

N-Queen Problem:

The N queen problem is to find all the ways to place N non-attacking queens in a $n \times n$ chess board.

The solution vector of N queen problem is represented by $\{x[1], x[2], \dots, x[n]\}$ in which $x[i]$ is the column of i^{th} row where i^{th} queen is placed.

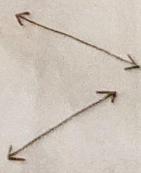
$x[i]$ will be distinct since no two queens can be placed in the same column or same diagonal. The column

The columnwise check of non-attacking queens can be done by placement of previous $k-1$ queens

The diagonalwise check of non-attacking queens can be done in two forms

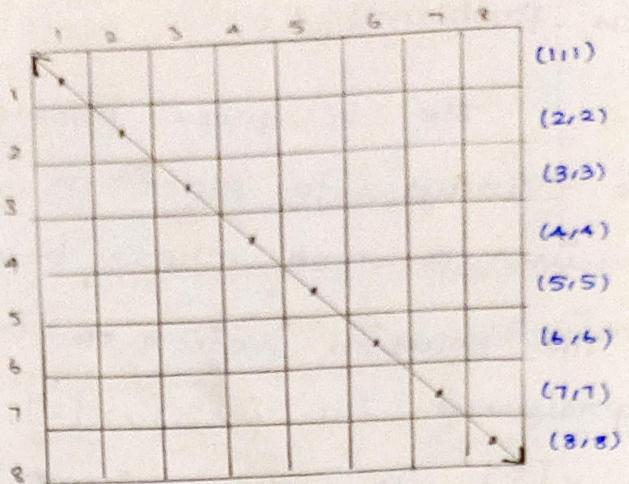
i) upper left to lower right

ii) lower left to upper right



upper left to lower right :

consider a 8×8 chess board



Let $[i, j]$ be the position of previous Queen placed in chess board.

Let $[k, l]$ be the position of new Queen ie) (current queen) yet to be placed, then $[i, j]$ and $[k, l]$ are on the same diagonal only if,

$$\begin{aligned} \boxed{j - i = l - k} \\ \Rightarrow \boxed{j - l = i - k} \quad \text{--- (1)} \end{aligned}$$

example:

Let $[2, 2]$ be the previous position of queen ie) $[i, j] = [2, 2]$

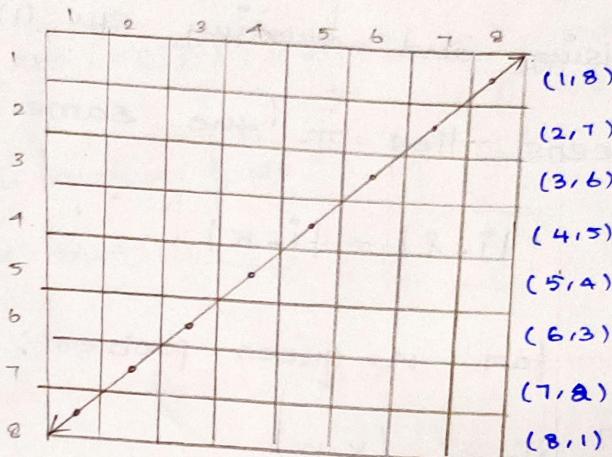
Let $[3, 3]$ be a position of new queen yet to be placed. ie) $[k, l] = [3, 3]$

If $[2, 2]$ and $[3, 3]$ lie on the same diagonal, then

$$\begin{aligned} j - l &= i - k \\ \Rightarrow 2 - 3 &= 2 - 3 \end{aligned}$$

$$\Rightarrow -1 = -1$$

lower left to upper right:



let $[i, j]$ be the position of previous queen placed in chess board.

let $[k, l]$ be the position of new queen (current queen) yet to be placed, then $[i, j]$ and $[k, l]$ lie on the same diagonal only if,

$$i + j = k + l$$

$$\Rightarrow j - l = k - i \quad \text{--- (2)}$$

example :

let $[2, 7]$ be the previous position of queen ie) $[i, j] = [2, 7]$

let $[3, 6]$ be the position of new queen yet to be placed. ie) $[k, l] = [3, 6]$

If $[2, 7] \neq [3, 6]$ lie on the same diagonal, then

$$j - l = k - i$$

$$\Rightarrow 7 - 6 = 3 - 2$$

$$\Rightarrow \boxed{1 = 1}$$

generalising and equating eqn (1) & (2)
two queens lies on the same diagonal or

if $|j - k| = |i - l|$

Algorithm for n-queen problem :

Algorithm NQueen (k, n)

// possible placement n -non attacking queens in

// a $n \times n$ chess board using Backtrack

{
for $i := 1$ to n do

{

if (place (k, i)) then

{

$x[k] := i;$

if ($k = n$) then write ($x[1:n]$);

else

NQueen ($k+1, n$);

}

3

3

Algorithm place (k, i)

// return true if queen can be placed in

// k^{th} row and i^{th} column.

// $x[]$ - a global array whose first $k-1$ values is
(known)

// Abs(j) - returns the Absolute value of j (only negative)

{

for $j := 1$ to $k-1$ do

{ if ($x[j] = i$) // two queens free in the same column.

or ($\text{Abs}(x[j]-i) = \text{Abs}(j-k)$) // two queens free
// in the same diagonal

then return false;

}

return true

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state space tree for 4 queen problem:

A 4 queen problem is to place 4

non-attacking queens in a 4×4 chess board such that no two queens should be placed in same row, column or same diagonal.

Two optimal solutions are provided for

4 queen Problem

solutions below are

solu 1:

		Q	
			Q
	Q		
			Q

$x_1 = 2$
 $x_2 = 4$
 $x_3 = 1$
 $x_4 = 3$

$i = j \leftarrow 1 = 1$

solu 2:

1st soln:

		Q	
			Q
	Q		
			Q

$x_1 = 3$
 $x_2 = 1$
 $x_3 = 4$
 $x_4 = 2$

$i = j \leftarrow 0 = 1$

2nd soln:

		Q	
			Q
	Q		
			Q

$x_1 = 2$
 $x_2 = 3$
 $x_3 = 1$
 $x_4 = 4$

$i = j \leftarrow 0 = 1$

at last, 3 level to start a root

