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COMPUTER Calisthenics & Orthodontia

Running Light Without Overbyte

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Volume 1, Number 1

A REFERENCE JOURNAL FOR USERS OF HOME COMPUTERS

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THE ORIGINAL

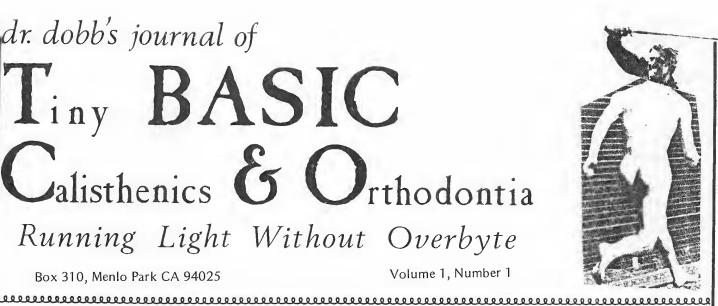
VOLUME 1, NUMBER 1

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Tiny BASIC Salisthenics & Orthodontia Running Light Without Overbyte

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STATUS LETTER

by Dennis Allison

The magic of a good language is the ease with which a particular idea may be expressed. The assembly language of most microcomputers is very complex, very powerful, and very hard to learn. The Tiny BASIC project at PCC represents our attempt to give the hobbyist a more human-oriented language or notation with which to encode his programs. This is done at some cost in space and/or time. As memory still is relatively expensive, we have chosen to trade features for space (and time for space) where we could.

Our own implementation of Tiny BASIC has been very slow. I have provided technical direction only on a sporadic basis. The real work has been done by a number of volunteers; Bernard Greening has left the project. As might be guessed, Tiny BASIC is a tiny part of what we do regularly. (And volunteer labor is not the way to run a software project with any kind of deadline!)

While we've been slow, several others have really been fast. In this issue we publish a version of Tiny BASIC done by Dick Whipple and John Arnold in Tyler, Texas. (And other versions can't be far behind.)

MY, HOW TINY BASIC GROWED!

Once upon a time, in PCC, Tiny BASIC started out to be: † a BASIC-like language for tiny kids, to be used for games, recreations, and the stuff you find in elementary school math books.

† an exercise in getting people together to develop FREE software.

- † portable-machine independent.
- † open-ended--a toy for software tinkerers.
- t small.

Then . . . (fanfare!) . . . along came Dick Whipple and John Arnold. They built Tiny BASIC Extended. It works. See pp 13-17 and 19 in this issue for more information. More next issue.

WANTED: More Tiny BASICs up and running. WANTED: More articles for this newsletter.

WANTED: Tiny other languages. I might be able to live with Tiny FORTRAN but, I implore you, no Tiny COBOL! How about Tiny APL? Or Tiny PASCAL (whatever that is)?

WANTED: Entirely new, never before seen, Tiny Languages, imported from another planet or invented here on Earth. Especially languages for kids using home computers that talk to tvs or play music or run model trains or . . .

BASIC

BASIC, Beginners' All-purpose Symbolic Instruction Code, was initially developed in 1963 and 1964 by Professors John Kemeny and Thomas Kurtz of Dartmouth College, with partial Isupport from the National Science Foundation under the terms of Grant NSF GE 3864. For information on Dartmouth BASIC publications, get Publications List (TM 086) from Documents Clerk, Kiewit Computation Center, Dartmouth College, Hanover NH 03755. Telephone 603-646-2643.

Try these: TM028 BASIC: A Specification \$3.15 TM075 BASIC \$4.50

******** It would help a lot if you would each send us a 3x5 card with your name, address (including zip), telephone number, and a rather complete description of your hardware.

DRAGON THOUGHTS

† We promised three issues. After these are done, shall we continue?

t If we do, we will change the name and include languages other than BASIC.

† This newsletter is meant to be a sharing experience, intended to disseminate FREE software. It's OK to charge a few bucks for tape cassettes or paper tape or otherwise recover the cost of sharing. But please make documentation essentially free, including annotated source code.

† If we do continue, we will have to charge about \$1 per issue to recover our costs. In Xeroxed form, we can provide about 20-24 pages per issue of tiny eye-strain stuff. If we get big bunches of subscriptions, we'll print it and expand the number of pages, depending on the number of subscribers.

† So, let us know . . . shall we continue?

man and a second

For our new readers, and those who have been following articles on Tiny BASIC as they appeared in People's Computer Company, we have reprinted on pages 3-12 the best of Tiny BASIC from PCC as an introduction, and as a reference.

t = v + v; **TECHNIQUES & PRACNIQUES** u = t + t; Computes 10 - x u = u + u + tby Dennis Allison, 12/1/75 Byte only as high order must be u = n - uegual (This will be a continuing column of tricks, algorithms, and if $u \ge 10$ then Perhaps one could use a decimal other good stuff everyone needs when writing software. Contribufeature here do; tions solicited.) u = u - 10;Corrects for case where [n/10] - 1 16-BIT BINARY TO DECIMAL CONVERSION ROUTINE n = v + 1;is computed and creates [n/10] and end t saves characters on stack n mod 10 † performs zero suppressed conversion n = v; t uses multiplication by 0.1 to obtain n/10 and n mod 10 Saves result on stack call push (u); define crutch = OFFH: Loop at least once until n = 0; declare n, u, v, t; BIT (16) These could be registers, or on the if n < 0 then stack ch = pop;do; do while < > crutch; Write result in reverse order n = -n;call outch (ch + 030H); Converts digits to ASCII call outch('-'); 0 = 030H 02 = 032H etc. end: ch = pop; Pop takes one word off the stack call push (crutch) The crutch marks the end of numend ber on the stack repeat; † Letters from readers are most welcome. Unless they v = shr(n,1);note otherwise, we will assume we are free to publish v = v + shr(v, 1);These all are 16 bit shifts and share them. Computes [n/10] or [n/10] - 1 by v = v + shr(v,4);v = v + shr(v.8): multiplication † We hereby assign reprint rights to all who wish to use v = shr(v,3);Call it x Tiny BASIC Calisthenics & Orthodontia for noncommercial purposes. † To facilitate connection between our subscribers, we will in subsequent issues publish our subscriber list (including addresses and equipment of access/interest). PCC Tiny BASIC Reorganizes... 12-15-75 I want to subscribe to DR. DOBB'S JOURNAL OF Bob Albrect TINY BASIC CALISTHENICS & ORTHODONTIA . Dennis Allisa (3 issues for \$3) Tuny Basic feels like a dead Albitross around my neck. I do not feel like working on it any NAME ADDRESS_ more. STATE. (If you would like us to publish your name, address, and equipment of access/interest in future issue(s), please indicate VERY ... and so we procede somewhat more slowly than some SPECIFICALLY: of our readers geeeeeeeeeeeeeeeeeeeeeeeeee EOUIPMENT Dennis Allison - technical editor OF ACCESS/INTEREST_ Bob Albrecht contributing John Arnold editors Dick Whipple Lois Britton circulation manager Please send check or money order (purchase order minimum: \$6) Rhoda Horse midwife-at-large to TINY BASIC CALISTHENICS & ORTHODONTIA Box 310, Menlo Park CA 94025. Thank you.

BUILD YOUR OWN BASIC

by Dennis Allison & Others (reprinted from *People's Computer Company* Vol. 3, No.4)

A DO IT YOURSELF KIT FOR BASIC??

Yes, available from PCC with this newspaper and a lot of your time. This is the beginning of a series of articles in which we will work our way through the design and implementation of a reasonable BASIC system for your brand X computer. We'll be working on computers based on the INTEL 8008 and 8080 microprocessors. But it doesn't make much difference — if your machine is the ZORT 9901 or ACME X you can still build a BASIC for it. But remember, it's a hard job and will take lots of time particularly if you haven't done it before. A good BASIC system could easily take one man six months!

We'd like everyone interested to participate in the design. While we could do it all ourselves, (we have done it before) your ideas may be better than ours. Maybe we can save you, or you can save us, a lot of work or problems. Write us and we'll publish your letter and comments.

WHICH BASIC?

There is not any one standard BASIC (yet). The question is which BASIC should we choose to implement. A smaller (fewer statements, fewer features) BASIC is easier to implement and (more important) takes less space in the computer. Memory is still expensive so the smaller the better. Yet maybe we can't give up some goodies like string variables, dynamic array allocation, and so on.

There is a standard version of BASIC which is to be the minimal language which can be called BASIC. It's a pretty big language with lots of goodies. Maybe too big. Is there any advantage to being compatible with, say, the EDU BASICS? We don't have to make any decision yet; but the time will come . . .

COMPILER OR INTERPRETER?

We favor using an interpreter. An interpreter is a program which will execute the BASIC program from its textual representation. The program you write is the one which gets executed. A compiler converts the BASIC program into the machine code for the machine it is to run on. Compiled code is a lot faster, but requires more space and some kind of mass storage device (tape or disk). Interpretative BASIC is the most common on small machines.

HOW MUCH MEMORY? AND ... WHAT KIND?

Can we make some guesses about how big the BASIC system will be? Only if you don't hold us to it. Suppose we want to be able to run a 50 line BASIC program. We need about 800 bytes to store the program, another 60 or so bytes for storing program values (all numeric) without leaving any space for the interpreter and its special data. Past experience has shown that something like 6 to 8 Kbytes are needed for a minimum BASIC interpreter and that at least 12K bytes are necessary for a comfortable system. That's a lot of memory, but not too much more than you need to run the assembler. A lot of BASIC could be put into ROM (Read Only Memory) once developed and checked out. ROM is a lot cheaper than RAM (Read and Write) memory, but you can't change it. It's lots better to make sure everything works first.

DIRECT MODE?

Some kind of "desk calculator" mode of opera-

tion would be nice. At least, we would like to

This feature makes it easier to find and gently

terminate the existence of "bugs."

be able to look at and set different variables in a

program and restart execution at any given point.

But... if we can agree on some chunks of code and get it properly checked out, some enterprising person out there might make a few thousand ROMs and save us all some \$\$\$. Let's see now...how about ROMs for floating point arithmetic, integer arithmetic, Teletype I/O...

DATA STRUCTURES

Data structures are places to put things so you can find them or use them later. BASIC has at least three important ones: a symbol table which looks up a program name, A or Z9 or A\$, with its value. If we had a big computer where space was not a huge problem, we could simply preallocate all storage since BASIC provides for only 312 different names excluding arrays. When memory is so costly this doesn't make much sense. Somewhere, also, we've got to store the names which BASIC is going to need to know, names like LET and GO TO and IF. This table gets pretty big when there are lots of statements.

Lastly, we need some information about what is a legal BASIC statement and which error to report when it isn't. These tables are called parsing tables since they control the decomposition of the program into its component parts.

STRATEGY

Divide and Conquer is the programmers maxim. BASIC will consist of a lot of smaller pieces which communicate with each other. These pieces themselves consist of smaller pieces which themselves consist of smaller pieces, and so forth down to the actual code. A large problem is made manageable by cutting it into pieces.

What are the pieces, the building blocks of BASIC? We see a bunch of them:

- *a supervisor which determines what is to be done next. It receives control when BASIC is loaded.
- *a program and line editor. This program collects lines as they are entered from the keyboard and puts them into a part of computer memory for later use.
- *a line executor routine which executes a single BASIC statement, whatever that is.
- *a line sequence which determines which line is to be executed next.
- *a floating point package to provide floating point on a machine without the hardware.
- *terminal I/O handler to input and output information from the Teletype and provide simple editing (backspace and line deletion).
- *a function package to provide all the BASIC functions (RND, INT, TAB, etc.) *an error handling routine (part of the supervisor).
- *a memory management program which provides dynamic allocation data objects.

These are the major ones. As we get futher into the system we'll begin to see others and we'll begin to be able to more fully define the function of each of these modules.

TINY BASIC

Pretend you are 7 years old and don't care much about floating point arithmetic (what's that?), logarithms, sines, matrix inversion, nuclear reactor calculations and stuff like that.

And . . . your home computer is kinda small, not too much memory. Maybe its a MARK-8 or an ALTAIR 8800 with less than 4K bytes and a TV typewriter for input and output.

You would like to use it for homework, math recreations and games like NUMBER, STARS, TRAP, HURKLE, SNARK, BAGELS, . . .

Consider then, TINY BASIC

- Integer arithmetic only 8 bits? 16 bits?
- 26 variables: A, B, C, D, . . . , Z
- The RND function of course!
 Seven BASIC statement types
 - INPUT
 - PRINT
 - LET
 - GO TO
 - GOSUB
 - RETURN
- Strings? OK in PRINT statements, not OK otherwise.

BUILD YOUR OWN BASIC-REVIVED

(reprinted from People's Computer Company Vol. 4, No. 1)

WHAT IS TINY BASIC???

TINY BASIC is a very simplified form of BASIC which can be implemented easily on a microcomputer. Some of its features are:

Integer arithmetic 16 bits only

26 variables (A, B, . . . , Z)

Seven BASIC statements

INPUT PRINT LET GOTO IF GOSUB RETURN

Strings only in PRINT statements

Only 256 line programs (if you've got that much memory)

Only a few functions including RND

It's not really BASIC but it looks and acts a lot like it. I'll be good to play with on your ALTAIR or whatever; better, you can change it to match your requirements and needs.

TINY BASIC LIVES!!!

We are working on a version of TINY BASIC to run on the INTEL 8080. It will be an interpretive system designed to be as conservative of memory as possible. The interpreter will be programmed in assembly language, but we'll try to provide adequate descriptions of our intent to allow the same system to be programmed for most any other machine. The next issue of PCC will devote a number of pages to this project.

★ In the meantime, read one of these.
Compiler Construction For Digital Computers, David Gries, Wiley, 1971
493 pages, \$14.95

Theory & Application of a Bottom-Up Syntax Directed Translator Harvey Abramson, Academic Press, 1973, 160 pages, \$11.00

Compiling Techniques, F.R.A. Hopgood, American Elsevier, I26 pages \$6.50

A BASIC Language Interpreter for the Intel 8008 Microprocessor A.C. Weaver, M.H. Tindall, R.L. Danielson. University of Illinois Computer Science Dept, Urbana IL 61801. June 1974. Report No. UIUCDCS-R-74-658. Distributed by National Technical Information Service, U.S. Commerce Dept, Springfield VA 22151. \$4.25.

A BASIC language interpreter has been designed for use in a microprocessor environment. This report discusses the development of 1) an elaborate text editor and 2) a table-driven interpreter. The entire system, including text editor, interpreter, user text buffer, and full floating point arithmetic routines fits in 16K 8-hit words.

The TINY BASIC proposal for small home computers was of great interest to me. The lack of floating point arithmetic however, tends to limit its usefulness for my objectives.

As a matter of a suggestion, consideration should be given to the optional inclusion of floating point arithmetic, logarithm and trigonmetric calculation capability via a scientific calculator chip interface.†

The inclusion of such an option would tend to extend

the interpreter to users who desire these complex calculation capabilities. A number of calculator chip proposals have been made, with the Suding unit being of the most interest.

Thank you for the note of 13 June, regarding my letter on the Tiny BASIC article (PCC Vol. 3 No. 4). It was with regret that I learned that the series was not continued in the next volume. Even though few responded to the article published, conceptually the knowledge and principles which would be disseminated regarding a limited lexicon, high level programming language are of importance to the *independent* avocational microcomputer community.

At this time, PCC may not have a wide distribution in the avocation microcomputer community. This could be possibly the cause for the low number of respondies. Never the less, this should not detract from the dissemination and importance of concepts and principles which are of significance.

The thrust of my letter of 15 April, 1975, was to suggest a mechanism for the inclusion of F.P. in a limited lexicon and memory consumptive BASIC. I hope that the implication that F.P. must be included was not read into my letter.

It is my interest that information, concepts and the principles of compiler/interpreter construction as it related to microcomputers be available to the limited budget avocational user. The MITS BASIC, which you brought up, appears from my viewpoint to be a licensed, hlackbox program which is not currently available to:
(a) 8008 users, (b) IMP-16 users, (c) independent 8080 users (except at a very large expense) or (d) MC6800 users who will shortly be on line.

Presently it appears that microcomputer compiler interpretor function languages will be coming available from a number of sources (MITS, NITS, Processor Technology and etc.). However, few will probably deal in the conceptualizations which are the basis of the interpreter. Information which will fill the void in the interpreter construction knowledge held by the avocation builder, should be made available.

I strongly urge that the series started with Vol. 3
No. 4 article be continued. Possibly the hardware, peripheral, machine programming difficulties incurred by the microcomputer builder, is prohibiting a major contribution at this time. However, I would expect that by Autumn a number of builders should have their construction and peripheral difficulties far enough along to start thinking about higher level languages. The previous objective for the article series sounds reasonable. It was not my purpose in submitting the letter to detract from the objective of a very limited lexicon BASIC, ie., to be attractive and usable by the young and beginner due to its simplicity.

If wives, children, neighbors or anyone who is not machine language or programming oriented is expected to use a home-base unit created under a restrained budget a high level language will be a necessity. It is with this foresight that I encourage the continuance of the "Build Your Own BASIC" series.

This issue aside, 1 would like to encourage the PCC to continue the quite creditable activities which have been its order of business with regard to avocational computing.

Michael Christoffer 4139 12th NE No. 400

Seattle, Wash. 98105

† Please see Dr Robert Suding's article on p. 18

DESIGN NOTES FOR TINY BASIC

by Dennis Allison, happy Lady, & friends (reprinted from *People's Computer Company* Vol. 4, No. 2)

SOME MOTIVATIONS

A lot of people have just gotten into having their own computer. Often they don't know too much about software and particularly systems software, but would like to be able to program in something other than machine language. The TINY BASIC project is aimed at you if you are one of these people. Our goals are very limited—to provide a minimal BASIC-like language for writing simple programs. Later we may make it more complicated, but now the name of the game is keep it simple. That translates to a limited language (no floating point, no sines and cosines, no arrays, etc.) and even this is a pretty difficult undertaking.

Originally we had planned to limit ourselves to the 8080, but with a variety of new machines appearing at very low prices, we have decided to try to make a portable TINY 3ASIC system even at the cost of some efficiency. Most of the language processor will be written in a pseudo language which is good for writing interpreters like TINY BASIC. This pseudo language (which interprets TINY BASIC) will then itself be implemented interpretively. To implement TINY BASIC on a new machine, one simply writes a simple interpreter for this pseudo language and not a whole interpreter for TINY BASIC.

We'd like this to be a participatory design project. This sequence of design notes follows the project which we are doing here at PCC. There may well be errors in content and concept. If you're making a BASIC along with us, we'd appreciate your help and your corrections.

Incidentally, were we building a production interpreter or compiler, we would probably structure the whole system quite differently. We chose this scheme because it is easy for people to change without access to specialized tools like parser generator programs.

THE TINY BASIC LANGUAGE

There isn't much to it. TINY BASIC looks like BASIC but all variables are integers. There are no functions yet (we plan to add RND, TAB, and some others later). Statement numbers must be between 1 and 255 so we can store them in a single byte. LIST only works on the whole program. There is no FOR-NEXT statement. We've tried to simplify the language to the point where it will fit into a very small memory so impecunious, tyros can use the system.

The boxes shown define the language. The guide gives a quick reference to what we will include. The formal grammar defines, exactly what is a legal TINY BASIC statement. The grammar is important because our interpreter design will be based upon it.

IT'S ALL DONE WITH MIRRORS----OR HOW TINY BASIC WORKS

All the variables in TINY BASIC: the control information as to which statement is presently being executed and how the next statement is to be found, the returnaddresses of active GOSUBS——all this information constitutes the state of the TINY BASIC interpreter.

There are several procedures which act upon this state. One procedure knows how to execute any TINY BASIC statement. Given the starting point in memory of a TINY BASIC statement, it will execute it changing the state of the machine as required. For example,

100 LETS = A+6 (cr) would change the value of S to the sum of the contents of the variable A and the interger 6, and sets the next line counter to whatever line follows 100, if the line exists.

A second procedure really controls the interpretation process by telling the line interpreter what to do. When TINY BASIC is loaded, this control routine performs some initialization, and then attempts to read a line of information from the console. The characters typed in are saved in a buffer, LBUF. It first checks to see if there is a leading line number. If there is, it incorporates the line into the program by first deleting the line with the same line number (if it is present) then inserting the new line if it is of nonzero length. If there is no line number present, it attempts to execute the line directly. With this strategy, all possible commands, even LIST and CLEAR and RUN are possible inside programs. Suicidal' programs are also certainly possible.

TINY BASIC GRAMMAR

The things in **bold face** stand for themselves. The names in lower case represent classes of things. '::=' is read 'is defined as'. The asterisk denotes zero or more occurances of the object to its immediate left. Parenthesis group objects.

• is the empty set. | denotes the alternative (the exclusive-or).

line::= number statement @ | statement constatement::= PRINT expr-list

1F expression relop expression THEN statement

GOTO expression INPUT var-list

LET var = expression GOSUB expression

RETURN CLEAR

RUN END

expr-list::= (string | expression) (, (string | expression) *)

var-list::= var (, var)*

expression::= (+|-|E|) term ((+|-|) term)*

term::= factor ((* | /) factor)*
factor::= var | number | (expression)

var::= A | B | C ... | Y | Z

number::= digit digit*

digit::= 0| 1|2|...|8|9

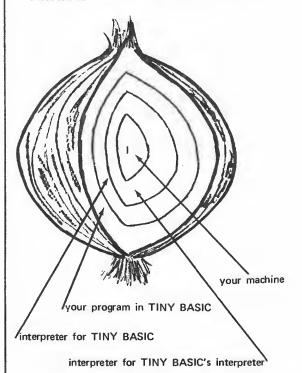
relop::=<(>|=|E|)|>(<|=|E|)|=

A BREAK from the console will interrupt execution of the program.

IMPLEMENTATION STRATGIES AND ONIONS

When you write a program in TINY BASIC there is an abstract machine which is necessary to execute it. If you had a compiler it would make in the machine language of your computer a program which emulates that abstract machine for your program. An interpreter implements the abstract machine for the entire language and rather than translating the program once to machine code it translates it dynamically as needed. Interpreters are programs and as such have their's as abstract machines. One can find a better instruction set than that of any general purpose computer for writing a particular interpreter. Then one can write an interpreter to interpret the instructions of the interpreter which is interpreting the TINY BASIC program. And if your machine is microprogrammed (like PACE), the machine which is interpreting the interpreter interpreting the interpreter interpreting BASIC is in fact interpreted.

This multilayered, onion-like approach gains two things: the interpreter for the interpreter is smaller and simpler to write than an interpreter for all of TINY BASIC, so the resultant system is fairly portable. Secondly, since the major part of the TINY BASIC is programmed in a highly memory efficient, tailored instruction set, the interpreted TINY BASIC will be smaller than direct coding would allow. The cost is in execution speed, but there is not such a thing as a free lunch.



LINE STORAGE

The TINY BASIC program is stored, except for line numbers, just as it is entered from the console. In some BASIC interpreters, the program is translated into an intermediate form which speeds execution and saves space. In the TINY BASIC environment, the code necessary to provide the

QUICK REFERENCE GUIDE FOR TINY BASIC

LINE FORMAT AND EDITING

- Lines without numbers executed immediately
- Lines with numbers appended to program
- Line numbers must be 1 to 255
- Line number alone (empty line) deletes line
- Blanks are not significant, but key words must contain no unneeded blanks
- deletes last character
- X^c deletes the entire line

EXECUTION CONTROL

CLEAR delete all lines and data RUN run program LIST list program

EXPRESSIONS

Operators

All arithmetic is modulo 2¹⁵ (±32762)

INPUT / OUTPUT

PRINT X,Y,Z PRINT 'A STRING' PRINT 'THE ANSWER IS' INPUT X INPUT X,Y,Z

ASSIGNMENT STATEMENTS

LET X=3 LET X= -3+5*Y

CONTROL STATEMENTS

GOTO X+10 GOTO 35 GOSUB X+35 GOSUB 50 RETURN IF X > Y THEN GOTO 30

transformation would easily exceed the space saved.

When a line is read in from the console device, it is saved in a 72-byte array called LBUF (Line BUFfer). At the same time, a pointer, CP, is maintained to indicate the next available space in LBUF. Indexing is, of course, from zero.

Delete the leading blanks. If the string matches the BASIC line, advance the cursor over the matched string and execute the next IL instruction. If the match fails, continue at the IL instruction labeled lbl.

The TINY BASIC program is stored as an array called PGM in order of increasing line numbers. A pointer, PGP, indicates the first free place in the array. PGP=0 indicates an empty program; PGP must be less than the dimension of the array PGM. The PGM array must be reorganized when new lines are added, lines replaced, or lines are deleted.

Insertion and deletion are carried on simultaneously. When a new line is to be entered, the PGM array searches for a line with a line number greater than or equal to that of the new line. Notice that lines begin at PGM (0) and at PGM

(j+1) for every j such that PGM (j)=[carriage return]. If the line numbers are equal, then the length of the existing line is computed. A space equal to the length of the new line is created by moving all lines with line numbers greater than that of the line being inserted up or down as appropriate. The empty line is handled as a special case in that no insertion is made.

TINY BASIC AS STORED IN MEMORY

byte in memory treated as an integer byte treated as a character 'o I N P U T L N @ '20 P JN * N 3 N * N * N @ 13 1 F THENLGOTOL 1 space

Two different things are going on at the same time. The routines must determine if the TINY BASIC line is a legal one and determine its form according to the grammar; secondly, it must call appropriate action routines to execute the line. Consider the TINY BASIC statement:

GOTO 100 At the start of the line, the interpreter looks for BASIC key words (LET, GO, IF, RETURN, etc.) In this case, it finds GO, and then finds TO. By this time it knows that it has found a GOTO statement. It then calls the routine EXPR to obtain the destination line number of the GOTO. The expression routine calls a whole bunch of other routines, eventually leaving the number 100 (the value of the expression) in a special place, the top of the arithmetic expression stack. Since everything is legal, the XFER operator is invoked to arrange for the execution of line 100 (if it exists) as the next line to be executed.

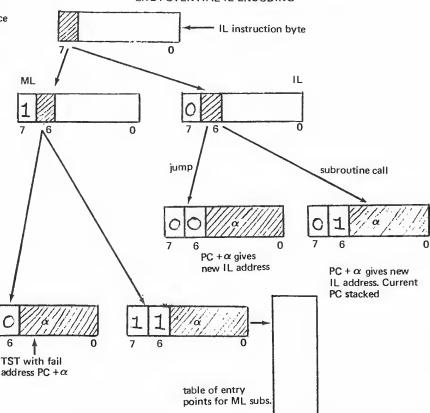
Each TINY BASIC statement is handled similarly. Some procedural section of an IL program corresponds to tests for the statement structure and acts to execute the statement.

ENCODING

There are a number of different considerations in the TINY BASIC design which fall in this general category. The problem is to make efficient use of the bits available to store information without loosing out by requiring a too complex decoding

In a number of places we have to indicate the end of a string of characters (or else we have to provide for its length somewhere). Commonly, one uses a special character (NUL = 00H for example) to indicate the end. This costs one byte per string but is easy to check. A better way depends upon the fact that ASCII code does not use the high order bit; normally it is used for parity

ONE POTENTIAL IL ENCODING



a carriage return symbol **ERRORS AND ERROR RECOVERY**

There are two places that errors can occur. If they occur in the TINY BASIC system, they must be captured and action taken to preserve the system. If the error occurs in the TINY BASIC program entered by the user, the system should report the error and allow the user to fix his problem. An error in TINY BASIC can result from a badly formed statement, an illegal action (attempt to divide by zero, for example), or the exhaustion of some resource such as memory space. In any case, the desired response is some kind of error message. We plan to provide a message of the form: ! mmm AT nnn

where mmm is the error number and nnn is the line number at which it occurs. For direct statements, the form will be: ! mmm

since there is no line number.

Some error indications we know we will need are:

- Syntax error Missing line
- 5 RETURN without GOSUB 6 Expression too complex
- Line number too large 7 Too many lines
- 4 Too many GOSUBs 8 Division by zero

THE BASIC LINE EXECUTOR

The execution routine is written in the interpretive language, IL. It consists of a sequence of instructions which may call subroutines written in IL, or invoke special instructions which are really subroutines written in machine language.

		January 1976 Tiny BASIC Calls	neme	5 X	Orthod	OIILIA I	BOX 510, Memo raik CA 54025 rage of
		se it to indicate the end			EXECUTOR WAIT		
		a string. When we process	statemer	nt, The c	m in fL will execut perators TST, TST lind characteristic	V, TSTN, and PI	as
the characters we must AND the character with			BASIC I	ine. Oth	or operations (NXT pints to another TI	, XPER) move	
07FH to scrub off the flag bit.			line.				
The inte	erpreter opco	odes can be encoded into	THE IL	CONTR	OL SECTION		
		fall into two distinct	START	NLIN	1E	: (NITIAL : WRITE (B/LF
classesthose	which call n	nachine language sub-	co:	TSTU	LINE	: TEST FO	PHOMPT & GET A LINE OR LINE NUMBER
		either call or transfer	STMT:	JMP XINI	T CO		IL (MAY BE OELETE)
		elf. The diagram indi-			ECUTOR	,	
		e. The CALL operations	STMT:	1.37	SITET	: IS STAT	EMENT A LET?
		the IL instruction set. relative to PC for IL		CALL	SIG EXPR	PLACE	ACE VAR ADDRESS ON AESTK.
		rent IL program size.		DON STOI TXN	E ₹ E	STORE	ERROR IF & NOT NEXT. RESULT: QUENCE TO NEXT.
•		is not, the address	St:	TST	\$3, "GO"	COTO O	R GOSUB?
could be used	to index an	array with the ML		DON XPE	S2, 'TO' L EXPR	GET LA	BEL. IF OF NOT NEXT.
class instructi	ons.		S2:	TS1 CAL	\$14, 'SUB'	:EHROR	ANO JUMP. If NO MATCH. STINATION.
	TINY BASIC IN	TERPHETIVE OPERATIONS		DON	E	; £RROR	FF W NOT NEXT.
	TST Itil, 'string'	delete leading blanks	53	TST	S8, PRINT	AND JUI	AP.
		If string matches the BASIC line, advance cursor over the matched string and execute the next IL instruction. If a	\$4. \$5:	121 PRS 121	\$7, ***: \$8, *, *	PRINTS	R OUDTE.
		match fails, execute the IE instruction at the labeled lbl.	30.	SPC	S4	, IS THEH , SPACE T ; YFS, JUI	O NEXT ZONE.
	CALL Ibi	Execute the fL subroutine starting at Itil. Save the IL ad-	Sō:	NUN	E	NO ERR	OR IF NO er.
		dress fullowing the CALL on the control stack.	\$7:	CALI		GET EX	PR VALUE.
	RTN	Return to the IL location specified by the top of the con- trol stack.	ça.	PAN PAP TST	55 \$9, '18"	IS THER	E MORE?
	DONE		S8	CALL	EXPR	DEFFRA	PRESSION. TINE OPE AND PUT ON STK.
	DONE	Report a syntax error if after deletion leading blanks the cursor is not positioned to read a carriage referen.		CALL	EXPR	; GLT EXF	PRESSION. MICOMPARISON - PERFORMS NEXT IF FALSE.
	JMP fbi	Continue execution of IL at the label specified.	59:	75T	STMT S12, 'INPUT	GET NE	KT STATEMENT. TATEMENT. 9 ADDINESS.
	PRS	Print characters from the BASIC text up to but not including the	\$10:	INNU	M.	: GLT VAI : MOVE N : STORE 1	UMBER FROM TTY TO AESTK."
	,5	closing quote mark. If a cr is found in the program text, report an		JMP	St 1. '." S10	15 THE B	F MORE?
		error. Move the cursor to the point following the closing quote.	S11:	DON		MUST BE SEQUEN	er, ce to next. Statement,
	PRN	Print number obtained by popping the top of the expres- sion strick.	S12:	DON!		: MUST RE	STATEMENT.
	SPC	finsert spaces to move the print head to next zone.	\$13:	NXT TST	S14, 'END'	SEQUEN	CE TO NEXT STATEMENT.
	NLINE	Output CRLF to Printer.	St 4:	FIN	SIS, LIST	; LIST CO	MAND.
				LST NXT			
	NXT	ff the present mode is direct (line number zero), then return to line collection. Otherwise, wheat the next	\$15:	TST	Ste, 'RUN'	; RUN CO	MAND.
		sequential line and begin interpretation.	S16:	NXT TST	S17, "CLEAR"	CLEAR C	OMMAND.
	XFER	Test value at the top of the AE stack to be within range. If not, report an error. If so, attempt to position cursor		JMP S	TART		
		at that line. If it exists, begin interpretation there; if not report an error.	\$17:	ERR		:SYNTAX	ERROR.
			EXPR:	CALL	EO."-"	TEST SO	R UNARY
	SAV	Place present line number on SBRS1K. Report overflow as error.		NEG	E1	:GET VAI	LUE.
	RSTR	Replace current line number with value on SBRSTK. If	E0:	CALL TST	£1,'+' TERM E2.'+'	: LOOK FO	OR MORE, R UNARY+,
		stack is empty, report error.		CALL	TERM	SUM TER	G TERM.
	CMPR	Compare AESTK(SP), the top of the stack, with	€2:	JAIP TST	E1 E3,'-*	:ANY MO	AE?
		AESTK(SP-2) as per the relation indicated by AESTK(SP-1). Defete all from stack. If condition specified did not match,		CALL SUB MIP		;OIFFERI	ENCE TERM.
		then pertorm NXT action.	E3; T2;	RIN	£3	; ANY MO	RE?
	INNUM	Read a number from the terminal and push its value onto the AESTK.	TERM	A:	CALL	FACT	
	FIN	Fieturn to the line collect routine.	TO:		TST	T1, '"'	DOODUUT TAOYOR
	ERR				CALL MPY	FACT	;PRODUCT FACTOR.
		Report syntax error and return to line collect routine.			JMP	TO	
	ADD	Replace top two elements of AES1K by their sum.	T1:		TST CALL	T2, '/' FACT	; ANY MORE?
	SUB	Replace top two elements of AESTK by their difference.			DIV	FACI	; QUOTIENT FACTOR.
	NEG	Replace top of AESTK with its negative.			JMP	TO	
	MUL	Replace top two elements of AESTK by their product.	FAC1	r.	TSTV	FO	· VARIARI E
	DIV	Replace top two elements of AESTK by their quotient.	. 70	•	IND		: VARIABLE. ; YES, GET THE VALUE.[
	STORE	Place the value at the top of the AESTK into the varieble			RTN	٠.	
		designified by the index specified by the value immediately below it. Delete both from the stack.	FO:		TSTN RTN	F1	; NUMBER, GET ITS VALUE.
	TSTV Ibi	Test for variable (i.e. letter) if present., Place its index	F1:		TST	F2, '('	; PARENTHESIZED EXPR.
		value onto the AESTK and continue execution at next suggested location. Otherwise, continue at lbl.			CALL	EXPR	
	TSTN Ibi	Test for number. If present, place its value onto the			RTN	F2,')'	; MATCHING PARENTHESIS.
	13114 101	AEST K aixf continue execution at next suggested location.	F2:		ERR		;ERROR.
		Otherwise, continue at Ibl.					
	IND	Replace top of stack by variable value if indexes.	RELO	OP:	TST	RO, '='	
	LST	fist the contents of the program area.			LIT	0	;=
	INIT	Performs global initialization	RO:		RTN TST	R4, '\'	
	CCTI INC	Clears program area, emptys GOSUB stack, etc.			TST	R1, '='	
	GETLINE	Input a line to LBUF.			LIT RTN	2	; <= R4: TST S17, ')' TST R5='
	TSTL Ibi	After editing leading blanks, look for a line number. Report error if invalid; transfer to lot if not present.	R1:		TST	R3, '>'	LIT 5 ;>=
	INSRT	finent line after deleting any line with same line			LIT	3	;(> RTN '
		number.	R3:		RTN LIT	1	R5: TST R6,'\(\alpha\)'
	XINIT	Perform initialization for each stated execution.	no.		RTN	•	R6: LiT4 ; >
		Empties AEXP stack.					RTN

TINY BASIC

by Dennis Allison, Bernard Greening, happy Lady, & lots

(reprinted from People's Computer Company Vol. 4, No. 3)

Dear People.

After a quick pique at TINY BASIC I have the following (possibly ill-considered) comments:

- 1. It looks useful for tiny computers, which is as intended.
- 2. Those accustomed to extended BASIC, or even the original Dartmouth BASIC, will be irked by its limitations. But then, that's how the bits byte!
- 3. How does the interpreter scan the word THEN in an IF statement?
- 4. Some of the comments for EXPR seem to be on the wrong line, or my reading is more biased than usual.
- 5. Users should note that arithmetic expressions are evaluated left-toright unless subexpressions are parenthesized (i.e., there is no implicit operator procedure).
- 6. Real numbers would be nice, but would take up a lot more space. Probably too much. Ditto for arrays and string variables.
- 7. Please consider adding semicolon (i.e., unzoned) PRINT format with a trailing semicolon inhibiting the CRLF. This would be very useful and would be easy to add.
- 8. If INPUT prompts with a question mark, please print a blank character after the question mark (for readability).
- 9. I suggest allowing THEN as a separator in any multi-statement line, not just in IF statements. Since lines like

IF 5(X THEN IF X(10 THEN GOSUB 100

are already legal, why not allow lines like

LET A=B THEN PRINT A

or any other combination, including silly ones like

GOTO 200 THEN INPUT Z

the second statement of which would never be executed. If THEN works for IF, it should be possible to make it work for anything. 10.1 also suggest allowing comments somehow. At present, comments must be held to a minimum

are possible via subterfuges such as

IF X > X THEN PRINT "THIS IS A COMMENT"

but that seems kind of gauche. Naturally, comments must be held to a minimum in TINY BASIC, but sometimes they may be vital. 11. Doing a

PRINT"

seems to be the only way to print a blank line. Well, all right.

12. Exponentiation via ** would seem fairly easy to add, and might be worthwhile.

13. By the way, all of this will execute in 1K, won't it?

Jim Day 17042 Gunther St. Granada Hills, CA 91344

Answering your Questions by number where appropriate:

- 3&4. Woops! There should be a TST instruction to scan the THEN. The comments are displaced a line. See the corrected IL listing in this issue.
- 5. Expressions are evaluated left-to-right with operator precedence. That is, 3+2*5 gives 13 and not 25. To see this, note that the routine EXPR which handles addition gets the operands onto the stack by calling TERM, and TERM will evaluate any product or quotient before returning.
- 7. Agreed, but this is intended as a minimal system.
- 9. One man's syntatic sugar is anothers poison. I don't like the idea. Incidentally, how would you interpret

LET A=B THEN GOSUB 200 THEN PRINT 'A'

The GOSUB then has to store a program address which botches up the line entry routine or one has to zap the GOSUB stack whenem error is found. Both are solved only by kludges. 10-12. See 7.

13. Maybe. But 2K certainly. See below.

Dear PCC.

I am thrilled with your idea of an IL but I think that if you intend only to write a BASIC interpreter that a good symbolic assembler would be appropriate. With an assembler similar to DEC's PAL 3 or PAL B the necessary routines could be written and used in nearly the same way without having to write the associated run time material that would be necessary for its use as an interpreter. A command decoder, a text buffer, and a line editor would be necessary and all of this uses up a good amount of space in memory.

If you are aware of all these things and still plan to develop an IL interpreter, then I suggest you start as DEC did with a simple symbolic editor as the backbone of the interpreter. In this way you allow a 2800% increase in development and debug ging speed (according to Datamation's comparison of interpreters and compilers whose fundamental difference is the on line editing capability). Once this has been implemented and IL is running on a particular system then the development of interpreters of all types is greatly simplified. By suggesting IL you have stumbled onto the most logical and easiest way to develop a complete library of interpreters. In addition to BASIC, it is very easy to write interpreters for: FOCAL, ALGOL, FORTRAN, PL 1, LISP, COBOL, SNOWBAL, PL/m, APL, and develop custom interpreters ters-with the ease with which one would write a long BASIC program!

As I pointed out earlier, all these features take up memory space and, as you have pointed out, run time is much slower. The way around this is to define the IL commands in assembly language subroutines then assemble the completed interpreter as calls to these subroutines. Thus the need for the IL interpreter as a run time space and time consumer is no longer necessary! (OK symbolic assembler haters, let's see you do this in machine language in less than ten man-years!)

In places where time and space are not so much of a problem, I suggest the addition of an interrupt handler and priority scheduler to allow IL to be used as a simplified and painless TIMESHARED system enabling many users to run in an interpreter and use more than one interpreter at once. Multi-lingual timeshare systems previously being available to those who have a highspeed swapping disk, drum, or virtual memory, are now avaliable to the user who has about 16K of memory and a method of equitably bringing interpreters in to main memory from the outside world (a paper tape reader or cassette system is the easiest to come by).

In short, IL as I suggested, in its minor stages would be a powerful software development aid; and in its final, most complex stages would provide a runtime system of unheard of

I have heard from unofficial sources that ordinarily an interpreter or compiler requires ten man-years to write and debug to the point of use (if one man works the job would require 10 years, if 10 men work it would take one year). Since this is to be expected as the initial development of IL. and since I have a general idea of the circulation of PCC, we should have IL up and running by the next issue of PCC!!

At this time I would like to request a few reprints of the article dealing with IL because I want-to get some help from others in my school in getting a timeshared version working on our 16K PDP 8/m with DECTAPE. I seem to have lent my copy of that issue to one of the people I had been trying to get on this project and be has not returned it to me. Meanwhile, I need the article to begin initial work on the interpreter to insure compatibility with the version coming across through PCC. I will keep you posted as with regards to the development.

> William Cattey 39 Pequet Road Wallingford, Ct. 06492

The IL approach to implementation is quite standard and dates back to Schorre's META II, Gleenie's Syntax Machine, and numerous early compilers. It was widely used in the Digitek FORTRAN systems. We did not "stumble" on to the technique, we chose it with some deliberation.

You are right that a symbolic assembler can be used either to assemble the pseudocode into an appropriate form or to

expand the pseudocode into actual machine instructions with the attendant cost in space (and decrease in execution time). Our goal is a small, easily transportable system. The interpretive approach seems consistant with this primary goal. We are using the Intel 8080 assembler's macro facility to assemble our pseudocode.

I certainly agree that it is relatively easy (but not simple!) to implement other languages using the IL approach. From the users standpoint, provided he is not compute bound, there is little difference. Interpreters are often a bit more forgiving of errors and can give better diagnostics.

In my experience, your figure of 10 man-years is high for some languages and low for others. A figure of two to four man-years is probably more accurate, and that includes documentation at both the implementation and user level. Good luck on your implementation.

....I have found in my adaptation of it (TINY BASIC IL) for full use that certain commands need strengthening, while some might be dropped. I will hopefully be coming out with these possible modifications. Concerning my ideas on space trade-offs; I think an assembled version would take less space, since each command is treated as a subroutine call in a program made up of routines, while the interpreter needs a run time system in the background which, since it is interpretive in itself, takes up space.

P.S. You missed my allusion to assembler over strictly octal or hexidecimal op codes (my meaning was twofold). In DEC's PAL8 assembler the following syntax is needed to make the most efficient use of routine calling:

TSTN=JMSI (jump to subroutine indirectly via this location) 10. XTSTN

The assembler shows the binary as if TSTN were like a JMSI 100/ JAP to subroutine indirectly via 100 (requiring very very little extra space per routine-one word, to be exact).

I would be happy to resolve any questions regarding compilers vs. interpreters. (Datamation did an article on the writing of a standard program in several languages then documented development and run time.) William Cattey

There are several different varieties of interpreters. One is simply a sequence of subroutine calls. Another is, as you suggest, a list of indirect references to subroutine calls. We are considering a different organization where the call address and some additional information is packed into a single byte. This is a good strategy vis a vis memory conservation only if the size of the code memory to decode the packed instruction plus the size of the encoded instructions is smaller than the size of a more straightforward encoding. This remains to be seen.

I guess I did miss your point on assemblers. However, let me assure you that I would never advocate making software by programming directly in hex or binary. Even an assembler seems cumbersome and difficult to me; I prefer a good systems language like PL/M!

Dear Dennis and other PCCers,

In my last crazily jumbled letter I made some comments about TINY BASIC. Here is the result of 2-3 days work and thinking about it. Instead of having an interpretive IL. I chose to set it up as detailed as possible, then have people with different machines code up subroutines to perform each IL instruction. I'm not convinced that this way would take more space, and I'm sure it would be faster

There are a couple of changes in the syntax from your published version: separate commands from statements, add terminal comma to PRINT, and restrict IF-THEN to a line number (implied GOTO).

The semantics are separated out from the syntax in IL as much as possible. This should make it easier to be clear about what the results of any given syntatic structure. This is most apparent in the TST instructions, and the elimination of the NXT instruction. That one in particular was a confusion.

Please let me know how this fits with what you're doing. I don't have a micro yet-time, not money, prevents it.

John Rible 51 Davenport St. Cambridge, MA 02140

Because of space limitations, we have not been able to publish all of John Rible's version (dialect) of TINY BASIC. We'll probably include it in the first issue of the TINY BASIC NEWSLETTER. Limited space requires it to be in 2nd issue.

By seperating the syntax from the semantics he has produced a larger and possibly simpler to understand IL. There are more IL instructions so, I believe, the resultant system will be larger; further, the speed of execution is roughly proportional to the number of IL instructions (decoding IL is costly), it will be slower.

EXTENDABLE TINY BASIC

JOHN RIBLE

INTERMEDIATE LANGUAGE PHILOSOPHY

Instead of IL being interpreted, my goal has been to describe IL well enough that almost anyone will be able to code the instructions as either single machine language instructions or small subroutines. Besides speed-ing up TINY BASIC, this should decrease its size. Most of the instructions are similar to those of Oennis' (PCC V4 no. 2), but the syntactical has been seperated from the active routines. This would be useful it you want the syntax errors to be printed while inputting the line, rather than when RUNning the program.

Most subroutines (STMT, EXPR, etc.) are recursively called, so in

addition to the return address being stacked, all the related data must be stacked. This can use up space quickly.



2

(A)

SYNTAX for John Rible's version of TINY BASIC

⟨PROGRAM⟩::-⟨PLINE>*¹ (PLINE):--(NUMBER) (STATEMENT) (ILINE > :: + (COMMANO) | (STATEMENT) ⟨COMMANO⟩ ::= CLEAR@|LIST@|RUN@

(STATEMENT) ::= **®**

LET(VAR) = (EXPR) GOSUB (EXPR)@

PRINT (EXPR-LIST) (, (E) (G) IF (EXPR) (RELOP) (EXPR) THEN(STATEMENT) 1 INPUT (VAR-LIST > 1991

ENO 🚱 CEXPR-LIST):= ((STRING)|(EXPR)) (((STRING)|(EXPR)))*2

CSTRING):= "(ANY CHAR) *2"

CANY-CHAR)::= any character except " or (**)

CEXPR)::= (+1-| (**)(TERM)((+1-)(TERM))*2

CEMPY::= (+1-| { } (TERM) { (+1-) (TERM) };

(TERM)::= (FACTOR) ((*|/) (FACTOR)) * 2

(FACTOR)::= (VAR) { (NUMBER) { (KEMPR) };

VAR-LIST)::= (VAR) { (VAR) } * 2

VAR)::= A | B | | Y | Z

(NUMBER)::= (OIGIT) { (OIGIT)**

(OIGIT)::= 0 | 1 | | B | 9

(RELOP)::= ((= | Y | E) |) (= | (| E) | =

notes: € is null character

actual characters are in bold face

*1 repeat limited by size of pro-

repeat limited by size of program memory space

*2 repeat limited by length of line *3 repeated 0 to 4 times

Dear Mr. Allison,

I was very interested in your Tiny BASIC article in PCC. Your ideas seem quite good. I have a few suggestions regarding your IL system. I hope I am not being presumptuous or premature with this. Unless I misunderstood you, your IL encoding scheme seems inadequate. For instance, IL JMPs must be capable of going up and down from the current PC. This means allotting one of the 6 remaining bits of the IL byte as a sign bit resulting in a maximum PC change of ±31 which is not adequate in some cases, ie. the JMP from just above S17 back to START. May I suggest the following scheme which is based on 2 bytes per IL instruction:

<u>IL</u>		ML	
JMP	CALL	TST	CALL
0XX ⁸	1XX ₈	2XX ₈	1XX ₈ (1st byte)
YYY8	YYY8	YYY8	YYY ₈ (2nd byte)

where XX= lower 6 bits of high part of address (assume upper 2 bits are 00)

YYY= all 8 hits of low part of address.

The complete address being 0XXYYY $_8$. These addresses represent the locations associated with the 1 L and ML instructions. Note that if $\mathfrak p$ points to a table with a stored address, you have 3 bytes used—my scheme uses only 2 bytes with the same basic information.

implementation I am using the following technique: the string follows the TST byte pair immediately with a bit 7 set in the last character.

I also wondered about the TST character string. In my

On the TSTL, TSTV, and TSTN IL's, it appears you need a ML address for the particular subroutine and 2 additional hytes for the fail address. At least this is how I am handling it.

I am looking forward to future articles in the series. Thanks again— keep up the good work!

P.S. I am co-owner of an Altair. We are writing our Tiny BASIC in Baudot to feed our Model 19's.

Richard Whipple 305 Clemson Dr. Tyler, Tx. 75701

We found the same problem with the published IL interpreter. We solved it by doing a bit of rearranging and introducing a new operations code which does jumps relative to the start of the program, but has the same basic encoding. Your mechanization will, of course, work, but requires one more byte per IL instruction, may be harder to implement on some machines, and takes more code.

We are using the same scheme of string termination (i.e., using the parity bit) as you are. It's simple, easy to test, and difficult to get into the assembler.

There are a few errors and oversights in the IL language and in the interpreter you didn't mention. See the new listing in this issue.

Good luck. Keep us informed of your progress.

Dear People at PCC,

I have a couple of comments on Tiny BASIC:

S4 says TST S7, but S7 got left out. T1 says TST on my paper which I suppose should be TST T2.

What is LIT and all these "or 2000"? When are we going to start putting some of this into machine code?

Sincerely,

BOB BEARD 2530 Hillegass, No. 109 Berkeley CA 94704

Soon! Ed.

Dear Tiny BASIC Dragon,

Please scratch my name onto your list for Tiny BASIC Vol. 1. Enclosed is a coupon for 3 chunks of fire.

I am really enjoying my subscription to PCC, especially the article on Tiny BASIC.

Someday I am going to build an extended Tiny BASIC that will take over the world.

Basically yours,

RON YOUNG 2505 Wilburn, No. 144 Bethany OK 73008 Since the last issue came out, the IL code, macro definitions for each IL instruction, a subroutine address table for the assembly language routines that execute the IL functions, the assembly language code that executes the IL functions (all except the 16-bit arithmetic ones), and the IL processor have been punched on paper tape in source form.

HOP, TST, TSTN, and TSTL now do branches +32 relative to the current position counter. If the relative branch field has a zero in it, indicating a branch to "here", the IL processor prints out the syntax error message with the line number. The ERR instruction that was in the old IL code no longer exists.

IJMP and ICALL are used because the Intel 8080 assembler uses JMP and CALL as mnemonics for 8080 instructions. IJMP and ICALL are followed by one byte with an unsigned number from 0 to 255. This is added to START to do an indexed jump or call.

Bernard

corrected INTERPRETIVE LANGUAGE SUBROUTINES EØ JTEST FO... *-* OR 2000 TERM JPUT TERM ON AESTK JNEGATE VALUE ON AESTK EI JGO GET A TERM TINY BASIC IL EXPR: TST DB ICALI. NEG HOP ;TEST FOR UNARY '+' STATEMENT EXECUTOR WRITTEN IN IL (INTERPRETIVE LANGUAGE) THIS IS WRITTEN IN MACROS FOR THE INTEL INTELEC BYMOD 86 SYSTEM USING INTEL'S ASSEMBLER. EØ 1 TST E01 ;TES DB JPUT TERM ON AESTK JTEST FOR ADDITION EØ1: ICALL TERM E2 ;1L: E1 1 TST CONTROL SECTION DB JGET SECOND TERM JPUT SUM OF TERMS ON AESTK JLOOP AROUND FOR MORE START: INIT ERRENT: NLINE #INITIALIZE JINITIALIZE JURITE A CR-LF JURITE PROMPT NO GET A LINE JIF NO LINE NUMBER GO EXECUTE IT JINSERT OR DELETE THE LINE JLOOP FOR ANDHER LINE JINITIALIZE FOR EXECUTION CALL TERM ADD COI GETLN TSTL HOP ΕĪ XEC INSRT CO E3 ;TE: ITEST FOR SUBTRACTION E2 : TST XINIT XEC : DB JGET SECOND TERM JPUT DIFFERENCE OF TERMS ON AESTK CALL TERM STATEMENT EXECUTOR SUB SI JCHECK FOR 'LET' 'LE', 'T' OR 2000 SEI JERROR IF NO VARIABLE! SE2 JERROR IF NO "=" '=' OR 2000 EXPR JPUT EXPRESSION ON AESTK JCHECK FOR CR LINE TERRINATOR JPUT VALUE OF EXPRESSION IT ITS CELL JCONTINUE NEXT LINE HOP 1LOOP AROUND FOR MORE TST STMT DB TSTV TST DB ICALL SEI 1 SE21 JTHIS CAN BE RECURSIVE E3 1 RTN DONE STORE NXT TERM: FACT JGET ONE FACTOR ICALL T1 ;TES '*' OR 2000 ;TEST FOR MULTIPLICATION TØ : TST DB TST DB TST DB ICALL DONE S3 JCHECK FOR "GO" "G"."O" OR 2000 S2 JCHECK FOR "GOTO" "I","O" OR 2000 EXPR JGET THE LABEL JCHECK FOR CR LINE TERMINATOR JDO A "GOTO" TO THE LABEL Śli ICALL FACT JGET A FACTOR PUT THE PRODUCT ON AESTK MPY HOP \$LOOP AROUND FOR MORE T1 TST STEST FOR DIVISION '/' OR 2000 FACT JGE DB JGET THE QUOTIENT JPUT QUOTIENT ON AESTK JLOOP FOR MORE ICALL 521 TST DIV DB TØ ICALL HOP XPER DONE T2 IRETURN TO CALLER RTN XFER 3 53 t TST DB TST FACTI TSTV FØ TEST FOR VARIABLE GET INDES OF THE VARIABLE 54: IND PR\$ TST RTN SSI FØ: TSTN FΙ STEST FOR NUMBER DB SPC HOP DONE NXT JERROR IF ITS NOT A '(' F1 1 TST (' OR 2000 JTHIS IS A RECURSIVE PROCESS ICALL. EXPR S9 JCHECK FOR 'IF' '1',F' OR 2020 EXPR JGET THE FIRST EXPRESSION RELOP JGET THE RELATIONAL OPPERATOR EXPR JGET THE SECOND EXPRESSION SBA JCHECK FOR 'THEN' 'THE','N' OR 2020 SIF NOT TRUE CONTINUE NEXT LINE SIMI JIF TRUE PROCESS THE REST OF THIS LINE 58: JEVERY '(' HAS TO HAVE A ')' E1: TST ')' OR 2000 ICALL ICALL ICALL TST DB DB RTN SBA CMPR RØ ;CHI RELOP: JCHECK FOR "=" IJMP DB LIT 591 TST JCHECK FOR 'INPUT' JCHECK FOR 'INPUT' ''' OR 2000 JGET THE VARIABLE'S INDEX JCET THE NUMBER FROM THE TELETYPE JPUT THE VALUE OF THE VARIABLE IN ITS CELL J'' MEANS MORE DATA INPU' T RTN SIE ICALL. VAR INNUM RØ i TST R4 ;CH: '<' OR 2000 JCHECK FOR '<' STORE DB TST S11 J., TST RI '=' OR 2000 JCHECK FOR CR LINE TERMINATOR JCONTINUE NEXT LINE 511: DONE NXT DB 2 RTN 3 5121 SI3 JCHECK FOR 'RETURN' 'RETUR','N' OR 2000 JCHECK FOR CR LINE TERMINATOR JRETURN TO CALLER TST DB R1 1 TST R3 JCH1 JCHECK FOR '>' DONE DB LIT 3 5131 S14 JCHECK FOR 'END' 'EN','D' OR 2000 JGO BACK TO CONTROL MODE RTN TST DB FIN R3: LIT 1 RTN SIS JCHECK FOR 'LIST' 'LIS','I' OR 2000 JCHECK FOR CR LINE TERMINATOR JTYPE OUT THE BASIC PROGRAM JCONTINUE NEXT LINE SIA TST DB R4 I TST R4 '>' OR 2000 DONE DB R5 *=* OR 2000 TST NXT DB S16 JCHECK FOR "RUN" "RU", 'N" OR 2000 " JCHECK FOR CR LINE TERMINATOR JCONTINUE NEXT LINE 51 51 TST RTN DB Done R6 R5: TST NXT ĎВ 1.17 3 CHECK FOR 'CLEAR', FAILURE IS AN ERRORI 'CLEA', 'R' OR 2000 START JREINITIALIZE EVEHYTH:NGI RTN R6: LIT 4 RTN

December 12, 1975

The Tyler Branch of the North Texas Computer Club is still having fun with Tiny BASIC as you can see by examining the print-out that follows. We are now calling it Tiny BASIC Extended after the addition of FOR-NXT loops, DIMension statements-arrays, and a few other goodies. The LIFE program was written by David Piper, a high school student of John's (he teaches at Robert E. Lee High School). David is working on KINGDOM now-we can hardly wait. Below are a few comments about our system and Tiny BASIC that may be of interest to your readers.

- 1. Our Altair 8800 is interfaced to a Model 19 Baudot Teletype at John's and via modems and a leased telephone line to a Model 15 Teletype at my house about 3/4 mile away. At present the system is strictly BAUDOT--no ASCII conversion whatsoever.
- 2. We use a Suding-type cassette interface that has been very reliable. 4K bytes load in about 1 minute 20 seconds.
- 3. The Tiny BASIC Extended takes about 2.9K bytes of memory.
- 4. The storage format for our Tiny BASIC is as follows: 2 byte statement label - 1 byte length of text - multibyte text - (cr) The statement label range is 1 to 65535. The "length of text byte" is used to speed up label searching in GOTO and other branching.
- 5. To conserve memory, we have shortened some commands to two or three letters (i.e., PR for PRINT, IN for INPUT, NXT for NEXT, etc.).
- 6. A "\$" is used to write multi-statement lines. A "!" is used to suppress new line output in a PR statement. This allows continuing the next PR on the same line. The ";" provides one skipped space in a PR statement.
 - 7. Functions currently on line are:

RN + generates random numbers between 0 and 10,000 decimal.

TB (exp) + TAB function in PR statement produces a number cipped spaces equal to the value of "exp," an arithmetic ession.

Memory for arrays is allotted from the top of memory down as the program builds from the bottom up. If they gross you get of skipped spaces equal to the value of "exp," an arithmetic expression.

- 8. Memory for arrays is allotted from the top of memory down while the program builds from the bottom up. If they cross, you get error message. Arrays may be 1 or 2 dimension. Max. size: 255 by 255
 - 9. Here are some BAUDOT equivalances:
 - = (equal to)
 -): > = (greater than equal to)
 - (: <= (less than equal to)
 -)(<> (not equal to)
 - + (plus)
 - * (times)

Parentheses are also used in arithmetic expressions. The system understands the difference by context.

10. FOR I=1,1000

NXTI **FND**

takes about 1.6 seconds to execute.

- 11. The colon is used as a Tiny BASIC prompt.
- 12. "?" is used as a rubout key and two LTR's keystrokes are used to begin a line over (LTR and FGS are keystrokes used to change case in Model 15/19 Teletypes)
- 13. Model 15/19 Teletypes are great machines and we have proved their worth to computer hobbyists!

Thanks again for your fine work at PCC, we remain Yours Truly,

DICK WHIPPLE JOHN ARNOLD 305 Clemson Dr Rt 4, Box 52A Tyler TX 75701 Tyler TX 75701

```
00090 PR "LIFE WITH TINY BASIC EXTENDED" 00100 PR "SIZE";1
 00105 LET F:0
 OO110 IN A OO112 PR "THE BEGINNING-WAIT"S PR"
00112 FR 3 FR 42

00115 LET B:A&2

00120 DIM G(B,B),H(B,B)

00130 FOR J:1 TO B

00140 FOR I:1 TO B

00150 LET G(I,J):0$ LET H(I,J):0
 00160 NXT I
                                                                RIIN
 00170 NXT J
                                                               LIFE WITH TINY BASIC EXTENDED SIZE 1 11
THE BEGINNING-WAIT
00175 LET M:A&1
00180 FOR J:2 TO M
00190 FOR I:2 TO M
00200 IN K
00210 IF K (: 1 GO TO 220
 00212 LET I:M
 00214 GO TO 230
00220 LET G(I,J):K
00225 IF K: 1 LET F:F&1
                                                                     10111111111111111111
 00230 NXT I
00240 PR**
00240 PR**
00250 NXT J
00260 PR **GENERATIONS**;1
00270 IN D
00280 PR**
00285 PR**
                                                               CENERATIONS ? 3
                                                               GENERATION O
                                                               POPULATION IS 7
O0285 PR""

O0287 LET S:0

O0290 FOR E:S TO D

O0300 PR "GENERATION";E$PR""

O0301 PR""

O0302 IF F ) 0 GO TO 305

O0303 PR "POPULATION IS ZERO"$PR""$END

O0305 PR "POPULATION IS ";F$PR""
                                                                                          . . . . . . . .
                                                                                     GENERATION 1
 00310 GO SUB 6000
 00315 LET F:0
00320 GO SUB :
             GO SUB 5000
                                                                                     POPULATION IS 15
00330 NXT E
00335 PR "HOW MANY MORE"; ISIN CSPR"
05030 LET N:G(I-1,J-1)&G(I,J-1)&G(I&1,J-1)&G(I-1,J)&G(I&1,J)
05040 LETN:N&C(I-1,J&1)&G(I,J&1)&G(I&1,J&1)
05110 IF G(I,J) )( 1 GO TO 5180
05120 IF N ) 1 GO TO 5150

LENERATION 2
05120 LTT H(I,J):0
05140 GO TO 5210
05150 IF N (: 3 GO TO 5200
05160 LET H(I,J):0
                                                                                                POPULATION IS 12
05160 LEF H(I,J):0
05170 GO TO 5210
05180 IF N )( 3 GO TO 5210
05200 LET H(I,J):1
05205 LET F:Fk1
05210 NXT
05220 NXT
05230 FOR I:1 TO B
05240 FOR J:1TO B
05250 LET G(I,J):H(I,J)
                                                                                                GENERATION 3
                                                                                                POPULATION IS 22
05260 LET H(I,J):0
05270 NXT J
05280 NXT I
05290 RET
06000 FOR J:2 TO M
06010 LET R:0
06020 FOR 0:1 TO M
06030 IF G(Q,J) : 1 LET R:1
06040 NXT Q
06050 IF R:0 GO TO 6120
06060 FOR I:2 TO M
                                                                                               HOW MANY MORE ? 1
06070 IF G(I,J) : 1 GO TO 6100
06080 PR " ";1
                                                                                               CENERATION 4
06080 PR "";1
06090 GO TO 6110
06100 PR "#";1
06110 NXT I
06120 PR ""
06130 NXT J
                                                                                               POPULATION IS 16
06150 RET
                                                                                               HOW MANY MORE ? O
```

TINY BASIC, EXTENDED VERSION

by Dick Whipple (305 Clemson Dr., Tyler TX 75701) & John Arnold (Route 4, Box 52-A, Tyler TX 75701)

INTRODUCTION

The version of TINY BASIC (TB) presented here is based on the design noted published in September 1975 PCC (Vol. 4, No. 2). The differences where they exist are noted below. In this issue we shall endeavor to present sufficient information to bring the system up on an Itel 8080-based computer such as the Altair 8800. Included is an octal listing of our ASCII version of TINY BASIC EXTENDED (TBX). In subsequent issues, structural details will be presented along with a source listing. A Suding-type cassette is now available from the authors (information to follow). We would greatly appreciate comments and suggestions from readers. Unlike some software people out there, we hope you will fiddle with TINY BASIC EXTENDED and make it less Tiny

ABBREVIATED COMMAND SET

TB AND TBX

LET PR GOTO GOSUB RET IF IN In In TB LST TBX RUN NEW * SZE DIM FOR NXT

STANDARD BASIC

LET PRINT GOTO GOSUB RETURN IF INPUT LIST RUN NEW SIZE DIMENSION FOR NEXT

*CLEAR in original TB

TEX - HOW IT DUFFERS FROM TB

- 1. TEX system prompt is a colon ":".
- 2. Statement labol values 1 to 65535.
- 3. Frror correction during line entry:
 - a) Rubout (ASCII 1778) to delete a character. Prints
 - b) Control L (Form Feed ASCII 0148) to delete full line.
- IN Statement: Tormination of numeric input is accomplished by SPACE keystroke. All other terminations use CR (Carriage Return) .
- PR Statement: A comma is used for zone epacing while a semicolon produces a single epace. A comma or semicolon at the end of e line surpresses CR and IF (Line Feed). To skip a line, use PR by itself.

DIM Statement: One or two dimensional arrays permitted. Array arguments can be expressions.

LET V = 10Example: 10

20 DIM A(10,10),B(2+V)

Array variables can be used in the same manner as ordinary variables.

 FOR end NYT Statemente: Step equal to 1 only. Iterative limits can be expressions. Nesting permitted. Care must be exercised when exiting a loop prior to completion of indexing. See Example.

Example:

LET X = 10 10 FOR I = 1 to X LET Y = 2 * A+B 30

1F Y= Z I = X\$NXT I\$GOTO 60 *

NXT I 60 LET Y=3

* For explanation of "}" eee no. 9.

8. Available Functions:

e) RN: Random number generator. Range 0≤RN≤10,000.

No argument permitted.

TB(E): Tab function. In a PR etatement, TB(E) prints a number of SPACE's equal to the value of expression "E".

9. The dollar sign can be used to write multiple etatement lines.

Example: 20 LET A=2*(B+1)\$PRASEND

When using an IF statement, a "false" condition transfere execution to the next <u>numbered</u> line. Thus in Line 40 of the example of no. 7, the chained statements will not be executed unless a "true" condition is encountered.

10. LST Command: Can take anyone of three forms:

LST CR- lists all statements in program

LST a CR — lists only statement labolled a LST a,b CR — lists all statemente between labels a and b c)

11. SZE Command: Printe two decimal numbers equal to:

Number of memory bytes used by current program.

Number of memory bytes remaining.

Note: Array storage included only after first execution of program.

12. Recording Progrems on Caseette: Core dumpe to caesette should begin at 033350 (polit octal) and continue through address etored at

033354 (low byte of address)
033355 (high byte of address)

Of course these cassette programs should be loaded . back at 033350.

IMPLEMENTING THX

Memory Allocation:

- I. Misc. Storege (I/O Routines) 000000 to 000377*
- II. TEX 020000 to 033377
- III. TEX Programs 034000 to upper limit of memory.
- * In our system we maintain a Monitor/Editor in the first 1K byte of memory. 3/4 K is protocted and 1/4 K can be used for system PAM. Such a configuration is useful but not necessary.

External Program Requirements:

1. System Entry Routine --

ADRS	INST	
000000	061)	
000001	377 1	LXI SP
000002	000	
000003	303)	
000004	254 7	JMP TEX Entry Point
000005	021)	

The stack pointer (SP) must not be in protected memory. If you desire to relocate the SP change the following locations accordingly:

- a) 000001 (SP low) and 000002 (SP high) b) 026301 (SP low) and 026302 (SP high)
- 2. System Recovery Routine -

ADRS 000070 000071	INSI		
000070	303		
000071	000		
000072	000		

- 3. Input Subroutine: Your input subroutine must begin at 000030. It should carry out the following functions:
 - a) Move an ASCII character from the input device to register A. The ASCII character should be right juetified in A with Parity bit equal to zero. Example: "B" keystroke should eet A to 102g.
 - b) Test for ESC keystroke (ASCII 177_{\odot}) and jump if true to 000000. Suggested instructions

- c) Output an echo check of the imputed character.
- d) No registers should be modified except A.
- 4. Output Subroutine: Your output subroutine ehould begin at 000050. It ehould move the ABCII character in register A to the output device. Parity bit is zore. No registers including A should be modified.
- CR-IF Subroutine: At 000020 you muet have n subroutine that will output a CR followed by a lF. Only register A may be modified.

LOADING TBX:

The octal listing of TBX is reproduced later in the text. Addressing is split octal and gives the address of the first byte of each line. An octal loader of some kind is almost a necessity. Loading by front panel switches would be a considerable chore. A Suding-type cassette is available for \$5, postpaid, from the authors. Send check or money order to: TBX Tape c/o John Arnold, Route 4, Box 52-A, Tyler TX 75701. If you are interested in a Baudot version of TBX, please inquire at the same address.

Use of a cassette tape to store TBX is virtually a necessity. Every effort has been made to protect TBX against self-destruction byt nothing is 100% sure!

The highest address available in your system for program storage must be loaded as follows:

026115 XXX₈ low part 026116 XXX₈ high part

Example: Suppose you have one 4K board: 026115 377

026116 037

EXECUTING TBX:

Simply examine 000000 and place the computer in the RUN mode. A colon indicates the system is operative.

ERROR MESSAGES

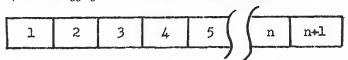
The form of error messages is: ERR α β where α is error number, and β is statement number where error was detected. Label 00000 indicates error occurred in direct execution.

ERROR NUMBER

- 1 Input line too long--exceeds 72 characters.
- 2 Numeric overflow on input.
- 3 Illegal character detected during execution.
- 4 No ending quotation mark in PR literal.
- 5 Arithmetic expression too complex.
- 6 Illegal arithmetic expression.
- 7 Label does not exist.
- 8 Division by zero not permitted.
- 9 Subroutine nesting too deep.
- 10 RET executed with no prior GOSUB
- 11 Illegal variable.
- 12 unrecognizable statement or command.
- 13 Error in use of parentheses.
- 14 Memory depletion.

EXAMPLE PROGRAM OF TBX

One example program written in TBX follows. It might assist you in debugging. A TBX line is structured as follows:



Byte No.

1 & 2 Binary value of label; most significant part in 1.

3 Length of text plus 2 in octal.

4 thru n Text of line. n + 1 CR (0158).

After the last line you should find two 377s. At the end of the example run is an octal dump of the program area of memory.

EXAMPLE PROGRAM IN TEX

020030 312 308 026 303 005 020 167 311 020040 053 094 076 077 357 303 005 020 020050 332 009 021 076 057 276 322 000 020060 021 303 371 920 000 000 000 020 020100 303 000 020 090 000 000 000 020110 000 114 123 124 040 066 060 060 020 120 042 124 020130 105 123 124 061 042 044 120 122 020 020 040 044 120 122 020 120 040 142 141 111 116 040 142 041 141 111 116 040 141 111 116 040 <t< th=""><th>OCTAL LISTING</th><th>023000</th></t<>	OCTAL LISTING	023000
021210		024210 031 315 044 023 247 311 315 071 024220 023 315 115 023 104 115 315 071 024220 023 315 115 023 104 115 315 071 024230 023 011 315 044 023 247 311 000 024240 325 006 000 315 071 023 174 267 024250 374 301 024 353 315 071 023 174 024260 267 374 301 024 315 306 024 005 024270 314 115 023 315 044 023 321 247 024300 311 004 315 115 023 311 305 104 024310 115 041 000 000 076 021 062 363 024320 033 170 037 107 171 037 117 322 024330 333 024 031 174 037 147 175 037

0260 16 026026 026036 026046 026056 026076 0261076 026116 026126 026126 026260 026360 026316	0 153 315 356 025 247 311 315 003 0 026 353 247 311 076 040 315 026 0 022 247 311 000 000 000 041 077 0 026 001 350 033 176 002 175 376 0 033 310 003 043 303 064 026 000 0 000 000 034 001 034 000 040 017 0 100 030 000 164 024 377 057 377 0 100 030 001 164 024 377 057 377 0 377 041 100 030 042 361 033 041 0 164 024 042 364 033 315 020 027 0 052 352 033 126 043 136 353 000 0 042 350 035 023 023 247 311 076 0 053 360 035 247 311 345 325 305 0 053 357 030 000 000 076 105 357 076 0 122 357 357 076 040 357 046 000 0 000 000 000 000 076 105 357 076 0 122 357 357 076 040 357 046 000 0 000 000 000 315 101 022 052 350 0 033 076 040 357 315 205 026 016 0 010 041 357 033 021 106 026 032 0 167 015 302 267 026 041 002 032 0 061 377 000 303 257 021 056 001 0 010 041 357 030 201 056 010 010 056 0 004 001 056 002 001 056 003 001 056 0 004 001 056 015 001 056 013 001 056 0 004 001 056 015 001 056 013 001 056 0 004 001 056 015 001 056 013 001 056 0 004 001 056 015 001 056 013 001 056 0 004 001 056 015 001 056 013 001 056 0 004 001 056 015 001 056 015 001 0 000 000 000 000 000 000 000 000 0	031000 052 354 033 053 104 115 052 376 031010 033 011 345 052 366 033 104 115 031020 052 352 033 011 301 315 060 051 031030 315 101 022 076 040 357 052 366 031040 033 104 115 052 354 033 053 315 031050 060 031 315 101 022 327 247 311 031060 171 225 157 170 234 147 311 052 031070 352 033 042 304 033 052 354 033 051100 042 306 033 247 311 315 165 031 042 304 033 043 043 043 076 015 043 031120 276 302 117 031 043 043 042 306 031130 031 247 311 000 315 165 031 043 043 043 043 043 043 043 043 043 043
027200 027210 027220 027220 027220 027226 027256 027256 027350 027336 02736 027	315 044 023 315 240 024 315 200 0 024 303 146 027 032 023 376 040 0 312 214 027 033 306 300 320 007 0 117 023 032 376 050 312 243 027 0 033 247 311 151 046 024 116 043 0 146 151 116 043 106 043 315 044 0 023 140 151 042 370 033 067 311 0 300 325 342 000 000 000 000 000 0 000 000 000 000	032200

the digital group

po hox 6528, denver, colorado 80206

December 14, 1975

Mr. Bob Albrecht & Bernard Greening People's Computer Company People's Computer Com PO Box 310 Menlo Park, CA 94025

Dear Bob and Bernard.

I am very interested in helping out with your Tiny BASIC (perhaps Micro BASIC might be more appropriate). Since my specialty is Hardware and the lowest level Software to interface this hardware to a system, I would like to suggest a simple hardware subsystem.

A scientific calculator IC can be easily interfaced to a micropro-cessor to provide all of the various mathematical operations very accurately with minimal software overhead. I am including a copy of some of the scientific calculator documentation out by the of some of the Digital Group.

This scientific calculator has been interfaced to an 8008 (Mark-8 modified) and MOS Technology 6501/2 system. The software can be easily modified to support an 8080 or 6800, thereby providing an eany access to building "Tiny BASIC" for 8008, 8080, 6800, 6501 or 6502 systems.

The major drawback of a calculator chip for math routines is that it is very slow compared to specialized hardware and software systems. The major advantages are:

- Low software overhead (about 300 bytes for interfacing)
 Low cost (around \$45 worth of parts 6 PC board)
 Quick way to develop Math routines with high accuracy.

I would be happy to assist PCC in developing Tiny BASIC using these Scientific Calculator IC's.

Dr. Robert Suding

c/o The Digital Group

SCIENTIFIC CALCULATOR

Here is a calculator circuit designed to be used with any computer of 8 bits or more capacity. I am presently using it with an 8008 system, approximately 300 bytes of storage being required to basically interface this circuit to my TV readout and keyboard. Only one 8-bit input port and one 8-bit output port is required.

The heart of the circuit is the 2529-103 calculator IC from Mos Technology. This is a simple IC which gives trig, log, memory, square root, etc., functions. The display is normally a 12-digit LED 7-sequent assembly. The sequent drivers are built into the 2529. The 12-digit outputs are usually fed to a pair of 75492's which serially scan each of the 12 digits at about a 60Hz cycle rate from an internal clock. A matrixed keybcard is normally attached between the 12 digit outputs of 2529 and 4 Keybcard inputs of the 2529, giving a potential 48-key input capability, 41 of which are actually used.

The design required efficient handling of the 12-digit outputs. Since it was necessary to utilize the digit outputs for both data entry and digit segment output, the design was centered on a controlled accessing of the asyncronously scanning 12 digits. The computer has 4 bits of an output port assigned to the duty of selecting a given digit by sending its binary equivalent to the inputs of a 74150 sixteen input selector. When the selected digit becomes present the output at pin 10 of the 74150 goes low as long as the digit is present. By combining this input with three more bits from the computer, the desired "keyboard" input is sent to the 2529. The computer word should be held for at least 40 ms to be certain that the asyncronously scanning digit has been accessed.

Likewise, the digit output must identify the digit to which the current segment outputs apply. By using the same coding scheme for the four inputs to the 74150, a computer controlled sampling system is established. The MSB output from the computer informs the calculator interface that a digit/segment output is desired. When the desired digit finally ripples by, a strobed MSB pulse appears on the interface output. This pulse then interrupts the computer to inform it that the segment data for the desired digit is present and valid as long as the MSB stays +.

Several considerations: First, only 5 of the 7 segments are needed to decode 0 through 9, minus, blank, and the error signs. Each digit may also have a decimal point attached to it, so the output becomes 6 bits, plus the MSB strobe bit. Be aware that these calculator chips are quite slow. When entering a data item or especially a function, the display will go blank up to 1/3 second while internal processing takes place. The result can take on any number of digits, but digit 9 is always used. By sampling for "digit 9 not blank," the end of internal processing can be detected. When this occurs, either further entries, or sampling of all 12 digits may proceed without data loss. 8008 programs have been written to handle simple keyboard entry and tv result display, and interactive calculation operations involving messages anf formula building and reiteration. These are available through the Digital Group.

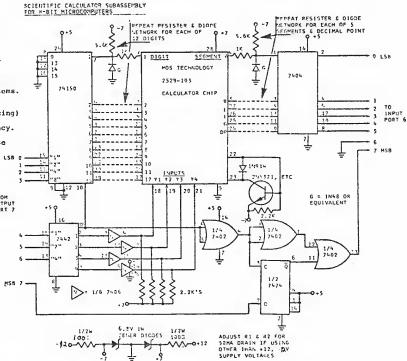
The 2529 is available from Mos Technology at \$27.50 apiece. Some newer scientific calculator chips have been announced by Mos Technology and are being presently sampled.

Other calculator chips could be used in similar circuits. However,

I would question the advisability of using these simpler chips with their much lower calculation power return. Mos Technology also makes an RPN format calculator IC, the 2529-106 for H.P. buffs. A metric conversion chip (2529-104) is also available from Mos Technology. These IC's have been tried in the circuit. They are directly usable in the enclosed circuit.

The basic functions are roughly equivalent to the TI SR-50, but the enhanced software version will be considerably better than the HP-65 programmable calculator due to its message display capacity and "almost" unlimited memory capacity.

-Dr Robert Suding WOLMD



						00
IN	PUT COD	ES POR	R FUNCTION E	NTRY		DIGIT ALONE
CTION	OCTAL	HEX	FUNCTION	OCTAL	HEX	0
0	021	11	ARC	033	18	1 2 3 4 5 6
1	022	12	SIN	061	31	-
2	023	13	cos	062	32	3
	024	1.4	TAN	063	33	9
4	025	15	LN	064	34	,
3 4 5 6 7	026	16	LOG	065	35	.7
6	027	17	RCL	067	37	
7	030	18	Σ	070	38	8
8	031	19	x+-y	071	39	9
9	032	1A	DGR	072	3A	-
	041	21	STO	073	38	ERROR
÷	042	22	CA/CE	074	3C	
	043	23	CHS	053	2B	D11-
×	044	24	-	054	2C	Blank
÷	045	25	ŽX.	046	26	
-	047	27	Ĩ	050	28	447 - 64
7	052	2A)	051	29	*fleft
_	066	36	10×	103	43	M INP
1/x	101	41	ex	104	44	COD
x2	102	42	N:	105	45	COL
NO Op	000	00	Restore-	034	10	* SEN
			Display			FOR

NO OP	000	••	Display	7	10	
INP	UT CODES	FOR	DIGIT DATA	REQUEST		
DIGIT	OCTAL	HEX	DIGIT	OCTAL	HEX	
1	201	81	7	207	87	
2	202	82	8	210	88	
2	203	83	9	211	89	
4	204	84	10	212	8A	
5	205	85	11	213	88	
6	206	86	12	214	8C 88	3

	SEGMENT	

	DIGIT ALONE	OCTAL	HEX	DIGIT AND DECIMAL PT	OCTAL	HEX
EΧ	0	260	во	0.	220	90
	1	275	BD	1.	235	9 D
В	ž	250	A8	2.	210	88
1	3	254	AC	3.	214	8C
2	4	245	AS	4.	205	85
2 3 4	Š	246	A6	5.	206	86
	6	243	A3	6.	203	83
S	.7	274	8C	7.	234	9C
7	8	240	A0	8.	200	80
8	9	244	λ4	9.	204	84
9	_	257	AF		217	8 F
A	ERROR*	262	B2	ERROR.*	222	92
В		or	or	2,4,0,1,0	or	or
С		242	A2		202	82
B C 6	Blank	277	BF	Blank.	237	9F
8 9	*flefts	nost di	git on:	ly)		

- IPUT CODES SENT TO ENTER DIGITS AND FUNCTIONS. DE MUST 8E HELD MORE THAN 40MS.
- SEND DIGIT 9 DATA REQUEST (211 OR 89) AND WAIT FOR MSB FROM 7402 TO GO +. THIS INDICATES INTERNAL CALCULATIONS FINISHED.
- SEND DIGITS 12 THROUGH I DATA REQUESTS, DECODE EACH WITH SEGMENT DECODE DATA TABLE AS DATA AVAILABLE MSB LINE FROM 7402 GOES + FOR EACH
- LONGEST CALCULATION DELAY APPEARS TO BE 69: (ABOUT 1/3 SEC).
- RPN (2529-106) AND METRIC-CONVERSION (2529-104) CALCULATOR IC'S FROM MOS TECHNOLOGY WILL WORK WITH THE SAME CIRCUIT.

DR. ROBERT SUDING WOLMD

SNOBOL FOR THE ALTAIR

Dear Dragons,

Thanks for the great publication and other nice things--like dragon shirts!. What a way to learn.

I have a problem. Without considering any possible consequences, I have committed myself to writing a SNOBOL Compiler (interpreter?) for an Altair 8800. My officemate has built the Altair for the college at which he teaches, and after many months of promising some kind of assistance, I finally offered to write a compiler.

To get to the point: does anyone out there have any experience in compiler writing, particularly in SNOBOL compiler writing? I know that some of the sharpest people in this field read *PCC*, so I'm really hoping to hear from someone.

Of course, once I get the compiler working, I will make it available to other Altair owners and users (for a nominal fee and a lot of glory).

(I realize all you people are heavily into BASIC, but SNOBOL is a pretty neat language for things like compiler writing, natural language translation, and general string manipulation.)

Also, since my friend's Altair is 75 miles away from my home, donations of Altairs will be accepted.

MAUREEN SUPPLE 828 S. Irving St Arlington VA 22204

(SNOBOL compilers are tough. An interpreter would be easier. A good place to start looking for information would be Griswold's book, *The Macro Implementation of SNOBOL*, W.H. Freeman, San Francisco, 1973; and Waite's book, *Implementing Software for Non-Numeric Applications*, Prentice Hall, Englewood Cliffs, New Jersey, 1971.)

FULL OF HOLES

I guess you know, Tiny BASIC as presented in its first chapter is full of holes. Look, for example, at what happens if you try to evaluate an expression without aunary plus or minus on the front. Ich. Also, I wonder if the interpreted interpreting interpreter interpreter executor is viable for a really small, slow system like an 8008 system. Talk about crunching! Anyway, I want to see more. I'm crazy, maybe? Who cares.

Sincerely,

FRITZ ROTH Rt 7

Carbondale IL 62901

A HIGH ORDER

Dear Bob Albrecht, I am writing this letter about many things I've read about in PCC. The Tiny BASIC project looks like something everyone would like to tackle. The interpreter idea is a little costly on time and storage, unless you plan to use it on many systems. Otherwise, it's a good idea. I'm interested in simulating languages using BASIC or FORTRAN as the "machine," so this type of thing is interesting. If only someone had the plans for ALGOL in IL . . .

If anyone has done any projects simulating languages/computers in a high order language, would they please contact me?!

Thanks for everything, *PCC*! Respectfully,

REED CHRISTIANSEN 2756 Fernwood No. Roseville MN 55113

TB CODE SHEET

by Dick Whipple

You may be interested in knowing that John Arnold and I write our programs (like TB) in machine language. We have found it to be less restrictive and more versatile although not having a source file of some kind is a disadvantage. We do keep a hand-generated source listing on coding sheets for our reference. A major program like TB requires a two-pass development: the first pass ends up with lots of "fixes" and "patches" to get the program to work; the second pass is then used to clean-up the mess produced in pass one. The coding sheets from pass two represent the nearest thing to source code we have. For your reference I have included a copy of one of our coding sheets from TB. The addresses are split octal.

Program Plane: Tiny BASIC (Rev 1) PCC Drages Copy B, RBW Date 11/10/25 H= 020 BUFFIN: 1 = 040 376 } CPI <?>+040, 000 0407 01 111 EXI H BUFSTRT 02 020) 42 3127 JPZ RUBOUT 03 10017 43 101 GY O40 LXI C: CASE 44 020 CONTY: 05 337 RST IN SP: 045 167 A -M 46 [043 INX H CONT 3: 007 3767 CPI (FGS) 47 303 IMP CONTY 50 006 51 020 11 3127 JPZ FGS 12 0527 FGS: 032 016 } mus c 046, 13 [020] 14 3767 15 037 54 3037 (PI (LTR) TIMP CONTY 55 006 55 020 16 3127 17 057 7 JPR LTR LTR: 0.7 016 3 MUI C 0, 20 [020] 61 337 RST IN 21 005 DIR B 22 312) 23 306 } JPZ ERR306 62 376 CPI (LTR) 63 037 24 026) 84 3027 25 376 CPI (CR) 65 1007 Y JNZ CONT3 26 010) 66 020 67 327 RST CRLF 27 3127 30 107 JPZ END GETLINE:070 076 } mus A <FGS> 31 050 32 13767 72 357 RST OUT CPI (SP) 33 004 73 (0767 MUE A (:> 351312 35 020 RST OUT 75 357 JP7 5F 75 303 37 201 ADR C JMP BUFFIN 77 000

Are you implementing Tiny BASIC or some other software. Let us know and we'll let others know. Let's stand on each others shoulders and not on each others toes (to paraphrase C. Strachey).