# Adaptive maximization of social welfare

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#### Introduction

#### How should a policymaker act,

- who aims to maximize social welfare,
  - Weighted sum of utility.
  - ⇒ Tradeoff redistribution vs. cost of behavioral responses.
- and needs to learn agent responses to policy choices?
  - Adaptively updated policy choices
  - ⇒ Tradeoff exploration vs. exploitation.

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   Adaptively updated policy choices.
  - $\Rightarrow$  Tradeoff exploration vs. exploitation.

#### Taxes and bandits

- Optimal tax theory
  - Mirrlees (1971); Saez (2001); Chetty (2009)
- Multi-armed bandits
  - Bubeck and Cesa-Bianchi (2012); Lattimore and Szepesvári (2020)
- This talk: Merging bandits and welfare economics.
  - Unobserved welfare, as in optimal taxation.
  - Unknown responses, as in multi-armed bandits.

#### Co-authors

- Nicolò Cesa-Bianchi and Roberto Colomboni, for the theory of adversarial and stochastic lower and upper bounds on regret.
- Frederik Schwertner, for implementation of an adaptive basic income experiment in Germany.

## Setup

Lower and upper bounds on regret

Comparison to related learning problems

Simulations

## Setup: Tax on a binary choice

Each time period  $i = 1, 2, \dots, T$ :

- One agent with willingness to pay  $v_i \in [0, 1]$ .
- Choices:
  - Tax rate  $x_i \in [0, 1]$ .
  - Individual response function:  $G_i(x) = \mathbf{1}(x \le v_i)$
  - Binary agent decision  $y_i = G_i(x_i)$ .
- Observability:
  - After period i, we observe  $y_i$ .
  - We do *not* observe welfare  $U_i(x_i)$ .

### Social welfare

Weighted sum of public revenue and private welfare:

$$U_i(x_i) = \underbrace{x_i \cdot \mathbf{1}(x_i \leq v_i)}_{ ext{Public revenue}} + \lambda \cdot \underbrace{\max(v_i - x_i, 0)}_{ ext{Private welfare}}.$$

We can rewrite private welfare as an integral (consumer surplus):

$$U_i(x) = \underbrace{x \cdot G_i(x)}_{ ext{Public revenue}} + \lambda \cdot \underbrace{\int_x^1 G_i(x') dx'}_{ ext{Private welfare}}.$$

## Cumulative demand, welfare and regret

Cumulative demand:

$$\mathbb{G}_T(x) = \sum_{i < T} G_i(x).$$

• Cumulative welfare for a constant policy x:

$$\mathbb{U}_{T}(x) = \sum_{i < T} \mathbb{U}_{i}(x) = x \cdot \mathbb{G}_{T}(x) + \lambda \int_{x}^{1} \mathbb{G}_{T}(x') dx'.$$

• Cumulative welfare for the policies  $x_i$  actually chosen:

$$\mathbb{U}_T = \sum_{i \leq T} \mathbb{U}_i(x_i).$$

Adversarial regret:

$$\mathcal{R}_{T}(\lbrace v_{i}\rbrace_{i=1}^{T}) = \sup_{x} E\left[\mathbb{U}_{T}(x) - \mathbb{U}_{T} \middle| \lbrace v_{i}\rbrace_{i=1}^{T}\right].$$

## The structure of observability

Choice  $x_i$  reveals  $G_i(x_i)$ . But

$$U_i(x) - U_i(x') = \left[x \cdot G_i(x) - x' \cdot G_i(x')\right] + \lambda \int_{x}^{x'} G_i(x'') dx''$$

depends on values of  $G_i(x'')$  for  $x'' \in [x, x']!$ 

Different from standard adaptive decision-making problems:

- Multi-armed bandits:
   Observe welfare for the choice made.
- Online learning:
   Observe welfare for all possible choices.
- Online convex optimization:
   Observe gradient of welfare for the choice made.

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## Lower bound on regret

#### Theorem

There exists a constant C > 0 such that, for any randomized algorithm for the choice of  $x_1, x_2, \ldots$  and any time horizon  $T \in \mathbb{N}$ :

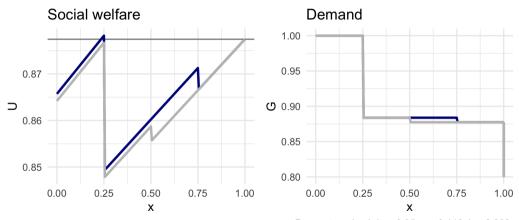
There exists a sequence  $(v_1, \ldots, v_T)$  for which

$$\mathcal{R}_T(\{v_i\}_{i=1}^T) \geq C \cdot T^{2/3}.$$

## Sketch of proof: Lower bound on regret

- Stochastic regret ≤ adversarial regret. (Since average ≤ maximum.)
- Construct a distribution for v with 4 points of support, e.g.  $(\frac{1}{4}, \frac{1}{2}, \frac{3}{4}, 1)$ .
- Choose the probability of each of these points such that
  - 1. The two middle points are far from optimal.
  - Learning which of the two end points is optimal requires sampling from the middle. (Because of the integral term.)

## Construction for the proof of the lower bound



Parameters: lambda = 0.95, a = 0.116, b = 0.003.

## Tempered Exp3 for social welfare

**Require:** Tuning parameters K,  $\gamma$  and  $\eta$ .

- 1: Set  $\tilde{x}_k = (k-1)/K$ , initialize  $\hat{\mathbb{G}}_k = 0$  for  $k = 1, \dots, K+1$ .
- 2: **for** individual i = 1, 2, ..., T **do**
- 3: **for** gridpoint k = 1, 2, ..., K + 1 **do**
- 4: Set

$$\widehat{\mathbb{U}}_{ik} = \widetilde{\mathbf{x}}_k \cdot \widehat{\mathbb{G}}_{ik} + \frac{\lambda}{K} \cdot \sum_{k' > k} \widehat{\mathbb{G}}_{ik'}, \quad p_{ik} = (1 - \gamma) \cdot \frac{\exp(\eta \cdot \widehat{\mathbb{U}}_{ik})}{\sum_{k'} \exp(\eta \cdot \widehat{\mathbb{U}}_{ik'})} + \frac{\gamma}{K + 1}.$$

- 5: end for
- 6: Choose  $k_i$  at random according to the probability distribution  $(p_1, \dots, p_{K+1})$ .
- 7: Set  $\mathbf{x}_i = \tilde{\mathbf{x}}_{k_i}$ , and query  $\mathbf{y}_i$  accordingly.
- 8: Update

$$\hat{\mathbb{G}}_{k_i} = \hat{\mathbb{G}}_{k_i} + \frac{y_i}{p_{ik_i}}.$$

9: **end for** 

## Upper bound on regret

#### Theorem

Consider the algorithm "Tempered Exp3 for social welfare." There exists a constant C' and choices for  $K, \gamma, \eta$  such that, for any sequence  $(v_1, \ldots, v_T)$ ,

$$\mathcal{R}_T(\{v_i\}_{i=1}^T) \leq C' \cdot \log(T)^{1/3} \cdot T^{2/3}.$$

⇒ Same rate as the lower bound, up to the logarithmic term!

Sketch of proof

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## Comparison to related learning problems

#### • Monopoly pricing:

Monopolist profits:

$$U_i^{MP}(x) = \underbrace{x \cdot G_i(x)}_{\text{Monopolist revenue}}$$

Easier – like a continuous multi-armed bandit.

#### Bilateral trade:

Buyer plus seller welfare:

$$U_i^{BT}(x) = G_i^b(x) \cdot \underbrace{\int_0^x G_i^s(x') dx'}_{ ext{Seller welfare}} + G_i^s(x) \cdot \underbrace{\int_x^1 G_i^b(x') dx'}_{ ext{Buyer welfare}}.$$

Harder – even gradients depend on global information.

## Comparison of regret rates

Model	Continuous	Discrete
Monopoly price setting Optimal tax Bilateral trade	T <sup>2/3</sup> T <sup>2/3</sup> T	$T^{1/2}$ $T^{2/3}$ $T^{2/3}$

- Rates are up to logarithmic terms.
- They reflect the different information structures in the three problems.

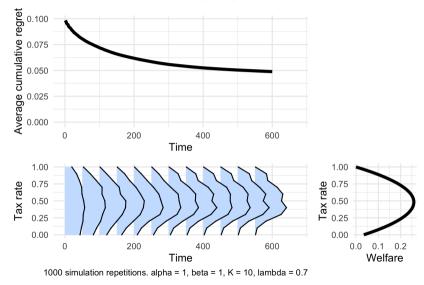
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Lower and upper bounds on regret

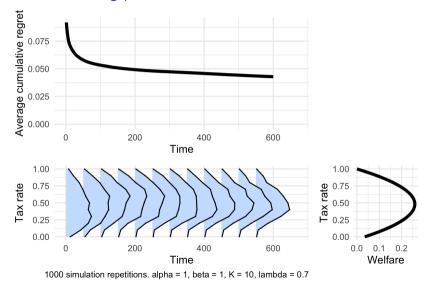
Comparison to related learning problems

#### Simulations

# Algorithm performance for $v \sim U[0,1]$



## Time-dependent tuning parameters



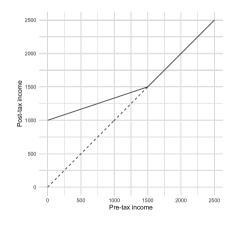
Setup

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Simulations

- Currently: Classic RCT evaluating a basic income, with the NGO "Mein Grundeinkommen" in Germany.
- In preparation: Adaptive follow-up.
  - Negative income tax: Basic income, taxed away until 0 transfer is reached.
  - Two policy parameters:
     Transfer size and tax rate.
     ⇒ Grid of possible combinations.



## Algorithm construction for the basic income experiment

- Structural model of labor supply:
  - Extensive and intensive margins.
  - Non-convex budget sets.
  - Measurement / optimization errors.
  - Observed and unobserved heterogeneity.
- Use MCMC (Metropolis-Hastings) to sample from the posterior for structural parameters.
- Map this into the posterior distribution of social welfare differences across policy choices.
- Assign policies using a version of tempered Thompson sampling.

# Thank you!

# Sketch of proof: upper bound on regret

- Discretize to balance the approximation error against the cost of having to learn G<sub>i</sub> on more points.
- $\widehat{\mathbb{G}}$  is an unbiased estimator for cumulative demand  $\mathbb{G}_i$ .  $\widehat{\mathbb{U}}$  is an unbiased estimator for cumulative discretized welfare.
- Consider  $W_i = \sum_k \exp(\eta \cdot \widehat{\mathbb{G}}_{ik})$ .
  - $E[\log W_T]$  is an bounded below by  $\eta$  times optimal constant policy welfare.
  - $E\left[\log\left(\frac{W_i}{W_{i-1}}\right)\right]$  is bounded above by a combination of expected  $\mathbb{U}_i$ , and a term based on the second moment of  $\widehat{\mathbb{U}}_i$ .
- Bounding this second moment, and optimizing tuning parameters, yields the bound on adversarial regret.