

Query Optimization With MySQL 8.0 and MariaDB 10.3: The Basics

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dbahire.com/pleu18

Agenda

1. Introduction	6. Joins
2. Query Profiling	7. Subqueries
3. Access Types and Basic Indexing Techniques	8. FULLTEXT, Geodata Search and JSON support
4. Multi-Column Indexing	9. Conclusions
5. Break	

Query Optimization with MySQL 8.0 and MariaDB 10.3: The Basics

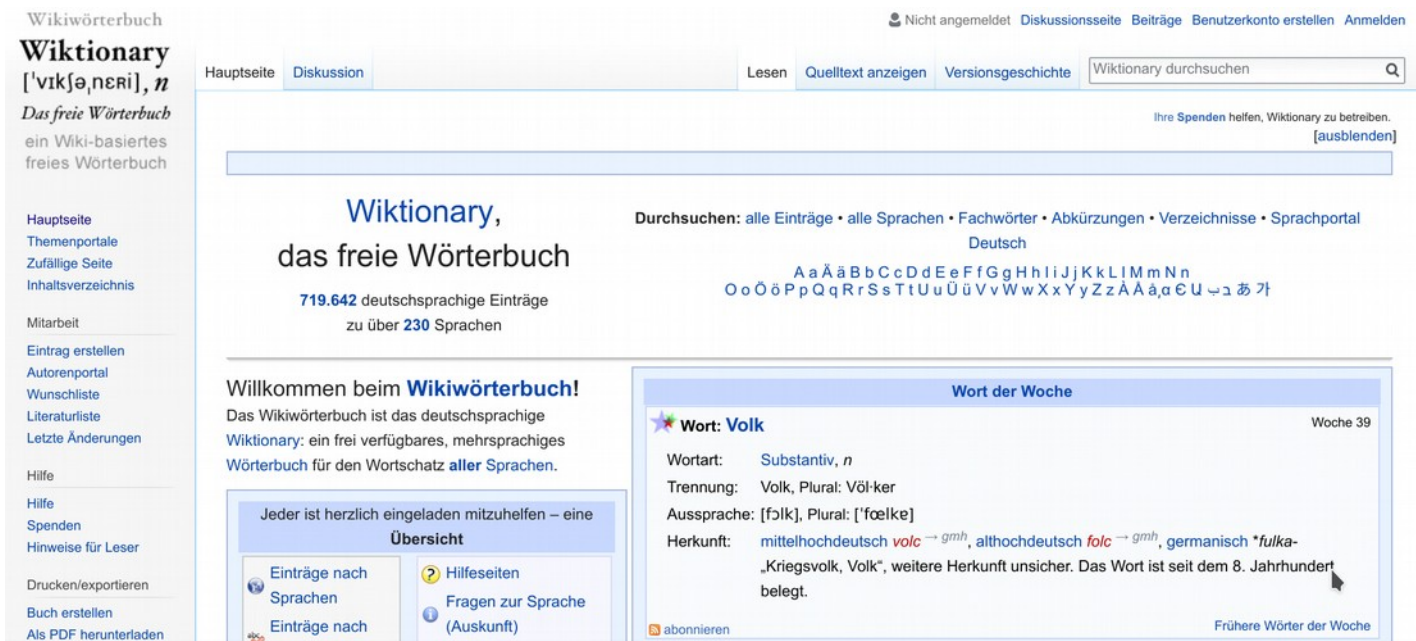
INTRODUCTION

This is me fighting bad query performance



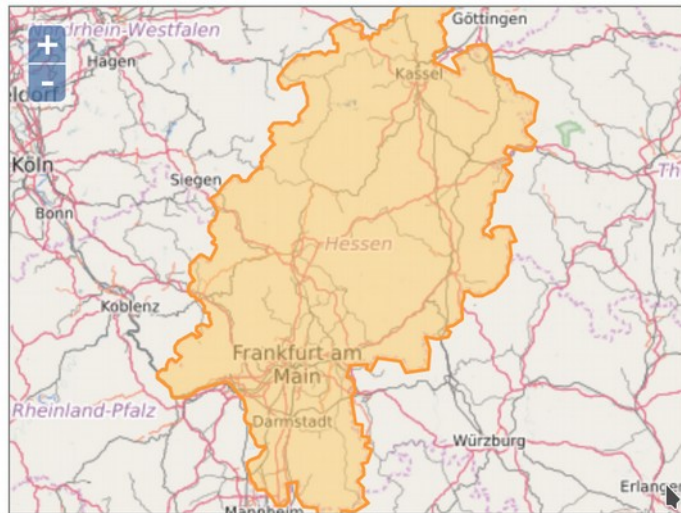
- Sr. Database Administrator at Wikimedia Foundation
- Used to work as a trainer for Oracle (MySQL), as a Consultant (Percona) and as a Freelance administrator (DBAHire.com)

Example Application (I)



- Wiktionary (and all Wikimedia projects' data) is licensed under the Creative Commons BY-SA-3.0 License and is Copyright its Contributors

Example Application (II)




- OSM Database is licensed under the Open DataBase License and is Copyright OpenStreetMap Contributors

Install the example databases

- Downloads and instructions at:
<http://dbahire.com/pleu18>
 - Requirements: a MySQL 8.0 or MariaDB 10.3 installation
 - The Wiktionary and OSM extracts
- Import them by doing:
`$ tar xf <file> && cd <dir> && myloader -d .`

Alternative: Query the Live Database

- <https://quarry.wmflabs.org/>

 Quarry Beta

[Home](#) [New Query](#) [Recent Queries](#) [Discuss](#) [Database tables](#) ▼

JCrespo (WMF) ▼

SQL

```
use enwiki_p;  
SELECT max(rev_id) FROM revision;
```

By running queries you agree to the [Labs ToS](#) and you irrevocably agree to release your SQL under [CC0 License](#).

Submit Query

Query status: **complete**

Executed in 0.13 seconds.

Resultset (1 rows)

Download data ▼

Search:

max(rev_id)
865217508

Query Optimization with MySQL 8.0 and MariaDB 10.3: The Basics

QUERY PROFILING

Which Queries Should I Optimize First?

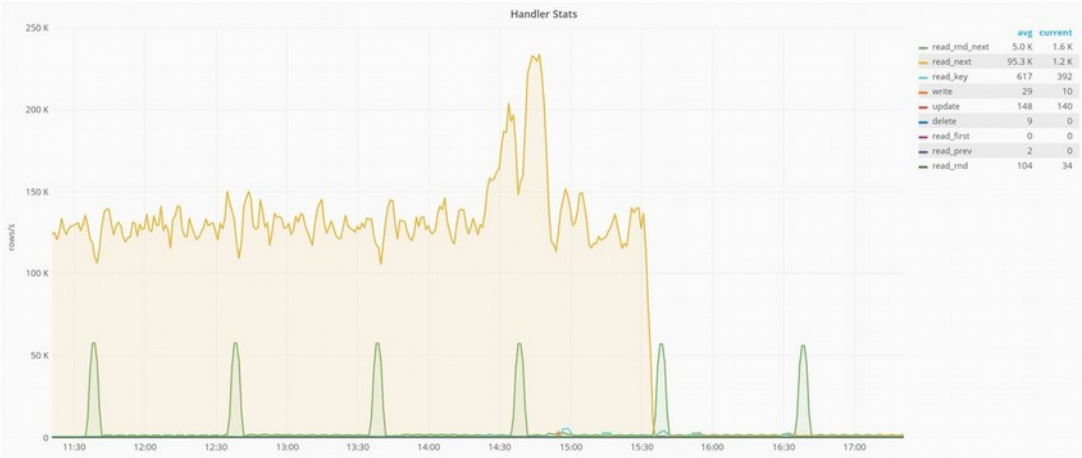
- Tools:
 - Application-level profiling
 - General monitoring metrics (CPU, IO, etc.)
 - Aggregated PROCESSLIST
 - pt-query-digest
 - PERFORMANCE_SCHEMA
- They are not exclusive to each other

Application-Level profiling

▶	2018-10-22T14:09:55	hhvm	ERROR	-	mw1309	SlowTimer [15002ms] at runtime/ext_mysql: slow query: SELECT /* LinksUpdate::acquirePageLock */ GET_LOCK('LinksUpdate:job:pageid:55791340', 15) AS lockstatus
---	---------------------	------	-------	---	--------	---

DBQuery						
▶	2018-10-15T05:16:42	10.64.48.172	FetchText::doGetText	MySQL server has gone away (10.64.48.172)	SELECT old_text,old_flags FROM `text` WHERE old_id = '750931741' LIMIT 1	
▶	2018-10-15T05:16:42	10.64.48.172	FetchText::doGetText	MySQL server has gone away (10.64.48.172)	SELECT old_text,old_flags FROM `text` WHERE old_id = '743182175' LIMIT 1	
▶	2018-10-15T05:16:42	10.64.48.172	FetchText::doGetText	MySQL server has gone away (10.64.48.172)	SELECT old_text,old_flags FROM `text` WHERE old_id = '754391637' LIMIT 1	

General Monitoring Metrics



pt-query-digest

- It is a 3rd party tool written in Perl, originally created by Baron Schwartz
- It requires activation of the slow log:
 - `SET GLOBAL slow_query_log = 1;`
 - `SET long_query_time = 0;`
- In Percona Server and MariaDB it can provide extra information:
 - `SHOW GLOBAL VARIABLES like 'log_slow_verbosity';`



**Be careful
with extra
IO and
latency!**

pt-query-digest Execution (I)

```
# 1094.7s user time, 9.4s system time, 141.22M rss, 205.21M vsz
# Current date: Wed Jul 1 07:32:28 2015
# Hostname: db1018
# Files: STDIN
# Overall: 4.66M total, 640 unique, 53.47 QPS, 0.02x concurrency _____
# Time range: 2015-06-30 07:00:10 to 2015-07-01 07:11:37
# Attribute          total      min      max      avg      95%    stddev  median
# =====
# Exec time          1320s      1us      3s      283us    332us     3ms    152us
# Lock time           238s        0     13ms     51us     93us    39us    52us
# Rows sent           5.02M        0    4.16k     1.13     1.96     8.69     0.99
# Rows examine        9.50M        0 232.93k     2.14     3.89   261.15     0.99
# Merge passes         0         0         0         0         0         0         0
# Query size          1.06G        17  67.89k   243.89   511.45   368.99   192.76
# Boolean:
# Filesort           8% yes,   91% no
# Full scan          94% yes,    5% no
# Priority que        3% yes,   96% no
# Tmp table          29% yes,   70% no
# Tmp table on        1% yes,   98% no
```

Actual execution
on Wikipedia
production
servers

pt-query-digest Execution (II)

# Profile									
#	Rank	Query ID	Response time		Calls	R/Call	V/M	Item	
#	====	=====	=====	=====	=====	=====	=====	=====	
#	1	0xSANITIZED	242.2765	18.4%	691005	0.0004	0.00	SELECT	revision page user
#	2	0xSANITIZED	204.7052	15.5%	80863	0.0025	0.01	SELECT	revision page user
#	3	0xSANITIZED	162.8476	12.3%	1025179	0.0002	0.00	SELECT	page
#	4	0xSANITIZED	68.1164	5.2%	93928	0.0007	0.01	SELECT	revision page user
#	5	0xSANITIZED	66.8302	5.1%	354562	0.0002	0.00	SELECT	page revision
#	6	0xSANITIZED	57.0374	4.3%	211631	0.0003	0.00	SELECT	page revision
#	7	0xSANITIZED	44.0751	3.3%	6925	0.0064	0.07	SELECT	page categorylinks
category									
#	8	0xSANITIZED	35.0655	2.7%	9689	0.0036	0.00	SELECT	text
#	9	0xSANITIZED	29.4363	2.2%	152259	0.0002	0.00	SELECT	page
#	10	0xSANITIZED	24.1864	1.8%	176927	0.0001	0.00	SELECT	msg_resource
#	11	0xSANITIZED	23.7016	1.8%	144807	0.0002	0.00	SELECT	page_restrictions
#	12	0xSANITIZED	16.6547	1.3%	10135	0.0016	0.03	SELECT	revision
#	13	0xSANITIZED	15.0564	1.1%	263809	0.0001	0.00	SET	

pt-query-digest Execution (III)

```
# Query 1: 7.93 QPS, 0.00x concurrency, ID 0xSANITIZED at byte 1553864032
# This item is included in the report because it matches -limit.
# Scores: V/M = 0.00
# Time range: 2015-06-30 07:00:10 to 2015-07-01 07:11:37
# Attribute      pct      total      min       max       avg       95%      stddev   median
# =====
# Count          14    691005
# Exec time      18     242s    163us     91ms    350us    348us    563us    301us
# Lock time      26      63s     47us      7ms     91us    103us    14us     84us
# Rows sent      12   657.18k      0        1     0.97    0.99    0.16     0.99
# Rows examine   6   657.18k      0        1     0.97    0.99    0.16     0.99
# Query size     31  345.42M     501      749   524.16  537.02    9.22   511.45
# String:
# Databases      itwiki (225976/32%), enwiktiona... (219461/31%)... 15 more
# Hosts
# Users          wikiuser
# Query_time distribution
#   1us
#  10us
# 100us #####
#   1ms #
#  10ms #
# 100ms
#   1s
#  10s+
# Tables
#   SHOW TABLE STATUS FROM `enwiktionary` LIKE 'revision'\G
#   SHOW CREATE TABLE `enwiktionary`.`revision`\G
#   SHOW TABLE STATUS FROM `enwiktionary` LIKE 'page'\G
#   SHOW CREATE TABLE `enwiktionary`.`page`\G
#   SHOW TABLE STATUS FROM `enwiktionary` LIKE 'user'\G
#   SHOW CREATE TABLE `enwiktionary`.`user`\G
# EXPLAIN /*!50100 PARTITIONS*/
SELECT /* Revision::fetchFromConds SANITIZED */ * FROM `revision`
INNER JOIN `page` ON ((page_id = rev_page)) LEFT JOIN `user` ON
((rev_user != 0) AND (user_id = rev_user)) WHERE page_namespace = '0' AND
page_title = 'SANITIZED' AND (rev_id=page_latest) LIMIT 1\G
```


PERFORMANCE_SCHEMA

- Monitoring schema (engine) enabled by default since MySQL 5.6
 - `performance_schema = 1` (it is not dynamic)
- Deprecates the old query profiling
- It is way more user-friendly when combined with the [SYS schema/ps_helper](#) (a set of views and stored procedures created by Mark Leith)
 - Included by default since 5.7.7

Installation of the SYS Schema

<5.7/MariaDB

```
$ git clone https://github.com/MarkLeith/mysql-sys.git
Cloning into 'mysql-sys'...
remote: Counting objects: 926, done.
remote: Compressing objects: 100% (73/73), done.
remote: Total 926 (delta 35), reused 6 (delta 2)
Receiving objects: 100% (926/926), 452.19 KiB | 225.00
KiB/s, done.
Resolving deltas: 100% (584/584), done.
$ cd mysql-sys/
$ mysql < sys_56.sql
```

Example Usage: Discovering Unused Indexes

```
mysql (osm) > SELECT * FROM sys.schema_unused_indexes LIMIT 5;
```

object_schema	object_name	index_name
osm	acls	acls_k_idx
osm	changeset_tags	changeset_tags_id_idx
osm	current_nodes	current_nodes_timestamp_idx
osm	current_nodes	current_nodes_tile_idx
osm	current_relations	current_relations_timestamp_idx

5 rows in set (0.04 sec)

```
mysql-5.7.8 (osm) > SELECT * FROM current_nodes WHERE tile = 100;
```

...

```
mysql-5.7.8 (osm) > SELECT * FROM sys.schema_unused_indexes LIMIT 5;
```

object_schema	object_name	index_name
osm	acls	acls_k_idx
osm	changeset_tags	changeset_tags_id_idx
osm	current_nodes	current_nodes_timestamp_idx
osm	current_relations	current_relations_timestamp_idx
osm	current_relations	changeset_id

5 rows in set (0.03 sec)



Example Usage: Slow Queries (ordered by server time)

```
mysql (osm) > SELECT * FROM sys.statement_analysis LIMIT 10\G
```

```
***** 1. row *****
      query: SELECT `way_id` AS
`id` , `v` FROM `way_tags` WHERE `v` LIKE ?
      db: osm
      full_scan: *
      exec_count: 15
      err_count: 0
      warn_count: 0
      total_latency: 7.83 s
      max_latency: 1.33 s
      avg_latency: 521.84 ms
      lock_latency: 17.94 ms
      rows_sent: 6779
      rows_sent_avg: 452
      rows_examined: 20152155
      rows_examined_avg: 1343477
      rows_affected: 0
      rows_affected_avg: 0
      tmp_tables: 0
      tmp_disk_tables: 0
      rows_sorted: 0
      sort_merge_passes: 0
      digest:
21f90695b1ebf20a5f4d4c1e5e860f58
      first_seen: 2014-11-01 17:04:51
      last_seen: 2014-11-01 17:05:22
```

Example Usage: Top Queries

Creating Temporary Tables

```
mysql (osm) > SELECT * FROM sys.statements_with_temp_tables  
WHERE db = 'osm' LIMIT 10\G
```

```
***** 1. row *****
```

```
query: SELECT ? AS TYPE , `node_id` A ... gs` WHERE `k` = ? AND `v` = ?
```

```
db: osm
```

```
exec_count: 11
```

```
total_latency: 7.57 s
```

```
memory_tmp_tables: 11
```

```
disk_tmp_tables: 0
```

```
avg_tmp_tables_per_query: 1
```

```
tmp_tables_to_disk_pct: 0
```

```
first_seen: 2014-11-01 17:33:55
```

```
last_seen: 2014-11-01 17:34:45
```

```
digest: 5e6e82799b7c7c0e5c57cfe63eb98d5d
```

Example Usage: Top Queries

Creating Temporary Tables (cont.)

```
mysql (osm) > SELECT DIGEST_TEXT FROM performance_schema.events_statements_summary_by_digest WHERE digest = '5e6e82799b7c7c0e5c57cfe63eb98d5d'\G
```

```
***** 1. row *****
DIGEST_TEXT: SELECT ? AS TYPE , `node_id` AS `id` FROM `node_tags` WHERE `k` = ? AND `v` = ? UNION
SELECT ? AS TYPE , `way_id` AS `id` FROM `way_tags` WHERE `k` = ? AND `v` = ? UNION SELECT ? AS
TYPE , `relation_id` AS `id` FROM `relation_tags` WHERE `k` = ? AND `v` = ?
1 row in set (0.00 sec)
```

```
mysql-5.7.8 (osm) > EXPLAIN SELECT 'node' as type, node_id as id FROM node_tags WHERE k='amenity'
and v='cafe' UNION SELECT 'way' as type, way_id as id FROM way_tags WHERE k='amenity' and v='cafe'
UNION SELECT 'relation' as type, relation_id as id FROM relation_tags WHERE k='amenity' and
v='cafe';
```

id	select_type	table	partitions	type	possible_keys	key	key_len	ref	rows	filtered	Extra
1	PRIMARY	node_tags	NULL	ALL	NULL	NULL	NULL	NULL	851339	0.00	Using where
2	UNION	way_tags	NULL	ALL	NULL	NULL	NULL	NULL	1331016	0.00	Using where
3	UNION	relation_tags	NULL	ALL	NULL	NULL	NULL	NULL	63201	0.00	Using where
NULL	UNION RESULT	<union1,2,3>	NULL	ALL	NULL	NULL	NULL	NULL	NULL	NULL	Using temporary

4 rows in set, 1 warning (0.01 sec)

Query Optimization with MySQL 8.0 and MariaDB 10.3: The Basics

ACCESS TYPES AND BASIC INDEXING TECHNIQUES

EXPLAIN

- Essential to understand the execution plan of our queries
 - Works on SELECTs, INSERTs, UPDATEs, REPLACEs, DELETEs and connections
 - Fully documented on:
<https://dev.mysql.com/doc/refman/8.0/en/explain-output.html>

EXPLAIN Example

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM page
WHERE page_title = 'German';
```

+-----+-----+-----+-----+-----+-----+-----						
+-----+-----+-----+-----+-----+-----+-----						
id	select_type	table	type	possible_keys		
key	key_len	ref	rows	Extra		
+-----+-----+-----+-----+-----+-----+-----						
+-----+-----+-----+-----+-----+-----+-----						
1	SIMPLE	page	ALL	NULL		
NULL	NULL	NULL	778885	Using		
+-----+-----+-----+-----+-----+-----+-----						
+-----+-----+-----+-----+-----+-----+-----						
1 row in set (0.00 sec)						

Difficult to see something

EXPLAIN Example (vertical format)

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM page WHERE  
page_title = 'German'\G
```

```
***** 1. row *****  
      id: 1  
  select_type: SIMPLE  
        table: page  
         type: ALL  
possible_keys: NULL  
          key: NULL  
     key_len: NULL  
         ref: NULL  
        rows: 778885  
   Extra: Using where  
1 row in set (0.00 sec)
```



**Use \G for
vertical
formatting**

EXPLAIN Example (id)

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM page WHERE  
page_title = 'German'\G
```

```
***** 1. row *****
```

```
    id: 1  
  select_type: SIMPLE  
    table: page  
    type: ALL  
possible_keys: NULL  
    key: NULL  
  key_len: NULL  
    ref: NULL  
   rows: 778885  
  Extra: Using where  
1 row in set (0.00 sec)
```



**Indicates nesting level,
not execution order**

EXPLAIN Example (select_type)

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM page WHERE
page_title = 'German'\G
***** 1. row *****
      id: 1
  select_type: SIMPLE
        table: page
         type: ALL
possible_keys: NULL
         key: NULL
      key_len: NULL
         ref: NULL
        rows: 778885
      Extra: Using where
1 row in set (0.00 sec)
```

Not a subquery or a UNION

EXPLAIN Example (table)

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM page WHERE  
page_title = 'German'\G
```

```
***** 1. row *****  
      id: 1  
  select_type: SIMPLE  
      table: page  
      type: ALL  
possible_keys: NULL  
      key: NULL  
   key_len: NULL  
      ref: NULL  
     rows: 778885  
   Extra: Using where  
1 row in set (0.00 sec)
```





Table scanned for
this step

EXPLAIN Example (type)

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM page WHERE  
page_title = 'German'\G
```

```
***** 1. row *****  
      id: 1  
    select_type: SIMPLE  
      table: page  
       type: ALL  
possible_keys: NULL  
        key: NULL  
     key_len: NULL  
        ref: NULL  
       rows: 778885  
    Extra: Using where  
1 row in set (0.00 sec)
```



**All rows are read for
this table (FULL
TABLE SCAN)**

EXPLAIN Example (rows)

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM page WHERE  
page_title = 'German'\G
```

```
***** 1. row *****
```

```
    id: 1  
  select_type: SIMPLE  
    table: page  
    type: ALL  
possible_keys: NULL  
    key: NULL  
  key_len: NULL  
    ref: NULL  
   rows: 778885  
Extra: Using where  
1 row in set (0.00 sec)
```

**Estimated number
of rows to be read
(all table rows)**

How to improve performance?

```
MariaDB [dewiktionary]> SELECT * FROM page  
WHERE page_title = 'German';  
1 row in set (0.10 sec)
```

- Let's add an index on page.page_title:

```
MariaDB [dewiktionary]> ALTER TABLE page ADD INDEX  
page_title (page_title);  
Query OK, 0 rows affected (22.70 sec)  
Records: 0 Duplicates: 0 Warnings: 0
```


Index creation results (type)

```

MariaDB [dewiktionary]> EXPLAIN SELECT * FROM page WHERE
page_title = 'German'\G
***** 1. row *****
      id: 1
  select_type: SIMPLE
        table: page
         type: ref
possible_keys: page_title
          key: page_title
        key_len: 257
         ref: const
         rows: 1
    Extra: Using index condition
1 row in set (0.00 sec)

```

type: ref means that an equality comparison will be checked against an index and several results could be returned

Index creation results (possible_keys and key)

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM page WHERE  
page_title = 'German'\G
```

```
***** 1. row *****
```

```
    id: 1  
  select_type: SIMPLE  
    table: page  
    type: ref  
possible_keys: page_title  
    key: page_title  
  key_len: 257  
    ref: const  
   rows: 1  
  Extra: Using index condition  
1 row in set (0.00 sec)
```

**index(es) that the
optimizer considered
potentially useful, and
final index chosen**

Index creation results (ref)

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM page WHERE
page_title = 'German'\G
```

```
***** 1. row *****
      id: 1
  select_type: SIMPLE
        table: page
         type: ref
possible_keys: page_title
          key: page_title
       key_len: 257
         ref: const
        rows: 1
   Extra: Using index
1 row in set (0.00 sec)
```

**Index is compared
with a constant, not
with another table**

Index creation results (rows)

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM page WHERE
page_title = 'German'\G
```

```
***** 1. row *****
      id: 1
  select_type: SIMPLE
        table: page
         type: ref
possible_keys: page_title
          key: page_title
       key_len: 257
         ref: const
        rows: 1
   Extra: Using index condit
1 row in set (0.00 sec)
```

Only 1 row read. In this case, estimation is exact (thanks to index dive)

Index creation results (query time)

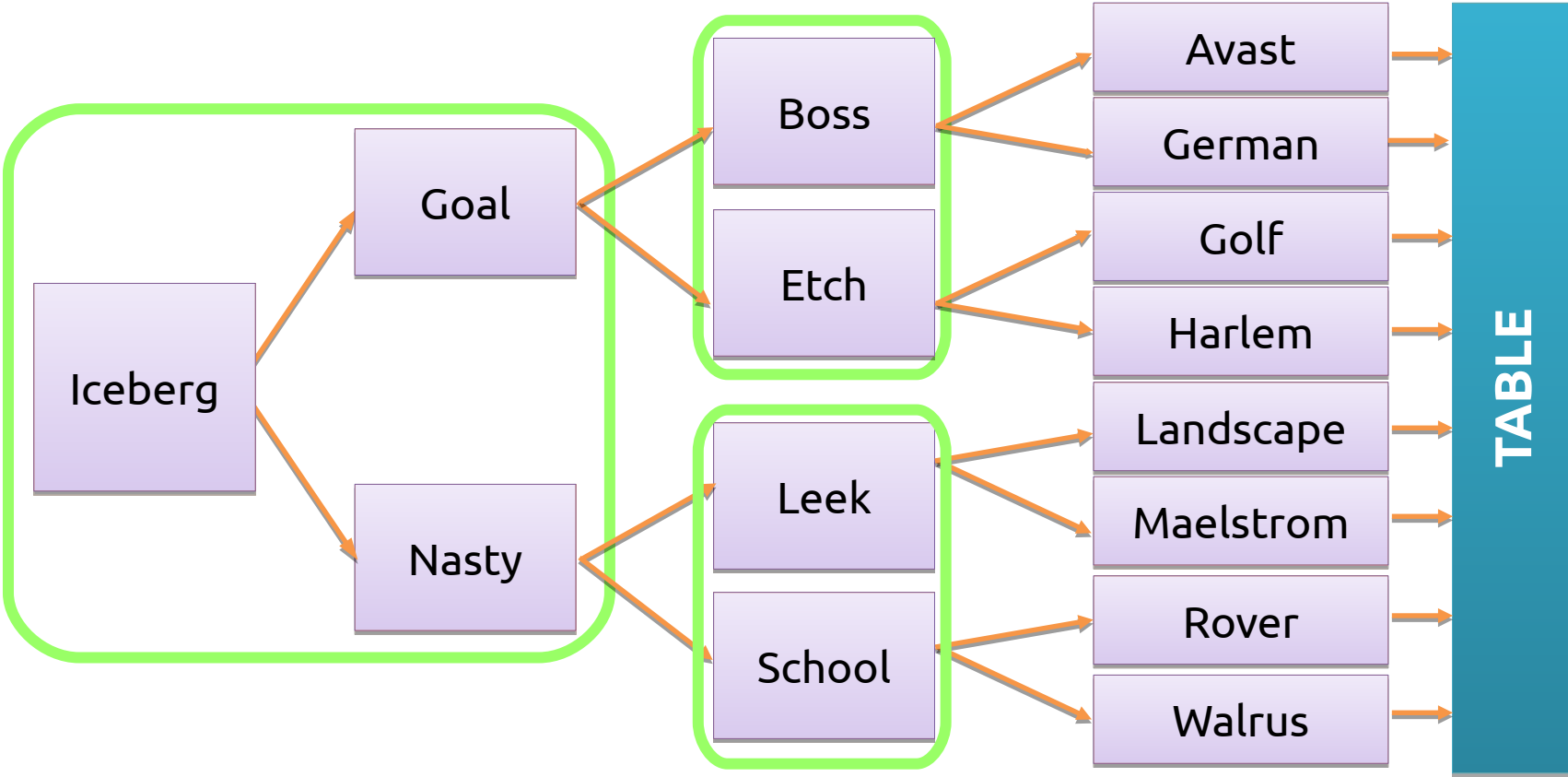
```
MariaDB [dewiktionary]> SELECT * FROM page  
WHERE page_title = 'German';  
1 row in set (0.00 sec)
```

**Query time has been
reduced substantially**

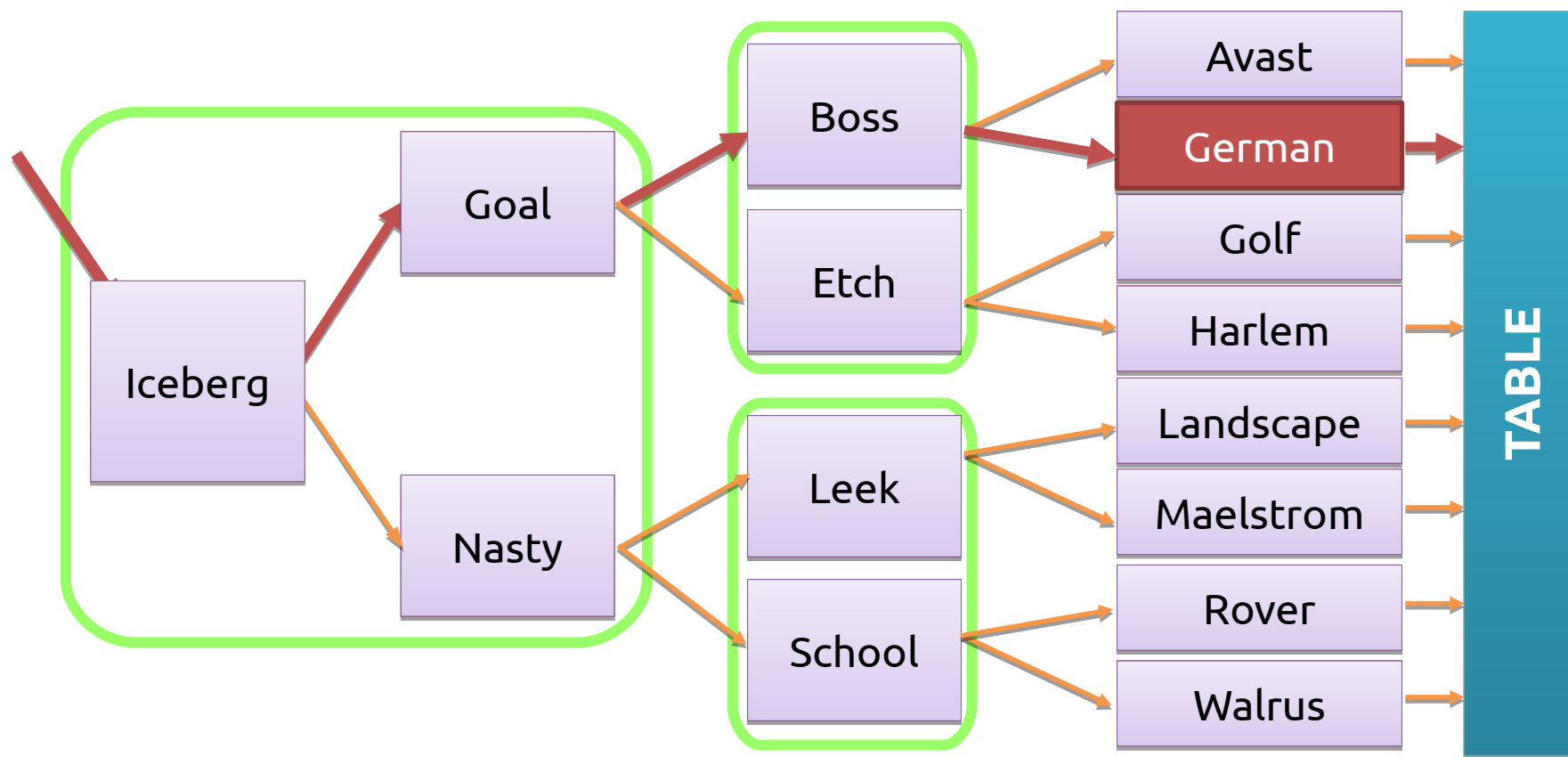
Types of indexes

- BTREE
 - B-TREE in MyISAM, B+TREE in InnoDB
- HASH
 - Only available for MEMORY and NDB
- FULLTEXT
 - Inverted indexes in MyISAM and InnoDB
- SPATIAL
 - RTREEs in MyISAM and InnoDB

Finding “German” with a BTREE



Finding “German” with a BTREE



Do indexes always work? (1/2)

- Can we use an index to make this query faster?

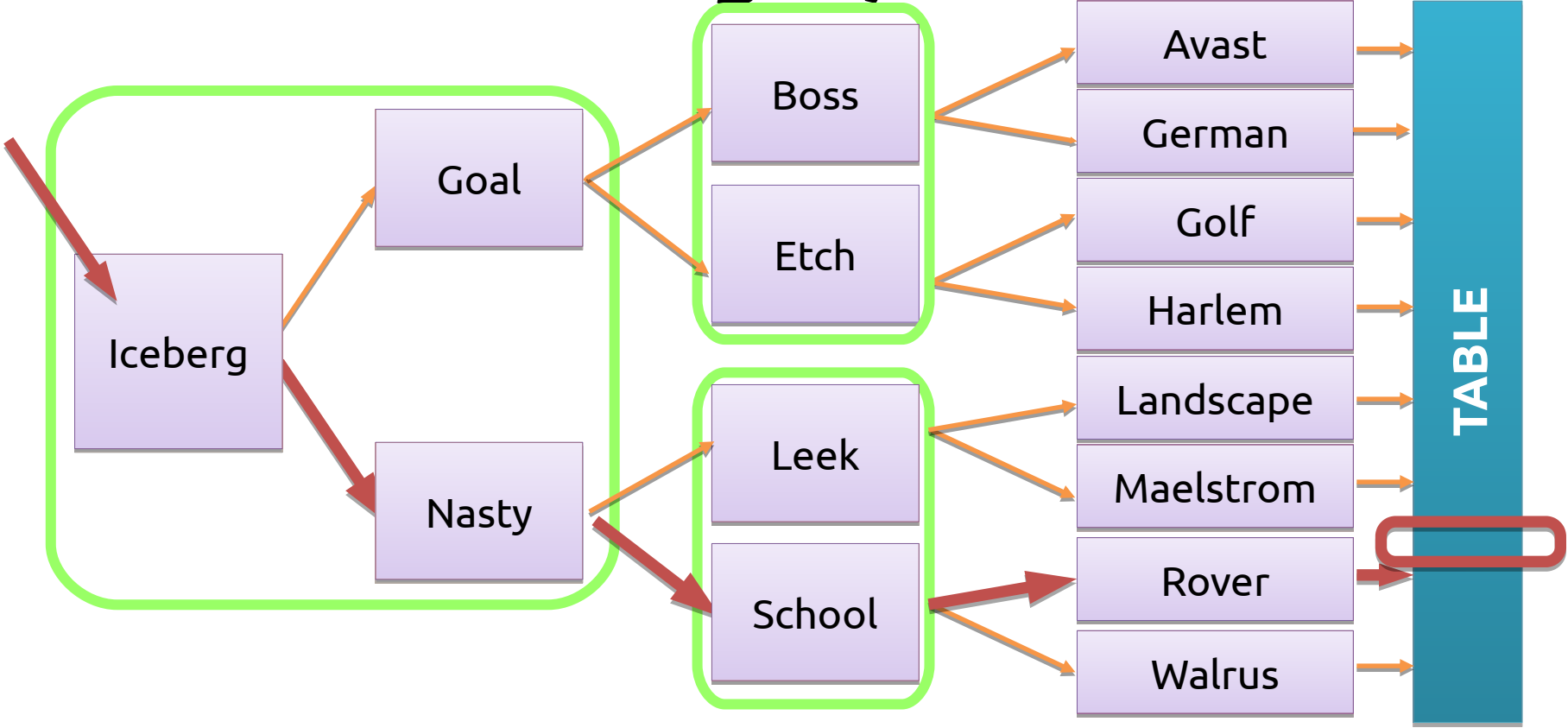
```
SELECT * FROM page WHERE page_title like 'Spa%';
```

It is a range

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM page WHERE
page_title like 'Spa%'\G
***** 1. row *****
      id: 1
  select_type: SIMPLE
        table: page
         type: range
possible_keys: page_title
         key: page_title
      key_len: 257
         ref: NULL
        rows: 563
   Extra: Using index condition
1 row in set (0.00 sec)
```

Despite not being an equality, we can use the index to find the values quickly

BTREE Indexes can be used for ranges, too



Do indexes always work? (2/2)

- What about this other query?

```
SELECT * FROM page WHERE page_title like '%erman';
```

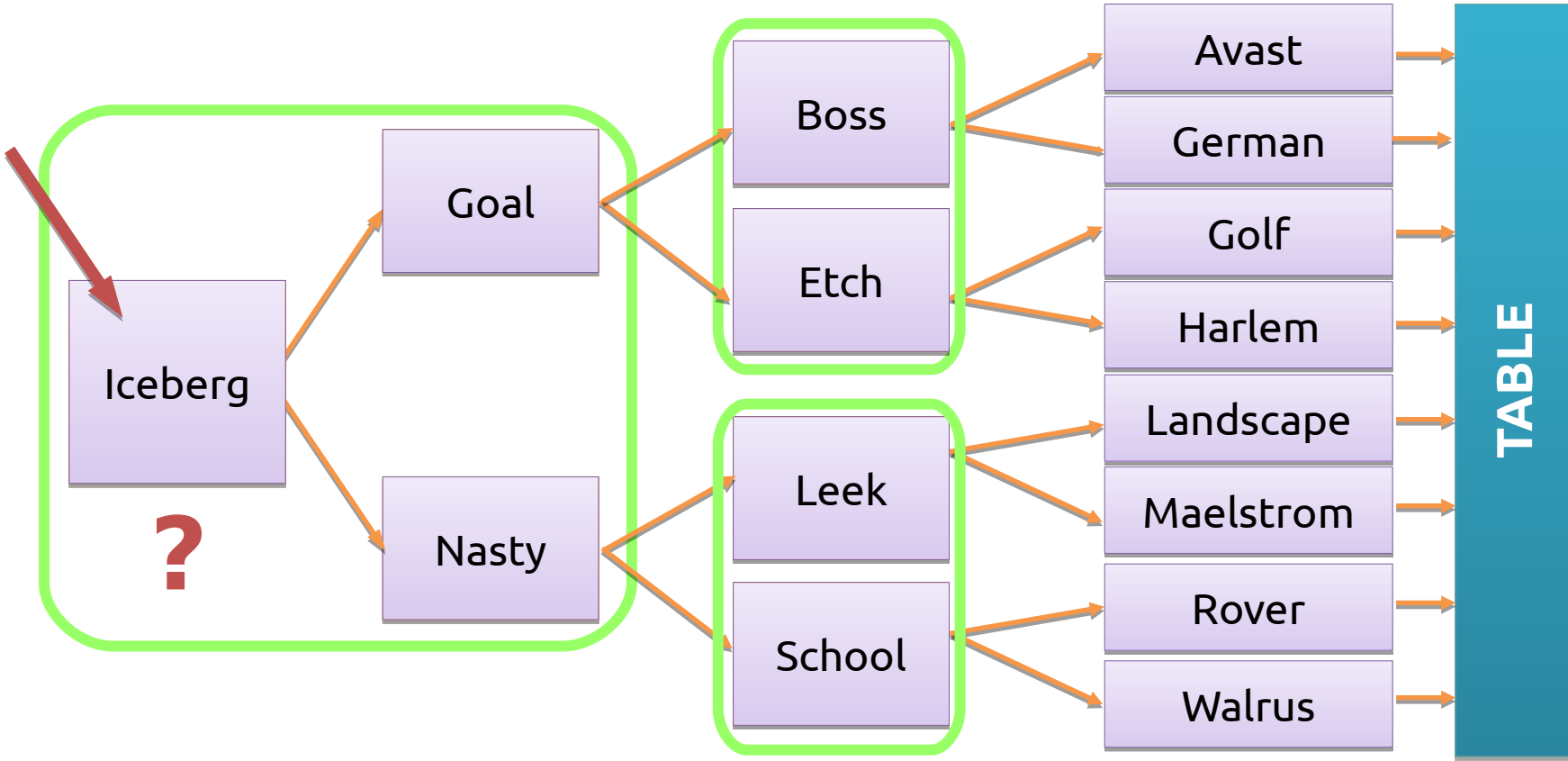
Let's check with EXPLAIN

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM page WHERE  
page_title like '%erman'\G
```

```
***** 1. row *****  
      id: 1  
  select_type: SIMPLE  
        table: page  
        type: ALL  
possible_keys: NULL  
         key: NULL  
    key_len: NULL  
         ref: NULL  
       rows: 778885  
  Extra: Using where  
1 row in set (0.00 sec)
```

**No index can be used
for filtering. A full
table scan is
performed.**

BTREE Index



Btree indexes usage (I)

- Filtering (“WHERE”)
 - Equality (operator '=')
 - Ranges (BETWEEN ... AND, >, <, >=, <=, like 'prefix%')
 - “EXISTS” operators: IN, OR on the same column
- Ordering
 - ORDER BY (indexed columns)
 - GROUP BY (indexed columns)

Btree indexes usage (II)

- Returning values directly from the index
 - Covering index
 - Functions like `max()`, `min()`, etc.
- Reduced contention (writes)
 - Locking less rows, increasing concurrency

type: const

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision WHERE  
rev_id = 2\G
```

```
***** 1. row *****  
      id: 1  
    select_type: SIMPLE  
      table: revision  
        type: const  
possible_keys: PRIMARY  
         key: PRIMARY  
    key_len: 4  
       ref: const  
      rows: 1  
    Extra:  
1 row in set (0.00 sec)
```

'const' is a special case of 'ref', when the index can assure that only 1 result can be returned (equality + primary key or unique key). It is faster.

type: NULL

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision WHERE  
rev_id = -1\G
```

```
***** 1. row *****
```

```
      id: 1  
    select_type: SIMPLE  
      table: NULL  
        type: NULL  
possible_keys: NULL  
          key: NULL  
    key_len: NULL  
         ref: NULL  
       rows: NULL
```

**'NULL' is not really a plan,
just an optimization that
allows immediately
discarding impossible
conditions**

```
      Extra: Impossible WHERE noticed after reading const  
tables  
1 row in set (0.00 sec)
```

type: ref_or_null

```
MariaDB [osm]> EXPLAIN SELECT * FROM nodes WHERE tile = 1 or  
tile is null\G
```

```
***** 1. row *****
```

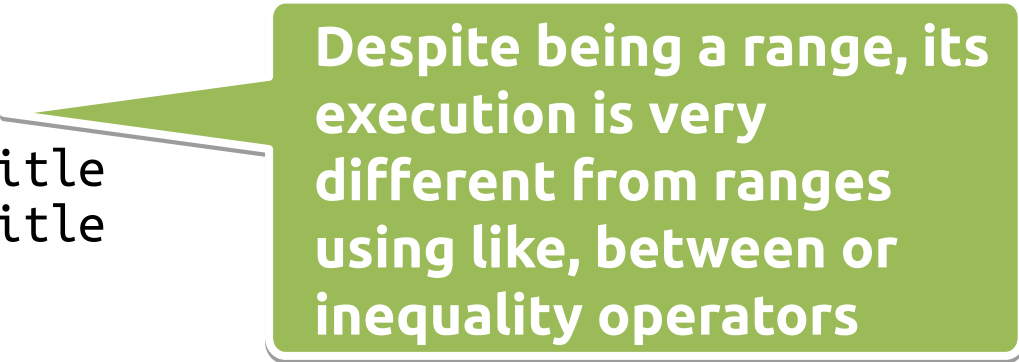
```
      id: 1  
    select_type: SIMPLE  
      table: nodes  
        type: ref_or_null  
possible_keys: nodes_tile_idx  
          key: nodes_tile_idx  
      key_len: 5  
        ref: const  
       rows: 2
```

```
      Extra: Using index condition; Using where  
1 row in set (0.00 sec)
```

**Equivalent to 'ref', but
also takes into account
NULL values**

type: range (using IN / OR)

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM page WHERE
page_title IN ('German', 'English', 'Spanish')\G
***** 1. row *****
      id: 1
    select_type: SIMPLE
      table: page
        type: range
possible_keys: page_title
         key: page_title
    key_len: 257
         ref: NULL
        rows: 4
    Extra: Using index condition
1 row in set (0.00 sec)
```



Despite being a range, its execution is very different from ranges using like, between or inequality operators

Is this a bug? (1/2)

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM page WHERE  
page_namespace > 200\G
```


```
***** 1. row *****  
      id: 1  
    select_type: SIMPLE  
      table: page  
       type: ref  
possible_keys: name_title  
       key: name_title  
    key_len: 4  
       ref: const  
      rows: 50  
    Extra:  
1 row in set (0.00 sec)
```

**An index is used to
return pages with
namespace>200**

Is this a bug? (2/2)

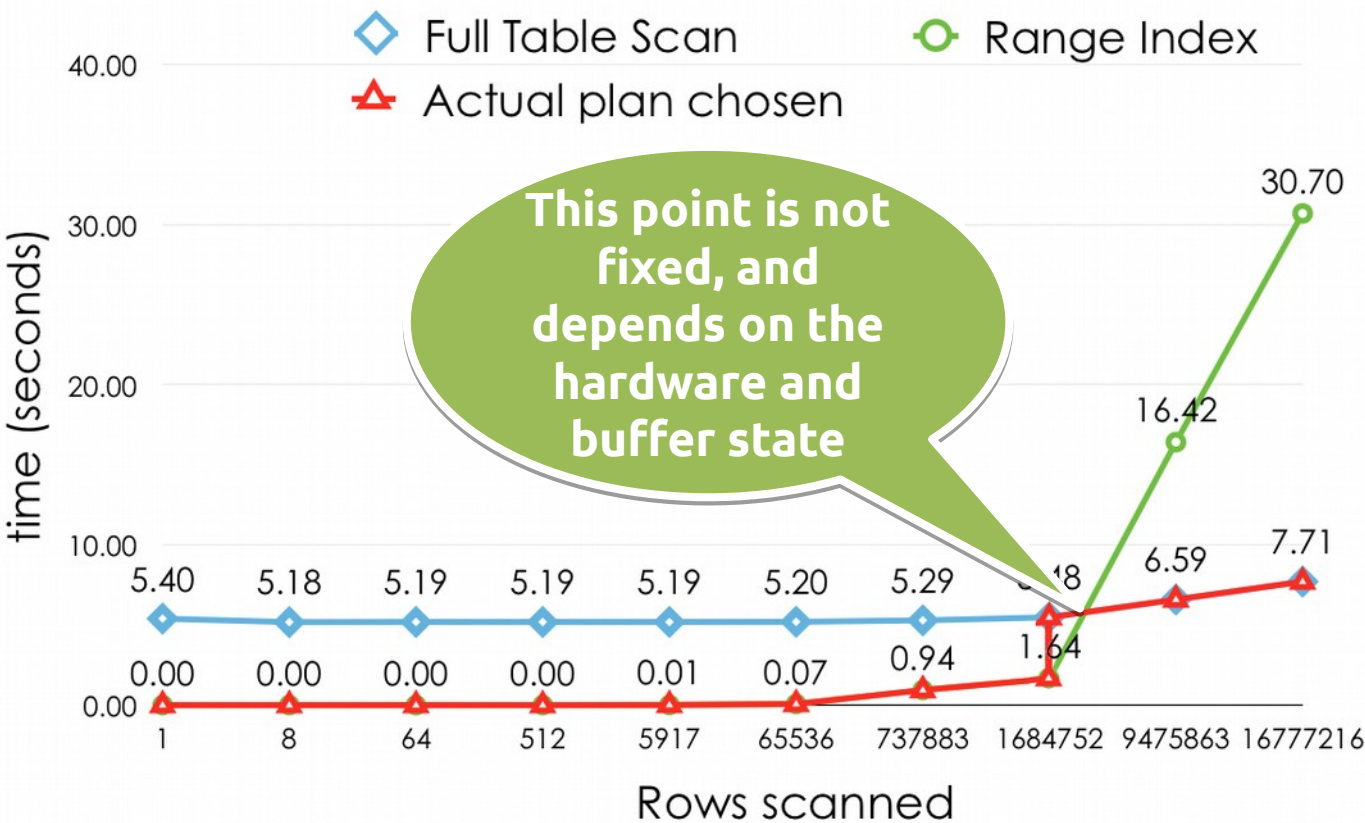
```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM page WHERE  
page_namespace < 200\G
```

```
***** 1. row *****  
      id: 1  
  select_type: SIMPLE  
        table: page  
         type: ALL  
possible_keys: name_title  
          key: NULL  
       key_len: NULL  
         ref: NULL  
        rows: 778885  
   Extra: Using where  
1 row in set (0.00 sec)
```



The index is not used
with ns=0

Using an index is sometimes suboptimal



What index should we add to make this query faster?

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision WHERE  
left(rev_timestamp, 6) = '201809'\G  
***** 1. row *****  
      id: 1  
  select_type: SIMPLE  
        table: revision  
         type: ALL  
possible_keys: NULL  
          key: NULL  
       key_len: NULL  
         ref: NULL  
        rows: 1376651  
   Extra: Using where  
1 row in set (0.00 sec)
```


The table has already an index on `rev_timestamp`

```
MariaDB [dewiktionary]> SHOW CREATE TABLE revision\G
```

```
***** 1. row *****
```

```
Table: revision
```

```
Create Table: CREATE TABLE `revision` (
```

```
...
```


```
KEY `rev_timestamp` (`rev_timestamp`),
```

We need to rewrite the query

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision WHERE  
rev_timestamp >= '201509' and rev_timestamp < '201510'\G  
***** 1. row *****  
      id: 1  
  select_type: SIMPLE  
        table: revision  
         type: range  
possible_keys: rev_timestamp  
          key: rev_timestamp  
       key_len: 14  
         ref: NULL  
        rows: 24588  
   Extra: Using index condition  
1 row in set (0.00 sec)
```

This transformation is not trivial or even possible in all cases

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision WHERE  
substr(rev_timestamp, 5, 2) = '09'\G  
***** 1. row *****  
      id: 1  
  select_type: SIMPLE  
        table: revision  
         type: ALL  
possible_keys: NULL  
         key: NULL  
      key_len: NULL  
         ref: NULL  
        Rows: 1376651  
      Extra: Using where  
1 row in set (0.00 sec)
```



Can you think of a way
to improve this query?

Indexes for Ordering

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM page ORDER BY  
page_touched DESC LIMIT 10\G
```

```
***** 1. row *****  
      id: 1  
    select_type: SIMPLE  
      table: page  
       type: ALL  
possible_keys: NULL  
       key: NULL  
    key_len: NULL  
       ref: NULL  
      rows: 778885  
    Extra: Using filesort  
1 row in set (0.00 sec)
```

**"Using filesort" indicates that
an ordering is needed before
returning the results**

If that is frequent, we can create an index on page_touched...

```
[dewiktionary]> ALTER TABLE page ADD INDEX page_page_touched(page_touched);  
Query OK, 0 rows affected (3.03 sec)  
Records: 0 Duplicates: 0 Warnings: 0
```

```
[dewiktionary]> EXPLAIN SELECT * FROM page ORDER BY page_touched DESC LIMIT  
10\G
```

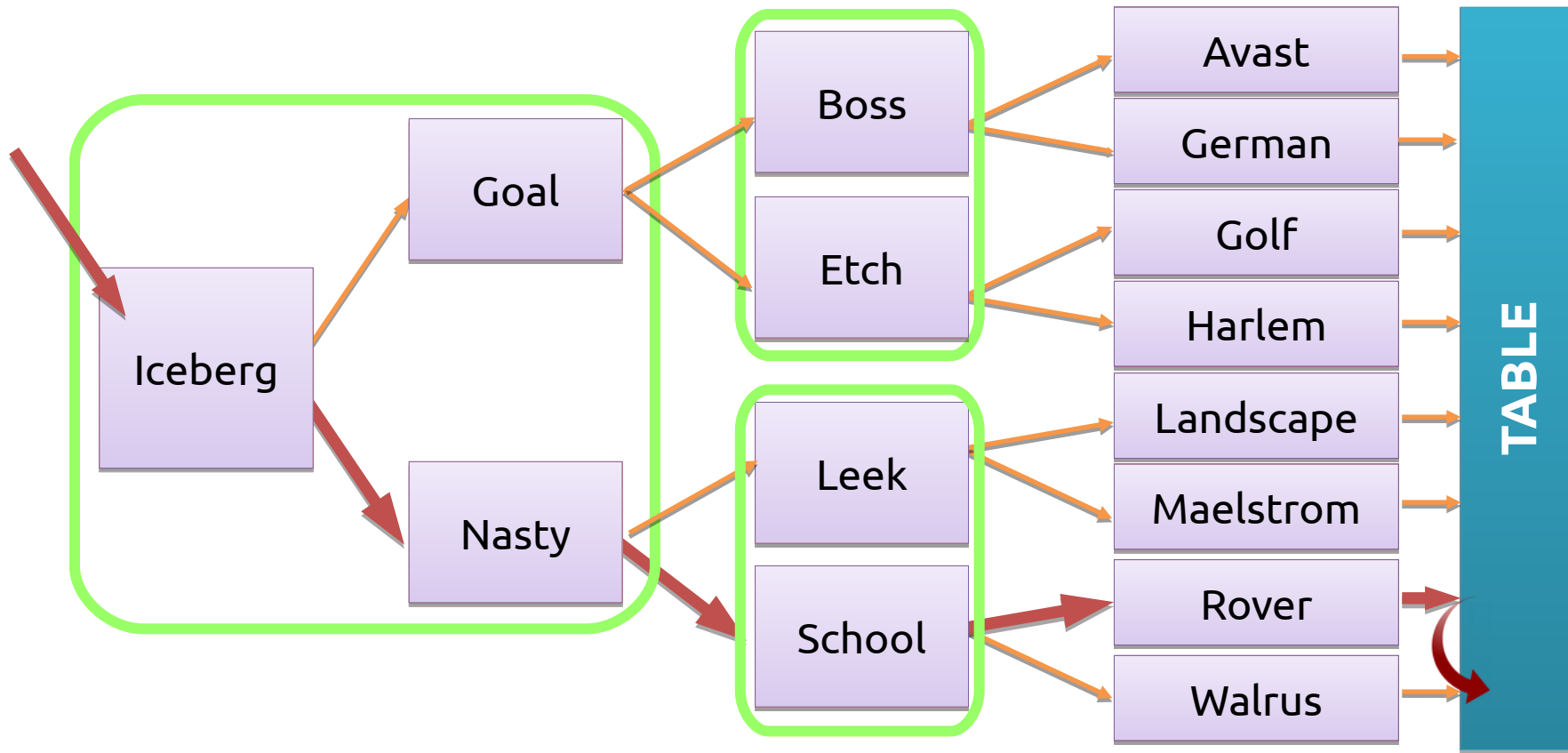
```
***** 1. row *****
```

```
id: 1  
select_type: SIMPLE  
table: page  
type: index  
possible_keys: NULL  
key: page_page_touched  
key_len: 16  
ref: NULL  
rows: 10  
Extra:  
1 row in set (0.00 sec)
```

The index does not produce any advantage for filtering

However, it is very effective by helping to avoid the sort phase

It can return data in index order faster



Trick: ORDER BY NULL avoids filesort

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision GROUP BY substr(rev_timestamp, 5, 2) = '09'\G
```

```
***** 1. row *****
```

```
id: 1
select_type: SIMPLE
table: revision
type: ALL
possible_keys: NULL
key: NULL
key_len: NULL
ref: NULL
rows: 1376651
Extra: Using temporary; Using filesort
1 row in set (0.00 sec)
```

There is no good index in this case

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM revision GROUP BY substr(rev_timestamp, 5, 2) = '09' ORDER BY NULL\G
```

```
***** 1. row *****
```

```
...
rows: 1376651
Extra: Using temporary
1 row in set (0.00 sec)
```

The advantage is not too big, but it avoids the filesort

Does LIMIT improve the performance? (I)

```
MariaDB [dewiktionary]> ALTER TABLE page ADD INDEX page_touched (page_touched); [...]  
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM page ORDER BY page_touched\G
```

```
***** 1. row *****
```

```
type: ALL  
possible_keys: NULL  
key: NULL  
key_len: NULL  
ref: NULL  
rows: 783231
```

```
Extra: Using filesort
```

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM page ORDER BY page_touched LIMIT 10\G
```

```
***** 1. row *****
```

```
type: index  
possible_keys: NULL  
key: page_page_touched  
key_len: 16  
ref: NULL  
rows: 10
```

```
Extra:
```

```
1 row in set (0.00 sec)
```

In some cases it can be essential to allow effective usage of the indexes

Does LIMIT improve the performance? (II)

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision ORDER BY rev_comment\G
***** 1. row *****
[...]
```

type:	ALL
possible_keys:	NULL
key:	NULL
key_len:	NULL
ref:	NULL
rows:	817636
Extra:	Using filesort

1 row in set (0.00 sec)

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision ORDER BY rev_comment LIMIT 10\G
***** 1. row *****
[...]
```

table:	revision
type:	ALL
possible_keys:	NULL
key:	NULL
key_len:	NULL
ref:	NULL
rows:	817636
Extra:	Using filesort

1 row in set (0.00 sec)

In other cases, it has no effect on the scanned rows (just on the returned ones)

Does LIMIT improve the performance? (III)

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM page ORDER BY page_title LIMIT 100\G
```

```
***** 1. row *****
```

```
    type: index
possible_keys: NULL
    key: page_title
   key_len: 257
    ref: NULL
   rows: 100
  Extra:
1 row in set (0.00 sec)
```

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM page ORDER BY page_title LIMIT 10000, 100\G
```

```
***** 1. row *****
```

```
    type: ALL
possible_keys: NULL
    key: NULL
   key_len: NULL
    ref: NULL
   rows: 90956
  Extra: Using filesort
1 row in set (0.00 sec)
```

**In this case,
performance will vary
depending on the
offset (not ideal)**

Indexes and GROUP BY (no indexes)

```
MariaDB [dewiktionary]> EXPLAIN SELECT rev_page, count(*) FROM  
revision IGNORE INDEX(rev_page_id, page_timestamp,  
page_user_timestamp) GROUP BY rev_page\G
```

```
***** 1. row *****
```

```
id: 1
```

```
select_type: SIMPLE
```

```
table: revision
```

```
type: ALL
```

```
possible_keys: NULL
```

```
key: NULL
```

```
key_len: NULL
```

```
ref: NULL
```

```
rows: 1376651
```

```
Extra: Using temporary; Using filesort
```

```
1 row in set (0.00 sec)
```

Without indexes, a temporary table is created to order results

Indexes and GROUP BY (rev_page_id)

```
MariaDB [dewiktionary]> EXPLAIN SELECT rev_page, count(*) FROM
revision GROUP BY rev_page\G
```

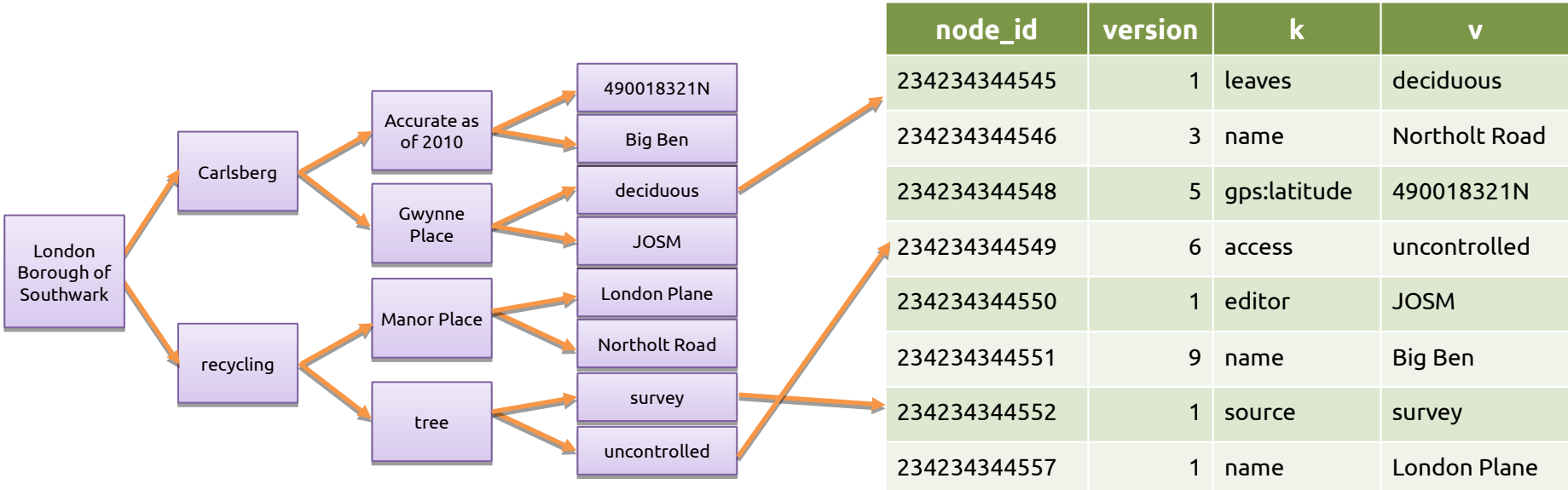
```
***** 1. row *****
```

```
id: 1
select_type: SIMPLE
table: revision
type: index
possible_keys: NULL
key: rev_page_id
key_len: 8
ref: NULL
rows: 1376651
Extra: Using index
1 row in set (0.00 sec)
```

The index does not produce any advantage for filtering (there is no WHERE clause)

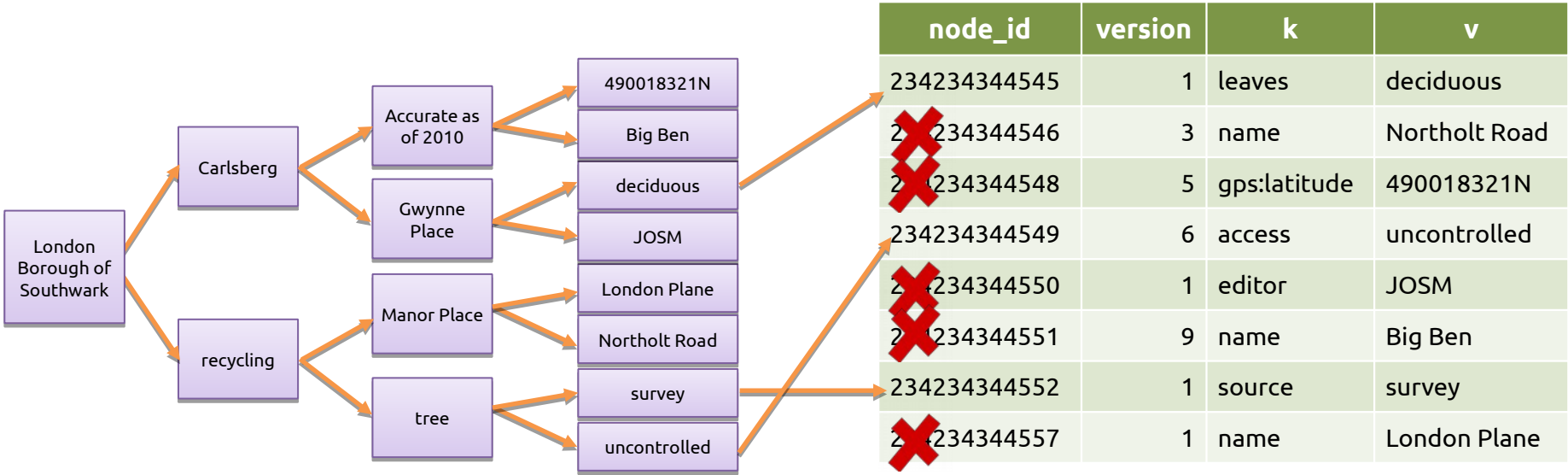
However, thanks to it we avoid a sort and a temporary table

type: index, loose index scan and covering index (1/3)



With 'type:index', all rows are read in index order (full index scan)

type: index, loose index scan and covering index (2/3)

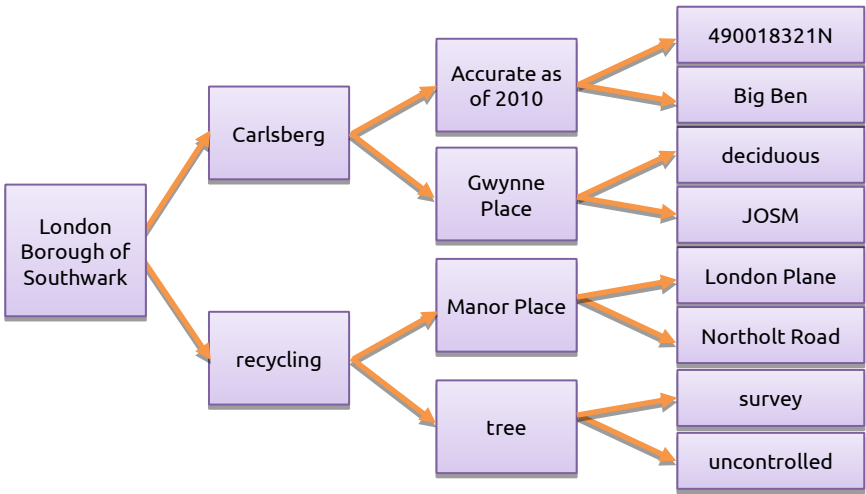


If we have in addition 'Using index for group-by' we have the loose index scan optimization

Loose Index Scan Example

```
MariaDB [dewiktionary]> EXPLAIN SELECT rev_page,  
max(rev_timestamp) FROM revision GROUP BY rev_page\G  
***** 1. row *****  
      id: 1  
  select_type: SIMPLE  
        table: revision  
         type: range  
possible_keys: NULL  
          key: page_timestamp  
     key_len: 4  
         ref: NULL  
        rows: 80980  
   Extra: Using index for group-by  
1 row in set (0.00 sec)
```

type: index, loose index scan and covering index (3/3)



node	version	v
23423	1	deciduous
23423434		Northolt Road
234234344548		490018321N
234234344549		uncontrolled
23423434455		JOSM
2342343	9	Ben
2342	1	source survey
23425	1	name London Plane

If we have in addition 'Using index' we have the covering index optimization

Covering Index Example (1/3)

```
MariaDB [dewiktionary]> ALTER TABLE revision DROP INDEX rev_page_id, drop index page_timestamp,
drop index page_user_timestamp;
Query OK, 0 rows affected (0.04 sec)
Records: 0 Duplicates: 0 Warnings: 0
```

Let's start with no indexes

```
MariaDB [dewiktionary]> EXPLAIN SELECT count(DISTINCT rev_user) FROM revision WHERE rev_page =
31579\G
```

```
***** 1. row *****
      id: 1
  select_type: SIMPLE
        table: revision
         type: ALL
possible_keys: NULL
         key: NULL
      key_len: NULL
         ref: NULL
        rows: 1376651
   Extra: Using where
1 row in set (0.00 sec)
```

```
MariaDB [dewiktionary]> SELECT count(DISTINCT rev_user) FROM revision WHERE rev_page = 31579\G
***** 1. row *****
count(DISTINCT rev_user): 1
1 row in set (3.57 sec)
```

Covering Index Example (2/3)

```
MariaDB [dewiktionary]> ALTER TABLE revision ADD INDEX revision_rev_page(rev_page);
Query OK, 0 rows affected (12.38 sec)
Records: 0 Duplicates: 0 Warnings: 0
```

```
MariaDB [dewiktionary]> EXPLAIN SELECT count(DISTINCT rev_user) FROM revision WHERE rev_page = 31579\G
```

```
***** 1. row *****
```

```
id: 1
select_type: SIMPLE
table: revision
type: ref
possible_keys: revision_rev_page
key: revision_rev_page
key_len: 4
ref: const
rows: 4038
Extra:
1 row in set (0.00 sec)
```

**Adding an index on rev_page
increases the speed due to
improved filtering**

```
MariaDB [dewiktionary]> SELECT count(DISTINCT rev_user) FROM revision WHERE rev_page = 31579\G
```

```
***** 1. row *****
```

```
count(DISTINCT rev_user): 1
1 row in set (0.04 sec)
```

Covering Index Example (3/3)

```
MariaDB [dewiktionary]> ALTER TABLE revision ADD INDEX revision_rev_page_rev_user(rev_page,
rev_user);
Query OK, 0 rows affected (13.20 sec)
Records: 0 Duplicates: 0 Warnings: 0
```

```
MariaDB [dewiktionary]> EXPLAIN SELECT count(DISTINCT rev_user) FROM revision WHERE rev_page = 31579\G
```

```
***** 1. row *****
```

```
id: 1
select_type: SIMPLE
table: revision
type: ref
possible_keys: revision_rev_page,revision_rev_page_rev_u
key: revision_rev_page_rev_user
key_len: 4
ref: const
rows: 4038
Extra: Using index
1 row in set (0.00 sec)
```

rev_page, rev_user does not increase the index selectiveness, but allows returning results directly from the index

```
MariaDB [dewiktionary]> SELECT count(DISTINCT rev_user) FROM revision WHERE rev_page = 31579\G
***** 1. row *****
count(DISTINCT rev_user): 1
1 row in set (0.00 sec)
```

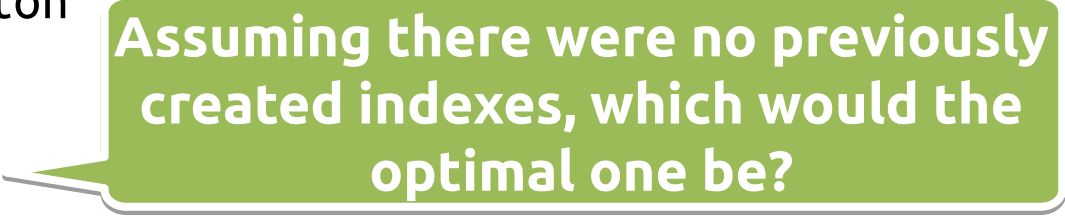
The speed difference can be huge

Query Optimization with MySQL 8.0 and MariaDB 10.3: The Basics

MULTI-COLUMN INDEXES

In many cases, conditions are applied on more than one column

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision WHERE  
rev_page = 31579 and rev_timestamp < '2008'\G  
***** 1. row *****  
      id: 1  
  select_type: SIMPLE  
        table: revision  
         type: ALL  
possible_keys: NULL  
         key: NULL  
      key_len: NULL  
         ref: NULL  
       rows: 1376651  
   Extra: Using where  
1 row in set (0.00 sec)
```



Assuming there were no previously created indexes, which would the optimal one be?

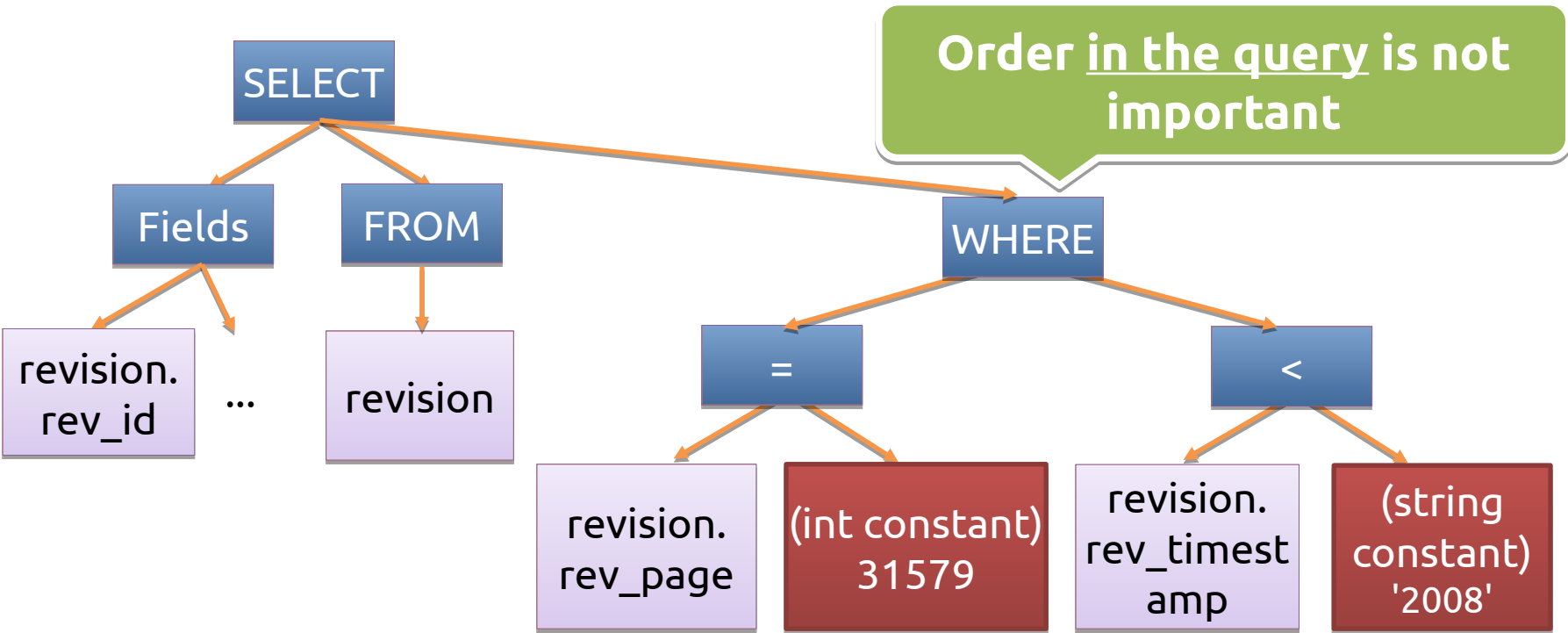
Options for indexes

- 1 index on column (rev_page)
- 1 index on column (rev_timestamp)
- 2 indexes, 1 on (rev_page) and another on (rev_timestamp)
- 1 multi-column index on (rev_page, rev_timestamp)
- 1 multi-column index on (rev_timestamp, rev_page)

**Are these last 2 different from each other?
Would it depend on the query order?**

A brief reminder about query parsing

```
SELECT * FROM revision WHERE rev_page = 31579 and rev_timestamp < '2008'
```



Index on (rev_page)

```
MariaDB [dewiktionary]> ALTER TABLE revision ADD INDEX rev_page (rev_page);
Query OK, 0 rows affected (11.87 sec)
Records: 0  Duplicates: 0  Warnings: 0
```

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision WHERE rev_page = 31579 and
rev_timestamp < '2008'\G
```

```
***** 1. row *****
      id: 1
  select_type: SIMPLE
        table: revision
         type: ref
possible_keys: rev_page
          key: rev_page
        key_len: 4
         ref: const
        rows: 4038
      Extra: Using where
1 row in set (0.00 sec)
```

Query time improves significantly with this index

Fewer rows are scanned

Adding (rev_timestamp)

```
MariaDB [dewiktionary]> ALTER TABLE revision ADD INDEX rev_timestamp  
(rev_timestamp);  
Query OK, 0 rows affected (17.01 sec)  
Records: 0 Duplicates: 0 Warnings: 0
```

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision WHERE rev_page =  
31579 and rev_timestamp < '2008'\G  
***** 1. row *****
```

```
      id: 1  
select_type: SIMPLE  
      table: revision  
      type: ref  
possible_keys: rev_page,rev_timestamp  
      key: rev_page  
      key_len: 4  
      ref: const  
      rows: 4038  
Extra: Using where  
1 row in set (0.01 sec)
```

**In general, only one index can
be used per table access**

**rev_page is preferred
over rev_timestamp**

Forcing the use of (rev_timestamp)

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision FORCE  
INDEX(rev_timestamp) WHERE rev_page = 31579 and rev_timestamp  
< '2008'\G
```

Forcing the index is worse than type:ALL!

```
id: 1  
select_type: SIMPLE  
table: revision  
type: range  
possible_keys: rev_timestamp  
key: rev_timestamp  
key_len: 14  
ref: NULL  
rows: 688325
```

**It is a range
access**

**A lot more accessed
rows**

```
Extra: Using index condition; Using where  
1 row in set (0.00 sec)
```

Adding (rev_page, rev_timestamp)

```
MariaDB [dewiktionary]> ALTER TABLE revision ADD INDEX  
rev_page_rev_timestamp(rev_page, rev_timestamp);  
Query OK, 0 rows affected (14.02 sec)  
Records: 0 Duplicates: 0 Warnings: 0
```

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision WHERE rev_page =  
31579 and rev_timestamp < '2008'\G  
***** 1. row *****
```

```
      id: 1  
  select_type: SIMPLE  
        table: revision  
         type: range  
possible_keys: rev_page,rev_timestamp,rev_page_rev_timestamp  
          key: rev_page_rev_timestamp  
     key_len: 18  
        ref: NULL  
       rows: 530  
  Extra: Using index condition  
1 row in set (0.00 sec)
```

**Reduced number of
rows scanned**

Is (rev_timestamp, rev_page) a better option?

```
MariaDB [dewiktionary]> ALTER TABLE revision ADD INDEX rev_timestamp_rev_page  
(rev_timestamp, rev_page);  
Query OK, 0 rows affected (16.80 sec)  
Records: 0 Duplicates: 0 Warnings: 0
```

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision WHERE rev_page = 31579 and  
rev_timestamp < '2008'\G
```

```
***** 1. row *****  
      id: 1  
  select_type: SIMPLE  
        table: revision  
         type: range  
possible_keys: rev_page,rev_timestamp,rev_page_rev_timestamp,rev_timestamp_rev_page  
          key: rev_page_rev_timestamp  
       key_len: 18  
         ref: NULL  
        rows: 530  
   Extra: Using index condition  
1 row in set (0.00 sec)
```

Previous index is still preferred, why?

Forcing (rev_timestamp, rev_page)

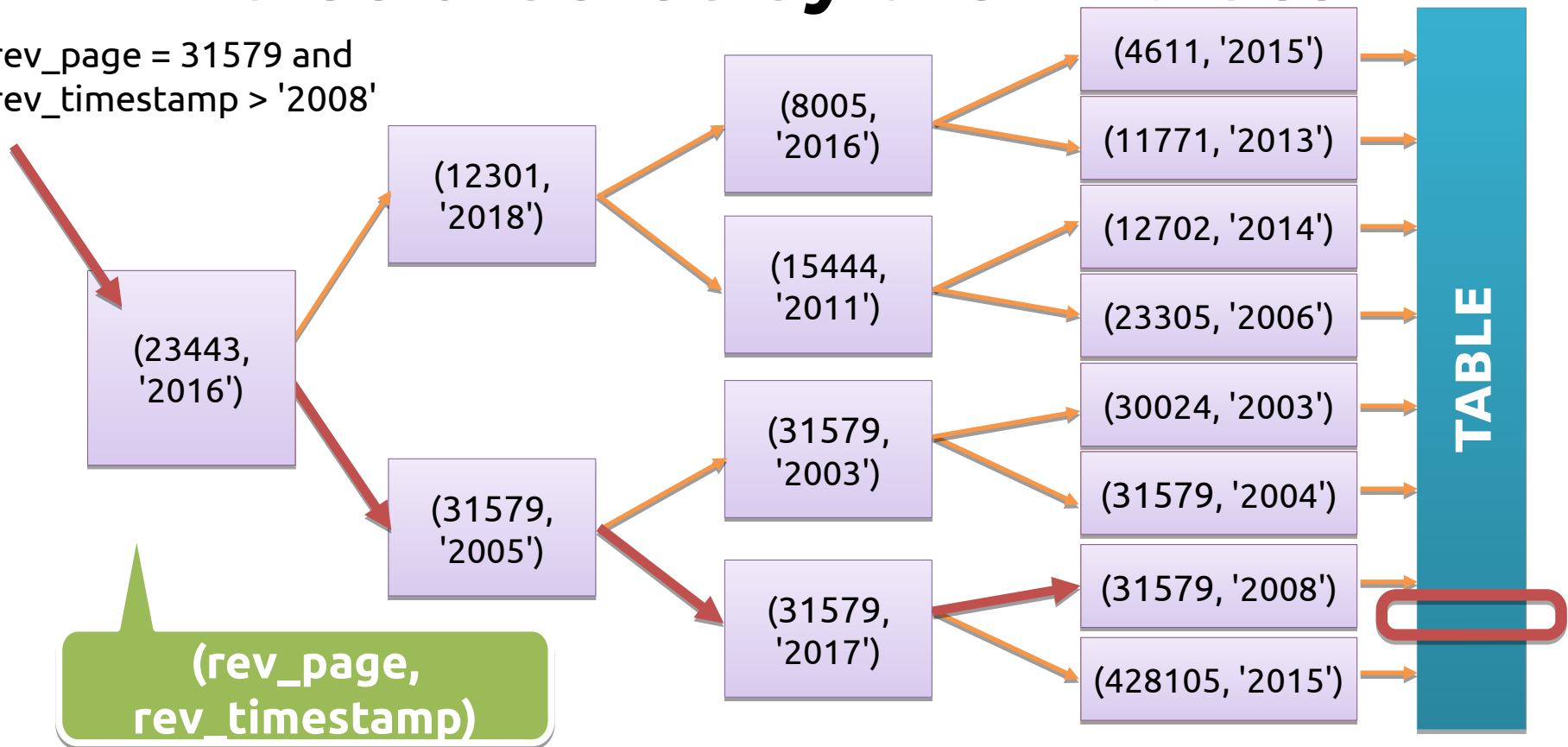
```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision FORCE
INDEX(rev_timestamp_rev_page) WHERE rev_page = 31579 and
rev_timestamp < '2008'\G
```

```
***** 1. row *****
      id: 1
  select_type: SIMPLE
        table: revision
         type: range
possible_keys: rev_timestamp_rev_page
          key: rev_timestamp_rev_page
       key_len: 18
         ref: NULL
        rows: 688325
   Extra: Using index condition
1 row in set (0.00 sec)
```

**Only the 1st column* is being
used effectively for filtering**

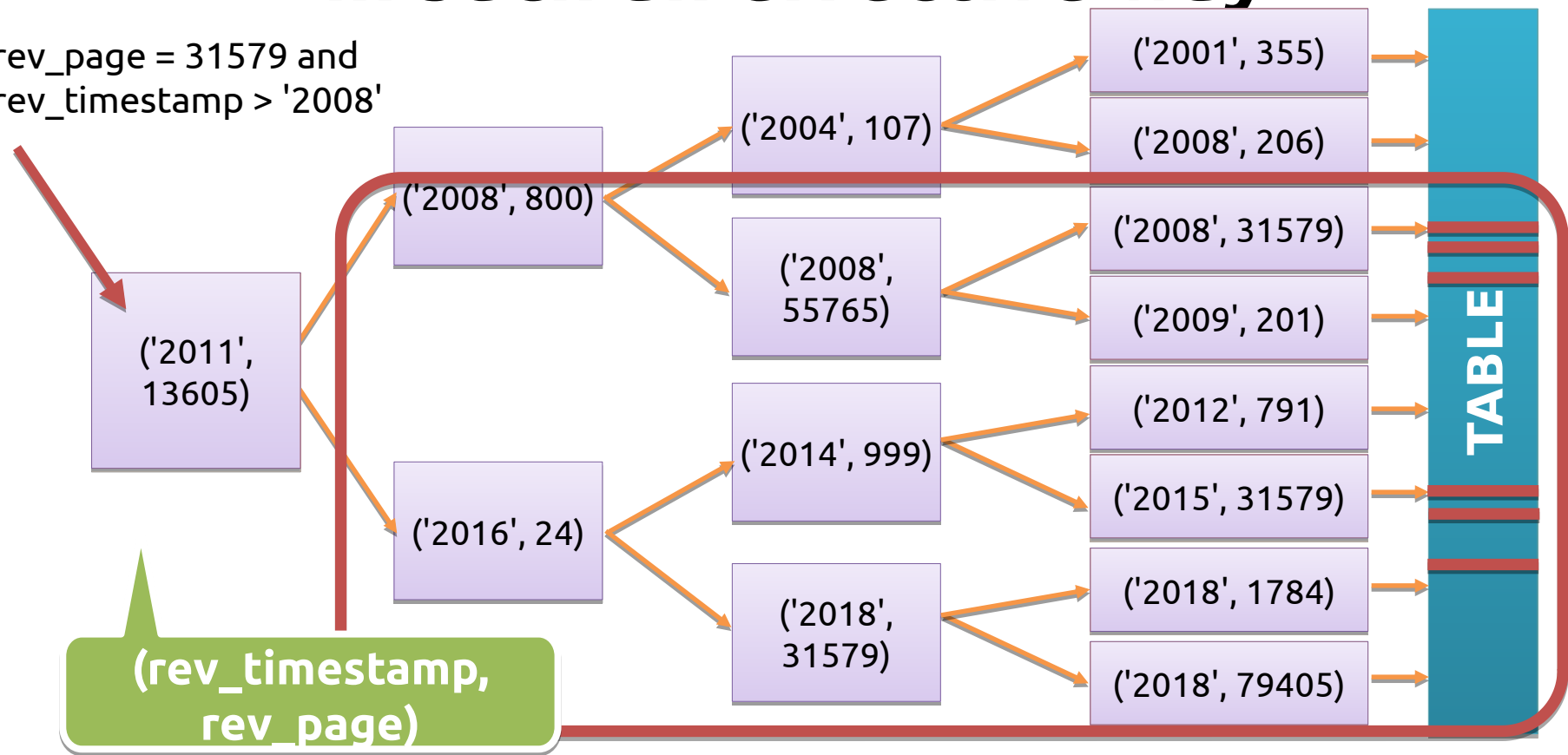
A compound index produces a single tree ordered by the 2 values

rev_page = 31579 and
rev_timestamp > '2008'



The alternative index cannot be used in such an effective way

rev_page = 31579 and
rev_timestamp > '2008'



Order and column selection

- Range access using $>$, $<$, $>=$, $<=$, BETWEEN can only be filtered once effectively, at the end of an index
- When selecting indexes, prefer columns with high cardinality (very selective)
 - The optimal index can depend on the constants used

Can we filter and sort at the same time using indexes?

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision WHERE  
rev_comment='' ORDER BY rev_timestamp ASC\G
```

```
***** 1. row *****
```

```
id: 1  
select_type: SIMPLE  
table: revision  
type: ALL  
possible_keys: NULL  
key: NULL  
key_len: NULL  
ref: NULL  
rows: 817636  
Extra: Using where; Using filesort  
1 row in set (0.00 sec)
```

**This query is slow
because a) the full
table scan**

**b) Required sort
after filtering**

Adding an index on (rev_comment, rev_timestamp)

```
MariaDB [dewiktionary]> ALTER TABLE revision ADD INDEX
rev_comment_rev_timestamp (rev_comment, rev_timestamp);
Query OK, 0 rows affected (16.19 sec)
Records: 0  Duplicates: 0  Warnings: 0
```

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision WHERE rev_comment=''
ORDER BY rev_timestamp ASC\G
```

```
***** 1. row *****
```

```
id: 1
select_type: SIMPLE
table: revision
type: ref
possible_keys: rev_comment_rev_timestamp
key: rev_comment_rev_timestamp
key_len: 769
ref: const
rows: 266462
Extra: Using index condition; Using where
1 row in set (0.00 sec)
```

Both 'type: ALL'
and 'filesort' have
disappeared

This is not always possible

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision WHERE  
rev_len > 5500 ORDER BY rev_timestamp ASC\G  
***** 1. row *****  
      id: 1  
  select_type: SIMPLE  
      table: revision  
      type: range  
possible_keys: rev_len_rev_timest  
      key: rev_len_rev_timest  
    key_len: 5  
      ref: NULL  
     rows: 38744  
  Extra: Using index condition; Using filesort  
1 row in set (0.00 sec)
```

**The range makes impossible
to use the index optimally
for the ORDER BY: either we
filter (rev_len) or sort
(rev_timestamp)**

A Strange Game. The Only Winning Move is Not to Play

```
mysql (osm) > SELECT * FROM nodes FORCE INDEX(version_idx)
WHERE version < 15 ORDER BY changeset_id;
/* type: range, Using filesort */
2859673 rows in set (30.58 sec)
```

```
mysql (osm) > SELECT * FROM nodes FORCE
INDEX(changeset_id_idx) WHERE version < 15 ORDER BY
changeset_id;
/* type: index */
2859673 rows in set (30.92 sec)
```

```
mysql (osm) > SELECT * FROM nodes WHERE version < 15 ORDER BY
changeset_id;
/* type: ALL, Using filesort */
2859673 rows in set (16.54 sec)
```

DESC Indexes (new in 8.0)

- Regular indexes could also be used to scan rows in inverse order
 - It had a penalty
 - It improves certain queries combining DESC and ASC order

```
mysql-8.0 [dewiktionary]> ALTER TABLE page ADD INDEX  
`page_redirect_namespace_len_DESC`  
(`page_is_redirect`, `page_namespace`, `page_len` DESC);  
Query OK, 0 rows affected, 1 warning (3.96 sec)  
Records: 0 Duplicates: 0 Warnings: 1
```

Descending Indexes in Action

```
mysql-8.0 [dewiktionary]> EXPLAIN SELECT * FROM page WHERE
page_is_redirect = 1 ORDER BY page_namespace, page_len DESC\G
***** 1. row *****
      id: 1
  select_type: SIMPLE
        table: page
         type: ref
possible_keys:
page_redirect_namespace_len,page_redirect_namespace_len_DESC
      key: page_redirect_namespace_len
    key_len: 1
       ref: const
      rows: 2404
  Extra: Using where
1 row in set (0.00 sec)
```



Index Condition Pushdown

```
MariaDB [dewiktionary]> ALTER TABLE revision ADD INDEX
rev_len_rev_timestamp(rev_len, rev_timestamp);
Query OK, 0 rows affected (18.11 sec)
Records: 0  Duplicates: 0  Warnings: 0
```

```
MariaDB [dewiktionary]> SET optimizer_switch='index_condition_pushdown=on'; EXPLAIN
SELECT * FROM revision WHERE rev_timestamp < '2008' AND rev_len > 10000\G
Query OK, 0 rows affected (0.00 sec)
***** 1. row *****
```

```
      id: 1
  select_type: SIMPLE
        table: revision
         type: range
possible_keys: rev_timestamp,rev_timestamp_
          key: rev_len_rev_timestamp
        key_len: 5
         ref: NULL
        rows: 180270
   Extra: Using index condition
1 row in set (0.00 sec)
```

Index condition pushdown (ICP) enables the engines to use extra parts of the index while avoiding costly row movements to and from the SQL layer

ICP optimizations

- Differences in execution time are more significant when the extra column condition is very selective (getting 5x the original performance)
- ICP is ignored when using covering Index, potentially making the performance worse

Use of "Handler_*" statistics

- They are post-execution statistics at row level
 - Unlike EXPLAIN's "rows" column, they are exact, not a guess
 - They allow comparison of query execution performance in a deterministic way, independently of the execution time

"Handler" Stats (indexed)

```
mysql> SHOW SESSION STATUS like 'Hand%';
```

Variable_name	Value
Handler_commit	80
Handler_delete	0
Handler_discover	0
Handler_external_lock	166
Handler_mrr_init	0
Handler_prepare	0
Handler_read_first	23
Handler_read_key	736212
Handler_read_last	0
Handler_read_next	22208001
Handler_read_prev	0
Handler_read_rnd	665215
Handler_read_rnd_next	14223297
Handler_rollback	0
Handler_savepoint	0
Handler_savepoint_rollback	0
Handler_update	66970
Handler_write	2869409

18 rows in set (0.00 sec)

Number of times that the first entry of an index was read. It may indicate the number of full index scans

Number of times a row has been retrieved using an index

Next row has been requested in index order (typical for index scans or ranges)

"Handler" Stats (unindexed)

```
mysql> SHOW SESSION STATUS like 'Hand%';
```

Variable_name	Value
Handler_commit	80
Handler_delete	0
Handler_discover	0
Handler_external_lock	166
Handler_mrr_init	0
Handler_prepare	0
Handler_read_first	23
Handler_read_key	736212
Handler_read_last	0
Handler_read_next	22208001
Handler_read_prev	0
Handler_read_rnd	665215
Handler_read_rnd_next	14223297
Handler_rollback	0
Handler_savepoint	0
Handler_savepoint_rollback	0
Handler_update	66970
Handler_write	2869409

18 rows in set (0.00 sec)

A row has been requested in a specific position (typical for joins or order by without indexes)

Request to read the next row in "table order" (typical for full table scans)

Insertions in SELECTS may indicate temporary tables

Comparing statistics of the previous indexes (no indexes)

MariaDB [dewiktionary]> FLUSH STATUS;
Query OK, 0 rows affected (0.00 sec)

MariaDB [dewiktionary]> SELECT * FROM revision IGNORE INDEX(rev_page, rev_timestamp, rev_page_rev_timestamp, rev_timestamp_rev_page) WHERE rev_page = 31579 and rev_timestamp < '2008';
530 rows in set (0.71 sec)

MariaDB [dewiktionary]> SHOW STATUS like 'Hand%';

Variable_name	Value
Handler_commit	1
Handler_delete	0
...	
Handler_read_first	0
Handler_read_key	0
Handler_read_last	0
Handler_read_next	0
Handler_read_prev	0
Handler_read_rnd	0
Handler_read_rnd_deleted	0
Handler_read_rnd_next	1430043
Handler_rollback	0
...	
Handler_update	0
Handler_write	0

27 rows in set (0.00 sec)

Typical result for a full table scan

Index on (rev_page)

MariaDB [dewiktionary]> SHOW STATUS like 'Hand%';

Variable_name	Value
Handler_commit	1
Handler_delete	0
...	
Handler_read_first	0
Handler_read_key	1
Handler_read_last	0
Handler_read_next	4038
Handler_read_prev	0
Handler_read_rnd	0
Handler_read_rnd_deleted	0
Handler_read_rnd_next	0
Handler_rollback	0
...	
Handler_update	0
Handler_write	0

27 rows in set (0.01 sec)

Using the index, request the first row with rev_page=31579

Then, scan them one by one in index order

Index on (rev_timestamp)

```
MariaDB [dewiktionary]> SHOW STATUS like 'Hand%';
```

Variable_name	Value
Handler_commit	1
Handler_delete	0
...	
Handler_read_first	0
Handler_read_key	1
Handler_read_last	0
Handler_read_next	199155
Handler_read_prev	0
Handler_read_rnd	0
Handler_read_rnd_deleted	0
Handler_read_rnd_next	0
Handler_rollback	0
...	

27 rows in set (0.00 sec)

ICP will be explained later, let's ignore it for now

Using the index, request the first row where rev_timestamp<2008

Then, scan them one by one in index order (more are matched)

Index on (rev_page, rev_timestamp)

```
MariaDB [dewiktionary]> SHOW STATUS like 'Hand%';
```

Variable_name	Value
Handler_commit	1
Handler_delete	0
...	
Handler_read_first	0
Handler_read_key	1
Handler_read_last	0
Handler_read_next	530
Handler_read_prev	0
Handler_read_rnd	0
Handler_read_rnd_deleted	0
Handler_read_rnd_next	0
Handler_rollback	0
...	
Handler_update	0
Handler_write	0

25 rows in set (0.00 sec)

With both conditions covered, we can find the actual first row that matches the condition using the index

Rows scanned == Rows returned

Index on (rev_timestamp, rev_page), no ICP

MariaDB [dewiktionary]> SHOW STATUS like 'Hand%';

Variable_name	Value
Handler_commit	1
Handler_delete	0
...	
Handler_read_first	0
Handler_read_key	1
Handler_read_last	0
Handler_read_next	452539
Handler_read_prev	0
Handler_read_rnd	0
Handler_read_rnd_deleted	0
Handler_read_rnd_next	0
Handler_rollback	0
...	
Handler_update	0
Handler_write	0

27 rows in set (0.00 sec)

Assuming no ICP, exact same results as with (rev_timestamp). The extra column does not help. Also, EXPLAIN's row count was very off.

Index on (rev_timestamp, rev_page), with ICP

MariaDB [dewiktionary]> SHOW STATUS like 'Hand%';

Variable_name	Value
Handler_commit	1
Handler_delete	0
...	
Handler_icp_attempts	452539
Handler_icp_match	530
...	
Handler_read_first	0
Handler_read_key	1
Handler_read_last	0
Handler_read_next	530
Handler_read_prev	0
Handler_read_rnd	0
Handler_read_rnd_deleted	0
Handler_read_rnd_next	0
Handler_rollback	0
...	
Handler_update	0
Handler_write	0

27 rows in set (0.00 sec)

ICP reduces the number of 'ENGINE API calls' significantly, although making it work more internally


Redundant Indexes

- Creating all 4 previous indexes in production is not a great idea
 - "Left-most index prefix" allows, for example (rev_page, rev_timestamp) doing everything you can do with (rev_page)
 - If two indexes have equal selectivity, MySQL chooses the shortest one

"Left-most index" Example

```

MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision FORCE
INDEX(rev_page_rev_timestamp) WHERE rev_page = 31579\G
***** 1. row *****
      id: 1
    select_type: SIMPLE
      table: revision
        type: ref
possible_keys: rev_page_rev_timestamp
         key: rev_page_rev_time
        key_len: 4
          ref: const
         rows: 4038
       Extra:
1 row in set (0.00 sec)
    
```



Only the first column is used

Duplicate Indexes

- It is very easy to create indexes with the same exact definition (same columns and ordering)
 - Set a convention for index naming (e.g `tablename_column1_column2_idx`) – MySQL does not allow 2 indexes with the same identifier
 - Since MySQL 5.6, an warning is thrown if a duplicate index is created:

Still a warning on 8.0

Duplicate index 'page_random2' defined on the table 'dewiktionary.page'. This is deprecated and will be disallowed in a future release.

pt-duplicate-index-checker

```
$ pt-duplicate-key-checker h=localhost,u=root,D=dewiktionary
[...]
# rev_timestamp is a left-prefix of rev_timestamp_rev_page
# Key definitions:
#   KEY `rev_timestamp` (`rev_timestamp`),
#   KEY `rev_timestamp_rev_page` (`rev_timestamp`,`rev_page`)
# Column types:
#   `rev_timestamp` binary(14) not null default '\0\0\0\0\0\0\0\0\0\0\0\0\0\0\0\0'
#   `rev_page` int(10) unsigned not null
# To remove this duplicate index, execute:
ALTER TABLE `dewiktionary`.`revision` DROP INDEX `rev_timestamp`;

# rev_page is a left-prefix of rev_page_rev_timestamp
# Key definitions:
#   KEY `rev_page` (`rev_page`),
#   KEY `rev_page_rev_timestamp` (`rev_page`,`rev_timestamp`),
# Column types:
#   `rev_page` int(10) unsigned not null
#   `rev_timestamp` binary(14) not null default '\0\0\0\0\0\0\0\0\0\0\0\0\0\0\0\0'
# To remove this duplicate index, execute:
ALTER TABLE `dewiktionary`.`revision` DROP INDEX `rev_page`;

# ##### Summary of indexes #####
# Size Duplicate Indexes      47836942
# Total Duplicate Indexes     4
# Total Indexes                285
```

Simple tool to check
redundant and
duplicate indexes

Invisible Indexes (new in 8.0)

- Before removing a (believed) redundant or duplicate index, make it “invisible” to prevent long-running outages:

```
mysql-8.0 [dewiktionary]> ALTER TABLE page ALTER INDEX  
page_len INVISIBLE;  
Query OK, 0 rows affected (0.04 sec)  
Records: 0 Duplicates: 0 Warnings: 0
```

```
mysql-8.0 [dewiktionary]> SELECT * FROM page FORCE  
INDEX(page_len);  
ERROR 1176 (42000): Key 'page_len' doesn't exist in table  
'page'
```

"OR"-style conditions over the same column

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision WHERE  
rev_page = 790 OR rev_page = 795 OR rev_page = 1024\G  
***** 1 row *****  
      id: 1  
  select_type: SIMPLE  
        table: revision  
         type: range  
possible_keys: rev_page,rev_page,rev_page  
          key: rev_page  
       key_len: 4  
         ref: NULL  
        rows: 201  
   Extra: Using index condition; Using where  
1 row in set (0.01 sec)
```

Equivalent to:
SELECT * FROM revision WHERE
rev_page IN (790, 795, 1024)

Handlers on "IN" / "OR" conditions over the same column

Variable_name	Value
Handler_commit	1
Handler_delete	0
...	
Handler_prepare	0
Handler_read_first	0
Handler_read_key	3
Handler_read_last	0
Handler_read_next	201
Handler_read_prev	0
Handler_read_rnd	0
Handler_read_rnd_deleted	0
Handler_read_rnd_next	0
Handler_rollback	0
...	
Handler_update	0
Handler_write	0

Despite identifying themselves as “range”s, the execution is slightly different, one index dive (similar to a ref) is done per value. This can be an issue in conditions with thousands of items.

25 rows in set (0.00 sec)

"OR"-style conditions over different columns

- We cannot use a single index efficiently for both conditions
 - We can scan both conditions separately and mix the results, discarding duplicates
 - Or use an index for one condition and not for the other
 - Index merge allows the use of two indexes for a single table access simultaneously

Index Merge Example

```

MariaDB [dewiktionary]> EXPLAIN SELECT * FROM revision WHERE
rev_page = 790 or rev_timestamp < '2004'\G
***** 1. row *****
      id: 1
  select_type: SIMPLE
        table: revision
        type: index_merge
possible_keys: rev_page,rev_timestamp,rev_page_rev_timestamp,
rev_timestamp_rev_page
          key: rev_page,rev_timestamp
      key_len: 4,14
         ref: NULL
        rows: 190
   Extra: Using sort_union(rev_page,rev_timestamp); Using
where
1 row in set (0.00 sec)
    
```

Both indexes are used, then combined using the "union" operation

INDEX_MERGE Issues

- Sometimes it is faster to execute the sentence using UNION:
 - This is especially true with (UNION ALL) since MySQL 5.7, if you do not care or expect duplicates
- There are also intersection merges, but multi-column indexes are preferred

Disabling optimizer features (I)

- The `optimizer_switch` variable allows enabling and disabling globally or per session many query optimizer features:

```
MariaDB [nlwiktionary]> SHOW VARIABLES like 'optimizer_switch'\G
***** 1. row *****
Variable_name: optimizer_switch
Value:
index_merge=on,index_merge_union=on,index_merge_sort_union=on,index_merge_intersection=
on,index_merge_sort_intersection=off,engine_condition_pushdown=off,index_condition_push
down=on,derived_merge=on,derived_with_keys=on,firstmatch=on,loosescan=on,materializatio
n=on,in_to_exists=on,semijoin=on,partial_match_rowid_merge=on,partial_match_table_scan=
on,subquery_cache=on,mrr=off,mrr_cost_based=off,mrr_sort_keys=off,outer_join_with_cache
=on,semijoin_with_cache=on,join_cache_incremental=on,join_cache_hashed=on,join_cache_bk
a=on,optimize_join_buffer_size=off,table_elimination=on,extended_keys=on,exists_to_in=on
n,orderby_uses_equalities=on,condition_pushdown_for_derived=on,split_materialized=on
1 row in set (0.00 sec)
```

Disabling optimizer features (II)

```
MariaDB [dewiktionary]> SET optimizer_switch='index_merge_sort_union=off';
Query OK, 0 rows affected (0.00 sec)
```

```
MariaDB [dewiktionary]> EXPLAIN SELECT * FROM
790 or rev_timestamp < '2004'\G
```

```
***** 1. row *****
```

```
id: 1
```

```
select_type: SIMPLE
```

```
table: revision
```

```
type: ALL
```

```
possible_keys:
```

```
rev_page,rev_timestamp,rev_page_rev_timestamp,rev_timestamp_rev_page
```

```
key: NULL
```

```
key_len: NULL
```

```
ref: NULL
```

```
rows: 1376651
```

```
Extra: Using where
```

```
1 row in set (0.00 sec)
```

This will only have effect for the current session.

What happens if we have two ranges?

- As seen previously, we cannot use efficiently two range types on the same table access. Alternatives:
 - Use only one index for the most selective column
 - Use index condition pushdown to get an advantage
 - Change one of the two ranges into a discrete "IN" comparison/bucketize with a new column
 - Use quadtrees or R-TREES (spatial indexing)

Example of Bucketizing (I)

```
MariaDB [dewiktionary]> ALTER TABLE revision DROP INDEX
rev_len_rev_timestamp;
MariaDB [dewiktionary]> EXPLAIN SELECT count(*) FROM revision
WHERE rev_timestamp < '2008' AND rev_len > 5500\G
***** 1. row *****
      id: 1
  select_type: SIMPLE
        table: revision
         type: ALL
possible_keys: rev_timestamp,rev_timestamp_rev_page
           key: NULL
        key_len: NULL
         ref: NULL
         rows: 1376651
      Extra: Using where
1 row in set (0.00 sec)
```

Looks like only an index on
(rev_timestamp) or (rev_len)
would be useful as we have 2
ranges.

Example of Bucketizing (II)

```
MariaDB [dewiktionary]> ALTER TABLE revision ADD  
rev_len_cat int;  
Query OK, 0 rows affected (0.01 sec)  
Records: 0 Duplicates: 0 Warnings: 0
```

**8.0/10.3
instant
alter-nice!**

```
MariaDB [dewiktionary]> UPDATE revision set rev_len_cat =  
IF(rev_len < 10000, rev_len div 1000, 10);  
Query OK, 1430042 rows affected (21.18 sec)  
Rows matched: 1430042 Changed: 1430042 Warnings: 0
```

```
MariaDB [dewiktionary]> ALTER TABLE revision ADD INDEX  
rev_len_cat_rev_timestamp (rev_len_cat, rev_timestamp);  
Query OK, 0 rows affected (7.77 sec)  
Records: 0 Duplicates: 0 Warnings: 0
```


Example of Bucketizing (III)

```
MariaDB [dewiktionary]> EXPLAIN SELECT count(*) FROM revision WHERE
rev_timestamp < '2008' AND rev_len > 5500 AND rev_len_cat IN (5, 6,
7, 8, 9, 10)\G
```

```
***** 1. row *****
      id: 1
  select_type: SIMPLE
        table: revision
         type: range
possible_keys:
rev_timestamp,rev_timestamp_rev_page,rev_len_cat_rev_timestamp
      key: rev_len_cat_rev_timestamp
     key_len: 19
        ref: NULL
       rows: 77601
  Extra: Using where
1 row in set (0.00 sec)
```

**We did some transformations
to both the structure and the
query.**

Even better: Functional Indexes

(I)

```
MariaDB [dewiktionary]> ALTER TABLE revision
DROP INDEX rev_len_cat_rev_timestamp,
DROP COLUMN rev_len_cat,
ADD COLUMN rev_len_cat_virtual int as (IF(rev_len < 10000,
rev_len div 1000, 10)) VIRTUAL,
ADD INDEX rev_len_cat_virtual_rev_timestamp
(rev_len_cat_virtual, rev_timestamp);
Query OK, 1430042 rows affected (1 min 42.35 sec)
Records: 1430042 Duplicates: 0 Warnings: 0
```

Indexing
virtual
columns,
new in 10.3

Storing the extra content
may be a waste of resources,
lets make it a virtual column

Even better: Functional Indexes

(II)

```
MariaDB [dewiktionary]> EXPLAIN SELECT count(*) FROM revision WHERE  
rev_timestamp < '2008' AND rev_len > 5500 AND rev_len_cat_virtual IN  
(5, 6, 7, 8, 9, 10)\G
```

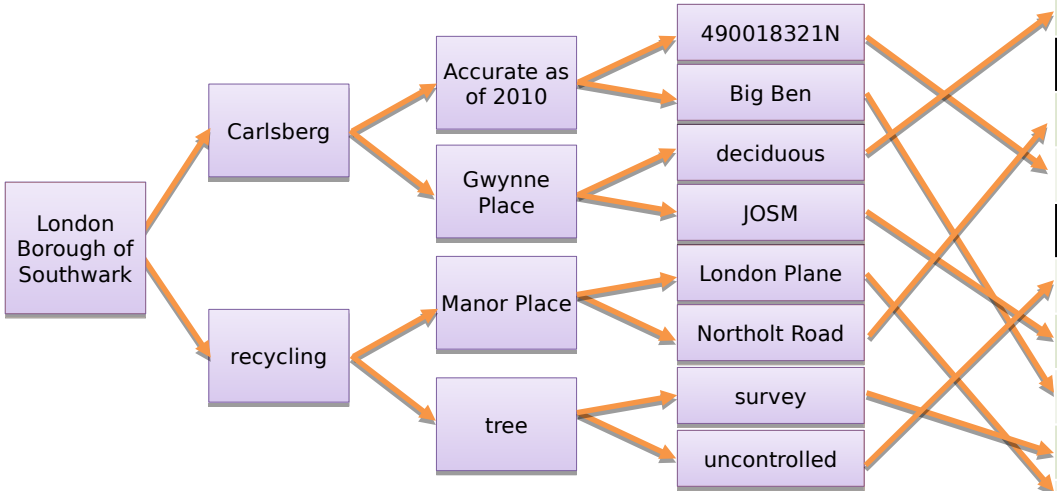
```
***** 1. row *****  
      id: 1  
  select_type: SIMPLE  
        table: revision  
        type: range  
possible_keys:  
rev_timestamp,rev_timestamp_rev_page,rev_len_cat_virtual_rev_timestamp  
      key: rev_len_cat_virtual_rev_timestamp  
    key_len: 19  
       ref: NULL  
     rows: 78649  
  Extra: Using where  
1 row in set (0.01 sec)
```

**We did some transformations
to both the structure and the
query.**

MyISAM Internals

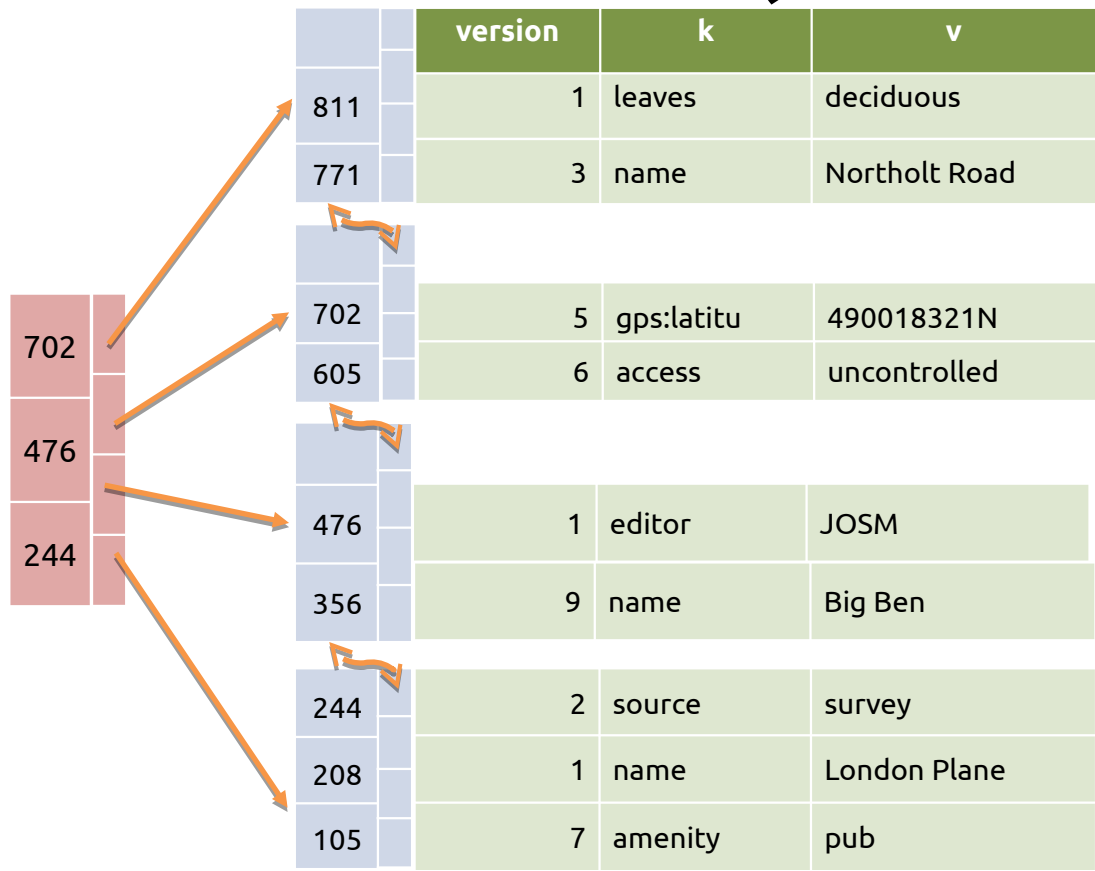
Index (part of revision.MYI)

Data (revision.MYD)



node_id	versio	k	v
234234344545	1	leaves	deciduous
(empty row)			
234234344548	5	gps:latitude	490018321N
234234344549	6	access	uncontrolled
(empty row)			
234234344551	9	name	Big Ben
234234344552	1	source	survey
234234344557	1	name	London Plane
234234344552	2	source	survey
234234344557	2	name	London Plane

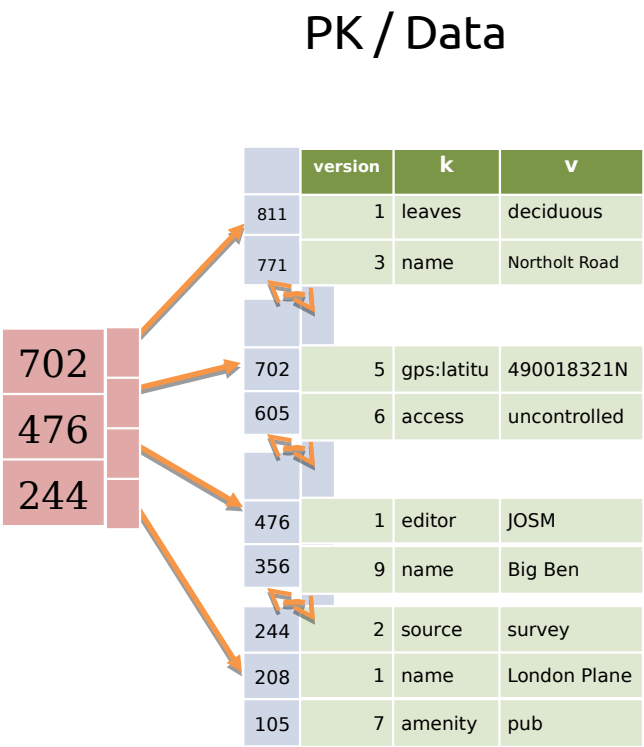
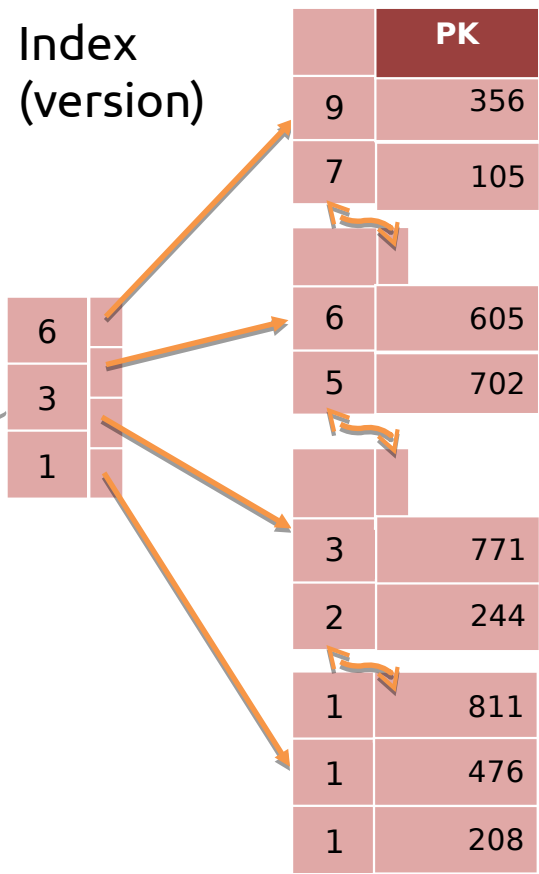
InnoDB Internals (PRIMARY)



Data clustered always using the primary key

InnoDB Internals (Secondary)

Secondary indexes contain the primary key value



Consequences of using InnoDB (I)

- Every table should have a primary key
 - If one is not defined, MySQL will choose an available NOT NULL unique key
 - If that is not possible, an internal 6-byte row identifier will be generated (not user-accessible)

Consequences of using InnoDB (II)

- Inserting in primary key order is much faster
 - Less fragmentation/page-split
 - Usage of "batch" mode, improving insert speed
- Using auto-increment keys as primary keys can be a good idea for InnoDB

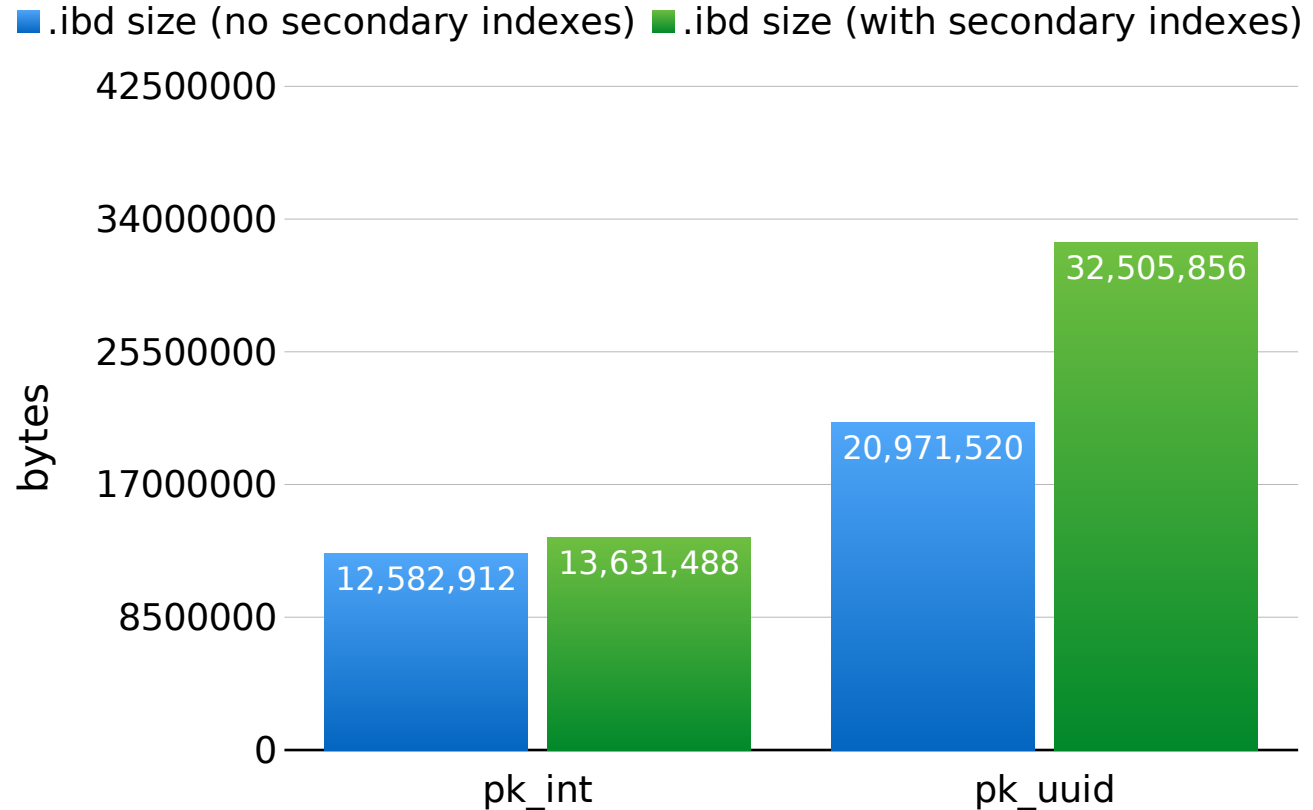
Consequences of using InnoDB (III)

- A very long primary key may increment substantially the size of secondary keys
 - Int or bigint types are recommended instead of UUIDs or other long strings

Differences in size

```
mysql (osm) > CREATE TABLE
pk_int (id int PRIMARY KEY
auto_increment,
a int,
b int,
c int,
d int);
Query OK, 0 rows affected
(0.16 sec)
```

```
mysql (osm) > CREATE TABLE
pk_uuid (id char(36)
PRIMARY KEY,
a int,
b int,
c int,
d int);
Query OK, 0 rows affected
(0.04 sec)
```



Extended primary key optimization

- As the primary key is part of all secondary keys, this can be used “for free”:
 - For row filtering (since MySQL 5.6)
 - To return results in primary key order
 - To avoid reading data from the table (covering index)

Extended Primary Key Example

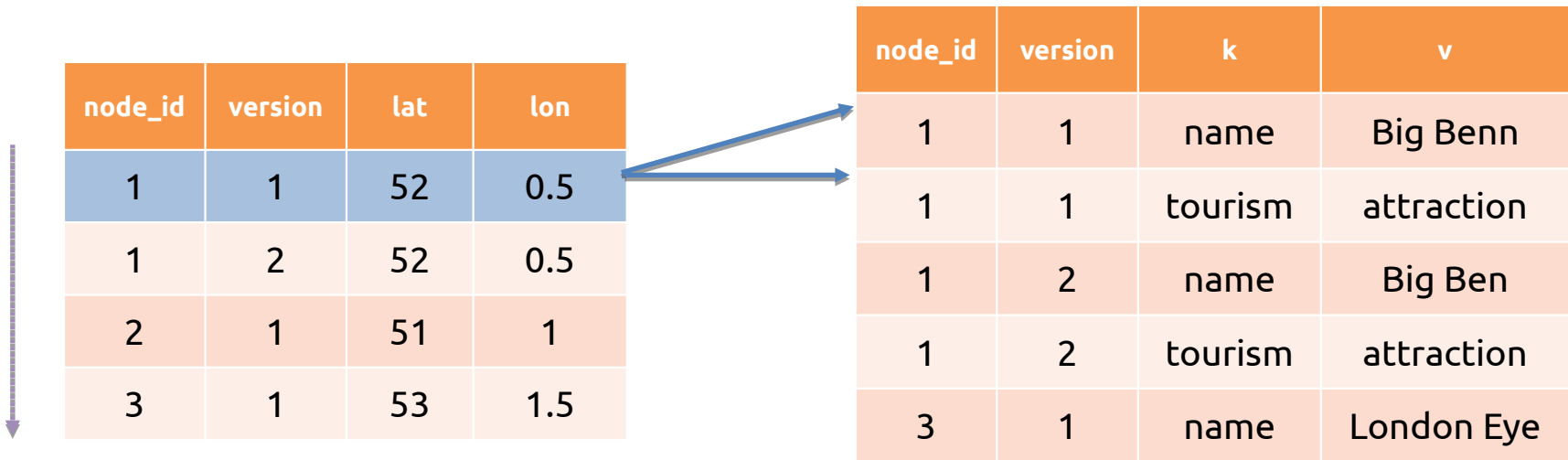
```
mysql (osm) > EXPLAIN SELECT node_id FROM nodes WHERE
changeset_id = 24284 and node_id <> 146472942\G
***** 1. row *****
      id: 1
  select_type: SIMPLE
        table: nodes
         type: range
possible_keys: PRIMARY,changeset_id_idx
          key: changeset_id_idx
       key_len: 16
         ref: NULL
        rows: 50
   Extra: Using where; Using index
1 row in set (0.07 sec)
```

Query Optimization with MySQL 8.0 and MariaDB 10.3: The Basics

JOINS

(Block) Nested Loop Join

- Until MySQL 5.5 there was only one algorithm to execute a JOIN:



Extra Access type: eq_ref

```
mysql (osm) > EXPLAIN SELECT * FROM nodes JOIN node_tags
                USING(node_id, version) WHERE node_tags.v= 'Big Ben'\G
***** 1. row *****
      id: 1
  select_type: SIMPLE
        table: node_tags
         type: ref
possible_keys: PRIMARY,v_idx
          key: v_idx
        key_len: 767
         ref: const
         rows: 1
      Extra: Using where; Using index
***** 2. row *****
      id: 1
  select_type: SIMPLE
        table: nodes
         type: eq_ref
possible_keys: PRIMARY,version_idx
          key: PRIMARY
        key_len: 16
         ref: osm.node_tags.node_id,osm.node_tags.version
         rows: 1
      Extra: NULL
2 rows in set (0.00 sec)
```

eq_ref is similar to ref, but allows faster JOINS because, by using a unique key, it only has to search one row for each previous result

JOIN Optimization

- Two main goals:
 - Perform an effective filtering on each table access, if possible using indexes
 - Perform the access in the most efficient table order
- When joining 3 or more tables in a star schema, the "covering index" strategy can have a huge impact

Example: optimize this JOIN (I)

```
SELECT n.node_id, n.latitude, n.longitude
FROM way_nodes w_n
JOIN way_tags w_t
  ON w_n.way_id = w_t.way_id and
    w_n.version = w_t.version
JOIN nodes n
  ON w_n.node_id = n.node_id
JOIN node_tags n_t
  ON n.node_id = n_t.node_id and
    n.version = n_t.version
WHERE w_t.k = 'building' and
      n_t.k = 'entrance' and
      n_t.v = 'main';
```



**We start without
secondary indexes**

Example: optimize this JOIN (II)


```
***** 1. row *****
id: 1
select_type: SIMPLE
table: w_t
type: index
possible_keys: PRIMARY
key: PRIMARY
key_len: 783
ref: NULL
rows: 1335702
Extra: Using where; Using index
***** 2. row *****
id: 1
select_type: SIMPLE
table: w_n
type: ref
possible_keys: PRIMARY
key: PRIMARY
key_len: 16
ref: osm.w_t.way_id,osm.w_t.version
rows: 3
Extra: NULL

***** 3. row *****
id: 1
select_type: SIMPLE
table: n_t
type: ref
possible_keys: PRIMARY
key: PRIMARY
key_len: 8
ref: osm.w_n.node_id
rows: 1
Extra: Using where
***** 4. row *****
id: 1
select_type: SIMPLE
table: n
type: eq_ref
possible_keys: PRIMARY
key: PRIMARY
key_len: 16
ref: osm.w_n.node_id,osm.n_t.version
rows: 1
Extra: Using index
4 rows in set (0.01 sec)
```

```
mysql (osm) > SELECT ...
858 rows in set (9.00 sec)
```

Example: optimize this JOIN (III)

```
mysql (osm) > ALTER TABLE way_tags  
               ADD INDEX k_idx(k);  
Query OK, 0 rows affected (4.80 sec)  
Records: 0 Duplicates: 0 Warnings: 0
```



Creating an index
on way_tags.k

Example: optimize this JOIN (IV)

```
***** 1. row *****
      id: 1
select_type: SIMPLE
      table: w_t
      type: ref
possible_keys: PRIMARY,k_idx
      key: k_idx
     key_len: 767
       ref: const
      rows: 452274
  Extra: Using where; Using index
***** 2. row *****
      id: 1
select_type: SIMPLE
      table: w_n
      type: ref
possible_keys: PRIMARY
      key: PRIMARY
     key_len: 16
       ref: osm.w_t.way_id,osm.w_t.version
      rows: 3
  Extra: NULL
```

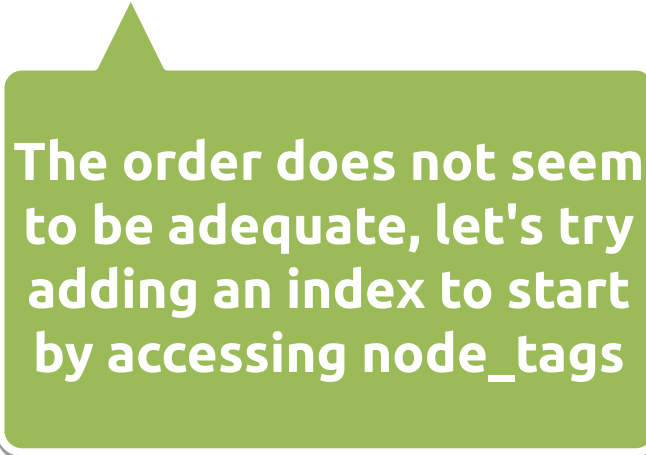
```
***** 3. row *****
      id: 1
select_type: SIMPLE
      table: n_t
      type: ref
possible_keys: PRIMARY
      key: PRIMARY
     key_len: 8
       ref: osm.w_n.node_id
      rows: 1
  Extra: Using where
***** 4. row *****
      id: 1
select_type: SIMPLE
      table: n
      type: eq_ref
possible_keys: PRIMARY
      key: PRIMARY
     key_len: 16
       ref: osm.w_n.node_id,osm.n_t.version
      rows: 1
  Extra: NULL
4 rows in set (0.00 sec)
```

It seems like the index is not very useful

```
mysql (osm) > SELECT ...
858 rows in set (8.58 sec)
```

Example: optimize this JOIN (V)

```
mysql (osm) > ALTER TABLE node_tags  
               ADD INDEX k_idx(k);  
Query OK, 0 rows affected (2.82 sec)  
Records: 0 Duplicates: 0 Warnings: 0
```



The order does not seem to be adequate, let's try adding an index to start by accessing node_tags

Example: optimize this JOIN (VI)

```

***** 1. row *****
      id: 1
    select_type: SIMPLE
      table: w_t
      type: ref
possible_keys: PRIMARY,k_idx
      key: k_idx
     key_len: 767
       ref: const
      rows: 452274
    Extra: Using where; Using index
***** 2. row *****
      id: 1
    select_type: SIMPLE
      table: w_n
      type: ref
possible_keys: PRIMARY
      key: PRIMARY
     key_len: 16
       ref: osm.w_t.way_id,osm.w_t.version
      rows: 3
    Extra: NULL

***** 3. row *****
      id: 1
    select_type: SIMPLE
      table: n_t
      type: ref
possible_keys: PRIMARY,k_idx
      key: PRIMARY
     key_len: 8
       ref: osm.w_n.node_id
      rows: 1
    Extra: Using where
***** 4. row *****
      id: 1
    select_type: SIMPLE
      table: n
      type: eq_ref
possible_keys: PRIMARY
      key: PRIMARY
     key_len: 16
       ref: osm.w_n.node_id,osm.n_t.version
      rows: 1
    Extra: NULL
4 rows in set (0.00 sec)

```

It keeps using the wrong order, even if we delete the w_t.k_idx index

```

mysql (osm) > SELECT ...
858 rows in set (7.33 sec)

```

Example: optimize this JOIN (VII)

```
SELECT STRAIGHT_JOIN n.node_id, n.latitude, n.longitude
  FROM node_tags n_t
  JOIN nodes n
    ON n.node_id = n_t.node_id and
       n.version = n_t.version
  JOIN way_nodes w_n
    ON w_n.node_id = n.node_id
  JOIN way_tags w_t
    ON w_n.way_id = w_t.way_id and
       w_n.version = w_t.version
 WHERE w_t.k = 'building' and
        n_t.k = 'entrance' and
        n_t.v = 'main';
```



Let's see why rewriting it
into this query

Example: optimize this JOIN (VIII)

```
***** 1. row *****
  id: 1
  select_type: SIMPLE
  table: n_t
  type: ref
  possible_keys: PRIMARY,k_idx
  key: k_idx
  key_len: 767
  ref: const
  rows: 2390
  Extra: Using index condition; Using where
***** 2. row *****
  id: 1
  select_type: SIMPLE
  table: n
  type: eq_ref
  possible_keys: PRIMARY
  key: PRIMARY
  key_len: 16
  ref: osm.n_t.node_id,osm.n_t.version
  rows: 1
  Extra: NULL
***** 3. row *****
  id: 1
  select_type: SIMPLE
  table: w_n
  type: ALL
  possible_keys: PRIMARY
  key: NULL
  key_len: NULL
  ref: NULL
  rows: 3597858
  Extra: Using where; Using join buffer (Block
Nested Loop)
***** 4. row *****
  id: 1
  select_type: SIMPLE
  table: w_t
  type: eq_ref
  possible_keys: PRIMARY
  key: PRIMARY
  key_len: 783
  ref: osm.w_n.way_id,osm.w_n.version,const
  rows: 1
  Extra: Using where; Using index
4 rows in set (0.00 sec)
```

There is no index on w_n that would allow efficient access

Example: optimize this JOIN (IX)

```
mysql (osm) > ALTER TABLE way_nodes  
               ADD INDEX node_id_idx(node_id);  
Query OK, 0 rows affected (17.77 sec)  
Records: 0 Duplicates: 0 Warnings: 0
```

Example: optimize this JOIN (X)

***** 1. row *****

id: 1
select_type: SIMPLE
table: n_t
type: ref
possible_keys: PRIMARY,k_idx
key: k_idx
key_len: 767
ref: const
rows: 2390

Now it starts with
the right table
(without
STRAIGHT_JOIN)

Extra: Using index condition; Using where

***** 2. row *****

id: 1
select_type: SIMPLE
table: n
type: eq_ref
possible_keys: PRIMARY
key: PRIMARY
key_len: 16
ref: osm.n_t.node_id,osm.n_t.version
rows: 1
Extra: NULL

***** 3. row *****

id: 1
select_type: SIMPLE
table: w_n
type: ref
possible_keys: PRIMARY,node_id_idx
key: node_id_idx
key_len: 8
ref: osm.n_t.node_id
rows: 1

Extra: Using index

***** 4. row *****

id: 1
select_type: SIMPLE
table: w_t
type: eq_ref
possible_keys: PRIMARY
key: PRIMARY
key_len: 783
ref: osm.w_n.way_id,osm.w_n.version,const
rows: 1

Extra: Using where; Using index

4 rows in set (0.04 sec)

mysql (osm) > SELECT ...
858 rows in set (0.73 sec)

Example: optimize this JOIN (XI)

```
***** 1. row *****
id: 1
select_type: SIMPLE
table: n_t
type: ref
possible_keys: PRIMARY,k_idx,k_v_i
key: k_v_idx
key_len: 1534
ref: const,const
rows: 900
Extra: Using where; Using index
***** 2. row *****
id: 1
select_type: SIMPLE
table: n
type: eq_ref
possible_keys: PRIMARY
key: PRIMARY
key_len: 16
ref: osm.n_t.node_id,osm.n_t.version
rows: 1
Extra: NULL

***** 3. row *****
id: 1
select_type: SIMPLE
table: w_n
type: ref
possible_keys: PRIMARY,node_id_idx
key: node_id_idx
key_len: 8
ref: osm.n_t.node_id
rows: 1
Extra: Using index
***** 4. row *****
id: 1
select_type: SIMPLE
table: w_t
type: eq_ref
possible_keys: PRIMARY
key: PRIMARY
key_len: 783
ref: osm.w_n.way_id,osm.w_n.version,const
rows: 1
Extra: Using where; Using index
4 rows in set (0.00 sec)
```

An index on (k,v) is even better

```
mysql (osm) > SELECT ...
858 rows in set (0.02 sec)
```

“New” JOIN methods

- MySQL 5.6 added:
 - Batch Key Access
- MariaDB has since 5.3:
 - Batch Key Access
 - Hash Joins
 - Slightly modified versions of the above ones (with “incremental” buffers to join 3 or more tables)

Multi-range read

- This optimization orders results obtained from a secondary key in primary key/physical order before accessing the rows
 - It may help execution time of queries when disk-bound
 - It requires tuning of the `read_rnd_buffer_size` (size of the buffer used for ordering the results)
- BKA JOINS are based on the mrr optimization

MRR Example (I)

```
mysql (osm) > EXPLAIN SELECT * FROM nodes WHERE timestamp >= '2013-07-01 00:00:00' AND timestamp < '2014-01-01 00:00:00'\G
```

```
***** 1. row *****
```

```
id: 1
```

```
select_type: SIMPLE
```

```
table: nodes
```

```
type: range
```

```
possible_keys: nodes_timestamp_idx
```

```
key: nodes_timestamp_idx
```

```
key_len: 5
```

```
ref: NULL
```

```
rows: 429684
```

```
Extra: Using index condition; Using MRR
```

```
1 row in set (0.02 sec)
```

MRR example (II)

[restart]

```
mysql> SET optimizer_switch='mrr=off';
```

```
mysql> SELECT * FROM nodes WHERE timestamp >= '2013-07-01 00:00:00' AND timestamp < '2014-01-01 00:00:00';
```

205617 rows in set (5.16 sec)

```
mysql> SELECT * FROM nodes WHERE timestamp >= '2013-07-01 00:00:00' AND timestamp < '2014-01-01 00:00:00';
```

205617 rows in set (0.60 sec)

[restart]

```
mysql> SET read_rnd_buffer_size=50 * 1024 * 1024;
```

```
mysql> SELECT * FROM nodes WHERE timestamp >= '2013-07-01 00:00:00' AND timestamp < '2014-01-01 00:00:00';
```

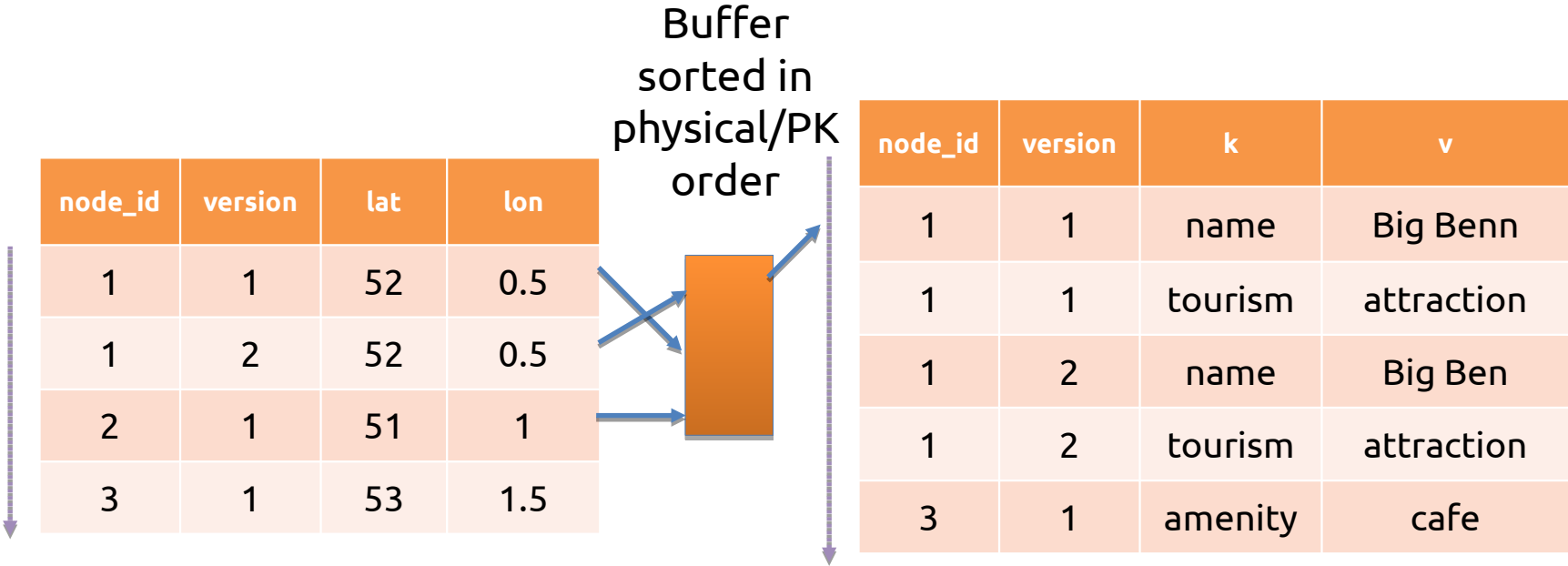
205617 rows in set (2.39 sec)

```
mysql> SELECT * FROM nodes WHERE timestamp >= '2013-07-01 00:00:00' AND timestamp < '2014-01-01 00:00:00';
```

205617 rows in set (0.73 sec)

“Cold” results are significantly better with mrr (but it can impact negatively, too)

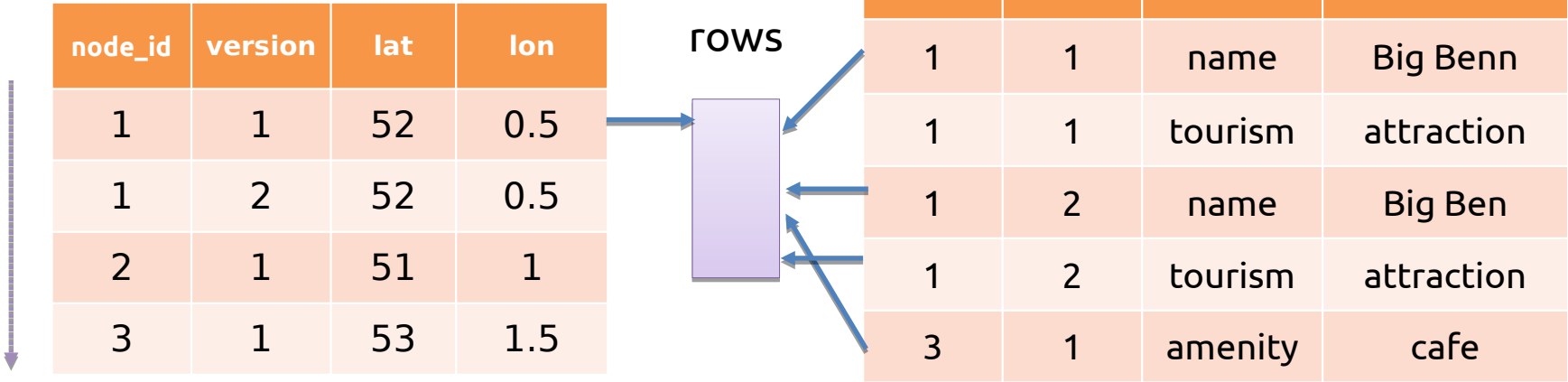
Batch Key Access



Hash Joins

- Only work for equi-joins

Hash table
for faster
access to
rows



MySQL Configuration

- BKA requires changes of default optimizer configuration:

```
mysql-8.0 (osm) > SET optimizer_switch= 'mrr=on';
```

```
mysql-8.0 (osm) > SET optimizer_switch= 'mrr_cost_based=off';
```

```
mysql-8.0 (osm) > SET optimizer_switch= 'batch_key_access=on';
```

- Additionally, configuring the `join_buffer_size` adequately

MariaDB configuration

```
mariadb-10.3 (osm) > SET optimizer_switch = 'join_cache_incremental=on';  
mariadb-10.3 (osm) > SET optimizer_switch = 'join_cache_hashed=on';  
mariadb-10.3 (osm) > SET optimizer_switch = 'join_cache_bka=on';
```

- Enabled by default

```
mariadb-10.3 (osm) > SET join_cache_level = 3 (for hash joins)  
mariadb-10.3 (osm) > SET join_cache_level = 5 (for BKA)
```

- Also, configure `join_buffer_size` appropriately.
- Hash joins, like BKA, are highly dependent on disk-bound DBs to be effective due to the extra overhead

Nested Loop Join (cold buffers buffer_pool=100MB, join_buffer=4M)

```
mariadb-10.3 (osm) > EXPLAIN SELECT changeset_id,
count(*) FROM changesets JOIN nodes on
changesets.id = nodes.changeset_id GROUP BY
visible\G
***** 1. row *****
      id: 1
  select_type: SIMPLE
        table: changesets
        type: index
possible_keys: PRIMARY
         key: PRIMARY
      key_len: 8
         ref: NULL
        rows: 69115
  Extra: Using index; Using temporary; Using
filesort
```

```
***** 2. row *****
      id: 1
  select_type: SIMPLE
        table: nodes
        type: ref
possible_keys: changeset_id
         key: changeset_id
      key_len: 8
         ref: osm.changesets.id
        rows: 19
  Extra:
2 rows in set (0.00 sec)
```

```
mariadb-10.3 (osm) > SELECT visible, count(*) FROM
changesets JOIN nodes on changesets.id =
nodes.changeset_id GROUP BY visible;
+-----+-----+
| visible | count(*) |
+-----+-----+
|      1 | 2865312 |
+-----+-----+
1 row in set (32.86 sec)
```

Hash Join (cold buffers, buffer_pool=100M, join_buffer=4M)

```
mariadb-10.3 (osm) > EXPLAIN SELECT
changeset_id, count(*) FROM changesets JOIN
nodes on changesets.id = nodes.changeset_id
GROUP BY visible\G
***** 1. row *****
      id: 1
  select_type: SIMPLE
        table: changesets
        type: index
possible_keys: PRIMARY
         key: PRIMARY
      key_len: 8
         ref: NULL
        rows: 69115
  Extra: Using index; Using temporary;
Using filesort

***** 2. row *****
      id: 1
  select_type: SIMPLE
        table: nodes
        type: hash_ALL
possible_keys: changeset_id
         key: #hash#changeset_id
      key_len: 8
         ref: osm.changesets.id
        rows: 2781732
  Extra: Using join buffer (flat, BNLH join)
2 rows in set (0.00 sec)

mariadb-10.3 (osm) > SELECT visible, count(*) FROM
changesets JOIN nodes on changesets.id =
nodes.changeset_id GROUP BY visible;
+-----+-----+
| visible | count(*) |
+-----+-----+
|      1 |   2865312 |
+-----+-----+
1 row in set (6.66 sec)
```

Query Optimization with MySQL 8.0 and MariaDB 10.3: The Basics

SUBQUERIES

Access types: unique_subquery/index_subquery

```
mysql-8.0 (osm) > EXPLAIN SELECT *
FROM node_tags WHERE v = 'Big Ben'
and node_id NOT IN (SELECT node_id
FROM nodes WHERE tile < 100000000)\G
***** 1. row *****
      id: 1
  select_type: PRIMARY
      table: node_tags
      type: ref
possible_keys: v_idx
      key: v_idx
    key_len: 767
      ref: const
      rows: 1
  Extra: Using where; Using
index
***** 2. row *****
      id: 2
  select_type: DEPENDENT SUBQUERY
      table: nodes
      type: index_subquery
possible_keys: PRIMARY,nodes_tile_idx
      key: PRIMARY
    key_len: 8
      ref: func
      rows: 1
  Extra: Using where
2 rows in set (0.00 sec)
```

Unique subquery is similar, but using a unique or primary key

Subqueries in MySQL

- MySQL versions traditionally had very bad press regarding subqueries
 - It was common to recommend rewriting them (when possible) into JOINS
- Since MySQL 5.6, its query execution plans have improved significantly

Lazy Materialization of derived tables

- Option available since MySQL 5.6
 - Improves the execution time of EXPLAIN (it no longer needs to execute subqueries)
 - Derived tables can be indexed automatically at execution time to improve its performance

Derived Table Example

```
mysql-5.5.40 (osm) > EXPLAIN SELECT count(*) FROM (SELECT * FROM nodes WHERE VISIBLE = 1) n JOIN changesets ON n.changeset_id = changesets.id;
```

id	select_type	table	type	possible_keys	key	key_len	ref	rows	Extra
1	PRIMARY	<derived2>	ALL	NULL	NULL	NULL	NULL	2865312	
1	PRIMARY	changesets	eq_ref	PRIMARY	PRIMARY	8	n.changeset_id	1	Using index
2	DERIVED	nodes	ALL	NULL	NULL	NULL	NULL	2865521	Using where

3 rows in set (1.42 sec)

```
mysql-8.0 (osm) > EXPLAIN SELECT count(*) FROM (SELECT * FROM nodes WHERE VISIBLE = 1) n JOIN changesets ON n.changeset_id = changesets.id;
```

id	select_type	table	type	possible_keys	key	key_len	ref	rows	Extra
1	PRIMARY	changesets	index	PRIMARY	PRIMARY	8	n.changeset_id	70917	Using index
1	PRIMARY	<derived>	ALL	NULL	NULL	NULL	NULL	40	NULL
2	DERIVED	nodes	ALL	NULL	NULL	NULL	NULL	2853846	Using where

3 rows in set (0.00 sec)

Subquery is not executed

Auto-generated index

A Common 5.5 Performance Problem

```
mysql-5.5.40 (osm) > EXPLAIN SELECT * FROM nodes
                        WHERE nodes.changeset_id IN (
                        SELECT changesets.id
                        FROM changesets
                        JOIN users
                        ON changesets.user_id = users.id and users.display_name = 'Steve');
```

id	select_type	table	type	possible_keys	key	key_len	ref	rows	Extra
1	PRIMARY	nodes	ALL	NULL	NULL	NULL	NULL	2865521	Using where
2	DEPENDENT SUBQUERY	users	const	PRIMARY,users_display_name_idx	users_display_name_idx	767	const	1	Using index
2	DEPENDENT SUBQUERY	changesets	eq_ref		PRIMARY	8	func	1	Using where

3 rows in set (0.00 sec)

```
mysql-5.5.40 (osm) > SELECT ...;
```

node_id	latitude	longitude	changeset_id	visi
99890	515276425	-1497621	552	
109174	515364532	-1457329	1875	
276538	515324296	-2094688	810	1
442987	515449207	-1275650	1941	1
442988	515449741	-1272860	1941	1
498803	515438432	-1269436	2171	1
138212838	513010180	-1699929	7757299	1

7 rows in set (2.60 sec)

This means that the subquery is executed almost 3 million times

Semijoin Optimization

- The only way to execute certain IN subqueries was to execute them with poor strategy
 - This forced rewriting of certain queries into JOINS or scalar subqueries, when possible
- There are now several additional options (many automatic):
 - Convert to a JOIN
 - Materialization (including index creation)
 - FirstMatch
 - LooseScan
 - Duplicate Weedout

The Previous Query is Not a Problem in 5.6+/MariaDB 5.3+

```
mysql-8.0 (osm) > EXPLAIN SELECT * FROM nodes
                    WHERE nodes.changeset_id IN (
                        SELECT changesets.id
                        FROM change sets
                        JOIN users
                        ON changesets.user_id = users.id and users.display_name = 'Steve');
```

id	select_type	table	type	possible_keys	key	key_len	ref	rows	Extra
1	SIMPLE	users	const	PRIMARY,users_display_name_idx	users_display_name_idx	767	const	1	Using index
1	SIMPLE	changesets	ALL	PRIMARY	NULL	NULL	NULL	70917	Using where
1	SIMPLE	nodes	ref	changeset_id	changeset_id	8	osm.changesets.id	21	NULL

3 rows in set (0.00 sec)

```
mysql-8.0 (osm) > SELECT ...;
```

node_id	latitude	longitude	changeset_id	visible	timestamp	tile	version
99890	515276425	-1497621	552	1	2005-10-25 00:35:24	2062268512	1
...							
138212838	513010180	-1699929	7757299	1	2011-04-03 18:14:14	2062220563	6

7 rows in set (0.02 sec)

Executed as
a regular
JOIN

First Match Strategy

```
mysql-8.0 (osm) > EXPLAIN SELECT * FROM changesets WHERE id IN (SELECT changeset_id FROM nodes)\G
***** 1. row *****
      id: 1
  select_type: SIMPLE
        table: changesets
  partitions: NULL
         type: ALL
possible_keys: PRIMARY
          key: NULL
        key_len: NULL
         ref: NULL
         rows: 70917
   filtered: 100.00
      Extra: NULL
***** 2. row *****
      id: 1
  select_type: SIMPLE
        table: nodes
  partitions: NULL
         type: ref
possible_keys: changeset_id
          key: changeset_id
        key_len: 8
         ref: osm.changesets.id
         rows: 33
   filtered: 100.00
      Extra: Using index; FirstMatch(changesets)
2 rows in set, 1 warning (0.00 sec)
```

It is converting the ref into an eq_ref, shot-circuiting the execution

Enabling and disabling materialization, semijoin, etc

```
mysql-8.0 (osm) > SHOW VARIABLES like 'optimizer_switch'\G
```

```
***** 1. row *****
```

```
Variable_name: optimizer_switch
```

```
Value:
```

```
index_merge=on,index_merge_union=on,index_merge_sort_union=on,index_merge_in  
tersection=on,engine_condition_pushdown=on,index_condition_pushdown=on,mrr=o  
n,mrr_cost_based=on,block_nested_loop=on,batched_key_access=off,materializat  
ion=on,semijoin=on,loosescan=on,firstmatch=on,duplicateweedout=on,subquery_m  
aterialization_cost_based=on,use_index_extensions=on,condition_fanout_filter  
=on,derived_merge=on
```

```
1 row in set (0.00 sec)
```

Query Optimization with MySQL 8.0 and MariaDB 10.3: The Basics

FULLTEXT, GEODATA SEARCH AND JSON SUPPORT

Fuzzy Search of “hotel/Grand Hotel/Hotel X/etc”

- “Typical” way to solve this:

```
mysql (osm) >  
SELECT way_id as id, v  
FROM way_tags  
WHERE v like '%otel%';  
2125 rows in set (3.19 sec)
```



Too slow

Let's Add an Index

```
mysql (osm) > ALTER TABLE way_tags ADD INDEX(v);  
Query OK, 0 rows affected (23.96 sec)  
Records: 0 Duplicates: 0 Warnings: 0
```

Still slow, why?

```
mysql (osm) > SELECT ...;  
2125 rows in set (2.75 sec)
```

```
mysql (osm) > EXPLAIN SELECT way_id as id, v FROM way_tags  
WHERE v like '%otel%';
```

id	type	possible_keys	key	key_len	ref	rows	filtered	Extra
1	index	NULL	v	769	NULL	8964000	11.11	Using where; Using index

1 row in set, 1 warning (0.01 sec)

Fulltext Index

```
mysql (osm) > ALTER TABLE way_tags add FULLTEXT index(v);
Query OK, 0 rows affected (53.53 sec)
Records: 0  Duplicates: 0  Warnings: 0
```

```
mysql (osm) > SELECT 'way' as type, way_id as id, v FROM
way_tags WHERE MATCH(v) AGAINST ('+hotel*' IN BOOLEAN MODE);
1851 rows in set (0.04 sec)
```

```
mysql (osm) > EXPLAIN ...;
```

id	type	possible_keys	key	key_len	ref	rows	filtered	Extra
1	fulltext	v	v	0		1	100.00	Using where; Ft_hints: no_ranking

1 row in set, 1 warning (0.00 sec)

Some fields are
useless for FULLTEXT

'Newer' Fulltext Optimizations

```
mysql-5.5 (osm) > EXPLAIN SELECT count(*) FROM way_tags_myisam WHERE MATCH(v) AGAINST('hotel');
```

id	select_type	table	type	possible_keys	key	key_len	ref	rows	Extra
1	SIMPLE	way_tags_myisam	fulltext	v	v	0		1	Using where

1 row in set (0.00 sec)

```
mysql-5.5 (osm) > SHOW STATUS like 'Hand%';
```

Variable_name	Value
Handler_commit	0
Handler_delete	0
...	
Handler_read_first	0
Handler_read_key	0
Handler_read_last	0
Handler_read_next	425
Handler_read_prev	0
Handler_read_rnd	0
Handler_read_rnd_next	0
...	
Handler_update	0
Handler_write	0

16 rows in set (0.00 sec)

Newer Fulltext Optimizations (cont.)

```
mysql-8.0 (osm) > EXPLAIN SELECT count(*) FROM way_tags WHERE MATCH(v) AGAINST('hotel');
+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+
| id | select_type | table | partitions | type | possible_keys | key | key_len | ref | rows | filtered | Extra |
+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+
| 1 | SIMPLE      | NULL | NULL        | NULL | NULL          | NULL | NULL    | NULL | NULL | NULL     | Select tables optimized away |
+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+
1 row in set, 1 warning (0.00 sec)
```

```
mysql-8.0 (osm) > SHOW STATUS like 'Hand%';
+-----+-----+
| Variable_name | Value |
+-----+-----+
| Handler_commit | 1     |
| Handler_delete | 0     |
...
| Handler_read_first | 0     |
| Handler_read_key   | 0     |
| Handler_read_last  | 0     |
| Handler_read_next  | 0     |
| Handler_read_prev  | 0     |
| Handler_read_rnd   | 0     |
| Handler_read_rnd_next | 0     |
...
| Handler_update     | 0     |
| Handler_write      | 0     |
+-----+-----+
18 rows in set (0.00 sec)
```

It's counting directly from the FULLTEXT index

Open Issues and Limitations

- No postfix support (wildcards)
- Simple Ranking (and different from MyISAM)
- No stemming support
- Some multi-language limitations

More on FULLTEXT InnoDB support:

<http://www.drdobbs.com/database/full-text-search-with-innodb/231902587>

Alternatives

- Apache Lucene
 - Solr
 - Elasticsearch
- Sphinx
 - SphinxSE

Find the Closest Starbucks (Spatial Search)



```
mysql (osm) > SET @lat:=50.11745057; SET @lon:= 8.62699463;
mysql (osm) > SELECT n.node_id,
    n.longitude/100000000 as longitude,
    n.latitude/100000000 as latitude,
    sqrt(pow((latitude/100000000 - @lat) * 111231.29, 2) +
        pow((longitude/100000000 - @lon) * 71520.91, 2))
    as `distance in metres`
FROM nodes n
JOIN node_tags n_t1
  ON n.node_id = n_t1.node_id
JOIN node_tags n_t2
  ON n.node_id = n_t2.node_id
WHERE
    n_t1.k = 'amenity' and
    n_t1.v = 'cafe' and
    n_t2.k = 'name' and
    n_t2.v like 'Starbucks%'
ORDER BY `distance in metres` ASC
LIMIT 1;

+-----+-----+-----+-----+
| node_id | longitude | latitude | distance in metres |
+-----+-----+-----+-----+
| 3136353405 | 8.6523 | 50.1093 | 2023.6465372574328 |
+-----+-----+-----+-----+
1 row in set, 36 warnings (0.81 sec)
```


This Query is Slow

```
mysql [osm]> EXPLAIN SELECT ... \G
```

```
***** 1. row *****
```

```

      Id: 1
  select_type: SIMPLE
        table: n_t1
         type: ALL
possible_keys: PRIMARY
          key: NULL
      key_len: NULL
         ref: NULL
        rows: 2210950
      Extra: Using where;
Using temporary; Using filesort

```

```
***** 2. row *****
```

```

      id: 1
  select_type: SIMPLE
        table: n_t2
         type: ref
possible_keys: PRIMARY
          key: PRIMARY
      key_len: 8
         ref: osm.n_t1.node_id
        rows: 1
      Extra: Using where

```

```
***** 3. row *****
```

```

      id: 1
  select_type: SIMPLE
        table: n
         type: ref
possible_keys: PRIMARY
          key: PRIMARY
      key_len: 8
         ref: osm.n_t1.node_id
        rows: 1
      Extra:
3 rows in set (0.00 sec)

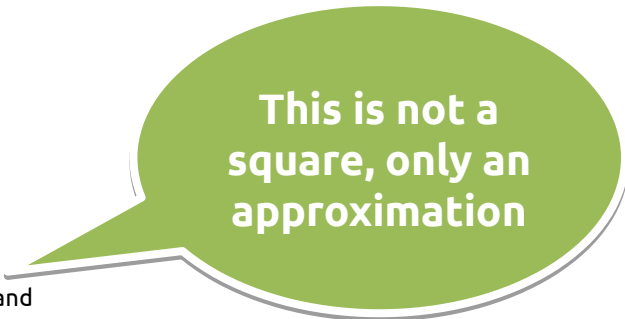
```

Can We Optimize it?

- We could add a bounding box:

```
mysql (osm) > ALTER TABLE nodes ADD INDEX (latitude, longitude);

mysql (osm) > SELECT n.node_id,
    n.longitude/100000000 as longitude,
    n.latitude/100000000 as latitude,
    sqrt(pow((latitude/100000000 - @lat) * 111231.29, 2) +
        pow((longitude/100000000 - @lon) * 71520.91, 2))
        as `distance in metres`
FROM nodes n
JOIN node_tags n_t1
    ON n.node_id = n_t1.node_id
JOIN node_tags n_t2
    ON n.node_id = n_t2.node_id
WHERE
    n_t1.k = 'amenity' and
    n_t1.v = 'cafe' and
    n_t2.k = 'name' and
    n_t2.v like 'Starbucks%' and
    n.latitude BETWEEN ((@lat - 1000/111231.29) * 100000000)
        AND ((@lat + 1000/111231.29) * 100000000) and
    n.longitude BETWEEN ((@lon - 1000/71520.91) * 100000000)
        AND ((@lon + 1000/71520.91) * 100000000)
ORDER BY `distance in metres` ASC
LIMIT 1;
```



This is not a square, only an approximation

We Create an Index... and Force It

mysql-8.0 (osm) > EXPLAIN SELECT ...;

id	select_type	table	partitions	type	pos_	key	key_	ref	rows	filtered	Extra
1	SIMPLE	n_t1	NULL	ALL	PRIMAR				2210950	0.00	Using where;
1	SIMPLE	n	NULL	ref	PRIMARY,	latitude	PRIMARY	8	osm.n_t1.node_id	1	Using temporary;
1	SIMPLE	n_t2	NULL	ref	PRIMARY	PRIMARY	8	osm.n_t1.node_id	1	1.41	Using where

3 rows in set, 1 warning (0.00 sec)

mysql-8.0 (osm) > EXPLAIN SELECT ... FROM nodes n FORCE INDEX(latitude) ...;

id	select_type	table	partitions	type	possible	key	key_	ref	rows	filtered	Extra
1	SIMPLE	n	NULL	range	latitude	latitude	8	NULL	949370	11.11	Using where;
1	SIMPLE	n_t1	NULL	ref	PRIMARY	PRIMARY	8	osm.n.n	1	1.41	Using index;
1	SIMPLE	n_t2	NULL	ref	PRIMARY	PRIMARY	8		1	1.41	Using index;

3 rows in set, 1 warning (0.00 sec)

mysql-8.0 (osm) > SELECT ...;
0 rows in set (0.26 sec)

MySQL ignores the newly created index, why?

Performance improvement is not great

Still many rows are examined

Most of the gain comes from the covering index, not the filtering

Multiple Range Scans Cannot Be Optimized with BTREE Indexes

- We need quadrees or R-TREE Indexes for indexing in multiple dimensions
 - The latter are implemented in MySQL with the name “SPATIAL indexes”, as they only apply to GIS types
- Spatial indexing is available on MySQL since 5.7.5

Creating a Spatial Index

```
mysql (osm) > ALTER TABLE nodes
                ADD COLUMN coord GEOMETRY NOT NULL SRID 4326;
Query OK, 0 rows affected (0.01 sec)
Records: 0 Duplicates: 0 Warnings: 0
```

Instant, yay!
(or create it
VIRTUAL)

```
mysql (osm) > UPDATE nodes
                SET coord = point(longitude/100000000,
                                latitude/100000000);
Query OK, 20216490 rows affected (7 min 10.28 sec)
Rows matched: 20216490 Changed: 20216490 Warnings: 0
```

```
mysql (osm) > ALTER TABLE nodes add SPATIAL index(coord);
Query OK, 0 rows affected (1 hour 7 min 13.22 sec)
Records: 0 Duplicates: 0 Warnings: 0
```

This is new since
5.7

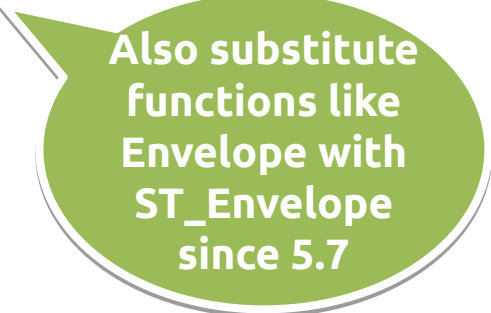
New Query

```
mysql> SET @area := ST_Envelope(linestring(POINT(@lon - 500/71520.91, @lat - 500/111231.29), POINT(@lon + 500/71520.91, @lat + 500/111231.29)));

mysql> SELECT n.node_id,
              x(n.coord) as longitude,
              y(n.coord) as latitude,
              st_distance(POINT(@lon, @lat), coord) as distance
FROM nodes n
JOIN node_tags n_t1
  ON n.node_id = n_t1.node_id
JOIN node_tags n_t2
  ON n.node_id = n_t2.node_id
WHERE
  n_t1.k = 'amenity' and
  n_t1.v = 'cafe' and
  n_t2.k = 'name' and
  n_t2.v like 'Starbucks%' and
  st_within(coord, @area)
ORDER BY st_distance(POINT(@lon, @lat), coord) ASC
LIMIT 1;
```



We can use any shape we want thanks to 5.6 improvements



Also substitute functions like Envelope with ST_Envelope since 5.7

Better Performance

```
mysql (osm) > SELECT ...;
```

```
+-----+-----+-----+-----+
| node_id   | longitude  | latitude  | distance   |
+-----+-----+-----+-----+
| XXXXXXXXXX | XXXXXXXXXX | XXXXXXXXXX | 0.0014631428672541478 |
+-----+-----+-----+-----+
```

Empty set (0.11 sec)

This field used to be almost useless (wait for it)

```
mysql (osm) > EXPLAIN SELECT ...;
```

id	select _type	table	parti tions	type	possible _keys	key	key _len	ref	rows	filtered	Extra
1	SIMPLE	n	NULL	range	PRIMARY ,coord	coord	34	NULL	2	100.00	Using where; Using filesort
1	SIMPLE	n_t1	NULL	ref	PRIMARY	PRIMARY	8	osm.n. node_id	3	1.41	Using where
1	SIMPLE	n_t2	NULL	ref	PRIMARY	PRIMARY	8	osm.n. node_id	3	1.41	Using where

```
3 rows in set (0.00 sec)
```

Better Filtering

mysql-8.0 (osm) > SHOW STATUS LIKE 'Hand%';

Not using the index: Using the BTREE index: Using the SPATIAL index:

Variable_name	Value
Handler_commit	1
Handler_delete	0
Handler_discover	0
Handler_external_lock	6
Handler_mrr_init	0
Handler_prepare	0
Handler_read_first	1
Handler_read_key	1914
Handler_read_last	0
Handler_read_next	1954
Handler_read_prev	0
Handler_read_rnd	1
Handler_read_rnd_next	833426
Handler_rollback	0
Handler_savepoint	0
Handler_savepoint_rollback	0
Handler_update	0
Handler_write	1

18 rows in set (0.00 sec)

Variable_name	Value
Handler_commit	1
Handler_delete	0
Handler_discover	0
Handler_external_lock	6
Handler_mrr_init	0
Handler_prepare	0
Handler_read_first	0
Handler_read_key	274
Handler_read_last	0
Handler_read_next	246540
Handler_read_prev	0
Handler_read_rnd	0
Handler_read_rnd_next	0
Handler_rollback	0
Handler_savepoint	0
Handler_savepoint_rollback	0
Handler_update	0
Handler_write	0

18 rows in set (0.00 sec)

Variable_name	Value
Handler_commit	1
Handler_delete	0
Handler_discover	0
Handler_external_lock	6
Handler_mrr_init	0
Handler_prepare	0
Handler_read_first	0
Handler_read_key	522
Handler_read_last	0
Handler_read_next	5254
Handler_read_prev	0
Handler_read_rnd	259
Handler_read_rnd_next	0
Handler_rollback	0
Handler_savepoint	0
Handler_savepoint_rollback	0
Handler_update	0
Handler_write	0

18 rows in set (0.00 sec)

Geohash Functions

```
mysql (osm) > SELECT ST_GeoHash(@lon, @lat, 10);
```

```
+-----+  
| ST_GeoHash(@lon, @lat, 10) |  
+-----+  
| u0yjh79gr9                |  
+-----+  
1 row in set (0.00 sec)
```

- Useful to index coordinates with a BTREE
 - It could be specially useful combined with indexed STORED columns (emulating quadtrees)

More on Geohashing: <http://mysqlserverteam.com/geohash-functions/>

GeoJSON Functions

```
mysql (osm) > SELECT nm.v, ST_AsGeoJson(n.coord)
                FROM node_tags n_t
                JOIN nodes n USING (node_id, version)
                JOIN node_tags nm USING (node_id, version)
                WHERE n_t.k='tourism' AND
                       n_t.v='attraction' AND
                       nm.k='name';
```

+-----+-----+	
v	ST_AsGeoJson(n.coord)
+-----+-----+	
Schießbuckel	{"type": "Point", "coordinates": [8.4938556, 49.6351754]}
AKW Informationszentrum	{"type": "Point", "coordinates": [8.4175585, 49.7060227]}
Burg Gleiberg	{"type": "Point", "coordinates": [8.6344314, 50.6150329]}
Brüderkirche	{"type": "Point", "coordinates": [9.5042975, 51.3149351]}
Römer	{"type": "Point", "coordinates": [8.6816587, 50.1104684]}

GeoJSON Functions (cont.)

```
$ mysql osm -B -e "SET NAMES utf8; SET SESSION group_concat_max_len = 10000; SELECT
CONCAT('{\"type\": \"FeatureCollection\", \"features\": [ ',
GROUP_CONCAT(CONCAT('{\"type\": \"Feature\", \"geometry\": ', ST_AsGeoJson(n.coord),
', \"properties\": {\"name\": \"', nm.v, '\"}}')), ' ]}') FROM node_tags n_t JOIN nodes n USING
(node_id, version) JOIN node_tags nm USING (node_id, version) WHERE n_t.k='tourism' and
n_t.v='attraction' AND nm.k='name'"
```

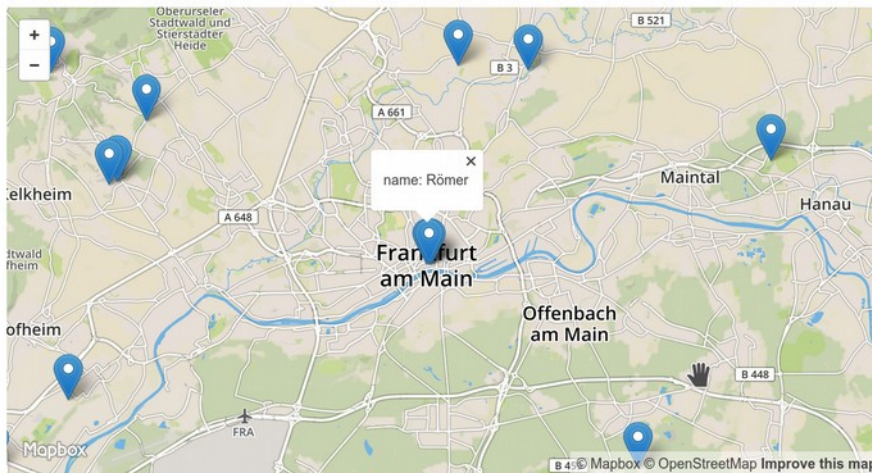
<http://geojsonlint.com/>

```
{
  "geometry": {
    "type": "Point",
    "coordinates": [
      9.1487472, 50.1716933
    ]
  },
  "properties": {
    "name": "Geografischer Mittelpunkt der EU von Jan. 2007 bis Juni 2013"
  },
  "type": "Feature",
  "geometry": {
    "type": "Point",
    "coordinates": [
      9.1785149, 50.1382509
    ]
  },
  "properties": {
    "name": "Mariengrotte"
  },
  "type": "Feature",
  "geometry": {
    "type": "Point",
    "coordinates": [
      8.2787808, 49.9999018
    ]
  },
  "properties": {
    "name": "Fischtor"
  },
  "type": "Feature",
  "geometry": {
    "type": "Point",
    "coordinates": [
      8.2828665, 49.9962757
    ]
  },
  "properties": {
    "name": "Dagobertor"
  },
  "type": "Feature",
  "geometry": {
    "type": "Point",
    "coordinates": [
      9.797159, 51.3216731
    ]
  },
  "properties": {
    "name": "Verlorener Bach"
  }
}]
```

Clear Current Features

Test GeoJSON

Clear



Older Issues

- Before MySQL 8.0, SRID could be set and retrieved, but all operations were done in squared euclidean coordinates:

```
mysql-5.7.5 (osm) > SET @p1 := GeomFromText('POINT(8 50)', 4326);
Query OK, 0 rows affected (0.00 sec)
```

```
mysql-5.7.5 (osm) > SET @p2 := GeomFromText('POINT(7 50)', 4326);
Query OK, 0 rows affected (0.00 sec)
```

```
mysql-5.7.5 (osm) > SET @p3 := GeomFromText('POINT(8 51)', 4326);
Query OK, 0 rows affected (0.00 sec)
```

```
mysql-5.7.5 (osm) > SELECT srid(@p1);
+-----+
| srid(@p1) |
+-----+
|      4326 |
+-----+
1 row in set (0.00 sec)
```

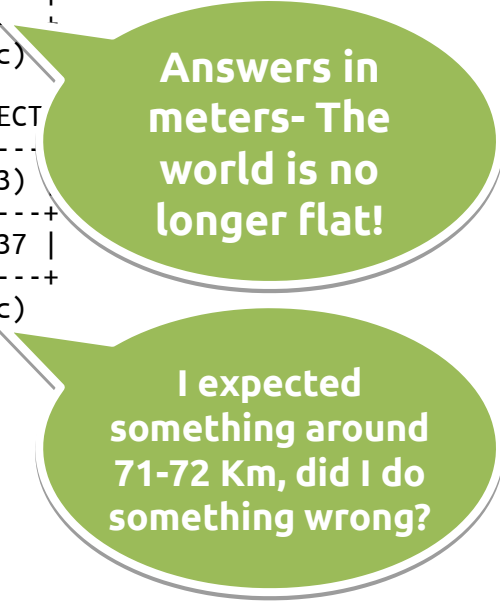
```
mysql-5.7.5 (osm) > SELECT st_distance(@p1, @p2);
+-----+
| st_distance(@p1, @p2) |
+-----+
|                      1 |
+-----+
1 row in set (0.00 sec)
```

```
mysql-5.7.5 (osm) > SELECT st_distance(@p1, @p3);
+-----+
| st_distance(@p1, @p3) |
+-----+
|                      1 |
+-----+
1 row in set (0.00 sec)
```

New 8.0 SRID Support !

```
mysql-8.0 (osm) > SET @p1 := ST_GeomFromText('POINT(8 50)',  
4326);  
Query OK, 0 rows affected (0.00 sec)  
  
mysql-8.0 (osm) > SET @p2 := ST_GeomFromText('POINT(7 50)',  
4326);  
Query OK, 0 rows affected (0.00 sec)  
  
mysql-8.0 (osm) > SET @p3 := ST_GeomFromText('POINT(8 51)',  
4326);  
Query OK, 0 rows affected (0.00 sec)  
  
mysql-8.0 (osm) > SELECT ST_srid(@p1);  
+-----+  
| ST_srid(@p1) |  
+-----+  
|          4326 |  
+-----+  
1 row in set (0.00 sec)
```

```
mysql-8.0 (osm) > SELECT st_distance(@p1, @p2);  
+-----+  
| st_distance(@p1, @p2) |  
+-----+  
| 110243.2689139409 |  
+-----+  
1 row in set (0.00 sec)  
  
mysql-8.0 (osm) > SELECT  
+-----+  
| st_distance(@p1, @p3) |  
+-----+  
| 110592.12831394037 |  
+-----+  
1 row in set (0.01 sec)
```



JSON Native Data Type

- Since 5.7.8, MySQL allows columns defined with the JSON data type:

```
mysql-8.0 > CREATE TABLE json_test(id int PRIMARY KEY  
content JSON);
```

```
Query OK, 0 rows affected (0.03 sec)
```

```
mysql-8.0 > INSERT INTO json_test (content) VALUES ('{"type":  
"correct_json"}');
```

```
Query OK, 1 row affected (0.00 sec)
```

```
mysql-8.0 > INSERT INTO json_test (content) VALUES ('{"type":  
"incorrect_json"}');
```

```
ERROR 3140 (22032): Invalid JSON text: "Missing a closing quotation  
mark in string." at position 24 in value (or column) '{"type":  
"incorrect_json"}'.
```

MariaDB
supports it
only as a
BLOB/LONGT
EXT alias

They get
validated on
insert

JSON functions

- MySQL includes almost all functions to manipulate JSON that you may think of:
 - Validation test: `JSON_TYPE`
 - Object creation: `JSON_ARRAY`, `JSON_MERGE`, ...
 - Searching: `JSON_EXTRACT`
 - Modifying: `JSON_SET`, `JSON_INSERT`, ...

Indexing JSON

- JSON Columns cannot be indexed:

```
mysql [localhost] {msandbox} (test) > ALTER TABLE  
json_test ADD INDEX(content);  
ERROR 3152 (42000): JSON column 'content' cannot be  
used in key specification.
```
- However, they can be compared with regular fields and use indexes thanks to virtual columns

Query Optimization with MySQL 8.0 and MariaDB 10.3: The Basics

CONCLUSIONS

8.0/10.3 Recently released

- They are mature enough to be used in production
- Latest release doesn't have as many new features related to query optimization
 - But they have a lot of other features and quality of life improvements
 - They also fix long outstanding query performance bugs
 - MariaDB also catches up with MySQL in some areas

MySQL 8.0 New Features

- MySQL 5.6 seemed Percona Server-inspired
- MySQL 5.7 seemed MariaDB-inspired
- MySQL 8.0 seems Galera and Oracle DB-inspired
 - Competition is [always good for consumer](#)



Modern SQL is now part of MySQL!

- Window functions, recursive CTEs:

```
WITH RECURSIVE cte (cl_from, cl_type) AS
(
  SELECT cl_from, cl_type FROM categorylinks WHERE cl_to = 'Database_management_systems'
  -- starting category
  UNION
  SELECT categorylinks.cl_from, categorylinks.cl_type FROM cte JOIN page ON cl_from =
  page_id JOIN categorylinks ON page_title = cl_to WHERE cte.cl_type = 'subcat' --
  subcat addition on each iteration
)
SELECT page_title FROM cte JOIN page ON cl_from = page_id WHERE page_namespace = 0
ORDER BY page_title; -- printing only articles in the end, ordered by title
```



Scan Wikipedia
categories,
page links
recursively!

Many Optimizer Advantages Have to Be Manually Enabled

- Modifying on a per-query basis:

```
SET optimizer_switch='batched_key_access=on';  
SET join_cache_level=8; # for MariaDB
```

in order to take advantage of them, making some of features useless unless you are fine-tuning

- I expect that to change in the future

I Herby Declare MyISAM as Dead

- All major MyISAM-only features are now on MySQL 5.7/8.0
 - FULLTEXT
 - GIS
 - Transportable tables
- Previous blockers such as grant tables (new InnoDB data dictionary) and temporary tables are now in InnoDB format

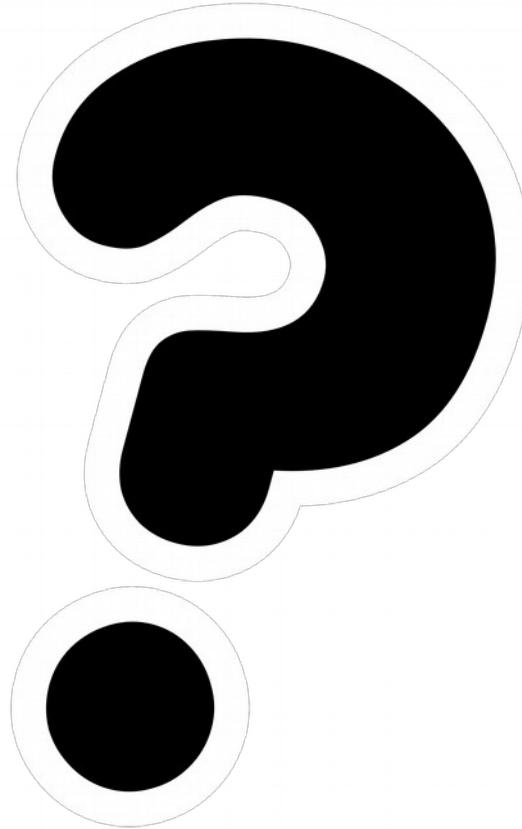


Benchmarks

- Do not trust first party benchmarks
 - In fact, do not trust 3rd party benchmarks either
- Only care about the performance of your application running on your hardware



Q&A



Not to Miss

- Operations track:
[TLS for MySQL at Large Scale:](#)
How we do relational data on-the-wire encryption at the Wikimedia Foundation
- Do you want to do query optimization for a website with 20 Billion views per month?
<https://wikimediafoundation.org/about/jobs/>



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