Errors in computer simulations

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Today's outline

- Introduction
- Roundoff and truncation errors
- 3 Break errors
- 4 Loss of digits
- **5** (Un)stable methods
- 6 Symbolic math
- Summary

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Start your spreadsheet program (Excel, ...)

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Enter:

Cell	Value
A1	0.1

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Enter:

Cell	Value
A1	0.1
A2	=(A1*10)-0.9

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Cell	Value
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A2	=(A1*10)-0.9
A3	=(A2*10)-0.9
A4	=(A3*10)-0.9

(repeat until A30)

Start your spreadsheet program (Excel, ...)

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A1	0.1
A2	=(A1*10)-0.9
A3	=(A2*10)-0.9
A4	=(A3*10)-0.9

(repeat until A30)

What's happening?

Start your spreadsheet program (Excel, ...)

Enter:

Cell	Value
A1	0.1
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A3	=(A2*10)-0.9
Α4	=(A3*10)-0.9

Enter:

Cell	Value
A1	2

(repeat until A30)

What's happening?

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Enter:

Cell	Value
A1	0.1
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А3	=(A2*10)-0.9
A4	=(A3*10)-0.9

Enter:

Cell	Value
A1	2
A2	=(A1*10)-18
A3	=(A2*10)-18
A4	=(A3*10)-18

(repeat until A30)

(repeat until A30)

What's happening?

Errors in computer simulations

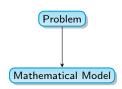
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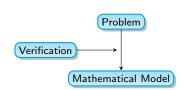
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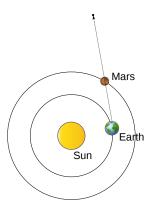
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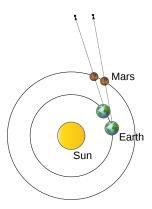
Problem

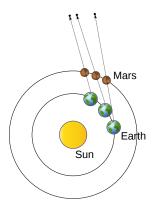


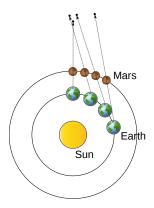
Verification

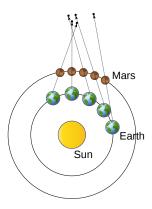


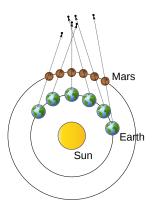


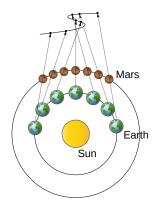


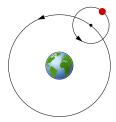


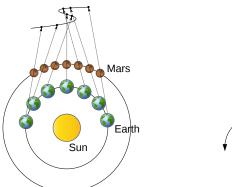


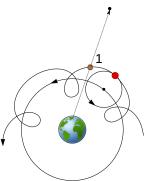


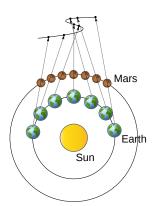


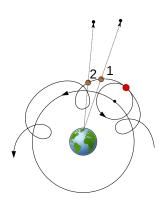


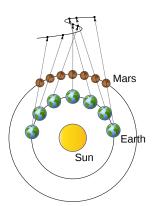


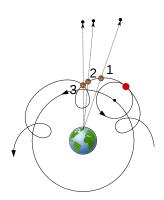


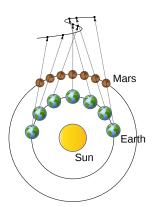


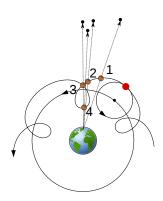


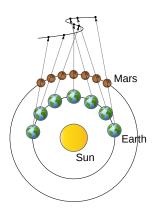


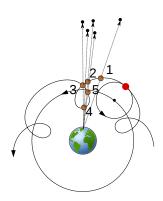


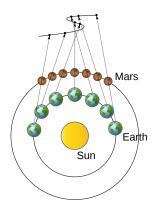


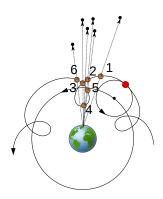


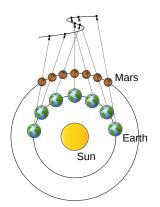


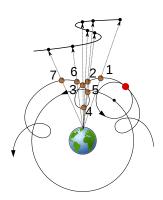












- The perceived orbit of Mars from Earth shows a zig-zag (in contrast to the Sun, Mercury, Venus)
- Even though they were not 'right', Earth-centered models (Ptolemy) were still valid

Be aware of your uncertainties

Aleatory uncertainty

Uncertainty that arises due to inherent randomness of the system, features that are too complex to measure and take into account

Epistemic uncertainty

Uncertainty that arises due to lack of knowledge of the system, but could in principle be known

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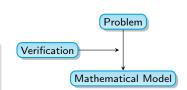
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Errors in computer simulations

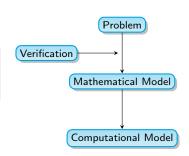
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- Errors in the mathematical model (physics)
- Errors in the program (implementation)

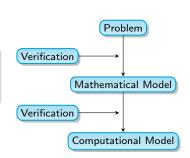
Verification



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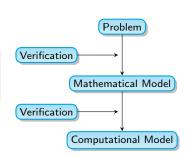
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- Errors in the entered parameters

Verification and validation

Verification

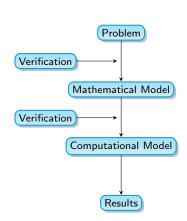
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Verification and validation

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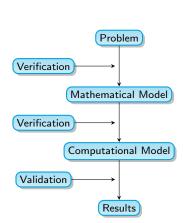
Verification and validation

Verification

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Validation

Validation is the process of determining the degree to which a model is an accurate representation of the real world from the perspective of the intended uses of the model



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A numerical result \tilde{x} is an approximation of the real value x.

Absolute error

$$\delta = |\tilde{x} - x|, x \neq 0$$

Relative error

$$\frac{\delta}{\tilde{x}} = |\frac{\tilde{x} - x}{\tilde{x}}|$$

• Error margin

$$\tilde{x} - \delta \le x \le \tilde{x} + \delta$$

$$x = \tilde{x} \pm \delta$$

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• \tilde{x} has m significant digits if the absolute error in x is smaller or equal to 5 at the (m+1)-th position:

$$10^{q-1} \le |\tilde{x}| \le 10^q$$

 $|x - \tilde{x}| = 0.5 \times 10^{q-m}$

For example:

$$x = \frac{1}{3}, \tilde{x} = 0.333 \Rightarrow \delta = 0.000333333...$$

3 significant digits

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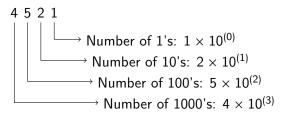
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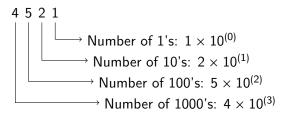
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$$(4521)_{10} = 4 \times 10^3 + 5 \times 10^2 + 2 \times 10^1 + 1 \times 10^0$$

$$(4521)_{10} =$$

$$(4521)_{10} = 1 \times 2^{12} +$$

$$=(1$$

$$(4521)_{10} = 1 \times 2^{12} + 0 \times 2^{11} +$$

$$=(10)$$

$$(4521)_{10} = 1 \times 2^{12} + 0 \times 2^{11} + 0 \times 2^{10} +$$

$$=(100)$$

$$(4521)_{10} = 1 \times 2^{12} + 0 \times 2^{11} + 0 \times 2^{10} + 0 \times 2^9 + 0$$

$$=(1000$$

$$(4521)_{10} = 1 \times 2^{12} + 0 \times 2^{11} + 0 \times 2^{10} + 0 \times 2^9 + 1 \times 2^8 + \dots$$

$$=(10001$$

$$(4521)_{10} = 1 \times 2^{12} + 0 \times 2^{11} + 0 \times 2^{10} + 0 \times 2^{9} + 1 \times 2^{8} + \dots$$
$$\dots 1 \times 2^{7} +$$
$$= (100011)$$

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$$= (1000110)$$

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$$\dots 1 \times 2^{7} + 0 \times 2^{6} + 1 \times 2^{5} + 0 \times 2^{4} + \dots$$
$$= (100011010$$

$$(4521)_{10} = 1 \times 2^{12} + 0 \times 2^{11} + 0 \times 2^{10} + 0 \times 2^{9} + 1 \times 2^{8} + \dots$$

$$\dots 1 \times 2^{7} + 0 \times 2^{6} + 1 \times 2^{5} + 0 \times 2^{4} + \dots$$

$$\dots 1 \times 2^{3} + \dots$$

$$= (1000110101)$$

$$(4521)_{10} = 1 \times 2^{12} + 0 \times 2^{11} + 0 \times 2^{10} + 0 \times 2^{9} + 1 \times 2^{8} + \dots$$

$$\dots 1 \times 2^{7} + 0 \times 2^{6} + 1 \times 2^{5} + 0 \times 2^{4} + \dots$$

$$\dots 1 \times 2^{3} + 0 \times 2^{2} + \dots$$

$$= (10001101010$$

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$$= (100011010100)$$

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$$\dots 1 \times 2^{3} + 0 \times 2^{2} + 0 \times 2^{1} + 1 \times 2^{0}$$

$$= (1000110101010_{1})_{2}$$

You could use another basis, computers often use the basis 2:

$$(4521)_{10} = 1 \times 2^{12} + 0 \times 2^{11} + 0 \times 2^{10} + 0 \times 2^{9} + 1 \times 2^{8} + \dots$$

$$\dots 1 \times 2^{7} + 0 \times 2^{6} + 1 \times 2^{5} + 0 \times 2^{4} + \dots$$

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$$= (1000110101001)_{2}$$

In general:

$$(c_m \dots c_1 c_0)_q = c_0 q^0 + c_1 q^1 + \dots + c_m q^m, c \in \{0, 1, 2, \dots, q-1\}$$

- Numbers are stored in binary in the memory of a computer, in segments of a specific length (called a word).
- We distinguish multiple types of numbers:
 - Integers: -301, -1, 0, 1, 96, 2293, . . .
 - Floating points: -301.01, 0.01, 3.14159265, 14498.2
- A binary integer representation looks like the following bit sequence:

$$z = \sigma \left(c_0 2^0 + c_1 2^1 + \ldots + c_{\lambda - 1} 2^{\lambda - 1} \right)$$

 σ is the sign of z (+ or -), and λ is the length of the word

• Endianness: the order of bits stored by a computer

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 - Octal: =DEC20CT(214)
 - Hexadecimal: =DEC2HEX(214)

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- Excel:
 - Decimal: =DEC2BIN(214)
 - Octal: =DEC20CT(214)
 - Hexadecimal: =DEC2HEX(214)
- Matlab:
 - Decimal: dec2bin(214)
 - Other base: dec2base(214, <base>)

$$0+0=0$$

 $0+1=1$
 $1+0=1$
 $1+1=0$
(carry one)

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 $1 4 5$
 $+ 2 3$

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 $0+1=1$
 $1+0=1$
 $1+1=0$
(carry one)
 1 4 5
 $+$ 2 3
 1 6 8

Addition:

Subtraction:

$$0-0=0$$

 $1-0=1$
 $1-1=0$
 $0-1=1$
(borrow one)

Addition:

Subtraction:

$$0-0=0$$

 $1-0=1$
 $1-1=0$
 $0-1=1$
1 4 5
- 2 3
1 2 2

(borrow one)

1

0

Addition:

(carry one)

Subtraction:

$$0-0=0$$
 $1-0=1$
 $1-1=0$
 $0-1=1$
(borrow one)
 $1 \quad 4 \quad 5$
 $- \quad 2 \quad 3$

Addition:

Subtraction:

Multiplication and division are more expensive, and more elaborate

Command	Result
intmin	-2147483648

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intmin	-2147483648
intmax	2147483647

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intmin	-2147483648
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<pre>i = int16(intmax)</pre>	i = 32767
whos i	int16 information
i = i + 100	i = 32767

Command	Result
intmin	-2147483648
intmax	2147483647
<pre>i = int16(intmax)</pre>	i = 32767
whos i	int16 information
i = i + 100	i = 32767
realmax	1.7977e+308

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intmin	-2147483648
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<pre>i = int16(intmax)</pre>	i = 32767
whos i	int16 information
i = i + 100	i = 32767
realmax	1.7977e+308
f = 0.1	

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intmin	-2147483648
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whos f	double information

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whos f	double information
format long e	

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format long e	
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Result
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int16 information
i = 32767
1.7977e+308
double information
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i = i + 100	i = 32767
realmax	1.7977e+308
f = 0.1	
whos f	double information
format long e	
realmax	1.797693134862316e+308
f	
<pre>fprintf("%0.16f",f)</pre>	0.1000000000000000

Command	Result
intmin	-2147483648
intmax	2147483647
<pre>i = int16(intmax)</pre>	i = 32767
whos i	int16 information
i = i + 100	i = 32767
realmax	1.7977e+308
f = 0.1	
whos f	double information
format long e	
realmax	1.797693134862316e+308
f	
fprintf("%0.16f",f)	0.1000000000000000
fprintf("%0.20f",f)	0.1000000000000000555

• In Matlab, integers of the type int32 are represented by 32-bit words ($\lambda = 31$).

¹Matlab does not perform actual integer overflows, it just stops at the maximum

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- If, during a calculation, an integer number becomes larger than $2^{\lambda}-1$, the computer reports an overflow¹
- How can a computer identify an overflow?

¹Matlab does not perform actual integer overflows, it just stops at the maximum

Representation of real (floating point) numbers

 Formally, a real number is represented by the following bit sequence

$$x = \sigma \left(2^{-1} + c_2 2^{-2} + \ldots + c_m 2^{-m}\right) 2^{e-1023}$$

Here, σ is the sign of x and e is an integer value.

 A floating point number hence contains sections that contain the sign, the exponent and the mantissa



Image: Wikimedia Commons CC by-SA

Representation of real (floating point) numbers

• Example: $\lambda = 3$, m = 2, $x = \frac{2}{3}$

$$x = \pm \left(2^{-1} + c_2 2^{-2}\right) 2^e$$

- $c_0 \in \{0, 1\}$
- $e = \pm a_0 2^0$
- $a_0 \in \{0, 1\}$
- Truncation: $fl(x) = 2^{-1} = 0.5$
- Round off: $fl(x) = 2^{-1} + 2^{-2} = 0.75$

Today's outline

- Introduction
- Roundoff and truncation errors
- Break errors
- 4 Loss of digits
- 6 (Un)stable methods
- 6 Symbolic math
- Summary

Trigonometric, Logarithmic, and Exponential computations

- Processors can do logic and arithmetic instructions
- Trigonometric, logarithmic and exponential calculations are "higher-level" functions:
 exp, sin, cos, tan, sec, arcsin, arccos, arctan, log, ln, ...
- Such functions can be performed using these "low level" instructions, for instance using a Taylor series:

$$\sin(x) = \sum_{n=0}^{\infty} \frac{(-1)^n}{(2n+1)!} x^{2n+1} = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots$$
$$e^x = \sum_{n=0}^{\infty} \frac{x^n}{x!} = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \frac{x^4}{4!} + \dots$$

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- These operations involve many multiplications and additions, and are therefore *expensive*
- Computations can only take finite time, for infinite series, calculations are interrupted at N

$$\sin(x) = \sum_{n=0}^{N} \frac{(-1)^n}{(2n+1)!} x^{2n+1} = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \dots + \frac{(-1)^N}{(2N+1)!} x^{2N+1}$$

$$e^x = \sum_{n=0}^{N} \frac{x^n}{x!} = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \frac{x^4}{4!} + \dots + \frac{x^N}{N!}$$

This results in a break error

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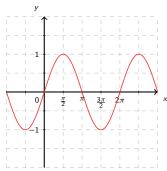
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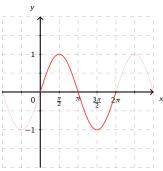
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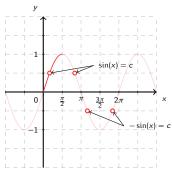


A computer may use a clever algorithm to limit the number of operations required to perform a higher-level function. A (fictional!) example for the computation of sin(x):

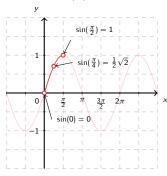
1 Use periodicity so that $0 < x < 2\pi$



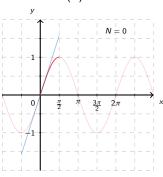
- 1 Use periodicity so that $0 \le x \le 2\pi$
- 2 Use symmetry $(0 \le x \le \frac{\pi}{2})$



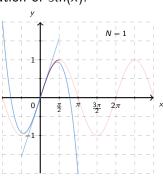
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- Use lookup tables for known values



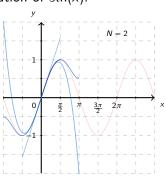
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- 4 Perform taylor expansion



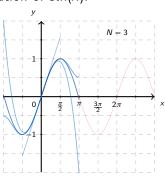
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Loss of digits

- During operations such as +, -, \times , \div , an error can add up
- Consider the summation of x and y

$$\tilde{x} - \delta \le x \le \tilde{x} + \delta$$
 and $\tilde{y} - \varepsilon \le y \le \tilde{y} + \varepsilon$

$$(\tilde{x} + \tilde{y}) - (\delta + \varepsilon) \le x + y \le (\tilde{x} + \tilde{y}) + (\delta + \varepsilon)$$

$$x = \pi, \tilde{x} = 3.1416$$

 $y = 22/7, \tilde{y} = 3.1429$

$$\begin{aligned}
x &= \pi, \tilde{x} = 3.1416 \\
y &= 22/7, \tilde{y} = 3.1429
\end{aligned}
\Rightarrow \begin{cases}
\delta &= \tilde{x} - x = 7.35 \times 10^{-6} \\
\varepsilon &= \tilde{y} - y = 4.29 \times 10^{-5}
\end{aligned}$$

$$x + y &= \tilde{x} + \tilde{y} \pm (\delta + \varepsilon) \approx 6.2845 - 5.025 \times 10^{-5} \\
x - y &= \tilde{x} - \tilde{y} \pm (\delta + \varepsilon) \approx -0.0013 + 3.55 \times 10^{-5}$$

- The absolute error is small ($\approx 10^{-5}$), but the relative error is much bigger (0.028).
- Adding up the errors results in a loss of significant digits!

- Calculate e^{-5}
 - Use the Taylor series
 - Calculate the first 26 terms (N = 26)
- Now repeat the calculation, but use for each calculation only 4 digits. What do you find?
- Without errors you would find: $e^{-5} = 0.006738$
- If you only use 4 digits in the calculations, you'll find 0.00998

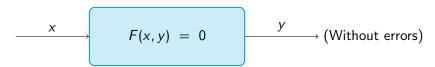
- Calculate e^{−5}
 - Use the Taylor series
 - Calculate the first 26 terms (N = 26)
- Now repeat the calculation, but use for each calculation only 4 digits. What do you find?
 - Use: str2double(sprintf('%.4g', term))
- Without errors you would find: $e^{-5} = 0.006738$
- If you only use 4 digits in the calculations, you'll find 0.00998

- Calculate e⁻⁵
 - Use the Taylor series
 - Calculate the first 26 terms (N = 26)
- Now repeat the calculation, but use for each calculation only 4 digits. What do you find?
- Use: str2double(sprintf('%.4g', term)) • Without errors you would find: $e^{-5} = 0.006738$
- If you only use 4 digits in the calculations, you'll find 0.00998

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Badly (ill) conditioned problems

We consider a system F(x, y) that computes a solution from input data. The input data may have errors:



Badly (ill) conditioned problems

We consider a system F(x, y) that computes a solution from input data. The input data may have errors:

$$F(x,y) = 0$$

$$x + \delta x$$

$$F(x + \delta x, y + \delta y) = 0$$

$$y + \delta y$$

$$(With errors)$$

Badly (ill) conditioned problems

We consider a system F(x, y) that computes a solution from input data. The input data may have errors:

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$$y(x + \delta x) - y(x) \approx y'(x)\delta x$$

Propagated error on the basis of Taylor expansion

$$C = \max_{\delta x} \left(\left| \frac{\delta y/y}{\delta x/x} \right| \right)$$

Condition criterion, C < 10 error development small

Solve the following linear system in Matlab using double and single precision:

$$A = \begin{bmatrix} 1.2969 & 0.8648 \\ 0.2161 & 0.1441 \end{bmatrix}, \quad x = \begin{bmatrix} 0.8642 \\ 0.1440 \end{bmatrix}$$

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Double precision

Single precision

```
>> clear; clc; format long e;
>> A = single(
   [[1.2969 0.8648];
   [0.2161 0.1441]] );
>> x = single(
   [0.8642; 0.1440] );
>> y = A\x
y =
   1.3331791e+00
   -1.0000000e+00
```

- Matlab already warned us about the bad condition number: Warning: Matrix is close to singular or badly scaled.
 Results may be inaccurate. RCOND = 1.148983e-08.
- The RCOND is the reciprocal condition number
- A small error in x results in a big error in y. This is called an ill conditioned problem.

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(Un)stable methods

- The condition criterion does not tell you anything about the quality of a numerical solution method!
- It is very well possible that a certain solution method is more sensitive for one problem than another
- If the method propagates the error, we call it an unstable method. Let's look at an example.

The Golden mean

• Let's evaluate the following recurrent relationship:

$$y_{n+1} = y_{n-1} - y_n$$

 $y_0 = 1, \quad y_1 = \frac{2}{1 + \sqrt{5}}$

• You can prove (by substitution) that:

$$y_n = x^{-n}$$
, $n = 0, 1, 2, ...$, $x = \frac{1 + \sqrt{5}}{2}$

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Recurrent version

```
% initialise
y(1) = 1;
y(2) = 2 / (1 + sqrt(5));

% Perform recurrent
    approach
for n = 2:39
    y(n+1) = y(n-1)-y(n);
end
```

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end
```

Powerlaw version

```
% initialise
x = (1 + sqrt(5))/2;
y2(1) = x^0; % n = 1

% Perform powerlaw apprach
for n = 0:39
    y2(n+1) = x^-n
end
```

n	Recurrent	Powerlaw
1	1.0000	1.0000
1	0.6180	0.6180
2	0.3820	0.3820
3	0.2361	0.2361
37	$3.080 \cdot 10^{-08}$	$2.995 \cdot 10^{-08}$
38	$1.714 \cdot 10^{-08}$	$1.851 \cdot 10^{-08}$
39	$1.366 \cdot 10^{-08}$	$1.144 \cdot 10^{-08}$
40	$3.485 \cdot 10^{-08}$	$7.071 \cdot 10^{-08}$

 The recurrent approach enlarges errors from earlier calculations!

n	Recurrent	Powerlaw
1	1.0000	1.0000
1	0.6180	0.6180
2	0.3820	0.3820
3	0.2361	0.2361
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Recall example 1, where the errors blew up our computation of 0.1, whereas they did not for 2. Why did we see these results?

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- The number 0.1 is not exactly represented in binary
 - A tiny error can accumulate up to catastrophic proportions!
- The number 2 does have an exact binary representation

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Start your calculation program of choice (Excel, Matlab, ...)

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Calculate the result of y:

$$y = e^{\pi} - \pi$$

Start your calculation program of choice (Excel, Matlab, ...)

Calculate the result of y:

$$y = e^{\pi} - \pi = 19.999099979$$

Start your calculation program of choice (Excel, Matlab, ...)

Calculate the result of y:

$$y = e^{\pi} - \pi = 19.999099979 \neq 20$$

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Calculate the result of y:

$$y = e^{\pi} - \pi = 19.999099979 \neq 20$$



DURING A COMPETITION, I TOLD THE PROGRAMMERS ON OUR TEAM THAT e^{π} - π WAS A STANDARD TEST OF FLOATING-POINT HANDLERS -- IT WOULD COME OUT TO 20 UNLESS THEY HAD ROUNDING ERRORS.





Image: xkcd

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Symbolic math packages

Definition

The use of computers to manipulate mathematical equations and expressions in symbolic form, as opposed to manipulating the numerical quantities represented by those symbols.

- Symbolic integration or differentiation, substitution of one expression into another
- Simplification of an expression, change of subject etc.
- Packages and toolboxes:

Symbolic math packages

- Mathematica Well known software package, license available via $\mathsf{TU/e}$
 - Maple Well known, license available via TU/e
- Wolfram Alpha Web-based interface by Mathematica developer.

 Less powerful in mathematical respect, but

 more accessible and has a broad application

 range (unit conversion, semantic commands).
 - Sage Open-source alternative to Maple, Mathematica, Magma, and MATLAB.
 - Matlab Symbolic math toolbox

$$f(x) = (x-1)(x+1)(x^2+1) + 1$$

$$f(x) = (x-1)(x+1)(x^2+1) + 1$$

```
>> syms x
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```

$$f(x) = \frac{1}{x^3 + 1}$$

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my_f_int = log(x + 1)/3 - log((x - 1/2)^2 + 3/4)/6 +

(3^(1/2)*atan((2*3^(1/2)*(x - 1/2))/3))/3
```

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   (3^{(1/2)}*atan((2*3^{(1/2)}*(x - 1/2))/3))/3
>> my_f_diff = diff(my_f_int)
my_f_diff = 1/(3*(x + 1)) + 2/(3*((4*(x - 1/2)^2)/3 +
   1)) - (2*x - 1)/(6*((x - 1/2)^2 + 3/4))
>> simplify(my_f_diff)
```

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   1)) - (2*x - 1)/(6*((x - 1/2)^2 + 3/4))
>> simplify(my_f_diff)
ans = 1/(x^3 + 1)
```

Exercise 1

Simplify the following expression:

$$f(x) = \frac{2\tan x}{(1 + \tan^2 x)}$$

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Simplify the following expression:

$$f(x) = \frac{2 \tan x}{(1 + \tan^2 x)} = \sin 2x$$

$$\Rightarrow \frac{1}{x} \sin x = \sin 2x$$

Exercise 2

Calculate the *value* of *p*:

$$p = \int_0^{10} \frac{e^x - e^{-x}}{\sinh x} dx$$

Exercise 1

Simplify the following expression:

$$f(x) = \frac{2 \tan x}{(1 + \tan^2 x)} = \sin 2x$$
>> simplify (2*tan(x)/(1 + tan(x)^2))

Exercise 2

Calculate the *value* of p:

$$p = \int_0^{10} \frac{e^x - e^{-x}}{\sinh x} dx$$
>> f = ((exp(x) - exp(-x))/sinh(x));
>> p = int(f,0,10)
p = 20

A root finding method searches for the values where a function reaches zero. We will cover the numerical methods later, here we show how to use root finding with symbolic math in Matlab.

Symbolic math function

$$f(x) = \frac{3}{x^2 + 3x} - 2$$

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Symbolic math function

$$f(x) = \frac{3}{x^2 + 3x} - 2$$

```
>> syms x
>> f = 3 / (x^2 + 3*x) - 2;
>> solve(f)
ans =
   15^(1/2)/2 - 3/2
- 15^(1/2)/2 - 3/2
```

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Symbolic math function

$$f(x) = \frac{3}{x^2 + 3x} - 2$$

Function as a string

$$f(x) = x^2 - 4x + 2$$

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Function as a string

$$f(x) = x^2 - 4x + 2$$

Symbolic math toolbox: variable precision arithmetic

Variable precision can be used to specify the number of significant digits.

Today's outline

- Introduction
- Roundoff and truncation errors
- Break errors
- 4 Loss of digits
- (Un)stable methods
- 6 Symbolic math
- Summary

- Numerical errors mar arise due to truncation, roundoff and break errors, which may seriously affect the accuracy of your solution
- Errors may propagate and accumulate, leading to smaller accuracy
- Ill-conditioned problems and unstable methods have to be identified so that proper measures can be taken
- Symbolic math computations may be performed to solve certain equations algebraically, bypassing numerical errors, but this is not always possible.

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