

# Proof Reconstruction in Classical Propositional Logic

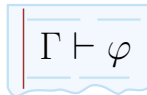
Work in Progress

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Agda Implementors' Meeting XXV  
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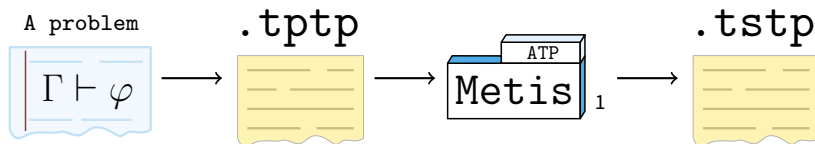
A problem


$$\Gamma \vdash \varphi$$

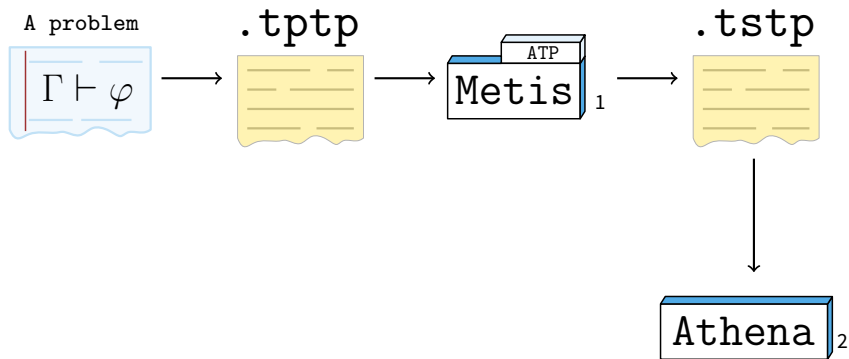




<sup>1</sup>It is available at <http://www.gilith.com/software/metis>

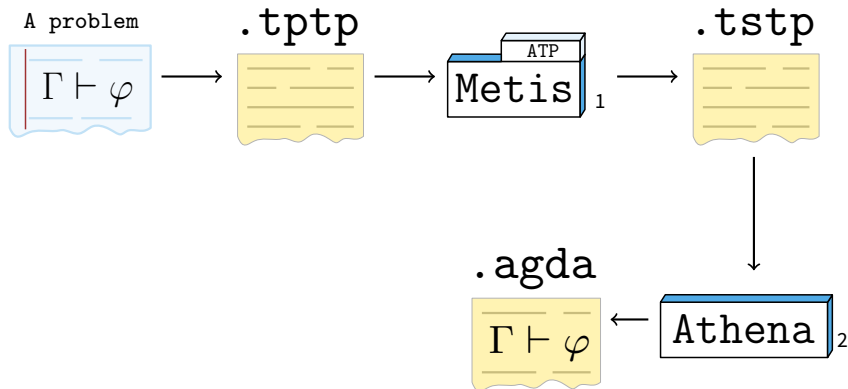


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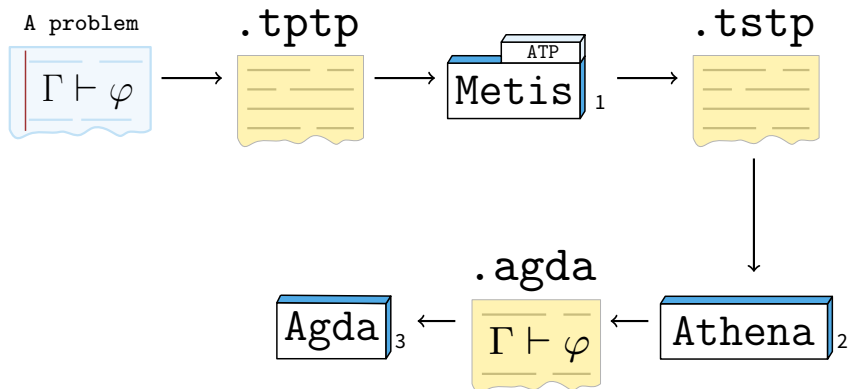
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# Proof Reconstruction: Overview

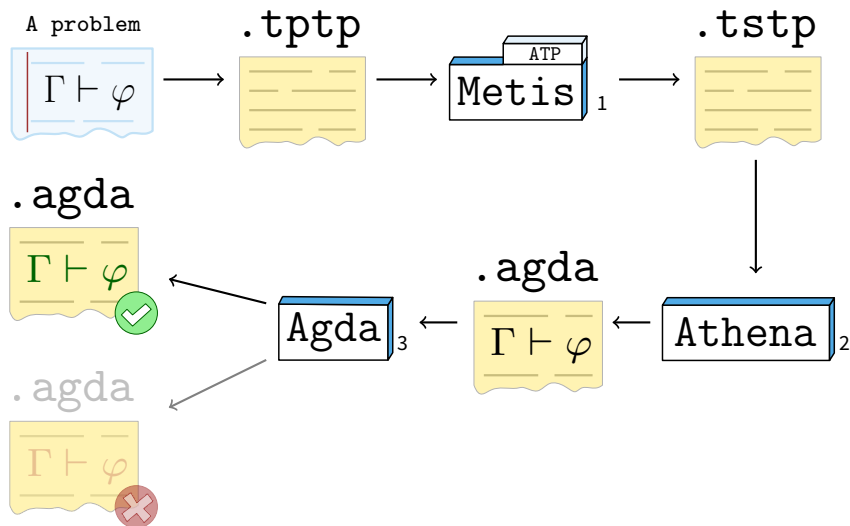


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### SledgeHammer

- ▶ Isabelle/HOL tool
- ▶ Metis ported within Isabelle
- ▶ Reconstruct proofs of well-known ATPs: EProver, Vampire, among others

### Integrating Waldmeister and Agda

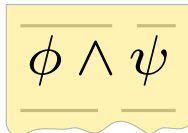
- ▶ Source code not available
- ▶ Equational Logic
- ▶ Reflection Layers

## Related Work: Apia

*Proving first-order theorems written in Agda using automatic theorem provers for first-order logic*

At the moment, the communication between Agda and the ATPs is unidirectional because the ATPs are being used as oracles (Sicard-Ramírez, 2015).

. agda



+ ATP-Pragma

```
$ cat Or.agda
module Or where

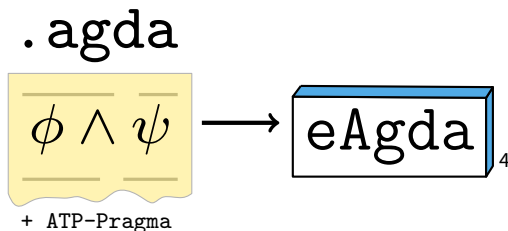
data _or_ (A B : Set) : Set where
  inj1  : A -> A or B
  inj2  : B -> A or B

postulate
  A B      : Set
  or-comm  : A or B -> B or A
  {-# ATP prove or-comm #-}
```

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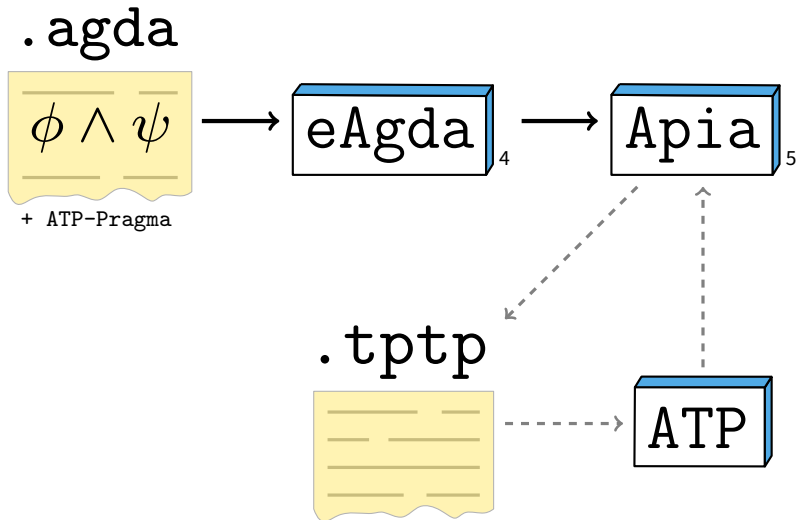
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<sup>4</sup>Development version of Agda in order to handle a new built-in ATP-pragma. <https://github.com/asr/eagda>

## Related Work: Apia

*Proving first-order theorems written in Agda using automatic theorem provers for first-order logic*



<sup>4</sup>Development version of Agda in order to handle a new built-in ATP-pragma. <https://github.com/asr/eagda>

<sup>5</sup>Haskell program for proving first-order theorems written in Agda using ATPs. <https://github.com/asr/apia>



.tptp



- ▶ Is a language<sup>6</sup> to encode problems
- ▶ Is the input of the ATPs
- ▶ TPTP format contains formulas with the form

```
language(name, role, formula).
```

**language** FOF or CNF

**name** to identify the formula within the problem

**role** axiom, definition, hypothesis, conjecture, among others

**formula** version in TPTTP format

<sup>6</sup>Is available at <http://www.cs.miami.edu/~tptp/TPTP/SyntaxBNF.html>

## Problems in Propositional Logic:

►  $p \vdash p$

```
$ cat basic-4.tptp
fof(a, axiom, p).
fof(goal, conjecture, p).
```

►  $p \wedge q \vdash q \wedge p$

```
$ cat conj-3.tptp
fof(a, axiom, p & q).
fof(goal, conjecture, q & p).
```

►  $\vdash \neg(p \wedge \neg p) \vee (q \wedge \neg q)$

```
$ cat neg-7.tptp
fof(goal, conjecture, ~ ((p & ~ p) | (q & ~ q))).
```



**.tstp**

A TSTP derivation<sup>8</sup>

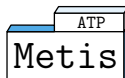
- ▶ Is a Directed Acyclic Graph where
  - leaf** is a formulae from the TPTP input
  - node** is a formulae inferred from parent formulae
  - root** the final derived formulae
- ▶ Is a list of annotated formulae:

```
language(name, role, formula, source [,useful info]).
```

where **source** typically is a inference record:

```
inference(rule, useful information, parents)
```

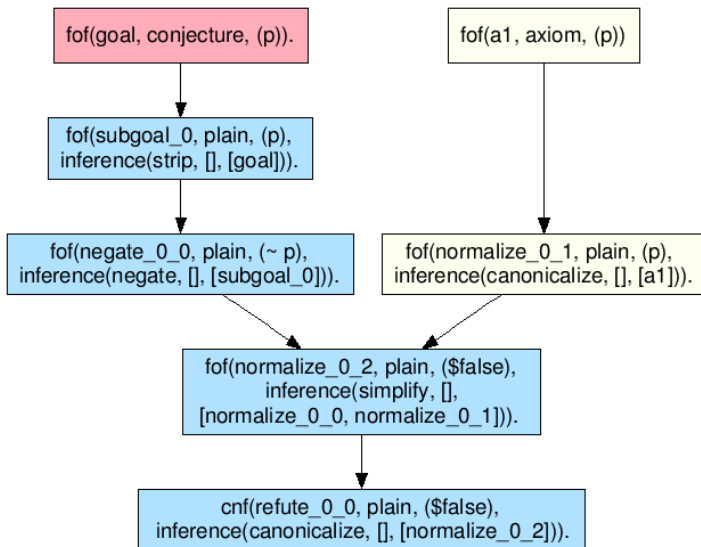
<sup>8</sup><http://www.cs.miami.edu/~tptp/TPTP/QuickGuide/Derivations.html>



- Proof found by **Metis** ATP for the problem  $p \vdash p$

```
$ metis --show proof basic-4.tptp
fof(a, axiom, (p)).
fof(goal, conjecture, (p)).
fof(subgoal_0, plain, (p),
    inference(strip, [], [goal])).
fof(negate_0_0, plain, (~ p),
    inference(negate, [], [subgoal_0])).
fof(normalize_0_0, plain, (~ p),
    inference(canonicalize, [], [negate_0_0])).
fof(normalize_0_1, plain, (p),
    inference(canonicalize, [], [a])).
fof(normalize_0_2, plain, ($false),
    inference(simplify, [],
        [normalize_0_0, normalize_0_1])).
cnf(refute_0_0, plain, ($false),
    inference(canonicalize, [], [normalize_0_2])).
```

## DAG for the previous TSTP derivation found by Meti's ATP

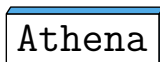


## Athena

Is a Haskell program that translates proofs given by Metis Prover in TSTP format to Agda code.

It depends on:

- ▶ **agda-prop** Classical Logic within Agda: Axioms + Theorems
- ▶ **agda-metis** Theorems of the inference rules of Metis Prover



## Haskell

- ▶ Parsing
- ▶ AST construction
- ▶ Creation and analysis of DAG derivations
- ▶ Analysis of inference rules used
- ▶ Generation of Agda code of the proof

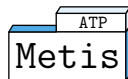


## Agda

Is a dependently typed functional programming language and it also a proof assistant.

We used it to type-check the proofs found by Metis Prover

- ▶ Agda-Prop Library: Logic framework for Classical Propositional Logic
- ▶ Agda-Metis Library: theorems based on the inference rules of Metis Prover



Metis is an automatic theorem prover for first order logic with equality

## Why Metis?

- ▶ Open source implemented in Standard ML
- ▶ Reads problem in **TPTP** format
- ▶ Outputs detailed proofs in **TSTP** format
- ▶ Each refutation step is one of **6 simple rules**

With Classical Propositional Logic, TSTP derivations exhibit these inferences:

- ▶ **Canonicalize** transforms formulas to CNF or NNF
- ▶ **Clausify** performs clausification
- ▶ **Conjunct** extracts a formula from a conjunction
- ▶ **Negate** applies negation to the formula
- ▶ **Resolve** applies theorems of resolution
- ▶ **Simplify** applies over a list of formula to simplify them
- ▶ **Strip** splits a formula into subgoals



canonicalize

clausify

## conjunct

```
$ cat ~/agda-metis/src/ATP/Metis/Rules/Conjunct.agd
```

```
...
```

```
conjunct : Prop → Prop → Prop
```

```
conjunct  $\phi$  (  $\omega \psi$  )  $\omega$  with  $\omega$  eq  $\phi \omega$  |  $\omega$  eq  $\psi \omega$ 
```

```
... | true | _ =  $\phi$ 
```

```
... | false | true =  $\psi$ 
```

```
... | false | false = conjunct  $\phi \omega$ 
```

```
conjunct  $\phi \omega$  =  $\phi$ 
```

```
atp-conjunct
```

```
:  $\omega \Gamma \{ \} \phi \{ \}$ 
```

```
→  $\omega$  ( : Prop )
```

```
→  $\Gamma \omega \phi$ 
```

```
→  $\Gamma \omega$  conjunct  $\phi \omega$ 
```

# Negate

*Inference Rule*

negate

# Resolve

*Inference Rule*

resolve

simplify

strip

Example goes here



Verified Example goes here

Failure Example goes here

- ▶ Add shallow embedding in order to work with Apia
- ▶ Support First-Order Logic
- ▶ Support another prover like EProver
- ▶ Support a Intuitionistic Prover



Sicard-Ramírez, Andrés (2015). *Reasoning about functional programs by combining interactive and automatic proofs*. PEDECIBA Informática, Universidad de la República.