

# Compiler Design Lab

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**SECTION:** K1

## Exercise 5: FIRST and FOLLOW

**AIM:** To compute the FIRST and FOLLOW of a grammar

### INTRODUCTION:

FIRST and FOLLOW are two functions associated with grammar that help us fill in the entries of an M-table.

FIRST() – It is a function that gives the set of terminals that begin the strings derived from the production rule.

A symbol  $c$  is in FIRST ( $\alpha$ ) if and only if  $\alpha \Rightarrow c\beta$  for some sequence  $\beta$  of grammar symbols.

A terminal symbol  $a$  is in FOLLOW ( $N$ ) if and only if there is a derivation from the start symbol  $S$  of the grammar such that  $S \Rightarrow \alpha N \alpha \beta$ , where  $\alpha$  and  $\beta$  are a (possibly empty) sequence of grammar symbols. In other words, a terminal  $c$  is in FOLLOW ( $N$ ) if  $c$  can follow  $N$  at some point in a derivation.

### Computation of FIRST

FIRST ( $\alpha$ ) is defined as the collection of terminal symbols which are the first letters of strings derived from  $\alpha$ .

$\text{FIRST}(\alpha) = \{a \mid \alpha \rightarrow^* a\beta \text{ for some string } \beta\}$

If  $X$  is Grammar Symbol, then First ( $X$ ) will be –

- If  $X$  is a terminal symbol, then  $\text{FIRST}(X) = \{X\}$
- If  $X \rightarrow \epsilon$ , then  $\text{FIRST}(X) = \{\epsilon\}$
- If  $X$  is non-terminal &  $X \rightarrow a \alpha$ , then  $\text{FIRST}(X) = \{a\}$
- If  $X \rightarrow Y_1, Y_2, Y_3$ , then  $\text{FIRST}(X)$  will be

(a) If  $Y$  is terminal, then

$\text{FIRST}(X) = \text{FIRST}(Y_1, Y_2, Y_3) = \{Y_1\}$

(b) If  $Y_1$  is Non-terminal and

If  $Y_1$  does not derive to an empty string i.e., If  $\text{FIRST}(Y_1)$  does not contain  $\epsilon$  then,  $\text{FIRST}(X) = \text{FIRST}(Y_1, Y_2, Y_3) = \text{FIRST}(Y_1)$

(c) If  $\text{FIRST}(Y_1)$  contains  $\epsilon$ , then.

$\text{FIRST}(X) = \text{FIRST}(Y_1, Y_2, Y_3) = \text{FIRST}(Y_1) - \{\epsilon\} \cup \text{FIRST}(Y_2, Y_3)$

Similarly,  $\text{FIRST}(Y_2, Y_3) = \{Y_2\}$ , If  $Y_2$  is terminal otherwise if  $Y_2$  is Non-terminal then

- $\text{FIRST}(Y_2, Y_3) = \text{FIRST}(Y_2)$ , if  $\text{FIRST}(Y_2)$  does not contain  $\epsilon$ .
- If  $\text{FIRST}(Y_2)$  contain  $\epsilon$ , then
- $\text{FIRST}(Y_2, Y_3) = \text{FIRST}(Y_2) - \{\epsilon\} \cup \text{FIRST}(Y_3)$

Similarly, this method will be repeated for further Grammar symbols, i.e., for Y4, Y5, Y6 ... . YK.

#### Computation of FOLLOW

Follow (A) is defined as the collection of terminal symbols that occur directly to the right of A.

$\text{FOLLOW}(A) = \{a \mid S \Rightarrow^* \alpha A a \beta \text{ where } \alpha, \beta \text{ can be any strings}\}$

Rules to find FOLLOW

- If S is the start symbol,  $\text{FOLLOW}(S) = \{\$ \}$
- If production is of form  $A \rightarrow \alpha B \beta$ ,  $\beta \neq \epsilon$ .

(a) If  $\text{FIRST}(\beta)$  does not contain  $\epsilon$  then,  $\text{FOLLOW}(B) = \{\text{FIRST}(\beta)\}$

Or

(b) If  $\text{FIRST}(\beta)$  contains  $\epsilon$  (i. e.,  $\beta \Rightarrow^* \epsilon$ ), then

$\text{FOLLOW}(B) = \text{FIRST}(\beta) - \{\epsilon\} \cup \text{FOLLOW}(A)$

∴ when  $\beta$  derives  $\epsilon$ , then terminal after A will follow B.

- If production is of form  $A \rightarrow \alpha B$ , then  $\text{Follow}(B) = \{\text{FOLLOW}(A)\}$ .

#### ALGORITHM FOR FIRST:

1. If X is a terminal then  $\text{FIRST}(X) = \{X\}$   
Example:  $F \rightarrow I \mid id$  We can write it as  $\text{FIRST}(F) \rightarrow \{ ( , id )$
2. If X is a non-terminal like  $E \rightarrow T$  then to get FIRST substitute T with other productions until you get a terminal as the first symbol.
3. If  $X \rightarrow \epsilon$  then add  $\epsilon$  to  $\text{FIRST}(X)$ .

#### ALGORITHM FOR FOLLOW:

1. Always check the right side of the productions for a non-terminal, whose FOLLOW set is being found. (never see the left side).
2. (a) If that non-terminal (S,A,B...) is followed by any terminal (a,b...,\*,+,(),...) , then add that terminal into the FOLLOW set.  
(b) If that non-terminal is followed by any other non-terminal then add FIRST of other non-terminal into the FOLLOW set.

#### PROGRAM:

```
#include<stdio.h>
#include<string.h>
#include<ctype.h>
int n,m=0,p,i=0,j=0;
char a[10][10],f[10];
void follow(char c);
void first(char c);
int main(){
int i,z;
```

```

char c,ch;
printf("Enter the no of prooductions:\n");
scanf("%d",&n);
printf("Enter the productions:\n");
for(i=0;i<n;i++)
scanf("%s%c",a[i],&ch);
do{
m=0;
printf("Enter the elemets whose fisrt & follow is to be found:");
scanf("%c",&c);
first(c);
printf("First(%c)={",c);
for(i=0;i<m;i++)
printf("%c",f[i]);
printf("}\n");
strcpy(f," ");
m=0;
follow(c);
printf("Follow(%c)={",c);
for(i=0;i<m;i++)
printf("%c",f[i]);
printf("}\n");
printf("Continue(0/1)?");
scanf("%d%c",&z,&ch);
}while(z==1);
return(0);
}

void first(char c){
int k;
if(!isupper(c))
f[m++]=c;
for(k=0;k<n;k++){
if(a[k][0]==c){
if(a[k][2]=='$')
follow(a[k][0]);
else if(islower(a[k][2]))
f[m++]=a[k][2];
else first(a[k][2]);
}
}

```

```

}
}
void follow(char c){
if(a[0][0]==c)
f[m++]='$';
for(i=0;i<n;i++){
for(j=2;j<strlen(a[i]);j++){
if(a[i][j]==c){
if(a[i][j+1]!='\0')
first(a[i][j+1]);
if(a[i][j+1]!='\0' && c!=a[i][0])
follow(a[i][0]);
}
}
}
}

```

**INPUT:**

E→E+T  
E→E+T

Manual Output:

Productions:

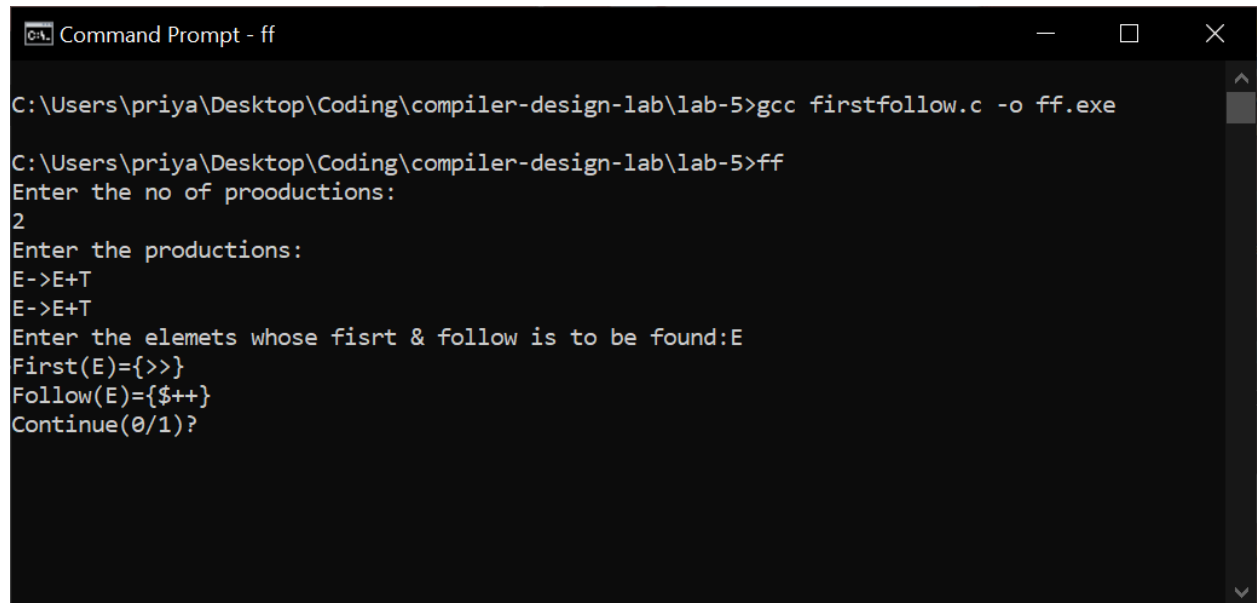
$E \rightarrow E+T$

$E \rightarrow E+T$

$FIRST(E) = \{ + \}$

$FOLLOW(E) = \{ \$, + \}$

## System Output:



```
Command Prompt - ff

C:\Users\priya\Desktop\Coding\compiler-design-lab\lab-5>gcc firstfollow.c -o ff.exe

C:\Users\priya\Desktop\Coding\compiler-design-lab\lab-5>ff
Enter the no of prooductions:
2
Enter the productions:
E->E+T
E->E+T
Enter the elemets whose fisrt & follow is to be found:E
First(E)={>+}
Follow(E)={$++}
Continue(0/1)?
```

## RESULT:

FIRST and FOLLOW for a set of productions has been computed.