Proof Reconstruction in Classical Propositional Logic

(Work in Progress)

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Agda Implementors' Meeting XXV May 9-15th

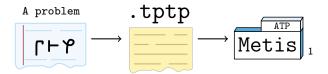




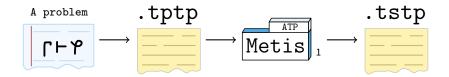
A problem



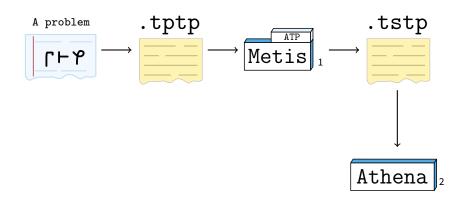




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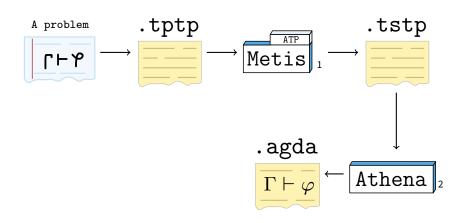


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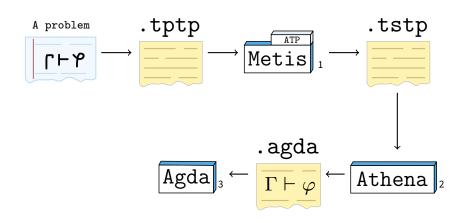
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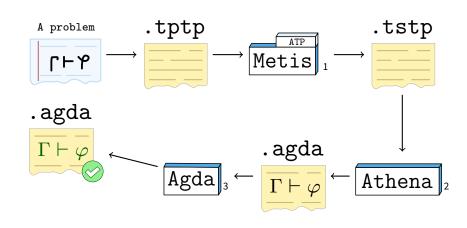
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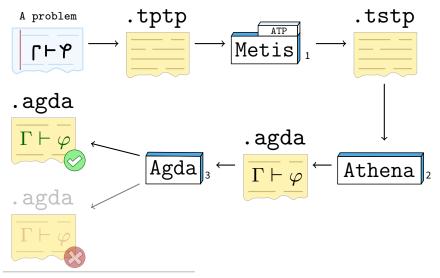
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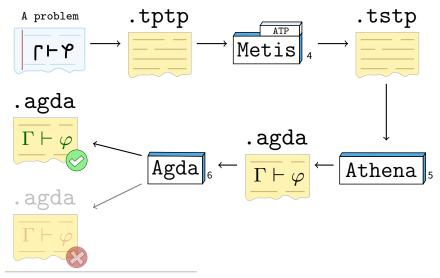
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Bonus Slides



- ▶ Is a language⁷ to encode problems (Sutcliffe, 2009)
- ▶ Is the input of the ATPs
- Annotated formulas with the form

```
language(name, role, formula).
```

```
language FOF or CNF
```

name to identify the formula within the problem

role axiom, definition, hypothesis, conjecture, among others

formula version in TPTTP format

⁷Is available at http://www.cs.miami.edu/~tptp/TPTP/SyntaxBNF.html

 $\triangleright p \vdash p$

fof(a, axiom, p).
fof(goal, conjecture, p).

ightharpoonup $\vdash \neg (p \land \neg p) \lor (q \land \neg q)$

fof(goal, conjecture, ~ ((p & ~ p) | (q & ~ q))).

⁸ Is available at http://github.com/jonaprieto/prop-pack

.tstp

A TSTP derivation 9

- Is a Directed Acyclic Graph where leaf is a formula from the TPTP input **node** is a formula inferred from parent formula root the final derived formula
- Is a list of annotated formulas with the form:

```
language(name, role, formula, source [,useful info]).
```

where **source** typically is an inference record:

```
inference(rule, useful info, parents)
```

⁹http://www.cs.miami.edu/~tptp/TPTP/QuickGuide/Derivations.html

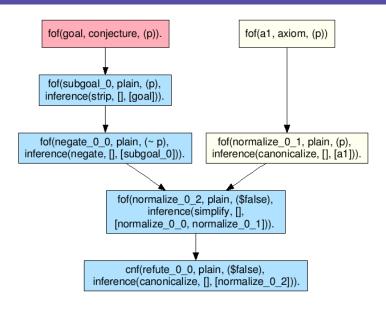
Metis

▶ Proof found by **Metis** ATP for the problem $p \vdash p$

```
$ metis --show proof basic-4.tptp
fof(a, axiom, (p)).
fof(goal, conjecture, (p)).
fof(subgoal_0, plain, (p),
  inference(strip, [], [goal])).
fof(negate_0_0, plain, (~ p),
  inference(negate, [], [subgoal_0])).
fof(normalize_0_0, plain, (~ p),
  inference(canonicalize, [], [negate_0_0])).
fof(normalize_0_1, plain, (p),
  inference(canonicalize, [], [a])).
fof(normalize_0_2, plain, ($false),
  inference(simplify, [],
    [normalize_0_0, normalize_0_1])).
cnf(refute_0_0, plain, ($false),
    inference(canonicalize, [], [normalize_0_2])).
```

Go Bacl

DAG for the previous TSTP derivation found by Metis ATP



Go Back

Athena

Athena

Is a Haskell program that translates proofs given by Metis Prover in TSTP format to Agda code.

It depends on:

- agda-prop Classical Logic within Agda: Axioms + Theorems
- agda-metis Theorems of the inference rules of Metis Prover

Design Decisions for the Reconstruction Tool

Athena

Haskell

- Parsing
- AST construction
- Creation and analysis of DAG derivations
- Analysis of inference rules used
- Generation of Agda code of the proof



Agda

Is a dependently typed functional programming language and it also a proof assistant.

We used it to type-check the proofs found by Metis Prover

- Agda-Prop Libary: Logic framework for Classical Propositional Logic
- Agda-Metis Library: theorems based on the inference rules of Metis Prover

Metis Theorem Prover

http://www.gilith.com/software/metis/



Metis is an automatic theorem prover for First-Order Logic with equality

Why Metis?

- Open source implemented in Standard ML
- Reads problem in TPTP format
- Outputs detailed proofs in TSTP format
- Each refutation step is one of 6 simple rules

TSTP derivations exhibit these inferences:

Task
transforms formulas to CNF, DNF or NNF
performs clausification
extracts a formula from a conjunction
applies negation to the formula
applies theorems of resolution
applies over a list of formula to simplify them
splits a formula into subgoals

Definition

$$conjunct(\phi_1 \land \phi_2 \land \cdots \land \phi_i \land \cdots \land \phi, \phi_i) \longrightarrow \phi_i$$

► Function¹⁰:

```
conjunct: Prop \rightarrow Prop conjunct (\varphi \land \psi) \omega with [eq \varphi \omega] | [eq \psi \omega]
... | true | _ = \varphi
... | false | true = \psi
... | false | false = conjunct \varphi \omega
conjunct \varphi \omega = \varphi
```

▶ Theorem¹⁰

```
atp-conjunct

: \forall \{\Gamma\} \{\varphi\}

\Rightarrow (\omega : Prop)

\Rightarrow \Gamma \vdash \varphi

\Rightarrow \Gamma \vdash \text{conjunct } \varphi \omega
```

 $^{^{10}}$ lt is available in ${\tt agda-metis}$ library in ATP.Metis.Rules.Conjunct module.

- ▶ The problem is $p \land q \vdash q \land p$
- TPTP derivation

```
fof(a, axiom, p & q).
fof(goal, conjecture, q & p).
```

A natural deduction proof

$$\frac{\frac{\phi \land \psi}{\phi} \land \text{-proj}_1 \quad \frac{\phi \land \psi}{\psi} \land \text{-proj}_2}{\psi \land \phi} \land \text{-intro}$$



```
proof_0 : \Gamma \vdash subgoal_0
proof_0 =
   (RAA
     (atp-canonicalize
        (atp-simplify
           (atp-canonicalize
             (atp-strip
                (assume \{\Gamma = \Gamma\}
                   (atp-negate subgoal<sub>0</sub>))))
           (atp-conjunct (q)
              (atp-canonicalize
                (weaken (atp-negate subgoal<sub>0</sub>)
                   (assume \{\Gamma = \emptyset\} a)))))))
```

```
proof_1 : \Gamma \vdash subgoal_1
proof<sub>1</sub> =
  (RAA
     (atp-canonicalize
        (atp-simplify
          (atp-conjunct (q)
             (atp-canonicalize
               (weaken (atp-negate subgoal,)
                  (assume \{\Gamma = \emptyset\} a))))
          (atp-simplify
             (atp-canonicalize
               (atp-strip
                  (assume \{\Gamma = \Gamma\}
                    (atp-negate subgoal,))))
             (atp-conjunct (p)
               (atp-canonicalize
                  (weaken (atp-negate subgoal,)
                    (assume \{\Gamma = \emptyset\} a))))))))
```

Type-checked Proof

Failure Example goes here

SledgeHammer

- Isabelle/HOL tool
- ► Metis ported within Isabelle
- Reconstruct proofs of well-known ATPs: EProver, Vampire, among others

Integrating Waldmeister and Agda

- Source code not available
- Equational Logic
- Reflection Layers

At the moment, the communication between Agda and the ATPs is unidirectional because the ATPs are being used as oracles (Sicard-Ramírez, 2015).

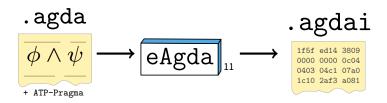


```
module Or where

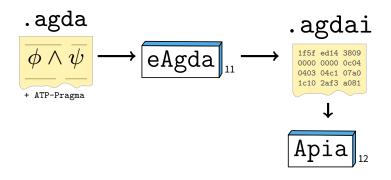
data _V_ (A B : Set) : Set where
  inj<sub>1</sub> : A → A ∨ B
  inj<sub>2</sub> : B → A ∨ B

postulate
  A B : Set
  ∨-comm : A ∨ B → B ∨ A
{-# ATP prove ∨-comm #-}
```

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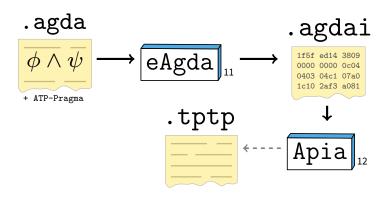


 $^{^{11} \}hbox{Development version of Agda in order to handle a new built-in ATP-pragma. } \textbf{https://github.com/asr/eagda}$



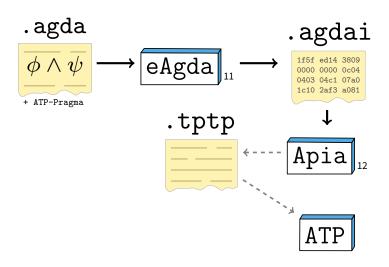
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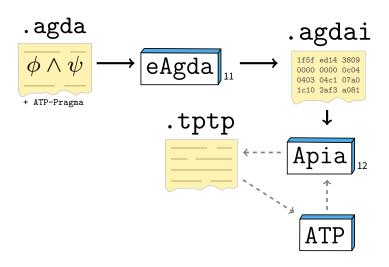
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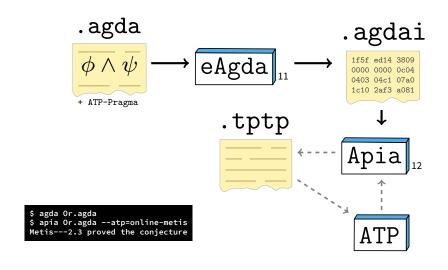
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Future Work

- Add shallow embedding in order to work with Apia
- Support First-Order Logic with Equality
- Support another prover like EProver

References



Sicard-Ramírez, Andrés (2015). Reasoning about functional programs by combining interactive and automatic proofs. PEDECIBA Informática, Universidad de la República.



Sutcliffe, G. (2009). "The TPTP Problem Library and Associated Infrastructure: The FOF and CNF Parts, v3.5.0". In: *Journal of Automated Reasoning* 43.4, pp. 337–362.