Communication Systems Laboratory Lab 4: Source Coding

(Report Due: 23:00, November 2, 2022)

1 Overview

After the journey of quantum information processing, now we are back to the fully classical (digital) zone (home sweet home?). In this lab, we are going to implement source coding for discrete memoryless sources. Here, the sources are classical information, which means that one can designate each possible outcome of such sources (e.g. it could be something like '010', •, or ©, etc.) with a probability value¹. Your goal in this lab is to implement a classical lossless compression scheme by e.g. Matlab to shown that we can indeed remove the redundancy hidden in the discrete sources.

2 Experiments

We are going to explore the Huffman coding from three perspectives. First, beginning with a toy example in Table 1, we analyze the entropy, the Huffman tree, the Huffman dictionary, and the average codeword length by hand. Next we implement three main blocks for Huffman coding: huffman_dict for building the Huffman dictionary from the probability distribution, huffman_enc for Huffman encoding, and huffman_dec for Huffman decoding. With these functions, we study the average codeword length of Huffman coding through Monte—Carlo simulations.

- 1. **(35 points) Information theory and Huffman coding:** First let us consider a random symbol *X* whose outcomes and the associated probabilities been given in Table 1. Show your answers to the following problems as what we did in the lecture.
 - (a) (4 points) Calculate the entropy of X, i.e. H[X].
 - (b) (10 points) Construct the Huffman tree and the Huffman dictionary for X.
 - (c) (4 points) Verify whether the codewords constructed by your Huffman tree satisfy the Kraft inequality or not.
 - (d) (4 points) Find the average codeword length \overline{L} for the dictionaries in Problem 1b. Do they satisfy the source-coding theorem?
 - (e) (4 points) Encode the sequence of symbols in (1) using the Huffman tree in Problem 1b.

$$\{g,a,c,a,b\}. \tag{1}$$

(f) (4 points) Decode the bitstream in Problem 1e using the Huffman tree in Problem 1b.

 $^{^{1}}$ At this point, I believe that you couldn't help but think about the scenarios that each outcome of the sources could be something like $|010\rangle$ (endowed with a probability value). Yes, the underlying source is now *quantum information*. If you really love it so much, please go ahead to explore *Schumacher's noiseless compression of quantum bits*.

- (g) (5 points) Let T_{ε}^n denote the typical set of X with $\varepsilon = 0.1$ and n = 10. Find 10 members in the set T_{ε}^n .
- 2. (35 points) Implementation of Huffman coding: In what follows, we implement the three main blocks of Huffman coding using Matlab.
 - (a) (15 points) Implement a Matlab function to construct Huffman dictionary with the following arguments:

```
dict = huffman_dict(symbols, prob)
```

where symbols is a vector consisting of the names of the symbols. The second argument prob is a vector containing the probabilities of these symbols. The output dict is a *cell* array with the following specifications

- Each row in dict corresponds to a node in the Huffman tree. The index of this node is row number.
- The first column of dict is the name of a node in the Huffman tree.
- The second column of dict corresponds to its probability.
- The third and the fourth columns of dict are the indices for the left child and the right child of this node.
- The fifth column of dict has the bits.

For the Huffman tree example given in class, the following sample codes are for your reference.

```
1 >> symbols = { 's0', 's1', 's2', 's3', 's4' };
2 >> prob = [ 0.26, 0.25, 0.20, 0.15, 0.14 ];
3 >> dict = huffman_dict( symbols, prob )
4
5 dict =
```

Table 1: A table of symbols in \mathcal{X} and their probabilities.

Symbol	Probability
a	0.2
b	0.05
c	0.005
d	0.2
e	0.3
f	0.05
g	0.045
h	0.15

```
9×5 cell array
7
8
       {'s0'
                   }
                        {[0.2600]}
                                     {0×0 double}
                                                    {0×0 double}
                                                                   {'01'
                                                                            }
       {'s1'
                        {[0.2500]}
                                                                   {'10'
                                                                            }
                    }
                                     {0×0 double} {0×0 double}
10
                   }
                      {[0.2000]}
                                     {0×0 double} {0×0 double}
                                                                   {'11'
        {'s2'
                                                                            }
11
       {'s3'
                   }
                      \{[0.1500]\} \{0\times0\ double\} \{0\times0\ double\}
                                                                   {'000'
                                                                            }
12
       {'s4'
                  } {[0.1400]}
                                     {0×0 double}
                                                    {0×0 double}
                                                                            }
                                                                   {'001'
13
       {'s3s4'
                   }
                        {[0.2900]}
                                     ]}
                                              4]}
                                                    {[
                                                             5]}
                                                                   {'00'
                                                                            }
14
                                                                   {'1'
15
       {'s1s2'
                  }
                      {[0.4500]}
                                     {[
                                              2]}
                                                    {[
                                                             3]}
                                                                            }
       {'s3s4s0' }
                        {[0.5500]}
                                     ]}
                                              6]}
                                                    ]}
                                                             1]}
                                                                   10'}
                                                                            }
16
       {'s3s4s0s1s2'}
                                     ]}
                                              8]}
                                                    ]}
                                                             7]}
                                                                   {0×0 char}
```

For example, the 8th row of dict (Line 16 in the above block) represents a node with name 's3s4s0', probability 0.55, left child with index 6 (node name 's3s4'), right child with index 1 (node name 's0'), and the bits '0'.

Print the Hufman dictionary dict associated with the symbols and the probabilities in Table 1.

(b) (10 points) Implement a Huffman encoding function which has the following arguments:

```
bin_seq = huffman_enc(sym_seq, dict);
```

where sym_seq is a sequence of symbols, dict is a Huffman dictionary with the same format in Problem 2a, and bin_seq is the corresponding bit strings for the encoded symbols.

Verify your results with Problem 1e.

(c) (10 points) Implement a Huffman decoder function which has the following arguments:

```
sym_seq = huffman_dec(bin_seq, dict);
```

where bin_seq is a sequence of bits after Huffman encoding, and sym_seq represents the decoded symbols.

Verify your results with Problem 1f.

Note: In addition to the verification in Problems 2a, 2b, and 2c, we will verify your implementation with extra test cases, which are not available to you.

3. (30 points) The average codeword length of Huffman coding: In this problem, we study the average codeword length of Huffman codes through the Monte–Carlo simulation. We consider the alphabet $\mathcal{X} = \{a, b, \dots, h\}$ in Table 1.

Warning: We suggest that you use huffman_dict and huffman_enc in Problem 2. If you use the built-in Matlab functions huffmandict and huffmanenco for Huffman coding, then the grade becomes $\times 50\%$ for Problem 3.

(a) (5 points) Generate a sequence of n = 10 symbols according to the probability in Table 1. List the symbol sequence and Huffman-encoded binary data, and find the length of the data (in bits) in your report.

(b) (5 points) Repeat the experiment in Problem 3a multiple times (say R = 200 times). Each experiment randomly generates a sequence of symbols. Denote the length of the encoded binary data in the r-th experiment by $L_n^{(r)}$. Plot the histogram of $L_n^{(1)}, L_n^{(2)}, \ldots, L_n^{(R)}$ and indicate the mean (denoted by $L_n(R)$) in the title according to

$$L_n(R) := \frac{1}{R} \sum_{r=1}^{R} L_n^{(r)}.$$
 (2)

You may use the Matlab command histogram for this problem. The parameter R is also known as the number of the Monte–Carlo simulations.

(c) (10 points) Based on R experiments, the average codeword length $\overline{L}_n(R)$ is defined as

$$\overline{L}_n(R) := \frac{1}{n} L_n(R). \tag{3}$$

Next we generate a plot with the following specifications.

• The horizontal axis corresponds to the number of the Monte-Carlo runs with

$$R = 10, 20, 50, 100, 200, 500, 1000,$$

in the logarithm scale.

- Two horizontal lines show the entropy H[X] and the average codeword length \overline{L} in Problems 1a and 1d.
- Three curves illustrate the *experimental* average codeword length $\overline{L}_n(R)$ for n = 10, 50, and 100.
- You can use Matlab commands semilogx, xlabel, ylabel, title, grid, legend for your plot. We suggest that you present these five curves with proper markers, line styles, and line colors.
- (d) (10 points) Comment on the plot in Problem 3c as much as possible.
- 4. (**Bonus 5 points**) Consider the *identical and independently-distributed* extension of the symbols X to $X_1 \times X_2 \times X_3$. Now do Problems 2 and 3 again ②. Note that you have to normalize the average length \bar{L} and the entropy $H[X_1X_2X_3]$ by factor 3. What do you observe?

3 Lab Report

Submission format: 請同學將實驗報告 (.pdf)、各題的程式碼 (.ipynb, .py, .m,.cpp) 以及題目要求的其他檔案 (如圖檔) 壓縮成一個.zip 檔,並將此壓縮檔上傳至作業繳交區。繳交的壓縮檔需以此格式命名: 學號 _ 姓名 _ Lab2.zip (Ex: b09900000 _ 王小明 _ Lab1.zip)。

There is no format/typesetting requirements for your lab report, but you have to make your report decent and looking nice. In the report, you should address the results of the exercises mentioned above. You should also include your simulation program in the appendix of the report. Include whatever discussions about the new findings during the lab exercise, or the problems encountered

and how are those solved. Please properly cite the literature if you referred to. Do not limit yourself to the exercises specified here. You are highly encouraged to play around with your simulation program on self-initiated extra lab exercises/discussions.

It is highly recommended to use Matlab to implement the Huffman coding. However, you are also free to use other software programming languages. For more references and instructions of Matlab, there are a bunch of online references such as:

- https://www.mccormick.northwestern.edu/documents/students/undergraduate/introduction-to-matlab.pdf
- https://blogs.mathworks.com/loren/2007/12/11/making-pretty-graphs/
- https://blogs.mathworks.com/loren/2007/12/11/making-pretty-graphs/
- https://web.eecs.umich.edu/~aey/eecs451/matlab.pdf
- http://staff.www.ltu.se/~damvar/matlab.html