

CS 503 Lab 2 Answers

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Files modified/added

lab2q3.c ; lab2q4.c ; lab2q5.c; iobound.c ; cpubound.c ;resched.c; ready.c; prototypes.h; initialize.c; process.h

Testing Methodology

The following script in the main.c is an example for carrying out the tests. No additional parameters are needed for the lab test functions. Please take a look at existing main.c for example execution of tests for each lab section.

```
/*  main.c  -  main */

#include <xinu.h>
#include <stdio.h>
int lab2flag=4; // set this!!!
process main(void)
{
    resume(create(lab2q4t1, 1024, 51, "lab2q4t1", 0, NULL));
    sleepms(6000);
    resume(create(lab2q4t2,1024,52, "lab2q4t2",0, NULL));
    sleepms(6000);
    resume(create(lab2q4t3,1024,52, "lab2q4t3",0, NULL));
    sleepms(3000);
    resume(create(shell, 8192, 50, "shell", 1, CONSOLE));
    /* Wait for shell to exit and recreate it */
    return OK;
}
```

The global variable lab2flag needs to be set prior to testing so as to enable the scheduling policy required to test the corresponding section of the lab, i.e set to 3/4/5. The sleeps in between resumes above are enabled for easy viewing(printing) to the console. To simulate these results I have 4 worker processes of priority of 54 and 1 monitor process of priority 57. Worker processes classify as either workerproceesTypeA or workerprocessTypeB. Type A processes

execute an infinite loop, and Type B is similar to Type A except that it will sleep on having counted upto a certain number(say 200), it will invoke a sleep instruction for 20ms and relinquish the processor.

Next up, we explain the monitor process. The monitor process is responsible for spawning the worker processes and then calling sleepms for a designated time. Since, monitor process has the highest priority it will take up the CPU exclusively. On relinquishing the processor with sleep for the designated time, the 4 worker processes get a chance to execute and we get a chance to test. But when the monitor process wakes up, it executes the process status shell xshps code to print the status of the processes and then kills the other 4 worker processes. Hence, this is the general methodology used to simulate the below results and by this we can visibly reason/ make inferences about the process execution.

Question 3: Monitoring CPU usage of processes

1 Test for fairness

In the image above, the upper half that describes lab2q3t2, I have set four workerprocessTypeA processes to have equal priority and they all share the CPU time, so as printed by the column CPU ms, they share the CPU time equitably. With worker process 1 being the last to execute right before the monitor process wakes up and takes over the cpu(after workerProc 1 time quantum completes of course).

2 Test for relinquishing CPU

In the middle part of the image above, I have set 4 processes to have equal priority. However, worker proce 1-3 are workerprocessTypeA and worker proc 4 is of type workerprocessTypeB. As mentioned earlier, workerprocessTypeB is set to have a sleepms() invocation after a certain duration(after it counts up to 2000000). Hence, we can assume this process sleeps a lot. Until then, however it whirs in its infinite loop. So, as we can tell, process 4 has received significantly lesser time as compared to the other 3 processes.

3 Test for highest priority process hogging CPU

In this test I have set the 3 workerprocessTypeA processes to have same priority 54 however worker proc 4 has a higher priority 55 (yet its priority is lower than that of monitor process whose priority is 57). As we can see for the duration that the monitor process sleeps, the worker proc 4 hogs the CPU without allowing the other 3 worker processes to execute and takes up all of the 375 ms allotted for worker processes to execute. In other words, monitor process was let to sleep only for 375 ms, and then after it wakes up it takes over the cpu and prints the above information for us to visually validate.

```

-----
In lab2q3t1 test | Note: sleeping for 5 seconds .. zzzzz ...
-----
Clock time in seconds since boot is 5
Clock time in milliseconds since boot is 5053
-----
In lab2q3t2 test | Note: 4 (workerProcessTypeA) equal CPU usage seen by monitorProcess
(Monitor process will have higher priority than workerProcessTypeA)
-----
Pid Name      State Prio Ppid Stack Base Stack Ptr  Stack Size CPU time  ReadyQ Posn
-----
0 prnull      ready  0    0 0x0EFDEFFC 0x0EFDEF1C 8192      0      6
1 rdsproc     wait  200   0 0x0EFDCCFFC 0x0EFDCAAC 16384     2
3 Main process ready  20    0 0x0EFD8FFC 0x0EFD8F44 65536     1      5
5 cputimetest1 ready  52    3 0x0FDEFFFC 0x0FDEFF40 1024      19      4
6 Worker Proc1 ready  54    5 0x0FDEFBFC 0x0FDEFB3C 1024      66      3
7 Worker Proc2 ready  54    5 0x0FDEF7FC 0x0FDEF76C 1024      61      0
8 Worker Proc3 ready  54    5 0x0FDEF3FC 0x0FDEF36C 1024      61      1
9 Worker Proc4 ready  54    5 0x0EFC8FFC 0x0EFC8F6C 1024      61      2
10 Monitor Proc curr  57    5 0x0EFC8BFC 0x0EFC8B5C 1024      61

-----
In lab2q3t3 test | Note: 3 (workerProcessTypeA) and 1 sleeper workerProcessTypeB
(Monitor process will have higher priority than workerProcessTypeA/B)
-----
Pid Name      State Prio Ppid Stack Base Stack Ptr  Stack Size CPU time  ReadyQ Posn
-----
0 prnull      ready  0    0 0x0EFDEFFC 0x0EFDEF1C 8192      0      6
1 rdsproc     wait  200   0 0x0EFDCCFFC 0x0EFDCAAC 16384     2
3 Main process ready  20    0 0x0EFD8FFC 0x0EFD8F44 65536     1      5
11 cputimetest2 ready  52    3 0x0FDEFFFC 0x0FDEFF40 1024      19      4
12 Worker Proc1 ready  54   11 0x0FDEFBFC 0x0FDEFB3C 1024      91      1
13 Worker Proc2 ready  54   11 0x0FDEF7FC 0x0FDEF76C 1024     121      2
14 Worker Proc3 ready  54   11 0x0FDEF3FC 0x0FDEF33C 1024     106      3
15 Worker Proc4 ready  54   11 0x0EFC8FFC 0x0EFC8F80 1024      61      0
16 Monitor Proc curr  57   11 0x0EFC8BFC 0x0EFC8B5C 1024      61

-----
In lab2q3t4 test | Note: 3 (workerProcessTypeA) and 1 higher priority workerProcessTypeA
(Monitor process will have higher priority than workerProcessTypeA)
-----
Pid Name      State Prio Ppid Stack Base Stack Ptr  Stack Size CPU time  ReadyQ Posn
-----
0 prnull      ready  0    0 0x0EFDEFFC 0x0EFDEF1C 8192      0      6
1 rdsproc     wait  200   0 0x0EFDCCFFC 0x0EFDCAAC 16384     2
3 Main process ready  20    0 0x0EFD8FFC 0x0EFD8F44 65536     1      5
17 cputimetest3 ready  52    3 0x0FDEFFFC 0x0FDEFF40 1024      19      4
18 Worker Proc1 ready  54   17 0x0FDEFBFC 0x0FDEFBCC 1024      1      1
19 Worker Proc2 ready  54   17 0x0FDEF7FC 0x0FDEF7CC 1024      1      2
20 Worker Proc3 ready  54   17 0x0FDEF3FC 0x0FDEF3CC 1024      1      3
21 Worker Proc4 ready  55   17 0x0EFC8FFC 0x0EFC8F3C 1024     376      0
22 Monitor Proc curr  57   17 0x0EFC8BFC 0x0EFC8B5C 1024      62

```

Figure 1: Experiment results using lab2q3.c tests

Question 4.2: Dynamic priority scheduling to achieve "fairness"

I used the same tests as for part 3 and these were the results. Right off the bat, we notice that the CPU time used by the CPU hogging processes and IO hogging processes is nearly indistinguishable in all three cases and all the values are 30ms within the range of each other. This shows that our dynamic priority scheduling strategy shares the CPU equitably among IO hogging and CPU hogging applications.

```

-- LAB 2.4.2 --

-----
In lab2q4t1 test | Note: 4 (workerProcessTypeA) equal CPU usage seen by monitorProcess
(Monitor process will have higher priority than workerProcessTypeA)
-----
Pid Name      State Prio Ppid Stack Base Stack Ptr  Stack Size CPU time  ReadyQ Posn
-----
0 prnull      ready  0    0 0x0EFDEFFC 0x0EFDEF50 8192 2147483647 4
1 rdsproc     wait   200  0 0x0EFDCEFC 0x0EFDCAAC 16384 2
3 Main process sleep  20    0 0x0EFD8FFC 0x0EFD8F74 65536 2
5 Worker Proc1 ready  54    4 0x0FDEFBFC 0x0FDEFB3C 1024 96
6 Worker Proc2 ready  54    4 0x0FDEF7FC 0x0FDEF76C 1024 91
7 Worker Proc3 ready  54    4 0x0FDEF3FC 0x0FDEF36C 1024 91
8 Worker Proc4 ready  54    4 0x0EFC8FFC 0x0EFC8F6C 1024 91
9 Monitor Proc curr   57    4 0x0EFC8BFC 0x0EFC8B5C 1024 61

-----
In lab2q4t2 test | Note: 3 (workerProcessTypeA) and 1 sleeper workerProcessTypeB
(Monitor process will have higher priority than workerProcessTypeA/B)
-----
Pid Name      State Prio Ppid Stack Base Stack Ptr  Stack Size CPU time  ReadyQ Posn
-----
0 prnull      ready  0    0 0x0EFDEFFC 0x0EFDEF1C 8192 2147483647 4
1 rdsproc     wait   200  0 0x0EFDCEFC 0x0EFDCAAC 16384 2
3 Main process sleep  20    0 0x0EFD8FFC 0x0EFD8F74 65536 2
11 Worker Proc1 ready  54   10 0x0FDEFBFC 0x0FDEFB6C 1024 120
12 Worker Proc2 ready  54   10 0x0FDEF7FC 0x0FDEF73C 1024 130
13 Worker Proc3 ready  54   10 0x0FDEF3FC 0x0FDEF33C 1024 118
14 Worker Proc4 ready  54   10 0x0EFC8FFC 0x0EFC8F80 1024 121
15 Monitor Proc curr   57   10 0x0EFC8BFC 0x0EFC8B5C 1024 61

-----
In lab2q4t3 test | Note: 3 (workerProcessTypeA) and 1 higher priority workerProcessTypeA
(Monitor process will have higher priority than workerProcessTypeA)
-----
Pid Name      State Prio Ppid Stack Base Stack Ptr  Stack Size CPU time  ReadyQ Posn
-----
0 prnull      ready  0    0 0x0EFDEFFC 0x0EFDEF1C 8192 2147483647 4
1 rdsproc     wait   200  0 0x0EFDCEFC 0x0EFDCAAC 16384 2
3 Main process sleep  20    0 0x0EFD8FFC 0x0EFD8F74 65536 2
17 Worker Proc1 ready  54   16 0x0FDEFBFC 0x0FDEFB3C 1024 136
18 Worker Proc2 ready  54   16 0x0FDEF7FC 0x0FDEF76C 1024 121
19 Worker Proc3 ready  54   16 0x0FDEF3FC 0x0FDEF36C 1024 121
20 Worker Proc4 ready  55   16 0x0EFC8FFC 0x0EFC8F6C 1024 121
21 Monitor Proc curr   57   16 0x0EFC8BFC 0x0EFC8B5C 1024 61

-- LAB 2.4.3 --

```

Figure 2: Dynamic priority scheduling

Question 4.3: Performance evaluation

1 All processes are CPU-bound

1.1 Nulluser runs only if no other process is ready

My design solution to this problem is to set the nulluser's `prcpumsec` to be `MAXINT32` when it's instantiated, as can be viewed in the image above the value for `prnull`'s CPU time usage is `MAXINT32 = 2147483647`. In addition, its `prcpumsec` field should not be updated when it is the oldprocess to be context switched out as we can't increment beyond the max value of the `int32`. This ensures that within the TS scheduling strategy the nulluser is always at the end of the readylist, and hence only executes if no other process is ready. The only caveat is that we can't deem this to be true for long running processes that

may be running for more than 24 days. Since, 24 days is roughly MAXINT32 milliseconds. In this lab since we are not dealing with long running instances, as specified in the document specifications, we can ignore this warning for now. As can be seen from the figure below, if we call ps as we execute the code the amount of time spent by each process is within 30ms of each other. So they work in lock-step where no process is left behind.

2 All processes are iobound

As can be seen from the figure below, if we call ps as we execute the code, all the processes work in lock-step. They are 30ms within the bound of each other. So this is the expected result. This is assuming the values for loop1 and loop2 are the same for all 6 processes. If they were different for different processes then time allotted to each process could be unpredictable since it could be possible that a process that has a high prcpumsec time could execute while processes with the lower prcpumsec time are sleeping, and since the higher process is the only ready process (apart from null) it will execute and further the gap between itself and the lower processes. If this is the case then it could be very difficult for these lower prcpumsec time processes to catch up to the higher one. This can be the case only for different values of loop1 loop2 for the 6 iobound processes. Intuitively it makes sense.

3 Half-half

As shown in image we go back to the last answer where we experimented with different values of loop1 and loop2 for the all processes are iobound case and notice that the same happens for cpubound processes. Every time all iobound processes are sleeping, the cpubound process takes over and gets more clock cycles and then it always stays in the lead and eventually the cpubound processes finish faster than the iobound processes. This is normal. Actually, it is the preferred way an OS must execute. What is the alternative to this? A cPU bound process not executing while all such IO bound processes are sleeping. Why waste the clock cycles. However, as soon as an iobound process wakes up it is given priority since it was given lower cpu clock cycles over the cpubound processes.

Question 5: Dynamic workloads

My solution to this problem involves a very intuitive yet effective method to manage dynamic workloads. To revisit the problem we have long standing process who gets very heavily demoted by a newly entered process. So the process would be allowed to execute until its cputime allotted rises up to the long standing process. But this is clearly leading to starvation. So the solution is to demote priority of process for cputime used and promote priority for the amount of time it has spent on the readylist. Hence, every time resched is

invoked all the processes on the ready list are given 6ms of promotion. Here we are using the same ready list as provided originally but the key of a process when inserted is the negation of its cputime used. Hence, the lower the cputime used by processes the better the chance of it to be executed next. By adding 6ms on each resched call to ready list processes that haven't won the ongoing time quantum, they slowly move up the process queue to get a chance to execute next, a chance they could've never got without this strategy. Hence, a priority of process is defined by $= \text{wait-time} - \text{cpu-time used}$, where wait time is $6\text{ms} * (\text{number of times resched was invoked while it was on ready list})$. The results are shown in the image below.

As we can see that process with pid 4 has CPU time 690 while processes with pid 8 and 9 are supposed to demote the other processes by a lot and hog the cpu until they get a similar amount of time as them 690ms. But after a while process 4's ReadyQ key is decremented because it has been rewarded for waiting. And even though pid 9 has cpu time of 481 and pid 4 process has 690 cpu time the position of process with pid 4 on readylist is higher than process with pid 9. Hence, we have shown that dynamic workloads can be scheduled fairly, to prevent starvation.

```

69 ps curr 20 28 0x0EFD8FFC 0x0EFD8DF8 8192 2
xsh $ ps
Pid Name State Prio Ppid Stack Base Stack Ptr Stack Size CPU Time
-----
0 prnull ready 0 0 0x0EFDEFFC 0x0EFDEF1C 8192 2147483647
1 rdsproc wait 200 0 0x0EFD8FFC 0x0EFDCAAC 16384 2
22 cpubound 1 ready 1 3 0x0FDEFFFC 0x0FDEFCEC 1024 7062
23 cpubound 2 ready 1 3 0x0FDEFBFC 0x0FDEF980 1024 7084
24 cpubound 3 ready 1 3 0x0FDEF7FC 0x0FDEF584 1024 7085
25 cpubound 4 ready 1 3 0x0FDEF3FC 0x0FDEF0EC 1024 7069
26 cpubound 5 ready 1 3 0x0EFC8FFC 0x0EFC8CF0 1024 7071
27 cpubound 6 ready 1 3 0x0EFC8BFC 0x0EFC88F0 1024 7072
28 shell recv 50 3 0x0EFC87FC 0x0EFC848C 8192 2
70 ps curr 20 28 0x0EFD8FFC 0x0EFD8DF8 8192 2
xsh $ ps
Pid Name State Prio Ppid Stack Base Stack Ptr Stack Size CPU Time
-----
0 prnull ready 0 0 0x0EFDEFFC 0x0EFDEF1C 8192 2147483647
1 rdsproc wait 200 0 0x0EFD8FFC 0x0EFDCAAC 16384 2
22 cpubound 1 ready 1 3 0x0FDEFFFC 0x0FDEFCEC 1024 7176
23 cpubound 2 ready 1 3 0x0FDEFBFC 0x0FDEF8F0 1024 7176
24 cpubound 3 ready 1 3 0x0FDEF7FC 0x0FDEF4F0 1024 7177
25 cpubound 4 ready 1 3 0x0FDEF3FC 0x0FDEF160 1024 7180
26 cpubound 5 ready 1 3 0x0EFC8FFC 0x0EFC8CF0 1024 7180
27 cpubound 6 ready 1 3 0x0EFC8BFC 0x0EFC88EC 1024 7180
28 shell recv 50 3 0x0EFC87FC 0x0EFC848C 8192 2
71 ps curr 20 28 0x0EFD8FFC 0x0EFD8DF8 8192 3
xsh $ ps
Pid Name State Prio Ppid Stack Base Stack Ptr Stack Size CPU Time
-----
0 prnull ready 0 0 0x0EFDEFFC 0x0EFDEF1C 8192 2147483647
1 rdsproc wait 200 0 0x0EFD8FFC 0x0EFDCAAC 16384 2
22 cpubound 1 ready 1 3 0x0FDEFFFC 0x0FDEFCEC 1024 7252
23 cpubound 2 ready 1 3 0x0FDEFBFC 0x0FDEF960 1024 7268
24 cpubound 3 ready 1 3 0x0FDEF7FC 0x0FDEF4EC 1024 7256
25 cpubound 4 ready 1 3 0x0FDEF3FC 0x0FDEF184 1024 7274
26 cpubound 5 ready 1 3 0x0EFC8FFC 0x0EFC8CF0 1024 7261
27 cpubound 6 ready 1 3 0x0EFC8BFC 0x0EFC8980 1024 7274
28 shell recv 50 3 0x0EFC87FC 0x0EFC848C 8192 2
72 ps curr 20 28 0x0EFD8FFC 0x0EFD8DF8 8192 3
xsh $ ps
Pid Name State Prio Ppid Stack Base Stack Ptr Stack Size CPU Time
-----
0 prnull ready 0 0 0x0EFDEFFC 0x0EFDEF1C 8192 2147483647
1 rdsproc wait 200 0 0x0EFD8FFC 0x0EFDCAAC 16384 2
22 cpubound 1 ready 1 3 0x0FDEFFFC 0x0FDEFD84 1024 7389
23 cpubound 2 ready 1 3 0x0FDEFBFC 0x0FDEF980 1024 7389
24 cpubound 3 ready 1 3 0x0FDEF7FC 0x0FDEF580 1024 7390
25 cpubound 4 ready 1 3 0x0FDEF3FC 0x0FDEF0EC 1024 7387
26 cpubound 5 ready 1 3 0x0EFC8FFC 0x0EFC8CF0 1024 7391
27 cpubound 6 ready 1 3 0x0EFC8BFC 0x0EFC88EC 1024 7391
28 shell recv 50 3 0x0EFC87FC 0x0EFC848C 8192 2
73 ps curr 20 28 0x0EFD8FFC 0x0EFD8DF8 8192 3
xsh $ ps

```

Figure 3: All processes are CPU bound

```

uter loop 36 | process priority is 1 | preempt is 30
iority is 1 | preempt is 30
ps
Pid Name          State Prio Ppid Stack Base Stack Ptr  Stack Size  CPU Time
-----
 0 prnull         ready  0   0 0x0EFDEFFC 0x0EFDEEB8    8192 2147483647
 1 rdsproc        wait  200  0 0x0EFD8FFC 0x0EFDCAAC   16384      2
22 shell          recv  50   3 0x0EFC8FFC 0x0EFC8C8C    8192      4
23 iobound 1      sleep  1   3 0x0FDEFFFC 0x0FDEFF58    1024     307
24 iobound 2      sleep  1   3 0x0FDEFBFC 0x0FDEFB58    1024     302
25 iobound 3      sleep  1   3 0x0FDEF7FC 0x0FDEF758    1024     304
26 iobound 4      sleep  1   3 0x0FDEF3FC 0x0FDEF358    1024     288
27 iobound 5      sleep  1   3 0x0EFC6FFC 0x0EFC6F58    1024     306
28 iobound 6      sleep  1   3 0x0EFC6BFC 0x0EFC6B58    1024     302
50 ps            curr  20  22 0x0EFD8FFC 0x0EFD8DF8    8192      2
xsh $
In iobound 4 |The PID is 26 | outer loop 37 | process priority is 1 | preempt is 30

In iobound 6 |The PID is 28 | outer loop 37 | process priority is 1 | preempt is 30

In iobound 2 |The PID is 24 | outer loop 37 | process priority is 1 | preempt is 30

In iobound 3 |The PID is 25 | out
In iobound 5 |The PID is 27 | outer loop 37 | process priority is 1 | preempt is 30

In iobound 1 |The PID ier loop 37 | process priority is 1 | preempt is 30
s 23 | outer loop 37 | process priority is 1 | preempt is 30
ps
Pid Name          State Prio Ppid Stack Base Stack Ptr  Stack Size  CPU Time
-----
 0 prnull         ready  0   0 0x0EFDEFFC 0x0EFDEEB8    8192 2147483647
 1 rdsproc        wait  200  0 0x0EFD8FFC 0x0EFDCAAC   16384      2
22 shell          recv  50   3 0x0EFC8FFC 0x0EFC8C8C    8192      4
23 iobound 1      sleep  1   3 0x0FDEFFFC 0x0FDEFF58    1024     316
24 iobound 2      sleep  1   3 0x0FDEFBFC 0x0FDEFB58    1024     309
25 iobound 3      sleep  1   3 0x0FDEF7FC 0x0FDEF758    1024     311
26 iobound 4      ready  1   3 0x0FDEF3FC 0x0FDEF1D0    1024     300
27 iobound 5      sleep  1   3 0x0EFC6FFC 0x0EFC6F58    1024     313
28 iobound 6      sleep  1   3 0x0EFC6BFC 0x0EFC6B58    1024     309
51 ps            curr  20  22 0x0EFD8FFC 0x0EFD8DF8    8192      2
xsh $ he PID is 26 | outer loop 38 | process priority is 1 | preempt is 30

```

Figure 4: All processes are IO bound


```

ps
Pid Name          State Prio Ppid Stack Base Stack Ptr  Stack Size  CPU Time
-----
0 prnull          ready  0    0 0x0EFDEFFC 0x0EFDEF1C    8192 2147483647
1 rdsproc         wait  200  0 0x0EFD8FFC 0x0EFDCAAC    16384      2
3 Main process    sleep  20  0 0x0EFD8FFC 0x0EFD8F74    65536      3
22 iobound 4      sleep  1   3 0x0FDEFFFC 0x0FDEFF58     1024     151
23 iobound 5      sleep  1   3 0x0FDEFBFC 0x0FDEFB58     1024      15
24 iobound 6      sleep  1   3 0x0FDEF7FC 0x0FDEF758     1024       8
25 cpubound 4     ready  1   3 0x0FDEF3FC 0x0FDEF0F0     1024     478
26 cpubound 5     ready  1   3 0x0EFC8FFC 0x0EFC8D54     1024     479
27 cpubound 6     ready  1   3 0x0EFC8BFC 0x0EFC8F0     1024     478
28 shell          recv  50   3 0x0EFC87FC 0x0EFC848C     8192       1
34 ps             curr  20  28 0x0EFC67FC 0x0EFC65B8     8192       2
xsh $ e PID is 24 | outer loop 1 | process priority is 1 | preempt is 30

In iobound 5 |The PID is 23 | outer loop 2 | process priority is 1 | preempt is 30

In iobound 4 |The PID is 22 | outer loop 3 | process priority is 1 | preempt is 30

In iobound 4 |The PID is 22 | outer loop 4 | process priority is 1 | preempt is 30

In iobound 5 |The PID is 23 | outer loop 3 | process priority is 1 | preempt is 30

In iobound 6 |The PID is 24 | outer loop 2 | process priority is 1 | preempt is 30
p
In iobound 4 |The PID is 22 | outer loop 5 | process priority is 1 | preempt is 30
s
Pid Name          State Prio Ppid Stack Base Stack Ptr  Stack Size  CPU Time
-----
0 prnull          ready  0    0 0x0EFDEFFC 0x0EFDEF1C    8192 2147483647
1 rdsproc         wait  200  0 0x0EFD8FFC 0x0EFDCAAC    16384      2
3 Main process    sleep  20  0 0x0EFD8FFC 0x0EFD8F74    65536      3
22 iobound 4      sleep  1   3 0x0FDEFFFC 0x0FDEFF58     1024     173
23 iobound 5      sleep  1   3 0x0FDEFBFC 0x0FDEFB58     1024      29
24 iobound 6      sleep  1   3 0x0FDEF7FC 0x0FDEF758     1024      63
25 cpubound 4     ready  1   3 0x0FDEF3FC 0x0FDEF0F0     1024     896
26 cpubound 5     ready  1   3 0x0EFC8FFC 0x0EFC8CF0     1024     897
27 cpubound 6     ready  1   3 0x0EFC8BFC 0x0EFC88EC     1024     900
28 shell          recv  50   3 0x0EFC87FC 0x0EFC848C     8192       1
35 ps             curr  20  28 0x0EFC67FC 0x0EFC65B8     8192       2

```

Figure 5: Half -Half

| -- LAB 2.5 -- | | | | | | | | | | | |
|---------------|--------------|-------|------|------|------------|------------|------------|------------|--------|------|-------------|
| Pid | Name | State | Prio | Ppid | Stack Base | Stack Ptr | Stack Size | CPU time | ReadyQ | Posn | ReadyQ Key |
| 0 | prnull | ready | 0 | 0 | 0x0EFDEFFC | 0x0EFDEF40 | 8192 | 2147483647 | | 6 | -2147483647 |
| 1 | rdspc | wait | 200 | 0 | 0x0EFD0FFC | 0x0EFDCA9C | 16384 | 2 | | | |
| 3 | Main process | curr | 20 | 0 | 0x0EFD8FFC | 0x0EFD8AD0 | 65536 | 32 | | | |
| 4 | cpubound 1 | ready | 1 | 3 | 0x0FDEF7FC | 0x0FDEFD74 | 1024 | 690 | | 3 | -624 |
| 5 | cpubound 2 | ready | 1 | 3 | 0x0FDEFBFC | 0x0FDEF970 | 1024 | 691 | | 2 | -619 |
| 6 | cpubound 3 | ready | 1 | 3 | 0x0FDEF7FC | 0x0FDEF544 | 1024 | 702 | | 5 | -648 |
| 7 | cpubound 4 | ready | 1 | 3 | 0x0FDEF3FC | 0x0FDEF174 | 1024 | 691 | | 4 | -631 |
| 8 | cpubound 5 | ready | 1 | 3 | 0x0EFC8FFC | 0x0EFC8D44 | 1024 | 131 | | 1 | -119 |
| 9 | cpubound 6 | ready | 1 | 3 | 0x0EFC8BFC | 0x0EFC8974 | 1024 | 91 | | 0 | -31 |
| 0 | prnull | ready | 0 | 0 | 0x0EFDEFFC | 0x0EFDEF40 | 8192 | 2147483647 | | 6 | -2147483647 |
| 1 | rdspc | wait | 200 | 0 | 0x0EFD0FFC | 0x0EFDCA9C | 16384 | 2 | | | |
| 3 | Main process | curr | 20 | 0 | 0x0EFD8FFC | 0x0EFD8AF0 | 65536 | 122 | | | |
| 4 | cpubound 1 | ready | 1 | 3 | 0x0FDEF7FC | 0x0FDEFD74 | 1024 | 690 | | 3 | -582 |
| 5 | cpubound 2 | ready | 1 | 3 | 0x0FDEFBFC | 0x0FDEF970 | 1024 | 691 | | 2 | -577 |
| 6 | cpubound 3 | ready | 1 | 3 | 0x0FDEF7FC | 0x0FDEF544 | 1024 | 702 | | 5 | -606 |
| 7 | cpubound 4 | ready | 1 | 3 | 0x0FDEF3FC | 0x0FDEF174 | 1024 | 691 | | 4 | -589 |
| 8 | cpubound 5 | ready | 1 | 3 | 0x0EFC8FFC | 0x0EFC8D74 | 1024 | 190 | | 1 | -160 |
| 9 | cpubound 6 | ready | 1 | 3 | 0x0EFC8BFC | 0x0EFC8974 | 1024 | 181 | | 0 | -115 |
| 0 | prnull | ready | 0 | 0 | 0x0EFDEFFC | 0x0EFDEF40 | 8192 | 2147483647 | | 6 | -2147483647 |
| 1 | rdspc | wait | 200 | 0 | 0x0EFD0FFC | 0x0EFDCA9C | 16384 | 2 | | | |
| 3 | Main process | curr | 20 | 0 | 0x0EFD8FFC | 0x0EFD8AD0 | 65536 | 212 | | | |
| 4 | cpubound 1 | ready | 1 | 3 | 0x0FDEF7FC | 0x0FDEFD74 | 1024 | 690 | | 3 | -534 |
| 5 | cpubound 2 | ready | 1 | 3 | 0x0FDEFBFC | 0x0FDEF970 | 1024 | 691 | | 2 | -529 |
| 6 | cpubound 3 | ready | 1 | 3 | 0x0FDEF7FC | 0x0FDEF544 | 1024 | 702 | | 5 | -358 |
| 7 | cpubound 4 | ready | 1 | 3 | 0x0FDEF3FC | 0x0FDEF174 | 1024 | 691 | | 4 | -341 |
| 8 | cpubound 5 | ready | 1 | 3 | 0x0EFC8FFC | 0x0EFC8D74 | 1024 | 250 | | 1 | -220 |
| 9 | cpubound 6 | ready | 1 | 3 | 0x0EFC8BFC | 0x0EFC8974 | 1024 | 301 | | 0 | -205 |
| 0 | prnull | ready | 0 | 0 | 0x0EFDEFFC | 0x0EFDEF40 | 8192 | 2147483647 | | 6 | -2147483647 |
| 1 | rdspc | wait | 200 | 0 | 0x0EFD0FFC | 0x0EFDCA9C | 16384 | 2 | | | |
| 3 | Main process | curr | 20 | 0 | 0x0EFD8FFC | 0x0EFD8AD0 | 65536 | 302 | | | |
| 4 | cpubound 1 | ready | 1 | 3 | 0x0FDEF7FC | 0x0FDEFD74 | 1024 | 690 | | 3 | -486 |
| 5 | cpubound 2 | ready | 1 | 3 | 0x0FDEFBFC | 0x0FDEF970 | 1024 | 691 | | 2 | -481 |
| 6 | cpubound 3 | ready | 1 | 3 | 0x0FDEF7FC | 0x0FDEF544 | 1024 | 702 | | 5 | -510 |
| 7 | cpubound 4 | ready | 1 | 3 | 0x0FDEF3FC | 0x0FDEF174 | 1024 | 691 | | 4 | -493 |
| 8 | cpubound 5 | ready | 1 | 3 | 0x0EFC8FFC | 0x0EFC8D70 | 1024 | 340 | | 0 | -274 |
| 9 | cpubound 6 | ready | 1 | 3 | 0x0EFC8BFC | 0x0EFC8974 | 1024 | 391 | | 1 | -301 |
| 0 | prnull | ready | 0 | 0 | 0x0EFDEFFC | 0x0EFDEF40 | 8192 | 2147483647 | | 6 | -2147483647 |
| 1 | rdspc | wait | 200 | 0 | 0x0EFD0FFC | 0x0EFDCA9C | 16384 | 2 | | | |
| 3 | Main process | curr | 20 | 0 | 0x0EFD8FFC | 0x0EFD8AD0 | 65536 | 392 | | | |
| 4 | cpubound 1 | ready | 1 | 3 | 0x0FDEF7FC | 0x0FDEFD74 | 1024 | 690 | | 2 | -414 |
| 5 | cpubound 2 | ready | 1 | 3 | 0x0FDEFBFC | 0x0FDEF974 | 1024 | 721 | | 1 | -409 |
| 6 | cpubound 3 | ready | 1 | 3 | 0x0FDEF7FC | 0x0FDEF544 | 1024 | 702 | | 5 | -438 |
| 7 | cpubound 4 | ready | 1 | 3 | 0x0FDEF3FC | 0x0FDEF174 | 1024 | 721 | | 3 | -421 |
| 8 | cpubound 5 | ready | 1 | 3 | 0x0EFC8FFC | 0x0EFC8D74 | 1024 | 430 | | 0 | -394 |
| 9 | cpubound 6 | ready | 1 | 3 | 0x0EFC8BFC | 0x0EFC8974 | 1024 | 481 | | 4 | -421 |
| 0 | prnull | ready | 0 | 0 | 0x0EFDEFFC | 0x0EFDEF40 | 8192 | 2147483647 | | 6 | -2147483647 |
| 1 | rdspc | wait | 200 | 0 | 0x0EFD0FFC | 0x0EFDCA9C | 16384 | 2 | | | |
| 3 | Main process | curr | 20 | 0 | 0x0EFD8FFC | 0x0EFD8AD0 | 65536 | 482 | | | |
| 4 | cpubound 1 | ready | 1 | 3 | 0x0FDEF7FC | 0x0FDEFD70 | 1024 | 720 | | 3 | -666 |
| 5 | cpubound 2 | ready | 1 | 3 | 0x0FDEFBFC | 0x0FDEF974 | 1024 | 721 | | 2 | -661 |
| 6 | cpubound 3 | ready | 1 | 3 | 0x0FDEF7FC | 0x0FDEF574 | 1024 | 731 | | 5 | -701 |
| 7 | cpubound 4 | ready | 1 | 3 | 0x0FDEF3FC | 0x0FDEF174 | 1024 | 721 | | 4 | -673 |
| 8 | cpubound 5 | ready | 1 | 3 | 0x0EFC8FFC | 0x0EFC8D74 | 1024 | 550 | | 0 | -484 |

Figure 6: Dynamic Workloads