
MODULE *JustInTimePaxos*

EXTENDS *Naturals, Sequences, FiniteSets, TLC*

The set of *Paxos* replicas

CONSTANT *Replicas*

The set of *Paxos* clients

CONSTANT *Clients*

The maximum clock interval

CONSTANT *MaxClockInterval*

An empty value

CONSTANT *Nil*

Client request/response types+

CONSTANTS

WriteRequest,
WriteResponse,
ReadRequest,
ReadResponse

Server request/response types

CONSTANTS

ViewChangeRequest,
ViewChangeResponse,
StartViewRequest

Replica roles

CONSTANTS

NormalStatus,
ViewChangeStatus,
RecoveringStatus

VARIABLE *replicas*

globalVars \triangleq $\langle replicas \rangle$

VARIABLE *messages*

messageVars \triangleq $\langle messages \rangle$

VARIABLE *time*

VARIABLE *requestID*

VARIABLE *responses*

VARIABLE *writes*
 VARIABLE *reads*
 $clientVars \triangleq \langle time, requestID, responses, writes, reads \rangle$
 VARIABLE *status*
 VARIABLE *log*
 VARIABLE *viewID*
 VARIABLE *lastNormalView*
 VARIABLE *viewChanges*
 $replicaVars \triangleq \langle status, log, viewID, lastNormalView, viewChanges \rangle$
 VARIABLE *transitions*
 $vars \triangleq \langle globalVars, messageVars, clientVars, replicaVars, transitions \rangle$

Helpers

RECURSIVE *SeqFromSet*($-$)
 $SeqFromSet(S) \triangleq$
 IF $S = \{\}$ THEN $\langle \rangle$
 ELSE LET $x \triangleq \text{CHOOSE } x \in S : \text{TRUE}$
 IN $\langle x \rangle \circ SeqFromSet(S \setminus \{x\})$
 $Max(s) \triangleq \text{CHOOSE } x \in s : \forall y \in s : x \geq y$
 $IsQuorum(s) \triangleq Cardinality(s) * 2 \geq Cardinality(Replicas)$
 $Quorums \triangleq \{r \in \text{SUBSET } Replicas : IsQuorum(r)\}$
 $Primary(v) \triangleq replicas[(v \% Len(replicas)) + (\text{IF } v \geq Len(replicas) \text{ THEN } 1 \text{ ELSE } 0)]$
 $IsPrimary(r) \triangleq Primary(viewID[r]) = r$

Messaging helpers

$Sends(ms) \triangleq messages' = messages \cup ms$
 $Send(m) \triangleq Sends(\{m\})$
 $Replies(req, resps) \triangleq messages' = (messages \cup resps) \setminus \{req\}$
 $Reply(req, resp) \triangleq Replies(req, \{resp\})$

$$\text{Discard}(m) \triangleq \text{messages}' = \text{messages} \setminus \{m\}$$

$$\begin{aligned} \text{Write}(c) &\triangleq \\ &\wedge \text{time}' = \text{time} + 1 \\ &\wedge \text{requestID}' = [\text{requestID} \text{ EXCEPT } ![c] = \text{requestID}[c] + 1] \\ &\wedge \text{Sends}(\{[src \mapsto c, \\ &\quad \quad \quad dest \mapsto r, \\ &\quad \quad \quad type \mapsto \text{WriteRequest}, \\ &\quad \quad \quad requestID \mapsto \text{requestID}'[c], \\ &\quad \quad \quad timestamp \mapsto \text{time}'] : r \in \text{Replicas}\}) \\ &\wedge \text{UNCHANGED } \langle \text{globalVars}, \text{replicaVars}, \text{responses}, \text{writes}, \text{reads} \rangle \end{aligned}$$

$$\begin{aligned} \text{Read}(c) &\triangleq \\ &\wedge \text{requestID}' = [\text{requestID} \text{ EXCEPT } ![c] = \text{requestID}[c] + 1] \\ &\wedge \text{Sends}(\{[src \mapsto c, \\ &\quad \quad \quad dest \mapsto r, \\ &\quad \quad \quad type \mapsto \text{ReadRequest}, \\ &\quad \quad \quad requestID \mapsto \text{requestID}'[c]] : r \in \text{Replicas}\}) \\ &\wedge \text{UNCHANGED } \langle \text{globalVars}, \text{replicaVars}, \text{time}, \text{responses}, \text{writes}, \text{reads} \rangle \end{aligned}$$

$$\begin{aligned} \text{ChecksumsMatch}(c1, c2) &\triangleq \\ &\wedge \text{Len}(c1) = \text{Len}(c2) \\ &\wedge \neg \exists i \in \text{DOMAIN } c1 : c1[i] \neq c2[i] \end{aligned}$$

$$\begin{aligned} \text{IsCommitted}(acks) &\triangleq \\ &\exists \text{msgs} \in \text{SUBSET } acks : \\ &\quad \wedge \{m.\text{src} : m \in \text{msgs}\} \in \text{Quorums} \\ &\quad \wedge \exists m1 \in \text{msgs} : \forall m2 \in \text{msgs} : m1.\text{viewID} = m2.\text{viewID} \wedge \text{ChecksumsMatch}(m1.\text{checksum}, m2.\text{checksum}) \\ &\quad \wedge \exists m \in \text{msgs} : m.\text{primary} \end{aligned}$$

$$\begin{aligned} \text{HandleWriteResponse}(c, r, m) &\triangleq \\ &\wedge \neg \exists w \in \text{writes}[c] : w.\text{requestID} = m.\text{requestID} \\ &\wedge \vee \wedge m.\text{requestID} \notin \text{DOMAIN } \text{responses}[c][r] \\ &\quad \wedge \text{responses}' = [\text{responses} \text{ EXCEPT } ![c] = [\text{responses}[c] \text{ EXCEPT } ![r] = \text{responses}[c][r] @@ (m.\text{requestID})] \\ &\quad \wedge \text{UNCHANGED } \langle \text{writes} \rangle \\ &\vee \wedge m.\text{requestID} \in \text{DOMAIN } \text{responses}[c][r] \\ &\quad \text{Do not overwrite a response from a newer view} \\ &\quad \wedge \text{responses}[c][r][m.\text{requestID}].\text{viewID} \leq m.\text{viewID} \\ &\quad \wedge \text{responses}' = [\text{responses} \text{ EXCEPT } ![c] = [\text{responses}[c] \text{ EXCEPT } ![r] = [\text{responses}[c][r] \text{ EXCEPT } ![m.\text{requestID}]] \\ &\quad \wedge \text{LET } \text{committed} \triangleq \text{IsCommitted}(\{\text{responses}'[c][x][m.\text{requestID}] : x \in \{x \in \text{Replicas} : m.\text{requestID} \leq x.\text{viewID}\}\}) \\ &\quad \text{IN} \\ &\quad \vee \wedge \text{committed} \\ &\quad \quad \wedge \text{writes}' = [\text{writes} \text{ EXCEPT } ![c] = \text{writes}[c] \cup \{m\}] \\ &\quad \vee \wedge \neg \text{committed} \\ &\quad \quad \wedge \text{UNCHANGED } \langle \text{writes} \rangle \end{aligned}$$

$$\begin{aligned}
& \wedge \text{Discard}(m) \\
& \wedge \text{UNCHANGED } \langle \text{globalVars}, \text{replicaVars}, \text{time}, \text{requestID}, \text{reads} \rangle \\
\text{HandleReadResponse}(c, r, m) & \triangleq \\
& \wedge \vee \wedge m.\text{primary} \\
& \quad \wedge m \notin \text{reads}[c] \\
& \quad \wedge \text{reads}' = [\text{reads} \text{ EXCEPT } ![c] = \text{reads}[c] \cup \{m\}] \\
& \vee \wedge \neg m.\text{primary} \\
& \quad \wedge \text{UNCHANGED } \langle \text{reads} \rangle \\
& \wedge \text{Discard}(m) \\
& \wedge \text{UNCHANGED } \langle \text{globalVars}, \text{replicaVars}, \text{time}, \text{requestID}, \text{responses}, \text{writes} \rangle
\end{aligned}$$

Server request/response handling

$$\begin{aligned}
\text{HandleWriteRequest}(r, c, m) & \triangleq \\
& \wedge \text{status}[r] = \text{NormalStatus} \\
& \wedge \vee \wedge \vee \text{Len}(\log[r]) = 0 \\
& \quad \vee \wedge \text{Len}(\log[r]) \neq 0 \\
& \quad \wedge m.\text{timestamp} > \log[r][\text{Len}(\log[r])].\text{timestamp} \\
& \wedge \text{LET } \text{checksum} \triangleq \text{Append}([i \in \text{DOMAIN } \log[r] \mapsto \log[r][i].\text{timestamp}], m.\text{timestamp}) \\
& \quad \text{entry} \triangleq [\text{client} \mapsto c, \\
& \quad \quad \text{requestID} \mapsto m.\text{requestID}, \\
& \quad \quad \text{timestamp} \mapsto m.\text{timestamp}, \\
& \quad \quad \text{checksum} \mapsto \text{checksum}] \\
& \text{IN} \\
& \wedge \log' = [\log \text{ EXCEPT } ![r] = \text{Append}(\log[r], \text{entry})] \\
& \wedge \text{Reply}(m, [\text{src} \mapsto r, \\
& \quad \text{dest} \mapsto c, \\
& \quad \text{type} \mapsto \text{WriteResponse}, \\
& \quad \text{requestID} \mapsto m.\text{requestID}, \\
& \quad \text{viewID} \mapsto \text{viewID}[r], \\
& \quad \text{primary} \mapsto \text{IsPrimary}(r), \\
& \quad \text{index} \mapsto \text{Len}(\log'[r]), \\
& \quad \text{checksum} \mapsto \log'[r][\text{Len}(\log'[r])].\text{checksum}, \\
& \quad \text{succeeded} \mapsto \text{TRUE}]) \\
& \vee \wedge \text{Len}(\log[r]) \neq 0 \\
& \wedge m.\text{timestamp} \leq \log[r][\text{Len}(\log[r])].\text{timestamp} \\
& \wedge \text{Reply}(m, [\text{src} \mapsto r, \\
& \quad \text{dest} \mapsto c, \\
& \quad \text{type} \mapsto \text{WriteResponse}, \\
& \quad \text{requestID} \mapsto m.\text{requestID}, \\
& \quad \text{viewID} \mapsto \text{viewID}[r], \\
& \quad \text{primary} \mapsto \text{IsPrimary}(r), \\
& \quad \text{index} \mapsto \text{Len}(\log[r]),
\end{aligned}$$

$$\begin{aligned}
& checksum \mapsto \log[r][\text{Len}(\log[r])].checksum, \\
& succeeded \mapsto \text{FALSE}] \\
& \wedge \text{UNCHANGED } \langle \log \rangle \\
& \wedge \text{UNCHANGED } \langle globalVars, clientVars, status, viewID, lastNormalView, viewChanges \rangle \\
\text{HandleReadRequest}(r, c, m) & \triangleq \\
& \wedge status[r] = \text{NormalStatus} \\
& \wedge \text{Len}(\log[r]) > 0 \\
& \wedge \text{Reply}(m, [src \mapsto r, \\
& \quad dest \mapsto c, \\
& \quad type \mapsto \text{ReadResponse}, \\
& \quad requestID \mapsto m.requestID, \\
& \quad viewID \mapsto viewID[r], \\
& \quad primary \mapsto \text{IsPrimary}(r), \\
& \quad index \mapsto \text{Len}(\log[r]), \\
& \quad checksum \mapsto \log[r][\text{Len}(\log[r])].checksum, \\
& \quad succeeded \mapsto \text{TRUE}]) \\
& \wedge \text{UNCHANGED } \langle globalVars, clientVars, status, log, viewID, lastNormalView, viewChanges \rangle \\
\text{ChangeView}(r) & \triangleq \\
& \wedge \text{Sends}(\{[src \mapsto r, \\
& \quad dest \mapsto d, \\
& \quad type \mapsto \text{ViewChangeRequest}, \\
& \quad viewID \mapsto viewID[r] + 1] : d \in \text{Replicas}\}) \\
& \wedge \text{UNCHANGED } \langle globalVars, clientVars, replicaVars \rangle \\
\text{HandleViewChangeRequest}(r, s, m) & \triangleq \\
& \wedge viewID[r] < m.viewID \\
& \wedge viewID' = [viewID \text{ EXCEPT } ![r] = m.viewID] \\
& \wedge status' = [status \text{ EXCEPT } ![r] = \text{ViewChangeStatus}] \\
& \wedge viewChanges' = [viewChanges \text{ EXCEPT } ![r] = \{\}] \\
& \wedge \text{Reply}(m, [src \mapsto r, \\
& \quad dest \mapsto \text{Primary}(m.viewID), \\
& \quad type \mapsto \text{ViewChangeResponse}, \\
& \quad viewID \mapsto m.viewID, \\
& \quad lastNormal \mapsto lastNormalView[r], \\
& \quad log \mapsto \log[r]]) \\
& \wedge \text{UNCHANGED } \langle globalVars, clientVars, log, lastNormalView \rangle \\
\text{HandleViewChangeResponse}(r, s, m) & \triangleq \\
& \wedge \text{IsPrimary}(r) \\
& \wedge viewID[r] = m.viewID \\
& \wedge status[r] = \text{ViewChangeStatus} \\
& \wedge viewChanges' = [viewChanges \text{ EXCEPT } ![r] = viewChanges[r] \cup \{m\}] \\
& \wedge \text{LET} \\
& \quad isViewQuorum(vs) \triangleq \text{IsQuorum}(vs) \wedge \exists v \in vs : v.src = r
\end{aligned}$$

$$\begin{aligned}
newViewChanges &\triangleq \{v \in viewChanges'[r] : v.viewID = viewID[r]\} \\
normalViews &\triangleq \{v.lastNormal : v \in newViewChanges\} \\
lastNormal &\triangleq \text{CHOOSE } v \in normalViews : \forall v2 \in normalViews : v2 \leq v \\
goodLogs &\triangleq \{n.log : n \in \{v \in newViewChanges : v.lastNormal = lastNormal\}\} \\
combineLogs(ls) &\triangleq \\
\text{LET} \\
\quad indexLogs(i) &\triangleq \{l \in ls : Len(l) \geq i\} \\
\quad indexEntries(i) &\triangleq \{l[i] : l \in indexLogs(i)\} \\
\quad quorumLogs(i) &\triangleq \{L \in \text{SUBSET } indexLogs(i) : IsQuorum(L)\} \\
\quad isCommittedEntry(i, e) &\triangleq \forall L \in quorumLogs(i) : \\
&\quad \exists l \in L : \\
&\quad \quad ChecksumsMatch(e.checksum, l[i].checksum) \\
\quad isCommittedIndex(i) &\triangleq \exists e \in indexEntries(i) : isCommittedEntry(i, e) \\
\quad commit(i) &\triangleq \text{CHOOSE } e \in indexEntries(i) : isCommittedEntry(i, e) \\
\quad maxIndex &\triangleq \text{Max}(\{Len(l) : l \in ls\}) \\
\quad committedIndexes &\triangleq \{i \in 1 \dots maxIndex : isCommittedIndex(i)\} \\
\quad maxCommit &\triangleq \text{IF } Cardinality(committedIndexes) > 0 \text{ THEN } \text{Max}(committedIndexes) \\
\text{IN} \\
\quad [i \in 1 \dots maxCommit \mapsto commit(i)] \\
\text{IN} \\
\quad \vee \wedge isViewQuorum(newViewChanges) \\
\quad \wedge Replies(m, \{src \mapsto r, \\
\quad \quad \quad dest \mapsto d, \\
\quad \quad \quad type \mapsto StartViewRequest, \\
\quad \quad \quad viewID \mapsto viewID[r], \\
\quad \quad \quad log \mapsto combineLogs(goodLogs)\} : d \in Replicas\}) \\
\quad \vee \wedge \neg isViewQuorum(newViewChanges) \\
\quad \wedge Discard(m) \\
\wedge \text{UNCHANGED } \langle globalVars, clientVars, status, viewID, log, lastNormalView \rangle \\
HandleStartViewRequest(r, s, m) \triangleq \\
\quad \wedge \vee viewID[r] < m.viewID \\
\quad \vee \wedge viewID[r] = m.viewID \\
\quad \wedge status[r] = ViewChangeStatus \\
\quad \wedge log' = [log \text{ EXCEPT } ![r] = m.log] \\
\quad \wedge status' = [status \text{ EXCEPT } ![r] = NormalStatus] \\
\quad \wedge viewID' = [viewID \text{ EXCEPT } ![r] = m.viewID] \\
\quad \wedge lastNormalView' = [lastNormalView \text{ EXCEPT } ![r] = m.viewID] \\
\quad \wedge Discard(m) \\
\quad \wedge \text{UNCHANGED } \langle globalVars, clientVars, viewChanges \rangle
\end{aligned}$$

$$\begin{aligned}
InitMessageVars &\triangleq \\
&\wedge messages = \{\}
\end{aligned}$$

$$\begin{aligned}
InitClientVars &\triangleq \\
&\wedge time = 0 \\
&\wedge requestID = [c \in Clients \mapsto 0] \\
&\wedge responses = [c \in Clients \mapsto [r \in Replicas \mapsto [s \in \{\} \mapsto [index \mapsto 0, checksum \mapsto Nil]]]] \\
&\wedge writes = [c \in Clients \mapsto \{\}] \\
&\wedge reads = [c \in Clients \mapsto \{\}]
\end{aligned}$$

$$\begin{aligned}
InitReplicaVars &\triangleq \\
&\wedge replicas = SeqFromSet(Replicas) \\
&\wedge status = [r \in Replicas \mapsto NormalStatus] \\
&\wedge log = [r \in Replicas \mapsto \langle \rangle] \\
&\wedge viewID = [r \in Replicas \mapsto 1] \\
&\wedge lastNormalView = [r \in Replicas \mapsto 1] \\
&\wedge viewChanges = [r \in Replicas \mapsto \{\}]
\end{aligned}$$

$$\begin{aligned}
Init &\triangleq \\
&\wedge InitMessageVars \\
&\wedge InitClientVars \\
&\wedge InitReplicaVars \\
&\wedge transitions = 0
\end{aligned}$$

The type invariant checks that no read ever reads a different value than a previous write

$$\begin{aligned}
Inv &\triangleq \\
&\wedge \forall c1, c2 \in Clients : \\
&\quad \neg \exists r \in reads[c1] : \\
&\quad \quad \exists w \in writes[c2] : \\
&\quad \quad \quad \wedge r.index = w.index \\
&\quad \quad \quad \wedge \neg ChecksumsMatch(r.checksum, w.checksum) \\
&\wedge \forall c1, c2 \in Clients : \\
&\quad \neg \exists r1 \in reads[c1] : \\
&\quad \quad \exists r2 \in reads[c2] : \\
&\quad \quad \quad \wedge r1.index = r2.index \\
&\quad \quad \quad \wedge \neg ChecksumsMatch(r1.checksum, r2.checksum)
\end{aligned}$$

$$Transition \triangleq transitions' = transitions + 1$$

$$\begin{aligned}
Next &\triangleq \\
&\vee \exists c \in Clients : \\
&\quad \wedge Write(c) \\
&\quad \wedge Transition \\
&\vee \exists c \in Clients : \\
&\quad \wedge Read(c) \\
&\quad \wedge Transition \\
&\vee \exists r \in Replicas : \\
&\quad \wedge ChangeView(r)
\end{aligned}$$

$$\begin{aligned}
& \wedge \textit{Transition} \\
\vee \exists m \in \textit{messages} : & \\
& \wedge m.type = \textit{WriteRequest} \\
& \wedge \textit{HandleWriteRequest}(m.dest, m.src, m) \\
& \wedge \textit{Transition} \\
\vee \exists m \in \textit{messages} : & \\
& \wedge m.type = \textit{WriteResponse} \\
& \wedge \textit{HandleWriteResponse}(m.dest, m.src, m) \\
& \wedge \textit{Transition} \\
\vee \exists m \in \textit{messages} : & \\
& \wedge m.type = \textit{ReadRequest} \\
& \wedge \textit{HandleReadRequest}(m.dest, m.src, m) \\
& \wedge \textit{Transition} \\
\vee \exists m \in \textit{messages} : & \\
& \wedge m.type = \textit{ReadResponse} \\
& \wedge \textit{HandleReadResponse}(m.dest, m.src, m) \\
& \wedge \textit{Transition} \\
\vee \exists m \in \textit{messages} : & \\
& \wedge m.type = \textit{ViewChangeRequest} \\
& \wedge \textit{HandleViewChangeRequest}(m.dest, m.src, m) \\
& \wedge \textit{Transition} \\
\vee \exists m \in \textit{messages} : & \\
& \wedge m.type = \textit{ViewChangeResponse} \\
& \wedge \textit{HandleViewChangeResponse}(m.dest, m.src, m) \\
& \wedge \textit{Transition} \\
\vee \exists m \in \textit{messages} : & \\
& \wedge m.type = \textit{StartViewRequest} \\
& \wedge \textit{HandleStartViewRequest}(m.dest, m.src, m) \\
& \wedge \textit{Transition}
\end{aligned}$$

$$Spec \triangleq Init \wedge \Box[Next]_{vars}$$

\ * Modification History
\ * Last modified Tue Sep 22 01:06:09 PDT 2020 by jordanhalterman
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