

# RSA PUBLIC-KEY ENCRYPTION AND SIGNATURE LAB

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## SEED Lab #7

### Identification

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### Task 1:

To calculate the private key  $d$ , we created the following script:

```
#include <stdio.h>
#include <openssl/bn.h>

void printBN(char *msg, BIGNUM * a) {
    char * number_str = BN_bn2hex(a);
    printf("%s %s\n", msg, number_str);
    OPENSSL_free(number_str);
}

int main () {
    BN_CTX *ctx = BN_CTX_new();

    BIGNUM *minusone = BN_new();

    BIGNUM *p = BN_new();
    BIGNUM *q = BN_new();
    BIGNUM *e = BN_new();

    BIGNUM *p1 = BN_new();
    BIGNUM *q1 = BN_new();
    BIGNUM *r = BN_new();

    BIGNUM *pkey = BN_new();

    BN_dec2bn(&minusone, "-1");

    BN_hex2bn(&p, "F7E75FDC469067FFDC4E847C51F452DF");
    BN_hex2bn(&q, "E85CED54AF57E53E092113E62F436F4F");
    BN_hex2bn(&e, "0D88C3");

    BN_add(p1, p, minusone);
    BN_add(q1, q, minusone);

    BN_mul(r, p1, q1, ctx);
```

```

    BN_mod_inverse(pkey, e, r, ctx);

    printBN("Private key:", pkey);

    return 0;
}

```

Basically, it starts by calculating the Euler's totient function using  $\Phi(n) = (p-1) * (q-1)$ . Then, we need to pick an  $e$  value. This  $e$  value can be any number  $e$  where  $1 < e < \Phi(n)$  and  $e$  are coprime to  $\Phi(n)$ , common values being the *Fermat Primes* (3, 17, and 65537). In this instance, it was already defined for us as 886979. Now we can compute the private key  $d$ , which is the modular multiplicative inverse of  $e \pmod{\Phi(n)}$ .

By executing that script, we obtained the value of the private key:  $d = 3587A24598E5F2A21DB007D89D18CC50ABA5075BA19A33890FE7C28A9B496AEB$ . We verified that this value is correct by entering the values into an online RSA calculator.

## Task 2:

The first step is to get the hex encoding of the secret message, which is  $4120746f702073656372657421$ , corresponding to our plaintext  $P$ . We did this by running the following python command:

```
python -c 'print("A top secret!".encode("hex"))'
```

Now, to encrypt the message, we generate the ciphertext using  $C = P^e \pmod{n}$ . To verify the encryption result, we decrypted the resulting ciphertext using  $P = C^d \pmod{n}$  and noticed that the plaintext remains the same. This procedure is done in the following script:

```

(...)

int main () {
    BN_CTX *ctx = BN_CTX_new();

    BIGNUM *n = BN_new();
    BIGNUM *e = BN_new();
    BIGNUM *M = BN_new();
    BIGNUM *d = BN_new();

    BIGNUM *C = BN_new();

    BIGNUM *secret = BN_new();

    BN_hex2bn(&n,
"DCBFFE3E51F62E09CE7032E2677A78946A849DC4CDDE3A4D0CB81629242FB1A5");
    BN_hex2bn(&e, "010001");
    BN_hex2bn(&M, "4120746f702073656372657421"); // Hex equivalent of "A top
secret!"
    BN_hex2bn(&d,
"74D806F9F3A62BAE331FFE3F0A68AFE35B3D2E4794148AACBC26AA381CD7D30D");
}

```

```

    BN_mod_exp(C, M, e, n, ctx);

    printBN("Ciphertext:", C);

    BN_mod_exp(secret, C, d, n, ctx);

    printBN("Decrypted message:", secret);

    return 0;
}

```

### Task 3:

In this task, the goal is to decrypt a ciphertext. For that, we repeat the procedure of the previous task, but now we just need to get the corresponding plaintext and convert it to an ASCII string. The following script does that:

```

(...)

int main () {
    BN_CTX *ctx = BN_CTX_new();

    BIGNUM *n = BN_new();
    BIGNUM *M = BN_new();
    BIGNUM *d = BN_new();

    BIGNUM *C = BN_new();

    BN_hex2bn(&n,
"DCBFFE3E51F62E09CE7032E2677A78946A849DC4CDDE3A4D0CB81629242FB1A5");
    BN_hex2bn(&C,
"8C0F971DF2F3672B28811407E2DABBE1DA0FEBBDFC7DCB67396567EA1E2493F");
    BN_hex2bn(&d,
"74D806F9F3A62BAE331FFE3F0A68AFE35B3D2E4794148AACBC26AA381CD7D30D");

    BN_mod_exp(M, C, d, n, ctx);

    printBN("Decrypted message:", M);

    return 0;
}

```

By running it, we discovered that the secret message is `50617373776F72642069732064656573`. This is the hexadecimal value, and converting it to ASCII, the message is revealed as being `Password is dees`.

### Task 4:

The first step is to get the hex encoding of both messages, the first one is

49206F776520796F752024323030302E, and the second one is 49206F776520796F752024333030302E. We did this by running the following python commands:

```
python -c 'print(" I owe you $2000.".encode("hex"))'
python -c 'print(" I owe you $3000.".encode("hex"))'
```

Now, we must generate the signatures for both messages, which is done by encrypting them with the private key, using  $S = P^d \pmod n$ :

```
(...)

int main () {
    BN_CTX *ctx = BN_CTX_new();

    BIGNUM *n = BN_new();
    BIGNUM *d = BN_new();
    BIGNUM *M1 = BN_new();
    BIGNUM *M2 = BN_new();

    BIGNUM *S1 = BN_new();
    BIGNUM *S2 = BN_new();

    BN_hex2bn(&n,
"DCBFFE3E51F62E09CE7032E2677A78946A849DC4CDDE3A4D0CB81629242FB1A5");
    BN_hex2bn(&d,
"74D806F9F3A62BAE331FFE3F0A68AFE35B3D2E4794148AACBC26AA381CD7D30D");

    BN_hex2bn(&M1, "49206F776520796F752024323030302E");
    BN_hex2bn(&M2, "49206F776520796F752024333030302E");

    BN_mod_exp(S1, M1, d, n, ctx);
    BN_mod_exp(S2, M2, d, n, ctx);

    printBN("Hex message 1:", M1);
    printBN("Hex message 2:", M2);

    printBN("Signed message 1:", S1);
    printBN("Signed message 2:", S2);

    return 0;
}
```

We obtained the following signatures:

- S1 = 55A4E7F17F04CCFE2766E1EB32ADDBA890BBE92A6FBE2D785ED6E73CCB35E4CB
- S2 = BCC20FB7568E5D48E434C387C06A6025E90D29D848AF9C3EBAC0135D99305822

As we can see, in the digital signatures above, even though the messages only differ on a single byte, their encrypted counterparts are extremely different due to the protection given by the modulus operation.

## Task 5:

To verify the signature, we must decrypt it using the public key, using  $P = S^e \pmod n$ . Then, we converted it to an ASCII string and noticed that the message is the same, meaning that the signature is indeed Alice's.

Then, we experimented corrupting the signature by changing its last byte from 2F to 3F:

- Decrypted message (hex): 4C61756E63682061206D697373696C652E
- Corrupted decrypted message (hex): 91471927C80DF1E42C154FB4638CE8BC726D3D66C83A4EB6B7BE0203B41AC294
- Decrypted message (ASCII): Launch a missile.
- Corrupted decrypted message (ASCII): □G□'Èñä,□O´c□è%rm=fÈ:N¶.¼□□´□Â□

As we can see, we got two entirely different results, where the plaintext of the corrupted signature is totally imperceptible. This happens because even though we're raising it to the same exponent, the base (the signature) is different, which leads to a completely different result.

The entire procedure is in the following script:

```
(...)  
  
int main () {  
    BN_CTX *ctx = BN_CTX_new();  
  
    BIGNUM *M = BN_new();  
    BIGNUM *bad_M = BN_new();  
    BIGNUM *S = BN_new();  
    BIGNUM *bad_S = BN_new();  
    BIGNUM *e = BN_new();  
    BIGNUM *n = BN_new();  
  
    BN_hex2bn(&S,  
"643D6F34902D9C7EC90CB0B2BCA36C47FA37165C0005CAB026C0542CBDB6802F");  
    BN_hex2bn(&bad_S,  
"643D6F34902D9C7EC90CB0B2BCA36C47FA37165C0005CAB026C0542CBDB6803F");  
    BN_hex2bn(&e, "010001");  
    BN_hex2bn(&n,  
"AE1CD4DC432798D933779FBD46C6E1247F0CF1233595113AA51B450F18116115");  
  
    BN_mod_exp(M, S, e, n, ctx);  
  
    printBN("Decrypted message:", M);  
  
    BN_mod_exp(bad_M, bad_S, e, n, ctx);  
  
    printBN("Corrupted decrypted message:", bad_M);
```

```
    return 0;
}
```

## Task 6:

### Step 1:

For this first step, we must save both certificates (the server CA and the intermediate CA) of a given web server. We chose to use the Twitter website ([www.twitter.com](http://www.twitter.com)).

Server certificate:

```
-----BEGIN CERTIFICATE-----
MIIGwzCCBaugAwIBAgIQAjf+8b3vOp6vWPt1+RUzwzANBgqhkiG9w0BAQsFADBP
MQswCQYDVQQGEwJVUzEVMBMGA1UEChMMRGlnaUNlcnQgSW5jMSkwJwYDVQQDEyBE
aWdpQ2VydCBUTFMgUlNBIFNIQTl1NiAyMDIwIENBMTAeFw0yMjAxMjQwMDAwMDBa
(...)
```

Intermediate certificate:

```
-----BEGIN CERTIFICATE-----
MIIEvjCCA6agAwIBAgIQBtjZBNVYQ0b2ii+nVCJ+xDANBgqhkiG9w0BAQsFADBh
MQswCQYDVQQGEwJVUzEVMBMGA1UEChMMRGlnaUNlcnQgSW5jMRkwFwYDVQQLExB3
d3cuZGlnaWNlcnQuY29tMSAwHgYDVQQDEXdEaWdpQ2VydCBHbG9iYWwgUm9vdCBD
(...)
```

### Step 2:

To extract the public key (e, n) from the issuer's certificate, we extract the value of **n** using the command:

```
openssl x509 -in c1.pem -noout -modulus
```

And we print out all the fields to find the exponent **e** using the command:

```
openssl x509 -in c1.pem -text -noout
```

- Modulus (n): 00:c1:4b:b3:65:47:70:bc:dd:4f:58:db:ec:9c:ed: c3:66:e5:1f:31:13:54:ad:4a:66:46:1f:2c:0a:ec: 64:07:e5:2e:dc:dc:b9:0a:20:ed:df:e3:c4:d0:9e:  
(...)
- Exponent (e): 65537 (0x10001)

**Step 3:**

To extract the signature, we must print all the fields using the command:

```
openssl x509 -in c0.pem -text -noout
```

- Signature Algorithm: sha384WithRSAEncryption 90:e8:32:79:96:50:59:39:06:18:bc:07:40:7e:2f:54:d7:9d:2d:2a:94:c0:08:7d:97:7f:d9:c0:86:d6:71:c0:2c:ef:b7:65:13:d4:2c:98:68:bf:9f:5e:32:be:2a:1b:58:b1:bc:cf:d8:2d: (...)

We then need to remove the spaces and colons from the data, to obtain a hex string that can be used in our program. We do that by running:

```
cat signature | tr -d '[:space:]:'
```

- Signature without the spaces and colons:  
90e83279965059390618bc07407e2f54d79d2d2a94c0087d977fd9c086d671c02cefb76513d42c9868bf9f5e32be2a1b58b1bccfd82d68be78a794f664b561dc1d9445a1bc470cee41a6cbd3e6 (...)

**Step 4:**

Now, we extract the body of the certificate by running the command:

```
openssl asn1parse -i -in c0.pem -strparse 4 -out c0_body.bin -noout
```

After this, we generate its hash by running:

```
sha256sum c0_body.bin
```

**Step 5:**

Now that we have all the information we need, we can verify whether the signature is valid or not by comparing the hash of the certificate body with the plaintext obtained through  $P = S^e \pmod n$ . The following program does that verification:

```
(...)  
  
int main () {  
    BN_CTX *ctx = BN_CTX_new();  
  
    BIGNUM *e = BN_new();  
    BIGNUM *n = BN_new();
```

}

[illegible]