## EXTENDS Naturals, Sequences, FiniteSets, TLC

The set of Paxos replicas CONSTANT Replicas

The set of *Paxos* clients CONSTANT *Clients* 

The maximum clock interval CONSTANT MaxClockInterval

An empty value CONSTANT Nil

Client request/response types+

CONSTANTS

WriteRequest, WriteResponse, ReadRequest,

ReadResponse

Server request/response types

CONSTANTS

 $ViewChangeRequest, \ ViewChangeResponse, \ StartViewRequest$ 

Replica roles

CONSTANTS

NormalStatus, ViewChangeStatus, RecoveringStatus

Variable replicas

 $globalVars \stackrel{\Delta}{=} \langle replicas \rangle$ 

Variable messages

 $messageVars \triangleq \langle messages \rangle$ 

 ${\tt VARIABLE}\ global Time$ 

Variable time

VARIABLE requestID

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VARIABLE responses
Variable writes
VARIABLE reads
clientVars \stackrel{\triangle}{=} \langle globalTime, time, requestID, responses, writes, reads \rangle
VARIABLE status
Variable log
VARIABLE viewID
Variable lastNormalView
VARIABLE viewChanges
replicaVars \triangleq \langle status, log, viewID, lastNormalView, viewChanges \rangle
Variable transitions
vars \triangleq \langle global Vars, message Vars, client Vars, replica Vars, transitions \rangle
 Helpers
RECURSIVE SeqFromSet(_)
SeqFromSet(S) \triangleq
  If S = \{\} Then \langle \rangle
   ELSE LET x \triangleq \text{CHOOSE } x \in S : \text{TRUE}
           IN \langle x \rangle \circ SeqFromSet(S \setminus \{x\})
Max(s) \stackrel{\triangle}{=} \text{ CHOOSE } x \in s : \forall y \in s : x \geq y
IsQuorum(s) \triangleq Cardinality(s) * 2 \geq Cardinality(Replicas)
Quorums \stackrel{\triangle}{=} \{r \in SUBSET \ Replicas : IsQuorum(r)\}
Primary(v) \triangleq replicas[(v\%Len(replicas)) + (\text{if } v \geq Len(replicas) \text{ THEN 1 ELSE 0})]
IsPrimary(r) \stackrel{\Delta}{=} Primary(viewID[r]) = r
 Messaging helpers
Sends(ms) \stackrel{\triangle}{=} messages' = messages \cup ms
Send(m) \triangleq Sends(\{m\})
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 $Replies(req, resps) \triangleq messages' = (messages \cup resps) \setminus \{req\}$ 

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Discard(m) \stackrel{\triangle}{=} messages' = messages \setminus \{m\}
AdvanceTime(c) \triangleq
           \land globalTime' = globalTime + 1
           \land IF time[c] < globalTime <math>\land globalTime - time[c] > MaxClockInterval THEN
                          time' = [time \ EXCEPT \ ![c] = globalTime' - MaxClockInterval]
                  ELSE
                          time' = [time \ EXCEPT \ ![c] = time[c] + 1]
CurrentTime(c) \stackrel{\triangle}{=} time'[c]
Write(c) \triangleq
           \wedge AdvanceTime(c)
           \land requestID' = [requestID \ EXCEPT \ ![c] = requestID[c] + 1]
           \land Sends(\{[src
                                                                   \mapsto c,
                                        dest
                                                                    \mapsto r,
                                       type
                                                                    \mapsto WriteRequest,
                                       requestID \mapsto requestID'[c],
                                       timestamp \mapsto CurrentTime(c) | : r \in Replicas \}
          ∧ UNCHANGED ⟨globalVars, replicaVars, responses, writes, reads⟩
Read(c) \triangleq
              \land requestID' = [requestID \ EXCEPT \ ![c] = requestID[c] + 1]
              \land Sends(\{[src
                                                                     \mapsto c,
                                           dest
                                                                     \mapsto r,
                                                                     \mapsto ReadRequest,
                                           type
                                           requestID \mapsto requestID'[c]]: r \in Replicas\}
              \land UNCHANGED \langle globalVars, replicaVars, globalTime, time, responses, writes, reads <math>\rangle
ChecksumsMatch(c1, c2) \triangleq
           \wedge Len(c1) = Len(c2)
           \land \neg \exists i \in \text{DOMAIN } c1: c1[i] \neq c2[i]
IsCommitted(acks) \triangleq
          \exists msqs \in \text{SUBSET} acks:
               \land \{m.src : m \in msgs\} \in Quorums
               \land \exists m1 \in msgs : \forall m2 \in msgs : m1.viewID = m2.viewID \land ChecksumsMatch(m1.checksum, m2.checksum, m2.checksum, m2.checksum, m3.checksum, m3.checksum,
               \wedge \exists m \in msgs : m.primary
HandleWriteResponse(c, r, m) \stackrel{\Delta}{=}
           \land \neg \exists \ w \in writes[c] : w.requestID = m.requestID
           \land \lor \land m.requestID \notin DOMAIN \ responses[c][r]
                        \land responses' = [responses \ EXCEPT \ ![c] = [responses[c] \ EXCEPT \ ![r] = responses[c][r] \ @@ (m.requestI)
                         \land UNCHANGED \langle writes \rangle
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 $Reply(req, resp) \stackrel{\Delta}{=} Replies(req, \{resp\})$ 

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\lor \land m.requestID \in DOMAIN \ responses[c][r]
                          Do not overwrite a response from a newer view
                        \land \ responses[c][r][m.requestID].viewID \ \leq m.viewID
                          \land \ responses' = [responses \ \ \texttt{EXCEPT} \ ! [c] = [responses[c] \ \ \texttt{EXCEPT} \ ! [r] \ \ = [responses[c] [r] \ \ \texttt{EXCEPT} \ ! [m. ] 
                         \land LET committed \stackrel{\triangle}{=} IsCommitted({responses'[c][x][m.requestID]}: x \in \{x \in Replicas : m.requestID\}
                              IN
                                       \vee \wedge committed
                                             \land writes' = [writes \ EXCEPT \ ![c] = writes[c] \cup \{m\}]
                                       \vee \wedge \neg committed
                                             \land UNCHANGED \langle writes \rangle
           \wedge Discard(m)
           ∧ UNCHANGED ⟨globalVars, replicaVars, globalTime, time, requestID, reads⟩
HandleReadResponse(c, r, m) \stackrel{\Delta}{=}
           \land \lor \land m.primary
                        \land m \notin reads[c]
                         \land reads' = [reads \ EXCEPT \ ![c] = reads[c] \cup \{m\}]
                  \lor \land \neg m.primary
                         \land UNCHANGED \langle reads \rangle
           \wedge Discard(m)
           \(\text{\text{UNCHANGED}}\)\(\langle \q \langle \langl
  Server request/response handling
Handle WriteRequest(r, c, m) \stackrel{\Delta}{=}
           \land status[r] = NormalStatus
           \wedge \vee \wedge \vee Len(log[r]) = 0
                               \vee \wedge Len(log[r]) \neq 0
                                      \land m.timestamp > log[r][Len(log[r])].timestamp
                        \land \text{ Let } checksum \stackrel{\triangle}{=} Append([i \in \text{DOMAIN } log[r] \mapsto log[r][i].timestamp), \ m.timestamp)
                                                                      \triangleq [client
                                                                                                           \mapsto c,
                                                                                 requestID \mapsto m.requestID,
                                                                                 timestamp \mapsto m.timestamp,
                                                                                 checksum \mapsto checksum
                                       \wedge log' = [log \ EXCEPT \ ![r] = Append(log[r], entry)]
                                       \land Reply(m, [src])
                                                                                                    \mapsto c,
                                                                                                    \mapsto WriteResponse,
                                                                         requestID \mapsto m.requestID,
                                                                         viewID \mapsto viewID[r],
                                                                         primary \mapsto IsPrimary(r),
                                                                                                    \mapsto Len(log'[r]),
                                                                         checksum \mapsto log'[r][Len(log'[r])].checksum,
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succeeded \mapsto TRUE
       \vee \wedge Len(log[r]) \neq 0
          \land m.timestamp \leq log[r][Len(log[r])].timestamp
          \land Reply(m, [src
                                     \mapsto r,
                         dest
                                     \mapsto c,
                                     \mapsto WriteResponse,
                          type
                         requestID \mapsto m.requestID,
                                     \mapsto viewID[r],
                         viewID
                         primary \mapsto IsPrimary(r),
                          index
                                     \mapsto Len(log[r]),
                         checksum \mapsto log[r][Len(log[r])].checksum,
                         succeeded \mapsto FALSE)
          \land UNCHANGED \langle log \rangle
    ∧ UNCHANGED ⟨qlobalVars, clientVars, status, viewID, lastNormalView, viewChanges⟩
HandleReadRequest(r, c, m) \stackrel{\Delta}{=}
    \wedge status[r] = NormalStatus
    \wedge Len(log[r]) > 0
    \land Reply(m, [src
                               \mapsto r,
                   dest
                               \mapsto c,
                   type
                               \mapsto ReadResponse,
                   requestID \mapsto m.requestID,
                   viewID
                              \mapsto viewID[r],
                   primary \mapsto IsPrimary(r),
                   index
                               \mapsto Len(log[r]),
                   checksum \mapsto log[r][Len(log[r])].checksum,
                   succeeded \mapsto TRUE
    ∧ UNCHANGED ⟨globalVars, clientVars, status, log, viewID, lastNormalView, viewChanges⟩
ChangeView(r) \triangleq
    \land Sends(\{[src
                          \mapsto r,
                 dest
                          \mapsto d.
                          \mapsto ViewChangeRequest,
                 viewID \mapsto viewID[r] + 1] : d \in Replicas\})
    ∧ UNCHANGED ⟨ global Vars, client Vars, replica Vars⟩
Handle View Change Request(r, s, m) \stackrel{\Delta}{=}
    \land viewID[r] < m.viewID
    \land viewID'
                       = [viewID \ EXCEPT \ ![r] = m.viewID]
                       = [status \ EXCEPT \ ![r] = ViewChangeStatus]
    \wedge status'
    \land viewChanges' = [viewChanges \ EXCEPT \ ![r] = \{\}]
    \land Reply(m, [src
                                 \mapsto r,
                                 \mapsto Primary(m.viewID),
                   dest
                                 \mapsto ViewChangeResponse,
                   type
                   viewID
                                 \mapsto m.viewID,
                   lastNormal \mapsto lastNormalView[r],
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\mapsto log[r]
     \land UNCHANGED \langle globalVars, clientVars, log, lastNormalView <math>\rangle
Handle View Change Response(r, s, m) \stackrel{\Delta}{=}
     \wedge IsPrimary(r)
     \land viewID[r]
                         = m.viewID
                         = ViewChangeStatus
     \wedge status[r]
     \land viewChanges' = [viewChanges \ EXCEPT \ ![r] = viewChanges[r] \cup \{m\}]
           isViewQuorum(vs) \triangleq IsQuorum(vs) \land \exists v \in vs : v.src = r
                                     \stackrel{\triangle}{=} \{ v \in viewChanges'[r] : v.viewID = viewID[r] \}
           newViewChanges
                                     \triangleq \{v.lastNormal : v \in newViewChanges\}
           normal Views \\
                                     \stackrel{\triangle}{=} Choose v \in normal Views : \forall v2 \in normal Views : v2 \leq v
           lastNormal
                                     \triangleq \{n.log: n \in \{v \in newViewChanges: v.lastNormal = lastNormal\}\}
           goodLogs
           combineLogs(ls)
               LET
                                                  \begin{array}{l} \triangleq \ \{l \in \mathit{ls} : \mathit{Len}(l) \geq i\} \\ \triangleq \ \{l[i] : l \in \mathit{indexLogs}(i)\} \end{array}
                  indexLogs(i)
                  indexEntries(i)
                                                  \triangleq \{L \in \text{SUBSET } indexLogs(i) : IsQuorum(L)\}
                  quorumLogs(i)
                  isCommittedEntry(i, e) \stackrel{\Delta}{=} \forall L \in quorumLogs(i) :
                                                        \exists l \in L:
                                                          ChecksumsMatch(e.checksum, l[i].checksum)
                                                  \triangleq \exists e \in indexEntries(i) : isCommittedEntry(i, e)
                  isCommittedIndex(i)
                                                  \stackrel{\triangle}{=} CHOOSE e \in indexEntries(i) : isCommittedEntry(i, e)
                  commit(i)
                                                  \stackrel{\Delta}{=} Max(\{Len(l): l \in ls\})
                  maxIndex
                                                  \triangleq \{i \in 1 ... maxIndex : isCommittedIndex(i)\}
                  committed Indexes
                                                  \stackrel{\triangle}{=} IF Cardinality(committedIndexes) > 0 THEN Max(committedIndexes)
                  maxCommit
              IN
                  [i \in 1 ... maxCommit \mapsto commit(i)]
       IN
            \lor \land isViewQuorum(newViewChanges)
               \land Replies(m, \{[src
                                   dest
                                            \mapsto d,
                                            \mapsto StartViewRequest,
                                  viewID \mapsto viewID[r],
                                            \mapsto combineLogs(goodLogs)]: d \in Replicas})
                                  loq
            \lor \land \neg isViewQuorum(newViewChanges)
               \wedge Discard(m)
     ∧ UNCHANGED ⟨globalVars, clientVars, status, viewID, log, lastNormalView⟩
HandleStartViewRequest(r, s, m) \triangleq
     \land \lor viewID[r] < m.viewID
        \lor \land viewID[r] = m.viewID
           \wedge status[r]
                             = ViewChangeStatus
     \wedge log'
                             = [log \ EXCEPT \ ![r] = m.log]
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\land \mathit{status'}
                               = [status \ EXCEPT \ ![r] = NormalStatus]
                               = [viewID \ EXCEPT \ ![r] = m.viewID]
     \land viewID'
     \land lastNormalView' = [lastNormalView \ EXCEPT \ ![r] = m.viewID]
     \wedge Discard(m)
     \land UNCHANGED \langle globalVars, clientVars, viewChanges \rangle
InitMessageVars \triangleq
     \land messages = \{\}
InitClientVars \triangleq
     \land globalTime = 0
     \land time
                       = [c \in Clients \mapsto 0]
     \land requestID = [c \in Clients \mapsto 0]
     \land responses = [c \in Clients \mapsto [r \in Replicas \mapsto [s \in \{\} \mapsto [index \mapsto 0, checksum \mapsto Nil]]]]
                       = [c \in Clients \mapsto \{\}]
     \land writes
                       = [c \in Clients \mapsto \{\}]
     \land reads
InitReplicaVars \triangleq
     \land replicas
                              = SeqFromSet(Replicas)
     \wedge status
                              = [r \in Replicas \mapsto NormalStatus]
     \land log
                              = [r \in Replicas \mapsto \langle \rangle]
                              = [r \in Replicas \mapsto 1]
     \land viewID
     \land lastNormalView = [r \in Replicas \mapsto 1]
     \land viewChanges
                              = [r \in Replicas \mapsto \{\}]
Init \triangleq
     \wedge InitMessageVars
     \land InitClientVars
     \land \ InitReplica Vars
     \land \ transitions = 0
 The type invariant checks that no read ever reads a different value than a previous write
Inv \triangleq
    \land \forall c1, c2 \in Clients:
         \neg \exists r \in reads[c1]:
             \exists w \in writes[c2]:
                \land r.index = w.index
                \land \neg ChecksumsMatch(r.checksum, w.checksum)
    \land \forall c1, c2 \in Clients:
         \neg \exists r1 \in reads[c1]:
             \exists r2 \in reads[c2]:
                \wedge r1.index = r2.index
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```
\land \neg ChecksumsMatch(r1.checksum, r2.checksum)
Transition \stackrel{\triangle}{=} transitions' = transitions + 1
Next \triangleq
     \vee \exists c \in Clients:
          \land Write(c)
          \land \ Transition
     \vee \exists c \in Clients:
          \land Read(c)
          \land Transition
     \vee \exists r \in Replicas :
          \land ChangeView(r)
          \land Transition
     \vee \exists m \in messages :
          \land m.type = WriteRequest
          \land Handle WriteRequest(m.dest, m.src, m)
          \land Transition
     \vee \exists m \in messages :
          \land m.type = WriteResponse
          \land Handle WriteResponse (m.dest, m.src, m)
          \land Transition
     \vee \exists m \in messages :
          \land m.type = ReadRequest
          \land HandleReadRequest(m.dest, m.src, m)
          \land Transition
     \vee \exists m \in messages :
          \land m.type = ReadResponse
          \land HandleReadResponse(m.dest, m.src, m)
          \land Transition
     \vee \exists m \in messages :
          \land \ m.type = \textit{ViewChangeRequest}
          \land Handle View Change Request (m.dest, m.src, m)
          \land Transition
     \vee \exists m \in messages :
          \land m.type = ViewChangeResponse
          \land Handle View Change Response (m.dest, m.src, m)
          \land \ \mathit{Transition}
     \vee \exists m \in messages :
          \land \ m.type = StartViewRequest
          \land HandleStartViewRequest(m.dest, m.src, m)
          \land Transition
Spec \triangleq Init \wedge \Box [Next]_{vars}
```

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