# Team Notebook: UPF Programming Force

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## 1 Basic

## 1.1 Auxiliar Comparer

```
// returns true if the first argument goes before the second argument
// in the strict weak ordering it defines, and false otherwise.
struct classcomp {
   bool operator() (const int& lhs, const int& rhs) const
      {return lhs > rhs;}
};

int main() {
   set<int> set1;
   set<int, classcomp> set2;
   set1.insert(26); set1.insert(93); set1.insert(42); // 26, 42, 93
   set2.insert(26); set2.insert(93); set2.insert(42); // 93, 42, 26
```

```
for (auto it=set1.begin(); it!=set1.end(); ++it) cout << *it << " ";
    cout << "\n";
    for (auto it=set2.begin(); it!=set2.end(); ++it) cout << *it << " ";
}</pre>
```

#### 1.2 Libraries

| algorithm | heap, sort   | map     | map <s, t=""></s,>           |
|-----------|--------------|---------|------------------------------|
| cfloat    | DBL_MAX      | queue   | priority_queue               |
| cmath     | pow, sqrt    | set     | set <s></s>                  |
| cstdlib   | abs, rand    | sstream | istringstream, ostringstream |
| iostream  | cin, cout    | string  | string                       |
| iomanip   | setprecision | utility | pair <s, t=""></s,>          |
| list      | list <t></t> | vector  | vector <t></t>               |

#### 1.3 Macros

```
#define X first
#define Y second
#define LI long long
#define MP make_pair
#define PB push_back
#define SZ size()
#define SQ(a) ((a)*(a))
#define MAX(a,b) ((a)>(b)?(a):(b))
#define MIN(a,b) ((a)<(b)?(a):(b))
#define FOR(i,x,y) for(int i=(int)x; i<(int)y; i++)
#define RFOR(i,x,y) for(int i=(int)x; i>(int)y; i--)
#define SORT(a) sort(a.begin(), a.end())
#define RSORT(a) sort(a.rbegin(), a.rend())
#define IN(a,pos,c) insert(a.begin()+pos,1,c)
#define DEL(a,pos,cant) erase(a.begin()+pos,cant)
```

#### 1.4 Permutations

```
int N = 3;
int a[] = {1,2,3};
do {
```

```
for (int i = 0; i < N; ++i) cout << a[i] << " ";
  cout << "\n";
}
while (next_permutation(a, a + N));</pre>
```

#### 1.5 Precision cout

```
cout.setf(ios::fixed);
cout.precision(8);
```

## 2 Data Structures

## 2.1 Binary Indexed Tree

```
/* Binary indexed tree. Supports cumulative sum queries in O(log n) */
#define N (1<<18)
LL bit[N];

void add(LL* bit,int x,int val) {
   for(; x<N; x+=x&-x)
        bit[x]+=val;
}

LL query(LL* bit,int x) {
   LL res=0;
   for(;x;x-=x&-x)
        res+=bit[x];
   return res;
}</pre>
```

## 2.2 Square Root Trick

```
/* Partitions an array in sqrt(n) blocks of size sqrt(n) to support
 * O(sqrt(n)) range sum queries, O(sqrt(n)) range sum updates, and O(1)
 * point updates */
void update(LL *S, LL *A, int i, int k, int x) {
    S[i/k] = S[i/k] - A[i] + x;
    A[i] = x;
```

```
}
LL query(LL *S, LL *A, int lo, int hi, int k) {
    int sum=0, i=lo;
    while((i+1)%k != 0 && i <= hi)
        sum += A[i++];
    while(i+k <= hi)
        sum += S[i/k], i += k;
    while(i <= hi)
        sum += A[i++];
    return sum;
}</pre>
```

## 3 Dynamic Programming

### 3.1 Change Making Problem

## 3.2 Knapsack Problem

```
int N = 8; // numero de objetos N <= 1000 int v[] = \{1,6,7,1,8,3,7,5\}; // valor de objetos int p[] = \{5,3,7,1,8,2,7,3\}; // peso de objetos int A[1001][1001]; // matriz de resultados
```

```
int main() {
   int C = 7; // capacidad C <= 1000

for (int j = 0; j <= C; j++)
        A[0][j] = 0;
   for (int i = 1; i <= N; i++) {
        A[i][0] = 0;
        for (int j = 1; j <= C; j++) {
             A[i][j] = A[i-1][j];
            if (p[i-1] <= j) {
                int r = A[i-1][j-p[i-1]] + v[i-1];
                A[i][j] = MAX(A[i][j], r);
            }
        }
    }
   cout << A[N][C] << endl; // output: 12</pre>
```

#### 3.3 TSP

```
// TSP in O(n^2 * 2^n). Subset is bitmask, Cost is cost.
// tsp_memoize[subset][i] stores the shortest tsp of the subset starting
// If you have a starting node, it 's not included in the search .you add
    the distance to it at the beginning
Cost distances[N][N], tsp_memoize[1 << (N+1)][N];</pre>
const sentinel=-0x3f3f3f3f;
#define TSP(subset, i) (tsp_memoize[subset][i] == sentinel ? \
                                             tsp(subset, i):
                                                 tsp_memoize[subset][i])
Cost tsp(const Subset subset, const int i) {
       Subset without = subset ^ (1 << i);
       Cost minimum = numeric_limits<Cost>::max();
       for(int j=0; j<n_operas; j++) {</pre>
              if(j==i || (without & (1 << j)) == 0)</pre>
                      continue;
              Cost v = TSP(without, j);
              v += distances[i][j];
              if(v < minimum)</pre>
                      minimum = v:
       }
       return tsp_memoize[subset][i] = minimum;
```

## 4 Geometry

#### 4.1 Convex Hull

```
typedef int T; // posiblemente cambiar a double
typedef pair<T,T> P;
T xp(P p, P q, P r) {
   return (q.X-p.X)*(r.Y-p.Y) - (r.X-p.X)*(q.Y-p.Y);
struct Vect {
   P p, q; T dist;
   Vect(P &a, P &b) {
       p = a; q = b;
       dist = SQ(a.X - b.X) + SQ(a.Y - b.Y);
   bool operator<(const Vect &v) const {</pre>
       T t = xp(p, q, v.p);
       return t < 0 || t == 0 && dist < v.dist;
   }
};
vector<P> convexhull(vector<P> v) { // v.SZ >= 2}
   sort(v.begin(), v.end());
   vector<Vect> u;
   for (int i = 1; i < (int)v.SZ; i++)</pre>
       u.PB(Vect(v[i], v[0]));
   sort(u.begin(), u.end());
   vector<P> w(v.SZ, v[0]);
   int j = 1; w[1] = u[0].p;
   for (int i = 1; i < (int)u.SZ; i++) {</pre>
       T t = xp(w[j-1], w[j], u[i].p);
       for (j--; t < 0 && j > 0; j--)
           t = xp(w[j-1], w[j], u[i].p);
       j += t > 0 ? 2 : 1;
       w[j] = u[i].p;
   }
```

```
w.resize(j+1);
    return w;
}

int main() {
    vector<P> v;
    v.PB(MP(0, 1));
    v.PB(MP(1, 2));
    v.PB(MP(3, 2));
    v.PB(MP(2, 1));
    v.PB(MP(3, 1));
    v.PB(MP(6, 3));
    v.PB(MP(7, 0));
    vector<P> w = convexhull(v);
} // resultado: (0,1) (7,0) (6,3) (1,2)
```

## 5 Graphs

## 5.1 Bellman-Ford (Shortest Path with Negative Weights)

```
// Complexity: E * V - Input: directed graph
typedef pair<pair<int,int>,int> P; // par de nodos + coste
int N;
                                // numero de nodos
vector<P> v;
                                // representacion aristas
int bellmanford(int a, int b) {
   vector<int> d(N, 1000000000);
   d[a] = 0;
   for (int i = 1; i < N; i++)</pre>
       for (int j = 0; j < (int)v.SZ; j++)
           if (d[v[j].X.X] < 1000000000 && d[v[j].X.X] + v[j].Y <
               d[v[j].X.Y])
              d[v[j].X.Y] = d[v[j].X.X] + v[j].Y;
   for (int j = 0; j < (int)v.SZ; j++)
       if (d[v[j].X.X] < 1000000000 && d[v[j].X.X] + v[j].Y < d[v[j].X.Y])</pre>
           return -1000000000; // existe ciclo negativo
   return d[b];
}
int main(){
       N=8;
```

```
v.PB(MP(MP(0, 1), +2)); v.PB(MP(MP(1, 2), -1)); v.PB(MP(MP(1, 3), +1));
v.PB(MP(MP(2, 3), +1)); v.PB(MP(MP(6, 4), -1)); v.PB(MP(MP(4, 5), -1));
v.PB(MP(MP(5, 6), -1));

// min distance, negative cycle, unreachable cout << bellmanford(0, 3) << " " << bellmanford(4, 6) << " " << bellmanford(0, 7) << endl;</pre>
```

#### 5.2 Bron-Kerbosch

```
#define U unsigned int
typedef vector<short int> V;
vector<vector<U> > graf; // vertices/aristas del grafo
U numv, kmax; // # conjuntos/tamano grupo independiente
int evalua(V &vec) {
   for (int n = 0; n < vec.size(); n++)</pre>
       if (vec[n] == 1) return n;
   return -1;
}
void Bron_i_Kerbosch() {
   vector<U> v;
   U i, j, aux, k = 0, bandera = 2;
   vector<V> I, Ve, Va;
   I.PB(V()); Ve.PB(V()); Va.PB(V());
   for (i = 0; i < numv; i++) {</pre>
       I[0].PB(0); // conjunto vacio
       Ve[0].PB(0); // conjunto vacio
       Va[0].PB(1); // contiene todos
   }
   while(true) {
       switch(bandera) {
       case 2: // paso 2
           v.PB(evalua(Va[k]));
          I.PB(V(I[k].begin(), I[k].end()));
           Va.PB(V(Va[k].begin(), Va[k].end()));
           Ve.PB(V(Ve[k].begin(), Ve[k].end()));
           aux = graf[v[k]].size();
```

```
I[k+1][v[k]] = 1; Va[k+1][v[k]] = 0;
      for (i = 0; i < aux; i++) {</pre>
         j = graf[v[k]][i]; Ve[k+1][j] = Va[k+1][j] = 0;
      k = k + 1; bandera = 3;
      break:
   case 3: // paso 3
      for (i = 0, bandera = 4; i < numv; i++) {</pre>
         if (Ve[k][i] == 1) {
            aux = graf[i].size();
            for (j = 0; j < aux; j++)
                if (Va[k][graf[i][j]] == 1)
                   break:
            if (j == aux) { i = numv; bandera = 5; }
         }
      }
      break:
   case 4: // paso 4
      if (evalua(Ve[k]) == -1 && evalua(Va[k]) == -1) {
         for (int n = 0; n < numv; n++)
            if (I[k][n] == 1) cout<< n << " ";</pre>
         cout << endl:</pre>
         if (k > kmax) kmax = k;
         bandera = 5:
      }
      else bandera = 2; // ir a paso 2
   case 5: // paso 5
      k = k - 1; v.pop_back(); I[k].clear();
      I[k].assign(I[k+1].begin(), I[k+1].end());
      I[k][v[k]] = 0; I.pop_back(); Ve.pop_back();
      Va.pop_back(); Ve[k][v[k]] = 1; Va[k][v[k]] = 0;
      if (k == 0) {
         if (evalua(Va[0]) == -1) return;
         bandera = 2; // ir a paso 2
      else bandera = 3; // ir a paso 3
   break;
   }
}
```

}

```
int main() {
   U idx, i; stringstream ss; string linea;
   while (cin >> numv) {
       getline(cin, linea);
       for (i = 0; i < numv; i++) { // Lectura del grafo</pre>
           // vertices adjacentes al i-esimo vertice
           vector<U> bb; graf.PB(bb);
           getline(cin, linea);
           ss << linea:
           while (ss >> idx) graf[i].PB(idx);
           ss.clear():
       // Llamada al algoritmo
       kmax = 0;
       cout << "Conjuntos independientes: "<< endl;</pre>
       if (numv > 0)
           Bron_i_Kerbosch();
       cout << "kmax: " << kmax << endl;</pre>
       // Limpieza variables
       for (i = 0; i < numv; i++) graf[i].clear();</pre>
       graf.clear();
   }
```

## 5.3 Dijkstra (Shortest Path)

```
// Complexity: ElogV - Input: undirected graph
                     // tipo de costes
typedef int V;
typedef pair<V,int> P; // par de (coste,nodo)
typedef set<P> S;
                     // conjunto de pares
                     // numero de nodos
int N;
vector<P> A[10001]; // listas adyacencia (coste,nodo)
// int prec[201]; // predecesores (nodes from s to t)
// another way to obtain a path (above all, if there is
// more than one, consists in using BFS from the target
// and add to the queue those nodes that lead to the
// minimum cost in the preceeding node)
V diikstra(int s. int t) {
   S m:
                              // cola de prioridad
   vector<V> z(N, 1000000000); // distancias iniciales
```

```
z[s] = 0:
                              // distancia a s es 0
                               // insertar (0,s) en m
   m.insert(MP(0, s));
   while (m.SZ > 0) {
       P p = *m.begin(); // p=(coste,nodo) con menor coste
       m.erase(m.begin()); // elimina este par de m
       if (p.Y == t) return p.X; // cuando nodo es t, acaba
       // para cada nodo adjacente al nodo p.Y
       for (int i = 0; i < (int)A[p.Y].SZ; i++) {</pre>
          // q = (coste hasta nodo adjacente, nodo adjacente)
          P q(p.X + A[p.Y][i].X, A[p.Y][i].Y);
          // si q.X es la menor distancia hasta q.Y
          if (q.X < z[q.Y]) {
              m.erase(MP(z[q.Y], q.Y)); // borrar anterior
              m.insert(q);
                                     // insertar q
              z[q.Y] = q.X;
                                      // actualizar distancia
                             // prec[q.Y] = p.Y;
                                                    // actualizar
                                 predecesores
          }
       }
   return -1;
int main() {
   N = 6:
                    // solucion 0-1-2-4-3-5, coste 11
   A[0].PB(MP(2, 1)); // arista (0, 1) con coste 2
   A[0].PB(MP(5, 2)); // arista (0, 2) con coste 5
   A[1].PB(MP(2, 2)); // arista (1, 2) con coste 2
   A[1].PB(MP(7, 3)); // arista (1, 3) con coste 7
   A[2].PB(MP(2, 4)); // arista (2, 4) con coste 2
   A[3].PB(MP(3, 5)); // arista (3, 5) con coste 3
   A[4].PB(MP(2, 3)); // arista (4, 3) con coste 2
   A[4].PB(MP(8, 5)); // arista (4, 5) con coste 8
   cout << dijkstra(0, 5) << endl;</pre>
```

## 5.4 Floyd-Warshall (All Pairs Shortest Path)

```
// Complexity: n^3
// A: matriz n*n de adyacencia con costes
// ausencia de arista representada por un numero grande
for (int k = 0; k < n; k++)
    for (int i = 0; i < n; i++)</pre>
```

```
for (int j = 0; j < n; j++)
    A[i][j] = MIN(A[i][j], A[i][k] + A[k][j]);</pre>
```

## 5.5 Kruskal (Minimum Spanning Tree)

```
// Complexity: ElogV - Input: undirected graph
typedef vector<pair<int,pair<int,int> > V;
int N, mf[2000]; // numero de nodos N <= 2000</pre>
V v;
               // vector de aristas
               // (coste, (nodo1, nodo2))
// vector< pair<long, int> > K[3001]; // kruskal tree
int set(int n) { // conjunto conexo de n
   if (mf[n] == n) return n;
   else mf[n] = set(mf[n]); return mf[n];
}
int kruskal() {
   int a, b, sum = 0;
   sort(v.begin(), v.end());
   for (int i = 0; i < N; i++)</pre>
       mf[i] = i; // inicializar conjuntos conexos
   for (int i = 0; i < (int)v.SZ; i++) {</pre>
       a = set(v[i].Y.X), b = set(v[i].Y.Y);
       if (a != b) { // si conjuntos son diferentes
           mf[b] = a; // unificar los conjuntos
           sum += v[i].X; // agregar coste de arista
                     // K[v[i].Y.X].PB(MP(v[i].X, v[i].Y.Y));
                      // K[v[i].Y.Y].PB(MP(v[i].X, v[i].Y.X));
       }
   return sum:
int main() {
   N = 5; // solucion 13 (0,3),(1,2),(2,3),(3,4)
   v.PB(MP(4,MP(0,1))); // arista (0,1) coste 4
   v.PB(MP(4,MP(0,2))); // arista (0,2) coste 4
   v.PB(MP(3,MP(0,3))); // arista (0,3) coste 3
   v.PB(MP(6,MP(0,4))); // arista (0,4) coste 6
   v.PB(MP(3,MP(1,2))); // arista (1,2) coste 3
```

```
v.PB(MP(7,MP(1,4))); // arista (1,4) coste 7
v.PB(MP(2,MP(2,3))); // arista (2,3) coste 2
v.PB(MP(5,MP(3,4))); // arista (3,4) coste 5
cout << kruskal() << endl;
}</pre>
```

### 5.6 Maximum Bipartite Matching

```
/* Input: VVI with 1 if connected, 0 if not. mr and mc have the matches
 * for each side. Complexity: E * V.
 * From Stanford University's notebook. */
typedef vector<int> VI;
typedef vector<VI> VVI;
bool FindMatch(int i, const VVI &w, VI &mr, VI &mc, VI &seen) {
   for (int j = 0; j < w[i].size(); j++) {</pre>
       if (w[i][j] && !seen[j]) {
           seen[i] = true;
           if (mc[i] < 0 || FindMatch(mc[i], w, mr, mc, seen)) {</pre>
               mr[i] = j;
              mc[j] = i;
               return true;
           }
       }
   return false;
}
int BipartiteMatching(const VVI &w, VI &mr, VI &mc) {
   mr = VI(w.size(), -1);
   mc = VI(w[0].size(), -1);
   int ct = 0:
   for (int i = 0; i < w.size(); i++) {</pre>
       VI seen(w[0].size());
       if (FindMatch(i, w, mr, mc, seen)) ct++;
   return ct;
```

#### 5.7 Min Cost Max Flow

```
/* From Stanford University's notebook.
* To perform minimum weighted bipartite matching:
* - Capacity between nodes = 1 (cost whatever given by the problem)
* - Capacity from source = 1 and cost = 0
* - Capacity to sink = 1 and cost = 0
* Output: <maximum flow value - minimum cost value>
* Complexity: O(|V|^2) per augmentation
             max flow: O(|V|^3) augmentations
             min cost max flow: O(|V|^4 * MAX\_EDGE\_COST) augmentations
*/
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef long long L;
typedef vector<L> VL;
typedef vector<VL> VVL;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
const L INF = numeric_limits<L>::max() / 4;
struct MinCostMaxFlow {
   int N:
   VVL cap, flow, cost;
   VI found:
   VL dist, pi, width;
   VPII dad;
   MinCostMaxFlow(int N) :
       N(N), cap(N, VL(N)), flow(N, VL(N)), cost(N, VL(N)),
       found(N), dist(N), pi(N), width(N), dad(N) {}
   void AddEdge(int from, int to, L cap, L cost) {
       this->cap[from][to] = cap;
       this->cost[from][to] = cost:
   }
   void Relax(int s, int k, L cap, L cost, int dir) {
       L val = dist[s] + pi[s] - pi[k] + cost;
       if (cap && val < dist[k]) {</pre>
          dist[k] = val;
          dad[k] = make_pair(s, dir);
           width[k] = min(cap, width[s]);
       }
```

```
}
   L Dijkstra(int s, int t) {
       fill(found.begin(), found.end(), false);
       fill(dist.begin(), dist.end(), INF);
       fill(width.begin(), width.end(), 0);
       dist[s] = 0;
       width[s] = INF;
       while (s !=-1) {
           int best = -1:
          found[s] = true;
          for (int k = 0; k < N; k++) {
              if (found[k]) continue;
              Relax(s, k, cap[s][k] - flow[s][k], cost[s][k], 1);
              Relax(s, k, flow[k][s], -cost[k][s], -1);
              if (best == -1 || dist[k] < dist[best]) best = k;</pre>
          s = best:
       }
       for (int k = 0; k < N; k++)
           pi[k] = min(pi[k] + dist[k], INF);
       return width[t]:
   pair<L, L> GetMaxFlow(int s, int t) {
       L totflow = 0, totcost = 0;
       while (L amt = Dijkstra(s, t)) {
           totflow += amt;
          for (int x = t; x != s; x = dad[x].first) {
              if (dad[x].second == 1) {
                  flow[dad[x].first][x] += amt;
                  totcost += amt * cost[dad[x].first][x];
              else {
                  flow[x][dad[x].first] -= amt;
                  totcost -= amt * cost[x][dad[x].first];
          }
       return make_pair(totflow, totcost);
};
```

### 5.8 Tarjan

```
int index, ct;
vector<bool> I;
vector<int> D, L, S;
vector<vector<int> > V; // listas de adyacencia
void tarjan (unsigned n) {
   D[n] = L[n] = index++;
   S.push_back(n);
   I[n] = true;
   for (unsigned i = 0; i < V[n].size(); ++i) {</pre>
       if (D[V[n][i]] < 0) {</pre>
           tarjan(V[n][i]);
           L[n] = MIN(L[n], L[V[n][i]]);
       }
       else if (I[V[n][i]])
           L[n] = MIN(L[n], D[V[n][i]]);
   }
   if (D[n] == L[n]) {
       ++ct;
       // todos los nodos eliminados de S pertenecen al mismo scc
       while (S[S.size() - 1] != n) {
           I[S.back()] = false;
           S.pop_back();
       }
       I[n] = false;
       S.pop_back();
void scc() {
   index = ct = 0;
   I = vector<bool>(V.size(), false);
   D = vector<int>(V.size(), -1);
   L = vector<int>(V.size());
   S.clear();
   for (unsigned n = 1; n <= V.size(); ++n)</pre>
       if (D[n] < 0)
         tarjan(n);
   // ct = numero total de scc
```

### 5.9 Topological Sort

```
vector<int> A[101]; // adjacency list (directed graph without cycles)
int inbound[101]; // number of nodes that point to each node
vector<int> fo; // final order
// M = number of nodes (there might be 'lonely' nodes)
void toposort(int M) {
   stack<int> order;
   int current;
   // Search for roots (identifiers might change between
   // problems (e.g. 1 to M))
   for(int m = 0; m < M; m++){
       if(inbound[m] == 0)
           order.push(m);
   }
   // Start topsort from roots
   while(!order.empty()){
       // Pop from stack
       current = order.top();
       order.pop();
       // Save order in fo
       fo.push_back(current);
       // Add childs only if inbound is 0
       for (int i = 0; i < A[current].size(); ++i)</pre>
           inbound[A[current][i]]--;
           if (inbound[A[current][i]] == 0)
               order.push(A[current][i]);
       }
}
int main() {
   A[0].push_back(1); A[0].push_back(2); A[2].push_back(1);
   inbound[0] = 0; inbound[1] = 2; inbound[2] = 1;
   toposort(3);
   for (int i = 0; i < fo.size(); ++i) cout << fo[i] << " ";</pre>
   // 0 2 1
```

### 6 Math

#### 6.1 Catalan Numbers

```
unsigned long long v[34]; // 1, 1, 2, 5, 14, 42, 132, 429, 1430, ...
void catalan(){
  v[0] = 1;
  for (int i = 1; i < 34; ++i){
     unsigned long long sum = 0;
     for (int j = 0; j < i; ++j){
         sum += v[j] * v[i-j-1];
     }
     v[i] = sum;
}</pre>
```

### 6.2 Complex Numbers

```
// Complex number class, from Stanford's Notebook. Required for FFT
struct cpx {
   cpx(){}
   cpx(double aa):a(aa){}
   cpx(double aa, double bb):a(aa),b(bb){}
   double a, b;
   double modsq(void) const { return a * a + b * b; }
   cpx bar(void) const { return cpx(a, -b); }
};
cpx operator +(cpx a, cpx b) { return cpx(a.a + b.a, a.b + b.b); }
cpx operator *(cpx a, cpx b) {
   return cpx(a.a * b.a - a.b * b.b, a.a * b.b + a.b * b.a);
cpx operator /(cpx a, cpx b) {
   cpx r = a * b.bar();
   return cpx(r.a / b.modsq(), r.b / b.modsq());
cpx EXP(double theta) { return cpx(cos(theta), sin(theta)); }
```

### 6.3 Exponent

```
template <typename T, typename U> T expo(T &t, U n) {
```

```
if (n == U(0)) return T(1);
else {
    T u = expo(t, n/2);
    if (n%2 > 0) return u*u*t;
    else return u*u;
}
```

### 6.4 FFT

```
// from Stanford's notebook:
    https://web.stanford.edu/~liszt90/acm/notebook.html
         input array
// out: output array
// step: {SET TO 1} (used internally)
// size: length of the input/output {MUST BE A POWER OF 2}
// dir: either plus or minus one (direction of the FFT)
// RESULT: out[k] = \sum_{j=0}^{size - 1} in[j] * exp(dir * 2pi * i * j *
    k / size)
const double two_pi = 4 * acos(0);
void FFT(cpx *in, cpx *out, int step, int size, int dir)
   if(size < 1) return;</pre>
   if(size == 1)
       out[0] = in[0];
       return:
   FFT(in, out, step * 2, size / 2, dir);
   FFT(in + step, out + size / 2, step * 2, size / 2, dir);
   for(int i = 0 ; i < size / 2 ; i++)</pre>
       cpx even = out[i];
       cpx odd = out[i + size / 2];
       out[i] = even + EXP(dir * two_pi * i / size) * odd;
       out[i + size / 2] = even + EXP(dir * two_pi * (i + size / 2) /
           size) * odd;
```

#### 6.5 Greatest Common Divisor

```
// in algorithm library: __gcd(a, b)
int gcd(int a, int b) {
   if (a < b) return gcd(b, a);
   else if (a%b == 0) return b;
   else return gcd(b, a%b);
}
gcd(a,b)*lcm(a,b) = a*b</pre>
```

#### 6.6 Newton Method

```
long double tolerance = 1E-6;
long double c0 = 1.0;
long double c1 = 1.0;
bool solutionFound = false;
// find the value of 'c' that makes the function equal to = 0
// might also be used in optimization problems setting y as
// the first derivative and yprime as the second
while (true)
       long double y = /* formula of the original function */;
       long double yprime = /* formula of the first derivative respect to
           c */:
       c1 = c0 - y / yprime;
       if ((fabs(c1 - c0) / fabs(c1)) < tolerance)</pre>
              solutionFound = true;
              break;
       }
       c0 = c1;
```

### 6.7 Primes

```
int v[10000]; // primes

void savePrimes()
{
   int k = 0;
```

```
v[k++] = 2:
    for (int i = 3; i <= 10010; i += 2) {
       bool b = true;
       for (int j = 0; b \&\& v[j] * v[j] <= i; j++)
           b = i \%v[i] > 0;
       if (b)
           v[k++] = i;
   }
}
bool isPrime(int x){
       bool prime = true;
       for (int j = 0; prime && v[j] * v[j] <= x; j++)</pre>
               prime = x%v[j] > 0;
       return prime;
}
// probar si un numero x <= 100000000 es primo
int main()
{
    savePrimes();
    cout << isPrime(4);</pre>
```

## 7 Sequences

## 7.1 Binary Search

```
// binary_search function can be found at algorithm library
// devuelve el i mas pequeno tal que t <= v[i]
// si no existe tal i, devuelve v.SZ
template<typename T> int bb(T t, vector<T> &v) {
   int a = 0, b = v.SZ;
   while (a < b) {
      int m = (a + b)/2;
      if (v[m] < t) a = m+1; else b = m;
   }
   return a;
}</pre>
```

### 7.2 Ternary Search

```
double E = 0.0000001; // tolerance
double L = 200000; // R and L are extreme possible values...
double R = -200000; // ... for the optimized parameter
while (1) {
    double dist = R - L;
    if (fabs(dist) < E) break;
    double leftThird = L + dist / 3;
    double rightThird = R - dist / 3;
    // f is the function which we are optimizing
    if (f(leftThird) < f(rightThird))
        R = rightThird;
    else
        L = leftThird;
}</pre>
```

#### 7.3 Vector Partition

```
bidirectional_iterator partition(bidirectional_iterator start,
                           bidirectional_iterator end,
                           Predicate p);
bool IsOdd(int i) {return (i%2==1);}
int main () {
   vector<int> myvector;
   vector<int>::iterator it, bound;
   // set some values:
   for (int i=1; i<10; ++i)</pre>
       myvector.push_back(i); // 1 2 3 4 5 6 7 8 9
   bound = partition(myvector.begin(), myvector.end(), IsOdd);
   // print out content:
   cout << "odd members:";</pre>
   for (it=myvector.begin(); it!=bound; ++it)
       cout << " " << *it;
   cout << "\neven members:";</pre>
   for (it=bound; it!=myvector.end(); ++it)
       cout << " " << *it;
   cout << endl;</pre>
```

# 8 Strings

#### 8.1 KMP

```
/*Search of substring in O(n+k) */
void TablaKMP(string T,vector<int> &F)
   int pos = 2; // posicion actual en F
   int cnd = 0; // ndice en T del siguiente carcter del actual candidato
        en la subcadena
   F[0] = -1:
   while(pos <= T.size())</pre>
       if(T[pos - 1] == T[cnd])
       {//siguiente candidato coincidente en la cadena
           cnd++;
          F[pos] = cnd;
          pos++;
       }else if(cnd > 0)
       {//sifallan coincidencias consecutivas entonces asignamos valor
           conocido la primera vez
           cnd = F[cnd];
       }else{
           F[pos] = 0;
          pos++;
       }
int KMPSearch(string T, string P)//T: texto donde se busca P: palabra a
    buscar salida: posicion del match (-1 si no encontrado
{
   int k = 0 ; //puntero de T
   int i = 0 ; //avance en P
   vector<int> F(T.size(),0);
   if(T.size() >= P.size())
```

```
TablaKMP(T,F);//optimizacin para no repetir busquedas de
           subcadenas que no hacen match
       while(k+i < T.size())</pre>
          if(P[i] == T[k+i])
              if(i == P.size()-1)
                  return k; //modificando el return podemos devolver
                      todos los matches
              }
              i++;
          }else{
              k += i-F[i];
              if(i > 0)
                  i = F[i];
          }
       }
   }
   return -1;
int main(){
   string T = "PARTICIPARIA CON MI PARACAIDAS PARTICULAR";
   string P = "RTI";
   cout<<KMPSearch(T,P)<<endl;</pre>
```