Proof Reconstruction in Classical Propositional Logic

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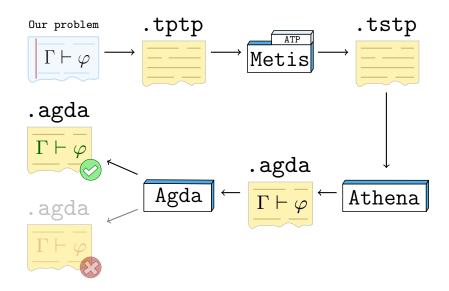
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Outline

Introduction

Motivation
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Proof Reconstruction



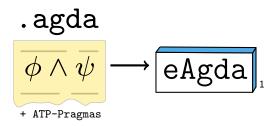
At the moment, the communication between Agda and the ATPs is unidirectional because the ATPs are being used as oracles (Sicard-Ramírez, 2015).



+ ATP-Pragmas

```
$ cat Or.agda
module Or where
data _or_ (A B : Set) : Set where
  inj1 : A \rightarrow A \text{ or } B
  inj2 : B \rightarrow A \text{ or } B
postulate
  AB : Set
  or-comm : A or B -> B or A
{-# ATP prove or-comm #-}
```

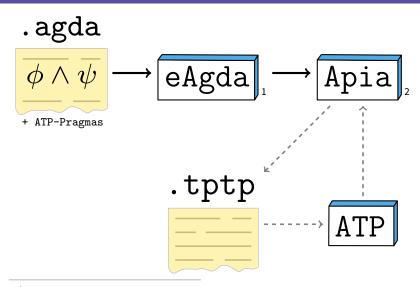
At the moment, the communication between Agda and the ATPs is unidirectional because the ATPs are being used as oracles (Sicard-Ramírez, 2015).



¹Development version of Agda in order to handle a new built-in ATP-pragma. https://github.com/asr/eagda

Previous Work: Apia

Proving first-order theorems written in Agda using automatic theorem provers for first-order logic



 $^{^1 {\}tt Development} \ {\tt version} \ {\tt of Agdain} \ {\tt order} \ {\tt to handle anew built-in ATP-pragma}. \ {\tt https://github.com/asr/eagdain} \ {\tt order} \ {\tt to handle anew built-in ATP-pragma}. \ {\tt https://github.com/asr/eagdain} \ {\tt order} \$

²Haskell program for proving first-order theorems written in Agda using ATPs. https://github.com/asr/apia

.tptp

- ▶ Is a language³ to encode problems in text files
- Is the input of the ATPs
- his problems contains formulas with the form

language(name, role, formula).

```
language FOF, or CNF
    name to identify the formula within the problem
    role axiom, definition, hypothesis, conjecture, among others
formula the logic formula in the language
```

Go to Reconstruction

³Is available at http://www.cs.miami.edu/~tptp/TPTP/SyntaxBNF.html

Problems in Propositional Logic:

 $\triangleright p \vdash p$

```
$ cat basic-4.tptp
fof(a, axiom, p).
fof(goal, conjecture, p).
```

 $\triangleright p \land q \vdash q \land p$

```
$ cat conj-3.tptp
fof(a, axiom, p & q).
fof(goal, conjecture, q & p).
```

 $\blacktriangleright \vdash \neg (p \land \neg p) \lor (q \land \neg q)$

```
$ cat neg-7.tptp
fof(goal, conjecture, ~ ((p & ~ p) | (q & ~ q))).
```

⁴Is available at http://github.com/jonaprieto/pro-pack

.tstp

A TSTP derivation 5

▶ Is a Directed Acyclic Graph where

leaf is a formulae from the TPTP input **node** is a formulae inferred from parent formulae

root the final derived formulae

Is a list of annotated formulae:

language(name, role, formula, source [,useful info]).

where **source** typically is a inference record:

inference(rule, useful information, parents)

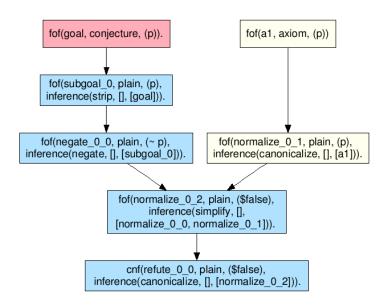
⁵http://www.cs.miami.edu/~tptp/TPTP/QuickGuide/Derivations.html



▶ Proof found by **Metis** ATP for the problem $p \vdash p$

```
$ metis --show proof basic-4.tptp
fof(a, axiom, (p)).
fof(goal, conjecture, (p)).
fof(subgoal 0, plain, (p),
 inference(strip, [], [goal])).
fof(negate_0_0, plain, (~ p),
 inference(negate, [], [subgoal_0])).
fof(normalize_0_0, plain, (~ p),
 inference(canonicalize, [], [negate_0_0])).
fof(normalize_0_1, plain, (p),
 inference(canonicalize, [], [a])).
fof(normalize 0 2, plain, ($false),
 inference(simplify, [],
   [normalize_0_0, normalize_0_1])).
cnf(refute_0_0, plain, ($false),
   inference(canonicalize, [], [normalize_0_2])).
```

DAG for the previous TSTP derivation found by Metis ATP



Design Decisions for the Reconstruction Tool

Athena

Haskell

- Parsing
- AST construction
- Creation and analysis of DAG derivations
- Analysis of inference rules used
- Generation of Agda code of the proof

Design Decisions for the Reconstruction Tool

Programming Languages

Agda is a dependently typed functional programming language and it also a proof assistant.

Our usages:

- Logic framework for Classical Propositional Logic
- Type-Checker for the Agda code of the proofs generated

References



Sicard-Ramírez, Andrés (2015). Reasoning about functional programs by combining interactive and automatic proofs. PEDECIBA Informática, Universidad de la República.