Reconstructing Propositional Proofs in Type Theory

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Research

Goal

Formalization in type theory, classical propositional derivations generated by the `Metis` theorem prover.

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Formalization in type theory, classical propositional derivations generated by the `Metis` theorem prover.

Topics

- Automatic reasoning using automatic theorem provers (ATPs) (e. g., `Metis`, `EProver`)
- Interactive proving using proof-assistants (e. g., `Agda`, `Coq`)
- Proof-reconstruction for proofs generated by ATPs in proof-assistants

Research Outcomes

Academic result: paper (work in progress)
Software related results:

- `Athena`[1]: a translator tool for `Metis` proofs to `Agda` in `Haskell`
- `Agda` libraries:
  - `Agda-Metis`[2]: `Metis` prover reasoning for propositional logic
  - `Agda-Prop`[3]: intuitionistic propositional logic with PEM
- Bugs found in `Metis`: see issues No. 2, No. 4, and commit `8a3f11e` in `Metis` official repository[4]

In parallel, we develop:

- `Online-ATPs`[5]: a client for the TPTP world in `Haskell` This tool allowed us to use `Metis` without installing it
- `Prop-Pack`[6]: compendium of TPTP problems in classical propositional logic used to test `Athena`

[1] https://github.com/jonaprieto/athena.

[2] https://github.com/jonaprieto/agda-metis.

[3] https://github.com/jonaprieto/agda-prop.

[4] https://github.com/gilith/metis.

[5] https://github.com/jonaprieto/online-atps.

[6] https://github.com/jonaprieto/prop-pack.

Bug in the Printing of the Proof

Fixed in `Metis` v2.3 (release 20161108)

$$\varphi := \neg p \wedge (\neg q \Leftrightarrow \neg r) \wedge (\neg p \Leftrightarrow (\neg q \Leftrightarrow \neg r))$$

$$\dfrac{\dfrac{\dfrac{\vdots}{\varphi}\ \text{canonicalize}}{\neg p \Leftrightarrow (\neg q \Leftrightarrow \neg r)}\ \text{conjunct} \qquad \dfrac{\dfrac{\vdots}{\varphi}\ \text{canonicalize}}{\neg q \Leftrightarrow \neg r}\ \text{conjunct} \qquad \dfrac{\dfrac{\vdots}{\varphi}\ \text{canonicalize}}{\neg p}\ \text{conjunct}}{\bot}\ \text{simplify}$$

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$$\varphi := \neg p \wedge (\neg q \Leftrightarrow \neg r) \wedge (\neg p \Leftrightarrow (\neg q \Leftrightarrow \neg r))$$

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The bug was caused by the conversion of `Xor` sets to `Iff` lists. After reporting this, `Metis` developer fixed the printing of `canonicalize` inference rule

$$\varphi := \neg p \wedge (\neg q \Leftrightarrow \neg r) \wedge (\neg p \Leftrightarrow (\neg q \Leftrightarrow {\color{red} r}))$$

Soundness Bug in Splitting goals

Fixed in `Metis` v2.3 (release 20170810)

Consider this TPTP problem

```
$ cat issue.tptp
fof(goal, conjecture, (~ (p <=> q)) <=> ((p => ~ q) & (q => ~p))).
```

`Metis` found a proof when other ATPs do not. Indeed, the problem is not a tautology.

```
$ metis issue.tptp
SZS status Theorem for issue.tptp
```

Testing with `EProver` with a client for `SystemOnTPTP` (`Online-ATPs`).

```
$ online-atps --atp=e issue.tptp
...
# No proof found!
# SZS status CounterSatisfiable
...
```

Soundness Bug in Splitting goals

Fixed in `Metis` v2.3 (release 20170810)

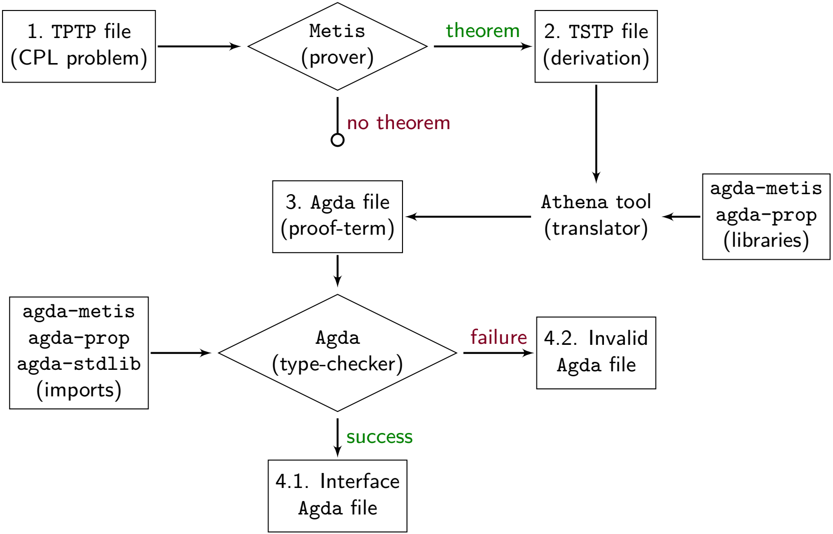
The bug was in the `strip` inference rule:

$$\neg\,(p \Leftrightarrow q) \Leftrightarrow ((p \Rightarrow \neg\, q) \wedge (q \Rightarrow \neg\, p))$$

Solved with:

$$\neg\,(p \Leftrightarrow q) \Leftrightarrow ((p \Rightarrow \neg\, q) \wedge (\neg\, q \Rightarrow p))$$

Proof Reconstruction: Overview



Inference Rules of `Metis`

`TSTP` derivations by `Metis` exhibit the following inferences:

`Metis` rule	Purpose
`strip`	Strip a goal into subgoals
`conjunct`	Takes a formula from a conjunction
`resolve`	A general form of the resolution theorem
`canonicalize`	Normalization of the formula
`clausify`	Performs clausification
`simplify`	Simplify definitions and theorems

Proposition Type in Agda

A data type for formulas

```
data PropFormula : Set where
  Var : Fin n → Prop
  ⊤   : Prop
  ⊥   : Prop
  _∧_ : (φ ψ : Prop) → Prop
  _∨_ : (φ ψ : Prop) → Prop
  _⇒_ : (φ ψ : Prop) → Prop
  _⇔_ : (φ ψ : Prop) → Prop
  ¬_  : (φ : Prop)   → Prop
```

Inference Rules For Propositional Logic I

Intuitionistic Propositional Logic + PEM ($\Gamma \vdash \varphi \vee \neg\varphi$)

$$\frac{}{\Gamma, \varphi \vdash \varphi}\ \textsf{assume} \qquad \frac{}{\Gamma \vdash \top}\ \top\textsf{-intro} \qquad \frac{}{\Gamma \vdash \varphi \vee \neg\, \varphi}\ \textsf{PEM}$$

$$\frac{\Gamma \vdash \bot}{\Gamma \vdash \varphi}\ \bot\textsf{-elim} \qquad \frac{\Gamma, \varphi \vdash \bot}{\Gamma \vdash \neg\, \varphi}\ \neg\textsf{-intro} \qquad \frac{\Gamma \vdash \neg\, \varphi \qquad \Gamma \vdash \varphi}{\Gamma \vdash \bot}\ \neg\textsf{-elim}$$

$$\frac{\Gamma \vdash \varphi \qquad \Gamma \vdash \psi}{\Gamma \vdash \varphi \wedge \psi}\ \wedge\textsf{-intro} \qquad \frac{\Gamma \vdash \varphi \wedge \psi}{\Gamma \vdash \varphi}\ \wedge\textsf{-proj}_1 \qquad \frac{\Gamma \vdash \varphi \wedge \psi}{\Gamma \vdash \psi}\ \wedge\textsf{-proj}_2$$

Inference Rules For Propositional Logic II

$$\frac{\Gamma \vdash \varphi}{\Gamma \vdash \varphi \vee \psi}\ \vee\text{-intro}_1 \qquad \frac{\Gamma \vdash \psi}{\Gamma \vdash \varphi \vee \psi}\ \vee\text{-intro}_2$$

$$\frac{\Gamma, \varphi \vdash \gamma \qquad \Gamma, \psi \vdash \gamma}{\Gamma, \varphi \vee \psi \vdash \gamma}\ \vee\text{-elim}$$

$$\frac{\Gamma, \varphi \vdash \psi}{\Gamma \vdash \varphi \Rightarrow \psi}\ \Rightarrow\text{-intro} \qquad \frac{\Gamma \vdash \varphi \Rightarrow \psi \qquad \Gamma \vdash \varphi}{\Gamma \vdash \psi}\ \Rightarrow\text{-elim}$$

Other Rules

- Weakening: to extend the hypotheses with additional formulas

$$\frac{\Gamma \vdash \varphi}{\Gamma, \psi \vdash \varphi}\ \text{weaken}$$

- The RAA rule is the formulation of the principle of proof by contradiction:

$$\frac{\Gamma, \neg\, \varphi \vdash \bot}{\Gamma \vdash \varphi}\ \text{RAA}$$

Syntactical Consequence Relation in Agda

- Inductive family $\_ \vdash \_$ with two indexes: a set of propositions Γ (the premises) and a proposition φ (the conclusion)

Example

In [8] we define $\_ \vdash \_$ as follows

```
data _⊢_ : (Γ : Ctxt)(φ : Prop) → Set
  ...
  ∧-intro
    : ∀ {Γ} {φ ψ}
    → Γ ⊢ φ → Γ ⊢ ψ
    → Γ ⊢ φ ∧ ψ

  ∧-proj₁
    : ∀ {Γ} {φ ψ}
    → Γ ⊢ φ ∧ ψ
    → Γ ⊢ φ

  ∧-proj₂
    : ∀ {Γ} {φ ψ}
    → Γ ⊢ φ ∧ ψ
    → Γ ⊢ ψ
  ...
```

Reconstructing Metis Rules in Type Theory

Let $\mathrm{metisRule}$ be a `Metis` inference rule. We define the function metisRule in type theory which has the following pattern[7]:

$$\text{metisRule} : \textsc{Premise} \to \textsc{Conclusion} \to \textsc{Prop}$$

$$\text{metisRule } \varphi \ \psi \quad = \begin{cases} \psi, & \text{if metisRule built } \psi \text{ by applying inference} \\ & \text{rules to } \varphi; \\ \varphi, & \text{otherwise;} \end{cases}$$

To justify all transformations done by the metisRule rule, we prove its soundness with a theorem like the following:

$$\text{If } \Gamma \vdash \varphi \text{ then } \Gamma \vdash \text{metisRule } \varphi \ \psi, \text{ where } \psi : \textsc{Conclusion}.$$

[7] Premise and Conclusion as synonyms of the Prop type to describe in the function types the role of the arguments

Reconstructing a `Metis` Inference Rule

The `clausify` rule transforms a formula into its clausal normal form.

Example

In the following `TSTP` derivation by `Metis`, we see how `clausify` transforms the $\texttt{norm}_0$ formula to get $\texttt{norm}_1$ formula.

$\texttt{fof(norm}_0\texttt{, } \neg\ \texttt{p} \vee \texttt{(q} \wedge \texttt{r) ...}$
$\texttt{fof(norm}_1\texttt{, (}\neg\ \texttt{p} \vee \texttt{q)} \wedge \texttt{(}\neg\ \texttt{p} \vee \texttt{r), inf(clausify, norm}_0\texttt{)).}$

Theorem

Let $\psi : \textsc{Conclusion}$. *If* $\Gamma \vdash \varphi$ *then* $\Gamma \vdash \text{clausify}\ \varphi\ \psi$, *where*

$$\text{clausify} : \textsc{Premise} \to \textsc{Conclusion} \to \textsc{Prop}$$

$$\text{clausify}\ \varphi\ \psi \quad = \begin{cases} \psi, & \textit{if } \varphi \equiv \psi; \\ \text{reorder}_{\wedge\vee}\ (\text{cnf}\ \varphi)\ \psi, & \textit{otherwise.} \end{cases}$$

The Intuition behind the Metis Algorithm

Algorithm 1 `Metis` refutation strategy

```
procedure METIS
input: the goal and a set of premises a_1, ⋯, a_n
output: maybe a derivation when a_1, ⋯, a_n ⊢ goal, otherwise nothing.

    strip the goal into a list of subgoals s_i
    for each subgoal s_i do
        try to find by a refutation for ¬s_i:
          apply clausification for the negated subgoal ¬ s_i
        if a premise a_j is relevant then
            apply clausification to a_j
        end if
          application of Metis inference rules
        if a contradiction can be derived from the assumptions then
            keep the refutation and continue with the others subgoals
        else
            exit without a proof.
        end if
    end for
    print the conjecture and the premises
    print each refutation for each negated subgoal
end procedure
```

Some Challenges

- Formalization
  - Understanding the `Metis` reasoning without a proper documentation or description from the `Metis` author
  - Terminating of functions that reconstruct `Metis` inference rules
  - Intuitionistic logic implementation
- Software related
  - Parsing of `TSTP` derivations
  - Printing valid `Agda` files

Complete Example

The problem[8]:

$$(p \Rightarrow q) \wedge (q \Rightarrow p) \vdash (p \vee q) \Rightarrow (p \wedge q)$$

In TPTP syntax:

```
fof(a₁, axiom, (p ⇒ q) ∧ (q ⇒ p)).
fof(goal, conjecture, (p ∨ q) ⇒ (p ∧ q)).
```

Its TSTP solution using Metis:

```
fof(a₁, axiom, (p ⇒ q) ∧ (q ⇒ p)).
fof(goal, conjecture, (p ∨ q) ⇒ (p ∧ q))).
fof(s₁, (p ∨ q) ⇒ p, inf(strip, goal)).
fof(s₂, ((p ∨ q) ∧ p) ⇒ q, inf(strip, goal)).
...
```

[8] Problem No. 13 in Disjunction Section in [7]

fof(s_1, $(p \vee q) \Rightarrow p$, inf(strip, goal)).
fof(s_2, $((p \vee q) \wedge p) \Rightarrow q$, inf(strip, goal)).
fof(neg_1, $\neg\ ((p \vee q) \Rightarrow p)$, inf(negate, s_1)).
fof(n00, $(\neg\ p \vee q) \wedge (\neg\ q \vee p)$, inf(canonicalize, a_1)).
fof(n01, $\neg\ q \vee p$, inf(conjunct, n00)).
fof(n02, $\neg\ p \wedge (p \vee q)$, inf(canonicalize, neg_1)).
fof(n03, $p \vee q$, inf(conjunct, n02)).
fof(n04, $\neg\ p$, inf(conjunct, n02)).
fof(n05, q, inf(simplify,[n03, n04])).
cnf(r00, $\neg\ q \vee p$, inf(canonicalize, n01)).
cnf(r01, q, inf(canonicalize, n05)).
cnf(r02, p, inf(resolve, q, [r01, r00])).
cnf(r03, $\neg\ p$, inf(canonicalize, n04)).
cnf(r04, $\perp$, inf(resolve, p, [r02, r03])).
fof(neg_2, $\neg\ (((p \vee q) \wedge p) \Rightarrow q)$, inf(negate, s2)).
fof(n10, $\neg\ q \wedge p \wedge (p \vee q)$, inf(canonicalize, neg_2)).
fof(n11, $(\neg\ p \vee q) \wedge (\neg\ q \vee p)$, inf(canonicalize, a_1)).
fof(n12, $\neg\ p \vee q$, inf(conjunct, n11)).
fof(n13, $\perp$, inf(simplify,[n10, n12])).
cnf(r10, $\perp$, inf(canonicalize, n13)).

TSTP Refutation of Subgoal No. 1

$$
\begin{array}{l}
\mathtt{fof(s_1,\ (p} \lor \mathtt{q)} \Rightarrow \mathtt{p,\ inf(strip,\ goal)).}\\
\mathtt{fof(neg_1,\ } \neg \mathtt{\ ((p} \lor \mathtt{q)} \Rightarrow \mathtt{p),\ inf(negate,\ s_1)).}\\
\mathtt{fof(n00,\ (}\neg\mathtt{\ p} \lor \mathtt{q)} \land \mathtt{(}\neg\mathtt{\ q} \lor \mathtt{p),\ inf(canonicalize,\ a_1)).}\\
\mathtt{fof(n01,\ } \neg \mathtt{\ q} \lor \mathtt{p,\ inf(conjunct,\ n00)).}\\
\mathtt{fof(n02,\ } \neg \mathtt{\ p} \land \mathtt{(p} \lor \mathtt{q),\ inf(canonicalize,\ neg_1)).}\\
\mathtt{fof(n03,\ p} \lor \mathtt{q,\ inf(conjunct,\ n02)).}\\
\mathtt{fof(n04,\ } \neg \mathtt{\ p,\ inf(conjunct,\ n02)).}\\
\mathtt{fof(n05,\ q,\ inf(simplify,[n03,\ n04])).}\\
\mathtt{cnf(r00,\ } \neg \mathtt{\ q} \lor \mathtt{p,\ inf(canonicalize,\ n01)).}\\
\mathtt{cnf(r01,\ q,\ inf(canonicalize,\ n05)).}\\
\mathtt{cnf(r02,\ p,\ inf(resolve,\ q,\ [r01,\ r00])).}\\
\mathtt{cnf(r03,\ } \neg \mathtt{\ p,\ inf(canonicalize,\ n04)).}\\
\mathtt{cnf(r04,\ } \bot \mathtt{,\ inf(resolve,\ p,\ [r02,\ r03])).}
\end{array}
$$

Tree for the Subgoal No. 1: $(p \vee q) \Rightarrow p$

```
fof(a₁, axiom, (p ⇒ q) ∧ (q ⇒ p)).
...
fof(n00, (¬ p ∨ q) ∧ (¬ q ∨ p), inf(canonicalize, a₁)).
fof(n01, ¬ q ∨ p, inf(conjunct, n00)).
...
```

$$(\mathcal{D}_1) \qquad \cfrac{\cfrac{\cfrac{\cfrac{}{\Gamma \vdash (p \Rightarrow q) \wedge (q \Rightarrow p)}\ \text{axiom } a_1}{\Gamma, \neg s_1 \vdash (p \Rightarrow q) \wedge (q \Rightarrow p)}\ \text{weaken}}{\Gamma, \neg s_1 \vdash (\neg p \vee q) \wedge (\neg q \vee p)}\ \text{canonicalize}}{\Gamma, \neg s_1 \vdash \neg q \vee p}\ \text{conjunct}$$

```
...
fof(s₁, (p ∨ q) ⇒ p, inf(strip, goal)).
fof(neg₁, ¬ ((p ∨ q) ⇒ p), inf(negate, s₁)).
...
fof(n02, ¬ p ∧ (p ∨ q), inf(canonicalize, neg₁)).
fof(n03, p ∨ q, inf(conjunct, n02)).
fof(n04, ¬ p, inf(conjunct, n02)).
...
```

$$(\mathcal{D}_2) \qquad \cfrac{\cfrac{\cfrac{}{\Gamma, \neg s_1 \vdash \neg s_1}\ \text{assume}}{\Gamma, \neg s_1 \vdash \neg p \wedge (p \vee q)}\ \text{canonicalize}}{\Gamma, \neg s_1 \vdash p \vee q}\ \text{conjunct}$$

$$(\mathcal{D}_3) \qquad \cfrac{\cfrac{\cfrac{}{\Gamma, \neg s_1 \vdash \neg s_1}\ \text{assume } \neg s_1}{\Gamma, \neg s_1 \vdash \neg p \wedge (p \vee q)}\ \text{canonicalize}}{\Gamma, \neg s_1 \vdash \neg p}\ \text{conjunct}$$

$$(\mathcal{D}_4) \qquad \dfrac{\dfrac{\mathcal{D}_2}{\Gamma, \neg s_1 \vdash p \vee q} \qquad \dfrac{\mathcal{D}_3}{\Gamma, \neg s_1 \vdash \neg p}}{\Gamma, \neg s_1 \vdash q}\;\textsf{simplify}$$

$$(\mathcal{R}_1) \qquad \dfrac{\dfrac{\dfrac{\dfrac{\mathcal{D}_1}{\Gamma, \neg s_1 \vdash \neg q \vee p} \qquad \dfrac{\mathcal{D}_4}{\Gamma, \neg s_1 \vdash q}}{\Gamma, \neg s_1 \vdash p}\;\textsf{resolve } q \qquad \dfrac{\mathcal{D}_3}{\Gamma, \neg s_1 \vdash \neg p}}{\Gamma, \neg s_1 \vdash \bot}\;\textsf{resolve } p}{\Gamma \vdash s_1}\;\textsf{RAA}$$

Tree for the Subgoal No. 2: $((p \vee q) \wedge p) \Rightarrow q$

```
fof(s₂, ((p ∨ q) ∧ p) ⇒ q, inf(strip, goal)).
fof(neg₂, ¬ (((p ∨ q) ∧ p) ⇒ q), inf(negate, s2)).
fof(n10, ¬ q ∧ p ∧ (p ∨ q), inf(canonicalize, neg₂)).
fof(n11, (¬ p ∨ q) ∧ (¬ q ∨ p), inf(canonicalize, a₁)).
fof(n12, ¬ p ∨ q, inf(conjunct, n11)).
fof(n13, ⊥, inf(simplify,[n10, n12])).
cnf(r10, ⊥, inf(canonicalize, n13)).
```

$$(\mathcal{R}_2) \qquad \cfrac{\cfrac{\cfrac{\cfrac{}{\Gamma, \neg s_2 \vdash \neg s_2}\;\text{assume } (\neg s_2)}{\Gamma, \neg s_2 \vdash \neg q \wedge p \wedge (p \vee q)}\;\text{canonicalize} \qquad \cfrac{\cfrac{\cfrac{\cfrac{}{\Gamma \vdash (p \Rightarrow q) \wedge (q \Rightarrow p)}\;\text{axiom } a_1}{\Gamma, \neg s_2 \vdash (p \Rightarrow q) \wedge (q \Rightarrow p)}\;\text{weaken}}{\Gamma, \neg s_2 \vdash (\neg p \vee q) \wedge (\neg q \vee p)}\;\text{canonicalize}}{\Gamma, \neg s_2 \vdash \neg p \vee q}\;\text{conjunct}}{\Gamma, \neg s_2 \vdash \bot}\;\text{simplify}}{\Gamma \vdash s_2}\;\text{RAA}$$

Summarizing the Example

The problem was:

$$(p \Rightarrow q) \wedge (q \Rightarrow p) \vdash (p \vee q) \Rightarrow (p \wedge q)$$

Its TSTP solution using Metis was:

$$\begin{array}{l}
\texttt{fof(a}_1\texttt{, axiom, (p} \Rightarrow \texttt{q)} \wedge \texttt{(q} \Rightarrow \texttt{p)).} \\
\texttt{fof(goal, conjecture, (p} \vee \texttt{q)} \Rightarrow \texttt{(p} \wedge \texttt{q))).} \\
\texttt{fof(s}_1\texttt{, (p} \vee \texttt{q)} \Rightarrow \texttt{p, inf(strip, goal)).} \\
\texttt{fof(s}_2\texttt{, ((p} \vee \texttt{q)} \wedge \texttt{p)} \Rightarrow \texttt{q, inf(strip, goal)).} \\
\dots
\end{array}$$

The proof is:

$$\dfrac{\dfrac{}{\Gamma \vdash (s_1 \wedge s_2) \Rightarrow \text{goal}}\,\text{strip} \qquad \dfrac{\dfrac{\mathcal{R}_1}{\Gamma \vdash s_1} \quad \dfrac{\mathcal{R}_2}{\Gamma \vdash s_2}}{\Gamma \vdash s_1 \wedge s_2}\,\wedge\text{-intro}}{\Gamma \vdash \text{goal}}\,\Rightarrow\text{-elim}$$

(Live example using Agda and Athena)

Future Work

Further research directions include, but are not limited to:

- improve the performance of the `canonicalize` rule
- extend the proof-reconstruction presented in this paper to
  - support the proposition logic with equality of `Metis`
  - support other ATPs for propositional logic like `EProver` or `Z3`. See Kanso's Ph.D. thesis [5]
  - support `Metis` first-order proofs

Related Work

In type theory:

- Kanso in [5] reconstructs in Agda propositional proofs generated by EProver and Z3
- Foster and Struth in [2] describe proof-reconstruction in Agda for equational logic of Waldmeister prover
- Bezem, Hendriks, and Nivelle in [1] transform a proof produced by the first-order prover Bliksem in a Coq proof-term

In classical logic:

- Paulson and Susanto in [6] introduce SledgeHammer, a tool ables to reconstructs proofs of well-known ATPs: EProver, Vampire, among others using SystemOnTPTP server
- Hurd in [3] integrates the first-order resolution prover Gandalf prover for HOL proof-assistant
- Kaliszyk and Urban in [4] reconstruct proofs of different ATPs for HOL Light

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BONUS SLIDES

TPTP Syntax

Thousands of Problems for Theorem Provers

- Is a language[9] to encode problems
- Is the input of the ATPs
- Annotated formulas with the form

  `language(name, role, formula).`

`language` FOF or CNF
`name` to identify the formula within the problem
`role` axiom, definition, hypothesis, conjecture
`formula` formula in TPTP format

[9] http://www.cs.miami.edu/~tptp/TPTP/SyntaxBNF.html.

`Metis` Theorem Prover

`Metis` is an automatic theorem prover for first-order logic with equality.

- Open source implemented
- Reads problems in `TPTP` format
- Outputs *detailed* proofs in `TSTP` format
- For the propositional logic, `Metis` has only three inference rules:

$$\frac{}{\Gamma \vdash \varphi_1 \vee \cdots \vee \varphi_n} \text{ axiom } \varphi_1, \cdots, \varphi_n$$

$$\frac{}{\Gamma \vdash \varphi \vee \neg\varphi} \text{ assume } \varphi$$

$$\frac{\Gamma \vdash \varphi_1 \vee \cdots \vee l \vee \cdots \vee \varphi_n \qquad \Gamma \vdash \psi_1 \vee \cdots \vee \neg\, l \vee \cdots \vee \psi_m}{\Gamma \vdash \varphi_1 \vee \cdots \vee \varphi_n \vee \psi_1 \vee \cdots \vee \psi_m} \text{ resolve } l$$

TSTP Syntax

A TSTP derivation[10]

- Is a **D**irected **A**cyclic **G**raph where

  `leaf` is a formula from the TPTP input
  `node` is a formula inferred from parent formula
  `root` the final derived formula

- Is a list of annotated formulas with the form

```
language(name, role, formula, source [,useful info]).
```

where `source` typically is an inference record

```
inference(rule, useful info, parents).
```

[10] http://www.cs.miami.edu/~tptp/TPTP/QuickGuide/Derivations.html.

Another TSTP Example

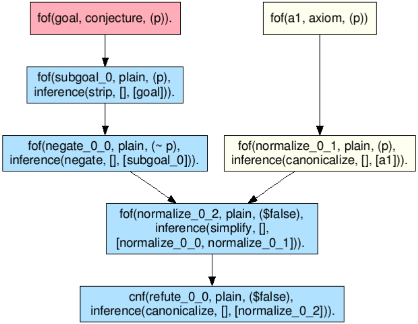
- Proof found by Metis for the problem $p \vdash p$

```
$ metis --show proof problem.tptp
fof(a, axiom, p).
fof(goal, conjecture, p).
fof(subgoal_0, plain, p),
  inference(strip, [], [goal])).
fof(negate_0_0, plain, ~ p,
  inference(negate, [], [subgoal_0])).
fof(normalize_0_0, plain, ~ p,
  inference(canonicalize, [], [negate_0_0])).
fof(normalize_0_1, plain, p,
  inference(canonicalize, [], [a])).
fof(normalize_0_2, plain, $false,
  inference(simplify, [],
    [normalize_0_0, normalize_0_1])).
cnf(refute_0_0, plain, $false,
    inference(canonicalize, [], [normalize_0_2])).
```

DAG Example

By refutation, we proved $p \vdash p$:

$$\dfrac{\dfrac{\dfrac{\dfrac{}{\neg p}\,\text{assume}}{\neg p}\,\text{strip} \qquad \dfrac{\dfrac{}{p}\,\text{axiom}}{p}\,\text{canonicalize}}{\bot}\,\text{simplify}}{\bot}\,\text{canonicalize}$$



Athena tool

Is a Haskell program that translates proofs given by Metis in TSTP format to Agda code

- Parsing of TSTP language
- Creation and analysis of **DAG** derivations
- Analysis of inference rules used in the TSTP derivation
- Agda code generation

Library	Purpose
Agda-Prop	axioms and theorems of classical propositional logic
Agda-Metis	versions of the inference rules used by Metis

Agda-Metis: Conjunct Inference [11]

- Definition

$$\mathsf{conjunct}(\overbrace{\varphi_1 \wedge \cdots \wedge \varphi_n}^{\varphi}, \psi) = \begin{cases} \varphi_i & \text{if } \psi \text{ is equal to some } \varphi_i \\ \varphi & \text{otherwise} \end{cases}$$

[11] `https://github.com/jonaprieto/agda-metis`.

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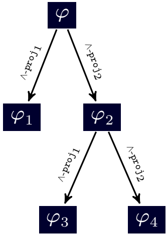
- Inference rules involved

$$\frac{\varphi_1 \wedge \varphi_2}{\varphi_1}\ \wedge\text{-}\mathtt{proj}_1 \qquad\qquad \frac{\varphi_1 \wedge \varphi_2}{\varphi_2}\ \wedge\text{-}\mathtt{proj}_2$$

- Example

$$\varphi := \varphi_1 \wedge \overbrace{(\varphi_3 \wedge \varphi_4)}^{\varphi_2}$$

  - conjunct $(\varphi, \varphi_3 \wedge \varphi_1) \equiv \varphi$
  - conjunct $(\varphi, \varphi_3) \equiv \varphi_3$
  - conjunct $(\varphi, \varphi_2) \equiv \varphi_2$



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