

Dissecting the Modern Android Data Encryption Scheme

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Who we are



- [@DamianoMelotti](#)
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- Interested in low-level mobile security and fuzzing

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- Security researcher and R&D leader @ Quarkslab
- Working on mobile and embedded software security



The trigger

> Hey! My device fell into water and the main SoC is dead. However the Titan M^{1,2} chip seems to be alive and well, do you think you would be able to help me recover my data on the phone?

[1]: [2021: A Titan M Odyssey](#) (Maxime Rossi Bellom, Philippe Teuwen, Damiano Melotti)

[2]: [Attack on Titan M, Reloaded: Vulnerability Research on a Modern Security Chip](#) (Damiano Melotti, Maxime Rossi Bellom)

The trigger

- Our answer: no, the main SoC is still essential for disk encryption/decryption
 - ... but up to what extent?
- Objective of this research: find out exactly
- Offensive approach:
 - What would a forensic analyst do?
 - Assuming infinite vulnerabilities, what can you do to get the secrets out?
 - Do you still need to bruteforce credentials?

Data Encryption at Rest 101

- Idea: no sensitive plaintext files in storage
 - Attackers must not find files in clear on disk
- Threat model: full physical access to powered-off device
- Data is automatically encrypted when written and automatically decrypted when read
- How?
 - Android: Full-Disk Encryption and File-Based Encryption (required from Android 10)
 - Underneath: dm-crypt for FDE, fscrypt for FBE

File-Based Encryption at Rest 101

- Relies on fscrypt, implemented in the Linux kernel
 - It supports Ext4, F2FS, and UBIFS
- Operates at the filesystem level
 - Allows files encrypted with different keys or unencrypted in a file system
- A master key is provided for directory tree
- Then derived per file keys (for regular file, directory, and symbolic link)
- Metadata are not encrypted by fscrypt

Android File-Based Encryption

- Each file has its own key
- Direct Boot and multi-user support
- Two encryption levels:
 - Credential Encrypted (CE), available only after authentication
 - Device Encrypted (DE), available also during boot
- In short, 2 “main” keys
 - DE key, for data decrypted at boot
 - CE key, available after authentication, protecting user data
- DE key is automatically decrypted using HW-backed keys

Android File-Based Encryption

```
a22x:/ # ls /data/data | head -n 20
+BWqxAAAAAQGI050Z57bZxMl6oTCNKC
+geIWCAAAAgIsCJB+mpPqIQYOH0FrnojC1KJ8e0lvGZWJYXWFTc0WD
+LE39AAAAAwNvineHXItKqcE1Glo9+EinbUaa3wvp,fXEjaJ3r4BiC
+ostbBAAA AwZ2g592ei8GG2DMfc4y6H94jxkEfoRzUDdlQZn0YZKqA5VCUgi89sf2JN8yCBFraC
,0wF+BAAA AgQtSCvHd5rRmVC2lVxEHt1Eo50M9kma1oe+vkWT176K8dxofVp5RcmnLv0xC7WMLJ
,WGCADAAA Ag66o3+mZ7f009fBNp8zQSdWBqRJIwPUZHafQiTu7khxB
,eGnGDAAA Aw+ZAptd14KRyth5ncJmJkYZBTBW7DoUNpMamRGj05MA
,vu29CAAAA AgKApb3amvoChi0pYILwr5xxvUtf0TpZ2h6Cr0wG4CR5D
0MQ5GCAAAQAtXNLN524L3GwuGAek+rVelNjgE5mwqljis0kydZa2FKJlVD6ezJJGkcjcRTTD7B
0QcmpBAAA AwMxohHqP+AT3ktRtIAJk9, bu2g04xwjLzf, vgN71QQ2B
0WZ9mBAAA A5KuJUTYZ61eUrgnVJAloJnZDm3b7JybYYPw8xhLrf0D
11ppkCAAAA I4M8ENP6k0t3rA9kIN9bepZ0FdpiYDp6bQ5Cj2IU2NB
13nD7DAAA Aw, NoobtAOXesI0kFqFC4MIwoKYAfdBQrGTm8rfbbdsNC
1K0haDAAA AibCc dZY7LE6W4+DAXnAiv, PEhu4CiUL6ofn9ZuL6buC
1V6luAAA AwmrNINeM84y83, LYzeNEkdCcACNx lAhJ+sI7VzbV51yuebjw9M0PTTrjRi9cv391kD
1fMeDAAA AgPnPYDGu0WWC8uWiojo+nhhnqWk+x0ZerDfrWY3ZVIbB
1hGhtCAAAA AwfxP3KKXNUYQ0SWoBzqZcsbu2g04xwjLzf, vgN71QQ2B
2d3fjDAAA Agf7ilkMftxkJiD0JbRx9, dtg04Rk9L1P5CCl7JuqqewNpAl07TQq0KJh2yXzcj9, 8JfbjIHwnpEmj7FQ0ixXBP
2v, UXBAAA QloHyVK08uify8onfmrQXJawy1Dqwq6kc2g2BKr0H48D
2wjJhBAAA AgQ0MvGF3NqTpssmh6XgAWw1jHW986vSwQMXIPQQz6ghD
```

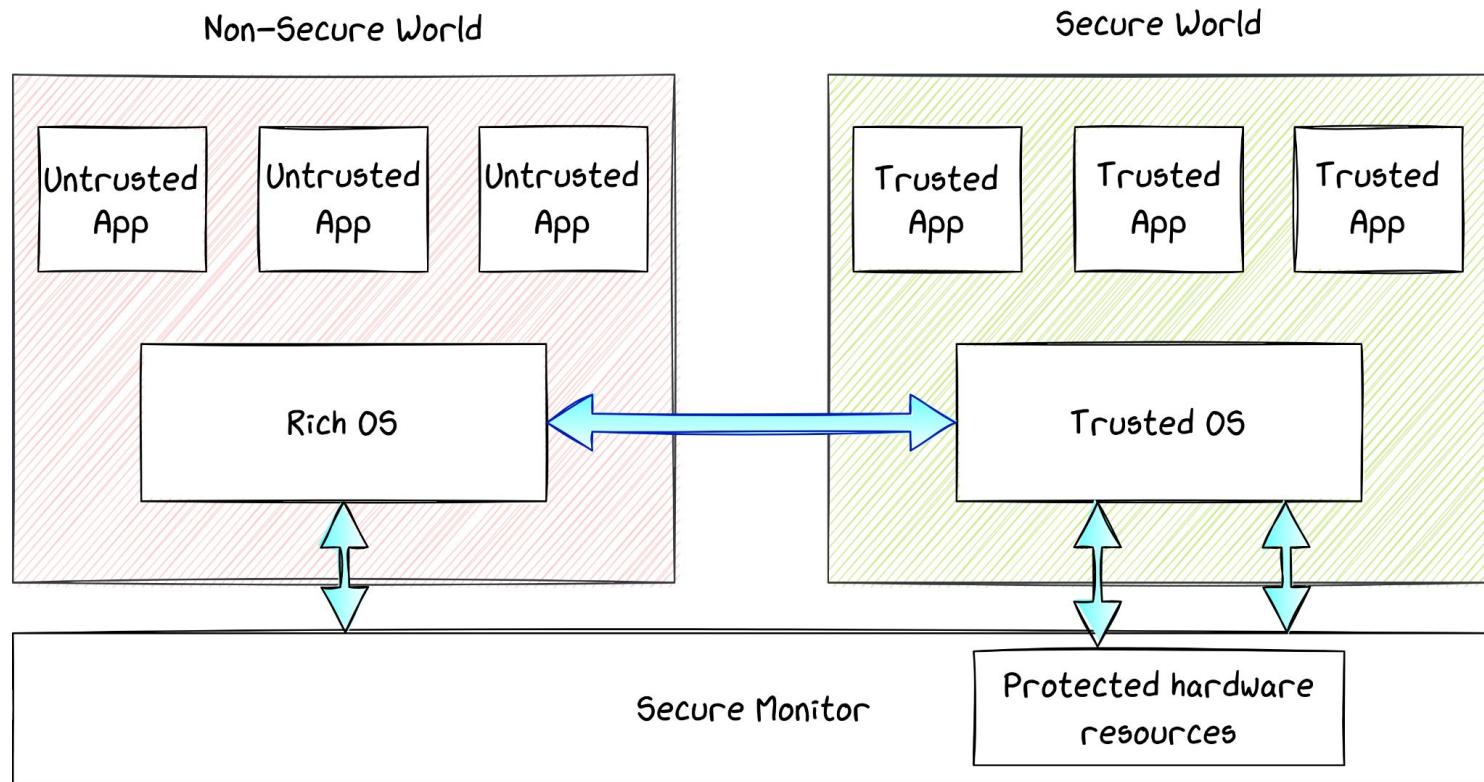
FBE key derivation

- We focus only on the CE key
- Complex derivation steps
 - Start from DE files owned by privileged users

```
a22x:/ # ls -l /data/system_de/0/spblob/
total 32
-rw----- 1 system system      58 2022-06-29 21:59 0000000000000000.handle
-rw----- 1 system system      72 2022-06-29 21:59 921e9ab09af8d9d.metrics
-rw----- 1 system system     93 2022-06-29 21:59 921e9ab09af8d9d.pwd
-rw----- 1 system system 16384 2022-06-29 21:59 921e9ab09af8d9d.secdis
-rw----- 1 system system    186 2022-06-29 21:59 921e9ab09af8d9d.spblob
```

- User credentials are used in the process
 - No matter how many bugs an attacker has, bruteforcing remains necessary!

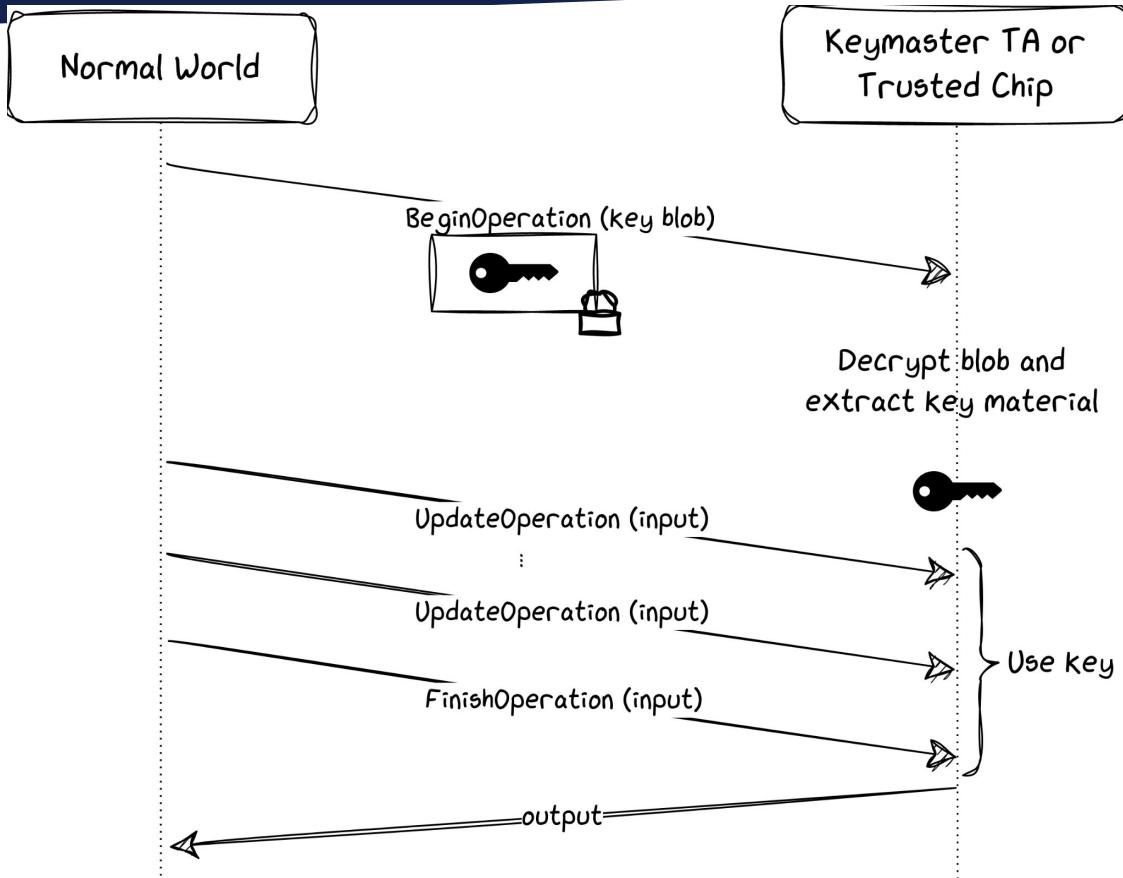
ARM TrustZone

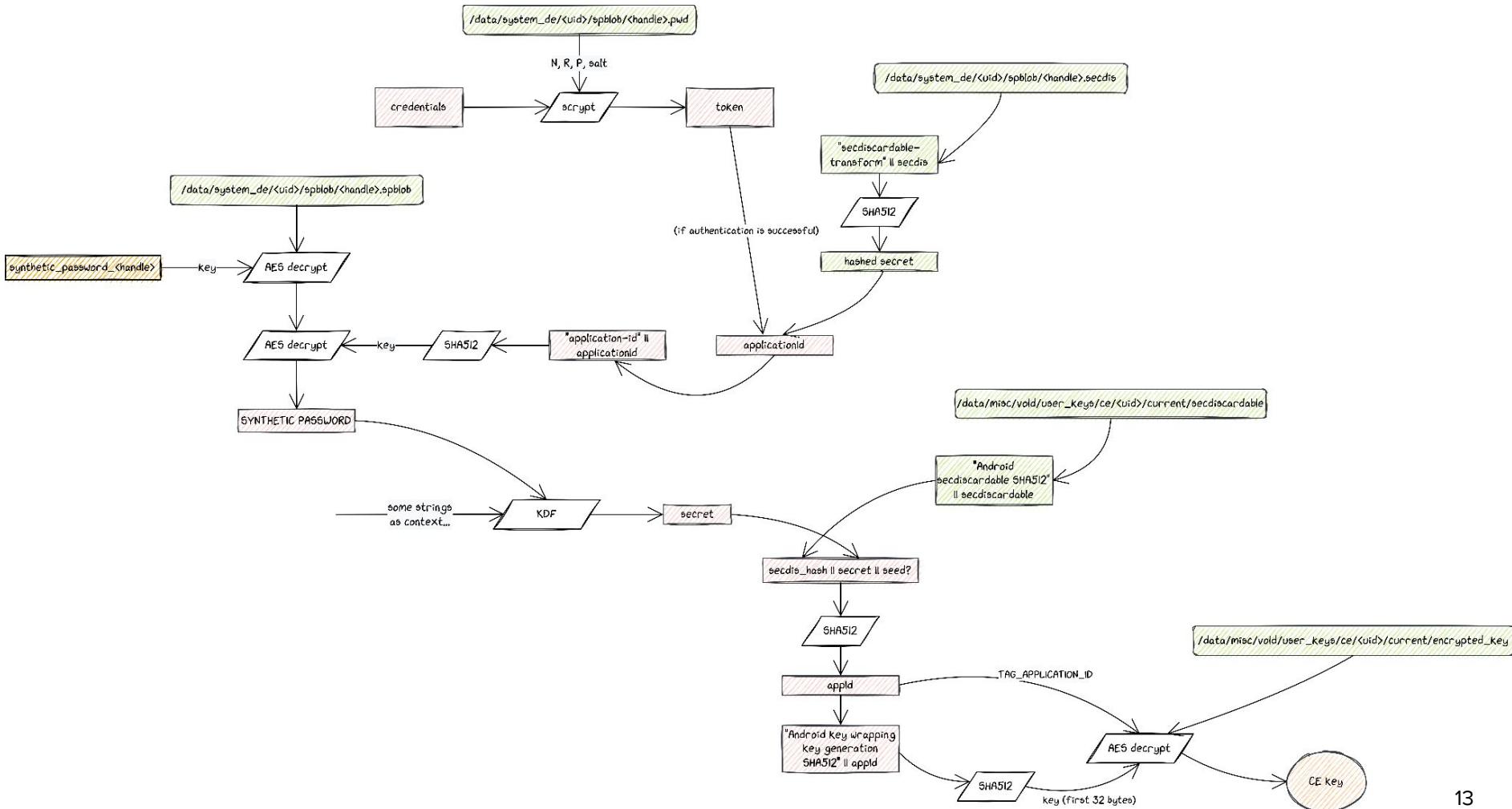


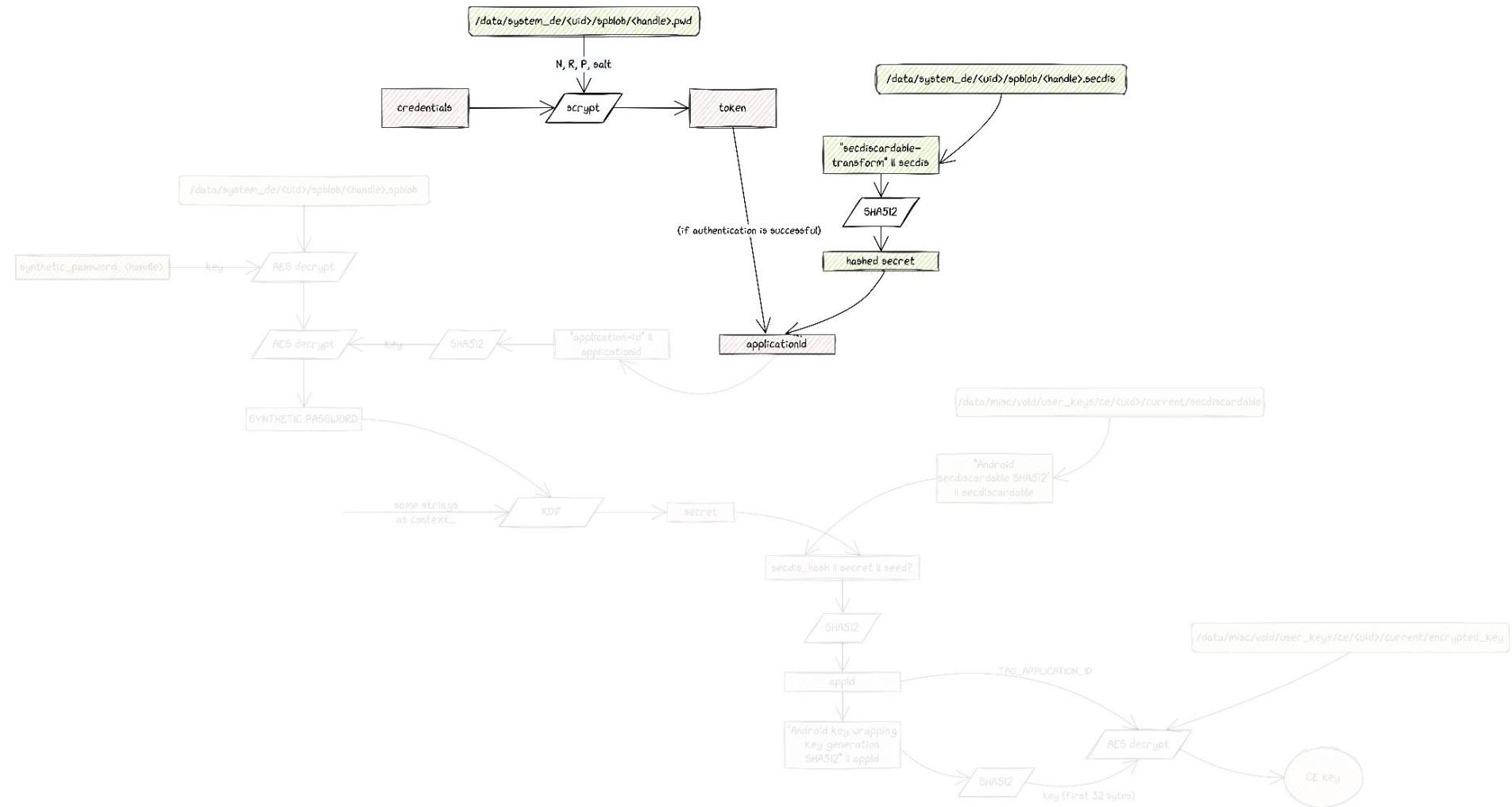
Android Keystore system

- Key storage and crypto services
- Keys are stored as *key blobs*
- Three protection levels:
 - Software only
 - TEE (default)
 - Hardware-backed (StrongBox)
- Raw key should never leave protected environment

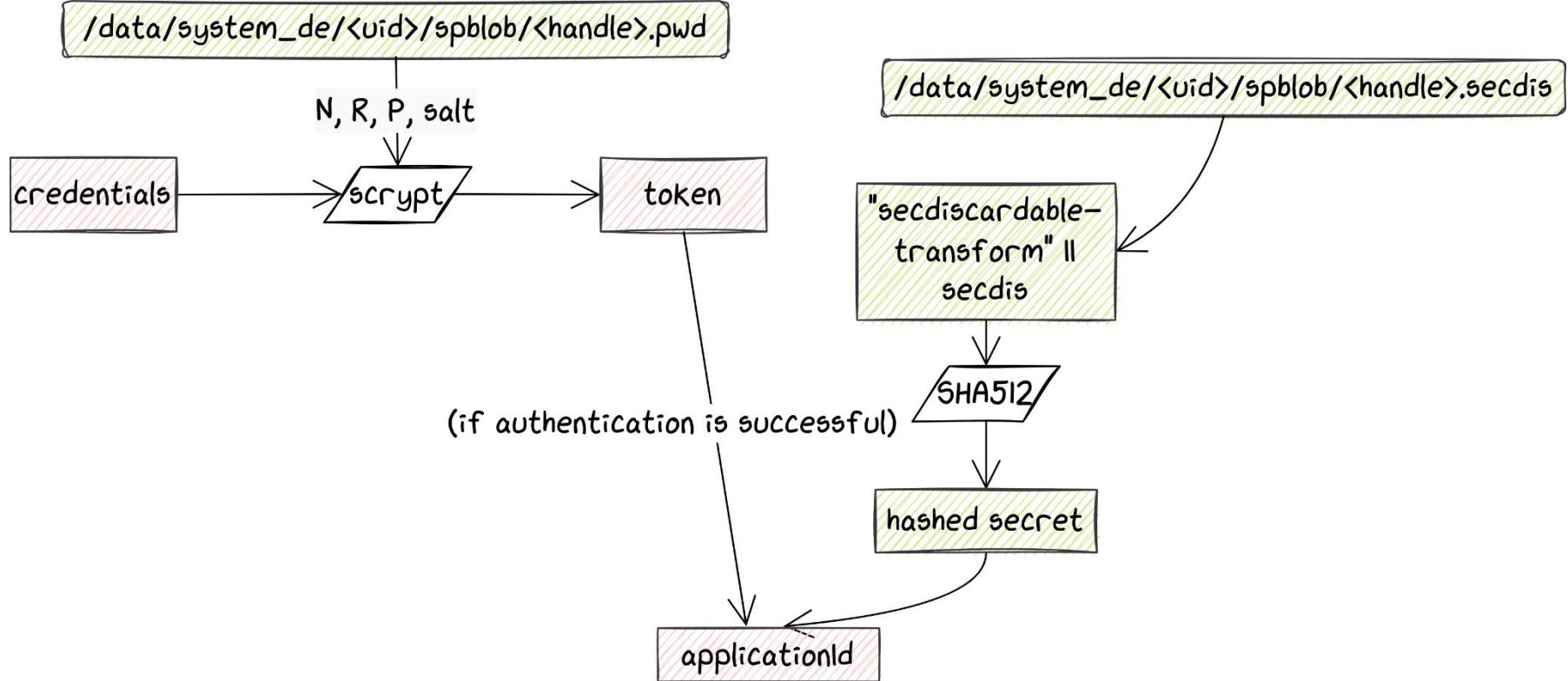
Android Keystore system







Credentials, scrypt, secdis

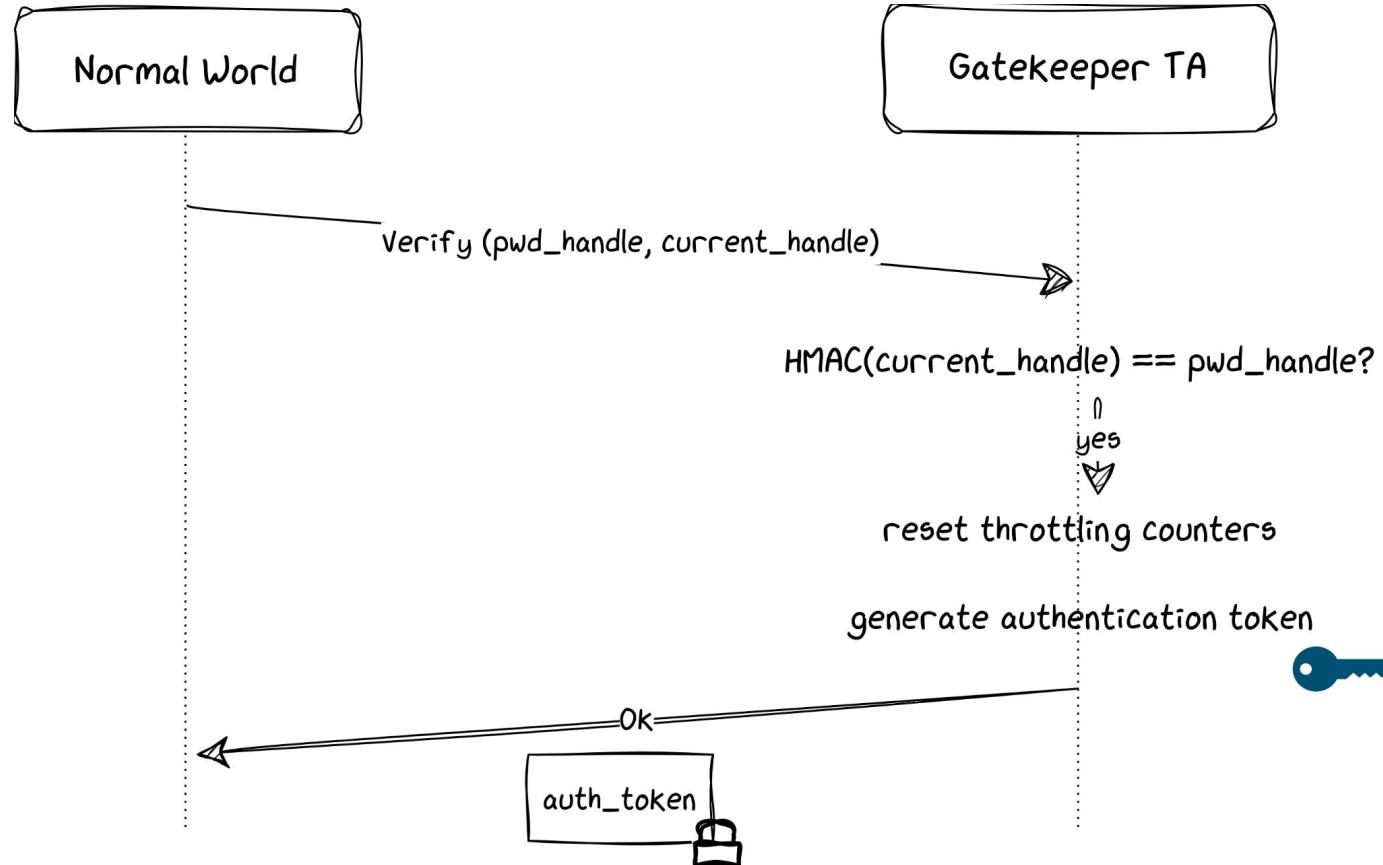


Authentication with Gatekeeper

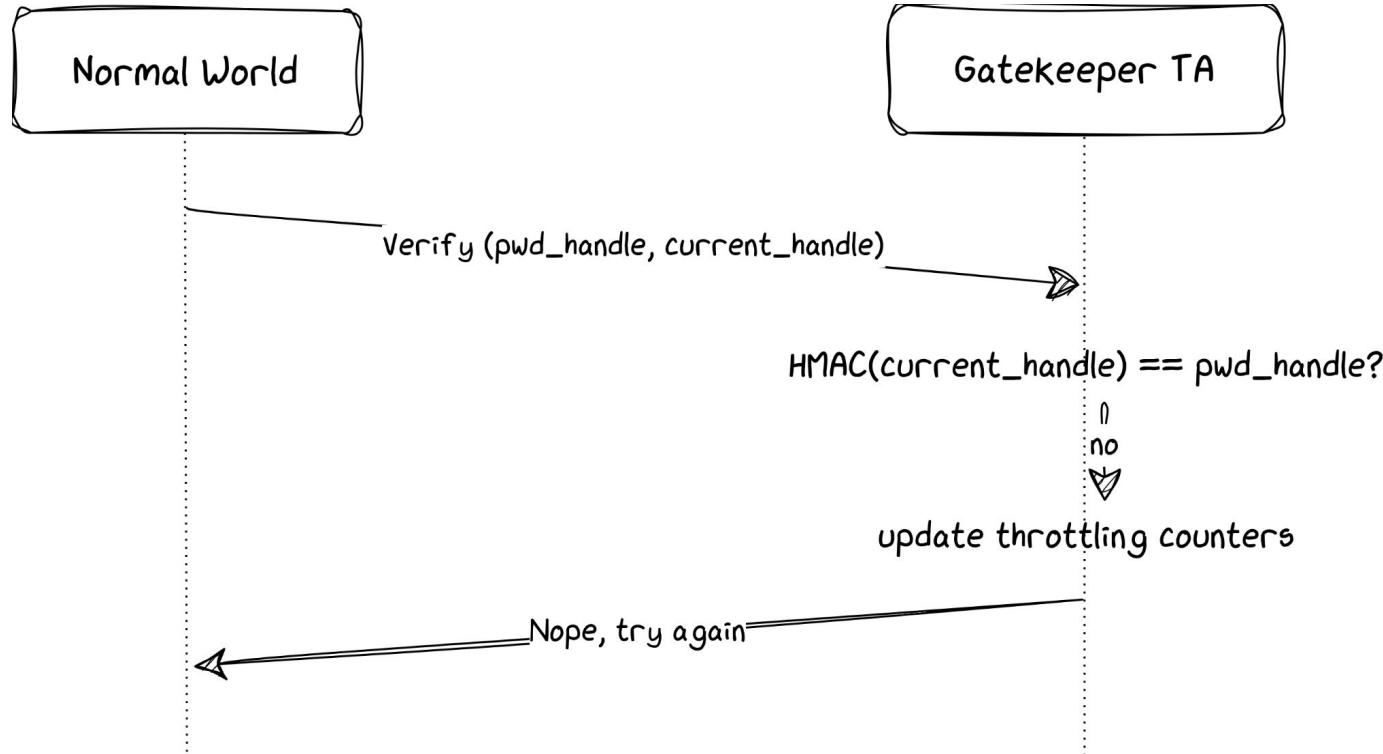
- The Gatekeeper TA verifies credentials from the TEE
- `/data/system_de/<uid>/spblob/<handle>.pwd`
 - scrypt parameters
 - *password handle*, i.e. `HMAC(SHA512("user-gk-authentication" || scrypt(credentials, params))`
- If successful, Gatekeeper returns an *authentication token*
 - Signed token to be used to prove successful authentication
 - Needed by Keymaster to use authentication-bound keys
 - Standard format, designed not to allow replay attacks³
- Gatekeeper implements throttling to prevent bruteforcing

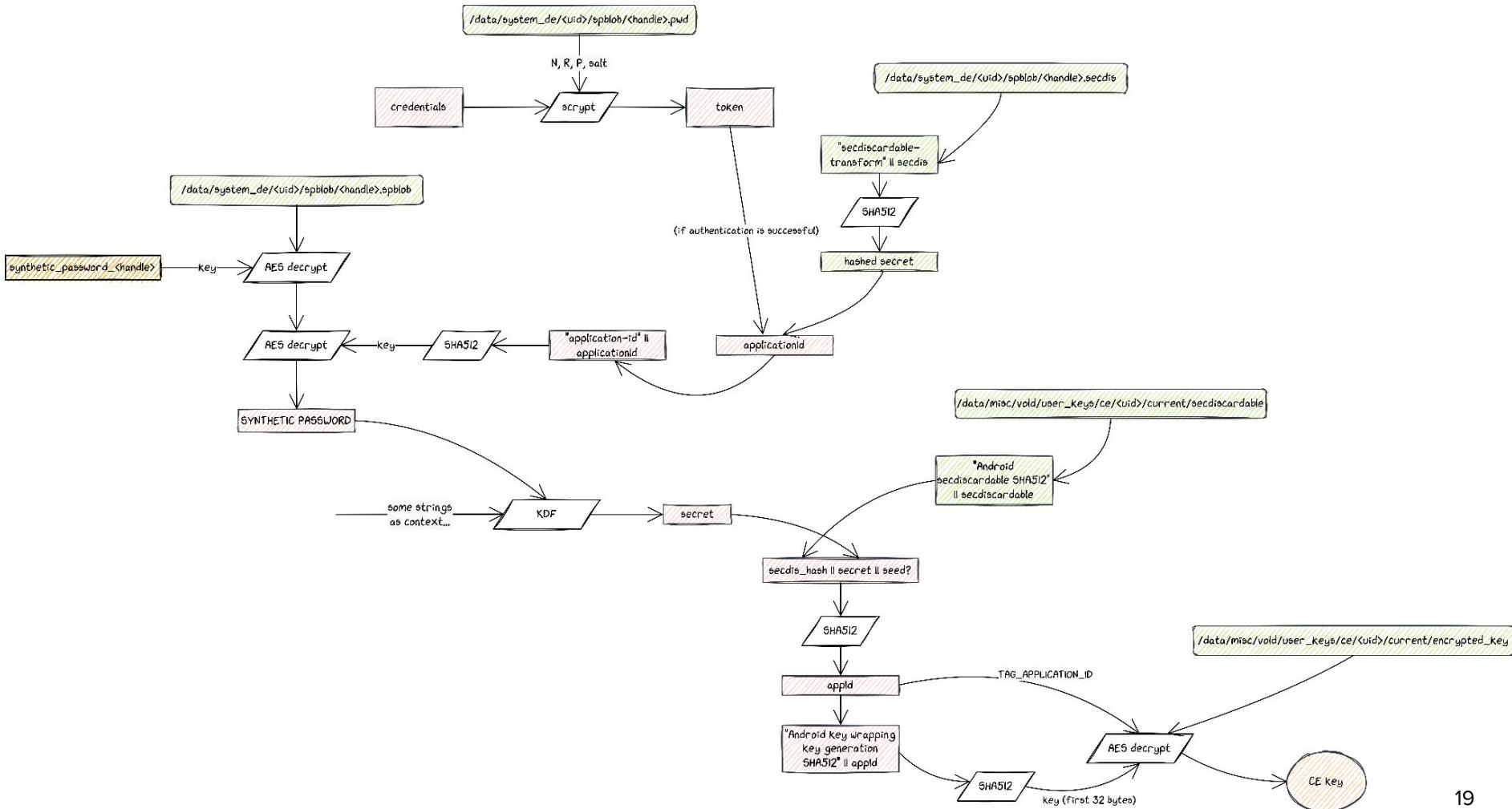
[3]: https://android.googlesource.com/platform/hardware/libhardware/+/master/include/hardware/hw_auth_token.h

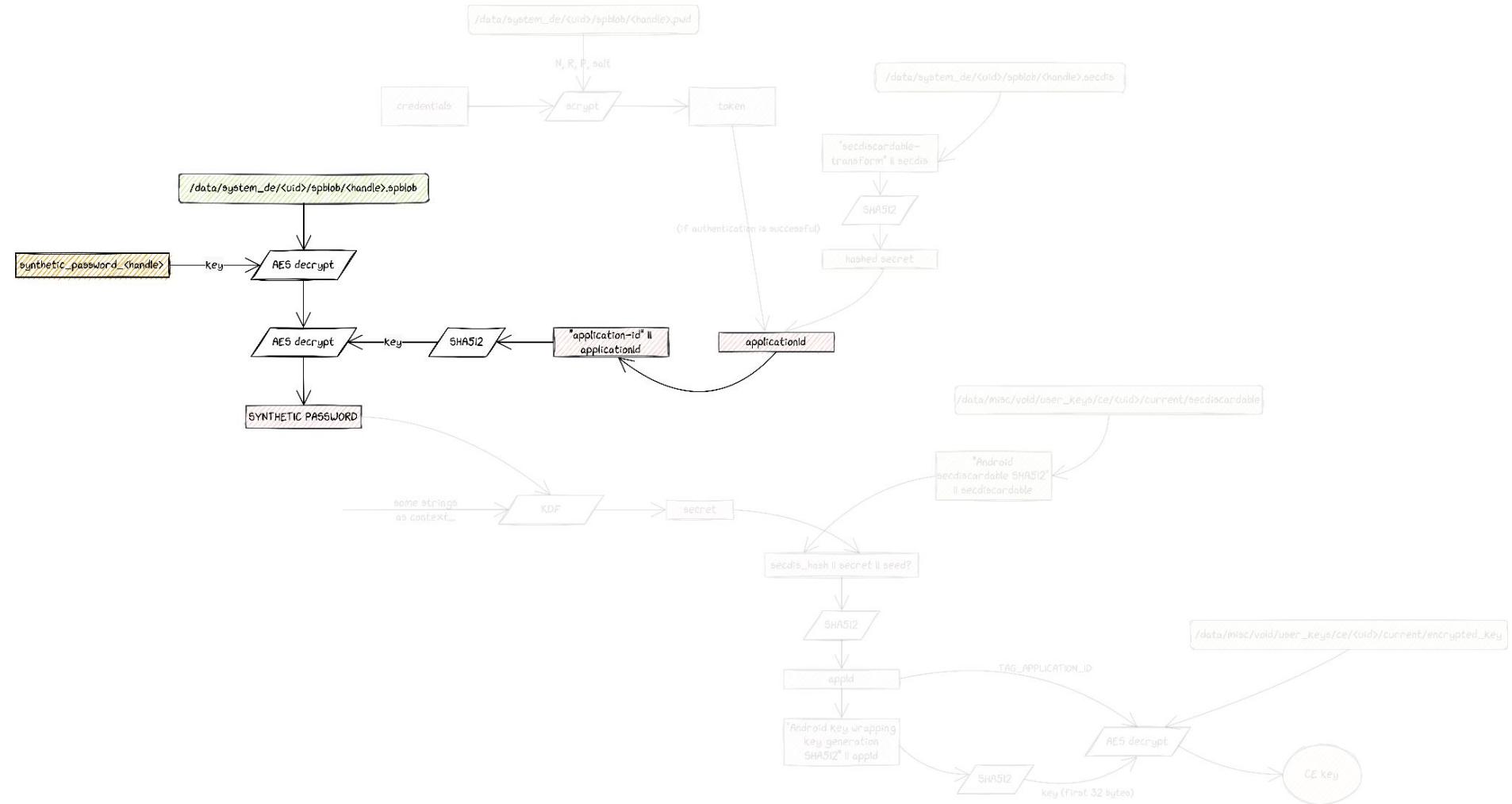
Successful authentication



Failed authentication

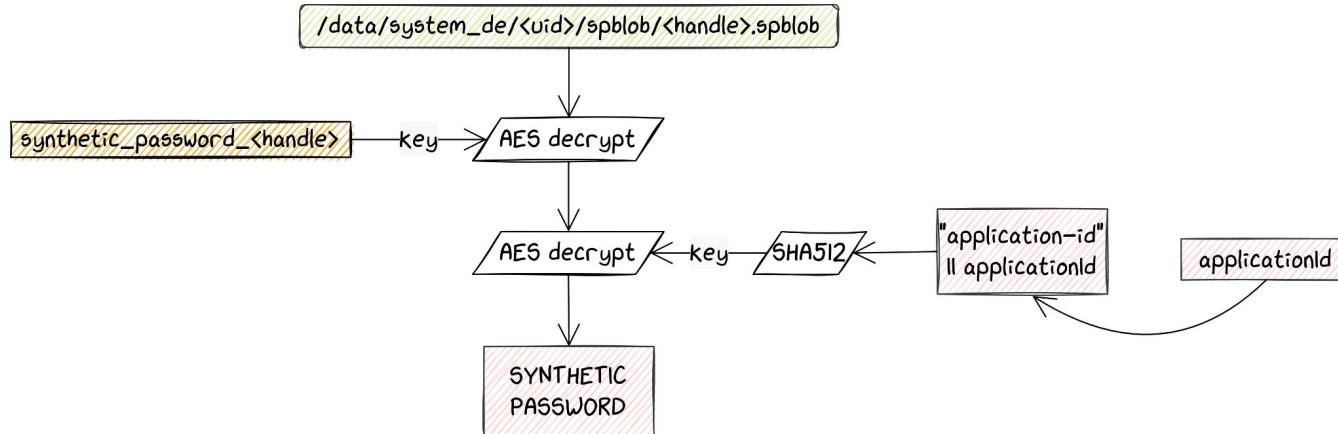






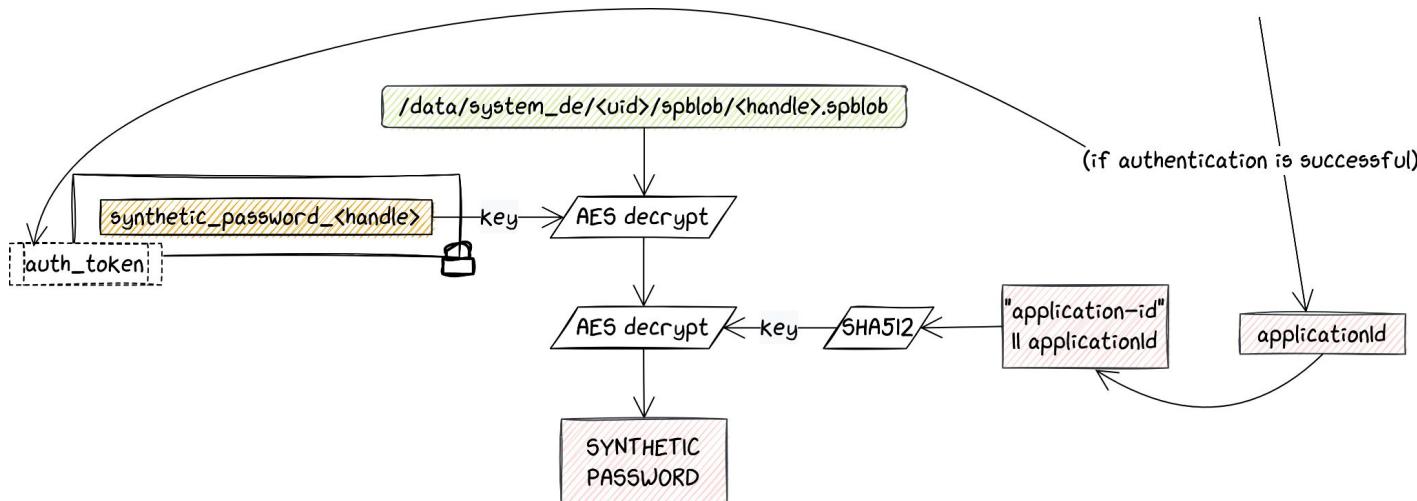
Synthetic Password

- Problem: credentials shouldn't be linked to the CE key
 - What if the user changes them?
- Solution: Synthetic Password
 - Key blob stored in /data/system_de/<uid>/spblob/<handle>.spblob
 - First, decrypted with an authentication-bound, TEE-protected key
 - Then, decrypted with the (hashed) applicationId



Attacking SP derivation

- Need to target the TEE
- Two alternatives
 - Keymaster TA (accessing the first AES key)
 - **Gatekeeper TA** (validating credentials and minting auth tokens)



Global strategy

- Our goal
 - Root the device and access all the device encrypted files
 - Patch the Gatekeeper trustlet to accept any credentials
- For that we need
 - Either multiple bugs (code exec, priv esc, etc)
 - Or one critical bug early in the boot process

PoC on Samsung Device

- Samsung A225f and A226b
 - Cheap (~250€)
 - Mediatek SoC MT6769V and MT6833V
 - No security chip
 - Mix of Mediatek and Samsung code
 - Trustzone OS: TEEGRIS
 - Known critical Boot ROM vulnerability



The Boot ROM Known Vulnerability

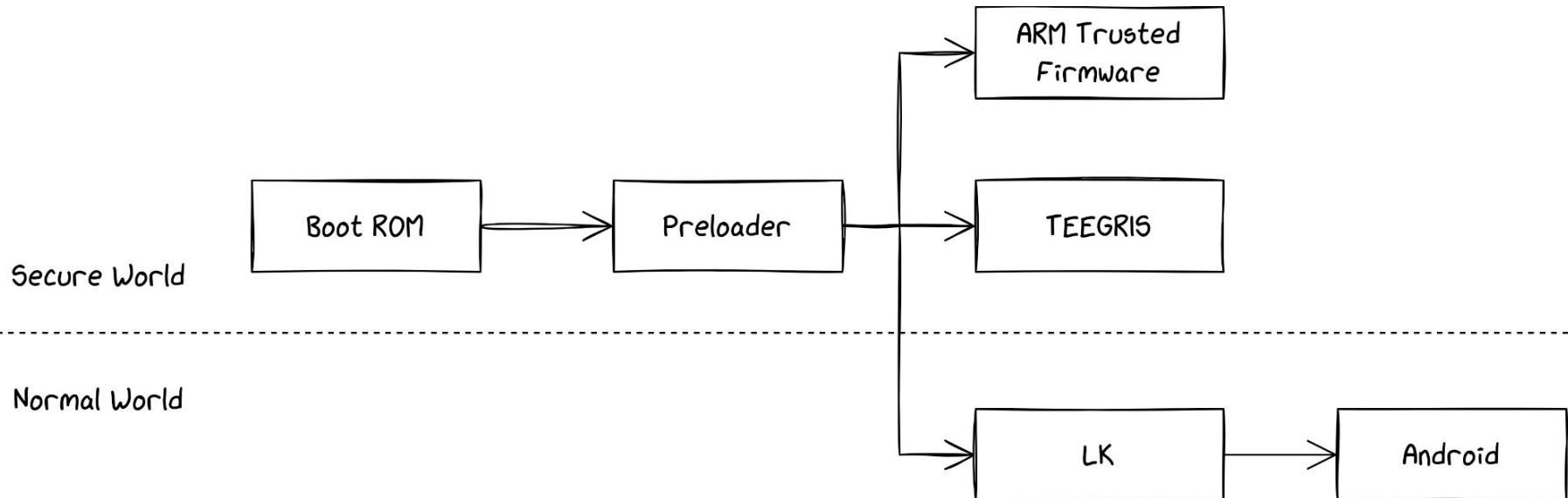
We use the project MTKClient⁴ (by Bjoern Kerler – [@viperbjk](#))

- Exploit boot ROM bugs impacting plenty of Mediatek SoC

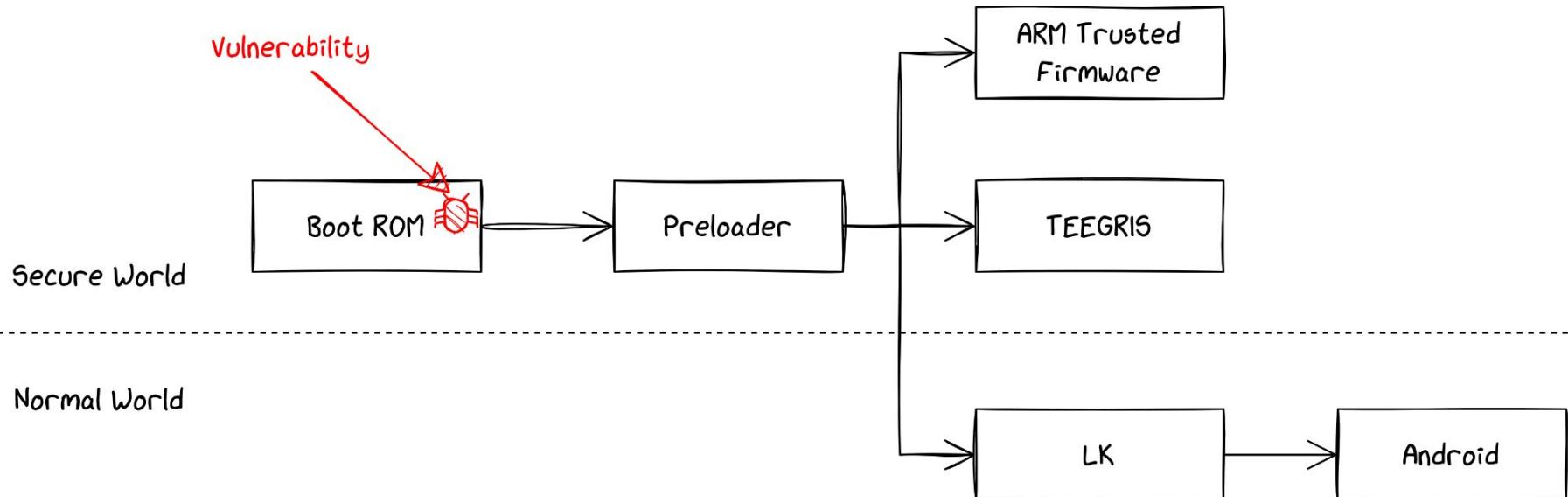
In short, we use it to

- Read/write all the partitions we need to patch
- Boot a patched preloader (BL2) image
- Bypass the secure boot checks done in boot ROM and preloader
- It just works :)

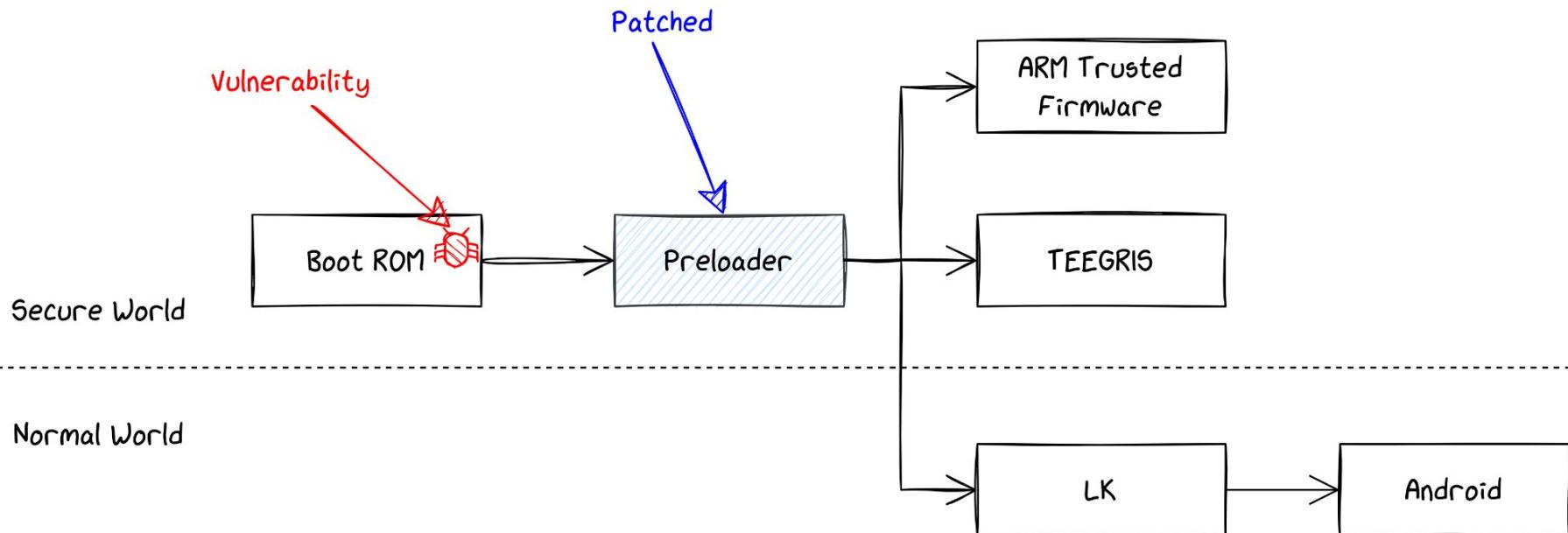
Boot Process



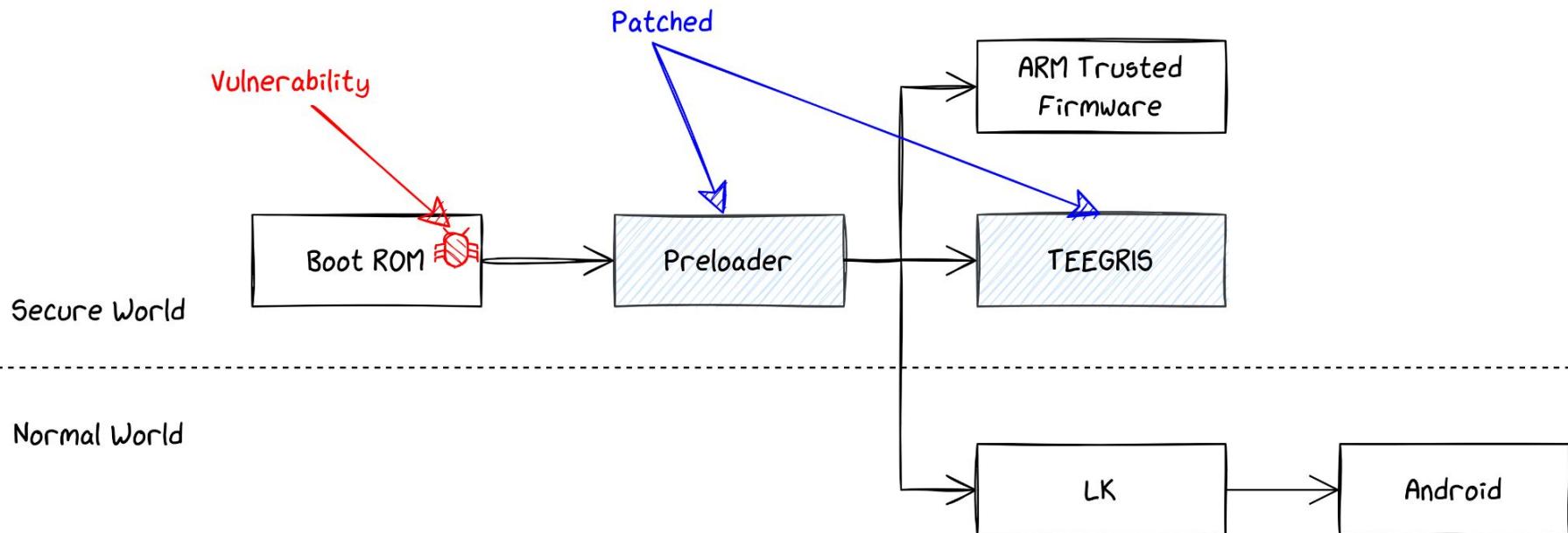
Boot Process



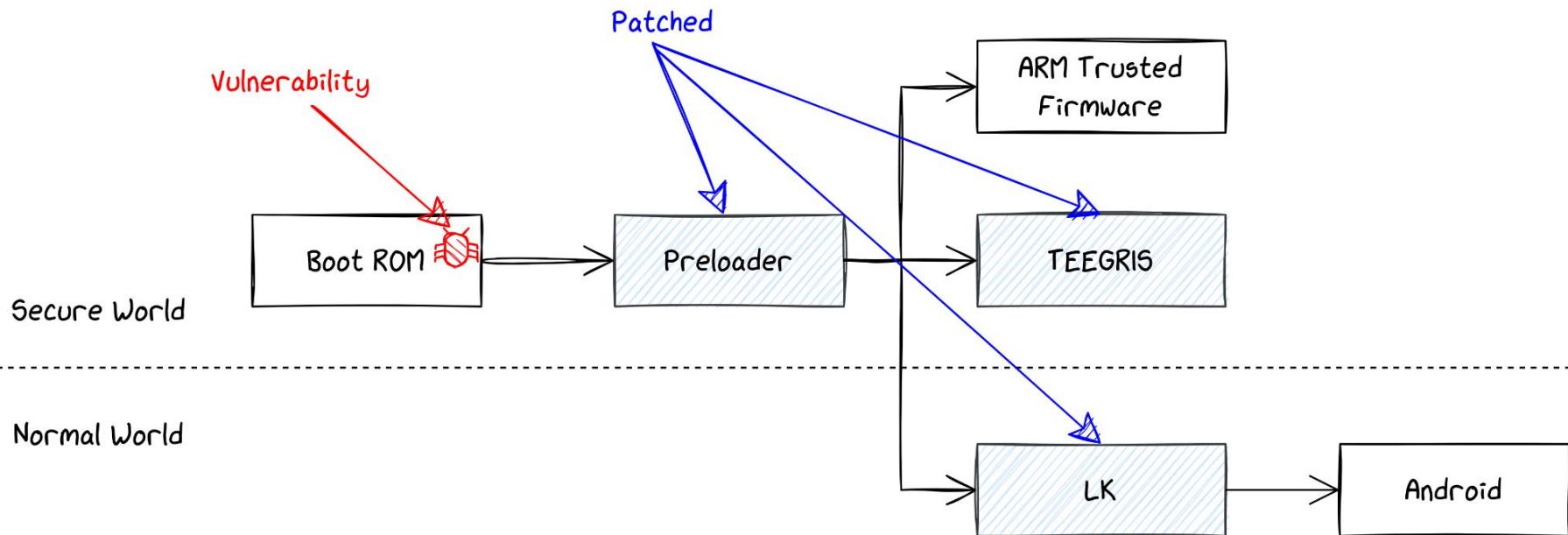
Boot Process



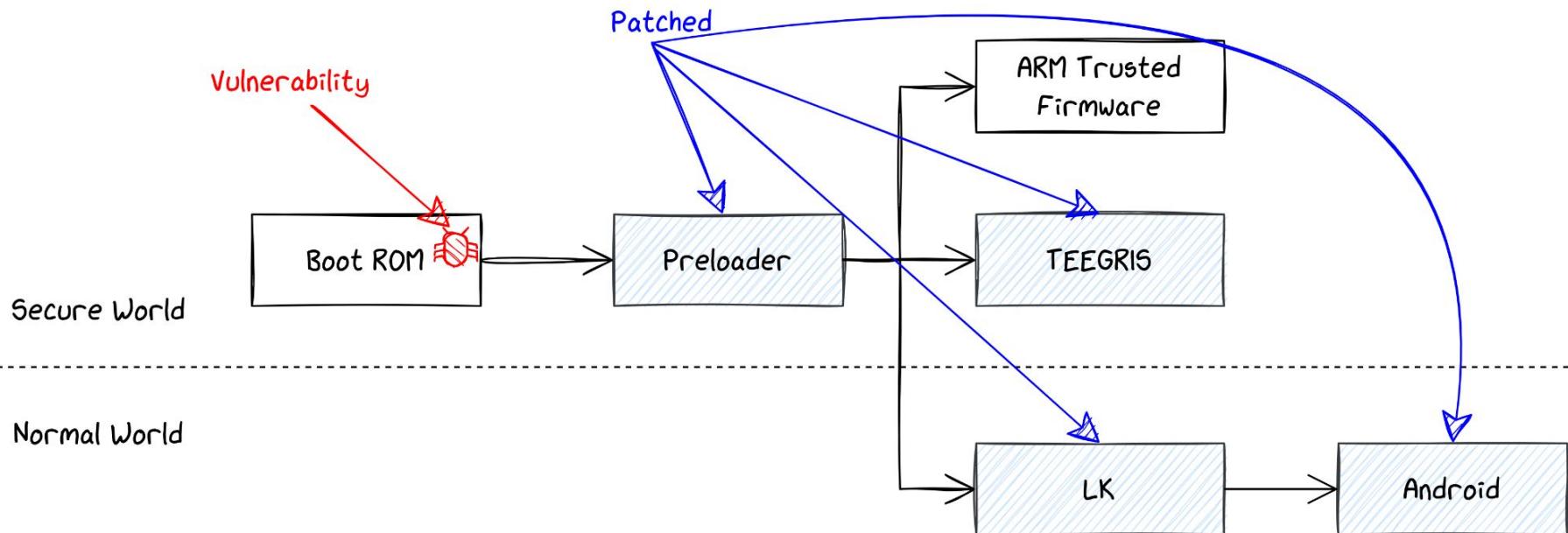
Boot Process



Boot Process



Boot Process



Little Kernel Patching

- Patching strategy: empirical approach
 - Reverse engineering and identify checks
 - Patch, test and repeat
- In the end we patch AVB to launch a modified boot image



Little Kernel Patching

```
26 iVar1 = do_hash(param_1,param_2,DAT_4c6463e0 - param_2,&hash,0x20);
27 if (iVar1 == 0) {
28     iVar2 = memcmp(&STORED_HASH,&hash,0x20);
29     if (iVar2 == 0) {
30         print("[%s][oem] img auth pass\n",&s_SBC_030151a8);
31         goto LAB_02ff82e0;
32     }
33     iVar1 = 0x7021;
34 }
35 print("[%s][oem] img auth fail (0x%x)\n",&s_SBC_030151a8,iVar1);
```

Little Kernel Patching

```
28 iVar1 = do_hash(param_1,param_2,_DAT_4c6463e0 - param_2,&hash,0x20);
29 if (iVar1 == 0) {
30     iVar2 = memcmp(&hash,&hash,0x20);
31     if (iVar2 == 0) {
32         print("[%s][oem] img auth pass\n",&DAT_030151a8);
33         goto LAB_02ff82e0;
34     }
35     iVar1 = 0x7021;
36 }
37 print("[%s][oem] img auth fail (0x%x)\n",&DAT_030151a8,iVar1);
```

Rooting Android

Main partitions used by Android: **boot** and **super**

- Boot contains the kernel and a ramdisk (only used for first boot stage)
- Super is a Dynamic Partition that contains 4 logical partitions
 - system, vendor, product, odm

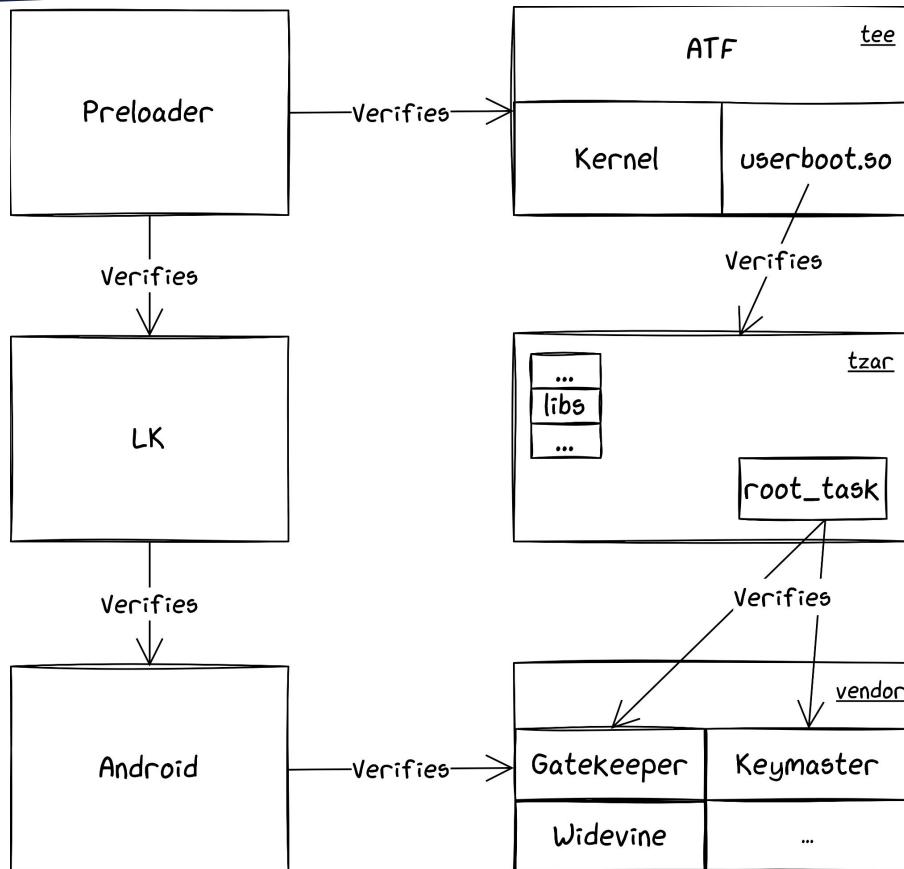
To root it

- Magisk⁵ to patch the boot image
- We made few modifications to su
- Plus other little tricks to patch the super partition

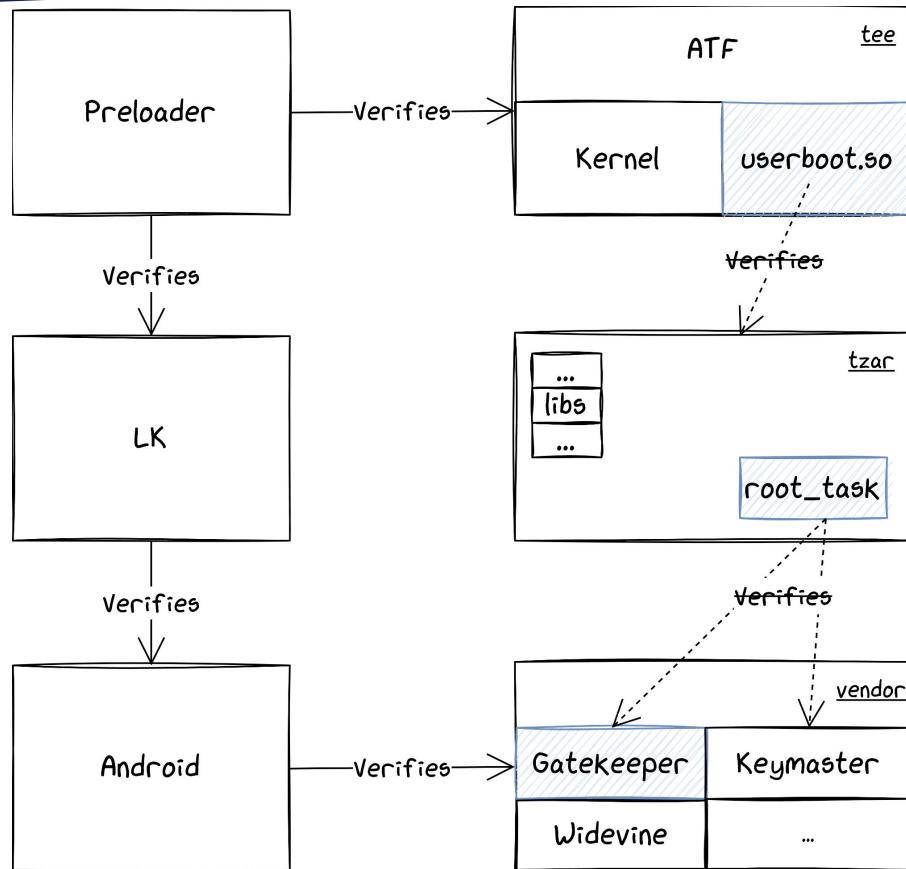
[5]: <https://github.com/topjohnwu/Magisk>

- Trustzone OS designed by Samsung
- For Mediatek and Exynos SoCs
- ROM images:
 - tee1.img: ATF, TEEGRIS kernel, userboot.so
 - tzar.img: TEE root filesystem
 - super.img: Android system, Trusted Applications and Drivers
- Excellent references online⁶

TEEGRIS Images Verification



Patching TEEGRIS



Reversing Gatekeeper

- TAs come in a slightly modified ELF format
 - 8-bytes header and footer with signature
 - Removing them allows to load a nice ELF in your favourite disassembler
- GlobalPlatform API
 - Standard API for TEEs (memory allocation, crypto operations, etc.)
 - Makes reversing easier
- Trusty reference implementation⁷
 - Suggests what to expect from a TA

Gatekeeper Reference Implementation

- 2 Gatekeeper commands: Enroll and Verify
- Verify does two things:
 - $\text{HMAC}(\text{pwd_handle}) == \text{expected}$?
 - If so, create new authentication token
- What if we can leak the key used by HMAC?
 1. $\text{pwd} = \text{generate new password}$
 2. $\text{Value} = \text{HMAC}(\text{pwd_handle})$
 3. $\text{Value} == \text{expected}$

Reversing & patching Gatekeeper

- 2 Gatekeeper commands: Enroll and Verify
- Verify does two things:
 - HMAC(pwd_handle) == expected?
 - If so, create new authentication token



Reversing & patching Gatekeeper

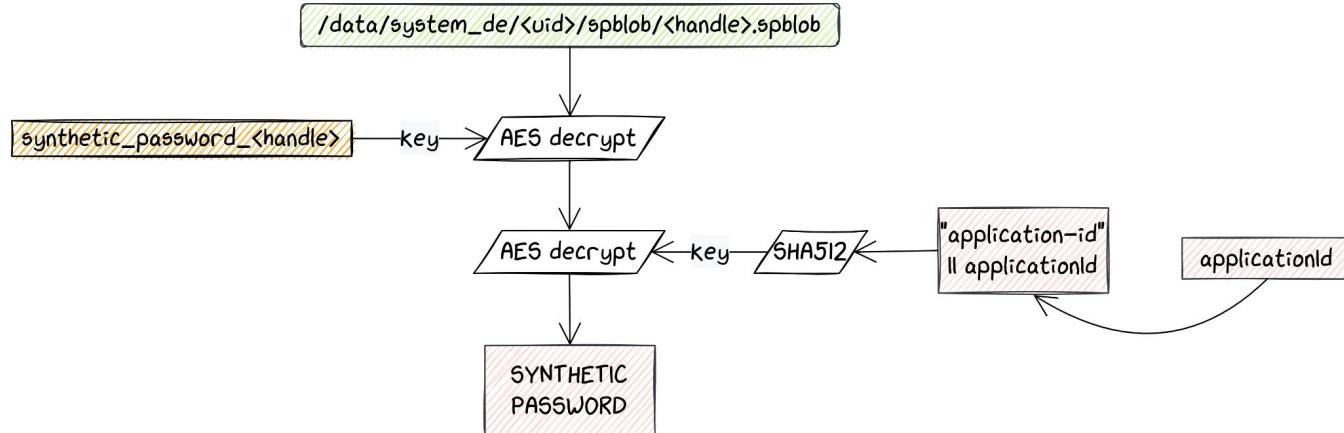
- This Gatekeeper implementation uses a KDF instead of a plain HMAC
 - KDF implemented in a library
 - which calls /dev/crypto
 - many steps to leak the key
- Simpler strategy: patch to accept any credentials
- Always return valid auth token to continue the process
 1. ~~KDF(pwd_handle) == expected?~~
 2. ~~If so,~~ create new auth_token

Reversing Gatekeeper

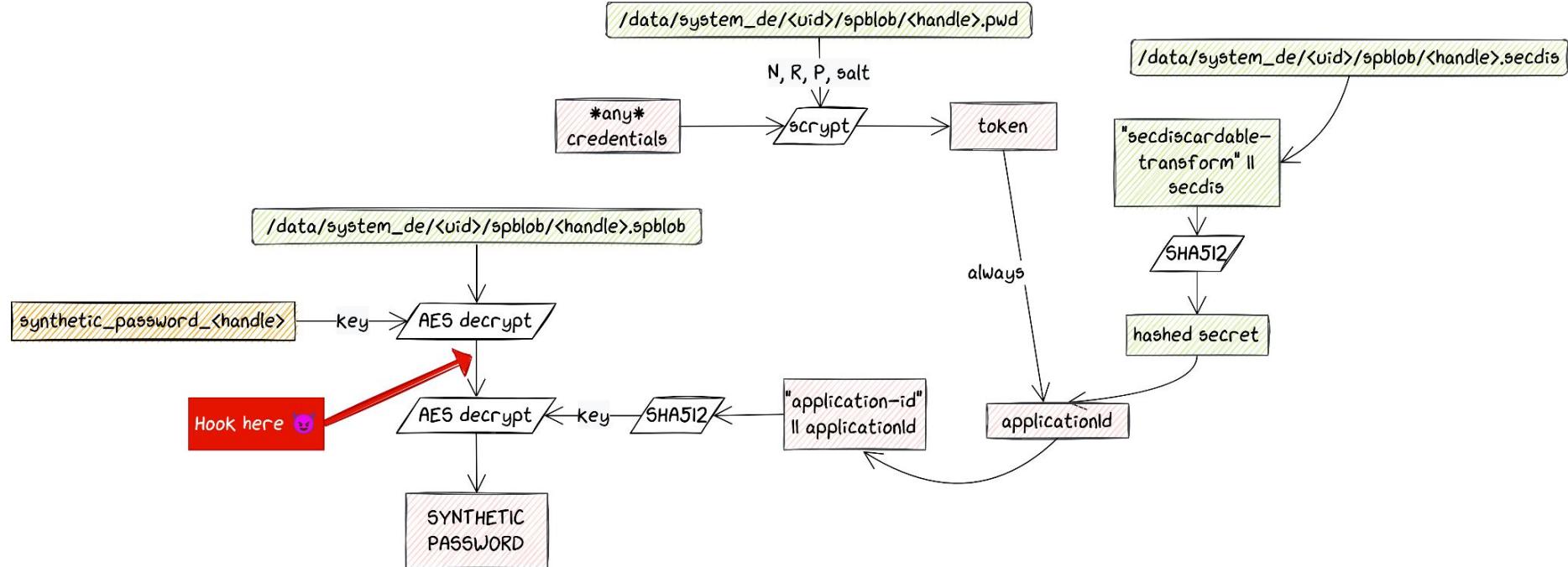
```
22 iVar1 = TEE_AllocateOperation(&local_30,0x50000004,5,0);
23 if (iVar1 == 0) {
24     iVar1 = TEE_DigestDoFinal(local_30,param_1,param_2,auStack_28,&local_38);
25     TEE_FreeOperation(local_30);
26     if (iVar1 == 0) {
27         uVar2 = TEE_AllocateTransientObject(0xa0000000,param_4 << 3,&local_30);
28         if (uVar2 == 0) {
29             uVar2 = TEES_DeriveKeyKDF(auStack_28,local_38,local_48,8,param_4,local_30);
30             if (uVar2 == 0) {
31                 uVar3 = 1;
32                 uVar2 = TEE_GetObjectBufferAttribute(local_30,0xc0000000,param_5,&iStack_34);
33                 if (uVar2 != 0) {
34                     uVar3 = 0;
35                     printf("gatekeeper [ERR] (%s:%u) failed to get object attribute: %x",
36                           "hal_pwd_hmac",
37                           0x12a,(ulong)uVar2);
38                 }
39             }
40         }
41     }
42 }
```

Attack strategy

- Read the output of the first AES decrypt
- Bruteforce credentials to generate applicationId
- Thanks to GCM mode, AES decrypt complains if the key is wrong



Hooking system_server

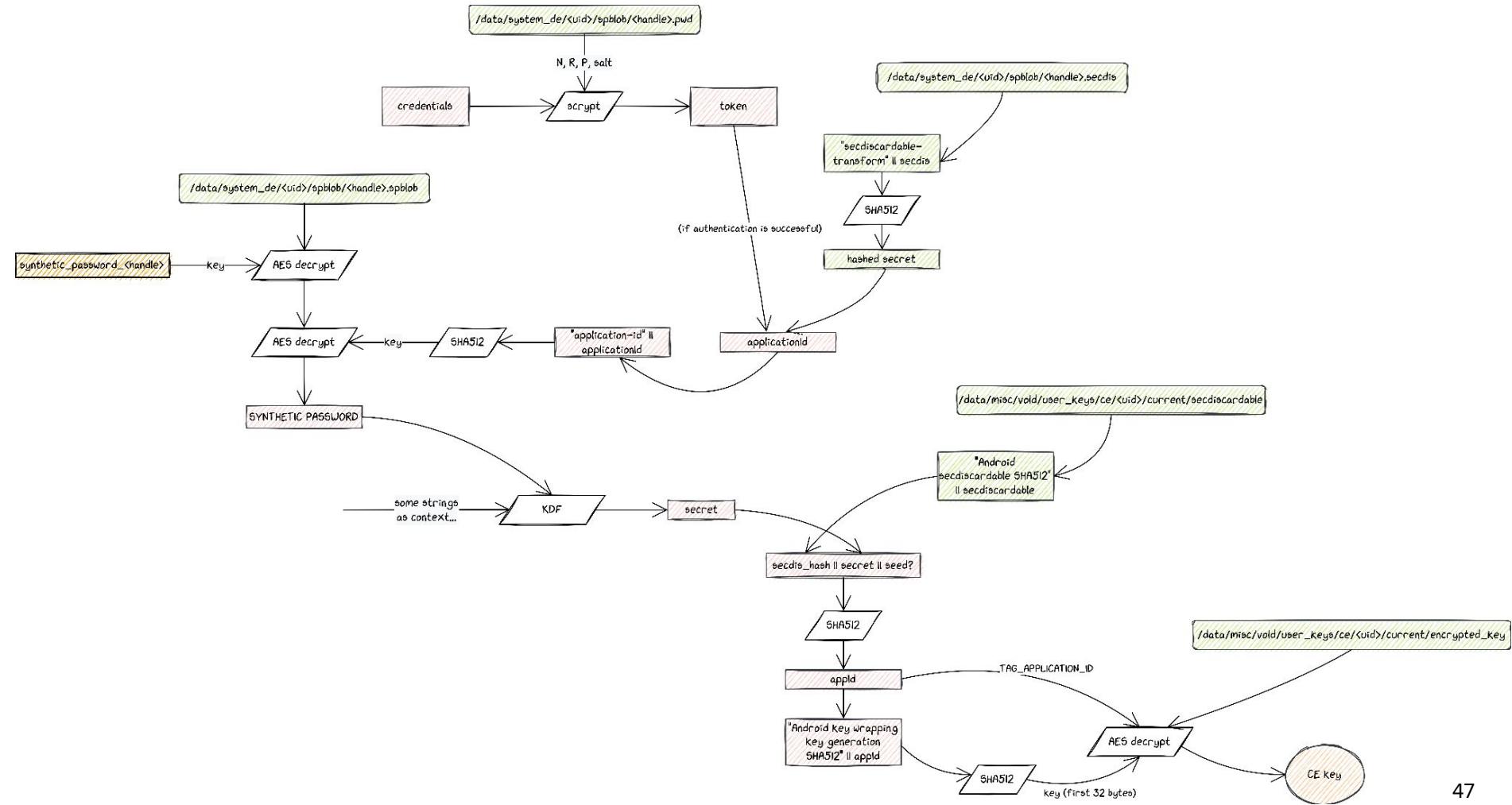


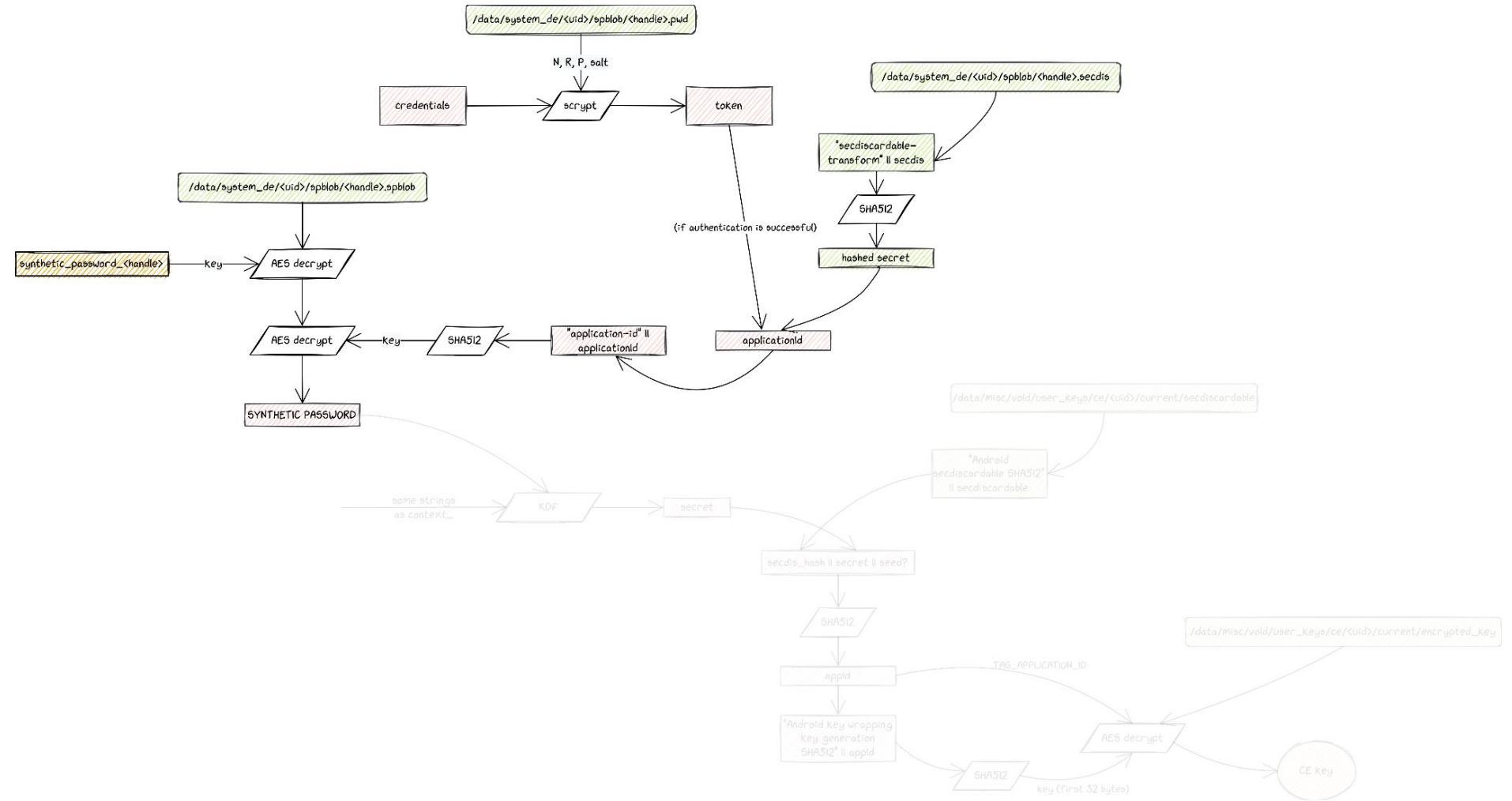
Retrieving intermediate key with Frida

- Use Frida to hook system_server
- Retrieve intermediate buffer decrypted by TEE
 - Possible thanks to the auth token

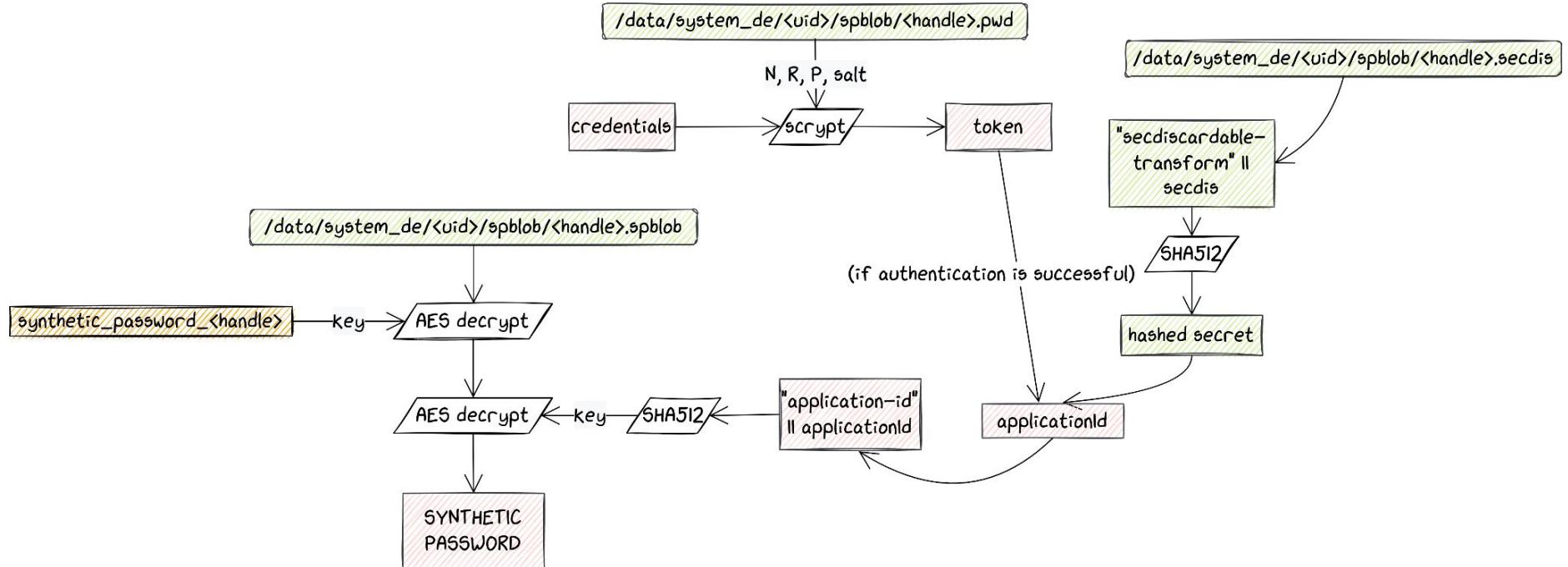
```
$ frida -U system_server -l system_server.js
    / _ _ |  Frida 16.0.19 - A world-class dynamic instrumentation toolkit
    | ( _| |
    > _ _| Commands:
    /_/_|_ help      -> Displays the help system
    . . . object?   -> Display information about 'object'
    . . . exit/quit -> Exit
    . . . .
    . . . More info at https://frida.re/docs/home/
    . . . .
    . . . Connected to SM A226BR (id=R9WTA0BYDPL)

[SM A226BR::system_server ]-> SyntheticPasswordCrypto.decrypt called!
ciphertext = 641a3ed0a68abdae99976b5aff32f8d5aa4d18127272af6ff638c1e88d57cd157fd6f75b4688465f
470bd4cc81081215e9f2085e4b8ea22e0e8f0ed32a381f641d5cd071d2e177c4a8a1b6e6824f52f251366ff730f66
b7cfdf72f11f9761efc5e0cf68bd7bdec00456e07dfb9f1a7f720e97aa262c0507bc87ef46e603a265c821cb1a1dc5
c6f6be6fd43ac3431d0d013de8c9
[SM A226BR::system_server ]->
[SM A226BR::system_server ]-> |
```

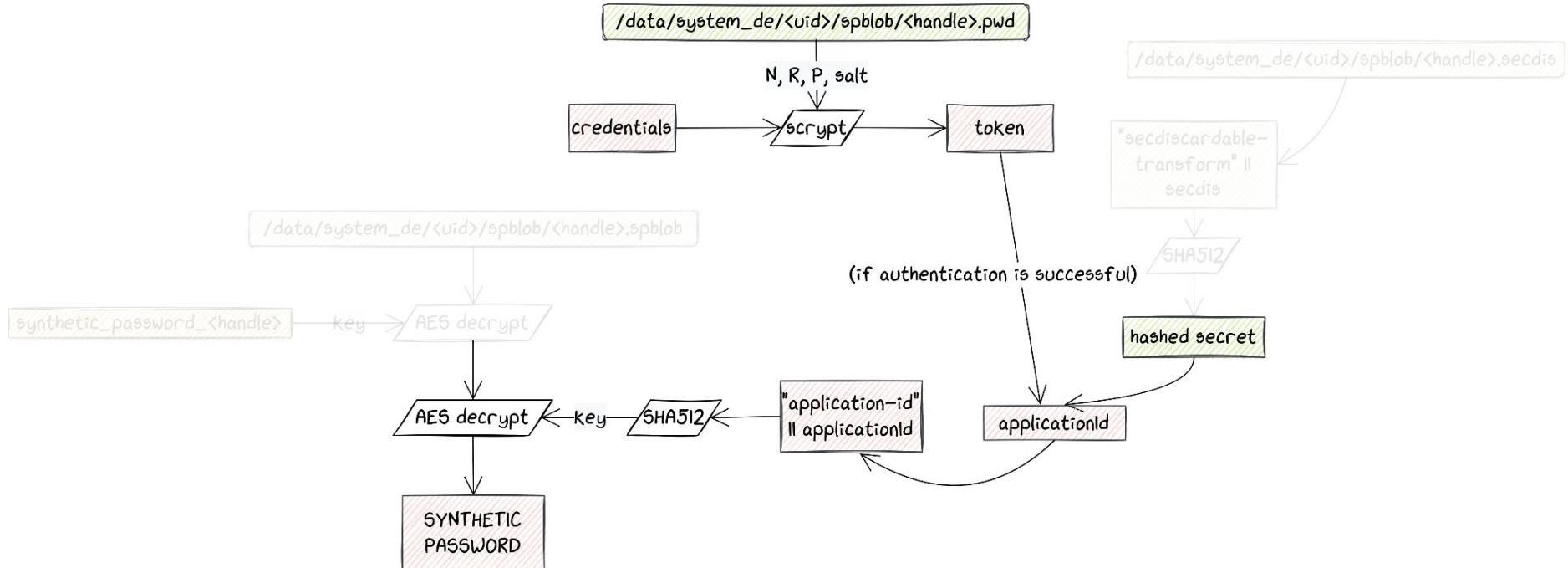




Bruteforce of the password



Bruteforce of the password



Bruteforce of the password

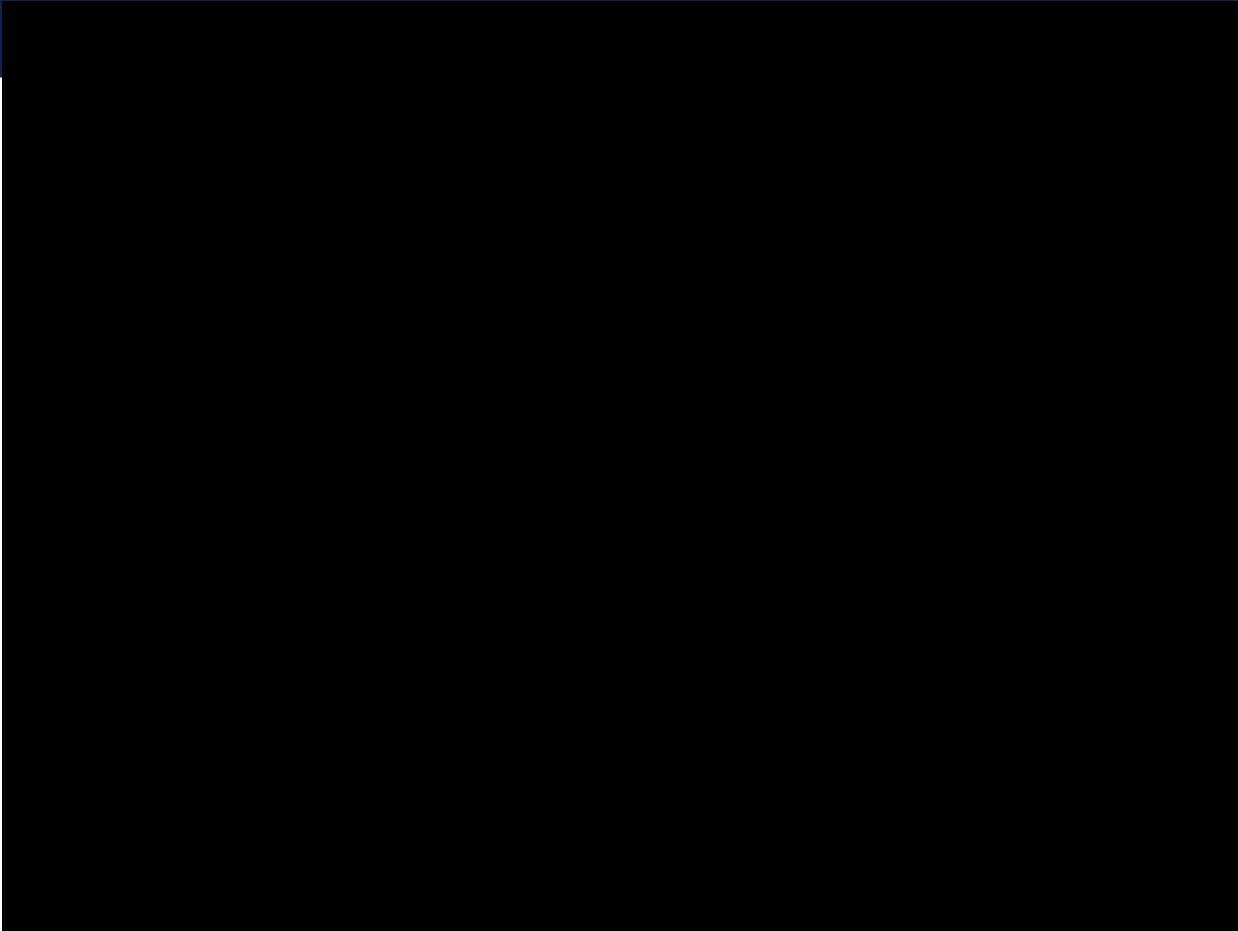
1. pwd = generate new password
2. token = **scrypt**(pwd, R, N, P, Salt)
3. Application_id = token || Prehashed value
4. Key = **SHA512**("application_id" || application_id)
5. **AES_Decrypt**(value_from_keymaster, key)

Bruteforce of the password

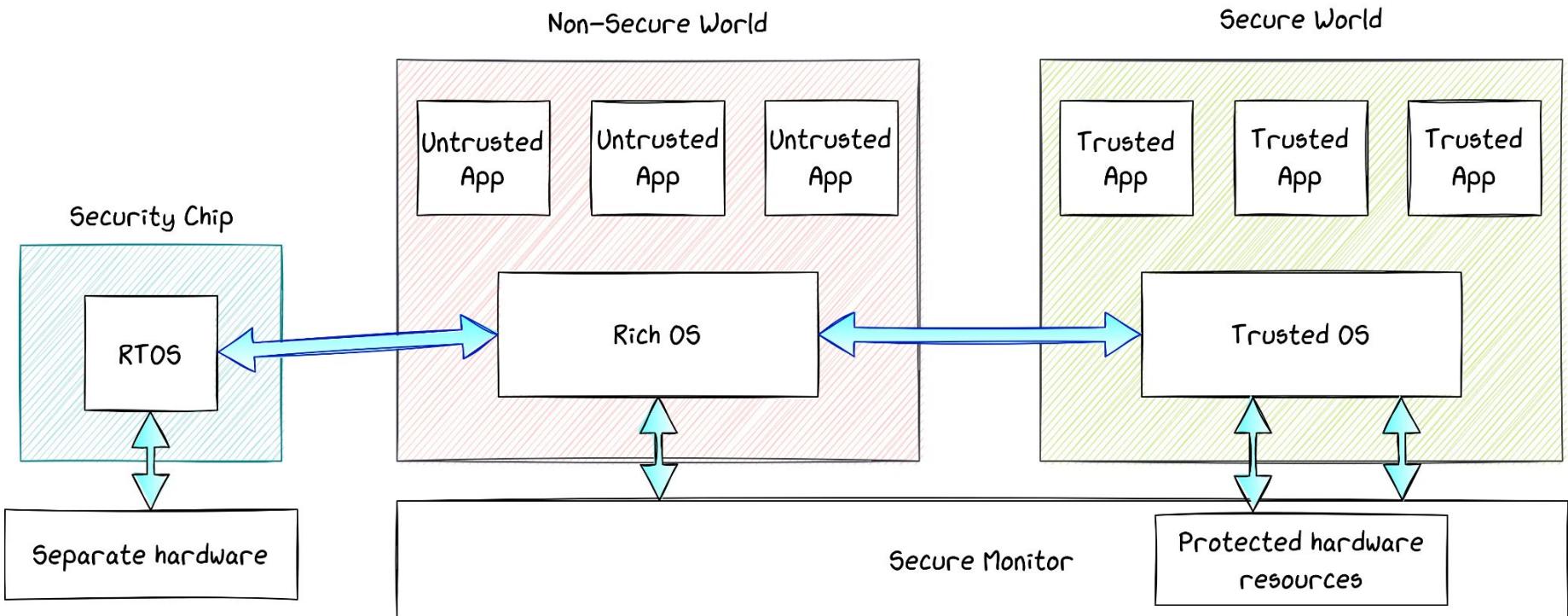
1. `pwd = generate new password`
2. `token = scrypt(pwd, R, N, P, Salt)`
3. `Application_id = token || Prehashed value`
4. `Key = SHA512("application_id" || application_id)`
5. `AES_Decrypt(value_from_keymaster, key)`

```
$ python3 bruteforce-tee.py
workers will cycle through the last 5 chars
Found it: 1234
the plaintext is '1234'
Done in 18.031058311462402s
Throughput: 1478.448992816657 tries/s
```

Demo 1



Architecture w/ Trusted Chip



The Titan M Chip

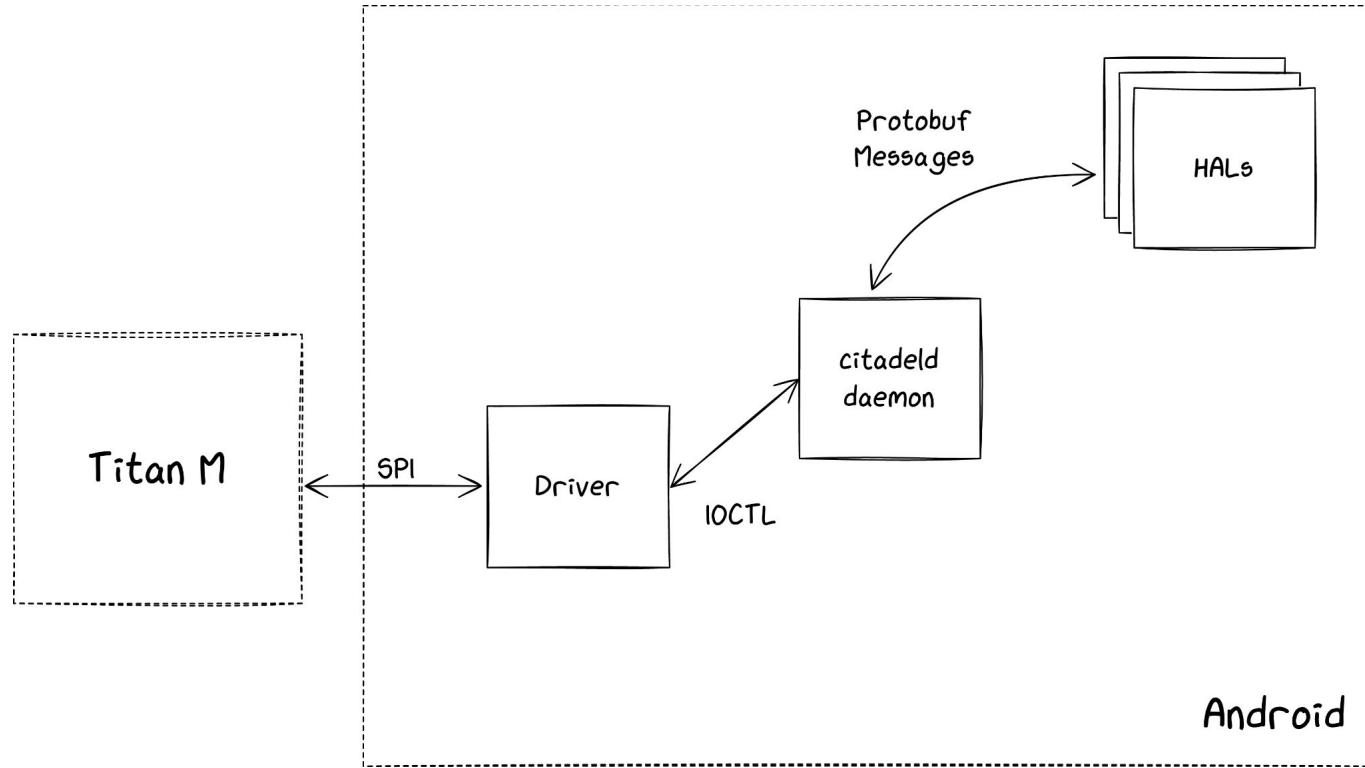
- Security chip made by Google for Pixel phones
- From Pixel 3 to Pixel 5a
 - In this PoC we use a Pixel 3a
 - Titan M2 introduced from Pixel 6
- Based on Arm Cortex-M3
- Most of the code is divided into tasks
 - Keymaster (Strongbox), **Weaver**, AVB, etc
- Separate memory and resources
 - Communicates with Android on SPI bus



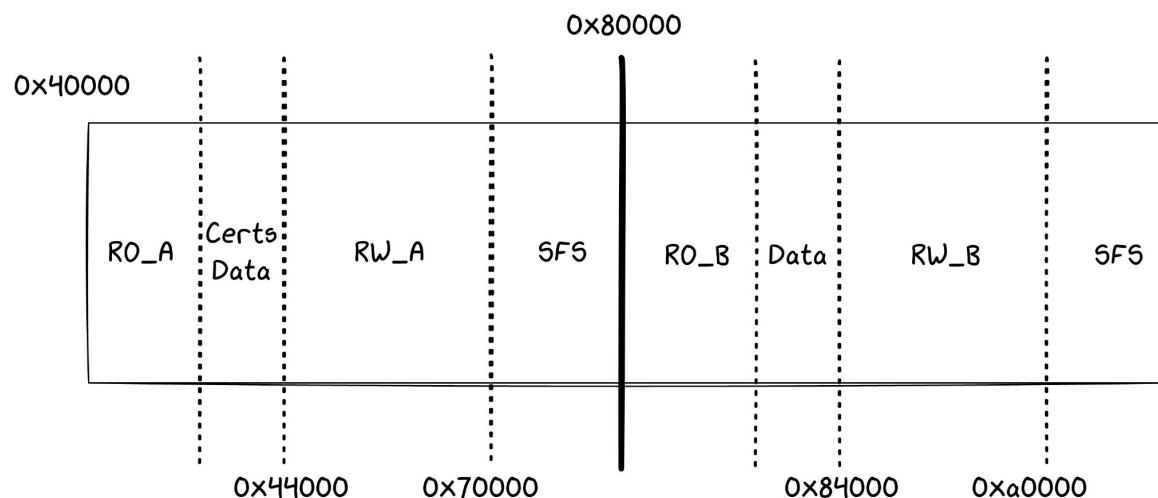
Trusted chip vs TrustZone

- In TrustZone, secure and normal world run on the same CPU
 - Shared hardware (cache, RAM)
 - Side-channel attacks are possible (e.g. Rowhammer)
- Titan M relies on tamper-resistant hardware
- Separate firmware
 - Limited in size
 - Conceptually simple
 - Isolated from the rest of the system

Communication with the chip



Memory Layout



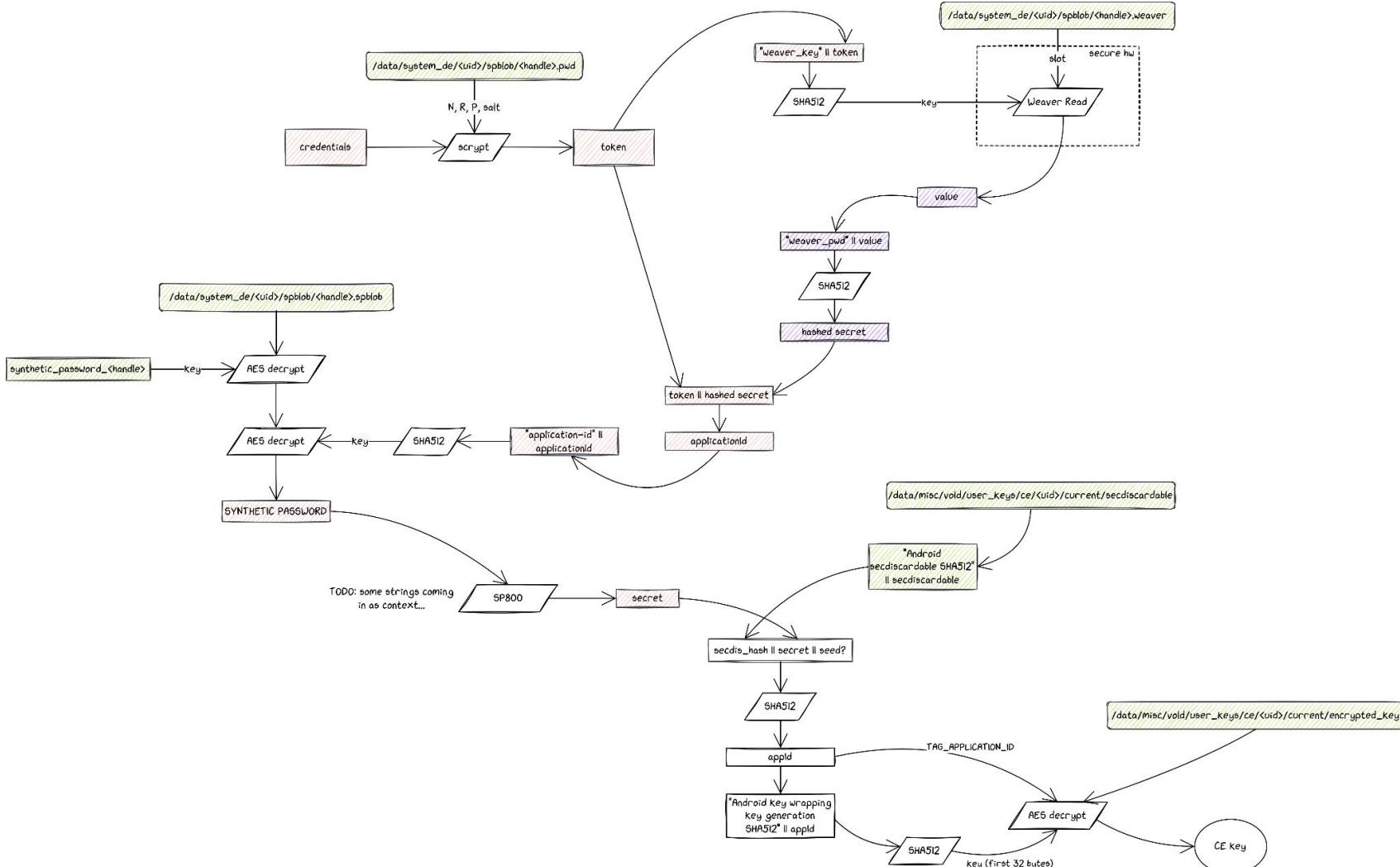
Weaver

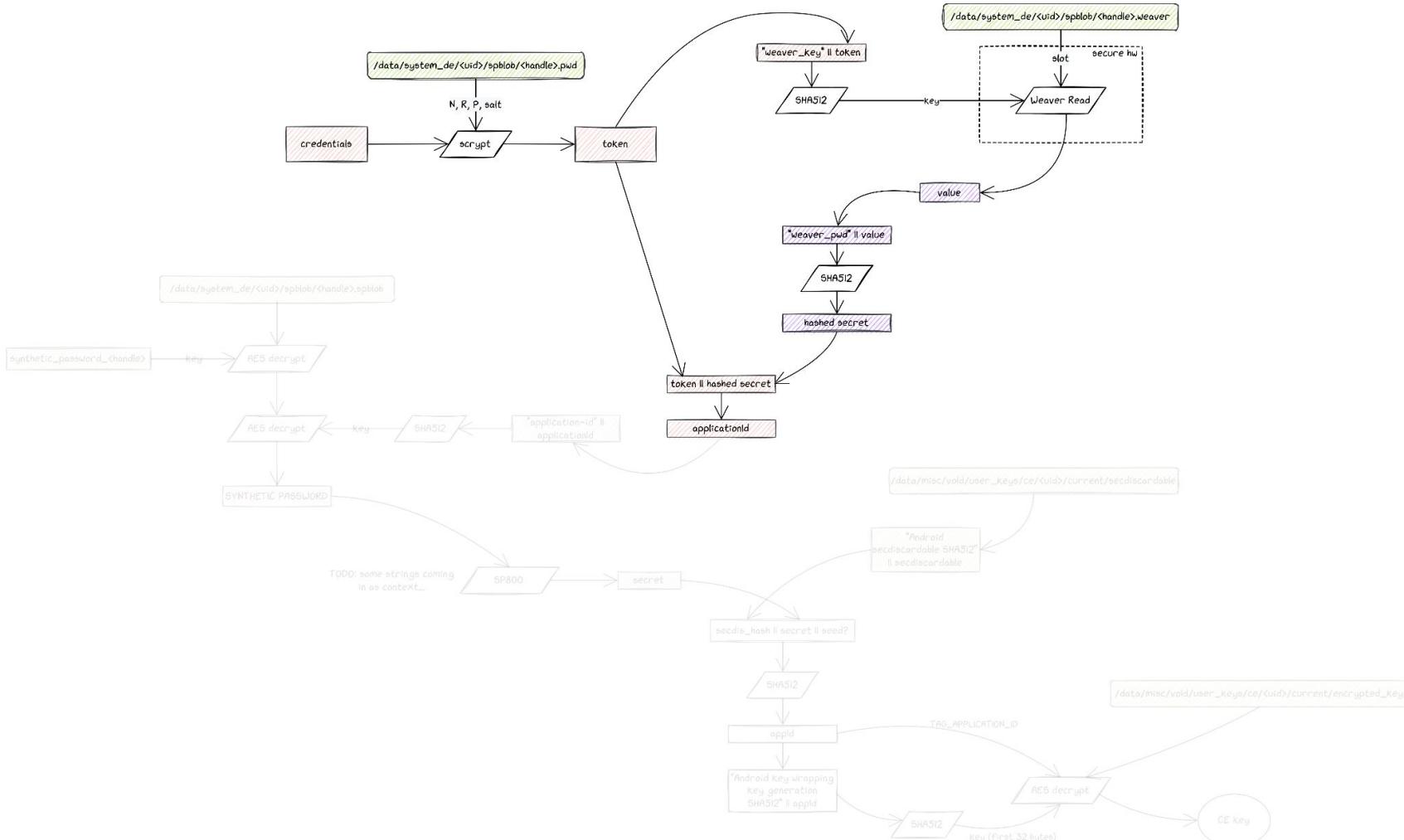
- Key/Value storage
 - Stored in slots
 - In two different places in the flash memory
- 4 commands: GetConfig, Read, Write, Erase
- Implements throttling as well

```
// Read

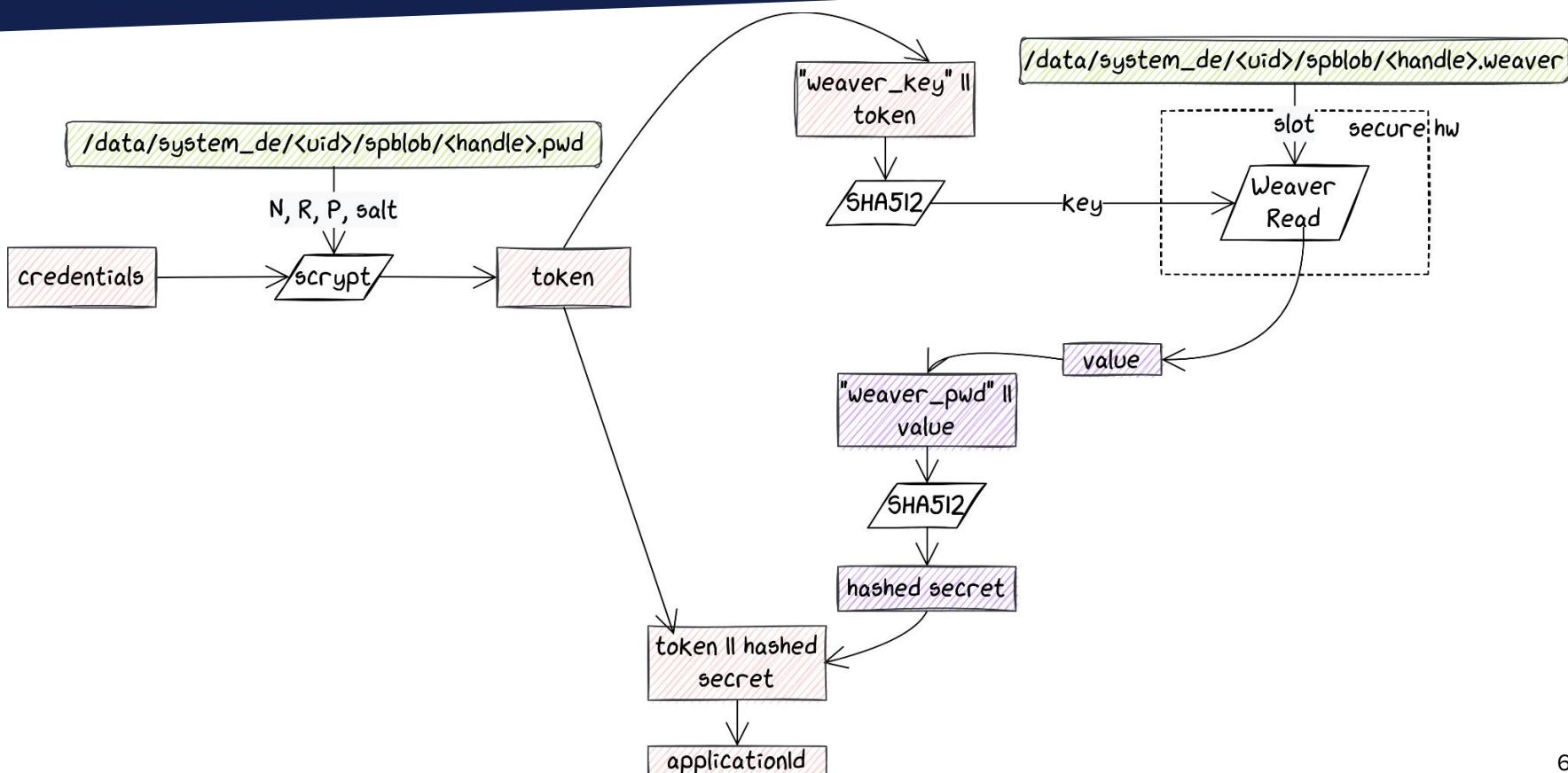
message ReadRequest {
    uint32 slot = 1;
    bytes key = 2;
}

message ReadResponse {
    Error error = 1;
    uint32 throttle_msec = 2;
    bytes value = 3;
}
```





CE key derivation with Weaver



PoC on Google Pixel

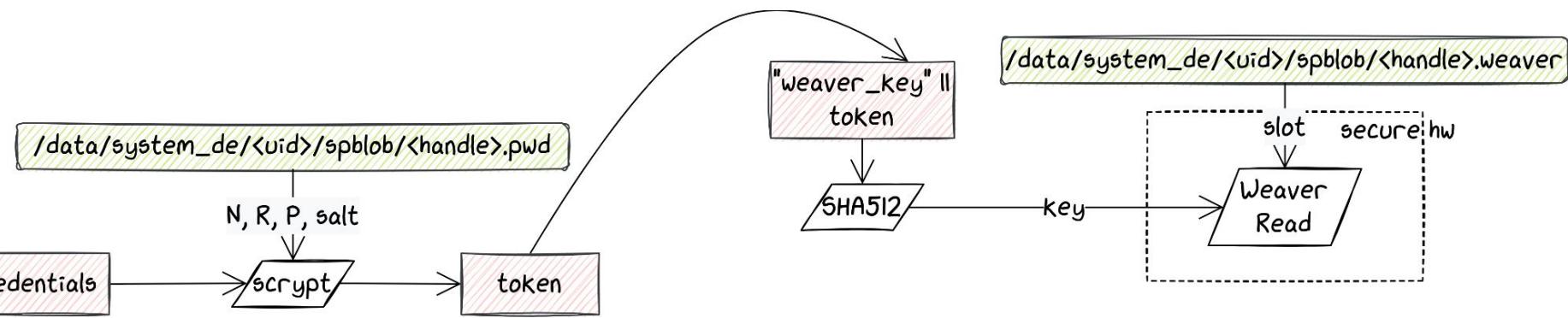
- We consider the device being already rooted
- Weaver relies on the security chip Titan M
- Here we exploit CVE-2022-20233 to execute code on the chip
- Out-of-bounds write of 1 byte to 0x1
 - Can be repeated multiple times
 - Huge constraints on the offset
 - We managed to overwrite a global field and cause another corruption
- Full exploit write-up in our blog⁸

[8]: <https://blog.quarkslab.com/attacking-titan-m-with-only-one-byte.html>

Nosclient and the leak command

- We built a client to communicate with Titan M, nosclient
- “Leak” feature:
 - `./nosclient leak <address> <size>`
 - Read `<size>` bytes from `<address>`
 - Arbitrary read in Titan M’s memory
- Weaver slots and values are stored in flash
 - Reverse engineering to understand a memory range
 - Then search for 16 bytes digests
 - Weaver Write and Read help out

Our Strategy



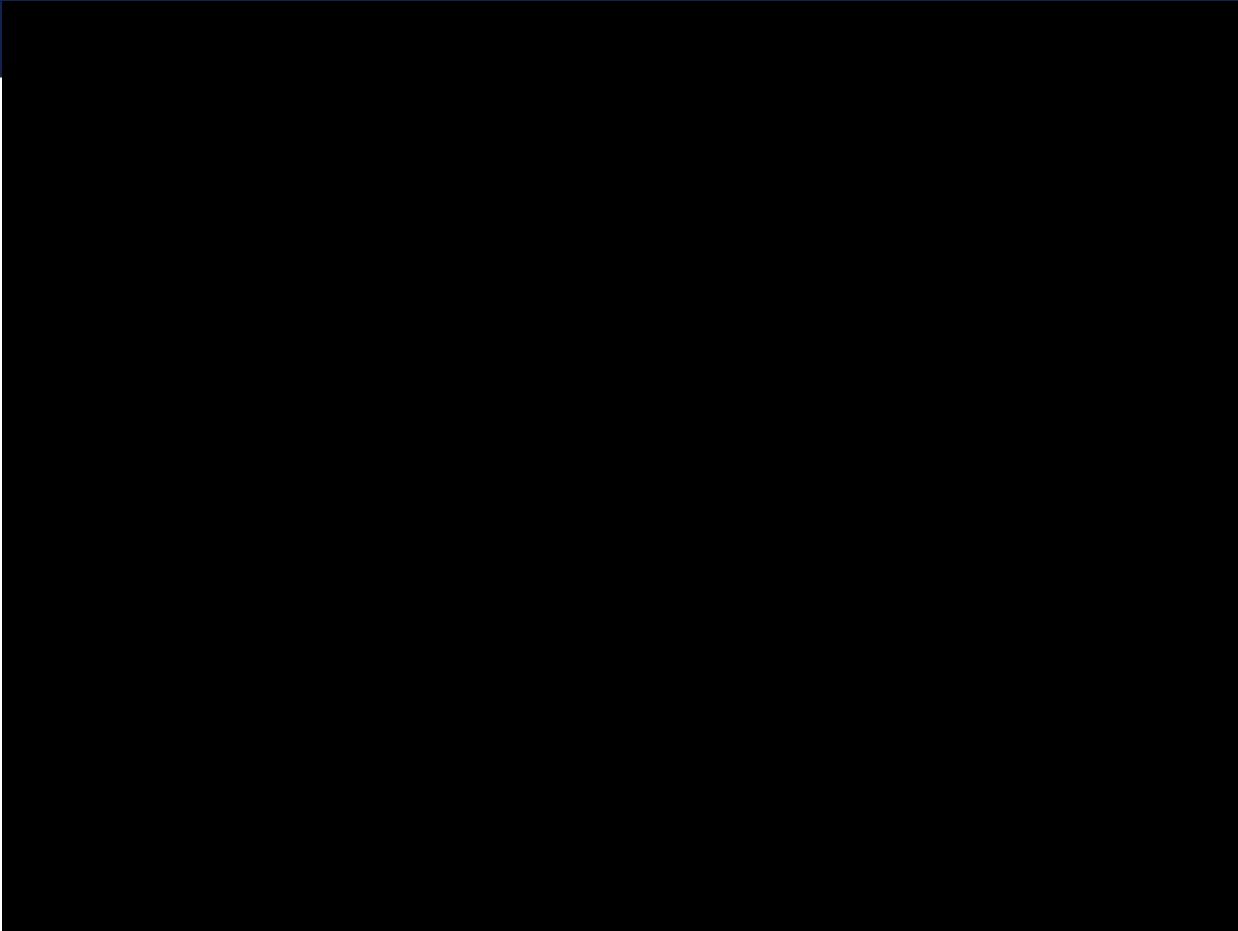
1. Leak the Weaver key
2. Use it to compare our generated credentials

Bruteforce of the password

1. `pwd = generate new password`
2. `token = scrypt(pwd, R, N, P, Salt)`
3. `key = SHA512("weaver_key" || token)`
4. Compare with leaked Weaver key

```
$ python3 bruteforce.py
workers will cycle through the last 5 chars
Found it: 1106
the plaintext is'1106'
Done in 15.063793659210205 s
Throughput: 1491.722504195411 tries/s
```

Demo 2



Conclusion

- FBE is very well designed
- Ingredients from “everywhere” are used to derive the key
 - Files owned by privileged users
 - TEE-protected keys
 - Weaver values (when available)
- Multiple bugs needed to break it
 - Or a very powerful one
- You still need to bruteforce credentials in the end
- “my very secret password example for REcon 2023” will be hard to guess :)

Thank you!

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Little Kernel

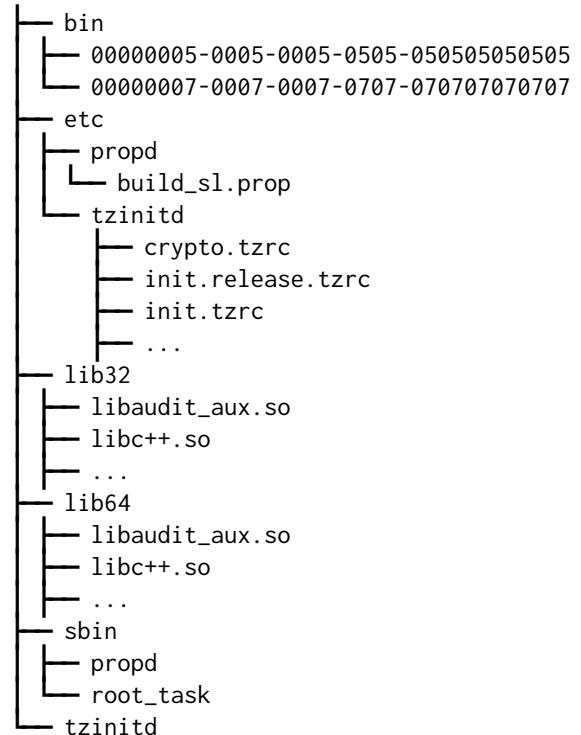
LK: Android bootloader based on Little Kernel

- Allows to boot Android or other modes (Recovery)
- Loads TZAR image in TEEGRIS
- Implements **Android Verified Boot v2**
 - Verification of Android images
 - Involving boot and vbmeta images
 - Anti-rollback

TZAR image

TrustZone ARchive: contains a root filesystem

- Shared libraries
- Binaries
- tzinitd (init binary)
- root_task



Patching TEEGRIS

Our final goal is to run a modified Gatekeeper TA

- Patch userboot.so from the tee1 partition
 - Disable verification of TZAR image
- Patch root_task from the TZAR image
 - Disable verification of TA
- Patch the Gatekeeper TA
 - Accept any credentials and return a valid auth token