Adaptive maximization of social welfare

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Introduction

How should a policymaker act,

- who aims to maximize social welfare,
 - Weighted sum of utility.
 - ⇒ Tradeoff redistribution vs. cost of behavioral responses.
- and needs to learn agent responses to policy choices?
 - Adaptively updated policy choices
 - ⇒ Tradeoff exploration vs. exploitation.

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 Adaptively updated policy choices.
 - \Rightarrow Tradeoff exploration vs. exploitation.

Taxes and bandits

- Optimal tax theory
 - Mirrlees (1971); Saez (2001); Chetty (2009)
- Multi-armed bandits
 - Bubeck and Cesa-Bianchi (2012); Lattimore and Szepesvári (2020)
- This talk: Merging bandits and welfare economics.
 - Unobserved welfare, as in optimal taxation.
 - Unknown responses, as in multi-armed bandits.

Co-authors

- Nicolò Cesa-Bianchi and Roberto Colomboni, for the theory of adversarial and stochastic lower and upper bounds on regret.
- Frederik Schwertner, for implementation of an adaptive basic income experiment in Germany.

Setup

Lower and upper bounds on regret

Comparison to related learning problems

Simulations

Setup: Tax on a binary choice

Each time period $i = 1, 2, \dots, T$:

- One agent with willingness to pay $v_i \in [0, 1]$.
- Choices:
 - Tax rate $x_i \in [0, 1]$.
 - Individual response function: $G_i(x) = \mathbf{1}(x \le v_i)$
 - Binary agent decision $y_i = G_i(x_i)$.
- Observability:
 - After period i, we observe y_i .
 - We do *not* observe welfare $U_i(x_i)$.

Social welfare

Weighted sum of public revenue and private welfare:

$$U_i(x_i) = \underbrace{x_i \cdot \mathbf{1}(x_i \leq v_i)}_{ ext{Public revenue}} + \lambda \cdot \underbrace{\max(v_i - x_i, 0)}_{ ext{Private welfare}}.$$

We can rewrite private welfare as an integral (consumer surplus):

$$U_i(x) = \underbrace{x \cdot G_i(x)}_{ ext{Public revenue}} + \lambda \cdot \underbrace{\int_x^1 G_i(x') dx'}_{ ext{Private welfare}}.$$

Cumulative demand, welfare and regret

Cumulative demand:

$$\mathbb{G}_T(x) = \sum_{i < T} G_i(x).$$

• Cumulative welfare for a constant policy x:

$$\mathbb{U}_{T}(x) = \sum_{i < T} \mathbb{U}_{i}(x) = x \cdot \mathbb{G}_{T}(x) + \lambda \int_{x}^{1} \mathbb{G}_{T}(x') dx'.$$

• Cumulative welfare for the policies x_i actually chosen:

$$\mathbb{U}_T = \sum_{i \leq T} \mathbb{U}_i(x_i).$$

Adversarial regret:

$$\mathcal{R}_{T}(\lbrace v_{i}\rbrace_{i=1}^{T}) = \sup_{x} E\left[\mathbb{U}_{T}(x) - \mathbb{U}_{T} \middle| \lbrace v_{i}\rbrace_{i=1}^{T}\right].$$

The structure of observability

Choice x_i reveals $G_i(x_i)$. But

$$U_i(x) - U_i(x') = \left[x \cdot G_i(x) - x' \cdot G_i(x')\right] + \lambda \int_{x}^{x'} G_i(x'') dx''$$

depends on values of $G_i(x'')$ for $x'' \in [x, x']!$

Different from standard adaptive decision-making problems:

- Multi-armed bandits:
 Observe welfare for the choice made.
- Online learning:
 Observe welfare for all possible choices.
- Online convex optimization:
 Observe gradient of welfare for the choice made.

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Lower bound on regret

Theorem

There exists a constant C > 0 such that, for any randomized algorithm for the choice of x_1, x_2, \ldots and any time horizon $T \in \mathbb{N}$:

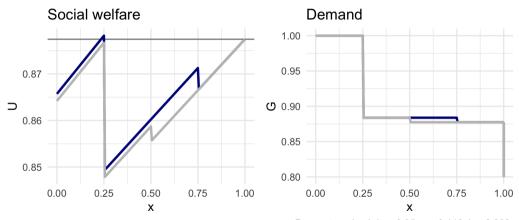
There exists a sequence (v_1, \ldots, v_T) for which

$$\mathcal{R}_T(\{v_i\}_{i=1}^T) \geq C \cdot T^{2/3}.$$

Sketch of proof: Lower bound on regret

- Stochastic regret ≤ adversarial regret. (Since average ≤ maximum.)
- Construct a distribution for v with 4 points of support, e.g. $(\frac{1}{4}, \frac{1}{2}, \frac{3}{4}, 1)$.
- Choose the probability of each of these points such that
 - 1. The two middle points are far from optimal.
 - Learning which of the two end points is optimal requires sampling from the middle. (Because of the integral term.)

Construction for the proof of the lower bound



Parameters: lambda = 0.95, a = 0.116, b = 0.003.

Tempered Exp3 for social welfare

Require: Tuning parameters K, γ and η .

- 1: Set $\tilde{x}_k = (k-1)/K$, initialize $\hat{\mathbb{G}}_k = 0$ for $k = 1, \dots, K+1$.
- 2: **for** individual i = 1, 2, ..., T **do**
- 3: **for** gridpoint k = 1, 2, ..., K + 1 **do**
- 4: Set

$$\widehat{\mathbb{U}}_{ik} = \widetilde{\mathbf{x}}_k \cdot \widehat{\mathbb{G}}_{ik} + \frac{\lambda}{K} \cdot \sum_{k' > k} \widehat{\mathbb{G}}_{ik'}, \quad p_{ik} = (1 - \gamma) \cdot \frac{\exp(\eta \cdot \widehat{\mathbb{U}}_{ik})}{\sum_{k'} \exp(\eta \cdot \widehat{\mathbb{U}}_{ik'})} + \frac{\gamma}{K + 1}.$$

- 5: end for
- 6: Choose k_i at random according to the probability distribution (p_1, \dots, p_{K+1}) .
- 7: Set $\mathbf{x}_i = \tilde{\mathbf{x}}_{k_i}$, and query \mathbf{y}_i accordingly.
- 8: Update

$$\hat{\mathbb{G}}_{k_i} = \hat{\mathbb{G}}_{k_i} + \frac{y_i}{p_{ik_i}}.$$

9: **end for**

Upper bound on regret

Theorem

Consider the algorithm "Tempered Exp3 for social welfare." There exists a constant C' and choices for K, γ, η such that, for any sequence (v_1, \ldots, v_T) ,

$$\mathcal{R}_T(\{v_i\}_{i=1}^T) \leq C' \cdot \log(T)^{1/3} \cdot T^{2/3}.$$

⇒ Same rate as the lower bound, up to the logarithmic term!

Sketch of proof

Setup

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Comparison to related learning problems

• Monopoly pricing:

Monopolist profits:

$$U_i^{MP}(x) = \underbrace{x \cdot G_i(x)}_{\text{Monopolist revenue}}$$

Easier – like a continuous multi-armed bandit.

Bilateral trade:

Buyer plus seller welfare:

$$U_i^{BT}(x) = G_i^b(x) \cdot \underbrace{\int_0^x G_i^s(x') dx'}_{ ext{Seller welfare}} + G_i^s(x) \cdot \underbrace{\int_x^1 G_i^b(x') dx'}_{ ext{Buyer welfare}}.$$

Harder – even gradients depend on global information.

Comparison of regret rates

Model	Continuous	Discrete
Monopoly price setting Optimal tax Bilateral trade	T ^{2/3} T ^{2/3} T	$T^{1/2}$ $T^{2/3}$ $T^{2/3}$

- Rates are up to logarithmic terms.
- They reflect the different information structures in the three problems.

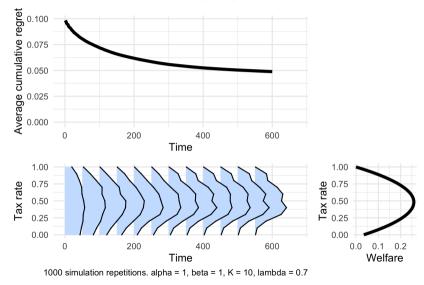
Setup

Lower and upper bounds on regret

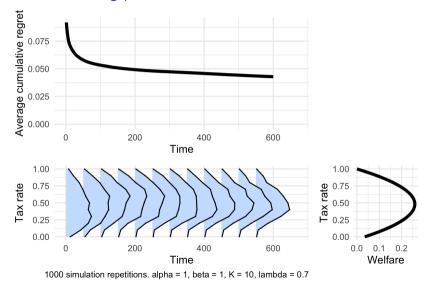
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Simulations

Algorithm performance for $v \sim U[0,1]$



Time-dependent tuning parameters



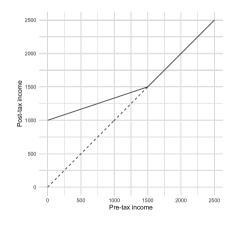
Setup

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Simulations

- Currently: Classic RCT evaluating a basic income, with the NGO "Mein Grundeinkommen" in Germany.
- In preparation: Adaptive follow-up.
 - Negative income tax: Basic income, taxed away until 0 transfer is reached.
 - Two policy parameters:
 Transfer size and tax rate.
 ⇒ Grid of possible combinations.



Algorithm construction for the basic income experiment

- Structural model of labor supply:
 - Extensive and intensive margins.
 - Non-convex budget sets.
 - Measurement / optimization errors.
 - Observed and unobserved heterogeneity.
- Use MCMC (Metropolis-Hastings) to sample from the posterior for structural parameters.
- Map this into the posterior distribution of social welfare differences across policy choices.
- Assign policies using a version of tempered Thompson sampling.

Thank you!

Sketch of proof: upper bound on regret

- Discretize to balance the approximation error against the cost of having to learn G_i on more points.
- $\widehat{\mathbb{G}}$ is an unbiased estimator for cumulative demand \mathbb{G}_i . $\widehat{\mathbb{U}}$ is an unbiased estimator for cumulative discretized welfare.
- Consider $W_i = \sum_k \exp(\eta \cdot \widehat{\mathbb{G}}_{ik})$.
 - $E[\log W_T]$ is an bounded below by η times optimal constant policy welfare.
 - $E\left[\log\left(\frac{W_i}{W_{i-1}}\right)\right]$ is bounded above by a combination of expected \mathbb{U}_i , and a term based on the second moment of $\widehat{\mathbb{U}}_i$.
- Bounding this second moment, and optimizing tuning parameters, yields the bound on adversarial regret.