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4 REFLECTION PRINCIPLES AND LARGE  
5 CARDINALS  
6 Bakalářská práce

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8 2015

<sup>10</sup> Prohlašuj, že jsem bakalářskou práci vypracoval samostatně a že jsem uvedl  
<sup>11</sup> všechny použité prameny a literaturu.

<sup>12</sup> V Praze 14. dubna 2015

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## Abstract

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Práce zkoumá vztah tzv. principů reflexe a velkých kardinálů. Lévy ukázal, že v ZFC platí tzv. věta o reflexi a dokonce, že věta o reflexi je ekvivalentní schématu nahrazení a axiomu nekonečna nad teorií ZFC bez axiomu nekonečna a schématu nahrazení. Tedy lze na větu o reflexi pohlížet jako na svého druhu axiom nekonečna. Práce zkoumá do jaké míry a jakým způsobem lze větu o reflexi zobecnit a jaký to má vliv na existenci tzv. velkých kardinálů. Práce definuje nedosažitelné, Mahlovy a nepopsatelné kardinály a ukáže, jak je lze zavést pomocí reflexe. Přirozenou limitou kardinálů získaných reflexí jsou kardinály nekonzistentní s L. Práce nabídne intuitivní zdůvodnění, proč tomu tak je.

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## Abstract

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This thesis aims to examine relations between so called "Reflection Principles" and Large cardinals. Lévy has shown that Reflection Theorem is a sound theorem of ZFC and it is equivalent to Replacement Scheme and the Axiom of Infinity. From this point of view, Reflection theorem can be seen a specific version of an Axiom of Infinity. This paper aims to examine the Reflection Principle and it's generalisations with respect to existence of Large Cardinals. This thesis will establish Inaccessible, Mahlo and Indescribable cardinals and their definition via reflection. A natural limit of Large Cardinals obtained via reflection are cardinals inconsistent with L. The thesis will offer an intuitive explanation of why this is the case.

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# 1 Introduction

## 1.1 Motivation and Origin

The Universe of sets cannot be uniquely characterized (i. e. distinguished from all its initial elements) by any internal structural property of the membership relation in it, which is expressible in any logic of finite or transfinite type, including infinitary logics of any cardinal order.

— Kurt Gödel [4]

To understand why we need reflection in the first place, let's think about infinity for a moment. In the intuitive sense, infinity is an upper limit of all numbers. But for centuries, this was merely a philosophical concept, closely bound to religious and metaphysical way of thinking, considered separate from numbers used for calculations or geometry. It was a rather vague concept. In ancient Greece, Aristotle's response to famous Zeno's paradoxes introduced the distinction between actual and potential infinity. He argued, that potential infinity is (in today's words) well defined, as opposed to actual infinity, which remained a vague incoherent concept. He didn't think it's possible for infinity to inhabit a bounded place in space or time, rejecting Zeno's thought experiments as a whole. Aristotle's thoughts shaped western thinking partly due to Aquinas, who himself believed actual infinity to be more of a metaphysical concept for describing God than a mathematical property attributed to any other entity. In his *Summa Theologica*<sup>1</sup> he argues:

A geometrician does not need to assume a line actually infinite, but takes some actually finite line, from which he subtracts whatever he finds necessary; which line he calls infinite.

Less than hundred years later, Gregory of Rimini wrote

If God can endlessly add a cubic foot to a stone—which He can—then He can create an infinitely big stone. For He need only add one cubic foot at some time, another half an hour later, another a quarter of an hour later than that, and so on ad infinitum. He would then have before Him an infinite stone at the end of the hour.

Which is basically a Zeno's Paradox made plausible with God being the actor. In contrast to Aquinas' position, Gregory of Rimini theoretically constructs an object with actual infinite magnitude that is essentially different from

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<sup>1</sup>Part I, Question 7, Article 3, Reply to Objection 1

94 God. Even later, in the 17th century, pushing the property of infinitness  
95 from the Creator to his creation, Nature, Leibniz wrote to Foucher in 1662:

96 I am so in favor of the actual infinite that instead of admitting  
97 that Nature abhors it, as is commonly said, I hold that Nature  
98 makes frequent use of it everywhere, in order to show more ef-  
99 fectively the perfections of its Author. Thus I believe that there  
100 is no part of matter which is not, I do not say divisible, but ac-  
101 tually divided; and consequently the least particle ought to be  
102 considered as a world full of an infinity of different creatures.

103 But even though he used potential infinity in what would become foundations  
104 of modern Calculus and argued for actual infinity in Nature, Leibniz refused  
105 the existence of an infinite, thinking that Galileo's Paradoxon<sup>2</sup> is in fact  
106 a contradiction. The so called Galileo's Paradoxon is an observation Galileo  
107 Galilei made in his final book "Discourses and Mathematical Demonstrations  
108 Relating to Two New Sciences". He states that if all numbers are either  
109 squares and non-squares, there seem to be less squares than there is all  
110 numbers. On the other hand, every number can be squared and every square  
111 has it's square root. Therefore, there seem to be as many squares as there  
112 are all numbers. Galileo concludes, that the idea of comparing sizes makes  
113 sense only in the finite realm.

114 Salviati: So far as I see we can only infer that the totality of all  
115 numbers is infinite, that the number of squares is infinite, and  
116 that the number of their roots is infinite; neither is the number  
117 of squares less than the totality of all the numbers, nor the lat-  
118 ter greater than the former; and finally the attributes "equal,"  
119 "greater," and "less," are not applicable to infinite, but only to  
120 finite, quantities. When therefore Simplicio introduces several  
121 lines of different lengths and asks me how it is possible that the  
122 longer ones do not contain more points than the shorter, I answer  
123 him that one line does not contain more or less or just as many  
124 points as another, but that each line contains an infinite number.

125 Leibniz insists in part being smaller than the whole saying

126 Among numbers there are infinite roots, infinite squares, infinite  
127 cubes. Moreover, there are as many roots as numbers. And there  
128 are as many squares as roots. Therefore there are as many squares  
129 as numbers, that is to say, there are as many square numbers as

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<sup>2</sup>zneni galileova paradoxu

there are numbers in the universe. Which is impossible. Hence it follows either that in the infinite the whole is not greater than the part, which is the opinion of Galileo and Gregory of St. Vincent, and which I cannot accept; or that infinity itself is nothing, i.e. that it is not one and not a whole.

TODO nejakej Hegel-strucne?

TODO Cantor

TODO mene teologie, vice matematiky

TODO definovat pojmy (trida etc)

In his work, he defined transfinite numbers to extend existing natural number structure so it contains more objects that behave like natural numbers and are based on an object (rather a meta-object) that doesn't explicitly exist in the structure, but is closely related to it. This is the first instance of reflection. This paper will focus on taking this principle a step further, extending Cantor's (or Zermelo–Fraenkel's, to be more precise) universe so it includes objects so big, they could be considered the universe itself, in a certain sense.

TODO dal asi smazat

The original idea behind reflection principles probably comes from what could be informally called “universality of the universe”. The effort to precisely describe the universe of sets was natural and could be regarded as one of the impulses for formalization of naive set theory. If we try to express the universe as a set  $\{x|x = x\}$ , a paradox appears, because either our set is contained in itself and therefore is contained in a set (itself again), which contradicts the intuitive notion of a universe that contains everything but is not contained itself.

TODO ???

If there is an object containing all sets, it must not be a set itself. The notion of class seems inevitable. Either directly the ways for example the Bernays–Gödel set theory, we will also discuss later in this paper, does in, or on a meta-level like the Zermelo–Fraenkel set theory, that doesn't refer to them in the axioms but often works with the notion of a universal class. duet Another obstacle of constructing a set of all sets comes from Georg Cantor, who proved that the set of all subsets of a set (let  $A$  be the set and  $\mathcal{P}()A$  its powerset) is strictly larger than  $A$ . That would turn every aspiration to finally establish an universal set into a contradictory infinite regression.<sup>3</sup> We will use  $V$  to denote the class of all sets. From previous thoughts we can

<sup>3</sup>An intuitive analogy of this *reductio ad infinitum* is the status of  $\omega$ , which was originally thought to be an unreachable absolute, only to become starting point of Cantor's hierarchy of sets growing beyond all boundaries around the end of the 19<sup>th</sup> century

167 easily argue, that it is impossible to construct a property that holds for  $V$   
 168 and no set and is neither paradoxical like  $\{x|x = x\}$  nor trivial. Previous  
 169 observation can be transposed to a rather naive formulation of the reflection  
 170 principle:

171 (Refl) Any property which holds in  $V$  already holds in some initial seg-  
 172 ment of  $V$ .

173 To avoid vagueness of the term "property", we could informally reformu-  
 174 late the above statement into a schema:

175 For every first-order formula<sup>4</sup>  $\varphi$  holds in  $V \leftrightarrow \varphi$  holds in some initial  
 176 segment of  $V$ .

177 Interested reader should note that this is a theorem scheme rather than  
 178 a single theorem.<sup>5</sup>

## 179 1.2 A few historical remarks on reflection

180 Reflection made it's first in set-theoretical appearance in Gödel's proof of  
 181 GCH in  $L$  (citace Kanamori ? Lévy and set theory), but it was around  
 182 even earlier as a concept. Gödel himself regarded it as very close to Russel's  
 183 reducibility axiom (an earlier equivalent of the axiom schema of Zermelo's  
 184 separation). Richard Montague then studied reflection properties as a tool  
 185 for verifying that Replacement is not finitely axiomatizable (citace?). a few  
 186 years later Lévy proved (citace? 1960a) equivalence of reflection with Axiom  
 187 of infinity together with Replacement in proof we shall examine closely in  
 188 chapter 2.

189 TODO co dal? recent results?

## 190 1.3 Notation (??) TODO

- 191 1. (*Refl*) je obecne reflexe (jaka presne)
- 192 2. (*Refl*)<sub>1</sub> je reflexe prvoradovych formulí
- 193 3. (*Refl*)<sub>2</sub> je reflexe druhoradovych formulí
- 194 4. etc...

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<sup>4</sup>this also works for finite sets of formulas [3, p. 168]

<sup>5</sup>If there were a single theorem stating "for any formula  $\varphi$  that holds in  $V$  there is an initial segment of  $V$  where  $\varphi$  also holds", we would obtain the following contradiction with the second Gödel's theorem: In ZFC, any finite group of axioms of ZFC holds in some initial segment of the universe. If we take the largest of those initial segments it is still strictly smaller than the universe and thus we have, via compactness, constructed a model of ZFC within ZFC. That is, of course a harsh contradiction. This also leads to an elegant way to prove that ZFC is not finitely axiomatizable.



## 2 Levy's first-order reflection

### 2.1 Introduction

This section will try to present Lévy's proof of a general reflection principle being equivalent to Replacement and Infinity under ZF minus Replacement and Infinity. We will first introduce a few axioms and definitions that were a different in Lévy's paper[2], but are equivalent to today's terms. We will write them in contemporary notation, our aim is the result, not history of set theory notation.

Please note that Lévy's paper was written in a period when Set theory was oriented towards semantics, which means that everything was done in a model. All proofs were theodel that of ZFC was  $V_\alpha$  (notated as  $R(\alpha)$  at the time) for some cardinal  $\alpha$ , which means that  $\alpha$  is a inaccessible cardinal. Please bear in mind that this is vastly different from saying that there is an inaccessible  $\alpha$  inside the model. This  $V_\alpha$  is also referred to as  $Scm^Q(u)$ , which means that  $u$  ( $u = V_\alpha$  in our case) is a standard complete model of an undisclosed axiomatic set theory  $Q$  formulated in the "non-simple applied first order functional calculus", which is second-order theory is today's terminology, we are allowed to quantify over functions and thus get rid of axiom schemes. (Note that Lévy always speaks of "the axiom of replacement"). Besides placeholder set theory  $Q$  and ZF, which the reader should be familiar with, theories  $Z$ ,  $S$ , and  $SF$  are used in the text.  $Z$  is ZF minus replacement,  $S$  is ZF minus replacement and infinity, and finally  $SF$  is ZF minus infinity. "The axiom of subsets" is an older name for the axiom scheme of specification (and it's not a scheme since we are now working in second order logic). Also note that universal quantifier does not appear,  $\forall x\varphi(x)$  would be written as  $(x)\varphi(x)$ , the symbol for negation is " $\sim$ ".

Lévy then mentions Mahlo's arithmetic construction of cardinals, noting, that he will use similar strategy to build higher levels of strong axioms of infinity.

TODO porovnani Mahlovy a Lévyho konstrukce, viz ref{mahlovy kardinaly}

TODO asi doplnit jak to souvisi se soucasnou definici slabe Mahlovych kardinalu pres stacionarni mnoziny?

### 2.2 Preliminaries

**Definition 2.1** *Relativization* *TODO (jech:161)*

## 2.3 Lévy's Original Proof From 1960

**Definition 2.2**  $N_0(\varphi)$

$$\exists u(Scm^S(u) \& x_1, \dots, x_n \in u \rightarrow \varphi \leftrightarrow \varphi^u) \quad (2.1)$$

where  $\varphi$  is a formula which does not contain free variables except  $x_1, \dots, x_n$ .

TODO muzu vyhodit

**Theorem 2.3** In  $S$ , the schema  $N_0$  implies the Axiom of Infinity.

*Proof.* For any  $\varphi$ ,  $N_0$  gives us  $\exists u Scm^S(u)$ , which means that there is a set  $u$  that is identical to  $V_\alpha$  for some alpha, so  $\exists \alpha Scm^S(V_\alpha)$ . We don't know the exact size of this  $\alpha$ , but we know that  $\alpha \geq \omega$ , otherwise  $\alpha$  would be finite, therefore not closed under the powerset operation, which would contradict the axiom of powersets. In order to prove that it is a model of  $S$ , we would need to verify all axioms of  $S$ . We have already shown that  $\omega$  is closed under the powerset operation. Foundation, extensionality and comprehension are clear from the fact that we work in  $ZF^6$ , pairing is clear from the fact, that given two sets  $A, B$ , they have ranks  $a, b$ , without loss of generality we can assume that  $a \leq b$ , which means that  $A \in V_a \in V_b$ , therefore  $V_b$  is a set that satisfies the paring axiom: it contains both  $A$  and  $B$ .

TODO vyhodit axiomy, staci vyrobit  $\omega$

We now want to prove that  $V_\alpha$  leads to existence of an inductive set, which is a set that satisfies  $\exists A(\emptyset \in A \& \forall x \in A((x \cup \{x\}) \in A))$ . If we can find a way to construct  $V_\omega$  from any  $V_\alpha$  satisfying  $\alpha \geq \omega$ , we are done. Since  $\omega$  is the least limit ordinal, all we need is the following

$$\bigcap \{V_\kappa \mid \forall \lambda(\lambda < \kappa \rightarrow \exists \mu(\lambda < \mu < \kappa))\} \quad (2.2)$$

because  $V_\kappa$  is a transitive set for every  $\kappa$ , thus the intersection is non-empty unless empty set satisfies the property or the set of  $V_\kappa$ s is itself empty.  $\square$

**Theorem 2.4** In  $S$ , the schema  $N_0$  implies Replacement schema.

*Proof.* TODO vysvetlit! (podle contemporary verze)

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<sup>6</sup>We only need to verify axioms that provide means of constructing larger sets from smaller to make sure they don't exceed  $\omega$ . Since  $\omega$  is an initial segment of  $ZF$ , the axiom scheme of specification can't be broken, the same holds for foundation and extensionality.

Let  $\varphi(v, w)$  be a formula wth no free variables except  $v, w, x_1, \dots, x_n$  where  $n$  is any natural number. Let  $\chi$  be an instance of replacement schema for this  $\varphi$ :

$$\chi = \forall r, s, t(\varphi(r, s) \& \varphi(r, t) \rightarrow s = t) \rightarrow \forall x \exists y \forall w (w \in y \leftrightarrow \exists v (v \in x \& \varphi(v, w))) \quad (2.3)$$

We can deduce the following from  $N_0$ :

- (i)  $x_1, \dots, x_n, v, w \in u \rightarrow (\varphi \leftrightarrow \varphi^u)$
- (ii)  $x_1, \dots, x_n, v \in u \rightarrow (\exists w \varphi \leftrightarrow (\exists w \varphi)^u)$
- (iii)  $x_1, \dots, x_n, x \in u \rightarrow (\chi \leftrightarrow \chi^u)$
- (iv)  $\forall x_1, \dots, x_n \forall x (\chi \leftrightarrow (\forall x_1, \dots, x_n \forall x \chi)^u)$

Note that (i), (ii), (iii) are obtained from instances of  $N_0$  for  $\varphi$ ,  $\exists w \varphi$  and  $\chi$  respectively. From relativization we also know that  $(\exists w \varphi)^u$  is equivalent to  $\exists w (w \in u \& \varphi^u)$ . Therefore (ii) is equivalent to  $x_1, \dots, x_n, v \in u \rightarrow (\exists w (w \in u \& \varphi^u))$ .

If  $\varphi$  is a function  $(\forall r, s, t(\varphi(r, s) \& \varphi(r, t) \rightarrow r = t))$ , then for every  $x \in u$ , which is also  $x \subset u$  by  $Scm^S(u)$ , it maps elements of  $x$  onto  $u$ . From the axiom scheme of comprehension<sup>7</sup>, we can find a set of all images of elements of  $x$ . Let's call it  $y$ . That gives us  $x_1, \dots, x_n, x \in u \rightarrow \chi$ . By (iii) we get  $x_1, \dots, x_n, x \in u \rightarrow \chi^u$ , closure of this formula is  $(\forall x_1, \dots, x_n \forall x \chi)^u$ , which together with (iv) yields  $\forall x_1, \dots, x_n \forall x \chi$ . By the means of specification we end up with  $\chi$ , which is all we need for now.

TODO btw co je x? nemela by tam tam byt nejaka volna promenna?

□

## 2.4 Contemporary restatement

TODO nejaký uvod.

TODO Levy rika ze existuje  $Scm^S(u)$  reflektujici varphi, coz uz nepotrebuje. atd.

TODO Ze prvoradova reflexe je theorem ZFC, vys uz max jako axiom/schema.

TODO ?

The following lemma is usually done in more parts, the first being with one formula and the other with  $n$ . We will only state and prove the generalised version for  $n$  formulas, knowing that  $n = 1$  is just a specific case and the proof is exactly the same.

<sup>7</sup>axiom of subsets in Levy's version

287 **Lemma 2.5** *Lemma Let  $\varphi_1, \dots, \varphi_n$  be any formulas with  $m$  parameters<sup>8</sup>.*

288 (i) *For each set  $M_0$  there is such  $M$  that  $M_0 \subset M$  and the following holds*  
 289 *for every  $i \leq n$ :*

$$\exists x \varphi_i(u_1, \dots, u_{m-1}, x) \rightarrow (\exists x \in M) \varphi_i(u_1, \dots, u_{m-1}, x) \quad (2.4)$$

290 *for every  $u_1, \dots, u_{m-1} \in M$ .*

291 (ii) *Furthermore there is an ordinal  $\alpha$  such that  $M_0 \subset V_\alpha$  and the following*  
 292 *holds for each  $i \leq n$ :*

$$\exists x \varphi_i(u_1, \dots, u_{m-1}, x) \rightarrow (\exists x \in V_\alpha) \varphi_i(u_1, \dots, u_{m-1}, x) \quad (2.5)$$

293 *for every  $u_1, \dots, u_{m-1} \in M$ .*

294 *Proof.* We will simultaneously prove statements (i) and (ii), denoting  $M^T$   
 295 the transitive set required by part (ii). Unless explicitly stated otherwise for  
 296 specific steps, it is thought to be equivalent to  $M$ .

297 Let us first define operation  $H(u_1, \dots, u_{m-1})$  that gives us the set of  
 298  $x$ 's with minimal rank satisfying  $\varphi_i(u_1, \dots, u_{m-1}, x)$  for given parameters  
 299  $u_1, \dots, u_{m-1}$  for every  $i \leq n$ .

$$H_i(u_1, \dots, u_n) = \{x \in C_i : (\forall z \in C)(\text{rank}(x) \leq \text{rank}(z))\} \quad (2.6)$$

300 for each  $i \leq n$ , where

$$C_i = \{x : \varphi_i(u_1, \dots, u_{m-1}, x)\} \text{ for } i \leq n \quad (2.7)$$

301 Next, let's construct  $M$  from given  $M_0$  by induction.

$$M_{i+1} = M_i \cup \bigcup_{j=0}^n \{H_j(u_1, \dots, u_{m-1}) : u_1, \dots, u_{m-1} \in M_i\} \quad (2.8)$$

302 In other words, in each step we add the elements satisfying  $\varphi(u_1, \dots, u_{m-1}, x)$   
 303 for all parameters that were either available earlier or were added in the  
 304 previous step. For statement (ii), this is the only part that differs from (i).

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<sup>8</sup>For formulas with different number of parameters take for  $m$  the highest number of parameters among given formulas. Add spare parameters to the other formulas so that  $x$  remains the last parameter. That can be done in a following manner: Let  $\varphi'_i$  be the a formula with  $k$  parameters,  $k < m$ . Let us set  $\varphi_i(u_1, \dots, u_{m-1}, x) = \varphi'_i(u_1, \dots, u_{k-1}, u_k, \dots, u_{m-1}, x)$ , notice that  $u_k, \dots, u_{m-1}$  are spare variables added just for formal simplicity.

Let us take for each step transitive closure of  $M_{i+1}$  from (i). In other words, let  $\gamma$  be the smallest ordinal such that

$$(M_i^T \cup \bigcup_{j=0}^n \{ \bigcup \{ H_j(u_1, \dots, u_{m-1}) : u_1, \dots, u_{m-1} \in M_i \} \}) \subset V_\gamma \quad (2.9)$$

Then the incremental step is like so:

$$M_{i+1}^T = V_\gamma \quad (2.10)$$

The final  $M$  is obtained by joining all incremental steps together.

$$M = \bigcup_{i=0}^{\infty} M_i, \quad M^T = \bigcup_{i=0}^{\infty} M_i^T \quad (2.11)$$

Let's try to construct a set  $M'$  that satisfies the same conditions like  $M$  but is kept as small as possible. Assuming the Axiom of Choice, we can modify the process so that cardinality of  $M'$  is at most  $|M_0| \cdot \aleph_0$ . Note that the size of  $M'$  is determined by the size of  $M_0$  and, most importantly, by the size of  $H_i(u_1, \dots, u_{m-1})$  for any  $i \leq n$  in individual levels of the construction. Since the lemma only states existence of some  $x$  that satisfies  $\varphi_i(u_1, \dots, u_{m-1}, x)$  for any  $i \leq n$ , we only need to add one  $x$  for every set of parameters but  $H_i(u_1, \dots, u_{m-1})$  can be arbitrarily large. Since Axiom of Choice ensures that there is a choice function, let  $F$  be a choice function on  $\mathcal{P}((M'))$ . Also let  $h_i(u_1, \dots, u_{m-1}) = F(H_i(u_1, \dots, u_{m-1}))$  for  $i \leq n$ , which means that  $h$  is a function that outputs an  $x$  that satisfies  $\varphi_i(u_1, \dots, u_{m-1}, x)$  for  $i \leq n$  and has minimal rank among all such witnesses. The induction step needs to be redefined to

$$M'_{i+1} = M'_i \cup \bigcup_j = 0^n \{ h_j(u_1, \dots, u_{m-1}) : u_1, \dots, u_{m-1} \in M'_i \} \quad (2.12)$$

In every step, the amount of elements added in  $M'_{i+1}$  is equivalent to the amount of sets of parameters the yielded elements not included in  $M'_i$ . So the cardinality of  $M'_{i+1}$  exceeds the cardinality of  $M'_i$  only for finite  $M'_i$ . It is easy to see that if  $M_0$  is finite,  $M'$  is countable because it was built from countable union of finite sets. If  $M_0$  is countable or larger, cardinality of  $M'$  is equal to the cardinality of  $M_0$ .<sup>9</sup> Therefore  $|M'| \leq |M_0| \cdot \aleph_0$

□

TODO proc  $\leq$  a ne =?

<sup>9</sup>It can not be smaller because  $|M'_{i+1}| \geq |M'_i|$  for every  $i$ . It may not be significantly larger because the maximum of elements added is the number of  $n$ -tuples in  $M'_i$ , which is of the same cardinality is  $M'_i$ . ((proc? Ramsey?))

330 **Theorem 2.6** *First-order Reflection*  $\varphi(x_1, \dots, x_n)$  is a first-order formula.

331 (i) For every set  $M_0$  there exists  $M$  such that  $M_0 \subset M$  and the following  
332 holds:

$$\varphi^M(x_1, \dots, x_n) \leftrightarrow \varphi(x_1, \dots, x_n) \quad (2.13)$$

333 for every  $x_1, \dots, x_n$ .

334 (ii) For every set  $M_0$  there is a transitive set  $M$ ,  $M_0 \subset M$  such that the  
335 following holds:

$$\varphi^M(x_1, \dots, x_n) \leftrightarrow \varphi(x_1, \dots, x_n) \quad (2.14)$$

336 for every  $x_1, \dots, x_n$ .

337 (iii) For every set  $M_0$  there is  $\alpha$  such that  $M_0 \subset V_\alpha$  and the following holds:

$$\varphi^{V_\alpha}(x_1, \dots, x_n) \leftrightarrow \varphi(x_1, \dots, x_n) \quad (2.15)$$

338 for every  $x_1, \dots, x_n$ .

339 (iv) Assuming the Axiom of Choice, for every set  $M_0$  there is  $M$  such that  
340  $M_0 \subset M$  and  $|M| \leq |M_0| \cdot \aleph_0$  and the following holds:

$$\varphi^M(x_1, \dots, x_n) \leftrightarrow \varphi(x_1, \dots, x_n) \quad (2.16)$$

341 for every  $x_1, \dots, x_n$ .

342 *Proof.* Let's prove (i) for one formula  $\varphi$  via induction by complexity first.  
343 We can safely assume that  $\varphi$  contains no quantifiers besides  $\exists$  and no logical  
344 connectives other than  $\neg$  and  $\&$ . Assume that this  $M$  is obtained from  
345 lemma 2.5. The fact, that atomic formulas are reflected in every  $M$  comes  
346 directly from definition of relativization and the fact that they contain no  
347 quantifiers.<sup>10</sup> The same holds for formulas in the form of  $\varphi = \neg\varphi'$ . Let us  
348 recall the definition of relativization for those formulas in .

$$(\neg\varphi_1)^M \leftrightarrow \neg(\varphi_1^M) \quad (2.17)$$

349 Because we can assume from induction that  $\varphi'^M \leftrightarrow \varphi'$ , the following holds:

$$(\neg\varphi')^M \leftrightarrow \neg(\varphi'^M) \leftrightarrow \neg\varphi' \quad (2.18)$$

350 The same holds for  $\varphi = \varphi_1 \& \varphi_2$ . From the induction hypothesis we know  
351 that  $\varphi_1^M \leftrightarrow \varphi_1$  and  $\varphi_2^M \leftrightarrow \varphi_2$ , which together with relativization for formulas  
352 in the form of  $\varphi_1 \& \varphi_2$  gives us

$$(\varphi_1 \& \varphi_2)^M \leftrightarrow \varphi_1^M \& \varphi_2^M \leftrightarrow \varphi_1 \& \varphi_2 \quad (2.19)$$

---

<sup>10</sup>Note that this does not hold generally for relativizations to  $M, E$ , but only for relativization to  $M, \in$ , which is our case.

Let's now examine the case when from the induction hypethesis,  $M$  reflects  $\varphi'(u_1, \dots, u_n, x)$  and we are interested in  $\varphi = \exists x \varphi'(u_1, \dots, u_n, x)$ . The induction hypothesis tells us that

$$\varphi'^M(u_1, \dots, u_n, x) \leftrightarrow \varphi'(u_1, \dots, u_n, x) \quad (2.20)$$

so, together with above lemma 2.5, the following holds:

$$\varphi(u_1, \dots, u_n, x) \quad (2.21)$$

$$\leftrightarrow \exists x \varphi'(u_1, \dots, u_n, x) \quad (2.22)$$

$$\leftrightarrow (\exists x \in M) \varphi'(u_1, \dots, u_n, x) \quad (2.23)$$

$$\leftrightarrow (\exists x \in M) \varphi'^M(u_1, \dots, u_n, x) \quad (2.24)$$

$$\leftrightarrow (\exists x \varphi'(u_1, \dots, u_n, x))^M \quad (2.25)$$

$$\leftrightarrow \varphi^M(u_1, \dots, u_n, x) \quad (2.26)$$

Which is what we have needed to prove:

So far we have proven part (i) of this theorem for one formula  $\varphi$ , we only need to verify that the same holds for any finite number of formulas. This has in fact been already done since lemma 2.5 gives us  $M$  for any (finite) amount of formulas. We can than use the induction above to verify that it reflects each of the formulas individually.

Now we want to verify other parts of our theorem. Since  $V_\alpha$  is a transitive set, by proving (iii) we also satisfy (ii). To do so, we only need to look at part (ii) of lemma 2.5. All of the above proof also holds for  $M = V_\alpha$ . To finish part (iv)

□

**Theorem 2.7** *(Refl) is equivalent to (Infinity) & (Replacement) under ZFC minus (Infinity) & (Replacement)*

*Proof.* Since 2.6 already gives one side of the implication, we are only interested in showing the converse:

(Refl)  $\rightarrow$  (Infinity)

Let us first find a formula to be reflected that requires a set  $M$  at least as large as  $V_\omega$ . Let us consider the following formula:

$$\varphi'(x) = \forall \lambda (\lambda < x \rightarrow \exists \mu (\lambda < \mu < x)) \quad (2.27)$$

Because  $\varphi$  says "there is a limit ordinal", if it holds for some  $x$ , the Infinity axiom is very easy to satisfy. But we know that there are limit ordinals in

382 ZF, therefore  $\varphi = \exists x\varphi'(x)$  is a valid statement. (Refl) then gives us a set  $M$   
 383 in which  $\varphi^M$  holds, that is, a set that contains a limit ordinal. So the set of  
 384 off limit ordinals is non-empty and because ordinals are well-founded, it has  
 385 a minimal element. Let's call it  $\mu$ .

$$\mu = \bigcap \{V_\kappa : \forall \lambda (\lambda < \kappa \rightarrow \exists \mu (\lambda < \mu < \kappa))\} \quad (2.28)$$

386 We can see that  $\mu$  is the least limit ordinal and therefore it satisfies (Infinity).

387 **(Refl)  $\rightarrow$  (Replacement)**

388 Given a formula  $\varphi(x, y, u_1, \dots, u_n)$ , we can suppose that it is reflected in  
 389 any  $M$ <sup>11</sup> What we want to obtain is the following:

$$\forall x, y, z (\varphi(x, y, u_1, \dots, u_n) \& \varphi(x, z, u_1, \dots, u_n) \rightarrow y = z) \rightarrow \quad (2.29)$$

$$\rightarrow \forall X \exists Y \forall y (y \in Y \leftrightarrow \exists x (\varphi(x, y, u_1, \dots, u_n) \& x \in X)) \quad (2.30)$$

391 We do also know that  $x, y \in M$ , in other words for every  $X, Y =$   
 392  $\{y \mid \varphi(x, y, u_1, \dots, u_n)\}$  we know that  $X \subset M$  and  $Y \subset M$ , which, together  
 393 with the comprehension schema<sup>12</sup> implies that  $Y$ , the image of  $X$  over  $\varphi$ , is  
 394 a set. Which is exactly the Replacement Schema we hoped to obtain.

395  $\square$  We have shown that (Refl) for first-order formulas is a theorem of  
 396 ZF, which means that it won't yield us any large cardinals. We have shown  
 397 that it can be used instead of the Axiom of Infinity and Replacement Scheme,  
 398 but ZF + (Refl) is a conservative extension of ZF. Besides being a starting  
 399 point for more general and powerful statements, it can be used to show that  
 400 ZF is not finitely axiomatizable. That is because (Refl) gives a model to  
 401 any finite number of (consistent) formulas. So if  $\varphi_1, \dots, \varphi_n$  for any finite  $n$   
 402 would be the axioms of ZF, (Refl) would always contain a model of itself,  
 403 which would in turn contradict the Second Gödel's Theorem<sup>13</sup>. Notice, that  
 404 reflection in a way counterpart to compactness. Compactness argues that  
 405 for an infinite<sup>14</sup> set of sentences, if every finite subset yields a model, then so  
 406 does the whole set. Reflection, on the other hand, says that while the whole  
 407 set has no model in the underlying theory, every finite subset does have a  
 408 model.

409 Also, notice how reflection can be used in ways similar to upward Löwenheim–Skolem  
 410 theorem. Since Reflection extends any set  $M_0$  into a model of given formulas  
 411  $\varphi_1, \dots, \varphi_n$ , we can choose  $M_0$  such that the final  $M$  is at least as big as we  
 412 need.<sup>15</sup>

<sup>11</sup>Which means that for  $x, y, u_1, \dots, u_n \in M$ ,  $\varphi^M(x, y, u_1, \dots, u_n) \leftrightarrow \varphi(x, y, u_1, \dots, u_n)$ .

<sup>12</sup>Called the axiom of subsets in Levy's proof.

<sup>13</sup>See chapter 3.3 for further details.

<sup>14</sup>Countable?

<sup>15</sup>Too vague?



413       TODO znacit  $(Refl)$  asi jako  $(Refl)_1$  pokud mluvim o prvoradovych  
414 formulich, nekde nahore zavest

415       TODO sjednotit kdy pisu Reflection a kdy  $(Refl)$

416       In the next section, we will try to generalize Reflection in a way that  
417 transcends ZF and finally yields us some large cardinals.

## 3 Large Cardinals and Higher-order Reflection

In this chapter we aim to explore possible generalisations of  $(Refl)$  for second- and higher-order formulas and use those to establish existence of various large cardinals. We will also argue whether there is a limit to the size of large cardinals accessible via generalised  $(Refl)$ .

### 3.1 Reflecting Second-order Formulas

To see that there is a way to transcend ZF, let us briefly show how a model of ZF can be obtained in  $ZF^2 + \text{"second - order reflection"}$ <sup>16</sup>. This will be more closely examined in section 3.3.

We know that ZF can not be finitely axiomatized in first-order formulas, however if Replacement and Comprehension schemes can be substituted by second-order formulas, ZF becomes  $ZF^2$ , which is finitely axiomatizable in second-order logic. Therefore if we take second-order reflection into consideration, we can obtain a set  $M$  that is a model of  $ZF^2$ . Therefore, second-order reflection. For now, we have left out the details of how exactly is first-order reflection generalised into stronger statements and how second-order axiomatization of ZF looks like. The purpose of this chapter is to try to answer these questions, as well as examine the relation of said reflection axioms to large cardinals.

We will now define reflection for second-order formulas.

**Definition 3.1** *Second-order reflection TODO*

### 3.2 Preliminaries

But first, let's establish some elementary terms that will allow us to define the relevant large cardinals.

**Definition 3.2** *(limit cardinal) kappa is a limit cardinal if it is  $\aleph_\alpha$  for some limit ordinal  $\alpha$ .*

**Definition 3.3** *(strong limit cardinal) kappa is a strong limit cardinal if for every  $\lambda < \kappa$ ,  $2^\lambda < \kappa$*

TODO budeme potrebovat club sety, stacionarni mnoziny? treba aby-chom dokazali ekvivalence?

---

<sup>16</sup>ZF<sup>2</sup> is an axiomatization of ZF in second-order formulas, to be more rigorously established later.

### 3.3 Inaccessibility

The inaccessible cardinal is the smallest of large cardinals<sup>17</sup>

**Definition 3.4** (*weak inaccessibility*) An uncountable cardinal  $\kappa$  is weakly inaccessible  $\leftrightarrow$  it is regular and limit.

**Definition 3.5** (*inaccessibility*) An uncountable cardinal  $\kappa$  is inaccessible  $\leftrightarrow$  it is regular and strongly limit.

Note that the above requirements could in fact be satisfied by  $V_\omega$ , except for the need for uncountability, which is being added exactly to leave out  $V_\omega$  for practical purposes.<sup>18</sup>

We will now show that the above notion is equivalent to the definition Levy uses in [2], which is, in more contemporary notation, the following:

**Theorem 3.6** *The following are equivalent:*

1.  $\kappa$  is inaccessible
2.  $\langle V_\kappa, \in \rangle \models \text{ZFC}$

*Proof.* Let's first prove that if  $\kappa$  is inaccessible, it is a model of ZFC. We will do that by verifying the axioms of ZFC just like Kanamori does it in 1.2 in [1]. Because  $\kappa$  is a limit ordinal, there's no need for us to verify the Powerset, Foundation, Extensionality, Subset, Pairing and Union axioms and the Specification scheme. Thus we only have the Replacement Scheme to verify.

Given arbitrary set  $x \in V_\kappa$  and a function  $F : x \rightarrow V_\kappa$ , we need to verify that  $y = F[x]$  is indeed a set and that it is an element of  $V_\kappa$ . The fact that  $F$  is a function implies that  $|y| \leq |x|$ . It follows from Specification that  $y \subset V_\kappa$ , which is still not exactly what we want. Let  $\alpha < \kappa$  be the least ordinal such that  $y \in V_\alpha$ <sup>19</sup>, since  $V_{\alpha+1} = \mathcal{P}(V_\alpha)$ ,  $y \in V_{\alpha+1}$ , together with  $\alpha + 1 < \kappa$  this means that  $y \in V_\kappa$ .

□

### 3.4 Mahlo cardinals

TODO reflektuji nedosazitelnost? TODO zminit Mahlovu konstrukci v Levym?  
 TODO zavest pomoci reflexe

<sup>17</sup>citation needed.

<sup>18</sup>Informally, this clearly illustrates the fact stressed in section 1, that large cardinals are obtained by similar way of thinning that lead mathematicians to establish  $\omega$  as an actual object.

<sup>19</sup>TODO jak vime ze takove alpha existuje?

### 3.5 Weakly Compact Cardinals and Higher-order Reflection

**Definition 3.7** *Weakly Mahlo Cardinals*  $\kappa$  is weakly Mahlo  $\leftrightarrow$  it is a limit ordinal and the set of all regular ordinals less than  $\kappa$  is stationary in  $\kappa$

TODO napsat co to znamena

**Definition 3.8** *Mahlo cardinals* The following definitions are equivalent:

- (i)  $\kappa$  is Mahlo
- (ii)  $\kappa$  is weakly Mahlo and strong limit
- (iii)  $\kappa$  is inaccessible and the regular cardinals below  $\kappa$  form a stationary subset of  $\kappa$ .
- (iv)  $\kappa$  is regular and the stationary sets below  $\kappa$  form a stationary subset of  $\kappa$ .

Note that Mahlo cardinals were first described in 1911, almost 50 years before Lévy's reflection, which was heavily inspired by them.

### 3.5 Weakly Compact Cardinals

TODO souvislost s reflexi! TODO co je "partition property"?

**Definition 3.9** *A cardinal  $\kappa$  is weakly compact if it is uncountable and satisfies the partition property  $\kappa \rightarrow (\kappa)^2$*

opsano z jecha!

### 3.6 Indescribable Cardinals

TODO uvod / intuice

TODO souvislost s reflexi

### 3.7 Bernays–Gödel Set Theory

TODO Plagiat – prepsat a vysvetlit

TODO

### 3.8 Reflection and the constructible universe

TODO reflektovat muzeme jenom kardinaly konzistentni s  $V=L$ , proc?

TODO Plagiat – prepsat a vysvetlit

$L$  was introduced by Kurt Gödel in 1938 in his paper *The Consistency of the Axiom of Choice and of the Generalised Continuum Hypothesis* and denotes a class of sets built recursively in terms of simpler sets, somewhat

508 similar to Von Neumann universe  $V$ . Assertion of their equality,  $V = L$ , is  
 509 called the *axiom of constructibility*. The axiom implies GCH and therefore  
 510 also AC and contradicts the existence of some of the large cardinals, our goal  
 511 is to decide whether those introduced earlier are among them.

512 On order to formally establish this class, we need to formalize the notion  
 513 of definability first:

514 TODO zduvodneni

515 TODO kratka diskuse jestli refl implikuje transcendenci na  $L$  - polemika,  
 516 nazor -  $V=L$  a slaba kompaktnost a dalsi

## 517 4 Higher-order reflection

518 TODO rict ze to je zobecneni a nejaky dalsi uvodni veci

### 519 4.1 Sharp

520 TODO je tohle higher-order vec?

### 521 4.2 Weleik: Global Reflection Principles

522 TODO

## 523 5 Conclusion

524 TODO na konec

## References

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