



SRM Institute of Science and Technology

College of Engineering and Technology

School of Computing

SRM Nagar, Kattankulathur – 603203, Chengalpattu District, Tamilnadu **SET-2**

Academic Year: 2022-23 (EVEN)

Test: CLA-T2 Date: 28-03-2023 Course Code & Title: **18CSC205J & Operating Systems**

Duration: 2 Hours Year & Sem: **II Year / V Sem** Max. Marks: 50


Course Articulation Matrix:

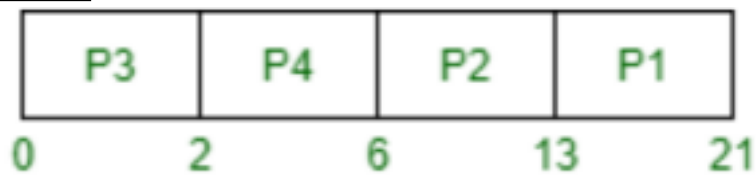
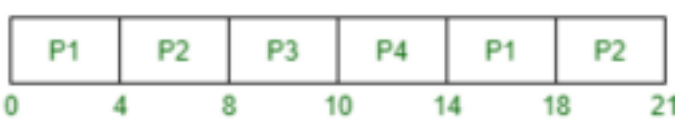
S.No.	Course Outcome	PO1	PO2	PO3	PO4	PO5	PO6	PO7	PO8	PO9	PO10	PO11	PO12
1	CO1	3		3									
2	CO2	2	1	3									
3	CO3	3	2	2									
4	CO4	3	2	2									
5	CO5	3		2	2								

Part – B
(5 x 10= 50 Marks)

Instructions: Answer any Five Questions

Q · N o	Question	Mar ks	B L	C O	PO	PI Cod e
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1	<p>Consider the following set of processes with their arrival and burst times that uses First-Come, First Served scheduling (FCFS) policy.</p> <p><i>Process Arrival Time Burst Time</i></p> <p>P1 0 20 P2 25 25 P3 30 105 P4 60 15 P5 100 10 P6 105 10</p> <p>i. Show the scheduling order of the processes using a Gantt chart. ii. Evaluate the average waiting time of processes. iii. Evaluate the average turnaround time of processes.</p> <p><u>Solution</u> Gantt chart (4)</p> <p>Show the scheduling order of the processes using a Gantt chart.</p>  <p>Find out the average waiting time of processes. (3)</p> <p>Waiting time for processes are - P1=0, P2=0, P3=20, P4=95, P5=70, P6=75 Therefore, average waiting time= $(0+0+20+95+70+75)/6 = 43.33$ ms</p> <p>Find out the average turnaround time of processes. (3)</p> <p>Turnaround time for processes is – P1=20, P2=25, P3=125, P4=110, P5=80, P6=85 Therefore, average turnaround time= $(20+25+125+110+80+85)/6 = 74.17$ ms</p>	10	4	2	PO 2	2.5. 3
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2	<p>Consider the following set of processes, assumed to have arrived at time 0. Consider the CPU scheduling algorithms Shortest Job First (SJF) and Round Robin (RR). For RR, assume that the processes are scheduled in the order P₁, P₂, P₃, P₄. If the time quantum for RR is 4 ms, then calculate the absolute value of the difference between the average turnaround times (in ms) of SJF and RR (round off to 2 decimal places).</p> <p>Processes P₁ P₂ P₃ P₄ Burst Time 8 7 2 4</p> <p>(in ms)</p> <p>Solution:</p>  <p>Therefore, Average Turn Around Time (TAT) is,</p> $= \{(21 - 0) + (13 - 0) + (6 - 0) + (2 - 0)\} / 4$ $= 10.5$  <p>Therefore, Average Turn Around Time (TAT) is,</p> $= \{(18 - 0) + (21 - 0) + (10 - 0) + (14 - 0)\} / 4$ $= 15.75$ <p>Hence,</p> $= \text{SJF (TAT)} - \text{RR(TAT)}$ $= 10.5 - 15.75$ $= 5.25$	10	4	2	PO 2	2.5. 3
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3	<p>A computer system uses the Banker's Algorithm to deal with deadlocks. Its current state is shown in the tables below, where P0, P1, P2 are processes and R0, R1, R2 are resource types.</p> <table><tr><th colspan="3">Maximum Need</th><th colspan="3">Current Allocation</th><th colspan="3">Available</th></tr><tr><th></th><th>R0</th><th>R1</th><th>R2</th><th></th><th>R0</th><th>R1</th><th>R2</th><th>R0</th><th>R1</th><th>R2</th></tr><tr><td>P0</td><td>4</td><td>1</td><td>2</td><td>P0</td><td>1</td><td>0</td><td>2</td><td>2</td><td>2</td><td>0</td></tr><tr><td>P1</td><td>1</td><td>5</td><td>1</td><td>P1</td><td>0</td><td>3</td><td>1</td><td></td><td></td><td></td></tr><tr><td>P2</td><td>1</td><td>2</td><td>3</td><td>P2</td><td>1</td><td>0</td><td>2</td><td></td><td></td><td></td></tr></table> <p>(a) Show that the system can be in safe state.</p> <p>(b) What will the system do on a request by process P0 for one unit of resource type R1?</p> <p><u>Solution:</u></p> <p>(a) System is in the safe state. The safe state is "P1,P2,P0" (b) By granting the request of P0, the system goes into unsafe state so system would delay the request</p>	Maximum Need			Current Allocation			Available				R0	R1	R2		R0	R1	R2	R0	R1	R2	P0	4	1	2	P0	1	0	2	2	2	0	P1	1	5	1	P1	0	3	1				P2	1	2	3	P2	1	0	2				10	4	2	PO 2	2.5. 3
Maximum Need			Current Allocation			Available																																																					
	R0	R1	R2		R0	R1	R2	R0	R1	R2																																																	
P0	4	1	2	P0	1	0	2	2	2	0																																																	
P1	1	5	1	P1	0	3	1																																																				
P2	1	2	3	P2	1	0	2																																																				

4	<p>A single processor system has three resource types X, Y and Z, which are shared by three processes. There are 5 units of each resource type. Consider the following scenario, where the column alloc denotes the number of units of each resource type allocated to each process, and the column request denotes the number of units of each resource type requested by a process in order to complete execution. Which of these processes will finish LAST?</p> <table> <tr> <th></th><th colspan="3">alloc</th><th colspan="3">request</th></tr> <tr> <th></th><th>X</th><th>Y</th><th>Z</th><th>X</th><th>Y</th><th>Z</th></tr> <tr> <td>P0</td><td>1</td><td>2</td><td>1</td><td>1</td><td>0</td><td>3</td></tr> <tr> <td>P1</td><td>2</td><td>0</td><td>1</td><td>0</td><td>1</td><td>2</td></tr> <tr> <td>P2</td><td>2</td><td>2</td><td>1</td><td>1</td><td>2</td><td>0</td></tr> </table> <p><u>Solution:</u></p> <p>A single processor system has three resource types X, Y and Z, which are shared by three processes. There are 5 units of each resource type.</p> <p>So, the resource instances which are left being unallocated = $\langle 0,1,2 \rangle$ { unallocated resources= total resources-allocated resources }</p> <p>now, from the request table, you can say that only request of P1 can be satisfied. So P1 can finish its execution first. Once P1 is done, it releases 2, 0 and 1 units of X, Y and Z respectively which were allocated to P1. So, Now unallocated resource instance are $= \langle 0,1,2 \rangle + \langle 2,0,1 \rangle = \langle 2,1,3 \rangle$</p> <p>Now Among P0 and P2, needs of P0 can only be satisfied. So P0 finishes its execution. $\langle 2,1,3 \rangle + \langle 1,2,1 \rangle = \langle 3,3,4 \rangle$.</p> <p>Finally, P2 finishes its execution. So, P2 is the process which finishes in end.</p>		alloc			request				X	Y	Z	X	Y	Z	P0	1	2	1	1	0	3	P1	2	0	1	0	1	2	P2	2	2	1	1	2	0	10	4	2	PO 2	2.5. 3
	alloc			request																																					
	X	Y	Z	X	Y	Z																																			
P0	1	2	1	1	0	3																																			
P1	2	0	1	0	1	2																																			
P2	2	2	1	1	2	0																																			

6	<p>The following program consists of 3 concurrent processes and 3 binary semaphores. The semaphores are initialized as $S_0 = 1$, $S_1 = 0$, $S_2 = 0$.</p> <table border="1"><thead><tr><th>Process P0</th><th>Process P1</th><th>Process P2</th></tr></thead><tbody><tr><td><pre>while (true) { wait (S0); print '0'; release (S1); release (S2); }</pre></td><td><pre>wait (S1); release (S0);</pre></td><td><pre>wait (S2); release (S0);</pre></td></tr></tbody></table> <p>Find,</p> <p>i) Minimum number of times P0 printing 0 (ZERO) ii) Maximum number of times P0 printing 0 (ZERO) Justify your Answer with explanation.</p> <p>Solution:</p> <p>a) Initially only P0 can go inside the while loop as $S_0 = 1$, $S_1 = 0$, $S_2 = 0$. Minimum no. of time 0 printed is twice when execute in this order (p0 -> p1 -> p2 -> p0)</p> <p>b) Maximum no. of time 0 printed is thrice when execute in this order (p0 -> p1 -> p0 -> p2 -> p0).</p>	Process P0	Process P1	Process P2	<pre>while (true) { wait (S0); print '0'; release (S1); release (S2); }</pre>	<pre>wait (S1); release (S0);</pre>	<pre>wait (S2); release (S0);</pre>	10	3	2	PO 1	1.6. 1
Process P0	Process P1	Process P2										
<pre>while (true) { wait (S0); print '0'; release (S1); release (S2); }</pre>	<pre>wait (S1); release (S0);</pre>	<pre>wait (S2); release (S0);</pre>										
7	<p>Consider a system with byte-addressable memory, 32-bit logical addresses, 4KB page size and page table entries of 4 bytes each. Calculate the process size and the size of page table.</p> <p>Solution Given:</p> <ul style="list-style-type: none">• Number of bits in logical address = 32 bits• Page size = 4KB• Page table entry size = 4 bytes <p>Process Size Number of bits in logical address = 32 bits Thus, Process size = 2^{32} B = 4 GB</p> <p>Number of Entries in Page Table Number of pages the process is divided = Process size / Page size size = 4 GB / 4 KB = 220 pages</p> <p>Thus, Number of entries in page table = 220 entries</p>	10	3	3	PO 2	2.5. 3						

<p>Page Table Size Page table size = Number of entries in page table x Page table entry size = 220×4 bytes = 4 MB</p>					
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*Program Indicators are available separately for Computer Science and Engineering in AICTE examination reforms policy.

Course Outcome (CO) and Bloom's level (BL) Coverage in Questions

Approved by the Audit Professor/Course Coordinator

