

Advanced System Software

(先端システムソフトウェア)

#10 (2018/11/15)

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Agenda

- Advanced Topics
 - Functional Reactive Programming for Embedded Systems

Time-Varying Values

- $TVV\ \alpha = Time \rightarrow \alpha$
 - Function from time to data of type α
 - First class entities representing (continuously) changing data (of type α) over time
 - ex) `tmp :: TVV float` (temperature sensor value)
 - aka. Signal α , Behavior α
- $Event\ \alpha = [(Time, \alpha)]$
 - List of pairs of time and data of type α
 - First class entities representing discrete events of type α
 - ex) `key :: Event Char` (keyboard event)

Time/Space-Leaks

- Unrestricted access to a time varying values leads to time/space-leaks
- $x :: TVV\ a$
- The value of x at time t (i.e., $x\ t$) depends on the values of x in the interval $[0, t)$.
- e.g., If $x\ t$ is allowed to be accessed and
 - If $x\ t$ is always calculated using $\{ x\ t' \mid t' \in [0, t) \}$, it takes unexpectedly long time (time-leak).
 - If $\{ x\ t' \mid t' \in [0, t) \}$ is stored to calculate $x\ t$, it requires unexpectedly large memory (space-leak)

To Avoid Time/Space-Leaks

- Restricted representation of TVV α
 - TVV α as a (primitive) type other than Time $\rightarrow \alpha$
 - ex) Emfrp
 - Nodes (TVVs in Emfrp) are not first class values
 - Only @last can be used to access the past value of a node
- Arrowization
 - Abandon time-varying values themselves and use "functions" on time-varying values
 - Arrow is used to construct the "functions"
 - ex) Yampa (FRP library for Haskell)

Signal Functions (Signal Transformers)

- $SF\ a\ b \approx Signal\ a \rightarrow Signal\ b$
 - Note: Here, $Signal\ a$ is the type of time-varying values (corresponds to $TVV\ a$)
- $SF\ a\ b$ is an abstract type
 - Operations on it provide a *disciplined* way to compose signals
 - The notion of *arrow* provides the discipline.

Arrow Operations (1)

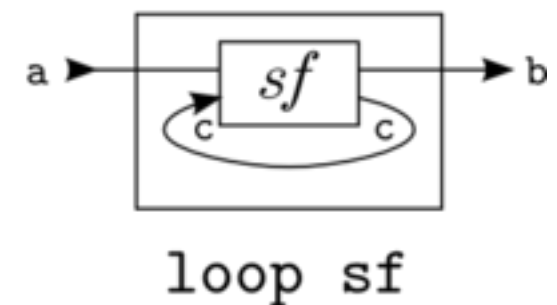
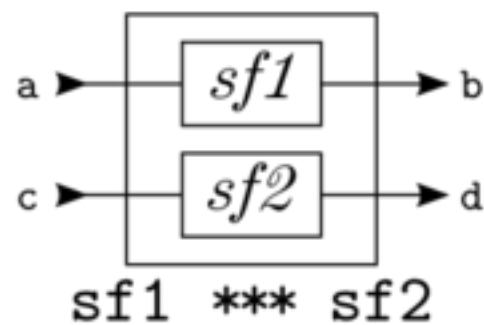
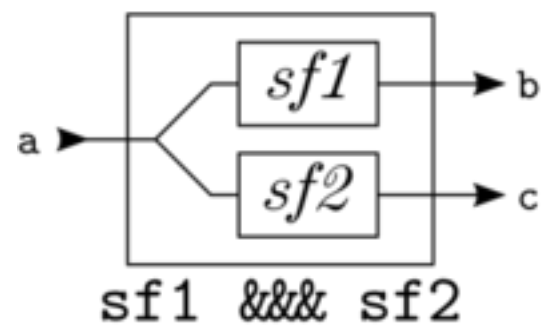
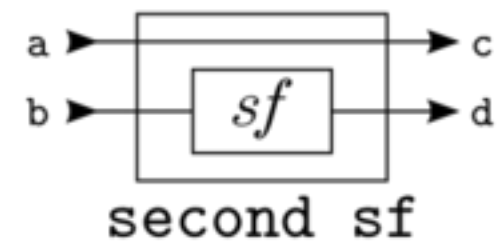
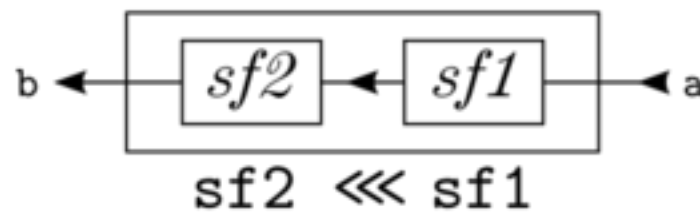
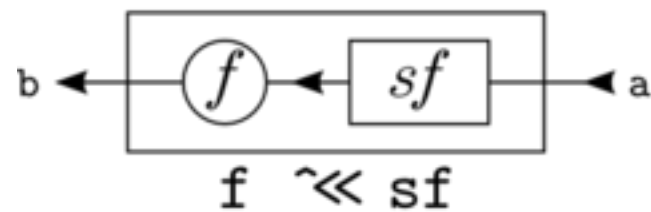
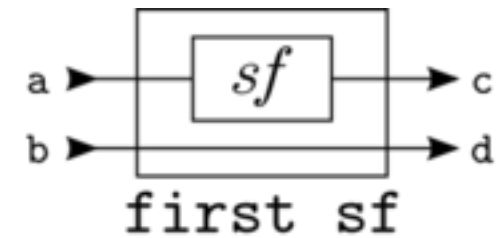
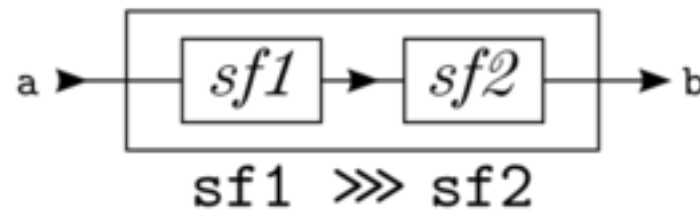
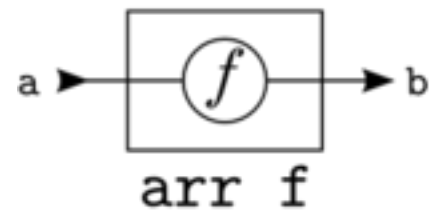
- Lifting
 - $\text{arr} :: (a \rightarrow b) \rightarrow \text{SF } a \ b$
- Composition
 - $(>>>) :: \text{SF } a \ b \rightarrow \text{SF } b \ c \rightarrow \text{SF } a \ c$
 - Note: $f >>> g = (>>>) f \ g$
- The following equations hold
 - $\text{arr } f >>> \text{arr } g = \text{arr } (g \cdot f)$
 - Note: $(g \cdot f) \ x = g \ (f \ x)$
 - $(f >> g) >> h = f >> (g >> h)$

Arrow Operations (2)

- $(\&\&\&) :: SF\ a\ b \rightarrow SF\ a\ c \rightarrow SF\ a\ (b, c)$
 - $arr\ f\ \&\&\&\ arr\ g = arr\ (f\ \&\ g)$
 - where $(f\ \&\ g)\ x = (f\ x, g\ x)$
- $(***) :: SF\ a\ b \rightarrow SF\ c\ d \rightarrow SF\ (a, c)\ (b, d)$
 - $f\ ***\ g = (arr\ fst\ >>>\ f)\ \&\&\&\ (arr\ snd\ >>>\ g)$
 - where $fst\ (x, y) = x$, $snd\ (x, y) = y$
- $first :: SF\ a\ b \rightarrow SF\ (a, c)\ (b, c)$
 - $first\ f = f\ ***\ arr\ id$
 - where $id\ x = x$
- $second :: SF\ a\ b \rightarrow SF\ (c, a)\ (c, b)$
 - $second\ f = arr\ id\ ***\ f$

Arrow Operations Visualized

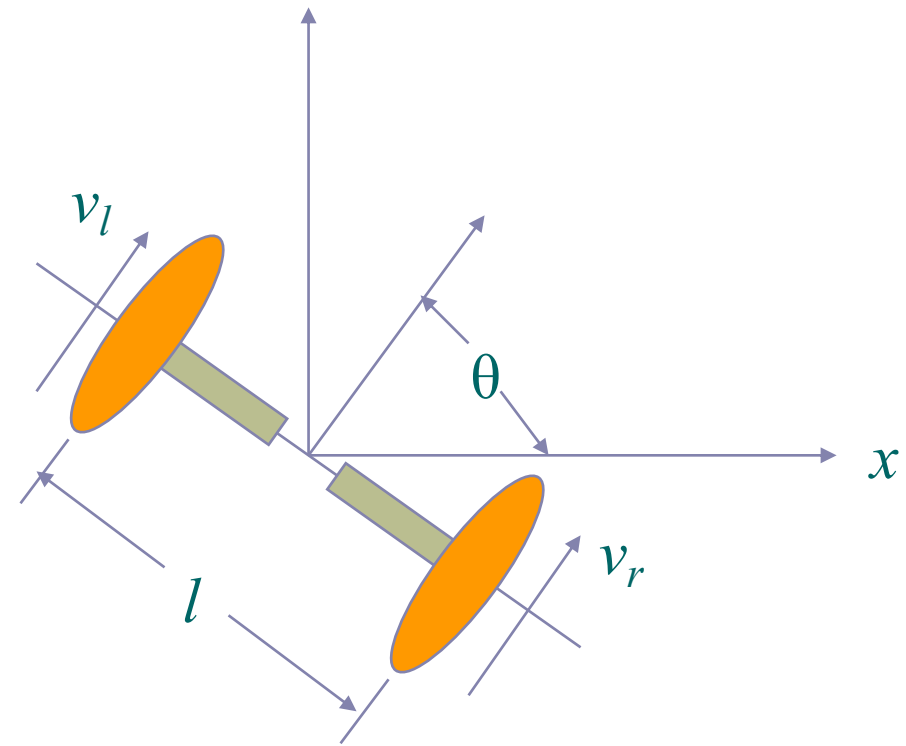
<https://wiki.haskell.org/Yampa>



Example

- Calculate the x position of a two-wheel robot

$$x(t) = \int_0^t \frac{vr(t) + vl(t)}{2} \cos \theta(t) dt$$



<https://wiki.haskell.org/Yampa>

```
vrSF, vlSF, thetaSF :: SF Input Float
```

```
x :: SF Input Float
```

```
x = let v = (vrSF &&& vlSF) >>> arr2 (+)
```

```
      t = thetaSF >>> arr cos
```

```
    in (v &&& t) >>> arr2 (*) >>> integral >>> arr (/ 2)
```

```
arr2 :: (a -> b -> c) -> SF (a, b) c
```

```
arr2 = arr . uncurry
```

```
uncurry :: (a -> b -> c) -> (a, b) -> c
```

```
uncurry f (x, y) = f x y
```

Using Arrow Syntax

```
vrSF, vlSF, thetaSF :: SF Input Float

x :: SF Input Float
x = proc input -> do
    vr <- vrSF -< input
    vl <- vlSF -< input
    theta <- thetaSF -< input
    i <- integral -< (vr + vl) * cos theta
    returnA -< (i / 2)
```

Example in Emfrp

```
module Robot
in   vl, vr, theta : Float,
      t : Int
out x : Float
use Std

node intv = (vl + vr) * cos(theta) / 2.0

node init[0] x = x@last +
                  (intv - intv@last) * (t - t@last)
```

Procedural Abstraction

- Synchronous (reactive) programming languages
 - Lustre, Esterel
 - Ceu
- Provides procedural abstractions for reactive event-based behaviors

Ceu

- Prints "Hello World!" every 250ms.

```
loop do
  await 250ms;
  _printf("Hello World!\n");
end
```

<http://www.ceu-lang.org>

- **await**
 - stops the execution until specified event (in this example, the expiration of 250ms timer) occurs

Input Events

```
input int MY_EVT;  
  
loop do  
    var int v = await MY_EVT;  
    _printf("MY_EVT=%d\n", v);  
    if v == 0 then  
        break; // escapes the loop when v==0  
    end  
end  
escape 0;
```

- MY_EVT
 - External event identifier (should be in uppercase)

Blinking an LED

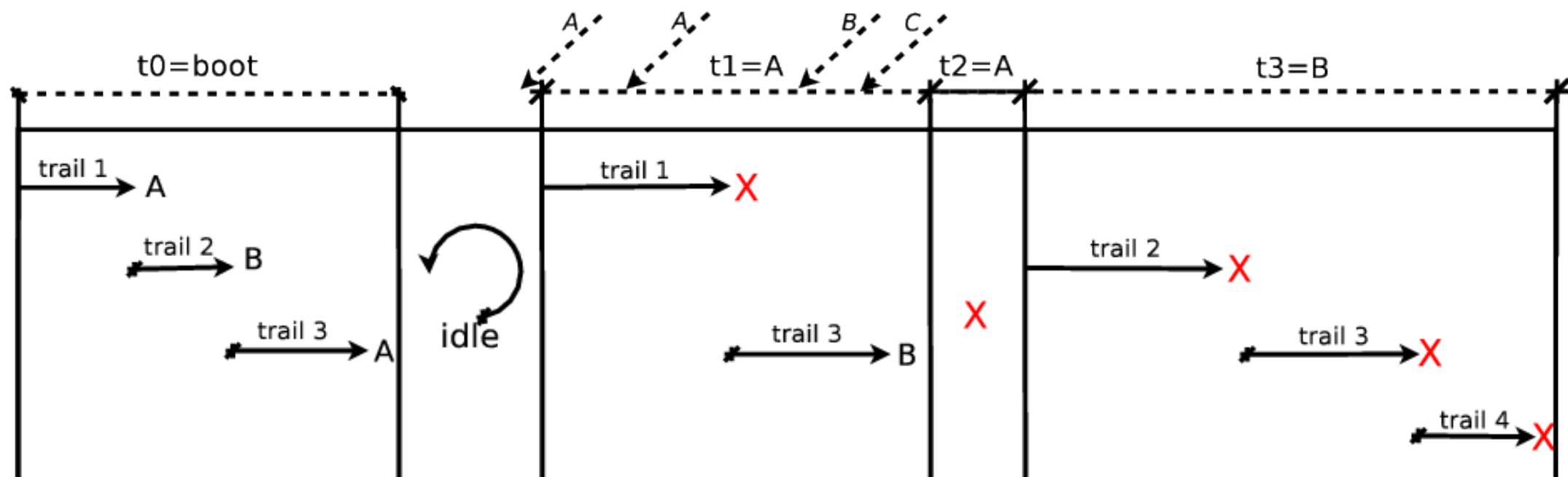
```
input  none  BUTTON;
output on/off LED;
par/or do
    await BUTTON;
with
    loop do
        await 1s;
        emit LED(on);
        await 1s;
        emit LED(off);
    end
end
```

- This program blinks an LED ~~after~~ pressing a button
until

Synchronous Execution Model of Céu

```
1:  input none A;  
2:  input none B;  
3:  input none C;  
4:  par/and do  
5:      // trail 1  
6:      <...>  
7:      await A;  
8:      <...>  
9:  with  
10:     // trail 2  
11:     <...>  
12:     await B;  
13:     <...>
```

```
14:  with  
15:      // trail 3  
16:      <...>  
17:      await A;  
18:      <...>  
19:      await B;  
20:      par/and do  
21:          // trail 3  
22:          <...>  
23:          with  
24:              // trail 4  
25:              <...>  
26:          end  
27:  end
```



Parallel Composition

- par/and
 - rejoins after all trails in parallel terminate
- par/or
 - rejoins after any trail in parallel terminates
 - aborts all other trails
- par
 - never rejoins

Parallel Compositions: par/and

```
loop do
  par/and do
    await 100ms;
    _printf("Hello ");
  with
    await 250ms;
    _printf("World!\n");
  end
end
end
```

- Repeatedly prints "Hello World!"

Parallel Compositions: par/or

```
loop do
  par/or do
    await 100ms;
    _printf("Hello ");
  with
    await 250ms;
    _printf("World!\n");           // never reached
  end
end
end
```

- Repeatedly prints "Hello " only

Parallel Compositions: par/or

```
input none TERM;
par/or do
  loop do          // par/and loop never terminates
    par/and do
      await 100ms;
      _printf("Hello ");
    with
      await 250ms;
      _printf("World!\n");
    end
  end
end
with
  await TERM;      // but par/or terminates on TERM
                  // and kills the original loop
end
```

- Repeatedly prints "Hello World!" until receives TERM

Waiting for Multiple Events

```
input none HELLO;
input none WORLD;
par do          // par/and, par/or would behave the same
  loop do
    await HELLO;
    _printf("Hello!\n");
  end
with
  loop do
    await WORLD;
    _printf("World!\n");
  end
end
```

- Waits for HELLO or WORLD
 - (cf. select)

Internal Events

```
// internal event begin in lowercase
event int e;
var    int i = 0;
par do
    loop do
        i = await e;    // waits for changes
        _printf("Hello World: %d!\n", i);
    end
with
    loop do
        await 250ms;
        emit e(i+1);    // triggers 'e'
    end
end
```

- Prints "Hello World: 1!", "Hello World: 2!", ...

Internal Events

```
event none e;          // an event carrying no values
var int v = 0;
par do
  loop do              // a simple loop that
    await e;           // when 'e' occurs
    v = v + 1;         // increments 'v'
  end
with
  // 1st trail is awaiting 'e'
  emit e;              // resumes after 1st trail halts
                      // v=0+1 => 1

  // 1st trail is awaiting 'e' again
  emit e;              // resumes after 1st trail halts
                      // v=1+1 => 1

  escape v;           // v=2
end
```