GC06 Database Design Report

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Group 37

Table of Contents

- 1. Project Overview
- 2. YouTube Video Link
- 3. Entity Relationship Diagram
- 4. Database Schema Listing
- 5. Application Architecture
- 6. Normalistation Analysis
- 7. Query Explainations

Project Overview

This report highlights the implementation of a peer-to-peer online auction system, called **Hashtagories**. The interaction model is based largely on Twitter; product descriptions are limited to 140 characters, and bids on an auction appear like Twitter @-replies. Descriptions include hashtags, which make up the main categorisation and search index.

The live hosted version of the application can be viewed at http://gc06.jasongwartz.com/app/.

Main Features

- Users can add an item to their profile, giving it a title and description (including hashtags)
- A registered item can be put up for auction, with a reserve price (hidden from other users) and end date
- Users can search for items up for auction based on their hashtags
- A user can place a bid on an item up for auction
- At the end time, the auction closes the seller and highest bidder are notified, and both are invited to leave feedback
- Users can 'watch' an auction, where they will be notified when new bids are placed

YouTube Video Link

https://www.youtube.com/watch?v=izUbymUykH4

Entity Relationship Diagram

Figure 1 shows the entity relationship diagram that was used in the final production of Hashtagories. It contains 8 tables.

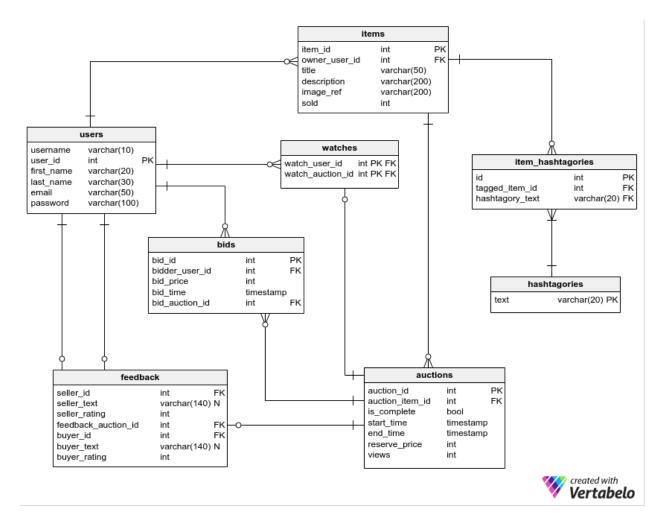


Figure 1: Auction Entity Relationship Diagram

Database Schema Listing

Tables

```
(Notes in brackets; \mathrm{PK} = \mathrm{Primary} \ \mathrm{Key}, \ \mathrm{FK} = \mathrm{Foreign} \ \mathrm{Key})
```

- users
 - user_id (PK)
 - username
 - first_name
 - last_name
 - email
 - password (hashed)
- items
 - item_id (PK)
 - owner_user_id (FK)
 - title
 - description
 - $-image_ref$
 - sold
- hashtagories
 - id (PK)
 - text (indexed)
- item_hashtagories
 - id (PK)
 - tagged_item_id (FK)
 - hashtagory_text (FK, FULLTEXT indexed)
- auctions
 - auction id (PK)
 - auction_item_id (FK)
 - is_complete
 - start_time
 - $\ \operatorname{end_time}$
 - $-\ {\bf reserve_price}$
 - views
- bids
 - bid id (PK)
 - bidder_user_id (FK)
 - bid_price
 - bid_time
 - bid_auction_id (FK)
- watches
 - watch_user_id (PK, FK)
 - watch_auction_id (PK, FK)
- feedback
 - feedback_auction_id (PK, FK)
 - seller_id (FK)
 - buyer_id (FK)
 - $-\ \mathrm{buyer_text}$
 - seller_text
 - buyer_ratingseller rating

Foreign Key Relationships

The database features a number of foreign key constraints to ensure referential integrity. The foreign key relationships are summarised in the following chart, a reduced ERD showing only the properties of each entity which are involved in constraints:

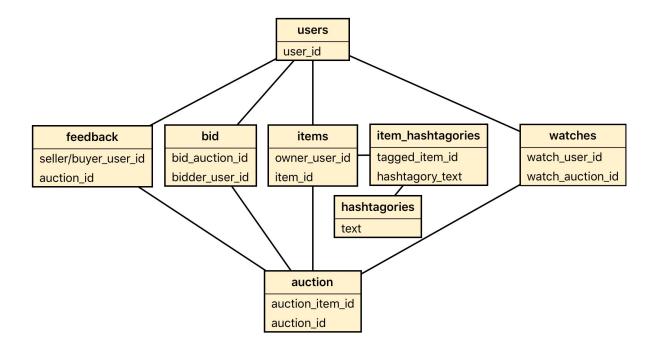


Figure 2: fk-chart

This chart shows the sequence of foreign key dependencies. As can be seen, the users and the hashtagories tables are the only table with no foreign key dependencies, and the item_hashtagories table is composed entirely of foreign key fields.

Additional Schema Details

In addition to the standard schema definitions, our schema also makes use of the following techniques:

- MySQL STORED PROCEDURES
 - Every query used in the application is implemented through one or more stored procedures. This provides additional safety from SQL injection, as well as allowing the database engine to optimise the procedure's query plan. Additionally, many API endpoints utilise the same stored procedure, meaning any query changes or adjustments can be made in a single location; this provides an efficient layer of abstraction.
- A MySQL VIEW, entitled auctions_retrieve_all
 - This view is used frequently to retrieve composited information about all current auctions, or a subset of all current auctions (a search)
- MySQL FULLTEXT binary searching, on the item_hashtagories table
 - This was employed to allow matching of multiple partial strings to multiple partial values (eg. a search for "table chair", modified to "table* chair*" will return any item tagged with "table" or "chair", as well as "tables", "tabletop", or "chairs". FULLTEXT searching required an additional index on the given column.

Application Architecture

This project takes advantage of several modern web architecture paradigms and industry best practices. While some were functionally unnecessary at the scale of this project, designing for expansion or scalability is

- Multi-level system achitecture
 - Our system is designed in several layers to provide appropriate abstraction and encapsulation, allowing for easier debugging and design changes. The application layers are:
 - * An AngularJS browser-based single-page application
 - * A PHP HTTP REST API (that's a lot of acronyms: PHP HyperText Preprocessor HyperText Transfer Protocol Representative State Transfer Application Programming Interface)
 - * A MySQL database with stored procedures
- Utilising cloud services
 - For the entirety of development, our app has been hosted on Amazon Web Services virtual servers. In order to aid in rapid deployment, a script was developed to automatically update the server-hosted application when a new commit was pushed on Github.
- Distributed functionality
 - In designing the system deployment, inspiration was taken from contemporary, scalable deployment practices. In many large web applications, the web server and database server are seperate physical or virtual machines, often replicated several times over, with traffic routed via a load balancer. We used the industry-standard Amazon Web Services tools replicate this structure in our application: the web application and REST API are hosted on an Amazon EC2 (Elastic Cloud Compute) virtual server, and the MySQL database is hosted on an Amazon RDS (Relational Database Service) virtual server.
- RESTful HTTP API design
 - Inspired by other prolific web APIs, we designed the PHP API in a REST standard format. This involves a series of nested endpoints, organised in a tree structure; each endpoint performs a given database operation, manipulating the retrieved data if necessary, and returning it to the caller. These endpoints are often parameterised (using URL parameters) to retrieve specific data sets. In our project, an example URL request might be: GET api/auctions/search?query=table
- Task automation using CRON
 - There are two tasks that need to be completed once per minute in order to maintain the system's core functionality: completing expired auctions (including emailing the seller and the winning bidder and opening the feedback portal), and notifying "watchers" of auctions about new bids. Both of these tasks can be triggered via an API endpoint the PHP script runs the necessary stored procedure (which performs most of the required logic), and then sends any required emails. In order to run these tasks regularly, the server uses the Unix application cron to, every minute, send an HTTP request via curl to each of these endpoints.

Normalisation Analysis

A database is in 3rd normal form if it meets 3 criteria: 1. It contains only atomic values and there are no repeating groups. 2. All non-key attributes are fully functional dependent on the primary key. 3. There is no transitive functional dependency.

1st Normal Form

All of our eight tables contain only atomic values meaning that there are no elements in any of the tables where the data can be further broken up. i.e. username is atomic, email is atomic, an item_description is atomic. The database is in first normal form. Also, any groups related to an entity have been separated in a separate table such as a the hashtagories table.

2nd Normal Form

This means that in every table in the database a value of a particular non-key field cannot be uniquely identified via another non-key or group of non-key fields. Attribute B is dependent on attribute A, but not on a proper subset of A, then B is fully functional dependent on A. For example, item_description can not be uniquely identified by an proper subset of the item_id such as title or owner_user_id.

3rd Normal Form

A transitive dependency is when a non-key attribute, C, is dependent on another attribute, A, via an attribute, B. Our database contains no transitive dependencies as we put all key attributes that are functional dependent on another attribute as primary keys in different tables. One transitive dependency may have occurred if we stored the hashtagories_text on the items table, as then the hashtagories_text would've been dependent on item_id via item_description and therefore transitively dependent. To avoid this we separated an items hashtagories tags into a separate table.

Table of normalisation

By analysing the attributes in tables of the database we ensured that none of the attributes break the rules of normalisation. As the table below indicates, all the tables in the database meet all the requirements of 3rd normalisation.

Attribute Name	1st Normal Form	2nd Normal Form	3rd Normal Form
users	yes	yes	yes
items	yes	yes	yes
hashtagories	yes	yes	yes
item_hashtagories	yes	yes	yes
auctions	yes	yes	yes
bids	yes	yes	yes
watches	yes	yes	yes
feedback	yes	yes	yes

Query Explainations

Below are all SQL queries used by the system. In writing them as stored procedures, we have encapsulated their functionality and hopefully improved readability and modularity/extensibility. Furthermore, because stored procedures can be parameterised and all parameters are passed as data, we are insulated against more common SQL injection attacts. Arguably, it gives better performance as the database does not have to interpret a SQL string from PHP every time it is called (they are already stored in executable form), however the speed impact of this is debated. Please refer to here and here for a more detailed discussion of stored procedures.

This also ensures that the middleware is very light-weight, simply passing data to and from the front-end via PHP scripts that call stored procedures.

Additionally, we make use of events for repeated tasks, to prevent the middleware from having to call the database regularly and hence consume the connections. We also make use of one view for the sake of performance (auctions retrieve all).

auctions_cancel

Cancels an auction by deleting it from the auction table.

```
PROCEDURE `auctions_cancel`(IN auction_id INT(11))

BEGIN

DELETE FROM `auctions` WHERE auctions.auction_id = auction_id;

END
```

auctions_create

Starts an auction, using timestamps generated in the middle-layer (could have also used NOW() function). Reserve price needs to be cast from string input.

```
PROCEDURE `auctions_create`(IN auction_item_id INT(11),
IN start_time timestamp,
IN end_time timestamp,
IN reserve_price varchar(12))
BEGIN
    INSERT INTO `auctions` (auctions.auction_item_id,
    auctions.start_time,
    auctions.end_time,
    auctions.reserve_price)
    VALUES(auction_item_id,
    start_time,
    end_time,
    CAST(reserve_price AS DECIMAL(10,2)));
    SELECT last_insert_id();
END
```

auctions retrieve all

This stored procedure shows a set of auction objects with user name, item details and auction details, as well as highest bid. This is actually achieved by selecting from a view (code below). Note: Highest bids needed to be selected into a subtable and then joined to cope with some auctions not having bids placed.

```
PROCEDURE `auctions_retrieve_all`()
BEGIN
    SELECT * FROM auctions_retrieve_all;
END
#View creation
CREATE TABLE `auctions_retrieve_all` (
   `auction_id` INT(11) NOT NULL DEFAULT '0',
   `auction_item_id` INT(11) NOT NULL,
   `is_complete` TINYINT(1) NOT NULL DEFAULT '0',
   `start_time` TIMESTAMP NOT NULL,
   `end_time` TIMESTAMP NOT NULL,
   `reserve_price` DECIMAL(10) NOT NULL,
   'views' INT(11) UNSIGNED ZEROFILL NOT NULL DEFAULT '00000000000',
   `current_bid` DECIMAL(10) NULL DEFAULT NULL,
   'item_id' INT(11) NULL DEFAULT '0',
   `owner_user_id` INT(11) NULL DEFAULT NULL,
   `title` VARCHAR(50) NULL DEFAULT '',
   `description` VARCHAR(200) NULL DEFAULT '',
   'image_ref' VARCHAR(255) NULL DEFAULT NULL,
   `sold` INT(1) NULL DEFAULT '0',
   `username` VARCHAR(10) NULL DEFAULT NULL
) ENGINE=MyISAM;
```

auctions search

Allows for a partial search of auctions, returning auction items with item and user details, sorted in a variety of modes. The stored procedure is also structured in such a way that it returns the full result set of all open auctions sorted in the way specified by the sort parameter. This allows us to use a single end-point to display both feed and search data sorted in the requisite order.

auctions_search_desc

Mirror of previous search procedure with descending sorting. This minimised the parameters that needed to be passed through the middleware (the PHP endpoints select which procedure to call based on a sorting parameter passed from the front-end).

```
PROCEDURE `auctions_search_desc`(IN str varchar(100), IN sort varchar(10))

BEGIN

SELECT DISTINCT *

FROM item_hashtagories AS ih, auctions_retrieve_all AS a

WHERE (str = '' OR MATCH(ih.hashtagory_text) AGAINST(str IN BOOLEAN MODE))

AND ih.tagged_item_id = a.item_id

GROUP BY a.auction_id

ORDER BY CASE sort

when 'start_time' then a.start_time
when 'end_time' then a.end_time
when 'views' then a.views
when 'title' then a.title

END
DESC
;
END
```

auctions_self

Returns the information about a particular auction (bids are pulled in seperately for display purposes). Also updates the view field of the auction record, as this procedure is always called every time an auction is viewed.

```
PROCEDURE auctions_self (IN auction_id INT(11))
BEGIN
```

```
UPDATE auctions
SET auctions.views = auctions.views+1
WHERE auctions.auction id = auction id;
SELECT items.*, users.username, auctions.* FROM `auctions`
    LEFT JOIN `items` ON auctions.auction_item_id = items.item_id
    LEFT JOIN 'users' ON items.owner user id = users.user id
    WHERE auctions.auction id = auction id;
END
auctions_user_auctions
Returns all auctions created by a particular user.
PROCEDURE `auctions_user_auctions`(IN user_id INT(11))
BEGIN
    SELECT * FROM `auctions` AS a, `items` AS i
    WHERE a.is complete = 0
    AND a.auction_item_id = i.item_id
    AND i.owner_user_id = user_id
    ORDER BY `end_time` ASC;
END
auctions_user_feed
Returns a feed of auctions relevant to the bids that they have made.
PROCEDURE `auctions_user_feed`(IN user_id INT(11))
    SELECT * FROM auctions as a
    LEFT JOIN items as i ON a.auction_item_id = i.item_id
    LEFT JOIN bids as b ON a.auction_id = b.bid_auction_id
    WHERE b.bidder_user_id = user_id AND a.is_complete = 0
    ORDER BY a.end_time ASC;
END
bids auction bids
Returns the set of bids on any given auction.
PROCEDURE `bids_auction_bids`(IN bid_auction_id INT(11))
BEGIN
    SELECT * FROM bids WHERE bids.bid auction id = bid auction id
    ORDER BY bids.bid_price DESC;
END
```

bids create

Procedure called when bid is placed on item. Bid value validation (that new bid is the higest) can be performed on the front-end and is also duplicated here. Here, the stored procedure first checks the highest bid against the attempted bid, and only inserts the bid if it is higher. It will also add the auction the the user's watchlist by calling the stored procedure.

```
PROCEDURE `bids_create`( IN bidder_user_id INT(11), IN bid_price VARCHAR(12), IN bid_auction_id INT)
```

```
BEGIN
    DECLARE highest_bid DECIMAL(10,2) DEFAULT 0;
    SELECT bid_price FROM `bids`
    WHERE bids.bid_auction_id = bid_auction_id
    ORDER BY bid price DESC LIMIT 1 INTO highest bid;
    SELECT bid_price_in, highest_bid;
    IF bid_price_in > highest_bid THEN
    INSERT INTO `bids` (bids.bidder_user_id,
    bids.bid_price,
    bids.bid_time,
    bids.bid_auction_id)
        VALUES(bidder_user_id,
        CAST(bid_price_in AS DECIMAL(10,2)),
        NOW(),
        bid auction id);
    SELECT last_insert_id();
ELSE
     SIGNAL SQLSTATE '45000' SET MESSAGE_TEXT = 'Bid price too low!';
IF;
    CALL watches_create(bidder_user_id, bid_auction_id);
END
bids_self
Returns the auction that a bid is related to.
PROCEDURE `bids_self`(IN bid_id INT(11))
BEGIN
    SELECT * FROM auctions AS a
    LEFT JOIN items AS i ON a.auction_item_id = i.item_id
    LEFT JOIN bids as b ON a.auction_id = b.bid_auction_id
    WHERE b.bid_id = bid_id;
END
bids user bids
Returns auction data objects on which a given user has bid (auction details, item deails and bid details).
PROCEDURE `bids_user_bids`(IN user_id INT(11))
BEGIN
    SELECT a.*, b.*, i.*, u.username FROM users AS u, auctions AS a
    LEFT JOIN items AS i ON a.auction_item_id = i.item_id
    LEFT JOIN bids AS b ON a.auction_id = b.bid_auction_id
    WHERE user_id = b.bidder_user_id AND i.owner_user_id = u.user_id
    ORDER BY b.bid_time DESC;
END
```

event_end_expired_auctions

Ends all expired auctions and returns the data on the auctions such as the buyer, seller, name, winning bid etc. By running the event on a timer, the database does not need to have its connection tied up by an automated call from the API.

```
PROCEDURE `event_end_expired_auctions`()
BEGIN
   DECLARE reserve_price_tmp INT DEFAULT 0;
   DECLARE highest bid tmp decimal(10,2) DEFAULT 0;
   DECLARE auction id tmp INT DEFAULT 0;
   DECLARE item_title_tmp varchar(200) DEFAULT 0;
   DECLARE item id tmp INT DEFAULT 0;
   DECLARE seller_url_tmp varchar(200) DEFAULT 0;
   DECLARE buyer_url_tmp varchar(200) DEFAULT 0;
   DECLARE seller_username_tmp varchar(200) DEFAULT 0;
   DECLARE buyer_username_tmp varchar(200) DEFAULT 0;
   DECLARE seller_email_tmp varchar(200) DEFAULT 0;
   DECLARE buyer_email_tmp varchar(200) DEFAULT 0;
   DECLARE seller_id_tmp INT DEFAULT 0;
   DECLARE buyer_id_tmp INT DEFAULT 0;
   DECLARE successful tmp INT DEFAULT 0;
   DECLARE n INT DEFAULT 0;
   DECLARE i INT DEFAULT 0;
        # Drops the temporary table if it exists. Then creates it.
       DROP TABLE IF EXISTS `tmp_end_expired_auctions`;
        CREATE TABLE `tmp_end_expired_auctions` (
          `auction_id` int(11) NOT NULL,
          `seller_username` varchar(200) DEFAULT NULL,
          `seller_email` varchar(200) DEFAULT NULL,
          `seller_feedback_url` varchar(50) DEFAULT NULL,
          `buyer_username` varchar(200) DEFAULT NULL,
          `buyer_email` varchar(200) DEFAULT NULL,
          `buyer_feedback_url` varchar(50) DEFAULT NULL,
          `item_title` varchar(200) DEFAULT NULL,
          `final_bid_price` varchar(200) DEFAULT NULL,
          `successful` varchar(200) DEFAULT NULL,
          PRIMARY KEY (`auction_id`)
        ) ENGINE=InnoDB DEFAULT CHARSET=latin1;
       SELECT count(*) FROM `auctions` WHERE end time < now() AND is complete=0 INTO n;
       # This loops through the unclosed auctions in the auctions table select values
        # from other tables. It evaluates whether an auction failed or was successful
        # every loop inserts a row to the tmp table.
       SET i=0;
        WHILE i<n DO
            # Auction table selects. Gets the auction_id, reserver_price and item_id.
            SELECT auction_id, reserve_price, auction_item_id FROM `auctions`
            WHERE end_time < now() AND is_complete=0 ORDER BY auction_id LIMIT i, 1
                INTO auction id tmp, reserve price tmp, item id tmp;
            # Bids table selects. Gets the highest bid on an item, and the user id
```

```
# of that bid.
SELECT bid_price, bidder_user_id FROM `bids`
WHERE bid_auction_id = auction_id_tmp ORDER BY bid_price DESC LIMIT 1
    INTO highest_bid_tmp, buyer_id_tmp;
# Items table selects. Gets the user id of the seller and the items title.
SELECT owner user id, title FROM `items`
WHERE item_id = item_id_tmp INTO seller_id_tmp, item_title_tmp;
# Users table selects seller. Gets seller username and email.
SELECT username, email FROM `users`
WHERE user_id = seller_id_tmp INTO seller_username_tmp, seller_email_tmp;
# If it was successful and there is a buyer user_id
IF buyer_id_tmp > 0 THEN
    # Users table selects buyer. Gets buyer username and email.
    SELECT username, email FROM `users`
    WHERE user_id = buyer_id_tmp INTO buyer_username_tmp, buyer_email_tmp;
END IF;
# If successful auction: create feedback, set successful to 1
IF highest bid tmp >= reserve price tmp AND highest bid tmp > 0 THEN
    # Create feedback
    INSERT IGNORE INTO `feedback`
        (`seller_id`,
        `feedback_auction_id`,
        `buyer id`)
        VALUES
        (seller_id_tmp,
        auction_id_tmp,
        buyer_id_tmp);
    # Sold field in items is set to the buyer id.
    UPDATE `items` SET sold = buyer_id_tmp WHERE item_id = item_id_tmp;
    # Sets successful.
    SET successful_tmp = 1;
    SET seller url tmp = CONCAT('#/feedback?', seller id tmp);
    SET buyer_url_tmp = CONCAT('#/feedback?', buyer_id_tmp);
ELSE
    SET successful_tmp = 0;
END IF;
# Inserts all the values into the tmp table.
INSERT INTO `tmp_end_expired_auctions`
(`auction_id`,
`seller_username`,
`seller email`,
```

```
`seller_feedback_url`,
            `buyer_username`,
            `buyer email`,
            `buyer_feedback_url`,
            `item_title`,
            `final_bid_price`,
            `successful`)
            VALUES
            (auction_id_tmp,
            seller_username_tmp,
            seller_email_tmp,
            seller_url_tmp,
            buyer_username_tmp,
            buyer_email_tmp,
            buyer_url_tmp,
            item_title_tmp,
            highest_bid_tmp,
            successful_tmp);
            SET i = i + 1;
       END WHILE;
        # Gets row count of tmp table.
        SELECT count(*) FROM `tmp_end_expired_auctions` INTO n;
        # Loops through the tmp table finally updating the is_complete
        # in auctions table
        # to 1.
        SET i=0;
        WHILE i<n DO
            SELECT auction_id FROM `tmp_end_expired_auctions` LIMIT i,1
                INTO auction_id_tmp;
            # Closes every expired auction
            UPDATE `auctions` SET is_complete = 1 WHERE auction_id = auction_id_tmp;
            SET i = i + 1;
        END WHILE;
        # Finally does an output select that is returned to the user.
        SELECT * FROM `tmp_end_expired_auctions`;
        # Drops the tmp table.
        DROP TABLE IF EXISTS `tmp_end_expired_auctions`;
End
event_retrieve_watches
```

Event that retrieves a user's watchlist with all relevant information (bid price, user id, item information). Used for periodic updates on a user's watchlist.

```
PROCEDURE `event_retrieve_watches`()
BEGIN
```

```
SELECT b.bid_price,
   b.bidder_user_id,
   w.watch_user_id,
   u.username,
   IF (b.bidder_user_id != w.watch_user_id,
   u.email, NULL) as email,
   a.end time,
    i.title,
    i.owner_user_id
   FROM bids b, watches w, users u, auctions a, items i
   WHERE b.bid_price = (SELECT max(bid_price)
   FROM bids where b.bid_auction_id = bid_auction_id)
   AND w.watch_auction_id = b.bid_auction_id
   AND w.watch_user_id = u.user_id
    AND a.auction_id = w.watch_auction_id
   AND a.auction_item_id = i.item_id
    AND bid_time >= NOW() - INTERVAL 1 MINUTE;
END
```

feedback_for_auction

Returns all fields for feedback on a given auction based on auction ID

```
PROCEDURE `feedback_for_auction`(IN feedback_auction_id INT(11))
BEGIN
SELECT * FROM feedback
WHERE feedback.feedback_auction_id = feedback_auction_id;
END
```

feedback_for_user

Returns all feedback for a user, both where they are a buyer or a seller. This was designed so that the system could display all of a user's feedback on one page.

```
PROCEDURE `feedback_for_user`(IN user_id INT(11))
BEGIN

SELECT feedback.*,

users.username as other_username
FROM feedback,

users

WHERE (users.user_id = feedback.seller_id

OR users.user_id = feedback.buyer_id)

AND users.user_id != user_id

AND (feedback.seller_id = user_id

OR feedback.buyer_id = user_id);
```

feedback_update

END

Updates the feedback table with feedback left by the users. Contains switching logic to allow feedback to be sent to a single PHP endpoint. Once this stored procedure is activated it first identifies whether the user is leaving feedback as a buyer or a seller, then selectively updates the feedback table. Also validates the feedback score to be within the 0-100 range

```
PROCEDURE `feedback_update` (IN feedback_text VARCHAR(140), IN
   feedback_rating DECIMAL(5,2), IN user_id INT(11), IN
   feedback auction id INT(11))
BEGIN
    # Feedback validation
    IF feedback rating <= 100 AND feedback rating >= 0 THEN
   #Check whether the feedback being left is by the buyer or the seller
            set @v1 = (SELECT seller_id
   FROM feedback
   WHERE feedback.feedback_auction_id = feedback_auction_id);
            IF @v1 = user_id THEN
                    UPDATE feedback
        SET feedback.seller_text = feedback_text,
        seller_rating = feedback_rating
        WHERE feedback.feedback_auction_id = feedback_auction_id;
            ELSE
                    UPDATE feedback
        SET feedback.buyer text = feedback text,
        buyer_rating = feedback_rating
        WHERE feedback.feedback_auction_id = feedback_auction_id;
            END IF;
   ELSE
            SIGNAL SQLSTATE '45000' SET MESSAGE_TEXT = 'Feedback rating
            out of range!';
   END IF;
END
hashtagories_all
```

Returns all possible hashtagories, for the front end to auto-suggest existing hashtagories.

```
PROCEDURE `hashtagories all`()
BEGIN
    SELECT text FROM hashtagories;
END
```

hashtagories_clear

Clears hashtagories associated with an item when item description is updated.

```
PROCEDURE `hashtagories_clear`(IN item_id INT(11))
   BEGIN
   DELETE FROM item_hashtagories
   WHERE tagged_item_id = item_id;
END
```

hashtagories_search

Searches the hashtagory table for hashtagories that partially match the input string, in order for the front-end to be able to search by hashtagory.

```
PROCEDURE `hashtagories_search`(IN str varchar(20))
BEGIN
```

```
SELECT text FROM hashtagories
WHERE INSTR(text, str);
END
```

hashtagories_self

Updates the hashtagory table if the user tags an item with a hashtagory that doesn't exist

```
PROCEDURE `hashtagories_self`(IN hashtext VARCHAR(20))
BEGIN
INSERT IGNORE INTO `hashtagories` values(hashtext);
END
```

hashtagories tag item

Adds a hashtagory to an item, allowing it to be searched at a later point. Structured so that if the hashtagory already exists, the operation will not crash and continue on to the second operation (actually tagging the item).

hashtagories_trending

Returns a ranked and numbered list of hastagories based on their popularity (numbers of items that are associated to them). Takes advantage of the GROUP BY command and COUNT aggregate function to return a ranked, depulicated list. Limits the trending hashtagories to 10.

```
PROCEDURE `hashtagories_trending`()

BEGIN

SELECT ih.hashtagory_text, COUNT(*) as count
FROM items i, item_hashtagories ih, auctions a

WHERE i.item_id = ih.tagged_item_id

AND a.`auction_item_id` = i.`item_id`

AND a.`is_complete` = 0

GROUP BY ih.hashtagory_text ORDER BY count DESC

LIMIT 10;

END
```

items_create

Creates an entry in the items table. Called every time an item is added to the system (note that hashtagorising has been encapsulated into a seperate stored procedure).

```
PROCEDURE `items_create`(IN owner_user_id int(11), IN title varchar(50), IN description varchar(200))
BEGIN
```

```
INSERT INTO `items` (`owner_user_id`, `title`, `description`)
VALUES(owner_user_id, title, description);
SELECT last_insert_id();
END
```

items_delete

Deletes an item from the database.

```
PROCEDURE `items_delete`(IN item_id INT(11))
BEGIN
    DELETE FROM `items` where items.item_id = item_id;
END
```

$items_self$

Returns all details regarding a single item. Uses GROUP_CONCAT to flatten the hashtagories (to which items have a 1..* relationship) into a single string for improved data transmission and display. Images are not hosted by Hashtagories, instead users submit a weblink to the photo of the item hosted by any third party or themselves.

```
PROCEDURE `items_self`(in item_id int(11))
BEGIN

SELECT `item_id`, `owner_user_id`, `title`,
  `description`, `image_ref`,
  GROUP_CONCAT(`hashtagory_text`
  ORDER BY `hashtagory_text` SEPARATOR ',') AS 'hashtagory_text'
  FROM `items`, `item_hashtagories`
  WHERE items.item_id = item_id
  AND item_hashtagories.tagged_item_id = item_id;
END
```

items_update

Updates description of item; note that hashtagory updates have been encapsulated into a seperate stored procedure. Note: hashtagories are updated by calling a seperate stored procedure from the middleware.

```
PROCEDURE `items_update`(IN item_id int(11),
IN title varchar(50),
IN description varchar(200))
BEGIN
    UPDATE `items`
    SET items.title = title, items.description = description,
    items.image_ref = image_ref WHERE items.item_id = item_id;
END
```

items_user_items

Returns all items that a user owns, in order for them to review, delete or put them up for auction. Flattens hashtagories (to which items have a 1..* relationship) into a single string for display purposes. Shows which items have been sold.

```
PROCEDURE `items_user_items`(IN owner_user_id INT(11))
BEGIN
SELECT `sold`, `item_id`, `owner_user_id`,
```

```
`title`, `description`, `image_ref`,
GROUP_CONCAT(`hashtagory_text`
ORDER BY `hashtagory_text`
SEPARATOR ',') AS 'hashtagory_text'
   FROM `items` as I LEFT OUTER JOIN `item_hashtagories` as IH
   ON I.item_id = IH.tagged_item_id
   WHERE I.owner_user_id = owner_user_id
   GROUP BY I.item_id ORDER BY I.item_id DESC;
END
```

users_authenticate

Authentictes users. Checks user name and hashed password against users table, only returns true if both match.

```
PROCEDURE `users_authenticate`(IN username varchar(20),
IN password varchar(20))
BEGIN
select user_id
FROM users where BINARY users.username = username
AND BINARY users.password = password;
END
```

users_change_password

Two-step password change. New password validation done in front end, if new password is entered correctly twice, this store procedure is triggered. Checks first if old user name / password pair exists, then updates password field. Also has secondary error-checking in the back-end.

```
PROCEDURE `users_change_password`(
IN userid int(11),
IN old password varchar(20),
IN new_password varchar(20))
BEGIN
    set @v1 = (select users.user_id
   FROM `users` where BINARY users.password = old_password
    AND users.user_id = userid);
   IF @v1 = userid THEN
        UPDATE `users` SET `password`= new_password WHERE `user_id`
         = userid;
   ELSE
       SIGNAL SQLSTATE '45000'
        SET MESSAGE_TEXT = 'INCORRECT USER NAME AND/OR PASSWORD';
   END IF;
END
```

users_create

Creates a new user. Validates that no fields are empty in the front end, then inserts into user table. Passwords are hashed using an MD5 function. Furthermore, password format is validated on the front-end (alphanumeric with at least one uppercase and one lowercase letter).

```
PROCEDURE `users_create`(
IN username varchar(10),
IN first name varchar(20),
```

```
IN last name varchar(20),
IN email varchar(50),
IN pass varchar(20)
)
BEGIN
    INSERT INTO `users` (`username`, `first_name`, `last_name`,
    `email`, `password`)
    values (username, first_name, last_name, email, pass);
END
users_rating
Displays the rating for a user based on their average feedback (both as buyer and as seller).
PROCEDURE `users_rating`(IN user_id INT(11))
BEGIN
    select IFNULL((IFNULL(s.seller rating,b.buyer rating)
    + IFNULL(b.buyer_rating,s.seller_rating))/2, 0) as rating FROM
    (select avg(buyer_rating) as seller_rating
    FROM `feedback`
    WHERE seller_id = user_id
    AND buyer_rating is not null) as s,
    (select avg(seller_rating) as buyer_rating
    FROM `feedback`
    WHERE buyer_id = user_id
    AND seller_rating is not null) as b;
END
users search
A search for all users whose usernames contain a substring that matches the input.
PROCEDURE `users_search` (in unstring varchar(20))
BEGIN
    SELECT username FROM users
WHERE INSTR(username, unstring);
END
users self
Returns all fields except for password of a particular user
PROCEDURE `users_self`(IN user_id int(11))
BEGIN
SELECT users.username,
users.user_id,
users.first_name,
users.last_name,
users.email
FROM `users` WHERE users.user_id = user_id;
END
users update
```

Updates a user. First checks that the system is attempting to update a valide user. Accepts empty strings and does not update those fields: uses an IF statement with IS NOT NULL to prevent overwriting of unsubmitted fields with empty strings.

```
PROCEDURE `users_update`(
IN user_id int,
IN username varchar(10),
IN first_name varchar(20),
IN last name varchar(30),
IN email varchar(50)
)
BEGIN
    IF user_id IS NULL THEN
        SIGNAL SQLSTATE '45000'
            SET MESSAGE_TEXT = 'No user provided';
    END IF;
    IF username IS NOT NULL THEN
        UPDATE `users` SET users.username = username
        WHERE users.user id = user id;
    END IF;
    IF first_name IS NOT NULL THEN
        UPDATE `users` SET users.first_name = first_name
        WHERE users.user_id = user_id;
    END IF;
    IF last_name IS NOT NULL THEN
        UPDATE `users` SET users.last name = last name
        WHERE users.user_id = user_id;
    END IF;
    IF email IS NOT NULL THEN
        UPDATE `users` SET users.email = email
        WHERE users.user_id = user_id;
    END IF;
END
users_username
Returns the user name of a user given their numerical ID.
PROCEDURE `users_username`(IN id INT(11))
BEGIN
    SELECT username FROM users
    WHERE user_id LIKE id;
END
watches create
Adds a 'watch' to the watches table when a user wishes to watch but not bid on an auction.
PROCEDURE `watches_create`(IN watch_user_id INT(11),
IN watch_auction_id INT(11))
BEGIN
    INSERT IGNORE INTO watches
    VALUES(watch_user_id, watch_auction_id);
    SELECT last_insert_id();
END
```

$watches_delete$

END

```
Removes a particular 'watch' when a user wishes to stop watching an item.
PROCEDURE `watches_delete`(IN watch_user_id INT(11),
IN watch_auction_id INT(11))
BEGIN
    DELETE FROM watches
    WHERE watches.watch_user_id = watch_user_id
    AND watches.watch_auction_id = watch_auction_id;
END
watches_user_watches
Returns all auctions that a user is watching.
PROCEDURE `watches_user_watches`(IN user_id INT(11))
BEGIN
    SELECT * FROM auctions AS a
    LEFT JOIN items AS i ON a.auction_item_id = i.item_id
    LEFT JOIN watches AS w ON a.auction_id = w.watch_auction_id
    WHERE w.watch_user_id = user_id;
```