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REFLECTION PRINCIPLES AND LARGE  
CARDINALS

Bakalářská práce

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### **Abstract**

Práce zkoumá vztah tzv. principů reflexe a velkých kardinálů. Lévy ukázal, že v ZFC platí tzv. věta o reflexi a dokonce, že věta o reflexi je ekvivalentní schématu nahrazení a axiomu nekonečna nad teorií ZFC bez axiomu nekonečna a schématu nahrazení. Tedy lze na větu o reflexi pohlížet jako na svého druhu axiom nekonečna. Práce zkoumá do jaké míry a jakým způsobem lze větu o reflexi zobecnit a jaký to má vliv na existenci tzv. velkých kardinálů. Práce definuje nedosažitelné, Mahlovy a nepopsatelné kardinály a ukáže, jak je lze zavést pomocí reflexe. Přirozenou limitou kardinálů získaných reflexí jsou kardinály nekonzistentní s L. Práce nabídne intuitivní zdůvodnění, proč tomu tak je.

### **Abstract**

This thesis aims to examine relations between the so called Reflection Principles and Large cardinals. Lévy has shown that the Reflection Theorem is a sound theorem of ZF and it is equivalent to the Replacement Scheme and the Axiom of Infinity. From this point of view, Reflection theorem can be seen a specific version of an Axiom of Infinity. This paper aims to examine the Reflection Principle and its generalisations with respect to the existence of Large Cardinals. This thesis will establish the Inaccessible, Mahlo and Indescribable cardinals and show how can those be defined via reflection. A natural limit of Large Cardinals obtained via reflection are cardinals inconsistent with L. This thesis will offer an intuitive explanation of why this holds.

# Contents

# 1 Introduction

## 1.1 Motivation and Origin

The Universe of sets cannot be uniquely characterised (i. e. distinguished from all its initial elements) by any internal structural property of the membership relation in it, which is expressible in any logic of finite or transfinite type, including infinitary logics of any cardinal order.

— Kurt Gödel [?]

## 1.2 Notation and Terminology

### 1.2.1 The Language of Set Theory

This text assumes the knowledge of basic terminology and some results from first-order predicate logic.<sup>1</sup>

We will now shortly review the basic notions that allow us to define *Zermelo–Fraenkel* set theory.

When we talk about *class*, we have the notion of definable class in mind. If  $\varphi(x, p_1, \dots, p_n)$  is a formula in the language of set theory, we call

$$A = \{x : \varphi(x)\} \tag{1.1}$$

a class of all sets satisfying  $\varphi(x)$  in a sense that

$$x \in A \leftrightarrow \varphi(x) \tag{1.2}$$

Given classes  $A, B$ , one can easily define the elementary set operations such as  $A \cap B, A \cup B, A \setminus C, \bigcup A$ , see the first part of [?] for details. The following axioms are the tools by which decide whether particular classes are in fact sets. A class that fails to be considered a set is called a *proper class*.

Speaking of formulas, we will often use syntax like ” $M$  is a limit ordinal”, it should be clear that this can be rewritten as a formula that was introduced earlier in the text.

### 1.2.2 The Axioms

**Definition 1.1** (*The Existence of a Set*)

$$\exists x(x = x) \tag{1.3}$$

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<sup>1</sup>todo odkaz na pripadny zdroj? svejdar? neco en?

The above axiom is usually not used because it can be deduced from the axiom of *Infinity* (see below), but since we will be using set theories that omit *Infinity*, this will be useful.

**Definition 1.2** (*Extensionality*)

$$\forall x, y (x = y \leftrightarrow \forall z (z \in x \leftrightarrow z \in y)) \quad (1.4)$$

**Definition 1.3** (*Specification*)

The following is a schema for every first-order formula  $\varphi(x, p_1, \dots, p_n)$  with no free variables other than  $x, p_1, \dots, p_n$ .

$$\forall x, p_1, \dots, p_n \exists y \forall z (z \in y \leftrightarrow z \in x \ \& \ \varphi(z, p_1, \dots, p_n)) \quad (1.5)$$

We will now provide two definitions that are not axioms, but will be helpful in establishing some of the other axioms in a more intuitive way.

**Definition 1.4** ( $x \subseteq y, x \subset y$ )

$$x \subseteq y \leftrightarrow (\forall z \in x) z \in y \quad (1.6)$$

$$x \subset y \leftrightarrow x \subseteq y \ \& \ x \neq y \quad (1.7)$$

**Definition 1.5** (*Empty Set*) Let  $\varphi = \neg(x = x)$ ,  $y$  is an arbitrary set, we there exists at least one set  $y$  from ?? or *Infinity*

$$\emptyset \stackrel{\text{def}}{=} \{x : x \in y \ \& \ \varphi(x)\} \quad (1.8)$$

We know that  $\emptyset$  is a set from specification and it is the same set for every  $y$  given from extensionality.

**Definition 1.6** (*Set Intersection*)

$$x \cap y = \{z : z \in x \ \& \ z \in y\} \quad (1.9)$$

**Definition 1.7** (*Set Union*)

$$x \cup y = \{z : z \in x \vee z \in y\} \quad (1.10)$$

Now we can introduce more axioms.

**Definition 1.8** (*Foundation*)

$$\forall x (x \neq \emptyset \rightarrow (\exists y \in x) (x \cap z = \emptyset)) \quad (1.11)$$

**Definition 1.9** (*Pairing*)

$$\forall x, y \exists z \forall q (q \in z \leftrightarrow q = x \vee q = y) \quad (1.12)$$

**Definition 1.10** (*Union*)

$$\forall x \exists y \forall z (z \in y \leftrightarrow \exists q (z \in q \ \& \ q \in x)) \quad (1.13)$$

**Definition 1.11** (*Powerset*)

$$\forall x \exists y \forall z (z \subseteq x \leftrightarrow z \in y) \quad (1.14)$$

**Definition 1.12** (*Infinity*)

$$\exists x (\emptyset \in x \ \& \ (\forall y \in x) (y \cup \{y\} \in x)) \quad (1.15)$$

Let us introduce a few more definitions that will make the two remaining axioms more comprehensible.

**Definition 1.13** (*Powerset function*)

Given a set  $x$ , the powerset of  $x$ , denoted  $\mathcal{P}(x)$  and satisfying ??, is defined as follows:

$$\mathcal{P}(x) \stackrel{\text{def}}{=} \{y : y \subseteq x\} \quad (1.16)$$

**Definition 1.14** (*Function*)

Given arbitrary first-order formula  $\varphi(x, y, p_1, \dots, p_n)$ , we say that  $\varphi$  is a function iff

$$\forall x, y, z, p_1, \dots, p_n (\varphi(x, y, p_1, \dots, p_n) \ \& \ \varphi(x, z, p_1, \dots, p_n) \rightarrow y = z) \quad (1.17)$$

When a  $\varphi(x, y)$  is a function, we also write the following:

$$f(x) = y \leftrightarrow \varphi(x, y) \quad (1.18)$$

<sup>2</sup> Alternatively,  $f = \{\langle x, y \rangle : \varphi(x, y)\}$  is a class.

**Definition 1.15** (*Dom(f)*)

Let  $f$  be a function. We read the following as "Dom( $f$ ) is the domain of  $f$ ".

$$\text{Dom}(f) \stackrel{\text{def}}{=} \{x : \exists y (f(x) = y)\} \quad (1.19)$$

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<sup>2</sup>This can also be done for  $\varphi$ s with more than two free variables by either setting  $f(x, p_1, \dots, p_n) = y \leftrightarrow \varphi(x, y, p_1, \dots, p_n)$  or saying that  $\varphi$  codes more functions, determined by the various parameters, so given  $t_1, \dots, t_n$ ,  $f(x) = y \leftrightarrow \varphi(x, y, t_1, \dots, t_n)$ .

We say " $f$  is a function on  $A$ ",  $A$  being a class, if  $A = \text{dom}(f)$ .

**Definition 1.16** (*Rng(f)*)

Let  $f$  be a function. We read the following as " $\text{Rng}(f)$  is the range of  $f$ ".

$$\text{Rng}(f) \stackrel{\text{def}}{=} \{x : \exists y(f(x) = y)\} \quad (1.20)$$

We say that  $f$  is a *function into*  $A$ ,  $A$  being a class, if  $\text{rng}(f) \subseteq A$ . We say that  $f$  is a *function onto*  $A$  if  $\text{rng}(f) = A$ , in other words,

$$(\forall y \in A)(\exists x \in \text{dom}(f))(f(x) = y) \quad (1.21)$$

We say a function  $f$  is a *one to one function*, iff

$$(\forall x_1, x_2 \in \text{dom}(f))(f(x_1) = f(x_2) \rightarrow x_1 = x_2) \quad (1.22)$$

$f$  is a bijection iff it is a one to one function that is onto.

Note that  $\text{Dom}(f)$  and  $\text{Rng}(f)$  are not definitions in a strict sense, they are in fact definition schemas that yield definitions for every function  $f$  given. Also note that they can be easily modified for  $\varphi$  instead of  $f$ , with the only difference being the fact that it is then defined only for those  $\varphi$ s that are functions, which must be taken into account. This is worth noting as we will sometimes interchange the notions of *function* and *formula*.

**Definition 1.17** (*Function Defined For All Ordinals*)

We say a function  $f$  is defined for all ordinals, this is sometimes written  $f : \text{Ord} \rightarrow A$  for any class  $A$ , if  $\text{Dom}(f) = \text{Ord}$ . Alternatively,

$$(\forall \alpha \in \text{Ord})(\exists y \in A)(f(\alpha) = y) \quad (1.23)$$

And now for the axioms.

**Definition 1.18** (*Replacement*)

The following is a schema for every first-order formula  $\varphi(x, p_1, \dots, p_n)$  with no free variables other than  $x, p_1, \dots, p_n$ .

$$"\varphi \text{ is a function}" \rightarrow \forall x \exists y \forall z (z \in y \leftrightarrow (\exists q \in x)(\varphi(x, y, p_1, \dots, p_n))) \quad (1.24)$$

**Definition 1.19** (*Choice*)

This is also a schema. For every  $A$ , a family of non-empty sets<sup>3</sup>, such that  $\emptyset \notin S$ , there is a function  $f$  such that for every  $x \in A$

$$f(x) \in x \quad (1.25)$$

Alternative:

$$\begin{aligned} \forall x \exists f ((f \text{ is a choice function with } \text{dom}(f) = x \setminus \{\emptyset\}) \\ \& \forall y ((y \in x \& y \neq \emptyset) \rightarrow f(y) \in y)) \end{aligned} \quad (1.26)$$

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<sup>3</sup>We say a class  $A$  is a "family of non-empty sets" iff there is  $B$  such that  $A \subseteq \mathcal{P}(B)$



We will refer the axioms by their name, written in italic type, e.g. *Foundation* refers to the Axiom of Foundation. Now we need to define some basic set theories to be used in the article.

**Definition 1.20 (S)**

We call **S** an axiomatic theory in the language  $\mathcal{L} = \{=, \in\}$  with exactly the following axioms:

- (i) Existence of a set (see ??)
- (ii) Extensionality (see ??)
- (iii) Specification (see ??)
- (iv) Foundation (see ??)
- (v) Pairing (see ??)
- (vi) Union (see ??)
- (vii) Powerset (see ??)

**Definition 1.21 (ZF)**

We call **ZF** an axiomatic theory in the language  $\mathcal{L} = \{=, \in\}$  that contains all the axioms of **S** in addition to the following

- (i) Replacement schema (see ??)
- (ii) Infinity (see ??)

Existence of a set is usually left out because it is implied by infinity.

**Definition 1.22 (ZFC)**

**ZFC** is an axiomatic theory in the language  $\mathcal{L} = \{=, \in\}$  that contains all the axioms of **ZF** plus Choice (??).

### 1.2.3 The Transitive Universe

**Definition 1.23 (Transitive Class)**

We say a class  $A$  is transitive iff

$$(\forall x \in A)(x \subseteq A) \tag{1.27}$$

**Definition 1.24 (Well Ordered Class)** A class  $A$  is said to be well ordered by  $\in$  iff the following hold:

- (i)  $(\forall x \in A)(x \not\in x)$  (*Antireflexivity*)
- (ii)  $(\forall x, y, z \in A)(x \in y \ \& \ y \in z \rightarrow x \in z)$  (*Transitivity*)
- (iii)  $(\forall x, y \in A)(x = y \vee x \in y \vee y \in x)$  (*Linearity*)
- (iv)  $(\forall x)(x \subseteq A \ \& \ x \neq \emptyset \rightarrow (\exists y \in x)(\forall z \in x)(z = y \vee z \in y))$  (*Existence of the least element*)

**Definition 1.25** (*Ordinal Number*)

A set  $x$  is said to be an ordinal number, also known as an ordinal, if it is transitive and well-ordered by  $\in$ .

For the sake of brevity, we usually just say " $x$  is an ordinal". Note that " $x$  is an ordinal" is a well-defined formula, since  $??$  is a formula and  $??$  is in fact a conjunction of four formulas. Ordinals will be usually denoted by lower case greek letters, starting from the beginning:  $\alpha, \beta, \gamma, \dots$ . Given two different ordinals  $\alpha, \beta$ , we will write  $\alpha < \beta$  for  $\alpha \in \beta$ , see [?] Lemma 2.11 for technical details.

**Definition 1.26** (*Non-Zero Ordinal*) We say an ordinal  $\alpha$  is non-zero iff  $\alpha \neq \emptyset$ .

**Definition 1.27** (*Successor Ordinal*)

Consider the following operation

$$\beta + 1 \stackrel{\text{def}}{=} \beta \cup \{\beta\} \quad (1.28)$$

An ordinal  $\alpha$  is called a successor ordinal iff there is an ordinal  $\beta$ , such that  $\alpha = \beta + 1$

**Definition 1.28** (*Limit Ordinal*)

A non-zero ordinal  $\alpha$  is called a limit ordinal iff it is not a successor ordinal.

**Definition 1.29** (*Ord*)

The class of all ordinal numbers, which we will denote  $\text{Ord}^4$  be the following class:

$$\text{Ord} \stackrel{\text{def}}{=} \{x : x \text{ is an ordinal}\} \quad (1.29)$$

The following construction will be often referred to as the *Von Neumann's Hierarchy*, sometimes also the *Von Neumann's Universe*.

**Definition 1.30** (*Von Neumann's Hierarchy*)

The Von Neumann's Hierarchy is a collection of sets indexed by elements of  $\text{Ord}$ , defined recursively in the following way:

(i)

$$V_0 = \emptyset \quad (1.30)$$

(ii)

$$V_{\alpha+1} = \mathcal{P}(V_\alpha) \text{ for any ordinal } \alpha \quad (1.31)$$

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<sup>4</sup>It is sometimes denoted  $\text{On}$ , but we will stick to the notation in [?]

(iii)

$$V_\lambda = \bigcup_{\beta < \lambda} V_\beta \text{ for a limit ordinal } \lambda \quad (1.32)$$

**Definition 1.31** (*Rank*)

Given a set  $x$ , we say that the rank of  $x$  (written as  $\text{rank}(x)$ ) is the least ordinal  $\alpha$  such that

$$x \in V_{\alpha+1} \quad (1.33)$$

Due to *Regularity*, every set has a rank.<sup>5</sup>

**Definition 1.32** ( $\omega$ )

$$\omega \stackrel{\text{def}}{=} \bigcap \{x : "x \text{ is a limit ordinal}"\} \quad (1.34)$$

$\omega$  is non-empty if *Infinity* or any equivalent holds.

**1.2.4 Cardinal Numbers****Definition 1.33** (*Cardinality*)

Given a set  $x$ , let the cardinality of  $x$ , written  $|x|$ , be defined as the smallest ordinal number such that there is a one to one mapping from  $x$  to  $\alpha$ .

For formal details as well as why every set can be well-ordered assuming *Choice*, see [?].

**Definition 1.34** (*Aleph function*)

Let  $\omega$  be the set defined by ???. We will recursively define the function  $\aleph$  for all ordinals.

- (i)  $\aleph_0 = \omega$
- (ii)  $\aleph_{\alpha+1}$  is the least cardinal larger than  $\aleph_\alpha$ <sup>6</sup>
- (iii)  $\aleph_\lambda = \bigcup_{\beta < \lambda} \aleph_\beta$  for a limit ordinal  $\lambda$

**Definition 1.35** (*Cardinal number*)

We say a set  $x$  is a cardinal number, usually shortened to a cardinal, if either  $x \in \omega$ , it is then called a finite cardinal, there is an ordinal  $\alpha$  such that  $\aleph_\alpha = x$ , then we call it an infinite cardinal

<sup>5</sup>See chapter 6 of [?] for details.

<sup>6</sup>"The least cardinal larger than  $\aleph_\alpha$ " is sometimes notated as  $\aleph_\alpha^+$

We say  $\kappa$  is an uncountable cardinal if it is an infinite ordinal and  $\aleph_0 > \kappa$ . Infinite cardinals will be notated by lowercase greek letters from the middle of the alphabet, e.g.  $\kappa, \mu, \nu, \dots$ <sup>7</sup>

**Definition 1.36** (*Cofinality of an Ordinal*)

Let  $\lambda$  be a limit ordinal. The cofinality of  $\lambda$ , written  $cf(\lambda)$ , is the smallest limit ordinal  $\alpha$ ,  $\alpha \leq \lambda$ , such that

$$(\forall x \in \lambda)(\exists y \in \alpha)(x < y) \quad (1.35)$$

**Definition 1.37** (*Regular Cardinal*)

We say a cardinal  $\kappa$  is regular iff  $cf(\kappa) = \kappa$

**Definition 1.38** (*Limit Cardinal*)

We say that a cardinal  $\kappa$  is a limit cardinal if

$$(\exists \alpha \in Ord)(\kappa = \aleph_\alpha) \quad (1.36)$$

**Definition 1.39** (*Strong Limit Cardinal*)

We say that an ordinal  $\kappa$  is a strong limit cardinal if it is a limit cardinal and

$$\forall \alpha (\alpha \in \kappa \rightarrow \mathcal{P}(\alpha) \in \kappa) \quad (1.37)$$

**Definition 1.40** (*Generalised Continuum Hypothesis*)

$$\aleph_{\alpha+1} = \mathcal{P}(\aleph_\alpha) \quad (1.38)$$

If *GCH* holds (for example in Gödel's *L*, see chapter 3), the notions of a limit cardinal and a strong limit cardinal are equivalent.

### 1.2.5 Relativisation and Absoluteness

**Definition 1.41** (*Relativization*)

Let  $M$  be a class,  $R \subseteq M \times M$  and let  $\varphi(p_1, \dots, p_n)$  be a first-order formula with no free variables besides  $p_1, \dots, p_n$ . The relativization of  $\varphi$  to  $M$  and  $R$  is the formula, written as  $\varphi^{M,R}(p_1, \dots, p_n)$ , defined in the following inductive manner:

- (i)  $(x \in y)^{M,R} \leftrightarrow R(x, y)$
- (ii)  $(x = y)^{M,R} \leftrightarrow x = y$

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<sup>7</sup> $\lambda$  is preferably used for limit ordinals, if it is ever used to denote an infinite cardinal, that should be contextually clear.

- (iii)  $(\neg\varphi)^{M,R} \leftrightarrow \neg\varphi^{M,R}$
- (iv)  $(\varphi \& \psi)^{M,R} \leftrightarrow \varphi^{M,R} \& \psi^{M,R}$
- (v)  $(\varphi \vee \psi)^{M,R} \leftrightarrow \varphi^{M,R} \vee \psi^{M,R}$
- (vi)  $(\varphi \rightarrow \psi)^{M,R} \leftrightarrow \varphi^{M,R} \rightarrow \psi^{M,R}$
- (vii)  $(\exists x\varphi(x))^{M,R} \leftrightarrow (\exists x \in M)\varphi^{M,R}(x)$
- (viii)  $(\forall x\varphi(x))^{M,R} \leftrightarrow (\forall x \in M)\varphi^{M,R}(x)$

When  $R = \in \cap (M \times M)$ , we usually write  $\varphi^M$  instead of  $\varphi^{M,R}$ . When we talk about  $\varphi^M(p_1, \dots, p_n)$ , it is understood that  $p_1, \dots, p_n \in M$ . We will also use  $M \models \varphi(p_1, \dots, p_n)$  and  $\varphi^M(p_1, \dots, p_n)$  interchangeably.

**Definition 1.42** (*Absoluteness*) Given a transitive class  $M$ , we say a formula  $\varphi$  is absolute in  $M$  if for all  $p_1, \dots, p_n \in M$

$$\varphi^M(p_1, \dots, p_n) \leftrightarrow \varphi(p_1, \dots, p_n) \quad (1.39)$$

**Definition 1.43** (*Hierarchy of First-Order Formulas*)

A first-order formula  $\varphi$  is  $\Delta_0$  iff it is logically equivalent to a first-order formula  $\varphi'$  satisfying any of the following:

- (i)  $\varphi'$  contains no quantifiers
- (ii)  $y$  is a set,  $\psi$  is a  $\Delta_0$  formula, and  $\varphi'$  is either  $(\exists x \in y)\psi(y)$  or  $(\forall x \in y)\psi(y)$ .
- (iii)  $\psi_1, \psi_2$  are  $\Delta_0$  formulas and  $\varphi'$  is any of the following:  $\psi_1 \vee \psi_2$ ,  $\psi_1 \& \psi_2$ ,  $\psi_1 \rightarrow \psi_2$ ,  $\neg\psi_2$ ,
- (I) If a formula is  $\Delta_0$  it is also  $\Sigma_0$  and  $\Pi_0$
- (II) A formula  $\varphi$  is  $\Pi_n + 1$  if it is logically equivalent to a formula  $\varphi'$  such that  $\varphi' = \forall x\psi$  where  $\psi$  is a  $\Sigma_n$ -formula for any  $n < \omega$ .
- (III) A formula  $\varphi$  is  $\Sigma_n + 1$  if it is logically equivalent to a formula  $\varphi'$  such that  $\varphi' = \exists x\psi$  where  $\psi$  is a  $\Pi_n$ -formula for any  $n < \omega$ .

Note that we can use the pairing function so that for  $\forall p_1, \dots, p_n \psi(p_1, \dots, p_n)$ , there a logically equivalent formula of the form  $\forall x\psi'(x)$ .

**Lemma 1.44** ( $\Delta_0$  absoluteness) Let  $\varphi$  be a  $\Delta_0$  formula, then  $\varphi$  is absolute in any transitive class  $M$ .

*Proof.* This will be proven by induction over the complexity of a given  $\Delta_0$  formula  $\varphi$ . Let  $M$  be an arbitrary transitive class. Suppose, that

Atomic formulas are always absolute by the definition of relativisation, see ???. Suppose that  $\Delta_0$  formulas  $\psi_1$  and  $\psi_2$  are absolute in  $M$ . Then

from relativization,  $(\psi_1 \& \psi_2)^M \leftrightarrow \psi_1^M \& \psi_2^M$ , which is, from the induction hypothesis, equivalent to  $\psi_1 \& \psi_2$ . The same holds for  $\vee, \rightarrow, \neg$ .

Suppose that a  $\Delta_0$  formula  $\psi$  is absolute in  $M$ . Let  $y$  be a set and let  $\varphi = (\exists x \in y)\psi(x)$ . From relativization,  $(\exists x\psi(x))^M \leftrightarrow (\exists x \in M)\psi^M(x)$ . Since the hypotheses makes it clear that  $\psi^M \leftrightarrow \psi$ , we get  $((\exists x \in y)\psi(x))^M \leftrightarrow (\exists x \in y \cap M)\psi(x)$ , which is the equivalent of  $\varphi^M \leftrightarrow \varphi$ . The same applies to  $\varphi = (\forall x \in y)\psi(x)$ .  $\square$

**Lemma 1.45** (*Downward Absoluteness*)

Let  $\varphi$  be a  $\Pi_1$  formula and  $M$  a transitive class. Then the following holds:

$$(\forall p_1, \dots, p_n \in M)(\varphi(p_1, \dots, p_n) \rightarrow \varphi(p_1, \dots, p_n)^M) \quad (1.40)$$

*Proof.* Since  $\varphi(p_1, \dots, p_n)$  is  $\Pi_1$ , there is a  $\Delta_0$  formula  $\psi(p_1, \dots, p_n, x)$  such that  $\varphi = \forall x\psi(p_1, \dots, p_n, x)$ . From relativization and lemma ??,  $\varphi^M(p_1, \dots, p_n) \leftrightarrow (\forall x \in M)\psi(p_1, \dots, p_n, x)$ .

Assume that for  $p_1, \dots, p_n \in M$  fixed, that  $\forall x\psi(p_1, \dots, p_n, x)$  holds, but  $(\forall x \in M)\psi(p_1, \dots, p_n, x)$  does not. Therefore  $\exists x\neg\psi(p_1, \dots, p_n, x)$ , which contradicts  $\forall x\psi(p_1, \dots, p_n, x)$ .  $\square$

**Lemma 1.46** (*Upward Absoluteness*)

Let  $\varphi$  be a  $\Sigma_1$  formula and  $M$  a transitive class. Then the following holds:

$$(\forall p_1, \dots, p_n \in M)(\varphi^M(p_1, \dots, p_n) \rightarrow \varphi(p_1, \dots, p_n)) \quad (1.41)$$

*Proof.* Since  $\varphi(p_1, \dots, p_n)$  is  $\Sigma_1$ , there is a  $\Delta_0$  formula  $\psi(p_1, \dots, p_n, x)$  such that  $\varphi = \exists x\psi(p_1, \dots, p_n, x)$ . From relativization and lemma ??,  $\varphi^M(p_1, \dots, p_n) \leftrightarrow (\exists x \in M)\psi(p_1, \dots, p_n, x)$ .

Assume that for  $p_1, \dots, p_n \in M$  fixed, that  $(\exists x \in M)\psi(p_1, \dots, p_n, x)$  holds, but  $\exists x\psi(p_1, \dots, p_n, x)$  does not. This is an obvious contradiction.  $\square$

### 1.2.6 More Functions

**Definition 1.47** (*Strictly Increasing Function*)

A function  $f : \text{Ord} \rightarrow \text{Ord}$  is said to be strictly increasing iff

$$\forall \alpha, \beta \in \text{Ord}(\alpha < \beta \rightarrow f(\alpha) < f(\beta)). \quad (1.42)$$

**Definition 1.48** (*Continuous Function*)

A function  $f : \text{Ord} \rightarrow \text{Ord}$  is said to be continuous iff

$$\alpha \text{ is limit} \rightarrow f(\alpha) = \bigcup_{\alpha < \lambda} f(\alpha). \quad (1.43)$$

**Definition 1.49** (*Normal Function*)

A function  $f : \text{Ord} \rightarrow \text{Ord}$  is said to be normal if it is strictly increasing and continuous.

**Definition 1.50** (*Fixed Point*)

We say  $x$  is a fixed point of a function  $f$  iff  $x = f(x)$ .

**Definition 1.51** (*Unbounded Class*)

We say a class  $A$  is unbounded if

$$\forall x(\exists y \in A)(x < y) \quad (1.44)$$

**Definition 1.52** (*Limit Point*)

Given a class  $x \subseteq \text{On}$ , we say that  $\alpha \neq \emptyset$  is a limit point of  $x$  iff

$$\alpha = \bigcup (x \cap \alpha) \quad (1.45)$$

**Definition 1.53** (*Closed Class*)

We say a class  $A \subseteq \text{Ord}$  is closed iff it contains all of its limit points.

**Definition 1.54** (*Club set*)

For a regular uncountable cardinal  $\kappa$ , a set  $x \subset \kappa$  is a closed unbounded subset, abbreviated as a club set, iff  $x$  is both closed and unbounded in  $\kappa$ .

**Definition 1.55** (*Stationary set*)

For a regular uncountable cardinal  $\kappa$ , we say a set  $A \subset \kappa$  is stationary in  $\kappa$  iff it intersects every club subset of  $\kappa$ .

**1.2.7 Structure, Substructure and Embedding**

Structures will be denoted  $\langle M, \in, R \rangle$  where  $M$  is a domain,  $\in$  stands for the standard membership relation, it is assumed to be restricted to the domain<sup>8</sup>,  $R \subseteq M$  is a relation on the domain. When  $R$  is not needed, we may as well only write  $M$  instead of  $\langle M, \in \rangle$ .

**Definition 1.56** (*Elementary Embedding*)

Given the structures  $\langle M_0, \in, R \rangle$ ,  $\langle M_1, \in, R \rangle$  and a one-to-one function  $j : M_0 \rightarrow M_1$ , we say  $j$  is an elementary embedding of  $M_0$  into  $M_1$ , we write  $j : M_0 \prec M_1$ , when the following holds for every formula  $\varphi(p_1, \dots, p_n)$  and every  $p_1, \dots, p_n \in M_0$ :

$$\langle M_0, \in, R \rangle \models \varphi(p_1, \dots, p_n) \leftrightarrow \langle M_1, \in, R \rangle \models \varphi(j(p_1), \dots, j(p_n)) \quad (1.46)$$

<sup>8</sup>To be totally correct, we should write  $\langle M, \in \cap M \times M, R \rangle$

**Definition 1.57** (*Elementary Substructure*)

Given the structures  $\langle M_0, \in, R \rangle$ ,  $\langle M_1, \in, R \rangle$  and a one-to-one function  $j : M_0 \rightarrow M_1$  such that  $j : M_0 \prec M_1$ , we say that  $M_0$  is an elementary substructure of  $M_1$ , denoted as  $M_0 \prec M_1$ , iff  $j$  is an identity on  $M_0$ . In other words

$$\langle M_0, \in, R \rangle \models \varphi(p_1, \dots, p_n) \leftrightarrow \langle M_1, \in, R \rangle \models \varphi(p_1, \dots, p_n) \quad (1.47)$$

for  $p_1, \dots, p_n \in M_0$

## 2 Lévy's First-Order Reflection

### 2.1 Lévy's Original Paper

This section is based on Lévy's paper *Axiom Schemata of Strong Infinity in Axiomatic Set Theory* [?] from 1960. It presents Lévy's general reflection principle and its equivalence to *Replacement* and *Infinity* under  $S^9$ .

When reading Lévy's article, one should bear in mind that while the author often speaks about a model of ZF, usually denoted  $u$ , this is equivalent to today's universal class  $V$ , so it doesn't necessarily mean that there is a set  $u$  that is a model of ZF<sup>10</sup>. We will review the notion of a standard complete model used by Lévy throughout the paper in a moment. The axioms used in what Lévy calls ZF are identical to those defined in (??), except for the *Axiom of Subsets*, which is just a different but equivalent formulation of *Specification*. Besides ZF and S, defined in (??) and (??) respectively, the set theories theories Z, and SF are used in the text. Z is ZF minus replacement, SF is ZF minus *Infinity*. Also note that the universal quantifier symbol does not appear,  $\forall x \varphi(x)$  would be written as  $(x)\varphi(x)$ . The symbol for negation is " $\sim$ ", implication is written as " $\supset$ " and equivalence is " $\equiv$ ", we will use standard notation with " $\neg$ ", " $\rightarrow$ " and " $\leftrightarrow$ " when presenting Lévy's results.

Next two definitions are not used in contemporary set theory, but they illustrate 1960's set theory mind-set and they are used heavily in Lévy's text, so we will include and explain them for clarity. Generally, in this chapter, Q stands for an arbitrary axiomatic set theory used for general definitions,  $u$  is usually a model of Q, counterpart of today's the universal class  $V$ .

This subsection uses ZF instead of the usual ZFC as the underlying theory.

**Definition 2.1** (*Standard Model of a Set Theory*)

Let Q be an arbitrary axiomatic set theory in first-order logic. We say the

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<sup>9</sup>see (??).

<sup>10</sup>This is indeed impossible to prove in ZF due to Gödel's Incompleteness.



the a class  $u$  is a standard model of  $\mathbf{Q}$  with respect to a membership relation  $E$ , written as  $Sm^{\mathbf{Q}}(u)$ , iff both of the following hold

- (i)  $(x, y) \in E \leftrightarrow y \in u \ \& \ x \in y$
- (ii)  $y \in u \ \& \ x \in y \rightarrow x \in u$

**Definition 2.2** (Standard Complete Model of a Set Theory)

Let  $\mathbf{Q}$  and  $E$  be like in (??). We say that  $u$  is a standard complete model of  $\mathbf{Q}$  with respect to a membership relation  $E$  iff both of the following hold

- (i)  $u$  is a transitive set with respect to  $\in$
- (ii)  $\forall E((x, y) \in E \leftrightarrow (y \in u \ \& \ x \in y) \ \& \ Sm^{\mathbf{Q}}(u, E))$

this is written as  $Scm^{\mathbf{Q}}(u)$ .

**Definition 2.3** (Inaccessible Cardinal With Respect to  $\mathbf{Q}$ )

Let  $\mathbf{Q}$  be an arbitrary axiomatic first-order set theory. We say that a cardinal  $\kappa$  is inaccessible with respect to  $\mathbf{Q}$ , we write  $In^{\mathbf{Q}}(\kappa)$ .

$$In^{\mathbf{Q}}(\kappa) \stackrel{\text{def}}{=} Scm^{\mathbf{Q}}(V_{\kappa}). \quad (2.48)$$

**Definition 2.4** (Inaccessible Cardinal With Respect to ZF)

When a cardinal  $\kappa$  is inaccessible with respect to  $\mathbf{ZF}$ , we only say that it is inaccessible. We write  $In(\kappa)$ .

$$In(\kappa) \stackrel{\text{def}}{=} In^{\mathbf{ZF}}(\kappa) \quad (2.49)$$

The above definition of inaccessibles is used because it doesn't require *Choice*.

For the definition of relativization, see (??). The syntax used by Lévy is  $Rel(u, \varphi)$ , we will use  $\varphi^u$ , which is more usual these days.

**Definition 2.5** ( $N$ )

The following is an axiom schema of complete reflection over  $\mathbf{ZF}$ , denoted as  $N$ :

$$\exists u (Scm^{\mathbf{ZF}}(u) \ \& \ \forall x_1, \dots, x_n (x_1, \dots, x_n \in u \rightarrow \varphi \leftrightarrow \varphi^u)) \quad (2.50)$$

where  $\varphi$  is a formula which contains no free variables except for  $x_1, \dots, x_n$ .

**Definition 2.6** ( $N_0$ )

With  $\mathbf{S}$  instead of  $\mathbf{ZF}$  we obtain what will now be called  $N_0$ :

$$\exists u (Scm^{\mathbf{S}}(u) \ \& \ \forall x_1, \dots, x_n (x_1, \dots, x_n \in u \rightarrow \varphi \leftrightarrow \varphi^u)) \quad (2.51)$$

where  $\varphi$  is a formula which contains no free variables except for  $x_1, \dots, x_n$ .

Now that we have established the basic terminology, we can review Lévy's proof that in a theory  $\mathbf{S}$ , which is defined in (??),  $N_0$  can be used to prove both *Replacement* and *Infinity*.

## 2.2 $S \vdash (N_0 \leftrightarrow \text{Replacement \& Infinity})$

Let  $S$  be a set theory as defined in (??). We will first prove a lemma to show what's mentioned as obvious in [?] and that is a fact, that any set  $u$  such that  $Scm^S(u)$  is a limit ordinal.

**Lemma 2.7** *The following holds for every  $u$ .*

$$"u \text{ is a limit ordinal}" \rightarrow Scm^S(u) \quad (2.52)$$

*Proof.* Let  $u$  be a standard complete model of  $S$ . We know that  $u$  is transitive from the definition of a standard complete model. To see that  $u$  is an ordinal, note that it is transitive and  $\emptyset \in u$  from *the existence of a set* (see (??)). To see that  $u$  is limit, consider that if  $u$  was a successor ordinal, there would be a set  $x \in u$  such that  $x \cup \{x\} = u$ , but then  $u \subset \mathcal{P}(x)$ , which contradicts the fact that  $(\forall x \in u)(\exists y \in u)(\mathcal{P}(x) = y)$  implied by *powerset* and it's not empty as stated earlier. We will now verify that all axioms of  $S$  are satisfied in a limit ordinal demoted  $u$ .

(i) *The existence of a set* comes from the fact that  $u$  is a non-empty set.

(ii) *Extensionality*:

(see (??))

$$\forall x, y (x = y \leftrightarrow \forall z (z \in x \leftrightarrow z \in y)) \quad (2.53)$$

Note that given arbitrary sets  $x, y$ , consider the formula  $\varphi$  defined as  $\varphi(x, y) = ((\forall z \in x)z \in y \ \& \ (\forall q \in y)q \in x) \leftrightarrow x = y$ . Because  $\varphi$  is  $\Delta_0$ ,  $\varphi \leftrightarrow \varphi^u$  by  $\Delta_0$ -absolutness lemma (??).

(iii) *Foundation*:

(see (??))

$$\forall x (x \neq \emptyset \rightarrow (\exists y \in x)(x \cap z = \emptyset)) \quad (2.54)$$

The formula  $\varphi(x) = x \neq \emptyset \rightarrow (\exists y \in x)(x \cap z = \emptyset)$  is  $\Delta_0$ , it is therefore absolute in  $u$  by lemma (??).

(iv) *Powerset*:

(see (??))

$$\forall x \exists y \forall z (z \subseteq x \leftrightarrow z \in y) \quad (2.55)$$

Given  $x \in u$ , we want to make sure that  $\mathcal{P}(x) \in u$ . Let  $\varphi$  denote the formula  $y \in \mathcal{P}(x) \leftrightarrow y \subset x$ . We know that  $y \subset x$  is  $\Delta_0$  according to definition (??). We also know that given  $x$ ,  $\varphi$  holds for every  $y$  due to the definition of  $\mathcal{P}(x)$ . That means that  $\varphi \leftrightarrow \varphi^u$  and therefore we can conclude that  $u \models \varphi$ .

(v) *Union*:

(see (??))

$$\forall x \exists y \forall z (z \in y \leftrightarrow \exists q (z \in q \ \& \ q \in x)) \quad (2.56)$$

Given any  $x \in u$ , we want verify that  $y = \bigcup x$  is also in  $u$ . Note that  $y = \bigcup x$  is also  $\Delta_0$ .

$$y = \bigcup x \leftrightarrow (\forall z \in y)(\exists q \in x)z \in q \ \& \ (\forall z \in x)(\forall q \in z)q \in y \quad (2.57)$$

So by lemma (??)

$$y = \bigcup x \leftrightarrow (y = \bigcup x)^u \quad (2.58)$$

(vi) *Pairing*:  
(see (??))

$$\forall x, y \exists z \forall q (q \in z \leftrightarrow q = x \vee q = y) \quad (2.59)$$

Given two sets  $x, y \in u$ , we want to show that  $z$ , defined as  $z = \{x, y\}$ , is also an element of  $u$ .

$$z = \{x, y\} \leftrightarrow x \in z \ \& \ y \in z \ \& \ (\forall q \in z)(q = x \vee q = y) \quad (2.60)$$

So  $(z = \{x, y\})$  is  $\Delta_0$ , and thus by lemma (??) it holds that

$$z = \{x, y\} \leftrightarrow (z = \{x, y\})^u \quad (2.61)$$

(vii) *Specification*:

Given a first-order formula  $\varphi$ , we want to show the following

$$u \models \forall x \forall p_1, \dots, p_n, \exists y \forall z (z \in y \leftrightarrow z \in x \ \& \ \varphi(z, p_1, \dots, p_n)) \quad (2.62)$$

Given any  $x$  along with parameters  $p_1, \dots, p_n$ , we set  $y = \{z \in x : \varphi^u(z, p_1, \dots, p_n)\}$ . From transitivity of  $u$  and the fact that  $y \subset x$  and  $x \in u$ , we can conclude that  $y \in u$ , so  $u \models \forall z (z \in y \leftrightarrow z \in x \ \& \ \varphi(z, p_1, \dots, p_n))$ .

□

Let  $N_0$  be defined as in (??), for *Infinity* see (??).

**Theorem 2.8** *In  $S$ , the schema  $N_0$  implies Infinity.*

*Proof.* Lévy skips this proof because it seems too obvious to him, but we will do it here for plasticity. For an arbitrary  $\varphi$ ,  $N_0$  gives us  $\exists u Scm^S(u)$ , but from lemma (??), we know that this  $u$  is a limit ordinal. This  $u$  already satisfies *Infinity*. □

Let  $N_0$  be defined as in (??), for *Replacement* see (??),  $S$  is again the set theory defined in (??).

**Theorem 2.9** *In  $S$ , the schema  $N_0$  implies Replacement.*

*Proof.* Let  $\varphi(x, y, p_1, \dots, p_n)$  be a formula with no free variables except  $x, y, p_1, \dots, p_n$  for an arbitrary natural number  $n$ .

$$\begin{aligned} \chi = & \forall x, y, z (\varphi(x, y, p_1, \dots, p_n) \ \& \ \varphi(x, z, p_1, \dots, p_n) \rightarrow y = z) \\ & \rightarrow \forall x \exists y \forall z (z \in y \leftrightarrow \exists q (q \in x \ \& \ \varphi(q, z, p_1, \dots, p_n))) \end{aligned} \quad (2.63)$$

Let  $\chi$  be an instance of *Replacement* schema for given  $\varphi$ . Let the following formulas be instances of the  $N_0$  schema for formulas  $\varphi$ ,  $\exists y \varphi$ ,  $\chi$  and  $\forall x, p_1, \dots, p_n \chi$  respectively:

We can deduce the following from  $N_0$ :

- (i)  $x, y, p_1, \dots, p_n \in u \rightarrow (\varphi \leftrightarrow \varphi^u)$
- (ii)  $x, p_1, \dots, p_n \in u \rightarrow (\exists y \varphi \leftrightarrow (\exists y \varphi)^u)$
- (iii)  $x, p_1, \dots, p_n \in u \rightarrow (\chi \leftrightarrow \chi^u)$
- (iv)  $\forall x, p_1, \dots, p_n (\chi \leftrightarrow (\forall x, p_1, \dots, p_n \chi)^u)$

From relativization, we also know that  $(\exists y \varphi)^u$  is equivalent to  $(\exists y \in u) \varphi^u$ . Therefore (ii) is equivalent to

$$x, p_1, \dots, p_n \in u \rightarrow (\exists y \in u) \varphi^u. \quad (2.64)$$

If  $\varphi$  is a function<sup>11</sup>, then for every  $x \in u$ , which is also  $x \subset u$  by the transitivity of  $\text{Scm}^S(u)$ , it maps elements of  $x$  onto  $u$ . From the axiom scheme of comprehension<sup>12</sup>, we can find  $y$ , a set of all images of elements of  $x$ . That gives us  $x, p_1, \dots, p_n \in u \rightarrow \chi$ . By (iii) we get  $x, p_1, \dots, p_n \in u \rightarrow \chi^u$ , the universal closure of this formula is  $(\forall x, p_1, \dots, p_n \chi)^u$ , which together with (iv) yields  $\forall x, p_1, \dots, p_n \chi$ . Via universal instantiation, we end up with  $\chi$ . We have inferred replacement for a given arbitrary formula.  $\square$

What we have just proven is just a single theorem from the above mentioned article by Lévy, we will introduce other interesting propositions, mostly related to the existence of large cardinals, later in their appropriate context in chapter 3.

## 2.3 Contemporary Restatement

We will now a theorem that is referred to as Lévy's Reflection in contemporary set theory. The only difference is that while Lévy reflects  $\varphi$  from  $V$  to a set  $u$  which is a *standard complete model of  $\mathbf{S}$* , we say that there is a  $V_\alpha$  for a limit  $\alpha$  that reflects  $\varphi$ . Those two conditions are equivalent due to lemma (??).

<sup>11</sup>See definition (??)

<sup>12</sup>Lévy uses its equivalent, axiom of subsets

**Lemma 2.10** *Let  $\varphi_1, \dots, \varphi_n$  be first-order formulas in the language of set theory, all with  $m$  free variables<sup>13</sup>.*

- (i) *For each set  $M_0$  there is such set  $M$  that  $M_0 \subset M$  and the following holds for every  $i$ ,  $1 \leq i \leq n$ :*

$$\exists x \varphi_i(p_1, \dots, p_{m-1}, x) \rightarrow (\exists x \in M) \varphi_i(p_1, \dots, p_{m-1}, x) \quad (2.65)$$

*for every  $p_1, \dots, p_{m-1} \in M$ .*

- (ii) *Furthermore, there is an ordinal  $\alpha$  such that  $M_0 \subset V_\alpha$  and the following holds for each  $i$ ,  $1 \leq i \leq n$ :*

$$\exists x \varphi_i(p_1, \dots, p_{m-1}, x) \rightarrow (\exists x \in V_\alpha) \varphi_i(p_1, \dots, p_{m-1}, x) \quad (2.66)$$

*for every  $p_1, \dots, p_{m-1} \in M$ .*

- (iii) *Assuming Choice, there is  $M$ ,  $M_0 \subset M$  such that (??) holds for every  $M$ ,  $i \leq n$  and  $|M| \leq |M_0| \cdot \aleph_0$ .*

*Proof.* We will simultaneously prove statements (i) and (ii), denoting  $M^T$  the transitive set required by part (ii). Steps in the construction of  $M^T$  that are not explicitly included are equivalent to steps for  $M$ .

Let us first define an operation  $H_i(p_1, \dots, p_{m-1})$  that yields the set of  $x$ 's with minimal rank<sup>14</sup> satisfying  $\varphi_i(p_1, \dots, p_{m-1}, x)$  for  $p_1, \dots, p_{m-1}$  and for every  $i$ ,  $1 \leq i \leq n$ .

$$H_i(p_1, \dots, p_n) = \{x \in C_i : (\forall z \in C)(\text{rank}(x) \leq \text{rank}(z))\} \quad (2.67)$$

for each  $1 \leq i \leq n$ , where

$$C_i = \{x : \varphi_i(p_1, \dots, p_{m-1}, x)\} \text{ for } 1 \leq i \leq n \quad (2.68)$$

Next, let's construct  $M$  from given  $M_0$  by induction.

$$M_{i+1} = M_i \cup \bigcup_{j=0}^n \{H_j(p_1, \dots, p_{m-1}) : p_1, \dots, p_{m-1} \in M_i\} \quad (2.69)$$

<sup>13</sup>For formulas with a different number of free variables, take for  $m$  the highest number of parameters among those formulas. Add spare parameters to every formula that has less than  $m$  parameters in a way that preserves the last parameter, which we will denote  $x$ . E.g. let  $\varphi'_i$  be the a formula with  $k$  parameters,  $k < m$ . Let us set  $\varphi_i(p_1, \dots, p_{m-1}, x) \stackrel{\text{def}}{=} \varphi'_i(p_1, \dots, p_{k-1}, x)$ , notice that the parameters  $p_k, \dots, p_{m-1}$  are not used.

<sup>14</sup>Rank is defined in (??)

In other words, in each step we include into the construction the elements satisfying  $\varphi(p_1, \dots, p_{m-1}, x)$  for  $p_1, \dots, p_{m-1}$  from the previous step. For statement (ii), this is the only part that differs from (i). To end up with a transitive  $M$ , we need to extend every step to its transitive closure transitive closure of  $M_{i+1}$  from (i). In other words, let  $\gamma$  be the smallest ordinal such that

$$(M_i^T \cup \bigcup_{j=0}^n \{ \bigcup \{ H_j(p_1, \dots, p_{m-1}) : p_1, \dots, p_{m-1} \in M_i \} \}) \subset V_\gamma \quad (2.70)$$

Then the incremental step is

$$M_{i+1}^T = V_\gamma \quad (2.71)$$

and the final  $M$  is obtained by joining the previous steps.

$$M = \bigcup_{i=0}^{\infty} M_i, \quad M^T = \bigcup_{i=0}^{\infty} M_i^T = V_\alpha \quad (2.72)$$

We have yet to finish part (iii). Let's try to construct a set  $M'$  that satisfies the same conditions like  $M$  but is kept as small as possible. Assuming the Axiom of Choice, we can modify the process so that the cardinality of  $M'$  is at most  $|M_0| \cdot \aleph_0$ . Note that the size of  $M$  in the previous construction is determined by the size of  $M_0$  and, most importantly, by the size of  $H_i(p_1, \dots, p_{m-1})$  for every  $i$ ,  $1 \leq i \leq n$  in individual iterations of the construction. Since (i) only ensures the existence of an  $x$  that satisfies  $\varphi_i(p_1, \dots, p_{m-1}, x)$  for any  $i$ ,  $1 \leq i \leq n$ , we only need to add one  $x$  for every set of parameters but  $H_i(u_1, \dots, u_{m-1})$  can be arbitrarily large. Let  $F$  be a choice function on  $\mathcal{P}(M')$ . Also let  $h_i(p_1, \dots, p_{m-1}) = F(H_i(p_1, \dots, p_{m-1}))$  for  $i$ , where  $1 \leq i \leq n$ , which means that  $h$  is a function that outputs an  $x$  that satisfies  $\varphi_i(p_1, \dots, p_{m-1}, x)$  for  $i$  such that  $1 \leq i \leq n$  and has minimal rank among all such sets. The induction step needs to be redefined to

$$M'_{i+1} = M'_i \cup \bigcup_{j=0}^n \{ H_j(p_1, \dots, p_{m-1}) : p_1, \dots, p_{m-1} \in M'_i \} \quad (2.73)$$

This way, the amount of elements added to  $M'_{i+1}$  in each step of the construction is the same as the amount of  $m$ -tuples of parameters that yielded elements not included in  $M'_i$ . It is easy to see that if  $M_0$  is finite,  $M'$  is countable because it was constructed as a countable union of at most count-

able sets. If  $M_0$  is countable or larger, the cardinality of  $M'$  is equal to the cardinality of  $M_0$ .<sup>15</sup> Therefore  $|M'| \leq |M_0| \cdot \aleph_0$   $\square$

**Theorem 2.11** (*Lévy's first-order reflection theorem*)

Let  $\varphi(p_1, \dots, p_n)$  be a first-order formula.

- (i) For every set  $M_0$  there exists  $M$  such that  $M_0 \subset M$  and the following holds:

$$\varphi^M(p_1, \dots, p_n) \leftrightarrow \varphi(p_1, \dots, p_n) \quad (2.74)$$

for every  $p_1, \dots, p_n \in M$ .

- (ii) For every set  $M_0$  there is a transitive set  $M$ ,  $M_0 \subset M$  such that the following holds:

$$\varphi^M(p_1, \dots, p_n) \leftrightarrow \varphi(p_1, \dots, p_n) \quad (2.75)$$

for every  $p_1, \dots, p_n \in M$ .

- (iii) For every set  $M_0$  there is  $\alpha$  such that  $M_0 \subset V_\alpha$  and the following holds:

$$\varphi^{V_\alpha}(p_1, \dots, p_n) \leftrightarrow \varphi(p_1, \dots, p_n) \quad (2.76)$$

for every  $p_1, \dots, p_n \in M$ .

- (iv) Assuming Choice, for every set  $M_0$  there is  $M$  such that  $M_0 \subset M$  and  $|M| \leq |M_0| \cdot \aleph_0$  and the following holds:

$$\varphi^M(p_1, \dots, p_n) \leftrightarrow \varphi(p_1, \dots, p_n) \quad (2.77)$$

for every  $p_1, \dots, p_n \in M$ .

*Proof.* Before we start, note that the following holds for any set  $M$  if  $\varphi$  is an atomic formula, as a direct consequence of relativisation to  $M$ ,  $\in$ <sup>16</sup>.

$$\varphi \leftrightarrow \varphi^M \quad (2.78)$$

Let's now prove (i) for given  $\varphi$  via induction by complexity. We can safely assume that  $\varphi$  contains no quantifiers besides " $\exists$ " and no logical connectives other than " $\neg$ " and " $\&$ ". Let  $\varphi_1, \dots, \varphi_n$  be all subformulas of  $\varphi$ . Then there is a set  $M$ , obtained by the means of lemma ??, for all of the formulas  $\varphi_1, \dots, \varphi_n$ .

<sup>15</sup>It can not be smaller because  $|M'_{i+1}| \geq |M'_i|$  for every  $i$ . It may not be significantly larger because the maximum of elements added is the number of  $n$ -tuples in  $M'_i$ , which is of the same cardinality is  $M'_i$ .

<sup>16</sup>See ??. Also note that this only holds for relativization to  $M, \in$ , not  $M, E$  for arbitrary  $E$ .

We know that  $\psi \leftrightarrow \psi^M$  for atomic  $\psi$ , we need to verify that it won't fail in the inductive step. Let us consider  $\psi = \neg\psi'$  along with the definition of relativization for those formulas in ??.

$$(\neg\psi')^M \leftrightarrow \neg(\psi'^M) \quad (2.79)$$

Because the induction hypothesis says that ?? holds for every subformula of  $\psi$ , we can assume that  $\psi'^M \leftrightarrow \psi'$ , therefore the following holds:

$$(\neg\psi')^M \leftrightarrow \neg(\psi'^M) \leftrightarrow \neg\psi' \quad (2.80)$$

The same holds for  $\psi = \psi_1 \ \& \ \psi_2$ . From the induction hypothesis, we know that  $\psi_1^M \leftrightarrow \psi_1$  and  $\psi_2^M \leftrightarrow \psi_2$ , which together with relativization for formulas in the form of  $\psi_1 \ \& \ \psi_2$  gives us

$$(\psi_1 \ \& \ \psi_2)^M \leftrightarrow \psi_1^M \ \& \ \psi_2^M \leftrightarrow \psi_1 \ \& \ \psi_2 \quad (2.81)$$

Let's now examine the case when, from the induction hypothesis,  $M$  reflects  $\psi'(p_1, \dots, p_n, x)$  and we are interested in  $\psi = \exists x\psi'(p_1, \dots, p_n, x)$ . The induction hypothesis tells us that

$$\varphi'^M(p_1, \dots, p_n, x) \leftrightarrow \psi'(p_1, \dots, p_n, x) \quad (2.82)$$

so, together with above lemma ??, the following holds:

$$\begin{aligned} & \psi(p_1, \dots, p_n, x) \\ & \leftrightarrow \exists x\psi'(p_1, \dots, p_n, x) \\ & \leftrightarrow (\exists x \in M)\psi'(p_1, \dots, p_n, x) \\ & \leftrightarrow (\exists x \in M)\psi'^M(p_1, \dots, p_n, x) \\ & \leftrightarrow (\exists x\psi'(p_1, \dots, p_n, x))^M \\ & \leftrightarrow \psi^M(p_1, \dots, p_n, x) \end{aligned} \quad (2.83)$$

Which is what we have needed to prove.

So far we have proven part (i) of this theorem for one formula  $\varphi$ , we only need to verify that the same holds for any finite number of formulas. This has in fact been already done since lemma ?? gives us  $M$  for any (finite) amount of formulas, we can find a set  $M$  for the union of all of their subformulas. We can then use the induction above to verify that  $M$  reflects each of the formulas individually iff it reflects all of its subformulas.



Since  $V_\alpha$  is a transitive set, by proving (iii) we also satisfy (ii). To do so, we only need to look at part (ii) of lemma ???. All of the above proof also holds for  $M = V_\alpha$ .

To finish part (iv), we take  $M$  of size  $\leq |M_0| \cdot \aleph_0$ , which exists due to part (iii) of lemma ??, the rest being identical.  $\square$

Let  $S$  be a set theory defined in ??, for ZFC see ??.

Let *Infinity* and *Replacement* be as defined in ?? and ?? respectively.

**Definition 2.12** (*First-Order Reflection Schema*)

For each set  $M_0$  there is such set  $M$  that  $M_0 \subset M$  and the following holds for every  $i$ ,  $1 \leq i \leq n$ :

$$\exists x \varphi_i(p_1, \dots, p_{m-1}, x) \rightarrow (\exists x \in M) \varphi_i(p_1, \dots, p_{m-1}, x) \quad (2.84)$$

for every  $p_1, \dots, p_{m-1} \in M$ .

**Theorem 2.13** *Reflection<sub>1</sub> is equivalent to Infinity & Replacement under S.*

*Proof.* Since ?? already gives us one side of the implication, we are only interested in showing the converse which we shall do in two parts:

*Reflection<sub>1</sub> → Infinity* From *Reflection<sub>1</sub>*, we know that for any first-order formula  $\varphi$  and a set  $M_0$ , there is a  $M$  such that  $M_0 \subseteq M$  and  $\varphi^M \leftrightarrow \varphi$ . Let's pick *Powerset* for  $\varphi$ , then by *Reflection<sub>1</sub>* there is a set that satisfies *Powerset*, ergo there is a strong limit cardinal, which in turn satisfies *Infinity*.

*Reflection → Replacement*

Given a formula  $\varphi(x, y, p_1, \dots, p_n)$ , we can suppose that it is reflected in any  $M$ <sup>17</sup> What we want to obtain is the following:

$$\begin{aligned} \forall x, y, z (\varphi(x, y, p_1, \dots, p_n) \ \& \ \varphi(x, z, p_1, \dots, p_n) \rightarrow y = z) \rightarrow \\ \rightarrow \forall X \exists Y \forall y (y \in Y \leftrightarrow \exists x (\varphi(x, y, p_1, \dots, p_n) \ \& \ x \in X)) \end{aligned} \quad (2.85)$$

We do also know that  $x, y \in M$ , in other words for every  $X$ ,  $Y = \{y \mid \varphi(x, y, p_1, \dots, p_n)\}$  and we know that  $X \subset M$  and  $Y \subset M$ , which, together with the specification schema implies that  $Y$ , the image of  $X$  over  $\varphi$ , is a set.  $\square$

We have shown that *Reflection* for first-order formulas, *Reflection<sub>1</sub>* is a theorem of ZF, which means that it won't yield us any large cardinals.

<sup>17</sup>Which means that for  $x, y, p_1, \dots, p_n \in M$ ,  $\varphi^M(x, y, p_1, \dots, p_n) \leftrightarrow \varphi(x, y, p_1, \dots, p_n)$ .

We have also shown that it can be used instead of the *Infinity* and *Replacement* scheme, but  $\mathbf{ZF} + \text{Reflection}_1$  is a conservative extension of  $\mathbf{ZF}$ . Besides being a starting point for more general and powerful statements, it can be used to show that  $\mathbf{ZF}$  is not finitely axiomatizable. That follows from the fact that *Reflection* gives a model to any finite number of (consistent) formulas. So if  $\varphi_1, \dots, \varphi_n$  for any finite  $n$  would be the axioms of  $\mathbf{ZF}$ , *Reflection* would always contain a model of itself, which would in turn contradict the Second Gödel's Theorem<sup>18</sup>. Notice that, in a way, reflection is complementary to compactness. Compactness argues that given a set of sentences, if every finite subset yields a model, so does the whole set. Reflection, on the other hand, says that while the whole set has no model in the underlying theory, every finite subset does have one.

Also, notice how reflection can be used in ways similar to upward Löwenheim–Skolem theorem. Since Reflection extends any set  $M_0$  into a model of given formulas  $\varphi_1, \dots, \varphi_n$ , we can choose the lower bound of the size of  $M$  by appropriately choosing  $M_0$ .

In the next section, we will try to generalize *Reflection* in a way that transcends  $\mathbf{ZF}$  and finally yields some large cardinals.

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<sup>18</sup>See chapter ?? for further details.

## **3 Conclusion**

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