## THE UNIVERSITY OF MELBOURNE SCHOOL OF COMPUTING AND INFORMATION SYSTEMS COMP90020 DISTRIBUTED ALGORITHMS

## **Tutorial Week 3: Logical Clocks**

## **Notes**

Clock Condition. For all events a, b: if  $a \to b$  then C(a) < C(B). This is satisfied if the following two conditions hold:

- C1. If a and b are events in process  $p_i$  and a comes before b, then  $C_i(a) < C_i(b)$
- C2. If a is the sending of a message by process  $P_i$  and b is the receipt of that message by process  $P_j$ , then  $C_i(a) < C_j(b)$

## **Exercises**

- 10. This question is from last weeks tutorial, please attempt if you haven't already. An NTP server B receives server A's message at 16:34:23.480 bearing a timestamp 16:34:13.430 and replies to it. A receives the message at 16:34:15.725, bearing B's timestamp 16:34:25.7. Estimate the offset between B and A and the accuracy of the estimate.
- 11. Consider the following sequence of events at processes  $p_0, p_1, p_2$  and  $p_3$ . Here  $s_i$  and  $r_i$  are corresponding send and receive events for all i, while a and b are internal events.

Use Lamport's logical clock to assign clock values to these events (A diagram may help).

- 12. By considering a chain of zero or more messages connecting events e and e' and using induction, show that  $e \to e' \Rightarrow L(e) < L(e')$
- 13. Using the same sequence of events from exercise 10, draw a send/receive diagram to give Vector timestamps to each of the events.
- 14. Show that:

(a) 
$$V_j[i] \leq V_i[i]$$
 for all  $i, j$ 

(b) 
$$e \rightarrow e' \Rightarrow V[e] < V[e']$$

(c) 
$$V(e) < V(e') \Rightarrow e \rightarrow e'$$