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- REFLECTION PRINCIPLES AND LARGE
- 5 CARDINALS

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Bakalářská práce

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 $^{10}\,$ Prohlašuj, že jsem bakalářkou práci vypracoval samostatně a že jsem uvedl $^{11}\,$ všechny použité prameny a literaturu.

12 V Praze 14. dubna 2015

13 Mikuláš Mrva

14 Abstract

Práce zkoumá vztah tzv. principů reflexe a velkých kardinálů. Lévy ukázal, že v ZFC platí tzv. věta o reflexi a dokonce, že věta o reflexi je ekvivalentní schématu nahrazení a axiomu nekonečna nad teorií ZFC bez axiomu nekonečna a schématu nahrazení. Tedy lze na větu o reflexi pohlížet jako na svého druhu axiom nekonečna. Práce zkoumá do jaké míry a jakým způsobem lze větu o reflexi zobecnit a jaký to má vliv na existenci tzv. velkých kardinálů. Práce definuje nedosažitelné, Mahlovy a nepopsatelné kardinály a ukáže, jak je lze zavést pomocí reflexe. Přirozenou limitou kardinálů získaných reflexí jsou kardinály nekonzistentní s L. Práce nabídne intuitivní zdůvodněn, proč tomu tak je.

26 Abstract

TODO Resumé práce v anglickém jazyce.

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1 Introduction

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1.1 Motivation and Origin

The Universe of sets cannot be uniquely characterized (i. e. distinguished from all its initial elements) by any internal structural property of the membership relation in it, which is expressible in any logic of finite of transfinite type, including infinitary logics of any cardinal order.

— Kurt Gödel [4]

To understand why do need reflection in the first place, let's think about 57 infinity for a moment. In the intuitive sense, infinity is an upper limit of all 58 numbers. But for centuries, this was merely a philosophical concept, closely 59 bound to religious and metaphysical way of thinking, considered separate from numbers used for calculations or geometry. It was a rather vague concept. In ancient Greece, Aristotle's response to famous Zeno's paradoxes 62 introduced the distinction between actual and potential infinity. He argued, 63 that potential infinity is (in today's words) well defined, as opposed to actual 64 infinity, which remained a vague incoherent concept. He didn't think it's pos-65 sible for infinity to inhabit a bounded place in space or time, rejecting Zeno's 66 thought experiments as a whole. Aristotle's thoughts shaped western think-67 ing partly due to Aquinas, who himself believed actual infinity to be more of a metaphysical concept for describing God than a mathematical property 69 attributed to any other entity. In his Summa Theologica ¹ he argues: 70

A geometrician does not need to assume a line actually infinite, but takes some actually finite line, from which he subtracts whatever he finds necessary; which line he calls infinite.

Less than hundred years later, Gregory of Rimini wrote

If God can endlessly add a cubic foot to a stone—which He can—then He can create an infinitely big stone. For He need only add one cubic foot at some time, another half an hour later, another a quarter of an hour later than that, and so on ad infinitum. He would then have before Him an infinite stone at the end of the hour.

Which is basically a Zeno's Paradox made plausible with God being the actor. In contrast to Aquinas' position, Gregory of Rimini theoretically constructs

¹Part I, Question 7, Article 3, Reply to Objection 1

an object with actual infinite magnitude that is essentially different from God. Even later, in the 17th century, pushing the property of infinitness from the Creator to his creation, Nature, Leibniz wrote to Foucher in 1962:

I am so in favor of the actual infinite that instead of admitting that Nature abhors it, as is commonly said, I hold that Nature makes frequent use of it everywhere, in order to show more effectively the perfections of its Author. Thus I believe that there is no part of matter which is not, I do not say divisible, but actually divided; and consequently the least particle ought to be considered as a world full of an infinity of different creatures.

But even though he used potential infinity in what would become foundations of modern Calculus and argued for actual infinity in Nature, Leibniz refused the existence of an infinite, thinking that Galileo's Paradoxon² is in fact a contradiction. The so called Galileo's Paradoxon is an observation Galileo Galilei made in his final book "Discourses and Mathematical Demonstrations Relating to Two New Sciences". He states that if all numbers are either squares and non-squares, there seem to be less squares than there is all numbers. On the other hand, every number can be squared and every square has it's square root. Therefore, there seem to be as many squares as there are all numbers. Galileo concludes, that the idea of comparing sizes makes sense only in the finite realm.

Salviati: So far as I see we can only infer that the totality of all numbers is infinite, that the number of squares is infinite, and that the number of their roots is infinite; neither is the number of squares less than the totality of all the numbers, nor the latter greater than the former; and finally the attributes "equal," "greater," and "less," are not applicable to infinite, but only to finite, quantities. When therefore Simplicio introduces several lines of different lengths and asks me how it is possible that the longer ones do not contain more points than the shorter, I answer him that one line does not contain more or less or just as many points as another, but that each line contains an infinite number.

Leibniz insists in part being smaller than the whole saying

Among numbers there are infinite roots, infinite squares, infinite cubes. Moreover, there are as many roots as numbers. And there are as many squares as roots. Therefore there are as many squares

²zneni galileova paradoxu

as numbers, that is to say, there are as many square numbers as there are numbers in the universe. Which is impossible. Hence it follows either that in the infinite the whole is not greater than the part, which is the opinion of Galileo and Gregory of St. Vincent, and which I cannot accept; or that infinity itself is nothing, i.e. that it is not one and not a whole.

TODO nejakej Hegel-strucne?

TODO Cantor

TODO mene teologie, vice matematiky

TODO definovat pojmy (trida etc)

In his work, he defined transfinite numbers to extend existing natural number structure so it contains more objects that behave like natural numbers and are based on an object (rather a meta-object) that doesn't explicitly exist in the structure, but is closely related to it. This is the first instance of reflection. This paper will focus on taking this principle a step further, extending Cantor's (or Zermelo–Fraenkel's, to be more precise) universe so it includes objects so big, they could be considered the universe itself, in a certain sense.

TODO dal asi smazat

The original idea behind reflection principles probably comes from what could be informally called "universality of the universe". The effort to precisely describe the universe of sets was natural and could be regarded as one of the impulses for formalization of naive set theory. If we try to express the universe as a set $\{x|x=x\}$, a paradox appears, because either our set is contained in itself and therefore is contained in a set (itself again), which contradicts the intuitive notion of a universe that contains everything but is not contained itself.

TODO ???

If there is an object containing all sets, it must not be a set itself. The notion of class seems inevitable. Either directly the ways for example the Bernays–Gödel set theory, we will also discuss later in this paper, does in, or on a meta–level like the Zermelo–Fraenkel set theory, that doesn't refer to them in the axioms but often works with the notion of a universal class. duet Another obstacle of constructing a set of all sets comes from Georg Cantor, who proved that the set of all subsets of a set (let A be the set and $\mathcal{P}(()A)$ its powerset) is strictly larger that A. That would turn every aspiration to finally establish an universal set into a contradictory infinite regression.³. We

³An intuitive analogy of this reductio ad infinitum is the status of ω , which was originally thought to be an unreachable absolute, only to become starting point of Cantor's hierarchy of sets growing beyond all boundaries around the end of the 19th century

 $_{\mbox{\tiny 156}}$ $\,$ will use V to denote the class of all sets.

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From previous thoughts we can easily argue, that it is impossible to construct a property that holds for V and no set and is neither paradoxical like $\{x|x=x\}$ nor trivial. Previous observation can be transposed to a rather naive formulation of the reflection principle:

(Refl) Any property which holds in V already holds in some initial segment of V.

To avoid vagueness of the term "property", we could informally reformulate the above statement into a schema:

For every first-order formula⁴ φ holds in $V \leftrightarrow \varphi$ holds in some initial segment of V.

Interested reader should note that this is a theorem scheme rather than a single theorem. ⁵

1.2 A few historical remarks on reflection

Reflection made it's first in set-theoretical appearance in Gödel's proof of GCH in L (citace Kanamori? Lévy and set theory), but it was around even earlier as a concept. Gödel himself regarded it as very close to Russel's reducibility axiom (an earlier equivalent of the axiom schema of Zermelo's separation). Richard Montague then studied reflection properties as a tool for verifying that Replacement is not finitely axiomatizable (citace?). A few years later Lévy proved (citace? 1960a) equivalence of reflection with Axiom of infinity together with Replacement in proof we shall examine closely in chaper 2.

TODO co dal? recent results?

⁴this also works for finite sets of formulas [3, p. 168]

⁵If there were a single theorem stating "for any formula φ that holds in V there is an initial segment of V where φ also holds", we would obtain the following contradiction with the second Gödel's theorem: In ZFC, any finite group of axioms of ZFC holds in some initial segment of the universe. If we take the largest of those initial segments it is still strictly smaller than the universe and thus we have, via compactness, constructed a model of ZFC within ZFC. That is, of course a harsh contradiction. This also leads to an elegant way to prove that ZFC is not finitely axiomatizable.

2 Levy's first-order reflection

2.1 Introduction

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This section will try to present Lévy's proof of a general reflection principle being equivalent to Replacement and Infinity under ZF minus Replacement and Infinity. We will first introduce a few axioms and definitions that were a different in Lévy's paper[2], but are equivalent to today's terms. We will write them in contemporary notation, our aim is the result, not history of set theory notation.

Please note that Lévy's paper was written in a period when Set theory was oriented towards semantics, which means that everything was done in a model. All proofs were theodel that of ZFC was V_{α} (notated as $R(\alpha)$ at the time) for some cardinal α , which means that α is a inaccessible cadinal. Please bear in mind that this is vastly different from saying that there is an inaccessible α inside the model. This V_{α} is also referred to as $Scm^{\mathbb{Q}}(u)$, which means that u ($u = V_{\alpha}$ in our case) is a standard complete model of an undisclosed axiomatic set theory Q formulated in the "non-simple applied first order functional calculus", which is second-order theory is today's terminology, we are allowed to quantify over functions and thus get rid of axiom schemes. (Note that Lévy always speaks of "the axiom of replacement"). Besides placeholder set theory Q and ZF, which the reader should be familiar with, theories Z, S, and SF are used in the text. Z is ZF minus replacement, S is ZF minus replacement and infinity, and finally SF is ZF minus infinity. "The axiom of subsets" is an older name for the axiom scheme of specification (and it's not a scheme since we are now working in second order logic). Also note that universal quantifier does not appear, $\forall x \varphi(x)$ would be written as $(x)\varphi(x)$, the symbol for negation is " \sim ".

Lévy then mentions Mahlo's arithmetic construction of cardinals, noting, that he will use similar strategy to build higher levels of strong axioms of infinity.

TODO porovnani Mahlovy a Lévyho konstrukce

TODO asi doplnit jak to souvisi se soucasnou definici slabe Mahlovych kardinalu pres stacionarni mnoziny?

2.2 Preliminaries

Definition 2.1 Relativization TODO (jech:161)

2.3 Lévy's Original Proof From 1960

Definition 2.2 $N_0(\varphi)$

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$$\exists u(Scm^{\mathsf{S}}(u)\&x_1,\ldots,x_n\in u\to\varphi\leftrightarrow\varphi^u)$$
 (2.1)

where φ is a formula which does not contain free variables except x_1, \ldots, x_n .

TODO muzu vyhodit

Theorem 2.3 In S, the schema N_0 implies the Axiom of Infinity.

Proof. For any φ , N_0 gives us $\exists uScm^{\mathsf{S}}(u)$, which means that there is a set u 219 that is identical to V_{α} for some alpha, so $\exists \alpha Scm^{\mathsf{S}}(V_{\alpha})$. We don't know the 220 exact size of this α , but we know that $\alpha \geq \omega$, otherwise α would be finite, 221 therefore not closed under the powerset operation, which would contradict 222 the axiom of powersets. In order to prove that it is a model of S, we would 223 need to verify all axioms of S. We have already shown that ω is closed under the powerset operation. Foundation, extensionality and comprehension are 225 clear from the fact that we work in ZF^6 , pairing is clear from the fact, that 226 given two sets A, B, they have ranks a, b, without loss of generality we can 227 assume that $a \leq b$, which means that $A \in V_a \in V_b$, therefore V_b is a set that 228 satisfies the paring axiom: it contains both A and B. 229

TODO vyhodit axiomy, staci vyrobit ω

We now want to prove that V_{α} leads to existence of an inductive set, which is a set that satisfies $\exists A(\emptyset \in A \& \forall x \in A((x \cup \{x\}) \in A))$. If we can find a way to construct V_{ω} from any V_{α} satisfying $\alpha \geq \omega$, we are done. Since ω is the least limit ordinal, all we need is the following

$$\bigcap \{V_{\kappa} \mid \forall \lambda (\lambda < \kappa \to \exists \mu (\lambda < \mu < \kappa))\}$$
 (2.2)

because V_{κ} is a transitive set for every κ , thus the intersection is non-empty unless empty set satisfies the property or the set of V_{κ} s is itself empty. \square

Theorem 2.4 In S, the schema N_0 implies Replacement schema.

8 Proof. TODO vysvetlit!

⁶We only need to verify axioms that provide means of constructing larger sets from smaller to make sure they don't exceed ω . Since ω is an initial segment of ZF, the axiom scheme of specification can't be broken, the same holds for foundation and extensionality.

Let $\varphi(v,w)$ be a formula wth no free variables except v,w,x_1,\ldots,x_n where n is any natural number. Let χ be an instance of replacement schema for this φ :

$$\chi = \forall r, s, t(\varphi(r, s)\&\varphi(r, t) \to s = t) \to \forall x \exists y \forall w (w \in y \leftrightarrow \exists v (v \in x\&\varphi(v, w)))$$
(2.3)

We can deduce the following from N_0 :

- (i) $x_1, \ldots, x_n, v, w \in u \to (\varphi \leftrightarrow \varphi^u)$
 - (ii) $x_1, \ldots, x_n, v \in u \to (\exists w \varphi \leftrightarrow (\exists w \varphi)^u)$
- 245 (iii) $x_1, \ldots, x_n, x \in u \to (\chi \leftrightarrow \chi^u)$

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(iv) $\forall x_1, \dots, x_n \forall x (\chi \leftrightarrow (\forall x_1, \dots, x_n \forall x \chi)^u)$

Note that (i), (ii), (iii) are obtained from instances of N_0 for φ , $\exists w\varphi$ and χ respectively. From relativization we also know that $(\exists w\varphi)^u$ is equivalent to $\exists w(w \in u\&\varphi^u)$. Therefore (ii) is equivalent to $x_1,\ldots,x_n,v\in u\to (\exists w(w\in u\&\varphi^u))$.

If φ is a function $(\forall r, s, t(\varphi(r, s)\&\varphi(r, t) \to r = t))$, then for every $x \in u$, which is also $x \subset u$ by $Scm^{\mathsf{S}}(u)$, it maps elements of x onto u. From the axiom scheme of comprehension⁷, we can find a set of all images of elements of x. Let's call it y. That gives us $x_1, \ldots, x_n, x \in u \to \chi$. By (iii) we get $x_1, \ldots, x_n, x \in u \to \chi^u$, closure of this formula is $(\forall x_1, \ldots, x_n \forall x_1, \ldots, x_n \forall x_n, x_n \forall$

TODO btw co je x? nemela by tam tam byt nejaka volna promenna?

2.9 2.4 Contemporary restatement

TODO nejaky uvod. Levy rika ze existuje $Scm^S(u)$ reflektujici varphi, coz uz nepotrebujeme. atd.

The following lemma is usually done in parts, the first being with one formula and the other with n. Will will only state and prove the generalised version for n formulas.

Lemma 2.5 Lemma Let $\varphi_1, \ldots, \varphi_n$ be any formulas with m parameters⁸.

⁷axiom of subsets in Levy's version

⁸For formulas with different number of parameters take for m the highest number of parameters among given formulas. Add spare parameters to the other formulas so that x remains the last parameter. That can be done in a following manner: Let φ_i' be the a formula with k parameters, k < m. Let us set $\varphi_i(u_1, \ldots, u_{m-1}, x) = \varphi_i'(u_1, \ldots, u_{k-1}, u_k, \ldots, u_{m-1}, x)$, notice that u_k, \ldots, u_{m-1} are spare variables added just for formal simplicity.

(i) For each set M_0 there is such M that $M_0 \subset M$ and the following holds for every $i \leq n$:

$$\exists x \varphi_i(u_1, \dots, u_{m-1}, x) \to (\exists x \in M) \varphi_i(u_1, \dots, u_{m-1}, x)$$
 (2.4)

(ii) Furthermore there is an ordinal α such that $M_0 \subset V_\alpha$ and the following holds for each $i \leq n$:

$$\exists x \varphi_i(u_1, \dots, u_{m-1}, x) \to (\exists x \in V_\alpha) \varphi_i(u_1, \dots, u_{m-1}, x)$$
 (2.5)

Proof. We will simultaneously prove statements (i) and (ii), denoting M^T the transitive model required by part (ii). Unless explicitly stated otherwise, it is equivalent to M.

Let us first define operation $H(u_1, \ldots, u_{m-1})$ that gives us the set of x's with minimal rank, satisfying $\varphi_i(u_1, ldots, u_{m-1}, x)$ for given parameteres u_1, \ldots, u_{m-1} for every $i \leq n$.

$$H_i(u_1, \dots, u_n) = \{x \in C_i : (\forall z \in C)(rank(x) \le rank(z))\} \text{ for } i \le n \quad (2.6)$$

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$$C_i = \{x : \varphi_i(u_1, \dots, u_{m-1}, x)\} \text{ for } i \le n$$
 (2.7)

Next, let's construct M from given M_0 by induction.

(2.8)

In other words, in each step we add the elements satisfying $\varphi(u_1, \ldots, u_{m-1}, x)$ for those parameters that were either available earlier or were added in the previous step.

For statement (ii), this is the only part that differs from (i). Let us take for each step transitive closure of M_{i+1} from (i). In other words, let γ be the smallest ordinal such that

$$(M_i^T \cup \bigcup_{j=0}^n \{ \bigcup \{ H_j(u_1, \dots, u_{m-1}) : u_1, \dots, u_{m-1} \in M_i \} \}) \subset V_\gamma$$
 (2.9)

Then the incremetal step is like so:

$$M_{i+1}^T = V_{\gamma} \tag{2.10}$$

The final M is obtained by joining all incremental steps together.

$$M = \bigcup_{i=0}^{\infty} M_i, \ M^T = \bigcup_{i=0}^{\infty} M_i^T$$
 (2.11)

Let's try to construct a set M' that satisfies the same conditions like 286 M but is kept as small as possible. Assuming the Axiom of Choice, we can 287 modify the process so that cardinality of M' is at most $|M_0| \cdot \aleph_0$. Note that the 288 size of M' is determined by the size of M_0 an, most importantly, by the size of 289 $H_i(u_1,\ldots,u_{m-1})$ for any $i\leq n$ in individual levels of the construction. Since 290 the lemma only states existence of some x that satisfies $\varphi_i(u_1,\ldots,u_{m-1},x)$ 291 for any i < n, we only need to add one x for every set of parameters but 292 $H_i(u_1,\ldots,u_{m-1})$ can be arbitrarily large. Since Axiom of Choice ensures 293 that there is a choice function, let F be a choice function on $\mathscr{P}(()M')$. Also 294 let $h_i(u_1,\ldots,u_{m-1})=F(H_i(u_1,\ldots,u_{m-1}))$ for $i\leq n$, which means that h is 295 a function that outputs an x that satisfies $\varphi_i(u_1,\ldots,u_{m-1},x)$ for $i\leq n$ and 296 has minimal rank among all such witnesses. The induction step needs to be redefined to 298

$$M'_{i+1} = M'_i \cup \bigcup_j = 0^n \{ h_j(u_1, \dots, u_{m-1}) : u_1, \dots, u_{m-1} \in M'_i \}$$
 (2.12)

In every step, the amount of elements added in M'_{i+1} is equivalent to the amount of sets of parameters the yielded elements not included in M'_i . So the cardinality of M'_{i+1} exceeds the cardinality of M'_i only for finite M'_i . It is easy to see that if M_0 is finite, M' is countable because it was built from countable union of finite sets. If M_0 is countable or larger, cardinaly of M' is equal to the cardinality of M_0 . Therefore $|M'| \leq |M_0| \cdot \aleph_0$

TODO proc \leq a ne =?

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TODO nemel bych neco rict o u_1, \ldots, u_n ? urcite se dostanou do nektereho M_i ?

Theorem 2.6 First-order Reflection $\varphi(x_1,\ldots,x_n)$ is a first-order formula.

(i) For every set M_0 there exists M such that $M_0 \subset M$ and the following holds:

$$\varphi^M(x_1,\ldots,x_n) \leftrightarrow \varphi(x_1,\ldots,x_n)$$
 (2.13)

for every x_1, \ldots, x_n .

(ii) For every set M_0 there is a transitive set M, $M_0 \subset M$ such that the following holds:

$$\varphi^M(x_1,\ldots,x_n) \leftrightarrow \varphi(x_1,\ldots,x_n)$$
 (2.14)

for every x_1, \ldots, x_n .

⁹It can not be smaller because $|M'_{i+1}| \ge |M'_i|$ for every i. It may not be significantly larger because the maximum of elements added is the number of n-tuples in M'_i , which is of the same cardinality is M'_i . ((proc? Ramsey?))

316 (iii) For every set M_0 there is α such that $M_0 \subset V_\alpha$ and the following holds:

$$\varphi^{V_{\alpha}}(x_1, \dots, x_n) \leftrightarrow \varphi(x_1, \dots, x_n)$$
 (2.15)

for every x_1, \ldots, x_n .

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(iv) Assuming the Axiom of Choice, for every set M_0 there is M such that $M_0 \subset M$ and $|M| \leq |M_0| \cdot \aleph_0$ and the following holds:

$$\varphi^M(x_1,\ldots,x_n) \leftrightarrow \varphi(x_1,\ldots,x_n)$$
 (2.16)

for every x_1, \ldots, x_n .)

TODO mozna vyhodit (iv) z vety, abychom mohli rict "the following are equivalent"

Proof. Let's prove (i) for one formula φ via induction by complexity first. We can safely assume that φ contains no quantifiers besides \exists and no logical connectives other than \neg and &. Assume that this M is obtained from lemma 2.5.

The fact, that atomic formulas are reflected in every M comes directly from definition of relativization and the fact that they contain no quantifiers.¹⁰

The same holds for formulas in the form of $\varphi = \neg \varphi'$. Let us recall the definition of relativization for those formulas in .

$$(\neg \varphi_1)^M \leftrightarrow \neg (\varphi_1^M) \tag{2.17}$$

Because we can assume from induction that $\varphi'^M \leftrightarrow \varphi'$, the following holds:

$$(\neg \varphi')^M \leftrightarrow \neg (\varphi'^M) \leftrightarrow \neg \varphi' \tag{2.18}$$

The same holds for $\varphi = \varphi_1 \& \varphi_2$. From the induction hypothesis we know that $\varphi_1^M \leftrightarrow \varphi_1$ and $\varphi_2^M \leftrightarrow \varphi_2$, which together with relativization for formulas of form $\varphi_1 \& \varphi_2$ gives up

$$(\varphi_1 \& \varphi_2)^M \leftrightarrow \varphi_1^M \& \varphi_2^M \leftrightarrow \varphi_1 \& \varphi_2$$
 (2.19)

Let's now examine the case when from the induction hypethesis, M reflects $\varphi'(u_1, \ldots, u_n, x)$ and we are interested in $\varphi = \exists x \varphi'(u_1, \ldots, u_n, x)$. The
induction hypothesis tells us, that

$$\varphi'^{M}(u_1,\ldots,u_n,x) \leftrightarrow \varphi'(u_1,\ldots,u_n,x)$$
 (2.20)

¹⁰Note that this does not hold generally for relativizations to M, E, but only for relativization to M, \in , which is our case.

Together with above lemma 2.5 the following holds:

$$\exists x \varphi'(u_1, \dots, u_n, x) \tag{2.21}$$

$$\leftrightarrow (\exists x \in M)\varphi'(u_1, \dots, u_n, x) \tag{2.22}$$

$$\leftrightarrow (\exists x \in M) \varphi'^{M}(u_1, \dots, u_n, x) \tag{2.23}$$

$$\leftrightarrow (\exists x \varphi'(u_1, \dots, u_n, x))^M \tag{2.24}$$

343 Which is what we needed to prove.

So far we have proven part (i) of this theorem for one formula φ , we only need to verify that the same holds for any finite number of formulas. This has in fact been already done sice lemma 2.5 gives us M for any (finite) amount of formulas. We can than use the induction above to verify that it reflects each of the formulas individually.

Now we want to verify other parts of our theorem. Since V_{α} is a transitive set, by proving (iii) we also satisfy (ii). To do so, we only need to look at part (ii) of lemma 2.5

TODO doplnit lemma

Theorem 2.7 (Refl) is equivalent to (Infinity) & (Replacement) under ZFC minus (Infinity) & (Replacement)

Proof. Since (Refl) is a sound theorem in ZFC, we are only interested in showing the converse:

 $(Refl) \rightarrow (Infinity)$

This is the easy part since Infinity says that there is an infinite set and (Refl) is just a stronger version that says "there is an inaccessible cardinal" which is all we need. ((PREFORMULOVAT))

$$(Refl) \rightarrow (Replacement)$$

TODO Plagiat – prepsat a vysvetlit

Definition 2.8 Let $\varphi(R)$ be a Π_m^n -formula which contains only one free variable R which is second-order. Given $R \subseteq V_{\kappa}$, we say that $\varphi(R)$ reflects in V_{κ} if there is some $\alpha < \kappa$ such that:

If
$$(V_{\kappa}, \in, R) \models \varphi(R)$$
, then $(V_{\alpha}, \in, R \cap V_{\alpha}) \models \varphi(R \cap V_{\alpha})$. (2.25)

3 Large Cardinals

TODO pozor na opsane definice, prefurmulovvat!!!

370 3.1 Preliminaries

- To avoid confusion¹¹, let's first define some basic terms.
- Definition 3.1 (weak limit cardinal) kappa is a weak limit cardinal if it is \aleph_{α} for some limit α .
- Definition 3.2 (strong limit cardinal) kappa is a strong limit cardinal if for every $\lambda < \kappa$, $2^{\lambda} < \kappa$

3.2 Inaccelssibility

- Definition 3.3 (weak inaccessibility) κ is weakly inaccessible \leftrightarrow it is regular and weakly limit.
- Definition 3.4 (inaccessibility) κ is inaccessible \leftrightarrow it is regular and strongly limit.
- Theorem 3.5 [Lévy] The following are equivalent:
- (i) κ is inaccessible.
- 383 (ii) For every $R \subseteq V_{\kappa}$ and every first-order formula $\varphi(R)$, $\varphi(R)$ reflects in V_{κ} .
- 385 (iii) For every $R \subseteq V_{\kappa}$, the set $C = \{ \alpha < \kappa \mid \langle V_{\alpha}, \in, R \cap V_{\alpha} \rangle \prec \langle V_{\kappa}, \in, R \rangle \}$ is closed unbounded.
- Proof. Let's start with (i) \rightarrow (iii) in a way similar to [1].
- The set $\{\alpha < \kappa \mid \langle V_{\alpha}, \in, R \cap V_{\alpha} \rangle \prec \langle V_{\kappa}, \in, R \rangle \}$ is clearly closed, it remains to
- show that it is also unbounded. To do so, let $\alpha < \kappa$ be arbitrary. Define
- $\alpha_n < \kappa \text{ for } n \in \omega \text{ by recursion as follows:}$
- Set $\alpha_0 = \alpha$. Given $\alpha_n < \kappa$ define α_{n+1} to be the least $\beta \ge \alpha_n$ such as when-
- ever $y_1, \ldots, y_k \in V_{\alpha_n}$ and $\langle V_{\kappa}, \in, R \rangle \models \exists v_0 \varphi[v_0, y_1, \ldots, y_k]$ for some formula
- φ , there is an $x \in V_{\beta}$ such that $\langle V_{\kappa}, \in, R \rangle \models \varphi[x, y_1, \dots, y_k]$.
- Since κ is inaccessible, $|V_{\alpha_n}| < \kappa$ and so $\alpha_{n+1} < \kappa$.
- Finally, set $\alpha = \sup(\alpha_n | n \in \omega)$. Then $\langle V_\alpha, \in, R \cap V_\alpha \rangle \prec \langle V_\kappa, \in, R \rangle$ by the
- usual (Tarski) criterion for elementary substructure.

¹¹While in most sources refer to weak limit cardinal as a limit cardinal and to strong limit cardinal, in some cases the distinction is weak limit cardinal and limit cardinal respectively. That's why I have decided to explicitly define those otherwise elementary terms.

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     The next part, proving (iii) \rightarrow (ii), should be elementary since C is closed
398
    unbounded, which means that it contains at least countably many elements
399
     but we need only one such \alpha to satisfy (2.8).
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     Finally, we shall prove that (ii) \to (i). Since it obviously holds that \kappa > \omega,
401
     we have yet to prove that \kappa is regular and a strong limit. Let's argue by
402
     contradiction that it is regular. If it wasn't, there would be a \beta < \kappa and a
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    function F:\beta \implies \kappa with range unbounded in \kappa. Set R=\{\beta\}\cup F. By
404
    hypothesis there is an \alpha < \kappa such that \langle V_{\alpha}, \in, R \cap V_{\alpha} \rangle \prec \langle V_{\kappa}, \in, R \rangle. Since
405
     \beta is the single ordinal in R, \beta \in V_{\alpha} by elementarity. This yields the desired
406
     contradiction since the domain if F \cap V_{\alpha} cannot be all of \beta.
407
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    Next, let's see whether \kappa is indeed a strong limit, again by contradiction.
409
     If not, there would be a \lambda < \kappa such that 2^{\lambda} > \kappa. Let G : \mathcal{P}(\lambda) \implies \kappa be
410
     surjective and set R = \{\lambda + 1\} \cup G. By hypothesis, there is an \alpha < \kappa such
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     that \langle V_{\alpha}, \in, R \cap V_{\alpha} \rangle \prec \langle V_{\kappa}, \in, R \rangle. \lambda + 1 \in V_{\alpha} and so \mathscr{P}(\lambda) \in V_{\alpha}, but this is
412
     again a contradiction.
413
     3.3
              Mahlo cardinals
```

- TODO reflektuji nedosazitelnost? TODO zminit Mahlovu konstrukci? 415
- **Definition 3.6** Weakly Mahlo Cardinals κ is weakly Mahlo \leftrightarrow it is a limit ordinal and the set of all regular ordinals less then κ is stationary in κ
- **Definition 3.7** Mahlo cardinals The following definitions are equivalent: 418
 - (i) κ is Mahlo

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- (ii) κ is weakly Mahlo and strong limit 420
- (iii) κ is inaccessible and the regular cardinals below κ form a stationary 421 subset of κ . 422
- (iv) κ is regular and the stationary sets below κ form a stationary subset of 423 κ . 424
- **Theorem 3.8** κ is Mahlo \leftrightarrow for any $R \subset V_{\kappa}$ there is an inaccessible cardinal 425 $\alpha < \kappa \text{ such that } \langle V_{\alpha}, \in, R \cap V_{\alpha} \rangle \prec \langle V_{\kappa}, \in, R \rangle.$ 426
- *Proof.* Start with the proof of (3.5) and add the following: 427
- κ is Mahlo by the following contradiction. If not, there would be a C closed 428
- unbounded in κ containing no inaccessible cardinals. By the hypothesis there 429
- is in inaccessible $\alpha < \kappa$ such that $\langle V_{\alpha}, \in, C \cap V_{\alpha} \rangle \prec \langle V_{\kappa}, \in, C \rangle$. By elemen-430
- tarity $C \cap \alpha$ is unbounded in α . But then, $\alpha \in C$, which is the contradiction 431
- we need. 432

Note that Mahlo cardinals were first described in 1911, almost 50 years before Lévy's reflection, which was heavily inspired by those.

3.4 Weakly Compact Cardinals

- 436 TODO souvislost s reflexi!
- In this section, we will introduce various well-known large cardinals and establish them via reflection.
- Definition 3.9 A cardinal κ is weakly compact if it is uncountable and satisfies the partition property $\kappa \to (\kappa)^2$
- 441 Lemma 3.10 Every weakly compact cardinal is inaccessible
- Proof. Let κ b a weakly compact cardinal. To show that κ is regular, let us assume that κ i the disjoint union $\bigcup \{A_{\gamma}: \gamma < \lambda\}$ such that $\lambda < \kappa$ and $|A_{\gamma}| < \kappa$ for each $\gamma < \lambda$. We define a partition $F: [\kappa]^2 \to \{0,1\}$ as follows: $F(\{\alpha,\beta\}) = 0$ just in cas α and β are the same size A_{γ} . Obviously, this partition does not have a homogenous set $H \subset \kappa$ of size κ . That κ is a strong limit cardinal follows from Lemma 9.4: (?? doplnit z jecha): If $\kappa \geq 2^{\lambda}$ for some $\lambda < \kappa$, the because $2^{\lambda} \leq (\lambda^{+})^{2}$, we have $\kappa \leq (\lambda^{+})^{2}$ and hence $\kappa \leq (\kappa)^{2}$.
- Theorem 3.11 Let κ be a weakly compact cardinal. Then for every stationary set $S \subset \kappa$ there is an uncountable regular cardinal $\lambda < \kappa$ such that the set $S \cap \lambda$ is stationary in λ .
- 453 Proof. TODO □

454 3.5 Indescribable Cardinals

- Definition 3.12 (Indescribability) For Q either Π_n^m or Σ_n^m
- 456 A cardinal κ is Q-indescribable if whenever $U \subseteq V_{\kappa}$ and φ is a Q sentence
- such that $\langle V_{\kappa}, \in, U \rangle \models \varphi$, then for some $\alpha < \kappa$, $\langle V_{\alpha}, \in, U \cap V_{\alpha} \rangle \models \varphi$.

3.6 Bernays-Gödel Set Theory

TODO Plagiat – prepsat a vysvetlit

Gödel–Bernays set theory, also known as Von Neumann–Bernays–Gödel set theory is an axiomatic set theory that explicitly talks about proper classes as well as sets, which allows it to be finitely axiomatizable, albeit our version stated below contains one schema. It is a conservative extension of Zermalo–Fraenkel set theory. Using forcing, one can prove equiconsistency of BGC and ZFC.

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Bernays–Gödel set theory contains two types of objects: proper classes and sets. The notion of set, usually denoted by a lower case letter, is identical to set in ZF, whereas proper classes are usually denoted by upper case letters. The difference between the two is in a fact, that proper classes are not members of other classes, sets, on the other hand, have to be members of classes.

Definition 3.13 (Gödel–Bernay set theory)

(i) extensionality for sets

$$\forall a \forall b [\forall x (x \in a \leftrightarrow x \in b) \to a = b] \tag{3.26}$$

(ii) pairing for sets

$$\forall x \forall y \exists z \forall w [w \in z \leftrightarrow (w = x \lor w = y)] \tag{3.27}$$

476 (iii) union for sets

$$\forall a \exists b \forall c [c \in b \leftrightarrow \exists d (c \in d \land d \in a)]$$
 (3.28)

(iv) powers for sets

$$\forall a \exists p \forall b [b \in p \leftrightarrow (c \in b \to c \in a)] \tag{3.29}$$

(v) infinity for sets

There is an inductive set.
$$(3.30)$$

(vi) Extensionality for classes

$$\forall x (x \in A \leftrightarrow x \in B) \to A = B \tag{3.31}$$

(vii) Foundation for classes

Each nonempty class is disjoint from each of its elements. (3.32)

(viii) Limitation of size for sets

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For any class
$$C$$
 a set x such that $x=C$ exists iff (3.33)

there is no bijection between C and the class V of all sets (3.34)

(ix) Comprehension schema for classes

For any formula φ with no quantifiers over classes, there is a class A such that $\forall x (x \in A \cdot (3.35))$

- The first five axioms are identical to axioms in ZF.
- Comprehension schema tells us, that proper classes are basically first-order
- 486 predicates. ... TODO Plagiat prepsat a vysvetlit
- Definition 3.14 We say that $\varphi(R)$ with a class parameter R reflects if there is α such that

$$\varphi(R) \to (V_{\alpha}, V_{\alpha+1}) \models \varphi(R \cap V_{\alpha}).$$
 (3.36)

Theorem 3.15 There is a second-order sentence φ which is provable in GB such that if φ reflects at α , i.e. if

$$\varphi \to (V_{\alpha}, V_{\alpha+1}) \models \varphi,$$
 (3.37)

- 491 then α is an inaccessible cardinal.
- 492 *Proof.* Take φ to say "there is no function from $\gamma \in ORD$ cofinal in ORD
- and for every $\gamma \in ORD$, $2^{\gamma} \in ORD$ ". Clearly, if φ reflects at some α ,
- then α is inaccessible (here we use that the second-order variable range over

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$$\mathscr{P}(V_{\alpha}) = V_{\alpha+1}$$
.

- As a corollary we obtain:
- Corollary 3.16 Second-order reflection in GB implies the existence of an inaccessible cardinal.

499 3.7 Morse–Kelley Set Theory

500 Axioms not

(i) Extensionality

$$\forall X \forall Y (\forall z (z \in X \leftrightarrow z \in Y) \to X = Y). \tag{3.38}$$

502 (ii) Pairing

$$asdfg$$
 (3.39)

503 (iii) Foundation For Classes

$$asdf$$
 (3.40)

504 (iv) Class Comprehension

$$\forall W_1, \dots, W_n \exists Y \forall x (x \in Y \leftrightarrow (\phi(x, W_1, \dots, W_n) \& set(x))). \tag{3.41}$$

Where set(x) is monadic predicate stating that class x is a set.

506 (v) Limitation Of Size For Classes

$$asdf$$
 (3.42)

507 (vi) Pairing

$$asdf$$
 (3.43)

508 (vii) Pairing

$$asdf$$
 (3.44)

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3.8 Reflection and the constructible universe

TODO reflektovat muzeme jenom kardinaly konzistentni s V=L, proc?

TODO Plagiat – prepsat a vysvetlit

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L was introduced by Kurt Gödel in 1938 in his paper The Consistency of the Axiom of Choice and of the Generalised Continuum Hypothesis and denotes a class of sets built recursively in terms of simpler sets, somewhat similar to Von Neumann universe V. Assertion of their equality, V=L, is called the axiom of constructibility. The axiom implies GCH and therefore also AC and contradicts the existence of some of the large cardinals, our goal is to decide whether those introduced earlier are among them.

On order to formally establish this class, we need to formalize the notion of definability first:

Definition 3.17 (Definable sets)

$$Def(X) := \{ \{ y | x \in X \land \langle X, \in \rangle \models \varphi(y, z_1, \dots, z_n) \} | \varphi \text{ is a first-order formula, } z_1, \dots, z_n \in X \}$$

$$(3.45)$$

Now we can recursively build L.

Definition 3.18 (The Constructible universe) (i)

$$L_0 := \emptyset \tag{3.46}$$

$$(ii) L_{\alpha+1} := Def(L_{\alpha}) (3.47)$$

(iii)
$$L_{\lambda} = \bigcup_{\alpha < \lambda} L_{\alpha} \text{ If } \lambda \text{ is a limit ordinal}$$
 (3.48)

$$L = \bigcup_{\alpha \in Ord} L_{\alpha} \tag{3.49}$$

TODO Plagiat – prepsat a vysvetlit

Fact 3.19 The reflection – constructed as explained in the previous paragraph (!!! preformulovat !!!) – with second-order parameters for higher-order formulas (even of transfinite type) does not yield transcendence over L.

TODO zduvodneni

TODO kratka diskuse jestli refl implikuje transcendenci na L - polemika, nazor - V=L a slaba kompaktnost a dalsi

⁵³⁰ 4 Higher-order reflection

TODO rict ze to je zobecneni a nejaky dalsi uvodni veci

532 **4.1** Sharp

TODO je tohle higher-order vec?

⁵³⁴ 4.2 Welek: Global Reflection Principles

535 TODO

5 Conclusion

TODO asi tady neco bude, ne?

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