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A few proofs regarding algorithms

Page 2: Bubble Sort (INCOMPLETE)

 $\overrightarrow{\mathrm{Page}}\ 3$: Knuth Volume 1. Chapter 1. Problem 2 (Euclidean Algorithm)

A proof on bubble sort

A sorted set is defined as follows (Where S_0 is the initial element):

$$S_k \leq S_{k+1} \vee$$

$$S_k = |S| - 1$$

$$\texttt{procedure bubbleSort}(\texttt{S}: \texttt{Ordered set})$$

$$\texttt{swapped} \leftarrow true$$

$$\texttt{while swapped}$$

$$\texttt{swapped} \leftarrow false$$

$$\texttt{for n} \rightarrow |\texttt{S}| - 2$$

$$\texttt{if } S_n > S_{n+1}$$

$$\texttt{Swap}(S_n, S_{n+1})$$

$$\texttt{swapped} \leftarrow true$$

A proof to solve a problem from Knuth regarding Euclidean Algorithm

2. [15] Prove that m is always greater than n at the beginning of step $\mathbf{E1}$, after it has already occurred once. Where $\mathbf{E1}$ is part of the algorithm, \mathbf{E} :

Algorithm E:

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E1: Divide m by n and let r be the remainder.
E2: If r = 0: return n
E3: m \leftarrow n, n \leftarrow r \text{ goto } \mathbf{E1}
Proof:
Prove that m_k > n_k where k > 0
Let:
        m_k = n_{k-1}
        n_k = m_{k-1} \bmod n_{k-1}
        n_k \neq 0
        k = 0 at the first time at E1
        k > 0
m_k > n_k \to
n_{k-1} > m_{k-1} \mod n_{k-1}
Case I m_{k-1} = n_{k-1}:
        Impossible:
        m_{k-1} = n_{k-1} \to m_{k-1} \bmod n_{k-1} = 0
        n_k = m_{k-1} \mod n_{k-1} = 0 \rightarrow
                                                                           Contradiction
        n_k = 0 \land n_k \neq 0
Case II m_{k-1} > n_{k-1}:
        1. m_{k-1} > n_{k-1} \to
        2. n_{k-1} > m_{k-1} \mod n_{k-1}
        3. m_k > n_k
Case III m_{k-1} < n_{k-1}:
         1. m_{k-1} < n_{k-1} \to
        2. m_{k-1} \mod n_{k-1} = m_{k-1} \to
                                                                      Definition of case
        3. m_k = n_{k-1} > m_{k-1} \wedge
        4. n_k = m_{k-1}
        5. m_k > n_k
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