WEEK-2 PRESENTED BY NOLLEH

LEARN U ERLANG

greet without patten matching

```
function greet(Gender,Name)

if Gender == male then

print("Hello, Mr. %s!", Name)

else if Gender == female then

print("Hello, Mrs. %s!", Name)

else

print("Hello, %s!", Name)

end
```

greet with patten matching

```
greet(male, Name) ->
io:format("Hello, Mr. ~s!", [Name]);
greet(female, Name) ->
io:format("Hello, Mrs. ~s!", [Name]);
greet(_, Name) ->
io:format("Hello, ~s!", [Name]).
```

disusing / using

```
function(Args)
if X then
Expression
else if Y then
Expression
else
Expression
```

```
function(X) ->
   Expression;
function(Y) ->
   Expression;
function(_) ->
   Expression.
```

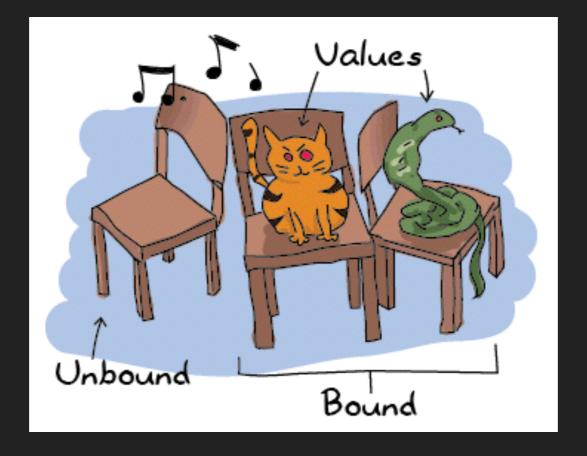
example

```
head([H|_]) -> H.
second([_,X|_]) -> X.
same(X,X) ->
true;
same(_,_) ->
false.
```

how's it works?

error occurs unless the new value is the same as the old

one



example 2

functions head's '=' operator

example 2

```
4> c(functions).
   {ok, functions}
5> functions:valid_time({{2011,09,06},{09,04,43}}).
   The Date tuple ({2011,9,6}) says today is: 2011/9/6,
   The time tuple ({9,4,43}) indicates: 9:4:43.
   ok
6> functions:valid_time({{2011,09,06},{09,04}}).
   Stop feeding me wrong data!
```

prob: It also recv just tuple! and too precise sometimes.

- needs expressive way on sometimes...
 - range of value
 - not limited as certain types of data

impractical vs practical

```
old_enough(0) -> false;
old_enough(1) -> false;
old_enough(2) -> false;
...
old_enough(14) -> false;
old_enough(15) -> false;
old_enough(_) -> true.
ok
```



```
old_enough(X) when X >= 16 -> true;
old_enough(_) -> false.
```

simillar with andalso (little diff aspect on exceptions)

```
right_age(X) when X >= 16, X =< 104 -> true;
right_age(_) -> false.
```



orelse

```
wrong_age(X) when X < 16; X > 104 ->
true;
wrong_age(_) ->
false.
```

- in guard, You will be
 - able to use functions like
 - A*B/C >= 0
 - is_integer/1, is_atom/1 ...
 - unable to use user defined function
 - because of side-effect

- similiar with guard but outside of function clauses head
- different from other language

```
-module(what_the_if).
-export([heh_fine/0]).
heh_fine() ->
 if 1 =:= 1 ->
  works
 end,
 if 1 =:= 2; 1 =:= 1 ->
   works
 end,
 if 1 =:= 2, 1 =:= 1 ->
   fails
 end.
```



```
1> c(what_the_if).
 ./what_the_if.erl:12: Warning: no clause will ever match
 ./what_the_if.erl:12: Warning: the guard for this clause evaluates to
 'false'
 {ok,what_the_if}
2> what_the_if:heh_fine().
 ** exception error: no true branch found when evaluating an if
 expression
 in function what_the_if:heh_fine/0
```

true branch

```
oh_god(N) ->
if N =:= 2 -> might_succeed;
true -> always_does %% this is Erlang's if's 'else!'
end.
```

```
4> what_the_if:oh_god(2).
might_succeed
5> what_the_if:oh_god(3).
always_does
```

WHAT IS IF!?

- why not else?
 - both branch should be avoided
 - if is usually easier
- guard has only limited set

IN CASE ... OF

example

```
insert(X,[]) ->
 [X];
insert(X,Set) ->
  case lists:member(X,Set) of
   true -> Set;
  false -> [X|Set]
  end.
```

IN CASE ... OF

pattern matching + guard

```
beach(Temperature) ->
 case Temperature of
   \{celsius, N\} \text{ when } N >= 20, N = < 45 ->
     'favorable';
   \{\text{kelvin, N}\}\ \text{when N} >= 293, N =< 318 ->
     'scientifically favorable';
   \{fahrenheit, N\}\ when N >= 68, N =< 113 ->
     'favorable in the US';
     'avoid beach'
 end.
```

IN CASE ... OF

instead, we can replace with bunch of functions.

```
beachf(\{\text{celsius}, N\}) when N \ge 20, N = < 45 \implies
 'favorable';
beachf(_) ->
  'avoid beach'.
```

WHICH TO USE?

- function call vs case of
 - same way at a lower level
 - only one difference when arg is more than one

```
case {A,B} of
Pattern Guards -> ...
end.
```

WHICH TO USE?

- function call vs case of
 - arguably cleaner

```
insert(X,[]) ->
  [X];
insert(X,Set) ->
  case lists:member(X,Set) of
   true -> Set;
  false -> [X|Set]
  end.
```

WHICH TO USE?

- if vs if through guard?
 - use where doesn't need whole pattern matching
- personal preference

DYNAMITE-STRONG TYPING

- as you've seen, no need to type Type!
- elang is dynamically typed
 - compiler won't always yell at you
- statically typed language is safer..?
 - elang is reported as nine nine (99.999 % available)
- strongly typed

```
1>6+"1".
```

DYNAMITE-STRONG TYPING

type conversion

<type>_to_<type>

```
1> erlang:list_to_integer("54").
54
2> erlang:integer_to_list(54).
"54"
3> erlang:list_to_integer("54.32").
** exception error: bad argument in function list_to_integer/1 called as list_to_integer("54.32")
```

TYPE CONVERSION

type conversion

<type>_to_<type>

```
4> erlang:list_to_float("54.32").
54.32
5> erlang:atom_to_list(true).
"true"
6> erlang:list_to_bitstring("hi there").
<<"hi there">>
7> erlang:bitstring_to_list(<<"hi there">>).
"hi there"
```

TO GUARD A DATA TYPE

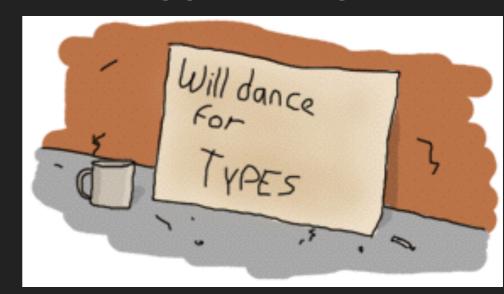
type check

```
is_atom/1 is_binary/1
is_bitstring/1 is_boolean/1 is_builtin/3
is_float/1 is_function/1 is_function/2
is_integer/1 is_list/1 is_number/1
is_pid/1 is_port/1 is_record/2
is_record/3 is_reference/1 is_tuple/1
```

- why not typeof
 - force user to make program that surely knowing the effect
- can used in guard expression

FOR TYPE JUNKIES

- briefly describe tools used to do static type analysis
 - first try in 1997



- success type
 - will not exact type every expression
 - type it infers are right, type errors it finds are really error

FOR TYPE JUNKIES

success type

```
and(false, _) -> false;
and(_, false) -> false;
and(true,true) -> true.
```

and(_,_) -> bool()

FOR TYPE JUNKIES

if you interested

\$ typer --help

\$ dialyzer --help

HELLO RECURSION!

functional programming do not offer loop



HELLO RECURSION!

factorial

```
-module(recursive).  n! = \begin{cases} 1 & \text{if } n = 0 \\ n((n-1)!) & \text{if } n > 0 \end{cases}  fac(N) when N == 0 -> 1;  fac(N) \text{ when N} > 0 \text{ -> N*fac(N-1)}.
```

```
fac(0) -> 1;
fac(N) when N > 0 -> N*fac(N-1).
```

HELLO RECURSION!

- recursion
 - function that calls itself
 - need to have stopping condition (base case)

LENGTH

we need

```
a base case;
a function that calls itself;
a list to try our function on.
```

simplest - empty list

```
fac(0) -> 1; fac(N) when N > 0 -> N*fac(N-1).
```

LENGTH

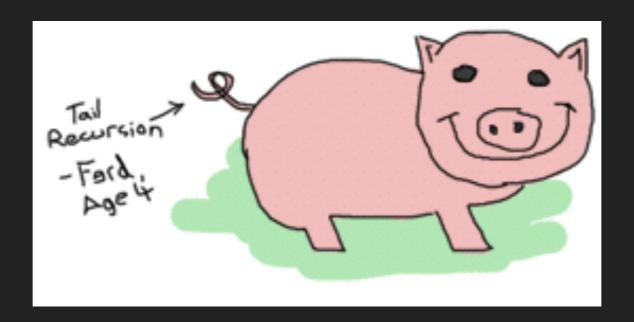
list is recursively

```
len([]) -> 0;
len([_|T]) -> 1 + len(T).
```

LENGTH

how it works?

- it is problematic
 - keep millions of numbers in memory for such a simple calculation.
- let's tail recursion



- linear -> iterative one
- need to be alone
- additional is stacked
- accumulator
 - need to extra variable to hold the intermediate result

using accumulator

```
tail_fac(N) -> tail_fac(N,1).

tail_fac(0,Acc) -> Acc;

tail_fac(N,Acc) when N > 0 -> tail_fac(N-1,N*Acc).
```

```
tail_fac(4) = tail_fac(4,1)

tail_fac(4,1) = tail_fac(4-1, 4*1)

tail_fac(3,4) = tail_fac(3-1, 3*4)

tail_fac(2,12) = tail_fac(2-1, 2*12)

tail_fac(1,24) = tail_fac(1-1, 1*24)

tail_fac(0,24) = 24
```

length tail recursion

```
len([]) -> 0;
len([_|T]) -> 1 + len(T).
```

```
tail_len(L) -> tail_len(L,0).
tail_len([], Acc) -> Acc;
tail_len([_|T], Acc) -> tail_len(T,Acc+1).
```

- After all, recursion being the only looping construct that exists in Erlang
 - except list comprehension
- duplicate

```
duplicate(0,_) ->
[];
duplicate(N,Term) when N > 0 ->
[Term|duplicate(N-1,Term)].
```

duplicate

true nightmare which is not tail recursion

```
reverse([]) -> [];
reverse([H|T]) -> reverse(T)++[H].
```

```
reverse([1,2,3,4]) = [4]++[3]++[2]++[1]

\uparrow \quad \leftarrow \downarrow
= [4,3]++[2]++[1]
\uparrow \uparrow \quad \leftarrow \downarrow
= [4,3,2]++[1]
\uparrow \uparrow \uparrow \quad \leftarrow \downarrow
= [4,3,2,1]
```

let's rescue

```
tail_reverse(L) -> tail_reverse(L,[]).
tail_reverse([],Acc) -> Acc;
tail_reverse([H|T],Acc) -> tail_reverse(T, [H|Acc]).
```

- sublist/2
 - a little different

```
sublist(_,0) -> [];
sublist([],_) -> [];
sublist([H|T],N) when N > 0 -> [H|sublist(T,N-1)].
```

```
tail_sublist(L, N) -> tail_sublist(L, N, []).

tail_sublist(_, 0, SubList) -> SubList;

tail_sublist([], _, SubList) -> SubList;

tail_sublist([H|T], N, SubList) when N > 0 ->

tail_sublist(T, N-1, [H|SubList]).
```

problems...

```
sublist([1,2,3,4,5,6],3)
```

solve

tail_sublist(L, N) -> reverse(tail_sublist(L, N, [])).

zip/2

```
1> recursive:zip([a,b,c],[1,2,3]).
[{a,1},{b,2},{c,3}]
```

implements

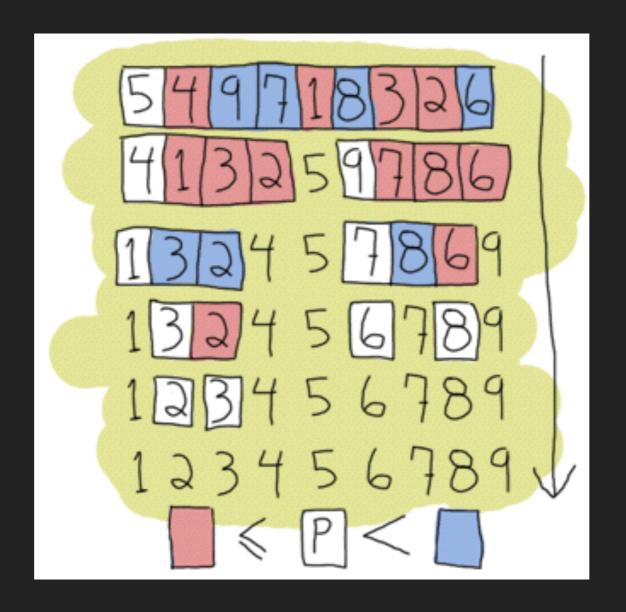
```
zip([],[]) -> [];
zip([X|Xs],[Y|Ys]) -> [{X,Y}|zip(Xs,Ys)].
```

```
lenient_zip([],_) -> [];
lenient_zip(_,[]) -> [];
lenient_zip([X|Xs],[Y|Ys]) -> [{X,Y}|lenient_zip(Xs,Ys)].
```

- TCO (tail call optimization)
 - vm does eliminate current stack frame
- LCO (last call optimization)
 - more general

- naive version
 - pivot
 - smallers; equals | lagers
 - until empty list to sort

naive version



- as two parts
 - partioning
 - apply the partitioning to each parts, and glue

```
quicksort([]) -> [];
quicksort([Pivot|Rest]) ->
{Smaller, Larger} = partition(Pivot,Rest,[],[]),
quicksort(Smaller) ++ [Pivot] ++ quicksort(Larger).
```

partitioning

```
partition(_,[], Smaller, Larger) -> {Smaller, Larger};
partition(Pivot, [H|T], Smaller, Larger) ->
  if H =< Pivot -> partition(Pivot, T, [H|Smaller], Larger);
  H > Pivot -> partition(Pivot, T, Smaller, [H|Larger])
  end.
```

- tree
 - key / two other node (smaller, larger)
 - also able to contain empty node



lets choice tuple!

```
{node, {Key, Value, Smaller, Larger}}
{node, nil}
```

empty

```
-module(tree).
-export([empty/0, insert/3, lookup/2]).
empty() -> {node, 'nil'}.
```

- base case is empty node
 - where to put content
- compare, larger / smaller

compare, larger / smaller

```
insert(Key, Val, {node, 'nil'}) ->
    {node, {Key, Val, {node, 'nil'}, {node, 'nil'}};
insert(NewKey, NewVal, {node, {Key, Val, Smaller, Larger}})
    when NewKey < Key ->
        {node, {Key, Val, insert(NewKey, NewVal, Smaller), Larger}};
insert(NewKey, NewVal, {node, {Key, Val, Smaller, Larger}})
    when NewKey > Key ->
        {node, {Key, Val, Smaller, insert(NewKey, NewVal, Larger)}};
insert(Key, Val, {node, {Key, _, Smaller, Larger}}) ->
    {node, {Key, Val, Smaller, Larger}}.
```

- returns completely new tree
 - sometimes shared by vm
- look up

```
lookup(_, {node, 'nil'}) ->
undefined;
lookup(Key, {node, {Key, Val, _, _}}) ->
{ok, Val};
lookup(Key, {node, {NodeKey, _, Smaller, _}}) when Key < NodeKey ->
lookup(Key, Smaller);
lookup(Key, {node, {_, _, _, Larger}}) -> lookup(Key, Larger).
```

using example

```
    1> T1 = tree:insert("Jim Woodland", "jim.woodland@gmail.com", tree:empty()).
    2> T2 = tree:insert("Mark Anderson", "i.am.a@hotmail.com", T1).
    3> Addresses = tree:insert("Anita Bath", "abath@someuni.edu", tree:insert("Kevin Robert", "myfairy@yahoo.com", tree:insert("Wilson Longbrow", "longwil@gmail.com", T2))).
```

```
{node,{"Jim Woodland","jim.woodland@gmail.com",
   {node,{"Anita Bath","abath@someuni.edu",
       {node,nil},
       {node,nil}}},
   {node,{"Mark Anderson","i.am.a@hotmail.com",
       {node,{"Kevin Robert","myfairy@yahoo.com",
          {node,nil},
          {node,nil}}},
       {node,{"Wilson Longbrow","longwil@gmail.com",
          {node,nil},
          {node,nil}}}}}}
```

using example

```
4> tree:lookup("Anita Bath", Addresses).
{ok, "abath@someuni.edu"}
5> tree:lookup("Jacques Requin", Addresses).
undefined
```

THINKING RECURSIVELY

- our approach is more declarative
- consise algorithm easy to understand
 - divide and conquer!
- you will learn how to abstract, next time.