

On-line Colouring of interval graphs

Introduction to research

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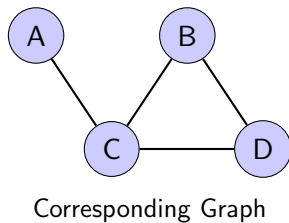
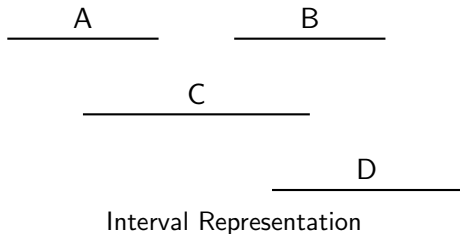
Supervised by: Dr. Grzegorz Gutowski

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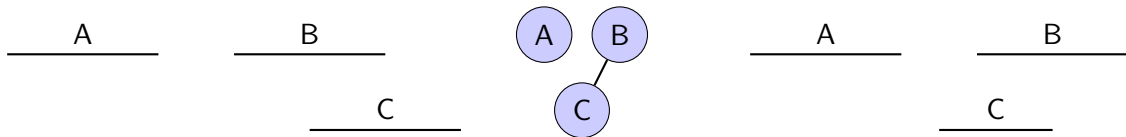
Overview

1. Interval graph representation
2. Problem definition
3. Status of research
4. Using a computer
5. Our work

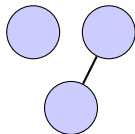
Interval Representation and Corresponding Graph



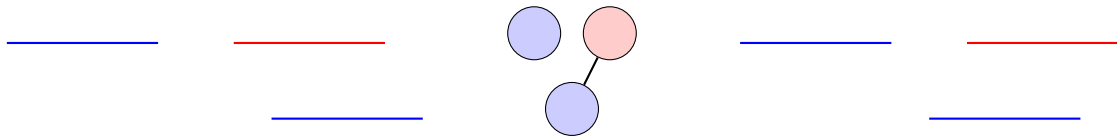
Graph with two Interval Representations

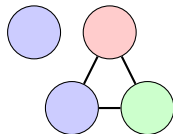
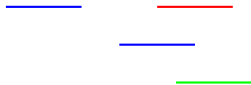
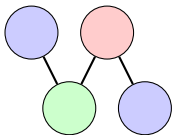


Are these representations equivalent?

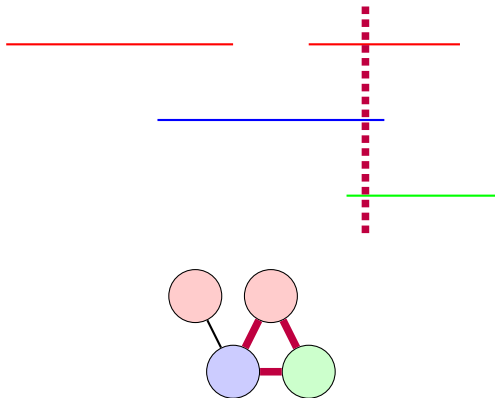


What if we introduce colours?

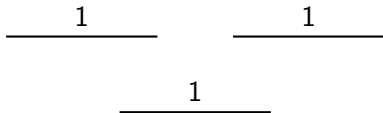




Max clique size



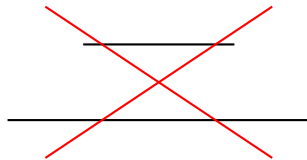
Classes of interval representations



Unit interval representation



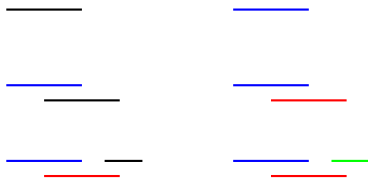
Proper interval representation



Problem definition

A **game** in rounds:

- The *builder* places an interval s.t. the max clique size is not exceeded
- The colouring *algorithm* chooses a colour for that interval



$R(\omega) :=$ what is the maximum number of colours the *builder* can force the *algorithm* to use without exceeding the max clique size ω ?

Where we are at?

$R(\omega)$:= what is the maximum number of colours the *builder* can force the *algorithm* to use without exceeding the max clique size ω ?

- $R(\omega) \leq 2\omega - 1$ Chrobak and Ślusarek, 1981
- $R(\omega) \geq \lfloor 3\omega/2 \rfloor$ Epstein and Levy, 2005
- $R(1) = 1, R(2) = 3$
- $R(3) \in \{4, 5\}$
 - $R(3) = 5$ Biró and Curbelo, 2022
- $R(4) \in \{6, 7\}$
 - $R(4) = 7$ Curbelo and Malko, 2024 [1]
- $R(\omega)$ unsolved for $\omega \geq 5$

Minimax

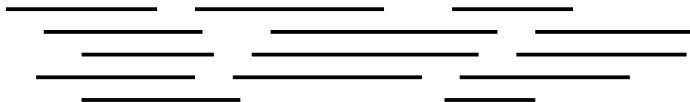
Maximin value

$$v_p = \max_{a_p} \min_{a_o} v_p(a_p, a_o)$$

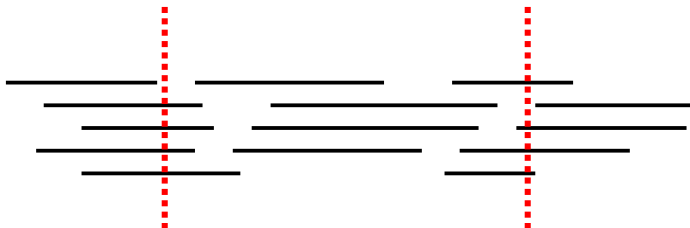
Where:

- p is the player of interest.
- o denotes the opponent.
- a_p is the action taken by player p .
- a_o is the action taken by opponent o .
- v_p is the value function of player p .

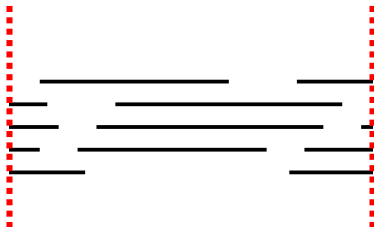
Confining the Space



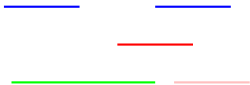
Confining the Space



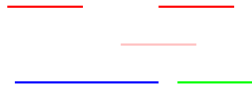
Confining the Space



State equivalence



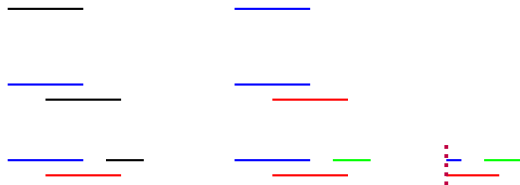
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Problem definition

A **game** in rounds:

- The *builder* places an interval s.t. the max clique size is not exceeded
- The colouring *algorithm* chooses a colour for that interval
- The *builder* **may** restrict the domain



$R(\omega) :=$ what is the maximum number of colours the *builder* can force the *algorithm* to use without exceeding the max clique size ω ?

State

$$E = \{1, 2, \dots, \omega, \neg\}$$

$$S = (e_1, e_2, \dots, e_k), e_i \in E$$

Our work

$$\lfloor 3\omega/2 \rfloor \leq R(\omega) \leq 2\omega - 1$$

- $R(1) = 1$
- $R(2) = 3$
- **$R(3) = 5$**
- $R(4) = 7$
- $R(5) = ?$
- $R(6) = ?$
- ...

Bibliography



Israel R. Curbelo, Hannah R. Malko, *On the on-line coloring of unit interval graphs with proper interval representation*, 2024