Numerical errors in computer simulations

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Numerical Methods (6BER03), 2024-2025

Today's outline

- Introduction
- Roundoff and truncation errors
- Break errors
- Loss of digits
- (Un)stable methods
- Symbolic math
- Summary



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Start your spreadsheet program (Excel, ...)



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Enter:

Cell	Value
A1	0.1



Start your spreadsheet program (Excel, ...)

Enter:

Cell	Value
A1	0.1
A2	=(A1*10)-0.9



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Enter:

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A1	0.1
A2	=(A1*10)-0.9
А3	=(A2*10)-0.9
A4	=(A3*10)-0.9

(repeat until A30)



Start your spreadsheet program (Excel, ...)

Enter:

Cell	Value
A1	0.1
A2	=(A1*10)-0.9
A3	=(A2*10)-0.9
A4	=(A3*10)-0.9

(repeat until A30)

What's happening?



Start your spreadsheet program (Excel, ...)

Enter:

Call

Cen	value
A1	0.1
A2	=(A1*10)-0.9
A3	=(A2*10)-0.9
A4	=(A3*10)-0.9

Value

Enter:

Cell	Value
A1	2

(repeat until A30)

What's happening?



Start your spreadsheet program (Excel, ...)

Enter:

Cell	Value
A1	0.1
A2	=(A1*10)-0.9
A3	=(A2*10)-0.9
A4	=(A3*10)-0.9

Enter:

Cell	Value
A1	2
A2	=(A1*10)-18
A3	=(A2*10)-18
A4	=(A3*10)-18

(repeat until A30)

(repeat until A30)

What's happening?



Start Python



Start Python

Investigate the result of np.sin(1e40 * np.pi)



Start Python

Investigate the result of np.sin(1e40 * np.pi)

Create a vector \mathbf{v} containing the powers of 10, e.g., from 10^0 up to 10^{40} , and solve $\mathbf{np.sin}(\mathbf{v} * \mathbf{np.pi})$:

```
import numpy as np
import matplotlib.pyplot as plt

v = np.logspace(0, 40, 41)
y = np.sin(v * np.pi)
plt.loglog(v, np.abs(y)) # Double log plot, values on y-axis must be positive.
plt.show()
```



Errors in computer simulations

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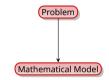


Verification and validation

Problem

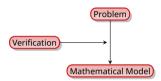


Verification and validation





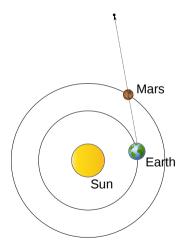
Verification and validation



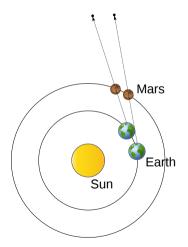
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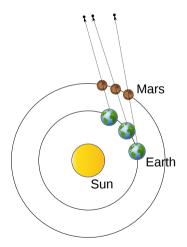




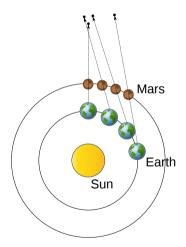




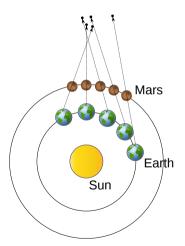




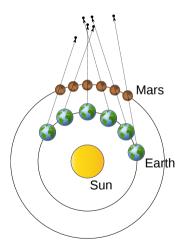




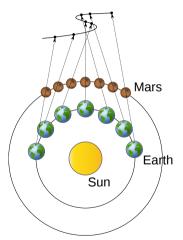


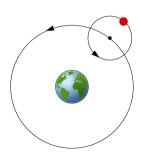




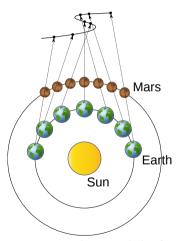


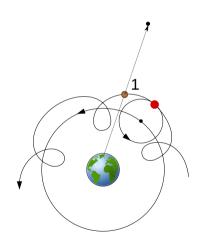




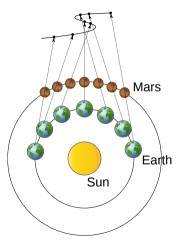


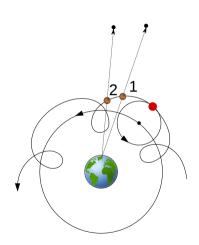




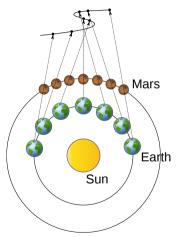


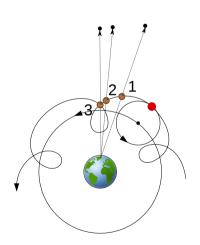




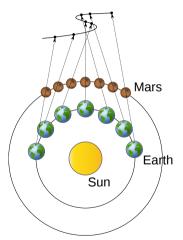


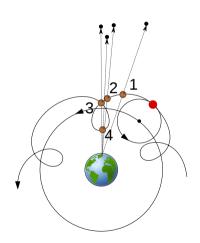




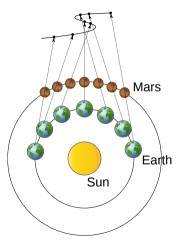


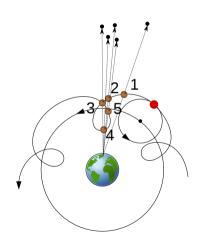




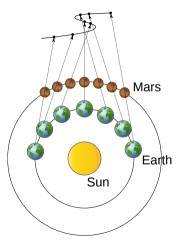


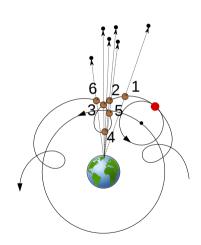




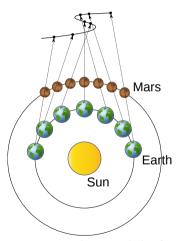


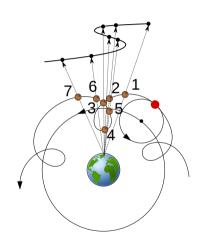












- The perceived orbit of Mars from Earth shows a zig-zag (in contrast to the Sun, Mercury, Venus)
- Even though they were not 'right', Earth-centered models (Ptolemy) were still valid



Be aware of your uncertainties

Aleatory uncertainty

Uncertainty that arises due to inherent randomness of the system, features that are too complex to measure and take into account

Epistemic uncertainty

Uncertainty that arises due to lack of knowledge of the system, but could in principle be known



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Errors in the mathematical model (physics)

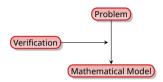


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- Errors in the mathematical model (physics)
- Errors in the program (implementation)

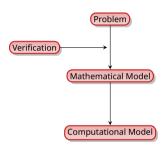




Verification

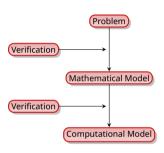


Verification





Verification





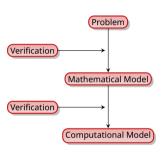
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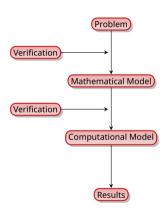


Verification





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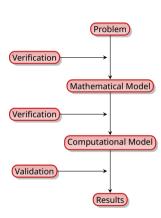


Verification

Verification is the process of mathematically and computationally assuring that the model computes the equations you intended to implement.

Validation

Validation is the process of determining the degree to which a model is an accurate representation of the real world from the perspective of the intended uses of the model





- Errors in the mathematical model (physics)
- Errors in the program (implementation)
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Introduction

A numerical result \tilde{x} is an approximation of the real value x.

Absolute error

$$\delta = |\tilde{x} - x|, x \neq 0$$

• Relative error

$$\frac{\delta}{\tilde{x}} = |\frac{\tilde{x} - x}{\tilde{x}}|$$

• Error margin

$$\tilde{x} - \delta \le x \le \tilde{x} + \delta$$



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• \bar{x} has m significant digits if the absolute error in x is smaller or equal to 5 at the (m+1)-th position:

$$10^{q-1} \le |\bar{x}| \le 10^q$$
$$|x - \bar{x}| = 0.5 \times 10^{q-m}$$

• For example:

$$x = \frac{1}{3}, \tilde{x} = 0.333 \Rightarrow \delta = 0.000333333...$$

3 significant digits



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```
4 5 2 1

Number of 1's: 1 \times 10^{(0)}

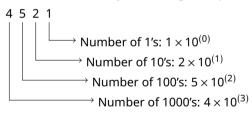
Number of 10's: 2 \times 10^{(1)}

Number of 100's: 5 \times 10^{(2)}

Number of 1000's: 4 \times 10^{(3)}
```



- Computers represent a number with a finite number of digits: each number is therefore an approximation due to roundoff and truncation errors.
- In the decimal system, a digit c at position n has a value of $c \times 10^{n-1}$



$$(4521)_{10} = 4 \times 10^3 + 5 \times 10^2 + 2 \times 10^1 + 1 \times 10^0$$



$$(4521)_{10} =$$

$$=($$



$$(4521)_{10} = 1 \times 2^{12} +$$

$$=(1$$



$$(4521)_{10} = 1 \times 2^{12} + 0 \times 2^{11} +$$

$$=(10$$



$$(4521)_{10} = 1 \times 2^{12} + 0 \times 2^{11} + 0 \times 2^{10} +$$



$$(4521)_{10} = 1 \times 2^{12} + 0 \times 2^{11} + 0 \times 2^{10} + 0 \times 2^9 + 0 \times 2^{11} + 0 \times 2$$

$$=(1000$$



$$(4521)_{10} = 1 \times 2^{12} + 0 \times 2^{11} + 0 \times 2^{10} + 0 \times 2^9 + 1 \times 2^8 + \dots$$

$$=(10001$$



$$(4521)_{10} = 1 \times 2^{12} + 0 \times 2^{11} + 0 \times 2^{10} + 0 \times 2^{9} + 1 \times 2^{8} + \dots$$

$$\dots 1 \times 2^{7} + \dots$$

$$= (100011)$$



$$(4521)_{10} = 1 \times 2^{12} + 0 \times 2^{11} + 0 \times 2^{10} + 0 \times 2^{9} + 1 \times 2^{8} + \dots$$

$$\dots 1 \times 2^{7} + 0 \times 2^{6} + \dots$$

$$= (1000110)$$



$$(4521)_{10} = 1 \times 2^{12} + 0 \times 2^{11} + 0 \times 2^{10} + 0 \times 2^{9} + 1 \times 2^{8} + \dots$$

$$\dots 1 \times 2^{7} + 0 \times 2^{6} + 1 \times 2^{5} + \dots$$

$$= (10001101)$$



$$(4521)_{10} = 1 \times 2^{12} + 0 \times 2^{11} + 0 \times 2^{10} + 0 \times 2^9 + 1 \times 2^8 + \dots$$
$$\dots 1 \times 2^7 + 0 \times 2^6 + 1 \times 2^5 + 0 \times 2^4 + \dots$$



$$(4521)_{10} = 1 \times 2^{12} + 0 \times 2^{11} + 0 \times 2^{10} + 0 \times 2^{9} + 1 \times 2^{8} + \dots$$

$$\dots 1 \times 2^{7} + 0 \times 2^{6} + 1 \times 2^{5} + 0 \times 2^{4} + \dots$$

$$\dots 1 \times 2^{3} + \dots$$

$$= (1000110101$$



$$(4521)_{10} = 1 \times 2^{12} + 0 \times 2^{11} + 0 \times 2^{10} + 0 \times 2^9 + 1 \times 2^8 + \dots$$

$$\dots 1 \times 2^7 + 0 \times 2^6 + 1 \times 2^5 + 0 \times 2^4 + \dots$$

$$\dots 1 \times 2^3 + 0 \times 2^2 + \dots$$

$$= (10001101010$$



$$(4521)_{10} = 1 \times 2^{12} + 0 \times 2^{11} + 0 \times 2^{10} + 0 \times 2^9 + 1 \times 2^8 + \dots$$

$$\dots 1 \times 2^7 + 0 \times 2^6 + 1 \times 2^5 + 0 \times 2^4 + \dots$$

$$\dots 1 \times 2^3 + 0 \times 2^2 + 0 \times 2^1 + \dots$$

$$= (100011010100$$



• You could use another basis, computers often use the basis 2:

$$\begin{aligned} (4521)_{10} &= 1 \times 2^{12} + 0 \times 2^{11} + 0 \times 2^{10} + 0 \times 2^9 + 1 \times 2^8 + \dots \\ &\dots 1 \times 2^7 + 0 \times 2^6 + 1 \times 2^5 + 0 \times 2^4 + \dots \\ &\dots 1 \times 2^3 + 0 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 \\ &= (1000110101001)_2 \end{aligned}$$



• You could use another basis, computers often use the basis 2:

$$(4521)_{10} = 1 \times 2^{12} + 0 \times 2^{11} + 0 \times 2^{10} + 0 \times 2^{9} + 1 \times 2^{8} + \dots$$

$$\dots 1 \times 2^{7} + 0 \times 2^{6} + 1 \times 2^{5} + 0 \times 2^{4} + \dots$$

$$\dots 1 \times 2^{3} + 0 \times 2^{2} + 0 \times 2^{1} + 1 \times 2^{0}$$

$$= (1000110101001)_{2}$$

In general:

$$(c_m \dots c_1 c_0)_q = c_0 q^0 + c_1 q^1 + \dots + c_m q^m, c \in \{0, 1, 2, \dots, q-1\}$$



- Numbers are stored in binary in the memory of a computer, in segments of a specific length (called a *word*).
- We distinguish multiple types of numbers:
 - Integers: -301, -1, 0, 1, 96, 2293, ...
 - Floating points: -301.01, 0.01, 3.14159265, 14498.2
- A binary integer representation looks like the following bit sequence:

$$z = \sigma \left(c_0 2^0 + c_1 2^1 + \dots + c_{\lambda - 1} 2^{\lambda - 1} \right)$$

 σ is the sign of z (+ or -), and λ is the length of the word



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$$214_{10} = 11010110_2$$



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- Excel:
 - Decimal: =DEC2BIN(214)
 - Octal: =DEC20CT(214)
 - Hexadecimal: =DEC2HEX(214)



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- Excel:
 - Decimal: =DEC2BIN(214)
 - Octal: =DEC20CT(214)
 - Hexadecimal: =DEC2HEX(214)
- Python:
 - Decimal: bin(214)
 - Octal: oct(214)
 - Hexadecimal: hex(214)



$$0+0=0$$

 $0+1=1$
 $1+0=1$
 $1+1=0$ (carry one)



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Addition:

$$0+0=0$$

 $0+1=1$
 $1+0=1$
 $1+1=0$ (carry one)

Subtraction:

$$0-0=0$$

 $1-0=1$
 $1-1=0$
 $0-1=1$ (borrow one)

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 $1-1=0$
 $0-1=1$ (borrow one)
 $1 4 5$
 $- 2 3$
 $1 2 2$

Multiplication and division are more expensive, and more elaborate



Command	Result
np.iinfo(np.int32).min	-2147483648



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np.iinfo(np.int32).min	-2147483648
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np.iinfo(np.int32).min	-2147483648
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<pre>i = np.int16(np.iinfo(np.int32).max)</pre>	i = 32767



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np.iinfo(np.int32).min	-2147483648
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<pre>i = np.int16(np.iinfo(np.int32).max)</pre>	i = 32767
type(i)	<pre><class 'numpy.int16'=""></class></pre>



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np.iinfo(np.int32).min	-2147483648
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<pre>i = np.int16(np.iinfo(np.int32).max)</pre>	i = 32767
type(i)	<class 'numpy.int16'=""></class>
i = i + 100	i = 32767



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np.finfo(np.float64).max	1.7976931348623157e+308



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f = 0.1	



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np.finfo(np.float64).max	1.7976931348623157e+308
f = 0.1	
type(f)	<class 'float'=""></class>



Command	Result
np.iinfo(np.int32).min	-2147483648
np.iinfo(np.int32).max	2147483647
<pre>i = np.int16(np.iinfo(np.int32).max)</pre>	i = 32767
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<pre>print(":.16f".format(f))</pre>	0.1000000000000000
print(":.20f".format(f))	0.10000000000000000555



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 automatically promotes it to a larger type if possible. In the case of numpy, an overflow
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 'np.int32' type.
- How can a computer identify an overflow? Usually through error reporting mechanisms
 within the language. Python's numpy library, for instance, will throw an OverflowError
 when an operation results in a number larger than what can be represented by the
 current data type.

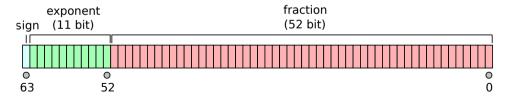
Representation of real (floating point) numbers

Formally, a real number is represented by the following bit sequence

$$x = \sigma (2^{-1} + c_2 2^{-2} + \dots + c_m 2^{-m}) 2^{e-1023}$$

Here, σ is the sign of x and e is an integer value.

 A floating point number hence contains sections that contain the sign, the exponent and the mantissa







Representation of real (floating point) numbers

• Example: $\lambda = 3$, m = 2, $x = \frac{2}{3}$

$$x = \pm \left(2^{-1} + c_2 2^{-2}\right) 2^e$$

- $c_0 \in \{0, 1\}$
- $e = \pm a_0 2^0$
- $a_0 \in \{0, 1\}$
- Truncation: $fl(x) = 2^{-1} = 0.5$
- Round off: $fI(x) = 2^{-1} + 2^{-2} = 0.75$



Introduction Roundoff and truncation errors Break errors Loss of digits (Un)stable methods Symbolic math Summar

Today's outline

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Trigonometric, Logarithmic, and Exponential computations

- Processors can do logic and arithmetic instructions
- Trigonometric, logarithmic and exponential calculations are "higher-level" functions: exp, sin, cos, tan, sec, arcsin, arccos, arctan, log, ln, ...
- Such functions can be performed using these "low level" instructions, for instance using a Taylor series:

$$\sin(x) = \sum_{n=0}^{\infty} \frac{(-1)^n}{(2n+1)!} x^{2n+1} = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots$$
$$e^x = \sum_{n=0}^{\infty} \frac{x^n}{x!} = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \frac{x^4}{4!} + \dots$$



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- Computations can only take finite time, for infinite series, calculations are interrupted at N

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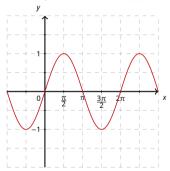
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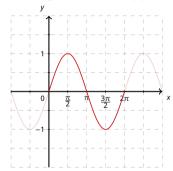






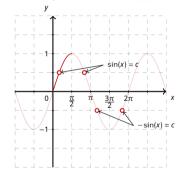
A computer may use a clever algorithm to limit the number of operations required to perform a higher-level function. A (fictional!) example for the computation of sin(x):

1 Use periodicity so that $0 \le x \le 2\pi$



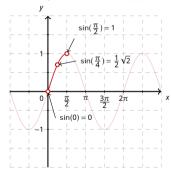


- **1** Use periodicity so that $0 \le x \le 2\pi$
- 2 Use symmetry $(0 \le x \le \frac{\pi}{2})$



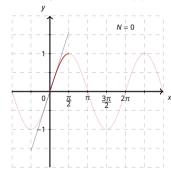


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- **3** Use lookup tables for known values



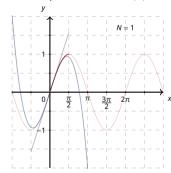


- **1** Use periodicity so that $0 \le x \le 2\pi$
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- Perform taylor expansion



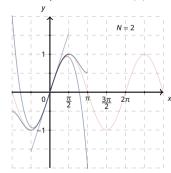


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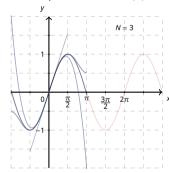


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Loss of digits

- During operations such as +, -, \times , \div , an error can add up
- Consider the summation of *x* and *y*

$$\tilde{x} - \delta \le x \le \tilde{x} + \delta$$
 and $\tilde{y} - \varepsilon \le y \le \tilde{y} + \varepsilon$

$$(\tilde{x} + \tilde{y}) - (\delta + \varepsilon) \le x + y \le (\tilde{x} + \tilde{y}) + (\delta + \varepsilon)$$



$$x = \pi, \tilde{x} = 3.1416$$

 $y = 22/7, \tilde{y} = 3.1429$



$$x = \pi, \tilde{x} = 3.1416 y = 22/7, \tilde{y} = 3.1429$$
 $\Rightarrow \delta = \tilde{x} - x = 7.35 \times 10^{-6}$
$$\varepsilon = \tilde{y} - y = 4.29 \times 10^{-5}$$



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$$x + y = \tilde{x} + \tilde{y} \pm (\delta + \varepsilon) \approx 6.2845 - 5.025 \times 10^{-5}$$

 $x - y = \tilde{x} - \tilde{y} \pm (\delta + \varepsilon) \approx -0.0013 + 3.55 \times 10^{-5}$

- The absolute error is small ($\approx 10^{-5}$), but the relative error is much bigger (0.028).
- Adding up the errors results in a loss of significant digits!



- Calculate e^{-5}
 - Use the Taylor series
 - Calculate the first 26 terms (N = 26)
- Now repeat the calculation, but use for each calculation only 4 digits. What do you find?
- Without errors you would find: $e^{-5} = 0.006738$
- If you only use 4 digits in the calculations, you'll find 0.00998



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$$y(x + \delta x) - y(x) \approx y'(x)\delta x$$

Propagated error on the basis of Taylor expansion

$$C = \max_{\delta x} \left(\left| \frac{\delta y/y}{\delta x/x} \right| \right)$$

Condition criterion, *C* < 10 error development small



Solve the following linear system in Python using double and single precision:

$$A = \begin{bmatrix} 1.2969 & 0.8648 \\ 0.2161 & 0.1441 \end{bmatrix}, \quad x = \begin{bmatrix} 0.8642 \\ 0.1440 \end{bmatrix}$$



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Double precision

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import numpy as np

# Double precision
A = np.array([[1.2969, 0.8648], [0.2161, 0.1441]], dtype=np.float64)
x = np.array([0.8642, 0.1440], dtype=np.float64)
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- Python does not automatically warn about bad condition number. You need to compute and check it manually using NumPy:
 - if np.linalg.cond(A)> threshold: raise("Ill conditioned error")
- The cond number is the condition number computed using NumPy.
- A small error in A (the matrix) results in a big error in the solution. This is called an ill-conditioned problem.



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(Un)stable methods

- The condition criterion does not tell you anything about the quality of a numerical solution method!
- It is very well possible that a certain solution method is more sensitive for one problem than another
- If the method propagates the error, we call it an *unstable method*. Let's look at an example.



The Golden mean

The Golden Mean is a well known ratio and one of the solutions of $\phi^2 = \phi + 1$, such that $\phi = \frac{1+\sqrt{5}}{2}$, and $\phi - 1 = \phi^{-1}$. Many relations to compute ϕ have been established (which pose often an interesting source for creating small test programs).

• Let's evaluate the following recurrent relationship:

$$y_{n+1} = y_{n-1} - y_n$$

$$y_0 = 1$$
, $y_1 = 1 - \phi = \phi^{-1} = \frac{2}{1 + \sqrt{5}}$

• Alternatively, a closed form power law relation exists, which again computes y_n :

$$y_n = x^{-n}$$
, $n = 0, 1, 2, ...$, $x = \frac{1 + \sqrt{5}}{2}$



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```
Recurrent version
  import numpy as np
  def golden_mean_recurrent(Ntot):
     # Initialize the series with the given
          initial conditions
     y = np.zeros(Ntot)
     v[0] = 1
     v[1] = 2 / (1 + np.sgrt(5))
     # Perform the recurrence to fill in the rest
          of the series
     for n in range(2, Ntot):
        v[n] = v[n-1] - v[n-2]
     return v
12
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Power law version

```
import numpy as np

def golden_mean_powerlaw(Ntot):
    # Initialize the constant value
    x = (1 + np.sqrt(5)) / 2

# Generate a range of values from 0 to Ntot
    and apply the power law
    y = x ** -np.arange(0, Ntot + 1)

return y
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- Compare the outcomes: plot(goldenMeanPowerlaw(40)- goldenMeanRecurrent(40))
- See what happens if you use single precision (add dtype=np.float32 to line 5 (np.zeros) and line 8 (np.arange)).

n	Recurrent	Power law
0	1.0000	1.0000
1	0.6180	0.6180
2	0.3820	0.3820
3	0.2361	0.2361
37	$1.714 \cdot 10^{-08}$	$1.851 \cdot 10^{-08}$
38	1.366 · 10 ⁻⁰⁸	$1.144 \cdot 10^{-08}$
39	3.485 · 10 ⁻⁰⁸	$7.071 \cdot 10^{-09}$
40	$1.017 \cdot 10^{-08}$	$4.370 \cdot 10^{-09}$

• The recurrent approach enlarges errors from earlier calculations!



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- The number 2 does have an exact binary representation



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Example 2 (large sine series)

The np.sin(1e40*np.pi) result gives poor results because 1e40 has an error margin on the order of floating-point machine epsilon, which is roughly 1×10^{-16} in Python (double-precision floating-point format). In Python, as in many computing environments, the number of $2\cdot\pi$ cycles is still much larger than $10^{40}\cdot 10^{-16}$. Also, π is not stored with an infinite number of digits, which further contributes to the imprecision.



Start your calculation program of choice (Excel, Python, ...)



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Calculate the result of *y*:

$$y = e^{\pi} - \pi$$



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Calculate the result of *y*:

$$y = e^{\pi} - \pi = 19.999099979$$



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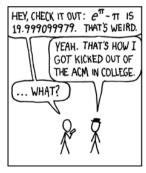
$$y = e^{\pi} - \pi = 19.999099979 \neq 20$$



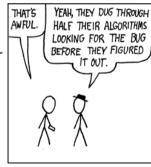
Start your calculation program of choice (Excel, Python, ...)

Calculate the result of *y*:

$$y = e^{\pi} - \pi = 19.999099979 \neq 20$$



DURING A COMPETITION, I TOLD THE PROGRAMMERS ON OUR TEAM THAT e^{π} - π WAS A STANDARD TEST OF FLOATING-POINT HANDLERS -- IT WOULD COME OUT TO 20 UNLESS THEY HAD ROUNDING ERRORS.





Today's outline

- Introduction
- Roundoff and truncation errors
- Break errors
- Loss of digits
- (Un)stable methods
- Symbolic math
- Summary



Symbolic math packages

Definition

The use of computers to manipulate mathematical equations and expressions in symbolic form, as opposed to manipulating the numerical quantities represented by those symbols.

- Symbolic integration or differentiation, substitution of one expression into another
- Simplification of an expression, change of subject etc.
- Packages and toolboxes:



Symbolic math packages

Mathematica Well known software package, license available via TU/e

Maple Well known, license available via TU/e

Wolfram Alpha Web-based interface by Mathematica developer. Less powerful in mathematical respect, but more accessible and has a broad application range (unit conversion, semantic commands).

Sage Open-source alternative to Maple, Mathematica, Magma, and MATLAB.

Matlab Symbolic math toolbox

Python SymPy



Symbolic math: simplify

$$f(x) = (x-1)(x+1)(x^2+1)+1$$



Symbolic math: simplify

$$f(x) = (x-1)(x+1)(x^2+1) + 1$$

```
from sympy import symbols, simplify
x = symbols('x') # Alt: from sympy.abc import x,y,z
f = (x - 1)*(x + 1)*(x**2 + 1) + 1
f_simplified = simplify(f)
print(f_simplified)
```



Symbolic math: simplify

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```
f_simplified = x**4
```



Symbolic math: integration and differentiation

$$f(x) = \frac{1}{x^3 + 1}$$



Symbolic math: integration and differentiation

$$f(x) = \frac{1}{x^3 + 1}$$

```
from sympy import symbols, simplify, integrate, diff

x = symbols('x')
f = 1/(x**3+1)
my_f_int = integrate(f, x)
my_f_diff = diff(my_f_int, x)
my_f_diff_simplified = simplify(my_f_diff)
print(my_f_diff_simplified)
```



Symbolic math: integration and differentiation

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smy_f_diff_simplified = simplify(my_f_diff)
sprint(my_f_diff_simplified)
```

 $my_f_diff_simplified = 1/(x**3 + 1)$



Exercise 1

Simplify the following expression:

$$f(x) = \frac{2 \tan x}{(1 + \tan^2 x)}$$



Exercise 1

Simplify the following expression:

$$f(x) = \frac{2\tan x}{(1 + \tan^2 x)} = \sin 2x$$

from sympy import symbols, trigsimp, tan

```
x = symbols('x')
expr = 2*tan(x)/(1 + tan(x)**2)
simplified_expr = trigsimp(expr)
```



Exercise 2

Calculate the *value* of *p*:

$$p = \int_0^{10} \frac{e^x - e^{-x}}{\sinh x} dx$$



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$$p = \int_0^{10} \frac{e^x - e^{-x}}{\sinh x} dx$$

```
from sympy import exp, sinh, integrate, symbols
x = symbols("x")
f = ((exp(x)- exp(-x))/sinh(x)).simplify()
p = integrate(f, (x, 0, 10))
print(p)
```



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$$p = 20$$



Symbolic math: root finding

A root finding method searches for the values where a function reaches zero. We will cover the numerical methods later, here we show how to use root finding with symbolic math in Python.

Symbolic math function

$$f(x) = \frac{3}{x^2 + 3x} - 2$$

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Symbolic math function

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```
from sympy import solve, symbols
x = symbols("x")
f = 3 / (x**2 + 3*x) - 2
solutions = solve(f, x)
print(solutions)
```

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```

$$solutions = [-3/2 + sqrt(15)/2, -sqrt(15)/2 - 3/2]$$

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- Numerical errors may arise due to truncation, roundoff and break errors, which may seriously affect the accuracy of your solution
- Errors may propagate and accumulate, leading to smaller accuracy
- Ill-conditioned problems and unstable methods have to be identified so that proper measures can be taken
- Symbolic math computations may be performed to solve certain equations algebraically, bypassing numerical errors, but this is not always possible.



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