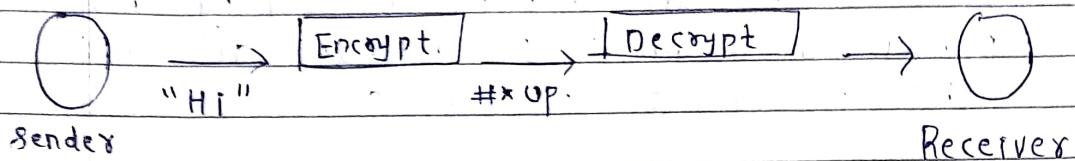


# Cryptography

## Cryptography



Cryptography : The act or art of writing in secret characters .

Cryptoanalysis : The analysis and de-ciphering of secret writings .

Cryptology : Scientific study of Cryptography and Cryptoanalysis .

- Cryptography :
- It is a technique of securing information and communications through use of codes so that only those person for whom the information is intended can understand it and process it .
- crypt = hidden
- graphy = writing

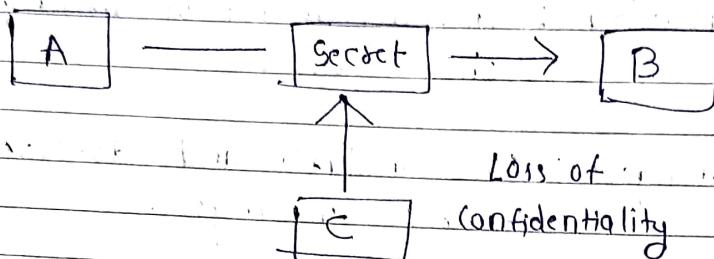
## Applications

- Secure communication
- File and database security
- Electronic funds transfer
- Digital cash
- Electronic mail
- Electronic voting
- Secure web browsing
- Digital currencies
- Secure protocols

## Security Services / Principles of Security / Features

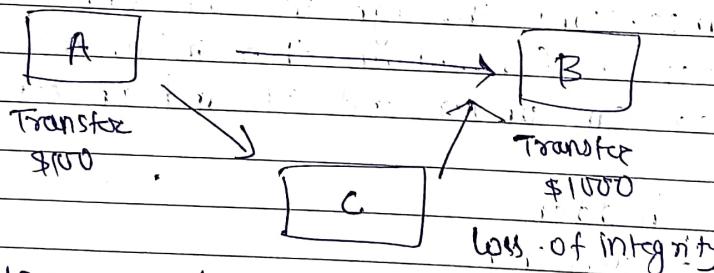
### (1) Confidentiality :-

- Information can only be accessed by the person for whom it is intended and no other person except him can access it.
- Only intended receiver understands the message.



### (2) Integrity :-

Info Ensure that their communication has not been altered, either maliciously or by accident during transmission.

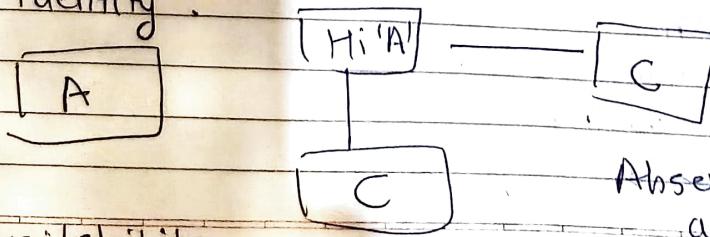


### (3) Non-repudiation :-

Sender should not be able to falsely deny that a message was sent

### (4) Authentication :-

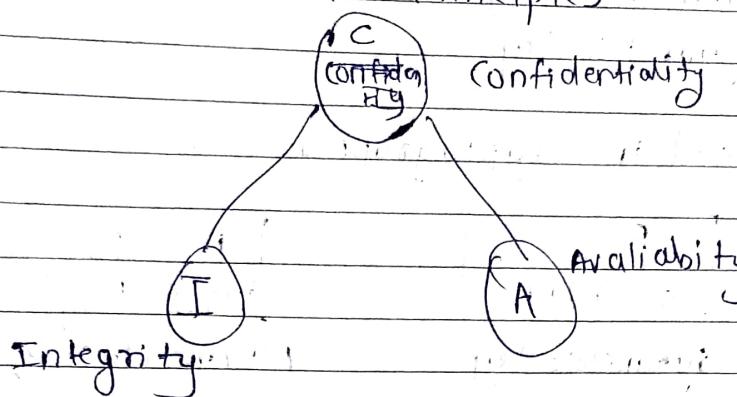
Sender & receiver need to confirm each others identity.



### (5) Availability

# The CIA triad in cryptography

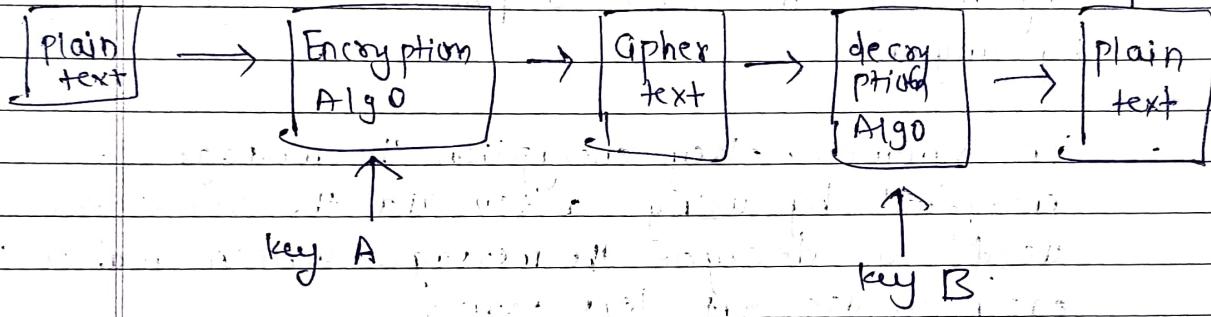
Three fundamental Principles



Cryptography components

Sender (Alice)

Receiver (Bob)



- Cipher is the method for encrypting messages .

The key which is an input to the Algo is Secret .

- key is a string of numbers or characters

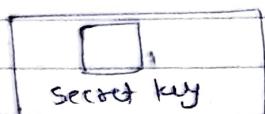
- If same key is used for encryption & decryption the algorithm is called symmetric

- If different keys are used for encryption & decryption the algorithm is called Asymmetric .

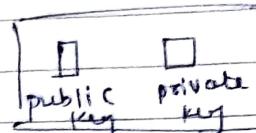
## Categories of Cryptography

- (1) Symmetric-key
- (2) Asymmetric key

key used in cryptography



Symmetric-key



Asymmetric-key

### (1) Symmetric key

- Encryption is a process to change the form of any message, in order to protect it from reading by anyone.
- In symmetric-key encryption the message is encrypted by using a key and the same key is used to decrypt the message which makes it easy to use but less secure.
- It also requires a safe method to transfer the key from one party to another.

### (2) Asymmetric key

Asymmetric key Encryption is based on public and private key encryption techniques.

It uses two different key to encrypt and decrypt the message.

It is more secure than the symmetric key encryption technique but is much slower.

## Symmetric-key Cryptography

- Traditional ciphers
- Simple Modern ciphers
- Modern Round Ciphers
- Mode of Operation.

Traditional  
ciphers

Substitution  
ciphers

Transposition  
ciphers

Mono  
alphabetic

poly-alphabetic

- A substitution cipher replaces one symbol with another.

### Caesar Ciphers

- Letters are replaced by other letters or symbols.

- Caesar cipher is a method in which each letter in the alphabet is rotated by three letters as shown.

Algo

for each plaintext letter ' $p$ ', substitute the cipher text letter ' $c$ '

$$c = E(p, k) \text{ mode } 26 = (p + k) \text{ mod } 26$$

$$p = D(c, k) \text{ mode } 26 = (c - k) \text{ mod } 26$$

Caesar Cipher - Example

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18
A	B	C	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s
+ u	v	w	x	y	z													

K	U		N		A		L
N	X		q		D		O

$$c = (p+k) \bmod 26$$

plain key

$$\begin{aligned} c &= (10+3) \bmod 26 \\ &= 13 \bmod 26 \end{aligned}$$

$$\begin{aligned} c &= 20+3 \% 26 \\ &= 23 \% 26 \\ &= 23 \end{aligned}$$

$$\begin{aligned} c &= (13+3) \bmod 26 \\ &= 16 \bmod 26 \\ &= 16 \end{aligned}$$

$$\begin{aligned} c &= (0+3) \% 26 \\ &= 3 \end{aligned}$$

$$\begin{aligned} 11+3 \bmod 26 \\ = 14 \% 26 \\ = 14 \end{aligned}$$

plain text = kuned

cipher text = NXqDO

Resultant output will be same as cipher text

2. (a) a + b = b + a

$$25 \bmod (3+7) = 25 \bmod 10$$

$$25 \bmod (3+3) = 25 \bmod 6$$

## Monoalphabetic cipher

- The "cipher" line can be any permutation of the 26 alphabetic characters.
- There are  $26!$  pairing of letters
- Statistical Analysis would make it feasible to crack the key.

## Polyalphabetic cipher (Vigenère cipher)

- Improve on the simple monoalphabetic technique
- 26 Caesar ciphers with shifts 0 through 25

Encryption process

$$C_i = (P_i + K_i \bmod m) \bmod 26$$

Decryption

$$P_i = (C_i - K_i \bmod m) \bmod 26$$

- Use table & key word to encipher a message
    - repeat keyword over text e.g. FACE  
F A C E F A C E F A C E F
- MY CAT HAS FLEAS

Using a key to shift alphabet

- Obtain a key for this algo; and then shift the alphabets;

## (e) Transposition Cipher

## Columnar Transposition

- This involves rearrangement of characters on the plain text into columns

pluviometer

## Cipher text

## Ciphers :- Shannon's characteristics of "Good" ciphers

- (1) The amount of Secrecy needed should determine the amount of labor appropriate for the encryption & decryption
- (2) The set of keys and enciphering algorithm should be free from complexity.
- (3) The implementation of the process should be as simple as possible
- (4) Errors in Ciphering should not propagate and cause corruption of further information in the message.
- (5) The size of the enciphered text should be no larger than the text of the original message.

## Encryption Systems :- properties of Trustworthy Systems

- It is based on sound Maths.
- It has been analyzed by competent experts & found to be sound.
- It has stood the "test of time"

## Cryptanalysis

### Techniques

Cryptanalysis is the process of breaking an encryption code.

Several techniques can be used to reduce the algorithm :-

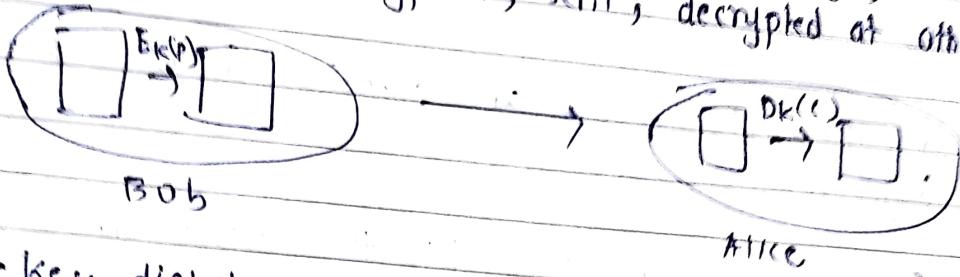
- (1) Attempt to recognize patterns in encrypted messages, to be able to break subsequent ones by applying a straightforward decryption algorithm.
- (2) Attempt to reduce the key, in order to break subsequent messages easily.
- (3) Attempt to find weaknesses in the implementation or env. of use of encryption.
- (4) Attempt to find general weaknesses in an encryption algorithm, without necessarily having intercepted any messages.

# Basic Terminology

- (1) Plain text - The original message
- (2) Ciphertext - the coded message
- (3) Cipher - algo. for transforming plaintext to ciphertext
- (4) key : info used in cipher known only to sender/receiver
- (5) encipher (encrypt)  
converting plaintext to ciphertext
- (6) decipher (decrypt)  
converting ciphertext to plaintext
- (7) Cryptography  
study of encryption principles / methods
- (8) Cryptanalysis :  
study of principles / methods of deciphering ciphertext without knowing key
- (9) Cryptography and its relationship with Cryptanalysis

# Symmetric Cryptography

- Both parties must agree on a secret key,  $K$
- message is encrypted, sent, decrypted at other side



- key distribution must be secret
  - otherwise messages can be decrypted
  - users can be impersonated

## key distribution

Sending this key directly makes it easy for hackers to capture it in transit

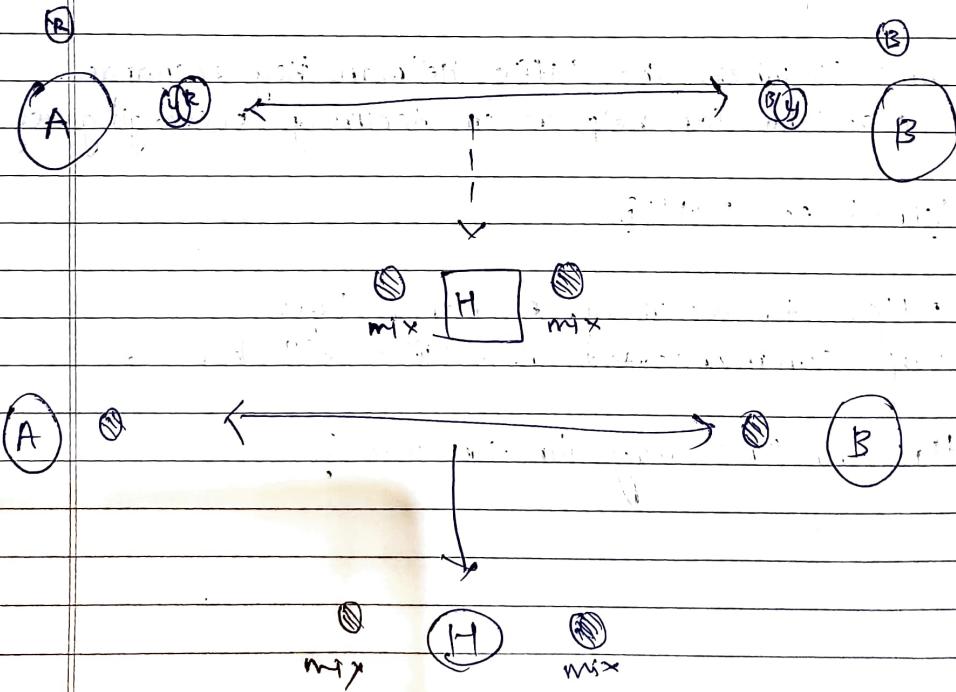
⇒ This is where the Diffie Hellman key exchange algorithm can help us in securely exchanging keys,

## What is DHKE

- Algo to securely exchange keys
- esp communicated over "insecure" channels

How do one-way functions work?

- Both users must have their private colors set, and must decide on a common public color to be used.
- The private colors and public colors are mixed together
- The color mixture is sent among the users.
- The private colors are now combined with the mixture to find actual secret color.
- Both will get single color



## Steps in key Exchange

### (1) choose $q \geqslant \alpha$

a. choose a prime number  $q$

b. select  $\alpha$  as a primitive root of  $q$

To be primitive root,

$$\alpha \text{ mod } q$$

$$\alpha^2 \text{ mod } q$$

$$\alpha^3 \text{ mod } q$$

:

$$\alpha^{q-1} \text{ mod } q$$

$$\leq p \mid q$$

$$(1, 2, 3, \dots, q-1)$$

### (2) Deriving the key Pair

Assume private key  $= X_a$   
where  $X_a < q$

public key ( $Y_a$ ) becomes:

$$Y_a = \alpha^{X_a} \text{ mod } q$$

key pair :  $\{X_a, Y_a\}$

Assume private key  $= X_b$   
where  $X_b < q$

public key ( $Y_b$ ) becomes:

$$Y_b = \alpha^{X_b} \text{ mod } q$$

key pair :  $\{X_b, Y_b\}$

### (3) key Generation

para  $X_a, Y_b, q$

secret key

$$K = (Y_b)^{X_a} \text{ mod } q$$

para :  $X_b, Y_a, q$

key

$$K = (Y_a)^{X_b} \text{ mod } q$$

$$q = 17$$

$$d = 9$$

$x_a$

$$x_a = 15$$

$(y_a)$

$$y_a = 3^{15} \bmod 17 = 6$$

key pair  $(15, 6)$

para :  $x_a, y_b, q$

$$k = 12^{15} \bmod 17 = 10$$

$$x_b = 13$$

$(y_b)$

$$(y_b) = 3^{13} \bmod 17 = 12$$

key pair  $(13, 12)$

$$k = 6^{13} \bmod 17 = 10$$

### Appl'n

- Public key Infrastructure
- SSL / TLS Handshake
- SSH (Secure Shell Access)