Parallel Query Service for Object-centric Data Management Systems

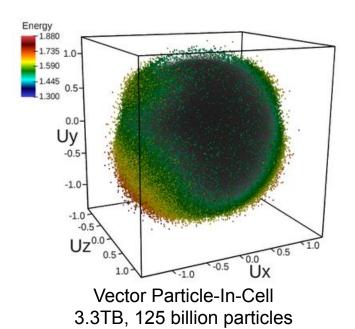
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Extract a **small fraction** of information from a **large amount** of data.



Baryon Oscillation Spectroscopic Survey 3.2 TB, 25 million objects





- DBMS, e.g. BerkeleyDB, PostgreSQL, MongoDB...
 - Efficient metadata queries.
 - Not optimized for multi-dimensional data queries.
- Multi-dimensional data indexing/querying system, e.g. SciDB, FastQuery
 - Targets large n-dimensional arrays, lack support for metadata queries.
 - Reading data may lead to significant overhead.

A unified data and metadata query system that provides elastic, efficient, and scalable query evaluations.



Current Data Management Systems

Hardware

Memory

Node-local storage

Shared burst buffer

Disk-based storage

Campaign storage

Archival storage (HPSS tape)

Software

High-level lib (HDF5, etc.)

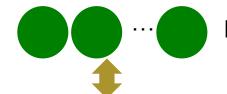
IO middleware (POSIX, MPI-IO)

IO forwarding

Parallel file systems

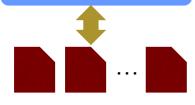
Usage

Applications



Data (in memory)

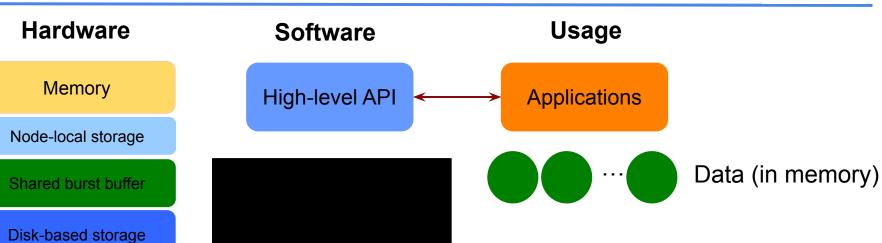
IO software



Files in file system



Object-centric Data Management Systems



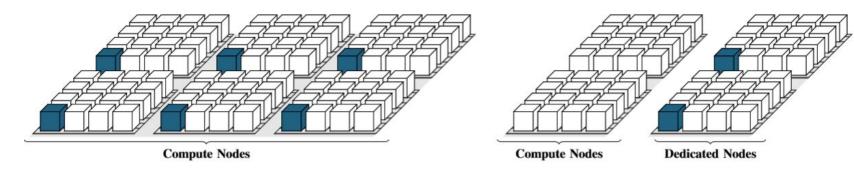
Campaign storage

Archival storage (HPSS tape)



Data management in PDC

- PDC servers run in background, manages data and metadata.
- Data objects can be stored on different layers of memory hierarchy.
- Large data objects are decomposed into smaller regions.
- Metadata is cached in server's memory and persisted to storage.
- Application send requests through linked PDC client library.



Queries in PDC

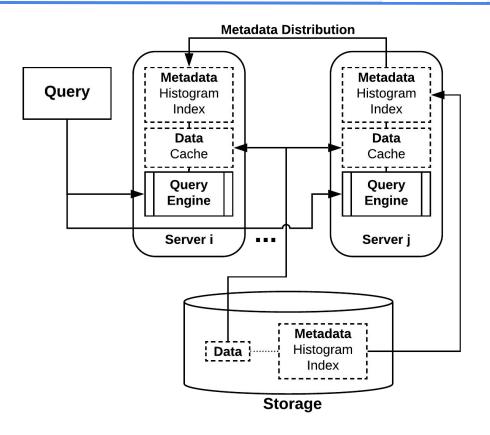


Metadata query

Previous paper: "SoMeta:
 Scalable Object-Centric Metadata
 Management for High
 Performance Computing"

Data query

- Single variable
- Multi variable
- Get number of hits
- Get selection
- Get value





PDC-query Interface

```
// Create a one-sided data query
pdcquery t *PDCquery_create(pdcid_t obj_id, pdcquery_op_t op, pdctype_t type,
void *value);
// Combine queries
pdcquery t *PDCquery_and(pdcquery_t *query1, pdcquery_t *query2);
pdcquery t *PDCquery or (pdcquery t *query1, pdcquery t *query2);
// Set query region constraint
perr t PDCquery set region (pdcquery t *query, pdcregion t *region);
// Query operations
perr_t PDCquery_get_nhits(pdcquery_t *query, uint64_t *n);
perr t PDCquery get selection (pdcquery t *query, pdcselection t *sel);
perr_t PDCquery_get_data(pdcid_t obj_id, pdcselection_t *sel, void *data);
perr t PDCquery get data batch (pdcid t obj id, pdcselection t *sel, uint64 t
batch size, void *data);
pdchistogram t * PDCquery_get_histogram(pdcid_t obj_id);
```





Full scan

- Straightforward parallel implementation.
- Go over all elements and check against query condition.
- Slow for single variable and simple query condition.

Data reorganization w/ sorting

- Requires data preparation, extra storage.
- Eliminates the need to go through all elements.
- Best performance for single variable query.

Bitmap index

- Requires index building in advance.
- Go through index instead of data.
- o Best performance if actual values are not required.

Optimization?



Full scan

Skip the inspection of some amount of data?

Data reorganization

Speedup the evaluation process for multivariate query conditions?

Index

- Skip the evaluation of some indexes?
- Evaluate the highly selective variable first?

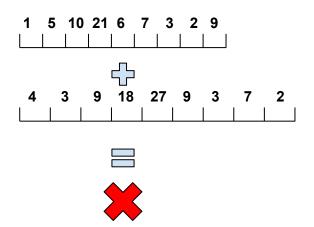
Histogram

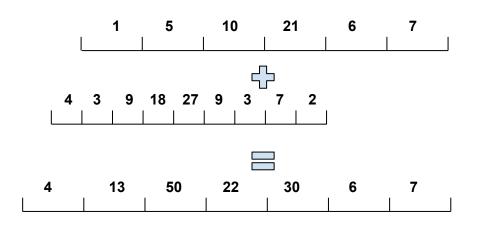


- Generate a histogram for each PDC region
 - Done at data creation time or during server "free" time asynchronously
- Use histogram to get max/min value of a region
- Use histogram to estimate the selectivity of each variable
 - Re-order the query evaluation, prune as many regions as possible.
- Generating a global one is costly, and needs coordination for updates.
 - Can we generate local region-specific histograms that can be easily merged into a global one?

Mergeable Histogram







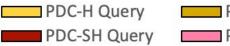
The bin width of different histograms must be same or divisible.

Use random sampling to get approximate min/max and make them aligned with bin boundaries of other histograms.

Both use values from pre-defined sets, 2^n and $N \pm 2^n$.



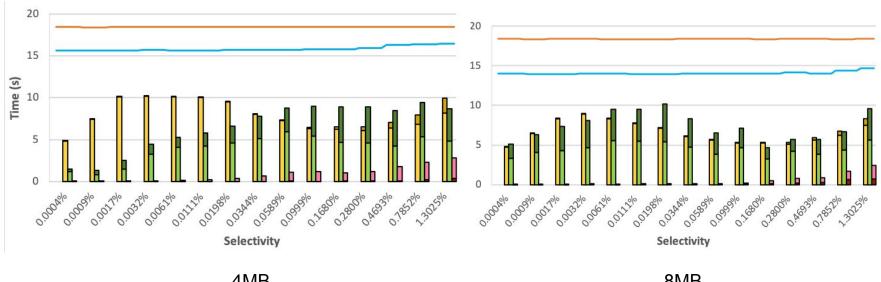
Region Size







H: histogram I: index S: sort F: full scan



4MB

8MB

Region Size

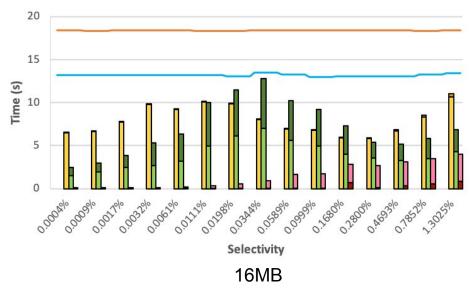


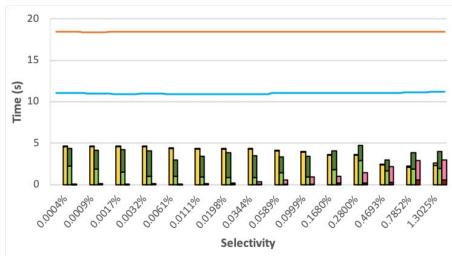












32MB

Region Size

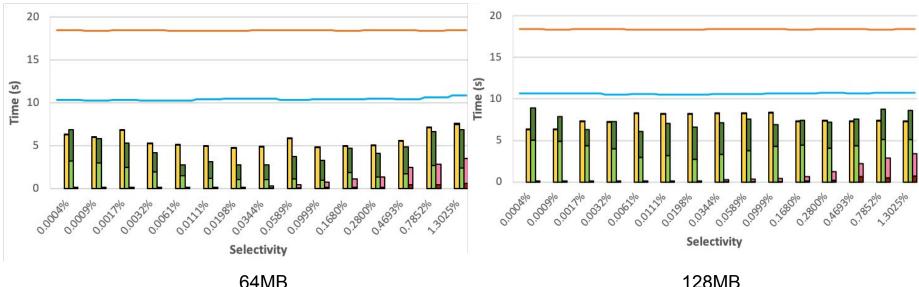












64MB 128MB



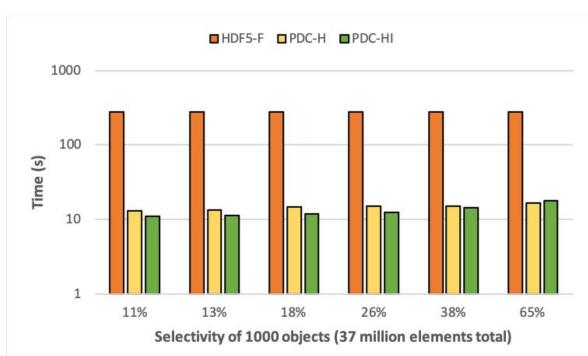
Results - Multivariate Query



PDC-H: PDC with **H**istogram only, **PDC-HI**: PDC with **H**istogram and Fastbit Index, **PDC-SH**: PDC with **S**orted data (sorted by the 'energy' object) and **H**istogram. **HDF5-F**: amortized time to evaluate the 6 queries with HDF5 **F**ull scan. **PDC-F**: amortized time to evaluate the 6 queries with PDC **F**ull scan.

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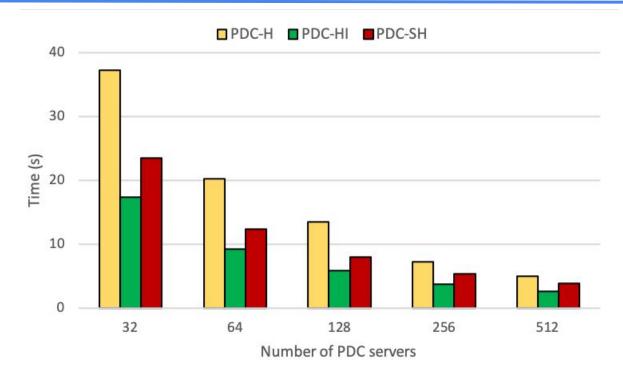
Results - Metadata + Data Queries



Comparison of queries with both metadata (fixed selectivity on 1000 objects) and data conditions (varied selectivity from 11% to 65%) on the H5BOSS dataset.



Results - Multivariate Scaling



Query time comparison for a multi-object query condition with 0:011% selectivity using different number of PDC servers.

Conclusion



- Data querying is a crucial tool for efficient information retrieval that enhances scientific productivity
- PDC-query provides a highly efficient and scalable query service
 - Designed for object-centric data management systems with simple APIs
 - Novel optimizations using mergeable histograms on top of existing approaches such as data reorganization and indexing.
 - Single variable queries on sorted data have the best performance, index with histogram good if not retrieving values.
 - Multivariate queries with indexes or histograms have similar performance when data needs to be retrieved.

Thanks!

Questions?