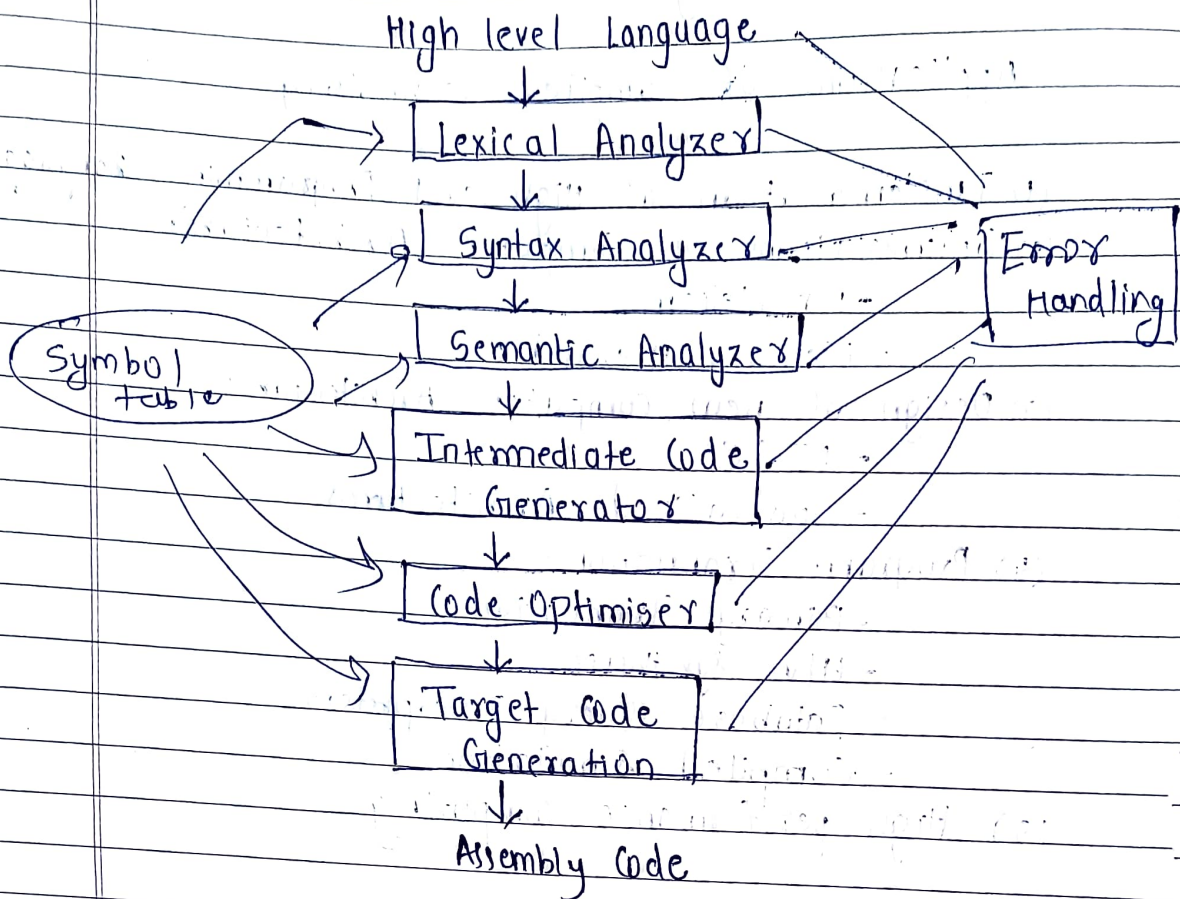


# Compiler Design

## Applications of Compiler Technology

- (1) Implementation of High-level Programming Languages
- (2) Optimizations of Computer Architectures
  - parallelism
  - Memory Hierarchies
- (3) Design of new Computer Architectures
  - RISC
  - Specialized Architectures
- (4) Program Translations
  - Binary Translation
  - HW synthesis
  - Database Query Interpreters
  - Compiled Simulation
- (5) High performance computing



There are two phases of Compilers

- (1) Analysis phase
- (2) Synthesis phase

(1) Analysis phase :-

The Analysis phase creates an intermediate representation from the given source code.

(2) Synthesis phase :-

The Synthesis phase creates an equivalent target program from the intermediate representation.

Symbol table : It is a data structure being used and maintained by the compiler.

## (1) Lexical Analysis :

- Also known as Scanning
- This phase reads the source code and break it into stream of tokens.
- Tokens are basic units of the programming language.

e.g

```
int main()
{
}
```

## (2) Syntax Analysis :

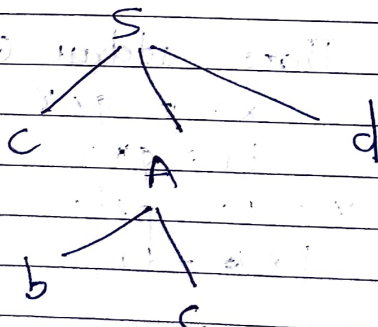
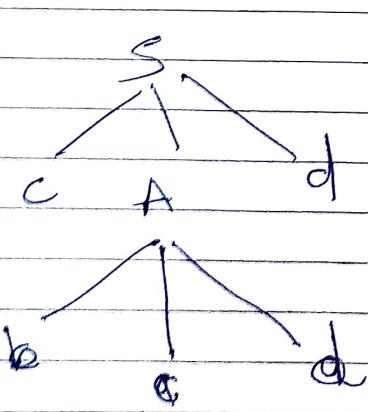
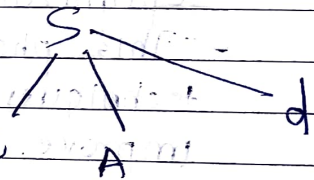
- The second phase of a compiler is syntax analysis.
- This phase takes the stream of tokens generated by the lexical analysis phase and checks whether the grammar correct or not.

e.g.

$S \rightarrow cAd$

$A \rightarrow bc \mid a$

IP  $\rightarrow cad$





### 3) Semantic Analysis

- The third phase of a compiler is semantic analysis.
- This phase checks whether the code is semantically correct.

e.g.

int - x

int - ✓

### 4) Intermediate Code Generation:

- The fourth phase of a compiler is intermediate code generation.
- This phase generates an intermediate representation of the src code that can be easily translated into machine code.

### (5) Code Optimization

- The fifth phase of a compiler is code optimization.
- This phase applies various optimization techniques to the intermediate code to improve the performance of the code.

e.g. Three address code

$x = p + q * r$

$t1 = q * r$

$x = t2 = p + t1$

$x = t2$

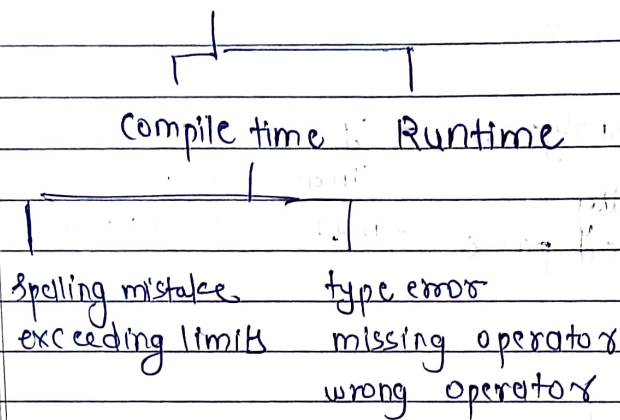
## 6) Code Generation :-

- The final phase of a compiler is code generation.
- This phase takes the optimized intermediate code and generates the actual 'machine' code that can be executed by the target hardware.

## Symbol Table

name	type	size	usage

## Error Handler



## LA / DFA / NFA / conversion

float x, y, z

x = y + z \* 60



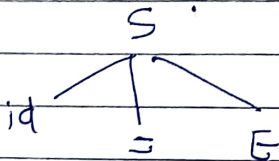
Lexical Analyzer

$\langle id, 1 \rangle = \langle id, 2 \rangle + \langle id, 3 \rangle * 60$



Syntax Analyzer

parse Tree:



Semantic Analyzer

= Type checking

- Undeclared variable
- multiple declaration

symantically  
verified parse  
tree



ICG

$t_1 = z * 60.0$

$t_2 = y + t_1$

$x = t_2$



Code Optimization

$t_1 = z * 60.0$

$x = y + t_1$

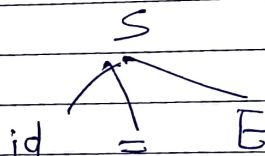
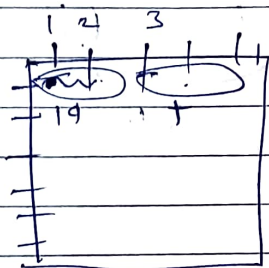
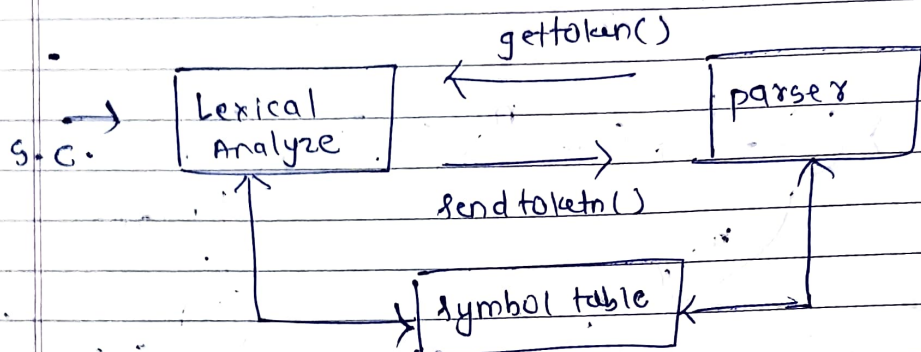


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## Lexical Analyzer

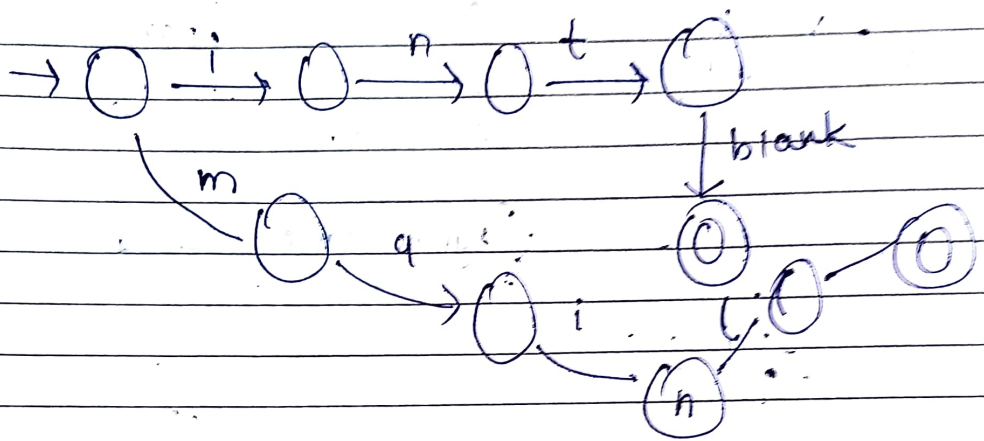
- Lexical analyzer divides the given program into meaningful words which is known as tokens
- Tokens are normally identifiers, keywords, operators, constant and special symbol



- LA helps in giving error messages providing row no. & column no
- LA eliminates comment lines from the given program
- LA eliminates white space character (blank, tab)

- LA uses DFA to do tokenization
- While doing tokenization LA always gives importance to longest matching

```
int main() - 4
{ - 5
    - 6
```



```
main ) ( => 3
{
```

```
x = a + b * c;
int x, a, b, c;
y = x + a;
```

```
{
```

(28)

Syntax error / Semantic error