
MODULE *JustInTimePaxos*

EXTENDS *Naturals, Sequences, FiniteSets, TLC*

The set of *Paxos* replicas
 CONSTANT *Replicas*

The set of *Paxos* clients
 CONSTANT *Clients*

The maximum clock interval
 CONSTANT *MaxClockInterval*

An empty value
 CONSTANT *Nil*

Client request/response types+
 CONSTANTS
 WriteRequest,
 WriteResponse,
 ReadRequest,
 ReadResponse

Server request/response types
 CONSTANTS
 SlotLookup
 SlotLookupRep
 GapCommit
 GapCommitRep
 ViewChangeRequest, *ViewChangeReq*
 ViewChangeResponse, *ViewChange*
 StartViewRequest, *StartView*
 SyncPrepareRequest, *SyncPrepare*
 SyncPrepareResponse, *SyncPrepareRep*
 SyncCommitRequest *SyncCommit*

Replica roles
 CONSTANTS
 NormalStatus,
 ViewChangeStatus,
 RecoveringStatus

VARIABLE *replicas*

globalVars \triangleq $\langle \textit{replicas} \rangle$

VARIABLE *messages*

$messageVars \triangleq \langle messages \rangle$
 VARIABLE $globalTime$
 VARIABLE $time$
 VARIABLE $requestID$
 VARIABLE $responses$
 VARIABLE $writes$
 VARIABLE $reads$
 $clientVars \triangleq \langle globalTime, time, requestID, responses, writes, reads \rangle$
 VARIABLE $status$
 VARIABLE $viewID$
 VARIABLE log
 $replicaVars \triangleq \langle status, viewID, log \rangle$
 VARIABLE $transitions$
 $vars \triangleq \langle globalVars, messageVars, clientVars, replicaVars, transitions \rangle$

Helpers

RECURSIVE $SeqFromSet(-)$
 $SeqFromSet(S) \triangleq$
 IF $S = \{\}$ THEN $\langle \rangle$
 ELSE LET $x \triangleq \text{CHOOSE } x \in S : \text{TRUE}$
 IN $\langle x \rangle \circ SeqFromSet(S \setminus \{x\})$
 $Quorums \triangleq \{r \in \text{SUBSET } Replicas : Cardinality(r) * 2 > Cardinality(Replicas)\}$
 $Primary(v) \triangleq replicas[(v \% Len(replicas)) + (\text{IF } v \geq Len(replicas) \text{ THEN } 1 \text{ ELSE } 0)]$
 $IsPrimary(r) \triangleq Primary(viewID[r]) = r$

Messaging helpers

$Sends(m) \triangleq messages' = messages \cup m$
 $Send(m) \triangleq Sends(\{m\})$
 $Reply(req, res) \triangleq messages' = messages \setminus \{req, res\}$

$Discard(m) \triangleq messages' = messages \setminus \{m\}$

$AdvanceTime(c) \triangleq$

$\wedge globalTime' = globalTime + 1$

$\wedge \text{IF } time[c] < globalTime \wedge globalTime - time[c] > MaxClockInterval \text{ THEN}$
 $time' = [time \text{ EXCEPT } ![c] = globalTime' - MaxClockInterval]$

ELSE

$time' = [time \text{ EXCEPT } ![c] = time[c] + 1]$

$CurrentTime(c) \triangleq time'[c]$

$Write(c) \triangleq$

$\wedge AdvanceTime(c)$

$\wedge requestID' = [requestID \text{ EXCEPT } ![c] = requestID[c] + 1]$

$\wedge Sends(\{[source \mapsto c,$
 $target \mapsto r,$
 $type \mapsto WriteRequest,$
 $requestID \mapsto requestID'[c],$
 $timestamp \mapsto CurrentTime(c)] : r \in Replicas\})$

$\wedge \text{UNCHANGED } \langle globalVars, replicaVars, responses, writes, reads \rangle$

$Read(c) \triangleq$

$\wedge requestID' = [requestID \text{ EXCEPT } ![c] = requestID[c] + 1]$

$\wedge Sends(\{[source \mapsto c,$
 $target \mapsto r,$
 $type \mapsto ReadRequest,$
 $requestID \mapsto requestID'[c]] : r \in Replicas\})$

$\wedge \text{UNCHANGED } \langle globalVars, replicaVars, globalTime, time, responses, writes, reads \rangle$

$IsCommitted(acks) \triangleq$

$\exists msgs \in \text{SUBSET } acks :$

$\wedge \{m.source : m \in msgs\} \in Quorums$

$\wedge \exists m1 \in msgs : \forall m2 \in msgs : m1.viewID = m2.viewID \wedge m1.checksum \setminus m2.checksum = \{\}$

$\wedge \exists m \in msgs : m.primary$

$HandleWriteResponse(c, r, m) \triangleq$

$\wedge \neg \exists w \in writes[c] : w.requestID = m.requestID$

$\wedge \vee \wedge m.requestID \notin \text{DOMAIN } responses[c][r]$

$\wedge responses' = [responses \text{ EXCEPT } ![c] = [responses[c] \text{ EXCEPT } ![r] = responses[c][r] @@ (m.requestID$

$\wedge \text{UNCHANGED } \langle writes \rangle$

$\vee \wedge m.requestID \in \text{DOMAIN } responses[c][r]$

Do not overwrite a response from a newer view

$\wedge responses[c][r][m.requestID].viewID \leq m.viewID$

$\wedge responses' = [responses \text{ EXCEPT } ![c] = [responses[c] \text{ EXCEPT } ![r] = [responses[c][r] \text{ EXCEPT } ![m$

$\wedge \text{LET } committed \triangleq IsCommitted(\{responses'[c][x][m.requestID] : x \in \{x \in Replicas : m.requestID$

$$\begin{aligned}
& \text{IN} \\
& \quad \vee \wedge \text{committed} \\
& \quad \quad \wedge \text{writes}' = [\text{writes} \text{ EXCEPT } ![c] = \text{writes}[c] \cup \{m\}] \\
& \quad \vee \wedge \neg \text{committed} \\
& \quad \quad \wedge \text{UNCHANGED } \langle \text{writes} \rangle \\
& \wedge \text{Discard}(m) \\
& \wedge \text{UNCHANGED } \langle \text{globalVars}, \text{messageVars}, \text{replicaVars}, \text{globalTime}, \text{time}, \text{requestID}, \text{reads} \rangle \\
\text{HandleReadResponse}(c, r, m) & \triangleq \\
& \wedge \vee \wedge m.\text{primary} \\
& \quad \wedge m \notin \text{reads}[c] \\
& \quad \wedge \text{reads}' = [\text{reads} \text{ EXCEPT } ![c] = \text{reads}[c] \cup \{m\}] \\
& \vee \wedge \neg m.\text{primary} \\
& \quad \wedge \text{UNCHANGED } \langle \text{reads} \rangle \\
& \wedge \text{Discard}(m) \\
& \wedge \text{UNCHANGED } \langle \text{globalVars}, \text{replicaVars}, \text{globalTime}, \text{time}, \text{requestID}, \text{writes} \rangle
\end{aligned}$$

Server request/response handling

$$\begin{aligned}
\text{HandleWriteRequest}(r, c, m) & \triangleq \\
& \wedge \text{status}[r] = \text{NormalStatus} \\
& \wedge \vee \wedge \vee \text{Len}(\log[r]) = 0 \\
& \quad \vee \wedge \text{Len}(\log[r]) \neq 0 \\
& \quad \quad \wedge m.\text{timestamp} > \log[r][\text{Len}(\log[r])].\text{timestamp} \\
& \wedge \log' = [\log \text{ EXCEPT } ![r] = \text{Append}(\log[r], m)] \\
& \wedge \text{Reply}(m, [\text{source} \mapsto r, \\
& \quad \text{target} \mapsto c, \\
& \quad \text{type} \mapsto \text{WriteResponse}, \\
& \quad \text{requestID} \mapsto m.\text{requestID}, \\
& \quad \text{viewID} \mapsto \text{viewID}[r], \\
& \quad \text{primary} \mapsto \text{IsPrimary}(r), \\
& \quad \text{index} \mapsto \text{Len}(\log'[r]), \\
& \quad \text{checksum} \mapsto \{\log'[r][i].\text{timestamp} : i \in \text{DOMAIN } \log'[r]\}, \\
& \quad \text{succeeded} \mapsto \text{TRUE}]) \\
& \vee \wedge \text{Len}(\log[r]) \neq 0 \\
& \quad \wedge m.\text{timestamp} \leq \log[r][\text{Len}(\log[r])].\text{timestamp} \\
& \wedge \text{Reply}(m, [\text{source} \mapsto r, \\
& \quad \text{target} \mapsto c, \\
& \quad \text{type} \mapsto \text{WriteResponse}, \\
& \quad \text{requestID} \mapsto m.\text{requestID}, \\
& \quad \text{viewID} \mapsto \text{viewID}[r], \\
& \quad \text{primary} \mapsto \text{IsPrimary}(r), \\
& \quad \text{index} \mapsto \text{Len}(\log[r]), \\
& \quad \text{checksum} \mapsto \{\log[r][i].\text{timestamp} : i \in \text{DOMAIN } \log[r]\},
\end{aligned}$$

$$\begin{aligned}
& \text{ succeeded } \mapsto \text{FALSE}] \\
& \wedge \text{ UNCHANGED } \langle \log \rangle \\
& \wedge \text{ UNCHANGED } \langle \text{globalVars}, \text{clientVars}, \text{status}, \text{viewID} \rangle \\
\text{HandleReadRequest}(r, c, m) & \triangleq \\
& \wedge \text{ status}[r] = \text{NormalStatus} \\
& \wedge \text{ Len}(\log[r]) > 0 \\
& \wedge \text{ Reply}(m, [\text{source} \mapsto r, \\
& \quad \text{target} \mapsto c, \\
& \quad \text{type} \mapsto \text{WriteResponse}, \\
& \quad \text{requestID} \mapsto m.\text{requestID}, \\
& \quad \text{viewID} \mapsto \text{viewID}[r], \\
& \quad \text{primary} \mapsto \text{IsPrimary}(r), \\
& \quad \text{index} \mapsto \text{Len}(\log[r]), \\
& \quad \text{checksum} \mapsto \{\log[r][i].\text{timestamp} : i \in \text{DOMAIN } \log[r]\}, \\
& \quad \text{succeeded} \mapsto \text{TRUE}]) \\
& \wedge \text{ UNCHANGED } \langle \text{globalVars}, \text{clientVars}, \text{status}, \text{viewID}, \log \rangle \\
\text{HandleViewChangeRequest}(r, s, m) & \triangleq \\
& \wedge \text{ FALSE} \\
& \wedge \text{ UNCHANGED } \langle \text{globalVars}, \text{messageVars}, \text{clientVars}, \text{replicaVars} \rangle \\
\text{HandleViewChangeResponse}(r, s, m) & \triangleq \\
& \wedge \text{ FALSE} \\
& \wedge \text{ UNCHANGED } \langle \text{globalVars}, \text{messageVars}, \text{clientVars}, \text{replicaVars} \rangle \\
\text{HandleStartViewRequest}(r, s, m) & \triangleq \\
& \wedge \text{ FALSE} \\
& \wedge \text{ UNCHANGED } \langle \text{globalVars}, \text{messageVars}, \text{clientVars}, \text{replicaVars} \rangle \\
\text{HandleSyncPrepareRequest}(r, s, m) & \triangleq \\
& \wedge \text{ FALSE} \\
& \wedge \text{ UNCHANGED } \langle \text{globalVars}, \text{messageVars}, \text{clientVars}, \text{replicaVars} \rangle \\
\text{HandleSyncPrepareResponse}(r, s, m) & \triangleq \\
& \wedge \text{ FALSE} \\
& \wedge \text{ UNCHANGED } \langle \text{globalVars}, \text{messageVars}, \text{clientVars}, \text{replicaVars} \rangle \\
\text{HandleSyncCommitRequest}(r, s, m) & \triangleq \\
& \wedge \text{ FALSE} \\
& \wedge \text{ UNCHANGED } \langle \text{globalVars}, \text{messageVars}, \text{clientVars}, \text{replicaVars} \rangle
\end{aligned}$$

$$\begin{aligned}
\text{InitMessageVars} & \triangleq \\
& \wedge \text{ messages} = \{\}
\end{aligned}$$

$$\begin{aligned}
InitClientVars &\triangleq \\
&\wedge globalTime = 0 \\
&\wedge time = [c \in Clients \mapsto 0] \\
&\wedge requestID = [c \in Clients \mapsto 0] \\
&\wedge responses = [c \in Clients \mapsto [r \in Replicas \mapsto [s \in \{\} \mapsto [index \mapsto 0, checksum \mapsto Nil]]]] \\
&\wedge writes = [c \in Clients \mapsto \{\}] \\
&\wedge reads = [c \in Clients \mapsto \{\}] \\
\\
InitReplicaVars &\triangleq \\
&\wedge replicas = SeqFromSet(Replicas) \\
&\wedge status = [r \in Replicas \mapsto NormalStatus] \\
&\wedge viewID = [r \in Replicas \mapsto 1] \\
&\wedge log = [r \in Replicas \mapsto \langle \rangle] \\
\\
Init &\triangleq \\
&\wedge InitMessageVars \\
&\wedge InitClientVars \\
&\wedge InitReplicaVars \\
&\wedge transitions = 0
\end{aligned}$$

The type invariant checks that no read ever reads a different value than a previous write

$$\begin{aligned}
Inv &\triangleq \forall c1, c2 \in Clients : \\
&\quad \neg \exists r \in reads[c1] : \\
&\quad \quad \exists w \in writes[c2] : \\
&\quad \quad \quad r.index = w.index \wedge r.requestID \neq w.requestID
\end{aligned}$$

$$Transition \triangleq transitions' = transitions + 1$$

$$\begin{aligned}
Next &\triangleq \\
&\vee \exists c \in Clients : \\
&\quad \wedge Write(c) \\
&\quad \wedge transitions' = transitions + 1 \\
&\vee \exists c \in Clients : \\
&\quad \wedge Read(c) \\
&\quad \wedge Transition \\
&\vee \exists m \in messages : \\
&\quad \wedge m.type = WriteRequest \\
&\quad \wedge HandleWriteRequest(m.target, m.source, m) \\
&\quad \wedge Transition \\
&\vee \exists m \in messages : \\
&\quad \wedge m.type = WriteResponse \\
&\quad \wedge HandleWriteResponse(m.target, m.source, m) \\
&\quad \wedge Transition \\
&\vee \exists m \in messages : \\
&\quad \wedge m.type = ReadRequest
\end{aligned}$$

$$\begin{aligned}
& \wedge \text{HandleReadRequest}(m.target, m.source, m) \\
& \wedge \text{Transition} \\
\vee \exists m \in \text{messages} : & \\
& \wedge m.type = \text{ReadResponse} \\
& \wedge \text{HandleReadResponse}(m.target, m.source, m) \\
& \wedge \text{Transition} \\
\vee \exists m \in \text{messages} : & \\
& \wedge m.type = \text{ViewChangeRequest} \\
& \wedge \text{HandleViewChangeRequest}(m.target, m.source, m) \\
& \wedge \text{Transition} \\
\vee \exists m \in \text{messages} : & \\
& \wedge m.type = \text{ViewChangeResponse} \\
& \wedge \text{HandleViewChangeResponse}(m.target, m.source, m) \\
& \wedge \text{Transition} \\
\vee \exists m \in \text{messages} : & \\
& \wedge m.type = \text{StartViewRequest} \\
& \wedge \text{HandleStartViewRequest}(m.target, m.source, m) \\
& \wedge \text{Transition} \\
\vee \exists m \in \text{messages} : & \\
& \wedge m.type = \text{SyncPrepareRequest} \\
& \wedge \text{HandleSyncPrepareRequest}(m.target, m.source, m) \\
& \wedge \text{Transition} \\
\vee \exists m \in \text{messages} : & \\
& \wedge m.type = \text{SyncPrepareResponse} \\
& \wedge \text{HandleSyncPrepareResponse}(m.target, m.source, m) \\
& \wedge \text{Transition} \\
\vee \exists m \in \text{messages} : & \\
& \wedge m.type = \text{SyncCommitRequest} \\
& \wedge \text{HandleSyncCommitRequest}(m.target, m.source, m) \\
& \wedge \text{Transition}
\end{aligned}$$

$$Spec \triangleq Init \wedge \Box[Next]_{vars}$$

\ * Modification History
\ * Last modified Sun Sep 20 19:08:16 PDT 2020 by jordanhalterman
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