:
$$\lambda = 3$$
, $m = 2$, $x = \frac{2}{3}$
 $x = \pm \left(2^{-1} + c_2 2^{-2}\right) 2^e$

- Truncation: $fl(x) = 2^{-1} = 0.5$
- Round off: $fl(x) = 2^{-1} + 2^{-2} = 0.75$

¹Matlab does not perform actual integer overflows, it just stops at the maximum.

- Break errors

Trigonometric, Logarithmic, and Exponential computations

- . These operations involve many multiplications and additions. and are therefore expensive
- . Computations can only take finite time, for infinite series, calculations are interrupted at N

$$\begin{split} \sin(x) &= \sum_{n=0}^{N} \frac{(-1)^n}{(2n+1)!} x^{2n+1} = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \dots + \frac{(-1)^N}{(2N+1)!} x^{2N+1} \\ e^x &= \sum_{n=1}^{N} \frac{x^n}{n!} = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \frac{x^4}{4!} + \dots + \frac{x^N}{N!} \end{split}$$

. This results in a break error

Trigonometric, Logarithmic, and Exponential computations

- · Processors can do logic and arithmetic instructions
- . Trigonometric, logarithmic and exponential calculations are "higher-level" functions:
 - exp, sin, cos, tan, sec, arcsin, arccos, arctan, log, ln, ...
- . Such functions can be performed using these "low level" instructions, for instance using a Taylor series:

$$\sin(x) = \sum_{n=0}^{\infty} \frac{(-1)^n}{(2n+1)!} x^{2n+1} = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots$$

$$e^x = \sum_{n=0}^{\infty} \frac{x^n}{x!} = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \frac{x^4}{4!} + \dots$$

Algorithm for sine-computation

A computer may use a clever algorithm to limit the number of operations required to perform a higher-level function. A (fictional!) example for the computation of sin(x):

- (1) Use periodicity so that $0 < x < 2\pi$
- ② Use symmetry $(0 < x < \frac{\pi}{4})$
- Use lookup tables for known values
- O Perform taylor expansion



- - O Loss of digits

- During operations such as +. −. ×. ÷. an error can add up
- · Consider the summation of x and v

$$\tilde{\mathbf{x}} - \delta \leq \mathbf{x} \leq \tilde{\mathbf{x}} + \delta \quad \text{and} \quad \tilde{\mathbf{y}} - \epsilon \leq \mathbf{y} \leq \tilde{\mathbf{y}} + \epsilon$$

$$(\tilde{x} + \tilde{y}) - (\delta + \varepsilon) \le x + y \le (\tilde{x} + \tilde{y}) + (\delta + \varepsilon)$$

Loss of digits (Un)stable 0000000 0000000

Loss of digits: Example 1

$$\begin{array}{l} x = \pi, \tilde{x} = 3.1416 \\ y = 22/7, \tilde{y} = 3.1429 \end{array} \Rightarrow \begin{array}{l} \delta = \tilde{x} - x = 7.35 \times 10^{-6} \\ \varepsilon = \tilde{y} - y = 4.29 \times 10^{-5} \end{array}$$

$$x + y = \tilde{x} + \tilde{y} \pm (\delta + \varepsilon) \approx 6.2845 - 5.025 \times 10^{-5}$$

 $x - y = \tilde{x} - \tilde{y} \pm (\delta + \varepsilon) \approx -0.0013 + 3.55 \times 10^{-5}$

- The absolute error is small ($\approx 10^{-5}$), but the relative error is much bigger (0.028).
- · Adding up the errors results in a loss of significant digits!

Loss of digits (Un)stable 0000000 0000000

Loss of digits: Example 2

- Calculate e^{−5}
 - · Use the Taylor series
 - Calculate the first 26 terms (N = 26)
- . Now repeat the calculation, but use for each calculation only 4 digits. What do you find? Use: str2double(sprintf('%,4g', term))
- Without errors you would find: e⁻⁵ = 0.006738
- . If you only use 4 digits in the calculations, you'll find 0.00998

Badly (ill) conditioned problems

We consider a system F(x, y) that computes a solution from input data. The input data may have errors:

$$\begin{array}{c|c} x & \hline \\ F(x,y) & = 0 \\ \hline \\ x + \delta x & \hline \\ F(x + \delta x, y + \delta y) & = 0 \\ \hline \end{array} \begin{array}{c} y + \delta y \\ \hline \\ \end{array} \begin{array}{c} (\text{With out errors}) \\ \end{array}$$

$$y(x + \delta x) - y(x) \approx y'(x)\delta x$$
 $C = \max_{\delta x} \left(\left| \frac{\delta y/y}{\delta x/x} \right| \right)$

Propagated error on the basis of Taylor expansion Condition criterion, C < 10 error development small

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k errors. Loss of digits. (Un)stable methods. Symbolic math. Summary

Badly (ill) conditioned problems: Example

- Matlab already warned us about the bad condition number: Warning: Matrix is close to singular or badly scaled.
 Results may be inaccurate. RCDND = 1.148983e-08.
- The RCOND is the reciprocal condition number
- A small error in x results in a big error in y. This is called an ill conditioned problem.

Badly (ill) conditioned problems: Example

Solve the following linear system in Matlab using double and single precision 1.2969 0.8648 0.2161 0.1441 Double precision Single precision >> clear; clc; format long e; >> clear; clc; format long e; >> A = [[1.2969 0.8648]: >> A = single([0.2161 0.1441]]: [[1.2969 0.8648]: >> x = [0.8642; 0.1440]; [0.2161 0.1441]]); >> y = A\x >> x = single([0.8642; 0.1440]); 2.000000002400302e+00 >> v = A\x -2.000000003599621e+00 у =

1.3331791e+00

-1.0000000e+00

Today's outline

- Introduction
- Roundoff and truncation errors
- Break e
- Loss of digit
- (Un)stable methods
- Symbolic ma
- Summan

Break errors Loss of digits (Un)stable methods Symbolic math Summary 0000 0000000 00000000 00000000 00

(Un)stable methods

- The condition criterion does not tell you anything about the quality of a numerical solution method!
- It is very well possible that a certain solution method is more sensitive for one problem than another
- If the method propagates the error, we call it an unstable method. Let's look at an example.

The Golden mean



```
Powerlaw version

% initialise
x = (1 + sqrt(5))/2;
y2(1) - x^0; % n = 1

% Perform powerlaw apprach
for n = 0:39
y2(n+1) = x^-n
end
```

The Golden mean

· Let's evaluate the following recurrent relationship:

Break errors Loss of digits (Un)stable m 0000 0000000 0000000

$$y_{n+1} = y_{n-1} - y_n$$

 $y_0 = 1, \quad y_1 = \frac{2}{1 + \sqrt{5}}$

. You can prove (by substitution) that:

$$y_n = x^{-n}$$
, $n = 0, 1, 2, ...$, $x = \frac{1 + \sqrt{5}}{2}$

The Golden mean

n	Recurrent	Powerlaw					
1	1.0000	1.0000					
1	0.6180	0.6180					
2	0.3820	0.3820					
3	0.2361	0.2361					
37	$3.080 \cdot 10^{-08}$	$2.995 \cdot 10^{-08}$					
38	$1.714 \cdot 10^{-08}$	$1.851 \cdot 10^{-08}$					
39	$1.366 \cdot 10^{-08}$	$1.144 \cdot 10^{-08}$					
40	$3.485 \cdot 10^{-08}$	$7.071 \cdot 10^{-08}$					

The recurrent approach enlarges errors from earlier calculations!

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Example 1: Explanation

Recall example 1, where the errors blew up our computation of 0.1, whereas they did not for 2. Why did we see these results?

Break errors Loss of digits (Un)stable 0000 0000000 0000000

- The number 0.1 is not exactly represented in binary
 A tiny error can accumulate up to catastrophic proportions!
- The number 2 does have an exact binary representation

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- (Un)stable methods
- Symbolic math
- Summary

Example 2

Start your calculation program of choice (Excel. Matlab. ...)

Calculate the result of v:

 $v = e^{\pi} - \pi = 19.999099979 \neq 20$





Image: xkcd

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Symbolic math packages

. .

Definition

The use of computers to manipulate mathematical equations and expressions in symbolic form, as opposed to manipulating the numerical quantities represented by those symbols.

- Symbolic integration or differentiation, substitution of one expression into another
- · Simplification of an expression, change of subject etc.
- · Packages and toolboxes:

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Symbolic math packages

via TU/e

Mathematica Well known software package, license available

Maple Well known, license available via TU/e

Wolfram Alpha Web-based interface by Mathematica developer. Less powerful in mathematical respect, but more accessible and has a broad application range (unit conversion, semantic commands).

> Sage Open-source alternative to Maple, Mathematica, Magma, and MATLAB.

Matlab Symbolic math toolbox

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Symbolic math: integration and differentiation

$$f(x) = \frac{1}{x^3 + 1}$$
>> syss x
>> f = 1/(x^3 + 1);
>> sys_t int = int(f)
sys_int = log(x + 1)/3 - log((x - 1/2)^2 + 3/4)/6 + (3^2 (1/2) + stan (2 * 3^2 (1/2) + (x - 1/2))/3)/3

>> my_f_diff = diff(my_f_int)

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Symbolic math: simplify

$$\begin{split} f(\mathbf{x}) &= (\mathbf{x} - 1)(\mathbf{x} + 1)(\mathbf{x}^2 + 1) + 1 \\ \\ >> & \text{ syms } \mathbf{x} \\ >> & f - (\mathbf{x} - 1) \cdot (\mathbf{x} + 1) \cdot (\mathbf{x}^{-2} + 1) + 1 \\ f - \\ & (\mathbf{x}^2 + 1) \cdot (\mathbf{x} - 1) \cdot (\mathbf{x} + 1) + 1 \\ \\ >> & f - \sin plify(f) \\ f 2 - \end{aligned}$$

Symbolic math: exercises

x^4

Simplify the following expression:

$$f(x) = \frac{2\tan x}{(1 + \tan^2 x)} = \sin 2x$$
>> simplify(2*tan(x)/(1 + tan(x)^2))

Exercise 2

Calculate the value of p:

$$\begin{split} \rho &= \int_0^{10} \frac{e^x - e^{-x}}{\sinh x} dx \\ >> f = \cdot ((\exp(x) - \exp(-x)) / \sinh(x)); \\ >> p = \cdot \inf(f, 0, 10) \\ p = 20 \end{split}$$

Symbolic math function

$$f(x) = \frac{3}{x^2 + 3x} - 2$$
>> syms x
>> f = 3 / (x^2 + 3*x) - 2;
>> solve(f)
ans =
15^*(1/2)/2 - 3/2
- 15^*(1/2)/2 - 3/2

Function as a string

```
f(x) = x^2 - 4x + 2
>> solve('x^2 - 4*x + 2')
ans =
 2^(1/2) + 2
 2 - 2^{(1/2)}
```

Today's outline

- Summary

Symbolic math toolbox: variable precision arithmetic

Variable precision can be used to specify the number of significant digits.

```
p = 0.33333333333333333
>> p = vpa(1/3,4)
p = 0.3333
>> a = vpa(0.1. 30)
a = 0.1
>> b = vpa(0.1, 5);
b = 0.1
33 a=b
```

>> p = vpa(1/3,16)

ans = 0.000000000000056843418860808014869689938467514

Summary

- · Numerical errors mar arise due to truncation, roundoff and break errors, which may seriously affect the accuracy of your solution
- . Errors may propagate and accumulate, leading to smaller accuracy
- . Ill-conditioned problems and unstable methods have to be identified so that proper measures can be taken
- Symbolic math computations may be performed to solve certain equations algebraically, bypassing numerical errors, but this is not always possible.