Numerical errors in computer simulations

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Today's outline

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Enter:

Cell	Value
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Give this a thought!

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Significant digits

A numerical result $\tilde{x}$ is an approximation of the real value x.

- Absolute error

  $\delta = \tilde{x} - x, x \neq 0$

- Relative error

  $\frac{\delta}{\tilde{x}} = \frac{\tilde{x} - x}{\tilde{x}}$

- Error margin

  $\tilde{x} - \delta \leq x \leq \tilde{x} + \delta$

  $x = \tilde{x} \pm \delta$

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Significant digits

- $\tilde{x}$ has m significant digits if the absolute error in x is smaller or equal to 5 at the $(m+1)$-th position:

  $$10^{q-1} \leq |\tilde{x}| \leq 10^{q}$$

  $$|x - \tilde{x}| = 0.5 \times 10^{q-m}$$

- For example:

  $$x = \frac{1}{3}, \tilde{x} = 0.333 \Rightarrow \delta = 0.00033333\ldots$$

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Number of 1's: $1 \times 10^{(0)}$

Number of 10's: $2 \times 10^{(1)}$

Number of 100's: $5 \times 10^{(2)}$

Number of 1000's: $4 \times 10^{(3)}$

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$$(4521)_{10} = 4 \times 10^3 + 5 \times 10^2 + 2 \times 10^1 + 1 \times 10^0$$

Representation of numbers

- You could use another basis, computers often use the basis 2:

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- In general:

$$(c_m \ldots c_1 c_0)_q = c_0 q^0 + c_1 q^1 + \ldots + c_m q^m, c \in \{0, 1, 2, \ldots, q-1\}$$

Representation of numbers

- Numbers are stored in binary in the memory of a computer, in segments of a specific length (called a *word*).
- We distinguish multiple types of numbers:
  - Integers: $-301, -1, 0, 1, 96, 2293, \ldots$
  - Floating points: $-301.01, 0.01, 3.14159265, 14498.2$
- A binary integer representation looks like the following bit sequence:

  $$z = \sigma \left(c_0 2^0 + c_1 2^1 + \ldots + c_{\lambda-1} 2^{\lambda-1} \right)$$

  σ is the sign of z (+ or -), and λ is the length of the word
- Endianness: the order of bits stored by a computer

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- Excel:
  - Decimal: `=DEC2BIN(214)`
  - Octal: `=DEC2OCT(214)`
  - Hexadecimal: `=DEC2HEX(214)`
- Matlab:
  - Decimal: `dec2bin(214)`
  - Other base: `dec2base(214,<base>)`

Arithmetic operations with binary numbers

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$0 + 0 = 0$

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- Multiplication and division are more expensive, and more elaborate

Excercise

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Command	Result
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`format long e`	

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`f`	
`fprintf("%0.16f",f)`	0.1000000000000000

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`f`	
`fprintf("%0.16f",f)`	0.1000000000000000
`fprintf("%0.20f",f)`	0.10000000000000000555

Representation of integer numbers

- In Matlab, integers of the type `int32` are represented by 32-bit words ($\lambda = 31$).

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- The set of numbers that an `int32` z can represent is:

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- If, during a calculation, an integer number becomes larger than $2^\lambda - 1$, the computer reports an overflow[1]

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- The set of numbers that an `int32` z can represent is:

$$-2^{31} \leq z \leq 2^{31} - 1 \approx 2 \times 10^{9}$$

- If, during a calculation, an integer number becomes larger than $2^{\lambda} - 1$, the computer reports an overflow[1]
- How can a computer identify an overflow?

[1] Matlab does not perform actual integer overflows, it just stops at the maximum.

Representation of real (floating point) numbers

- Formally, a real number is represented by the following bit sequence

$$x = \sigma \left(2^{-1} + c_2 2^{-2} + \ldots + c_m 2^{-m}\right) 2^e$$

Here, σ is the sign of x and e is an integer value.

- A floating point number hence contains sections that contain the sign, the exponent and the mantissa

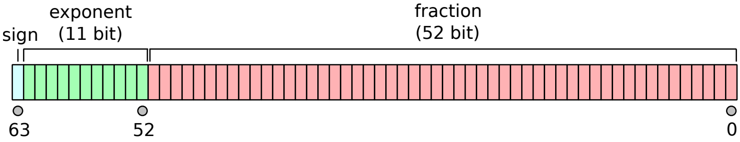

Image: Wikimedia Commons CC by-SA

Representation of real (floating point) numbers

- Example: $\lambda = 3,\ m = 2,\ x = \frac{2}{3}$

  $$x = \pm\left(2^{-1} + c_2 2^{-2}\right) 2^e$$

  - $c_0 \in \{0, 1\}$
  - $e = \pm a_0 2^0$
  - $a_0 \in \{0, 1\}$
- Truncation: $fl(x) = 2^{-1} = 0.5$
- Round off: $fl(x) = 2^{-1} + 2^{-2} = 0.75$

Today's outline

Trigonometric, Logarithmic, and Exponential computations

- Processors can do logic and arithmetic instructions
- Trigonometric, logarithmic and exponential calculations are "higher-level" functions:
  exp, sin, cos, tan, sec, arcsin, arccos, arctan, log, ln, ...
- Such functions can be performed using these "low level" instructions, for instance using a Taylor series:

$$\sin(x) = \sum_{n=0}^{\infty} \frac{(-1)^n}{(2n+1)!} x^{2n+1} = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \ldots$$

$$e^x = \sum_{n=0}^{\infty} \frac{x^n}{x!} = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \frac{x^4}{4!} + \ldots$$

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- These operations involve many multiplications and additions, and are therefore *expensive*
- Computations can only take finite time, for infinite series, calculations are interrupted at N

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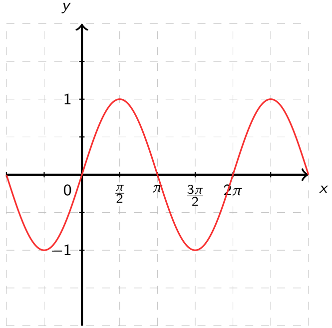
Algorithm for sine-computation

A computer may use a clever algorithm to limit the number of operations required to perform a higher-level function. A (fictional!) example for the computation of $\sin(x)$:



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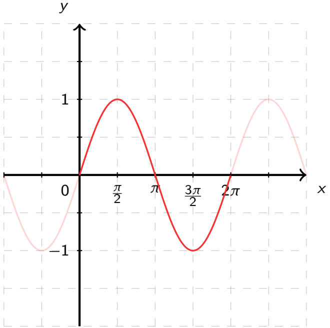

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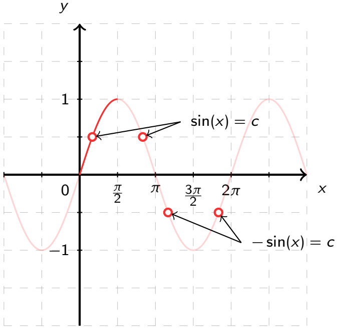

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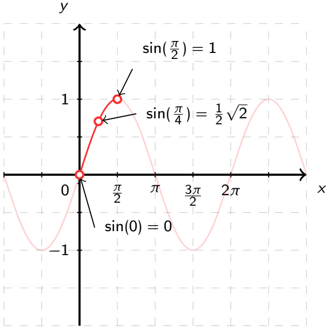

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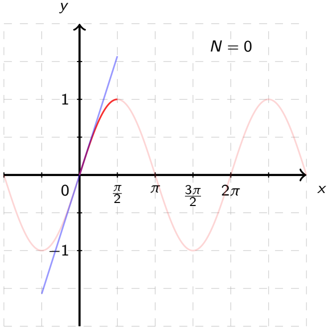

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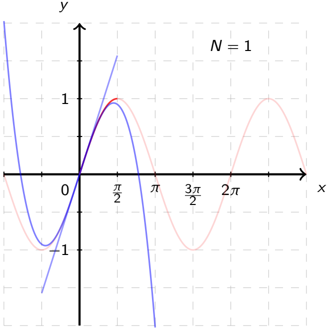

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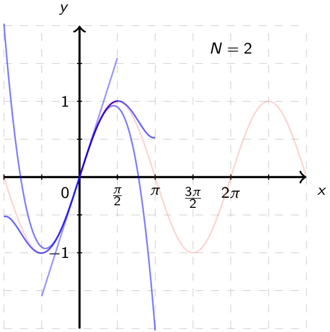

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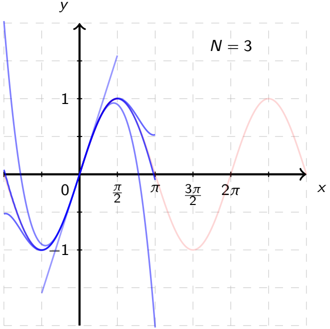

Today's outline

Loss of digits

- During operations such as $+$, $-$, $\times$, $\div$, an error can add up
- Consider the summation of x and y

$$\tilde{x} - \delta \leq x \leq \tilde{x} + \delta \quad \text{and} \quad \tilde{y} - \varepsilon \leq y \leq \tilde{y} + \varepsilon$$

$$(\tilde{x} + \tilde{y}) - (\delta + \varepsilon) \leq x + y \leq (\tilde{x} + \tilde{y}) + (\delta + \varepsilon)$$

Loss of digits: Example 1

$x = \pi, \tilde{x} = 3.1416$

$y = 22/7, \tilde{y} = 3.1429$

Loss of digits: Example 1

$$\left.\begin{array}{l} x = \pi, \tilde{x} = 3.1416 \\ y = 22/7, \tilde{y} = 3.1429 \end{array}\right\} \Rightarrow \left.\begin{array}{l} \delta = \tilde{x} - x = 7.35 \times 10^{-6} \\ \varepsilon = \tilde{y} - y = 4.29 \times 10^{-5} \end{array}\right\}$$

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$$x + y = \tilde{x} + \tilde{y} \pm (\delta + \varepsilon) \approx 6.2845 - 5.025 \times 10^{-5}$$

$$x - y = \tilde{x} - \tilde{y} \pm (\delta + \varepsilon) \approx -0.0013 + 3.55 \times 10^{-5}$$

- The absolute error is small ($\approx 10^{-6}$), but the relative error is much bigger (0.028).
- Adding up the errors results in a loss of significant digits!

Loss of digits: Example 2

- Calculate e^{-5}
  - Use the Taylor series
  - Calculate the first 26 terms ($N = 26$)
- Now repeat the calculation, but use for each calculation only 4 digits. What do you find?

- Without errors you would find: $e^{-5} = 0.006738$
- If you only use 4 digits in the calculations, you'll find 0.00998

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  Use: `str2double(sprintf('%.4g', term))`
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Badly (ill) conditioned problems

We consider a system $F(x, y)$ that computes a solution from input data. The input data may have errors:

$$\xrightarrow{\ x\ } \boxed{\quad F(x, y) \ = \ 0 \quad} \xrightarrow{\ y\ } \text{(Without errors)}$$

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$$\xrightarrow{x} \boxed{F(x, y) = 0} \xrightarrow{y} \text{(Without errors)}$$

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$$\xrightarrow{x+\delta x} \boxed{F(x+\delta x, y+\delta y) = 0} \xrightarrow{y+\delta y} \text{(With errors)}$$

$$y(x+\delta x) - y(x) \approx y'(x)\delta x$$

Propagated error on the basis of Taylor expansion

$$C = \max_{\delta x}\left(\left|\frac{\delta y/y}{\delta x/x}\right|\right)$$

Condition criterion, $C < 10$ error development small

Badly (ill) conditioned problems: Example

Solve the following linear system in Matlab using double and single precision:

$$A = \begin{bmatrix} 1.2969 & 0.8648 \\ 0.2161 & 0.1441 \end{bmatrix}, \quad x = \begin{bmatrix} 0.8642 \\ 0.1440 \end{bmatrix}$$

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Double precision

```
>> clear;clc;format long e;
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    [0.2161 0.1441]];
>> x = [0.8642; 0.1440];
>> y = A\x
y =
      2.000000002400302e+00
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```

Single precision

```
>> clear;clc;format long e;
>> A = single(
  [[1.2969 0.8648];
  [0.2161 0.1441]] );
>> x = single(
  [0.8642; 0.1440] );
>> y = A\x
y =
    1.3331791e+00
   -1.0000000e+00
```

Badly (ill) conditioned problems: Example

- Matlab already warned us about the bad condition number:

```
Warning: Matrix is close to singular or badly scaled.
Results may be inaccurate. RCOND = 1.148983e-08.
```

- The `RCOND` is the reciprocal condition number
- A small error in x results in a big error in y. This is called an ill conditioned problem.

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Today's outline

(Un)stable methods

- The condition criterion does not tell you anything about the quality of a numerical solution method!
- It is very well possible that a certain solution method is more sensitive for one problem than another
- If the method propagates the error, we call it an *unstable method*. Let's look at an example.

The Golden mean

- Let's evaluate the following recurrent relationship:

$$y_{n+1} = y_{n-1} - y_n$$

$$y_0 = 1, \quad y_1 = \frac{2}{1+\sqrt{5}}$$

- You can prove (by substitution) that:

$$y_n = x^{-n}, \quad n = 0, 1, 2, \ldots, \quad x = \frac{1+\sqrt{5}}{2}$$

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The Golden mean

Recurrent version

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% initialise
y(1) = 1;
y(2) = 2 / (1 + sqrt(5));

% Perform recurrent
    approach
for n = 2:39
    y(n+1) = y(n-1)-y(n);
end
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The Golden mean

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% Perform recurrent
    approach
for n = 2:39
    y(n+1) = y(n-1)-y(n);
end
```

Powerlaw version

```
% initialise
x = (1 + sqrt(5))/2;
y2(1) = x^0; % n = 1

% Perform powerlaw apprach
for n = 0:39
    y2(n+1) = x^-n
end
```

The Golden mean

n	Recurrent	Powerlaw
1	1.0000	1.0000
1	0.6180	0.6180
2	0.3820	0.3820
3	0.2361	0.2361
...	...	...
37	$3.080 \cdot 10^{-08}$	$2.995 \cdot 10^{-08}$
38	$1.714 \cdot 10^{-08}$	$1.851 \cdot 10^{-08}$
39	$1.366 \cdot 10^{-08}$	$1.144 \cdot 10^{-08}$
40	$3.485 \cdot 10^{-08}$	$7.071 \cdot 10^{-08}$

- The recurrent approach enlarges errors from earlier calculations!

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- The number 0.1 is not exactly represented in binary
    - A tiny error can accumulate up to catastrophic proportions!
- The number 2 does have an exact binary representation

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Start your calculation program of choice (Excel, Matlab, ...)

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Calculate the result of y:

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Calculate the result of y:

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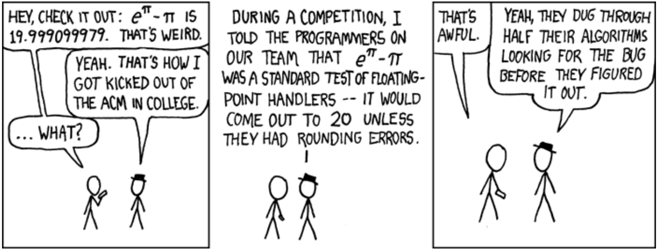

Image: xkcd

Today's outline

Symbolic math packages

Definition

The use of computers to manipulate mathematical equations and expressions in symbolic form, as opposed to manipulating the numerical quantities represented by those symbols.

- Symbolic integration or differentiation, substitution of one expression into another
- Simplification of an expression, change of subject etc.
- Packages and toolboxes:

Symbolic math packages

Mathematica	Well known software package, license available via TU/e
Maple	Well known, license available via TU/e
Wolfram|Alpha	Web-based interface by Mathematica developer. Less powerful in mathematical respect, but more accessible and has a broad application range (unit conversion, semantic commands).
Sage	Open-source alternative to Maple, Mathematica, Magma, and MATLAB.
Matlab	Symbolic math toolbox

Symbolic math: simplify

$$f(x) = (x-1)(x+1)(x^2+1)+1$$

Symbolic math: simplify

$$f(x) = (x-1)(x+1)(x^2+1)+1$$

```
>> syms x
>> f = (x - 1)*(x + 1)*(x^2 + 1) + 1

f =

(x^2 + 1)*(x - 1)*(x + 1) + 1

>> f2 = simplify(f)

f2 =

x^4
```

Symbolic math: integration and differentiation

$$f(x) = \frac{1}{x^3 + 1}$$

Symbolic math: integration and differentiation

$$f(x) = \frac{1}{x^3 + 1}$$

```
>> syms x
>> f = 1/(x^3+1);
>> my_f_int = int(f)

my_f_int = log(x + 1)/3 - log((x - 1/2)^2 + 3/4)/6 +
    (3^(1/2)*atan((2*3^(1/2)*(x - 1/2))/3))/3

>> my_f_diff = diff(my_f_int)

my_f_diff = 1/(3*(x + 1)) + 2/(3*((4*(x - 1/2)^2)/3 +
    1)) - (2*x - 1)/(6*((x - 1/2)^2 + 3/4))

>> simplify(my_f_diff)

ans = 1/(x^3 + 1)
```

Symbolic math: exercises

Exercise 1

Simplify the following expression:

$$f(x) = \frac{2\tan x}{(1 + \tan^2 x)}$$

Symbolic math: exercises

Exercise 1

Simplify the following expression:

$$f(x) = \frac{2\tan x}{(1+\tan^2 x)} = \sin 2x$$

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>> simplify(2*tan(x)/(1 + tan(x)^2))
```

Symbolic math: exercises

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>> simplify(2*tan(x)/(1 + tan(x)^2))
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Exercise 2

Calculate the *value* of p:

$$p = \int_0^{10} \frac{e^x - e^{-x}}{\sinh x}$$

Symbolic math: exercises

Exercise 1

Simplify the following expression:

$$f(x) = \frac{2\tan x}{(1 + \tan^2 x)} = \sin 2x$$

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>> simplify(2*tan(x)/(1 + tan(x)^2))
```

Exercise 2

Calculate the *value* of p:

$$p = \int_0^{10} \frac{e^x - e^{-x}}{\sinh x}$$

```
>> f = ((exp(x)- exp(-x))/sinh(x));
>> p = int(f,0,10)
p = 20
```

Symbolic math: root finding

A root finding method searches for the values where a function reaches zero. We will cover the numerical methods later, here we show how to use root finding with symbolic math in Matlab.

Symbolic math function

$$f(x) = \frac{3}{x^2 + 3x} - 2$$

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```
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Function as a string

$$f(x) = x^2 - 4x + 2$$

```
>> solve('x^2 - 4*x + 2')
ans =
 2^(1/2) + 2
 2 - 2^(1/2)
```

Symbolic math toolbox: variable precision arithmetic

Variable precision can be used to specify the number of significant digits.

```
>> p = vpa(1/3,16)
p = 0.3333333333333333
>> p = vpa(1/3,4)
p = 0.3333
>> a = vpa(0.1, 30)
a = 0.1
>> b = vpa(0.1, 5);
b = 0.1
>> a-b

ans = 0.00000000000000568434188608080148696899384675140
```

Today's outline

Summary

- Numerical errors mar arise due to truncation, roundoff and break errors, which may seriously affect the accuracy of your solution
- Errors may propagate and accumulate, leading to smaller accuracy
- Ill-conditioned problems and unstable methods have to be identified so that proper measures can be taken
- Symbolic math computations may be performed to solve certain equations algebraically, bypassing numerical errors, but this is not always possible.

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