

# Banker's Algorithm

①

— Used in Banking to ensure that bank never allocated its available cash in such a way that it could no longer satisfy the needs of all customers.

Data Structure Used:

- Available — indicates available res. of each type
- Max — max demand of each process
- Allocation — allocated to each process
- Need — remaining resources needed by each process

$$\text{Need}_i = \text{Max}_i - \text{Allocation}_i$$

Safety Algorithm:

i) Let  $work = \text{Available}$

$finish[i] = \text{false}$   $i = 0, 1, 2, \dots, n$

2) Find an  $i$  such that

$\# / finish[i] = \text{false}$

$need[i] \leq work$

If no such  $i$  exists go to 4

3.  $work = work + \text{Allocation}_i$

$finish = \text{true}$

go to 2



4. If  $\text{finish}[i] = \text{true}$  for all  $i$  then system is in safe state.

Example problem: 5 processes  $P_0 - P_4$ . Three resource types of resources  $A=10, B=5, C=7$

Process	Allocation A B C	Max A B C	(none) Available <del>13</del>	Need
$P_0$	0 1 0	7 5 3	<del>3 3 2</del>	7 4 3
$P_1$	2 0 0	3 2 2	5 3 2	1 2 2
$P_2$	3 0 2	9 0 2		6 0 0
$P_3$	2 1 1	2 2 2		0 1 1
$P_4$	0 0 2	4 3 3		4 3 1

Available = total - Allocation



# Apply Safety Algorithm

(3)

Initially  $work = Available$

$$work = 3 \quad 3 \quad 2$$

$$finish[i] = false \quad \forall i$$

P<sub>0</sub>

$$finish[0] = false$$

$$need[0] \leq work.$$

$$7 \ 4 \ 3 \leq 3 \ 3 \ 2$$

~~X~~

P<sub>1</sub>

$$finish[0] = false$$

$$need[1] \leq work$$

$$\cancel{7 \ 4 \ 3} \leq \cancel{3 \ 3 \ 2} \quad \checkmark$$

$$work = work + Allocation$$

$$= 3 \ 3 \ 2 + 2 \ 0 \ 0$$

$$= 5 \ 3 \ 2$$

$$finish[0] = True$$

P<sub>2</sub>

$$need[2] \leq work$$

$$6 \ 0 \ 0 \leq 5 \ 3 \ 2$$

X

$$finish[0] = false$$



P<sub>3</sub>  $fun[3] = false$

$head[3] \leq wone$

$011 \leq 532$

$fun[3] = True$

$wone = 532 + 211 = 743$

P<sub>4</sub>

$fun[4] = false$

$head[4] \leq wone$

$431 \leq \underline{743}$

$wone = 743 + 002$

$= 745$

$fun[4] = True$

P<sub>5</sub>

$head[5] \leq wone$

$743 \leq 745$

$wone = wone + 110$

$= 745 + 010$

$= 755$

P<sub>6</sub>

$head[6] \leq wone$

$600 \leq 755$

$wone = wone + Allocation$

$= 755 + 302 = 1057$



## Resource Request Algorithm:

- If a process made an additional request
- Check it & granted by this algorithm

work

Process	Allocation	Max	Available	Need
	A B C	A B C	A B C	A B C
P0	0 1 0	7 5 3	<del>3 3 2</del>	7 4 3
P1	<del>2 0 0</del> 3 0 2	3 2 2	2 3 0	<del>1 0 2</del> 0 2 0
P2	3 0 2	9 0 2		6 0 0
P3	2 1 1	2 2 2		0 1 1
P4	0 0 2	4 3 3		4 3 1

P1 makes an additional request (1, 0, 2)

Check Request  $\leq$  need & Req  $\leq$  Avail

$$P_1 \quad (1 \ 0 \ 2) \leq (1 \ 2 \ 2) \quad \text{Yes}$$

$$(1 \ 0 \ 2) \leq (3 \ 3 \ 2) \quad \text{Yes}$$

So allocate the requested resource.

$$\therefore \text{Avail} = \text{Avail} - \text{Req}$$

$$(\text{---}) = 3 \ 3 \ 2 - 1 \ 0 \ 2$$

$$= 2 \ 3 \ 0$$

update it in table



$$\text{Allocation} = \text{Alloc} + \text{Req}$$

$$(P1) = 200 + 102$$

$$= 302 //$$

$$\text{Need} = \text{Need} - \text{Req}$$

$$(P1) = 122 - 102$$

$$= 20 //$$

(2.2) Request for resource allocation

Request for resource allocation is given as follows

$$(2.2) \geq (2.1)$$

Yes

$$\text{Request for resource allocation} = \text{Request for resource allocation}$$

$$201 - 200 = 1$$

$$= 1$$